

Lecture 32

Binary Search Trees

FIT 1008
Introduction to Computer Science



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Objectives

- To understand **Binary Search Trees**
- **Implement** Binary Search Trees:
 - ➔ search
 - ➔ insert
- **Advantages and disadvantages** of Binary Search Trees over sorted lists.

insert

```
>>> a = dict()
>>> a[123465] = "Julian"
>>> a[133123] = "Nicole"
>>> a[982211] = "David"
>>>
>>> a
{123465: 'Julian', 133123: 'Nicole', 982211: 'David'}
>>>
>>>
>>> a[133123]
'Nicole'
```

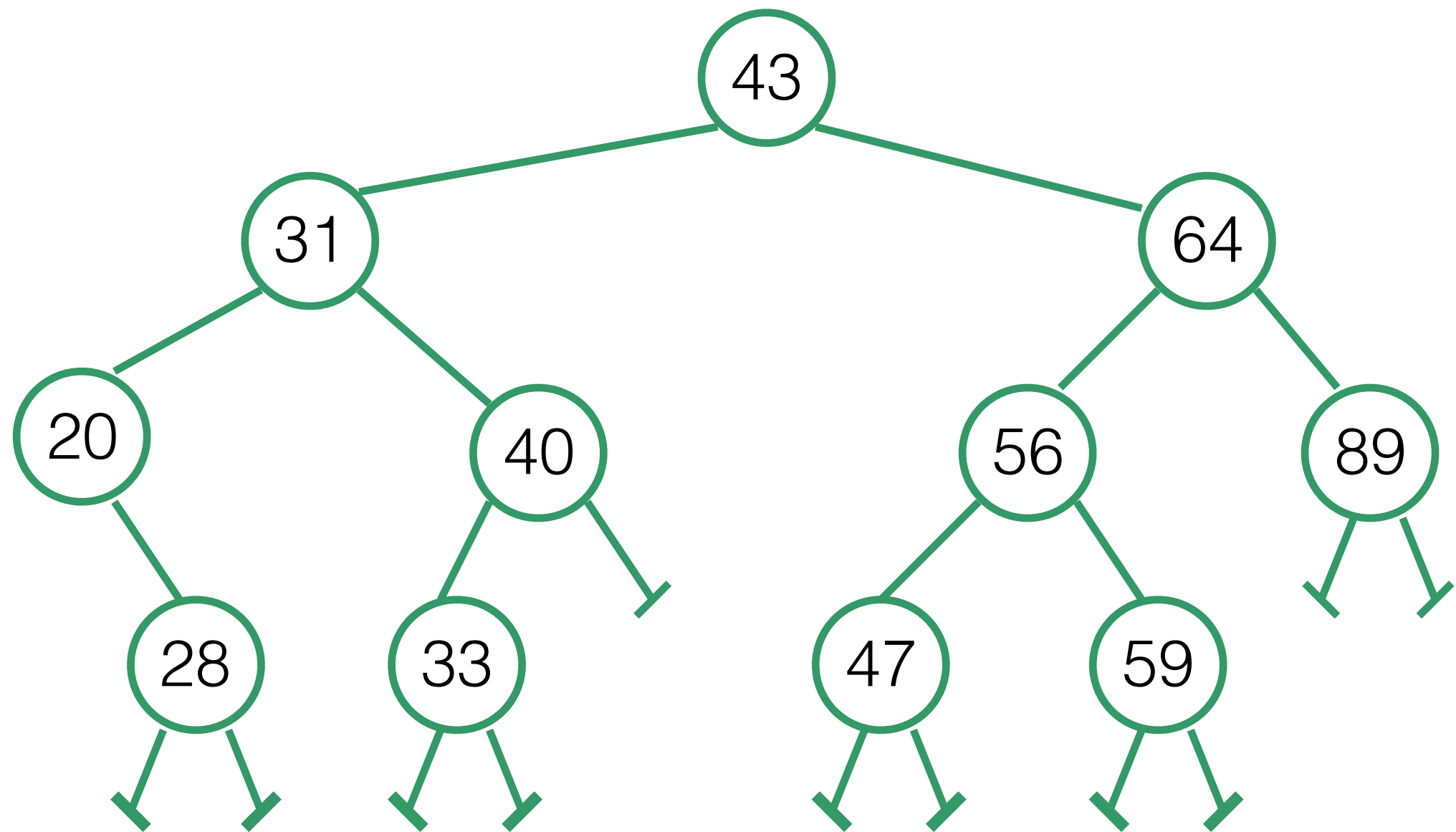
search

Python dictionaries are implemented using Hash Tables
You can also use a Binary Search Tree!

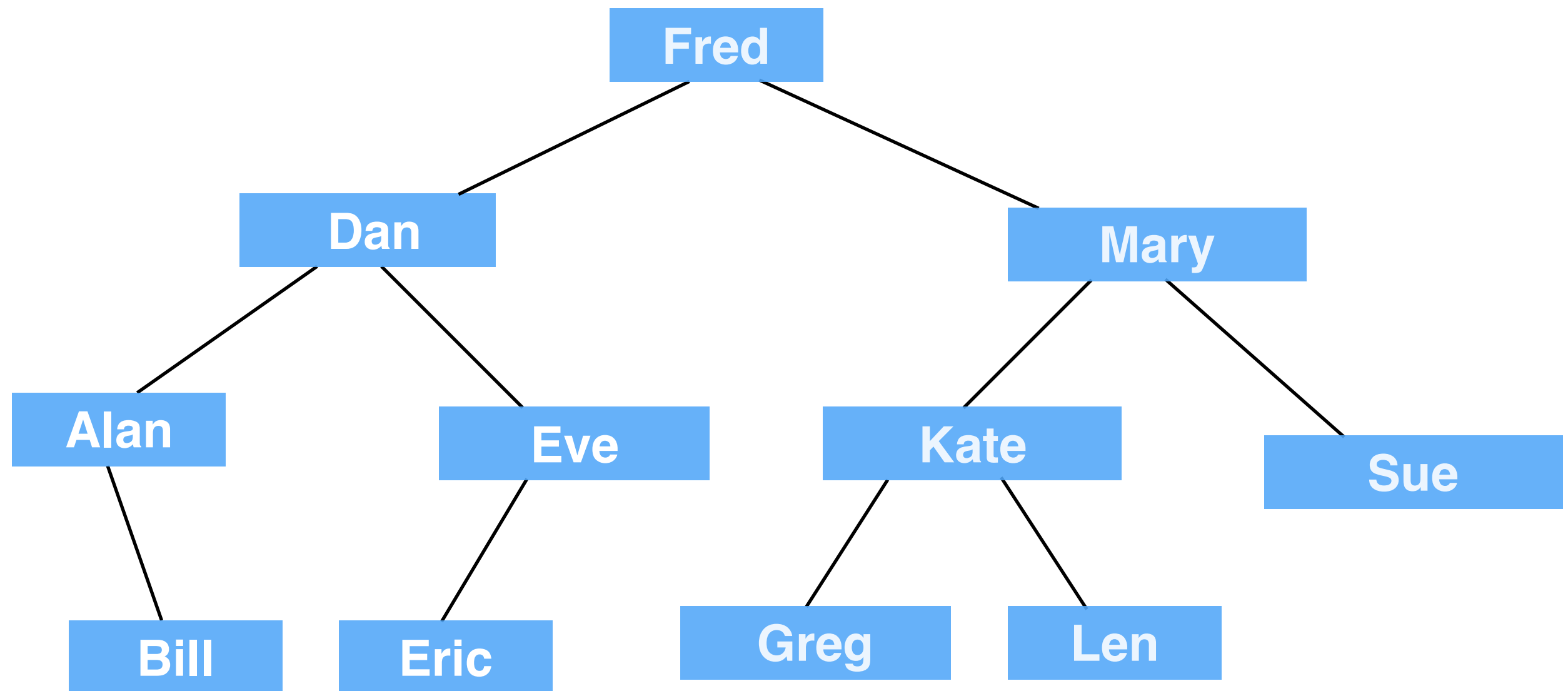
Binary Search Tree

A Binary Tree such that:

- Every node entry has a **key**
 - ➔ All **keys** in the **left subtree** of a node are **smaller than** the key of the node
 - ➔ All **keys** in the **right subtree** of a node are **greater than** the key of the node



key is an integer.



key is a string

(here not showing the associated items)

Why Binary Search Trees?

