

# 15-451 Assignment 05

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## 1: Cell Towers

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## 2: Eliminating Negative Edges

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(a) (b) (c)

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## 3: Color Me Red, Color Me Blue

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For convenience, we will call a  $(k, n - k)$  partition which minimizes split edges a “minimizing k-partition”.

We will consider a function `kPart` as follows:

Given a node  $n$ , value  $i$  ( $1 \leq i \leq k$ ) and bit  $b$ ,

The function returns `NONE`, or `SOME (A, B, S)`.

`kPart(n, k, b) = NONE` signifies that no k-partition of vertices in the subtree rooted at  $n$  exists, which happens when  $(k > |n|)$ .

`kPart(n, k, b) = SOME(A, B, S)` signifies that  $A, B$  is some minimizing k-partition the nodes of the subtree rooted at  $n$  such that  $n \in A \iff b = 1$ ,  $S$  is the set of split edges

We create a memo table for the computation which we assume is available to the function.

The table can be accessed by  $T[n][k][b]$  where  $n$  is a node,  $k$  is an int from 0 to  $k$ ,  $b$  is a bit.

We initialize the table with `NULL`s, and compute the values in the following recursive function.

Note that the table has  $n * (k + 1) * 2$  cells.

The procedure to compute `kPart(r, k, b)` is as follows: If  $T[r][k][i] \neq \text{NULL}$ , return  $T[r][k][i]$  If  $k = 0$ ,  $T[r][k][i] = (, \text{all\_nodes}(b), )$  If  $k > n$ ,  $T[r][k][i] = \text{DNE}$

If both  $r\_i\text{left}$  and  $r\_i\text{right}$  are `NULL`, If  $k = 1$   $T[r][k][i] = (r, , )$  Else

If only  $r\_i\text{left}$  is `NULL`,  $T[r][k][i] =$  The  $(A, B, S)$  with the minimum  $|S|$  over the following answers computed for all possible colorings of  $r$  and  $r\_i\text{right}$  that might yield a minimizing k-partition:  $r \in A, r\_i\text{right} \in A \Rightarrow A_r, B_r, S_r = \text{kPart}(r\_i\text{right}, k-1, 1)$  answer is  $(r \cup A_r, B_r, S_r)$   $r \in A, r\_i\text{right} \in B \Rightarrow A_r, B_r, S_r = \text{kPart}(r\_i\text{right}, k-1, 0)$  answer is  $(r \cup A_r, B_r, S_r \cup (r, r\_i\text{right}))$   $r \in B, r\_i\text{right} \in A \Rightarrow A_r, B_r, S_r = \text{kPart}(r\_i\text{right}, k, 1)$  answer is  $(A_r, r \cup B_r, S_r \cup (r, r\_i\text{right}))$   $r \in B, r\_i\text{right} \in B \Rightarrow 1 + T[r\_i\text{right}][k][0]$   $A_r, B_r, S_r = \text{kPart}(r\_i\text{right}, k, 0)$  answer is  $(A_r, r \cup B_r, S_r)$

If only  $r.\text{right}$  is NULL, Logic is symmetric to above case.

If  $r.\text{left}$  and  $r.\text{right}$  are not NULL,  $T[r][k][i] = \text{The } (A, B, S) \text{ with the minimum } |S| \text{ over the following answers computed for all possible colorings of } r, r.\text{right}, \text{ and } r.\text{left} \text{ that might yield a minimizing } k\text{-partition: Left and right subtrees can be partitioned in at most } k+1 \text{ ways such that the sum of their } A \text{ sets should be } k-1 \text{ if } r \text{ in } A \text{ or } k \text{ otherwise. We must minimize the number of split edges across all variations of } A\text{-set-size.}$

$r \in A, r- > \text{left} \in A, r- > \text{right} \in A$

$\text{minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k - 1$

$A_l, B_l, S_l = \text{kPart}(r- > \text{left}, i, 1)$

$A_r, B_r, S_r = \text{kPart}(r- > \text{right}, j, 1) \text{ answer is } (A_l \cup A_r \cup r, B_l \cup B_r, S_l \cup S_r)$

$r \in A, r- > \text{left} \in A, r- > \text{right} \in B \text{ minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k - 1$   
 $\text{kPart}(r- > \text{left}, i, 1) + 1 + \text{kPart}(r- > \text{right}, j, 0) \text{ answer is } (A_l \cup A_r \cup r, B_l \cup B_r, S_l \cup S_r \cup (r, r- > \text{right}))$

$r \in A, r- > \text{left} \in B, r- > \text{right} \in A \text{ minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k - 1$   
 $\text{kPart}(r- > \text{left}, i, 0) + \text{kPart}(r- > \text{right}, j, 1) \text{ answer is } (A_l \cup A_r \cup r, B_l \cup B_r, S_l \cup S_r \cup (r, r- > \text{left}))$

$r \in A, r- > \text{left} \in B, r- > \text{right} \in B \text{ minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k - 1$   
 $\text{kPart}(r- > \text{left}, i, 0) + 1 + \text{kPart}(r- > \text{right}, j, 0) \text{ answer is } (A_l \cup A_r \cup r, B_l \cup B_r, S_l \cup S_r \cup (r, r- > \text{left}), (r, r- > \text{right}))$

$r \in B, r- > \text{left} \in A, r- > \text{right} \in A \text{ minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k$   
 $\text{kPart}(r- > \text{left}, i, 1) + 1 + \text{kPart}(r- > \text{right}, j, 1) \text{ answer is } (A_l \cup A_r, B_l \cup B_r \cup r, S_l \cup S_r \cup (r, r- > \text{left}), (r, r- > \text{right}))$

$r \in B, r- > \text{left} \in A, r- > \text{right} \in B \text{ minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k$   
 $\text{kPart}(r- > \text{left}, i, 1) + \text{kPart}(r- > \text{right}, j, 0) \text{ answer is } (A_l \cup A_r, B_l \cup B_r \cup r, S_l \cup S_r \cup (r, r- > \text{left}))$

$r \in B, r- > \text{left} \in B, r- > \text{right} \in A \text{ minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k$   
 $\text{kPart}(r- > \text{left}, i, 0) + 1 + \text{kPart}(r- > \text{right}, j, 1) \text{ answer is } (A_l \cup A_r, B_l \cup B_r \cup r, S_l \cup S_r \cup (r, r- > \text{right}))$

$r \in B, r- > \text{left} \in B, r- > \text{right} \in B \text{ minimize } |S|_{\text{overall}}, i, j \geq 0, \text{st } i + j = k$   
 $\text{kPart}(r- > \text{left}, i, 0) + \text{kPart}(r- > \text{right}, j, 0) \text{ answer is } (A_l \cup A_r, B_l \cup B_r \cup r, S_l \cup S_r)$

return  $T[r][k][i]$