15-451 Assignment 01

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1a.

Compute $n^3/4$ in constant time.

Use DeterministicSelect to select the $n^{3/4}$ th largest number in O(n) time.

Then filter out the elements greater than or equal to it in O(n) time.

Now sort the $n^{3/4}$ numbers using mergesort in $O(n^{3/4}log(n^{3/4}))$ time.

The algorithm seems to be dominated by the latter expression, but it can be reduced to O(n) as follows:

$$O(n^{3/4}log(n^{3/4}))$$

 $\leq O(\frac{3}{4}n^{3/4}log(n))$

Notice that $O(n^{1/4}) \ge O(\log(n))$ so we can make the following substitution:

$$\leq O(\frac{3}{4}n^{3/4}n^{1/4})$$

$$\leq O(n)$$

1b.

Pair up the elements, and for each pair, compare the elements. $(\frac{n}{2}$ comparisons)

Call the larger element a "winner" and the smaller element a "loser".

Among the $\frac{n}{2}$ winners, find the max by going one by one keeping track of the max so far. $(\frac{n}{2} - 1$ comparisons)

This is the max of all the elements.

Among the $\frac{n}{2}$ losers, find the min by going one by one keeping track of the min so far. $(\frac{n}{2} - 1$ comparisons)

The sum of all the comparisons is $(\frac{n}{2}) + (\frac{n}{2} - 1) + (\frac{n}{2} - 1) = \frac{3n}{2} - 2$

 $\mathbf{2}$

Lorem ipsum