Sorted List Array-based:

- Good for search — $O(\log N)$ [binary search]
- Bad for inserting/deleting — $O(N)$ [shuffling things around]

Linked Sorted List:

- Good for inserting/deleting — $O(1)$ [modifying links]
- Bad for searching — $O(N)$ [steps through the list]

Binary Search Trees:

good for searching **and** good for inserting/deleting

```
class BinarySearchTreeNode:
    def __init__(self, key, item=None, left=None, right=None):
        self.key = key
        self.item = item
        self.left = left
        self.right = right
```

```
class BinarySearchTree:
    def __init__(self):
        self.root = None

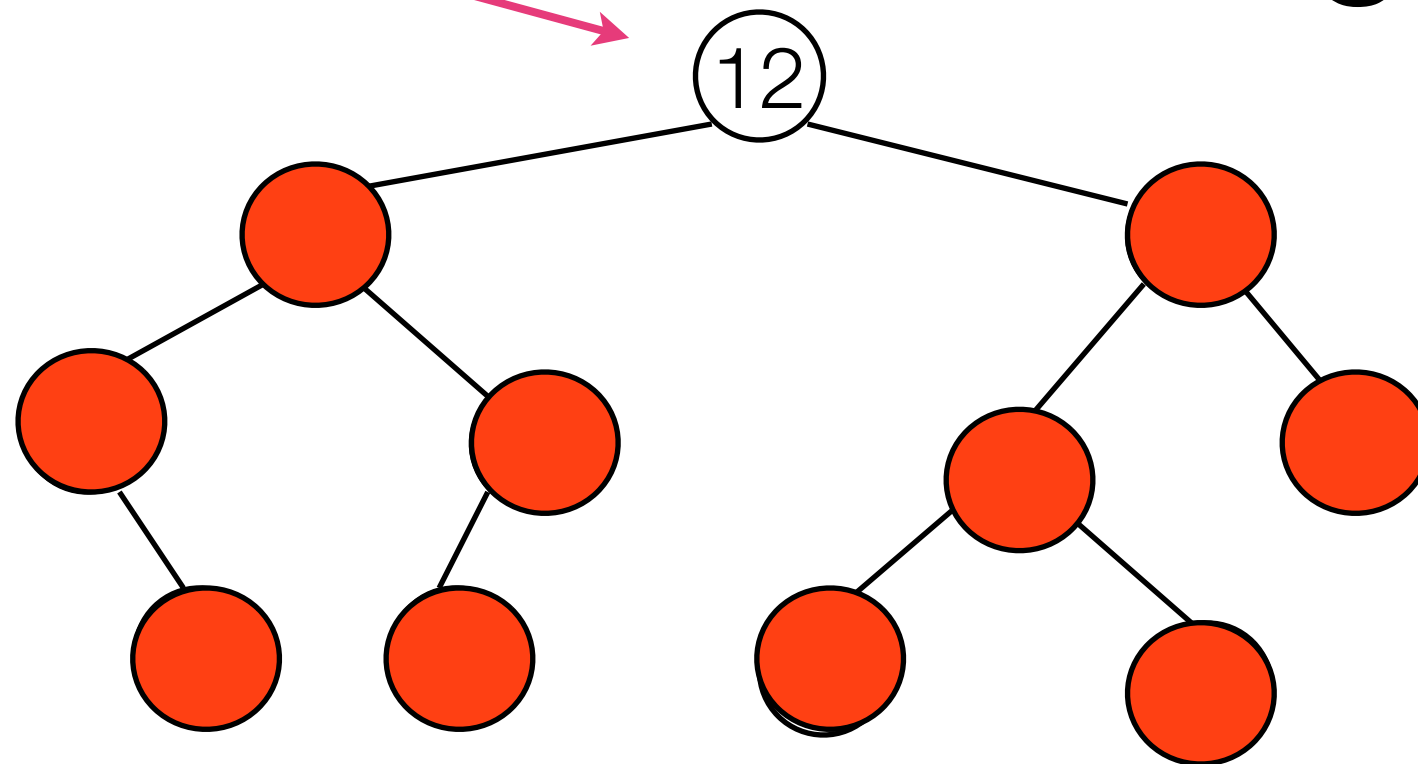
    def is_empty(self):
        return self.root is None
```


Motivation: Search

Searching

root

target = 13

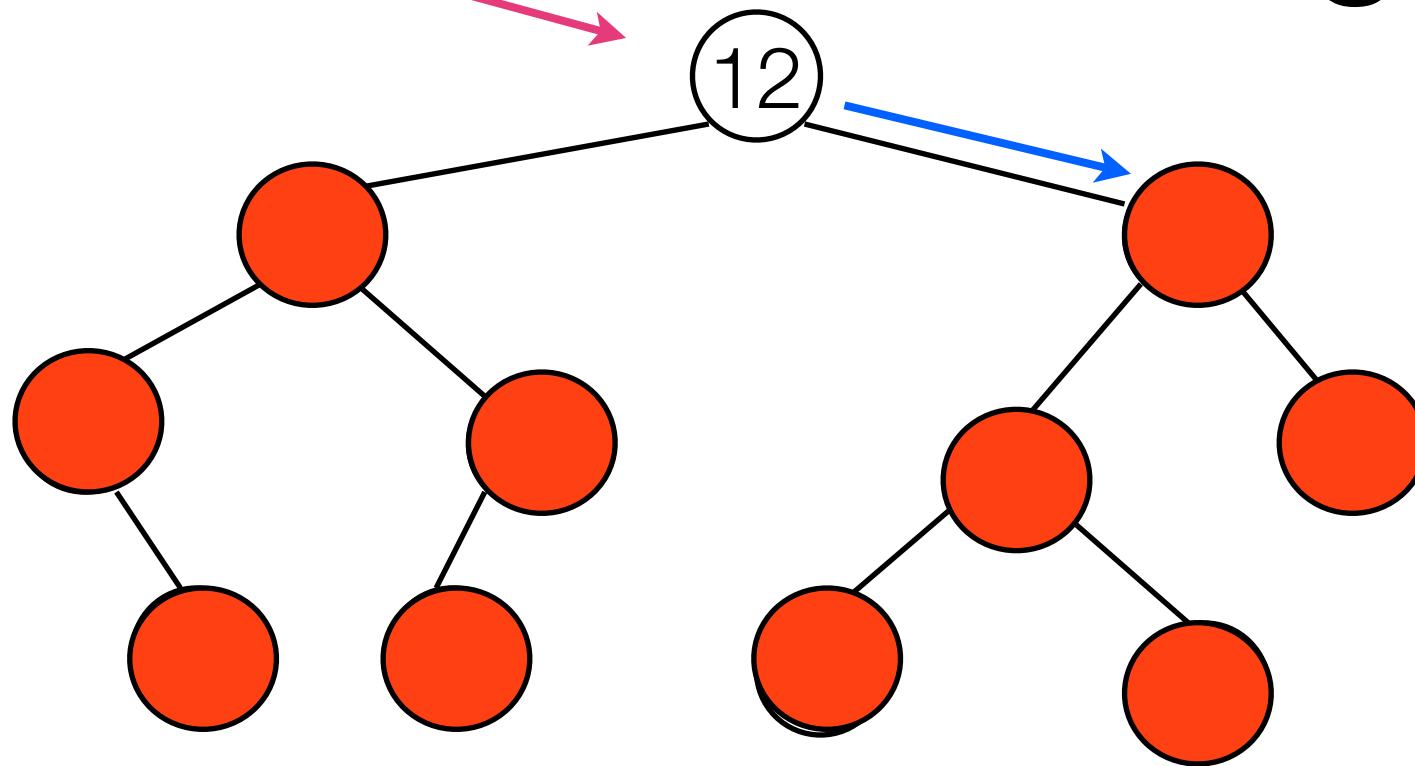


Only showing keys!

