## CSE 547 - Assignment 1

Philip Pham

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#### Problem 0

**List of collaborators:** I have not collaborated with anyone.

List of acknowledgements: None.

Certify that you have read the instructions: I have read and understood these policies.

### Problem 1: Gaussian Random Projections and Inner Products

Let  $\phi(x) = \frac{1}{\sqrt{m}}Ax$  represent our random projection of  $x \in \mathbb{R}^d$ , with A an  $m \times d$  projection matrix with each entry sampled i.i.d from N(0,1). Note that each row of A is a random projection vector,  $v^{(i)}$ .

The norm preservation theorem states that for all  $x \in \mathbb{R}^d$ , the norm of the random projection  $\phi(x)$  approximately maintains the norm of the original x with high probability:

$$\mathbb{P}\left((1-\epsilon)\|x\|^2 \le (1+\epsilon)\|x\|^2 \le 1 - 2\exp\left(-\left(\epsilon^2 - \epsilon^3\right)m/4\right)\right),\tag{1}$$

where  $\epsilon \in (1, 1/2)$ .

Using the norm preservation theorem, prove that for any  $u, v \in \mathbb{R}^d$  such that  $||u|| \leq 1$  and  $||v|| \leq 1$ ,

$$\mathbb{P}\left(\left|u \cdot v - \phi\left(u\right) \cdot \phi\left(v\right)\right| \ge \epsilon\right) \le 4\exp\left(-\left(\epsilon^2 - \epsilon^3\right)m/4\right) \tag{2}$$

Proof.

# Problem 2: Locality-Sensitive Hashing (LSH) for Angle Similarity

Suppose our set of n points  $D = \{p_1, \ldots, p_n\}$  are vectors in d dimensions. Our problem is: given a query point q find a point  $p \in D$ , which has a small angle with q. Recall that the angle between two vectors a and b is  $\cos^{-1}\left(\frac{a \cdot b}{\|a\| \|b\|}\right)$ .

As doing this exactly may be computationally expensive, let us try to do this approximately with a fast algorithm. The approximate objective is as follows: suppose there exists a point  $p \in D$  which has cosine similarity larger than  $\theta$ , then our goal is return a point with cosine similarity greater than  $c\theta$ . As doing this exactly may be computationally expensive, let us try to do this approximately with a fast algorithm. The approximate objective is as

follows: suppose there exists a point  $p \in D$  which has cosine similarity larger than  $\theta$ , then our goal is return a point with cosine similarity greater than  $c\theta$ .

Let us try to do this with LSH. Let us consider the a family of hash functions, where  $h(p) = \text{sign}(u \cdot p)$  where we will sample u uniformly at random from a Gaussian (or from a unit sphere).

1. Provide an exact expression for  $\mathbb{P}\left(h\left(p\right)=h\left(p'\right)\right)$  based on some geometric relation between p and p'.

#### Solution

Define

$$angle (u, v) = \cos^{-1} \left( \frac{u \cdot v}{\|u\| \|v\|} \right), \tag{3}$$

which is the angle between two vectors.

Then,

$$\mathbb{P}\left(h\left(p\right) = h\left(p'\right)\right) = 1 - \frac{\operatorname{angle}\left(p, p'\right)}{\pi}.$$
(4)

2. Provide an expression for  $P_1$  and  $P_2$  in terms of  $\theta$  and  $c\theta$ . Note that since we want a large angle

3.