

# 15-150 Assignment 2

Karan Sikka

ksikka@andrew.cmu.edu

Section G but im not sure

January 25, 2012

---

## 1: Task 2.1

---

No, this does not type check. When `smlnj` tries to evaluate the expression `square 7.0`, it fails to do so, because `7.0` is not an `int` but `square` is a function of type `int -> int`.

---

## 2: Task 2.2

---

- a. The value `3 : int` is put in the place of the `x` in line 5. This is because of the `val` declaration in line 4 takes precedence over the `x` defined in the arguments of the function, due to the scope properties of the `let` clause.
- b. The value of the expression `5.2 * (real x)` is substituted for `p`, where `x` is `3`, as said in part a. This evaluates to `5.2 * (real 3)`, then `5.2 * 3.0`, then `15.6`. So ultimately the value which replaces `p` is `15.6`. This is because `p` is within the `in` block so it takes the value prescribed to it in the `let` block, specifically line 5. Also, SML is pass-by-value, so the value is passed to the `in` block, not the expression.
- c. The value of the first argument passed to the function replaces `x`, because the declaration of the function names the first argument `x`. The `x` in line 4 or 7 do not substitute for the `x` in line 11, because those are within the scope of the inner `let` statement, and they disappear after the end on line 9.
- d. `assemble (x, 2.0)` evaluates to `58`.

---

## 3: Task 2.3

---

```
let val x : real = real (square 6)
in 3 + (trunc x)
end
```

$\mapsto$

```
let val x : real = real (6 * 6)
in 3 + (trunc x)
end
```

$\mapsto$

```
let val x : real = real (36)
in 3 + (trunc x)
end
```

$\mapsto$ 

```
let val x : real = 36.0
in 3 + (trunc x)
end
```

 $\mapsto 3 + (\text{trunc } 36.0)$  $\mapsto 3 + 36$  $\mapsto 39$ 

---

#### 4: Task 2.4

---

They are equivalent. The fact function steps to a case statement, which eventually steps to  $\sim 1 * \text{fact}(\sim 2)$ . This repeats for quite a long time, and assuming integers do not overflow, and that our heap is infinitely large, this function will infinite loop.

In the case of f 10, this function steps to f 10. It steps to the same thing consistently, using no additional memory. This will also infinite loop.

Therefore, the expressions are equivalent.

---

#### 5: Task 2.5

---

Consider the expression `fact 3`.

```
fact(3)
 $\mapsto$  case 3 of 0 => 1 | _ => 3 * fact(3-1)
 $\mapsto$  3 * fact(3-1)
 $\mapsto$  3 * fact(2)
 $\mapsto$  3 * case 2 of 0 => 1 | _ => 2 * fact(2-1)
 $\mapsto$  3 * 2 * fact(2-1)
 $\mapsto$  3 * 2 * fact(1)
 $\mapsto$  6 * fact(1)
 $\mapsto$  6 * case 1 of 0 => 1 | _ => 1 * fact(1-1)
 $\mapsto$  6 * 1 * fact(1-1)
 $\mapsto$  6 * 1 * fact(0)
 $\mapsto$  6 * fact(0)
 $\mapsto$  6 * case 0 of 0 => 1 | _ => 0 * fact(0-1)
 $\mapsto$  6 * 1
 $\mapsto$  6
```

Thus, `fact(3)` evaluates to 6.

Now consider:

---

```

    fact(fact(3))
  ≅ fact(6)                                     equivalence proven above
  ≅ case 6 of 0 => 1 | _ 6 * fact(6-1) step
  ≅ 6 * fact(6-1)                               step
  ≅ fact(3) * fact(fact(3) - 1)               step, equiv. is symmetric, fact(3)=6 as shown above.

```

Therefore:  
 $\text{fact}(\text{fact}(3)) \cong \text{fact}(3) * \text{fact}(\text{fact}(3) - 1)$

---

### 6: Task 3.2

---

$H_0$  is 0.  $H_1$  is 1.  $H_2$  is  $\frac{3}{2}$ .  $H_3$  is  $\frac{11}{6}$ .  $H_4$  is  $\frac{25}{12}$ .

---

### 7: Task 3.4

---

$I_0$  is 0.  $I_1$  is 1.  $I_2$  is  $\frac{1}{2}$ .  $I_3$  is  $\frac{5}{6}$ .  $I_4$  is  $\frac{7}{12}$ .

---

### 8: Task 4.1

---

Claim: For all natural numbers  $n$ ,  $2(n+1) \cong \text{double}(n+1)$ .

We proceed by induction on  $n$ . Base Case:

```

    double n where n = 0
  ≅ double 0                                     substituting
  ≅ case 0 of 0 => 0 | _ 2 + double(n-1) step
  ≅ 0                                           step
  ≅ 2*0                                         math, also step is symmetric
  ≅ 2n where n = 0                             step

```

Inductive Hypothesis: Assume  $\text{double } k \cong 2 * k$  for some  $k \in \mathbb{N}$

Inductive Step: Consider:

```

    2(k+1)
  ≅ 2k + 2                                     math, distributive property
  ≅ 2 + double(k)                             IH
  ≅ 2 + double(k + 1 - 1)                     math, step, equiv. is sym.
  ≅ case 0 of 0 => 0 | 2 + double(k + 1 - 1) step, equiv. is symmetric
  ≅ double(k+1)                               step, equiv. is symmetric

```

Therefore  $2(k+1) \cong \text{double}(k+1)$ . By induction, the claim holds.

---

### 9: Task 4.2

---

Claim: For all natural numbers  $n$ ,  $\text{summ } n \cong (n*(n+1)) \text{ div } 2$ .

We proceed by induction on  $n$ .

Base Case:

	<code>summ n where n = 0</code>	
$\cong$	<code>summ 0</code>	substitution
$\cong$	<code>case 0 of 0 =&gt; 0   _ 0 + (summ(0-1))</code>	step
$\cong$	<code>0</code>	step
$\cong$	<code>0*1</code>	math
$\cong$	<code>0*(0+1)</code>	math
$\cong$	<code>0*(0+1) div 2</code>	math
$\cong$	<code>n*(n+1) div 2 where n = 0</code>	substitution

Inductive Hypothesis:

Assume `summ n  $\cong$  n*(n+1) div 2` for some  $n \in \mathbb{N}$ .

Inductive Step:

First, we make the assumption that `n+1` evaluates to a value that we will call `n+1`. Consider:

$\cong$	<code>summ(n+1)</code>	
$\cong$	<code>case n+1 of 0 =&gt; 0   _ n+1 + summ((n+1)-1)</code>	step
$\cong$	<code>n+1 + summ((n+1)-1)</code>	step
$\cong$	<code>n+1 + summ(n)</code>	math, addition is associative
$\cong$	<code>n+1 + n*(n+1) div 2</code>	IH
$\cong$	<code>((n+1)(n+2)) div 2</code>	algebra shown below

The algebra:

$$\frac{n^2 + n}{2} + \frac{2(n+1)}{2}$$

Distribute the right fraction:

$$\frac{n^2 + n}{2} + \frac{2n + 2}{2}$$

Add the fractions:

$$\frac{n^2 + 3n + 2}{2}$$

Factor:

$$\frac{(n+1)(n+2)}{2}$$

By induction, the claim holds.