Searching

root

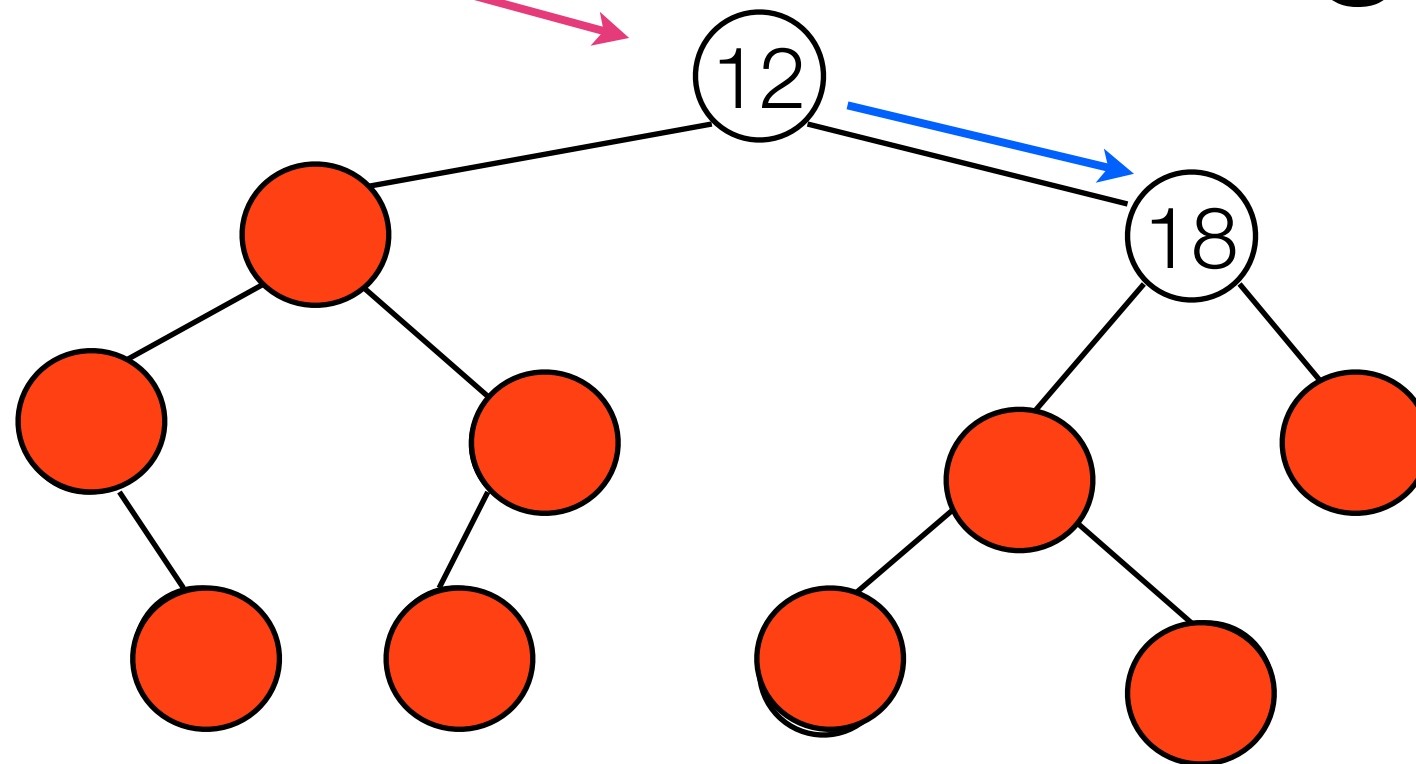
target = 13



Searching

root

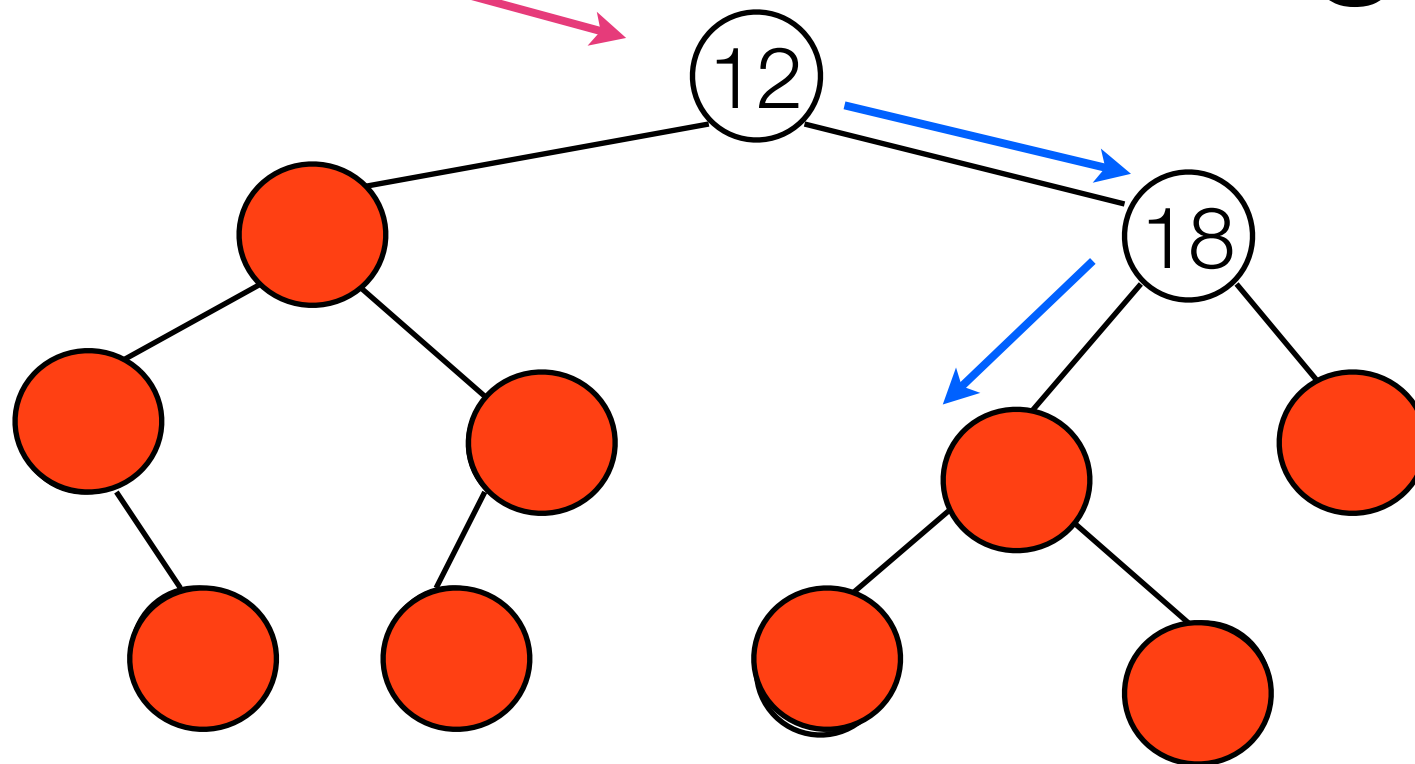
target = 13



Searching

root

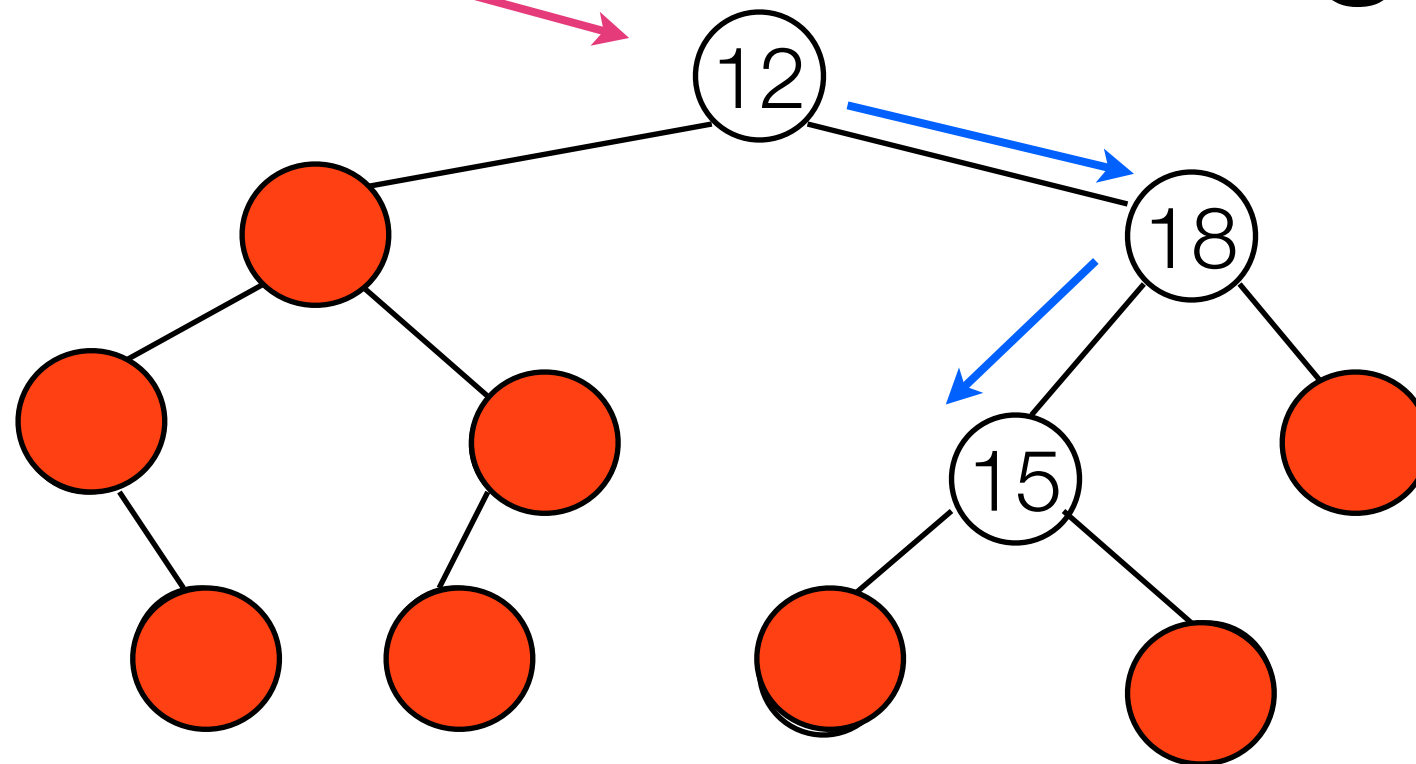
target = 13



Searching

root

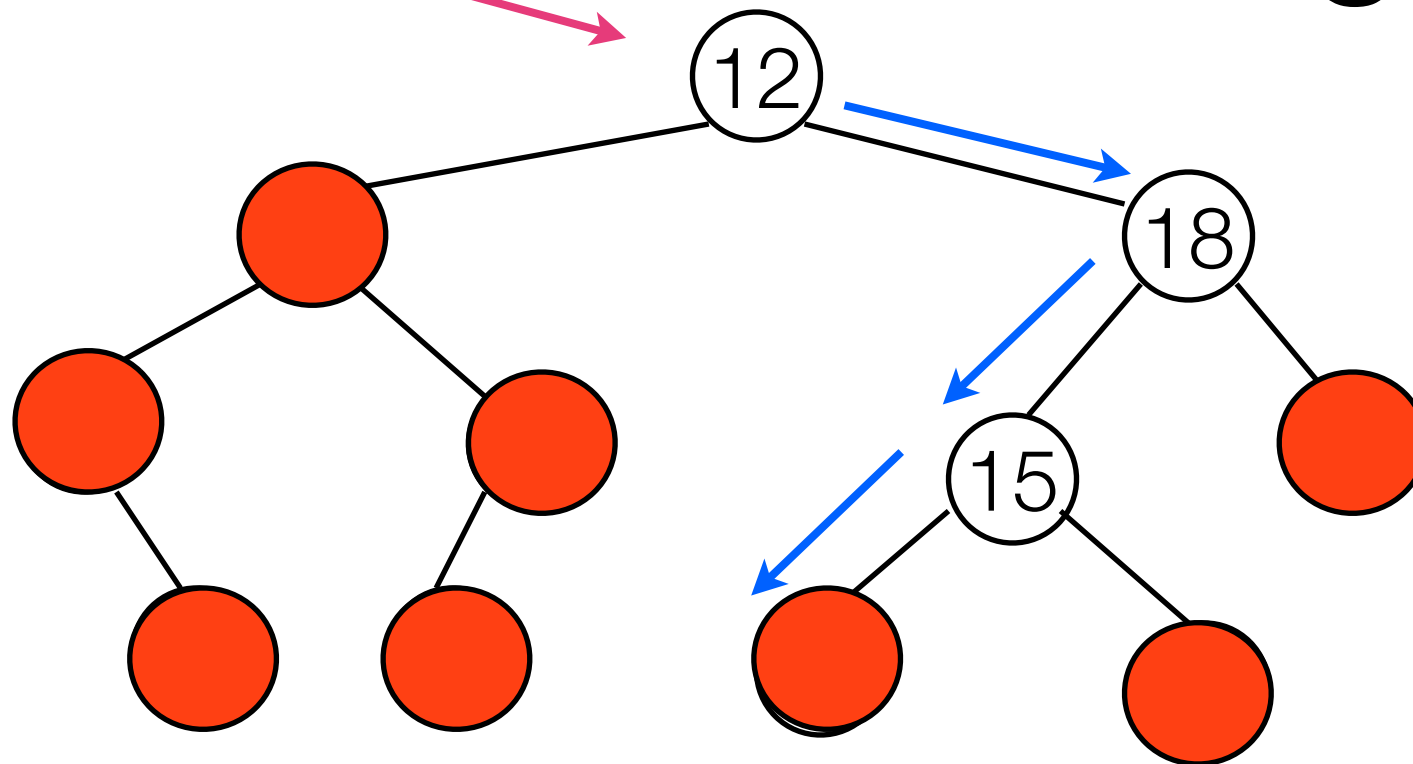
target = 13



Searching

root

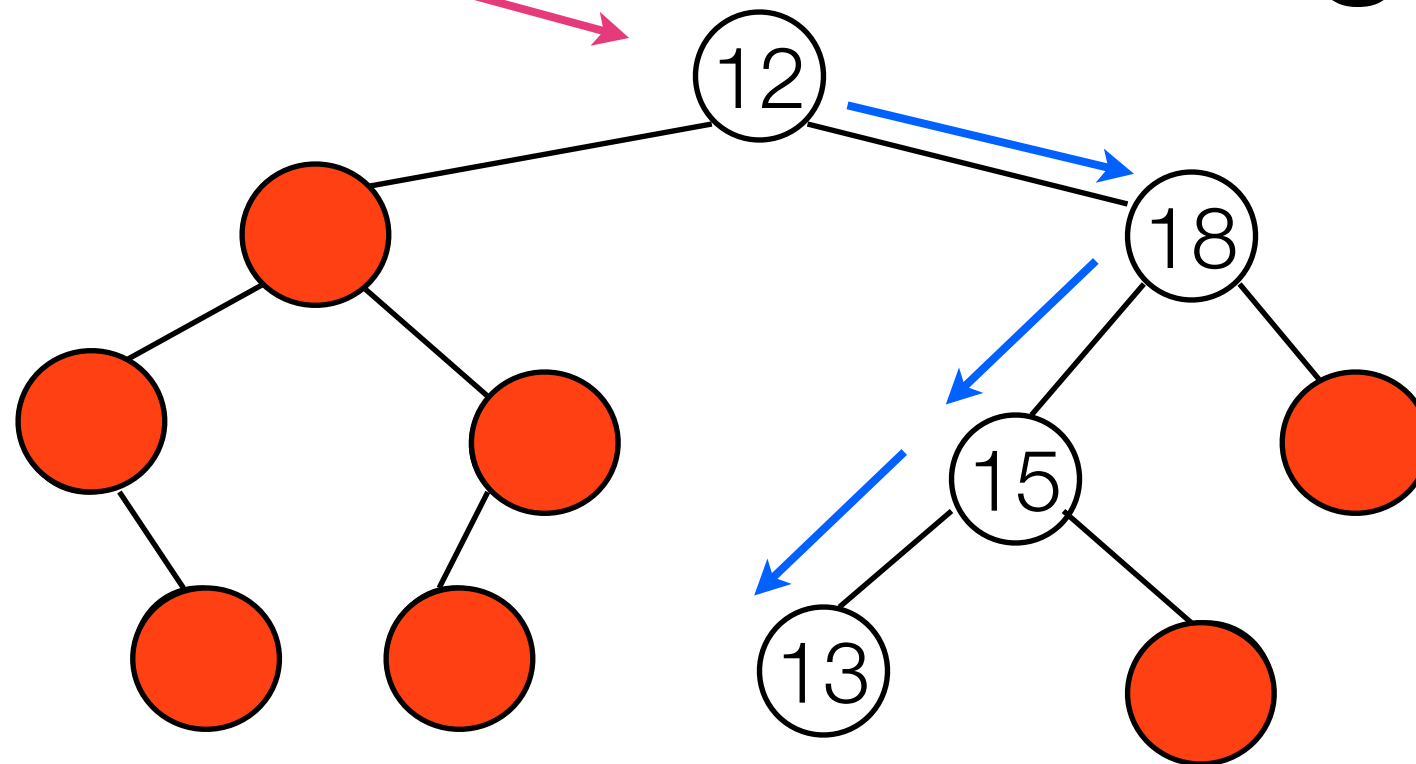
target = 13



Searching

root

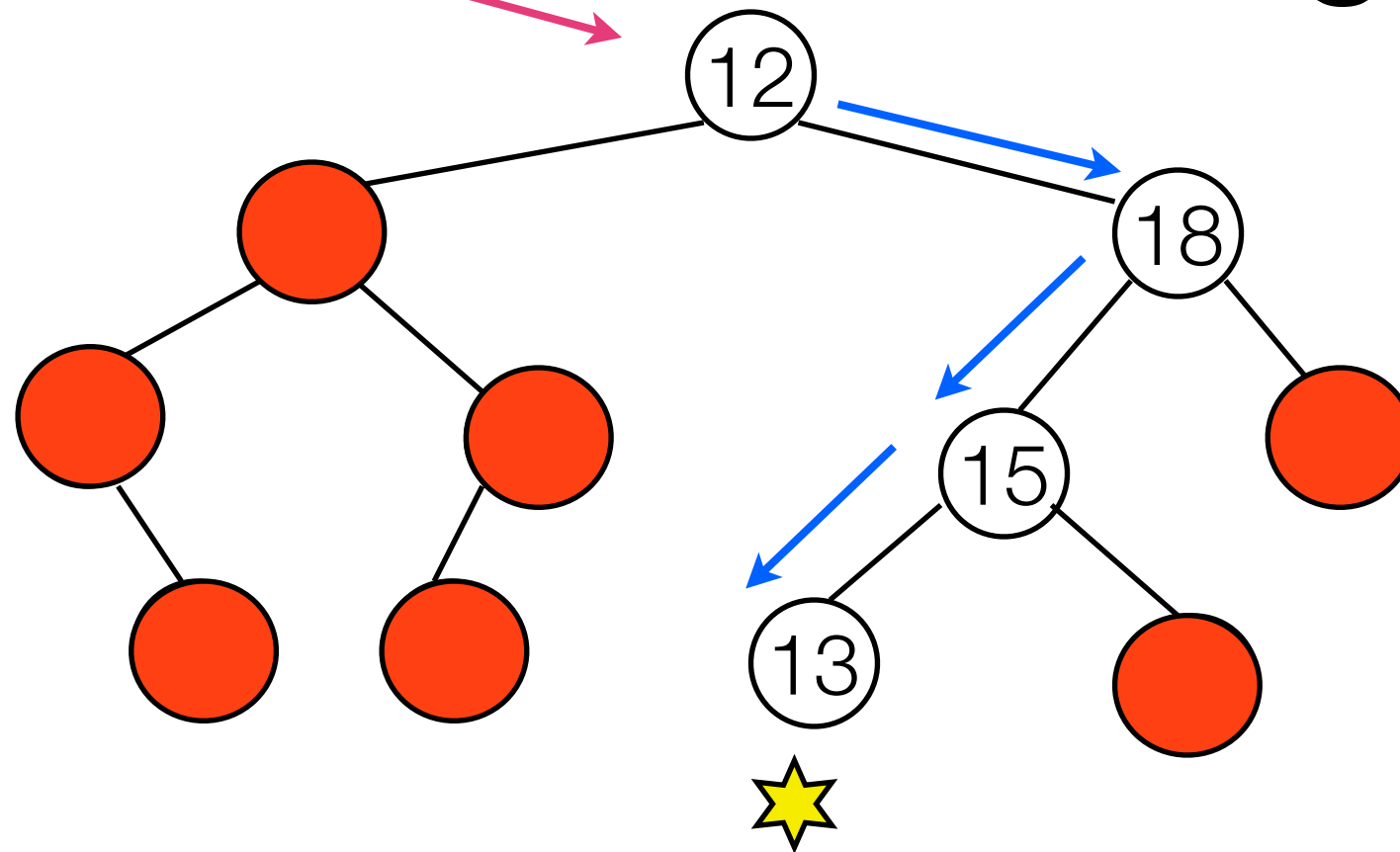
target = 13



Searching

root

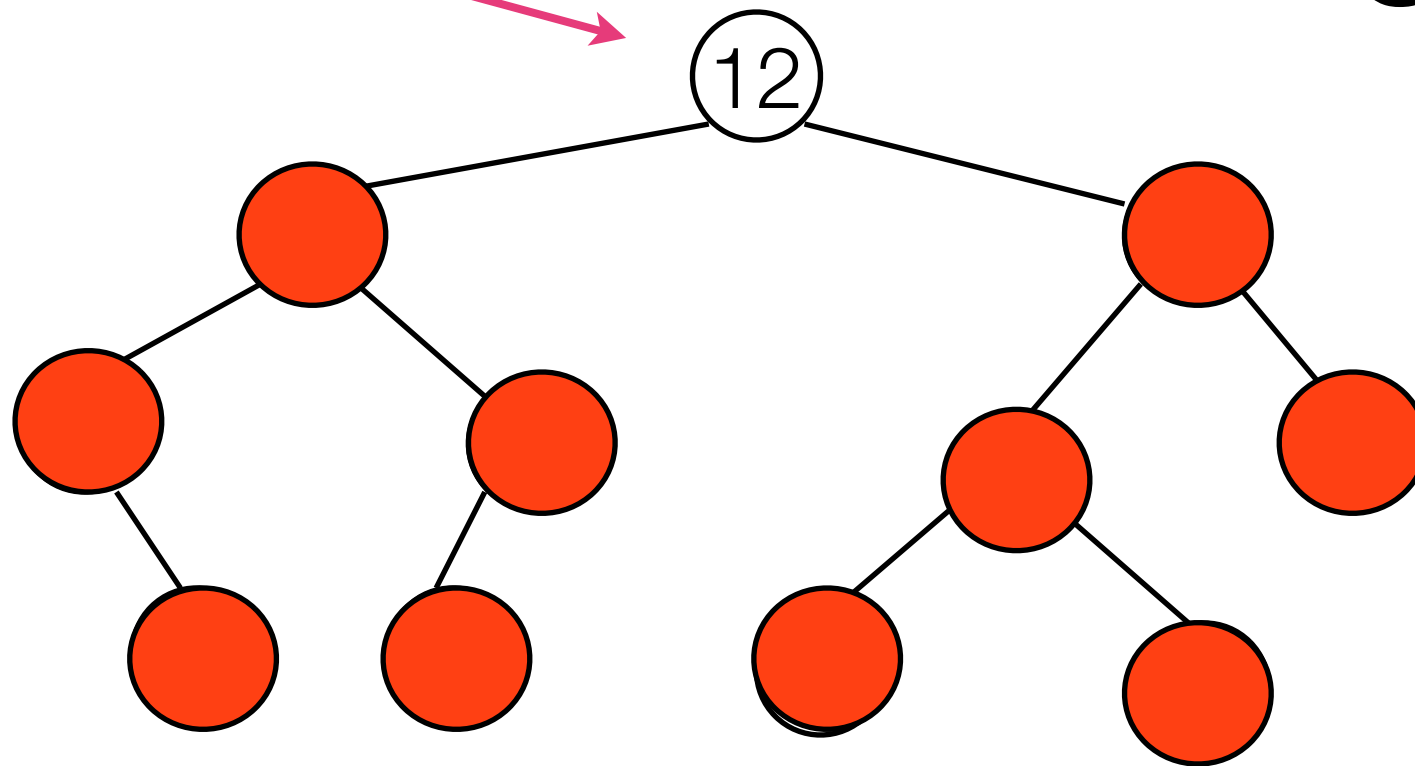
target = 13



Searching

root

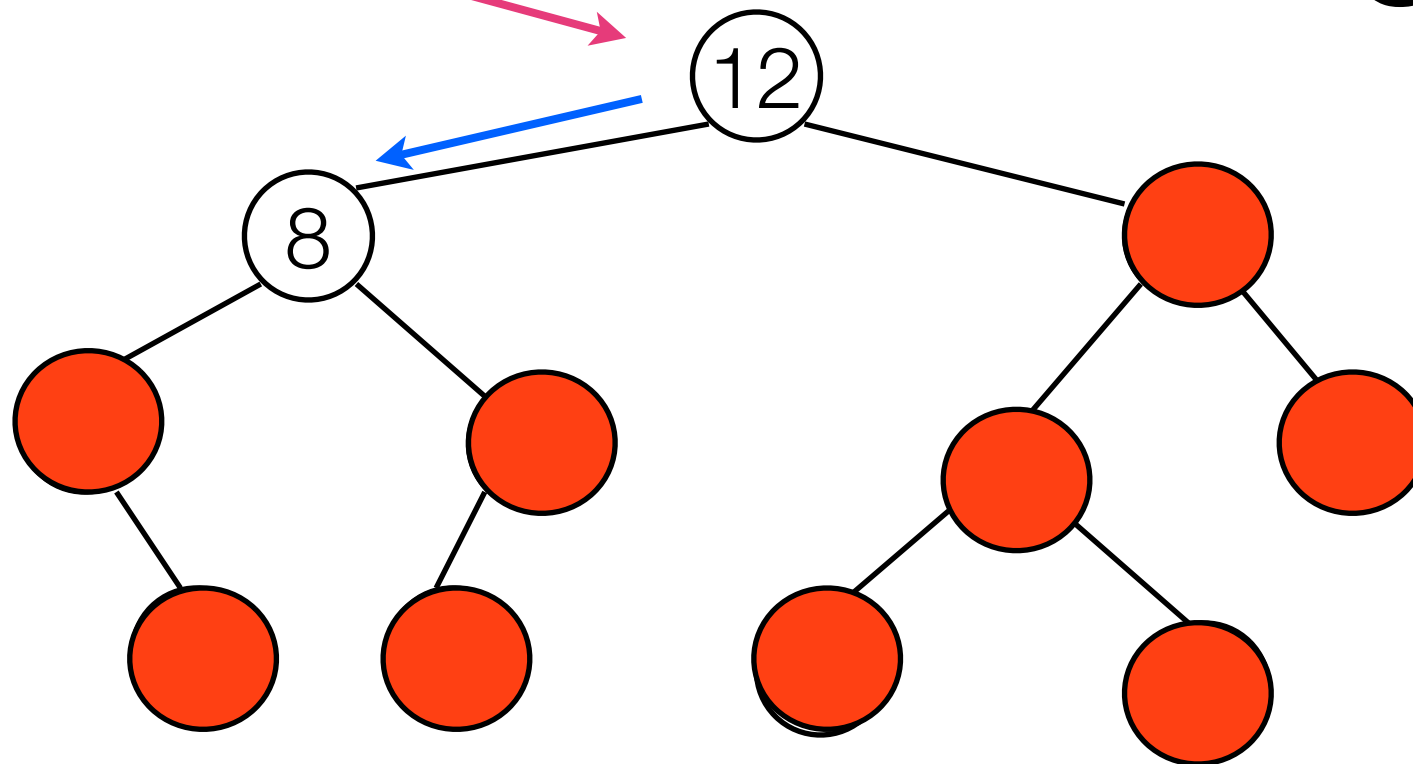
target = 2



Searching

root

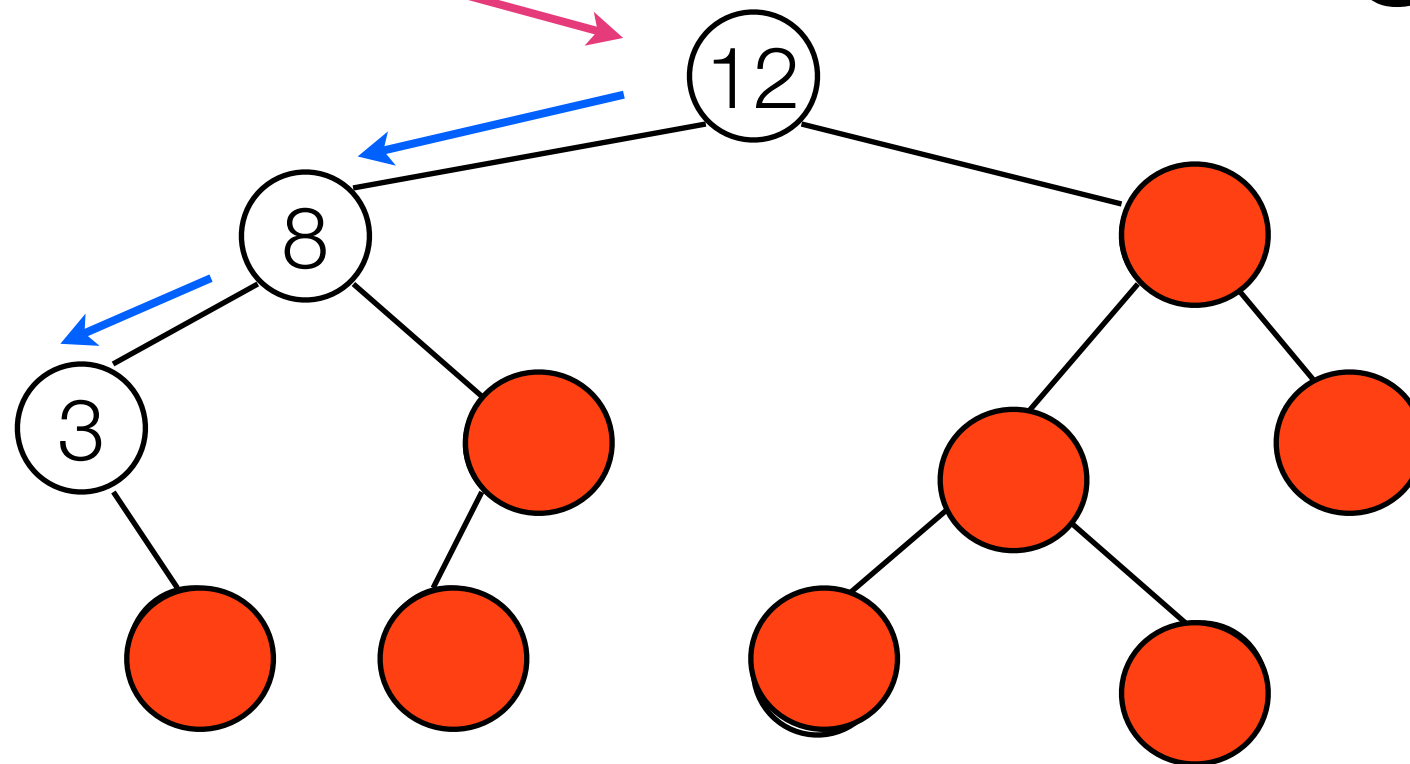
target = 2



Searching

root

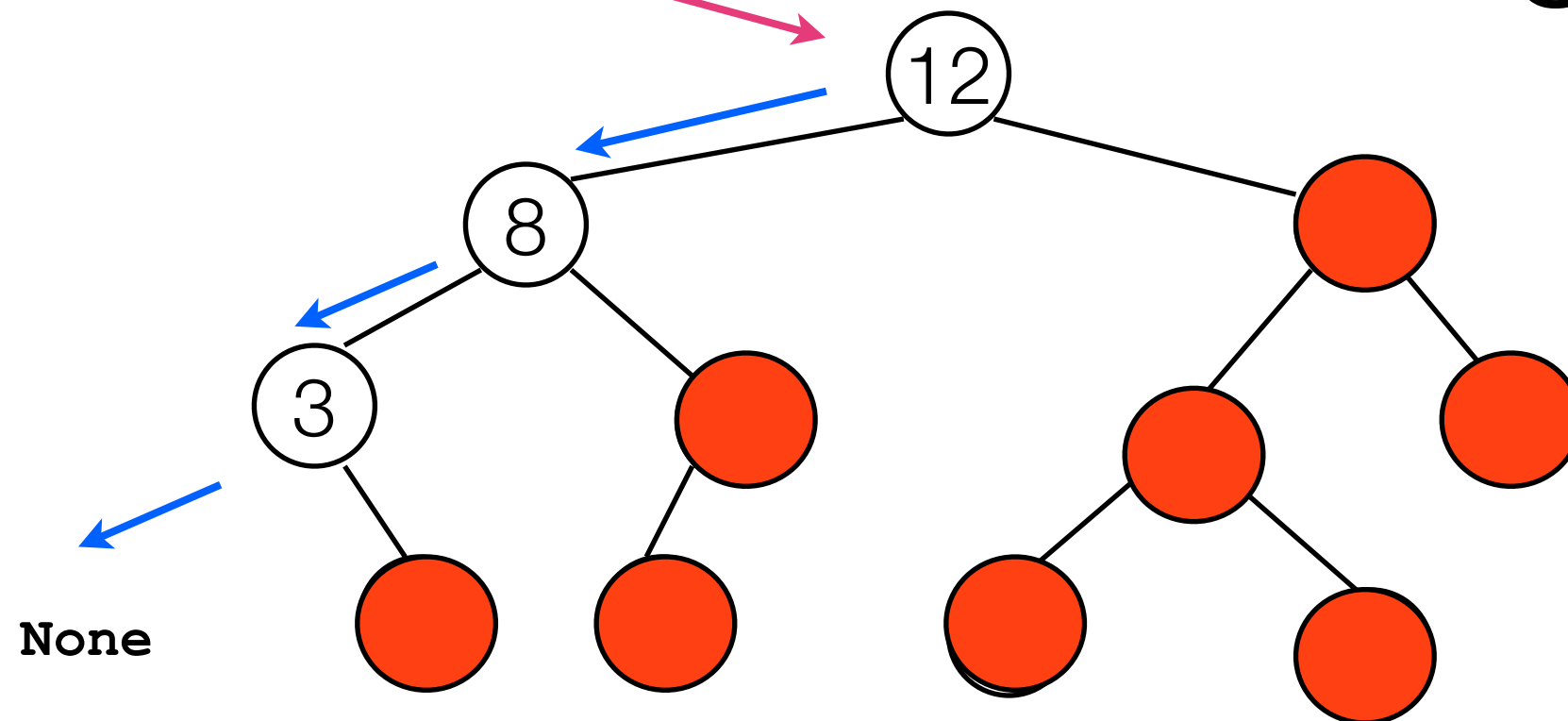
target = 2



Searching

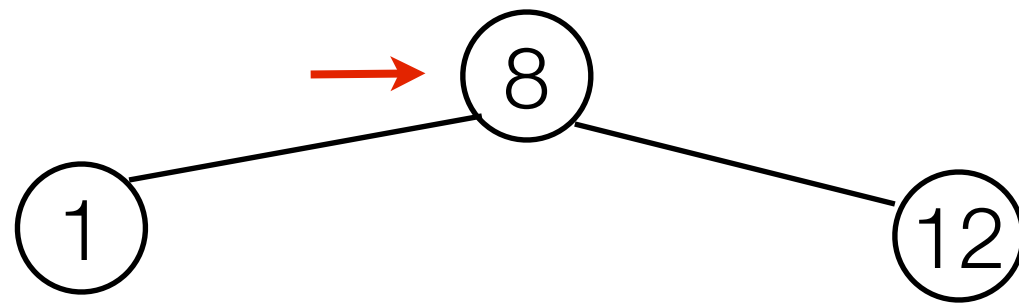
root

target = 2

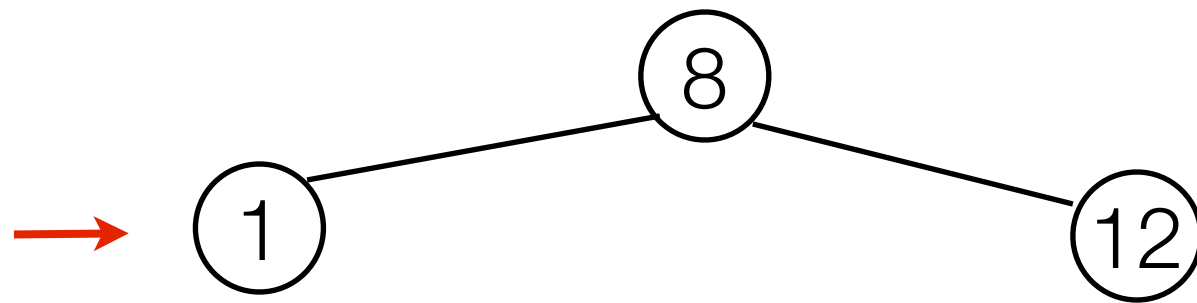


Insert

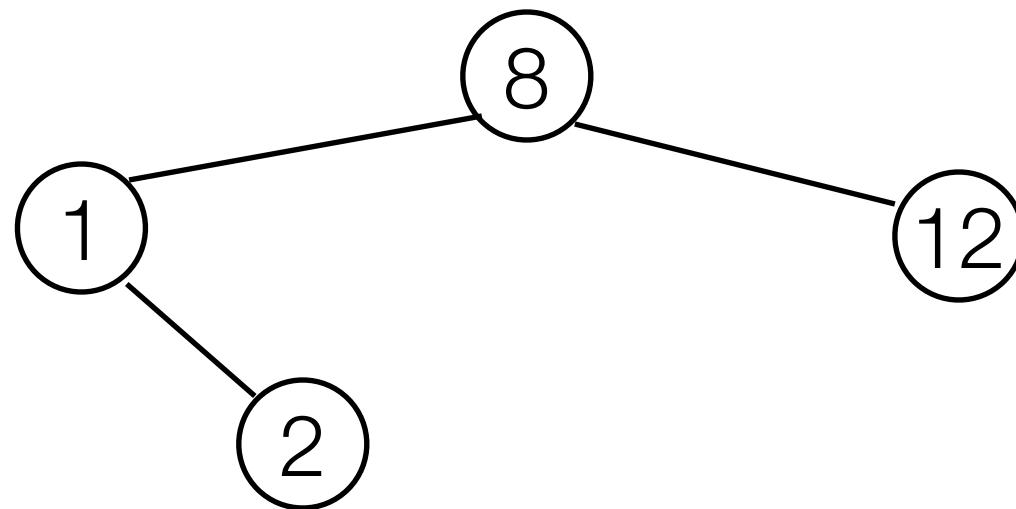
Insert 2



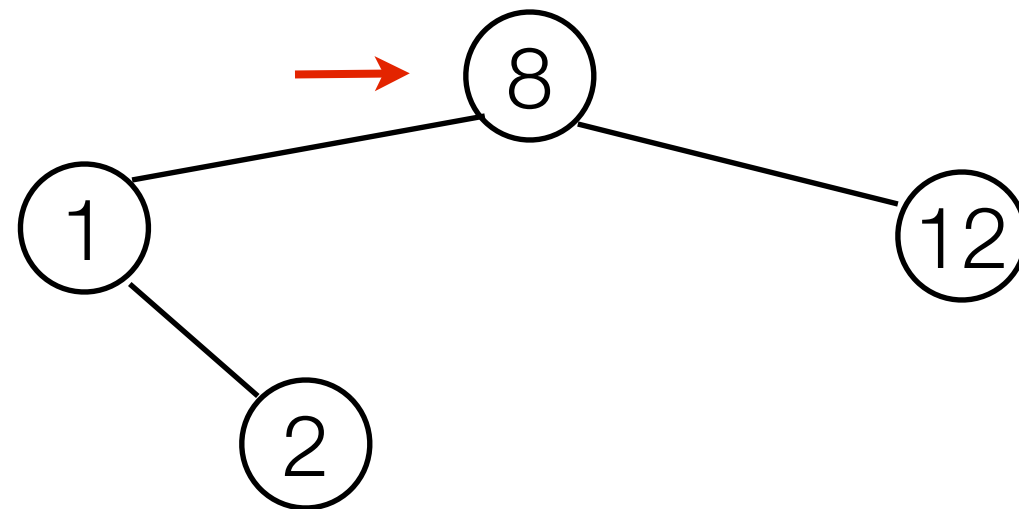
Insert 2



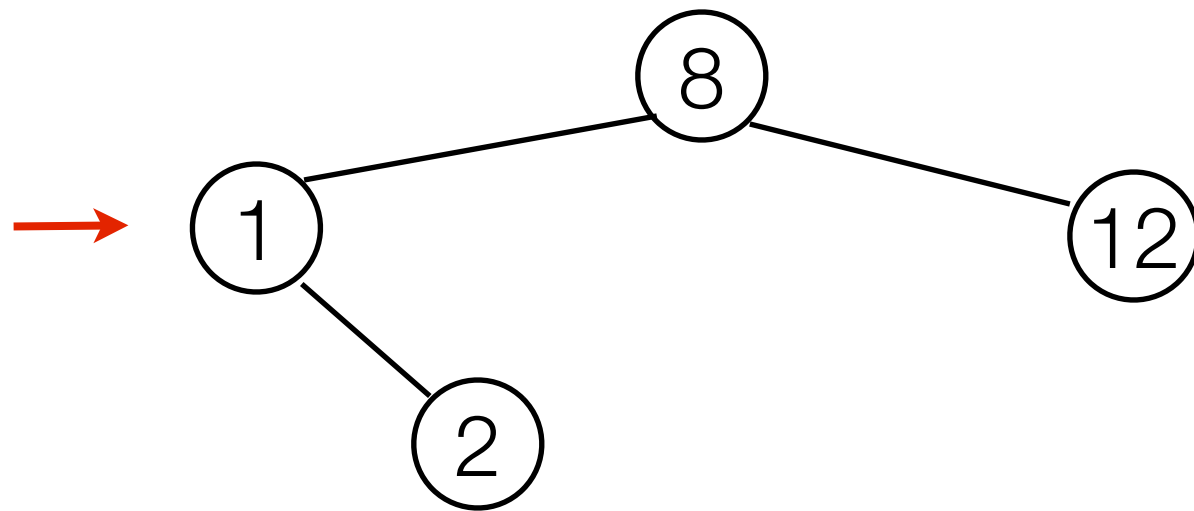
Insert 2



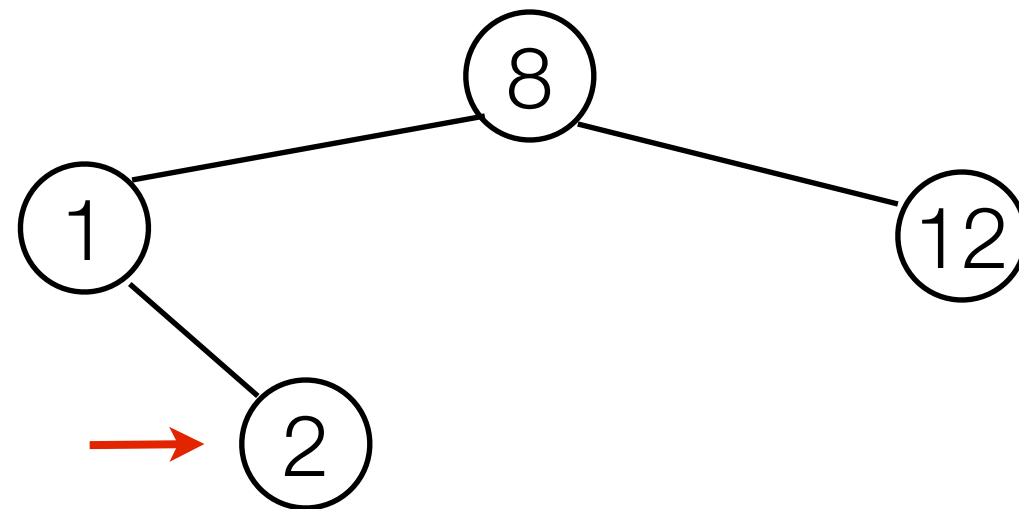
Insert 7



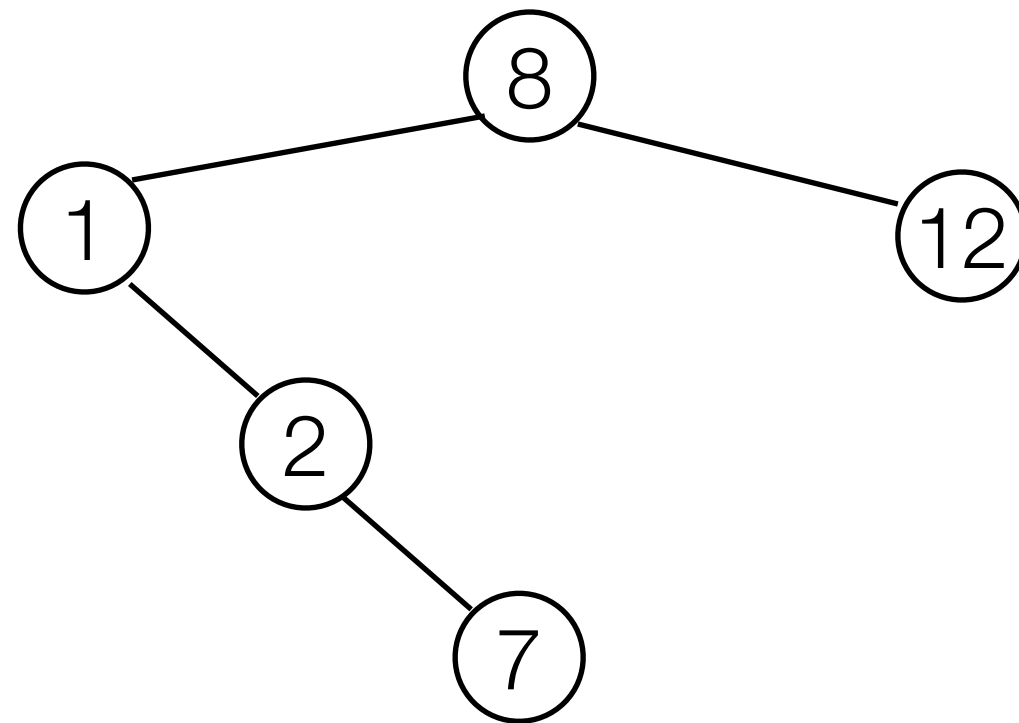
Insert 7



Insert 7



Insert 7



Our BST does not allow for duplicates, so we need to do something if we find the key in the tree...

Insert algorithm

Input:

key and associated value to insert.

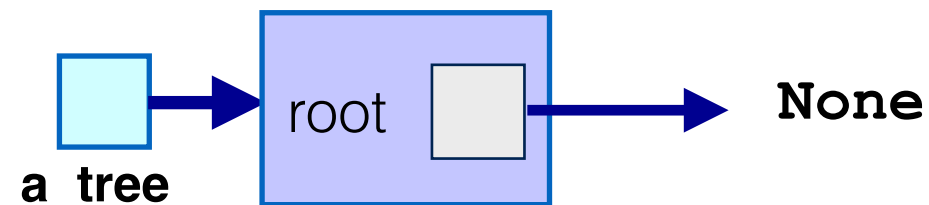
Idea:

Find the right spot (search) then create new node.

- Try to find the key...
 - ➔ **Found**? Raise an exception, keys must be unique....or replace value.
 - ➔ **Not found**? parent of **None** should be the parent of new node, which needs to be created.

```
def insert(self, key, item):  
    self.root = self._insert_aux(self.root, key, item)
```

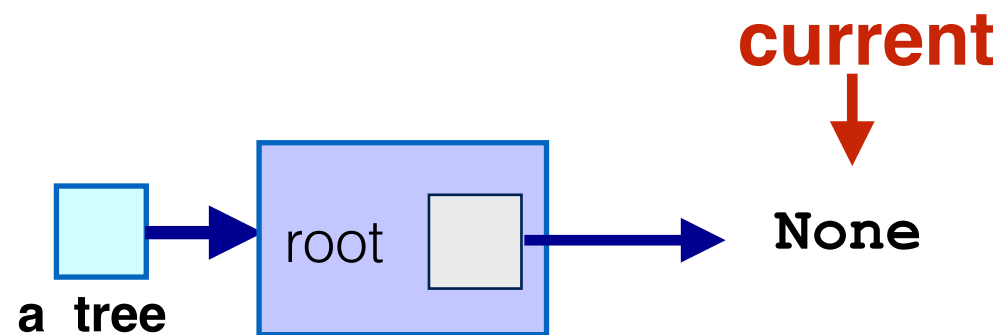
```
def _insert_aux(self, current, key, item):  
    if current is None: # base case: at the leaf  
        current = BinarySearchTreeNode(key, item)  
    elif key < current.key:  
        self._insert_aux(current.left, key, item)  
    elif key > current.key:  
        self._insert_aux(current.right, key, item)  
    else: # key == current.key  
        raise ValueError("Duplicate Item")
```



`a_tree.insert(57, "Coco")`

```
def insert(self, key, item):  
    self.root = self._insert_aux(self.root, key, item)
```

```
def _insert_aux(self, current, key, item):  
    if current is None: # base case: at the leaf  
        current = BinarySearchTreeNode(key, item)  
    elif key < current.key:  
        self._insert_aux(current.left, key, item)  
    elif key > current.key:  
        self._insert_aux(current.right, key, item)  
    else: # key == current.key  
        raise ValueError("Duplicate Item")
```



key → 57

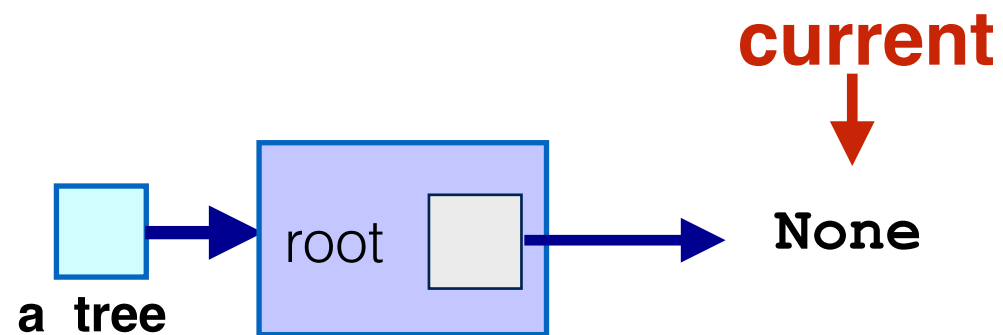
item → "Coco"

```
a_tree.insert(57, "Coco")
```



```
def insert(self, key, item):  
    self.root = self._insert_aux(self.root, key, item)
```

```
def _insert_aux(self, current, key, item):  
    if current is None: # base case: at the leaf  
        current = BinarySearchTreeNode(key, item)  
    elif key < current.key:  
        self._insert_aux(current.left, key, item)  
    elif key > current.key:  
        self._insert_aux(current.right, key, item)  
    else: # key == current.key  
        raise ValueError("Duplicate Item")
```



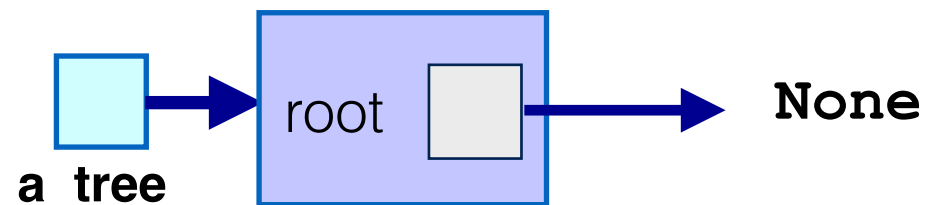
key → 57

item → "Coco"

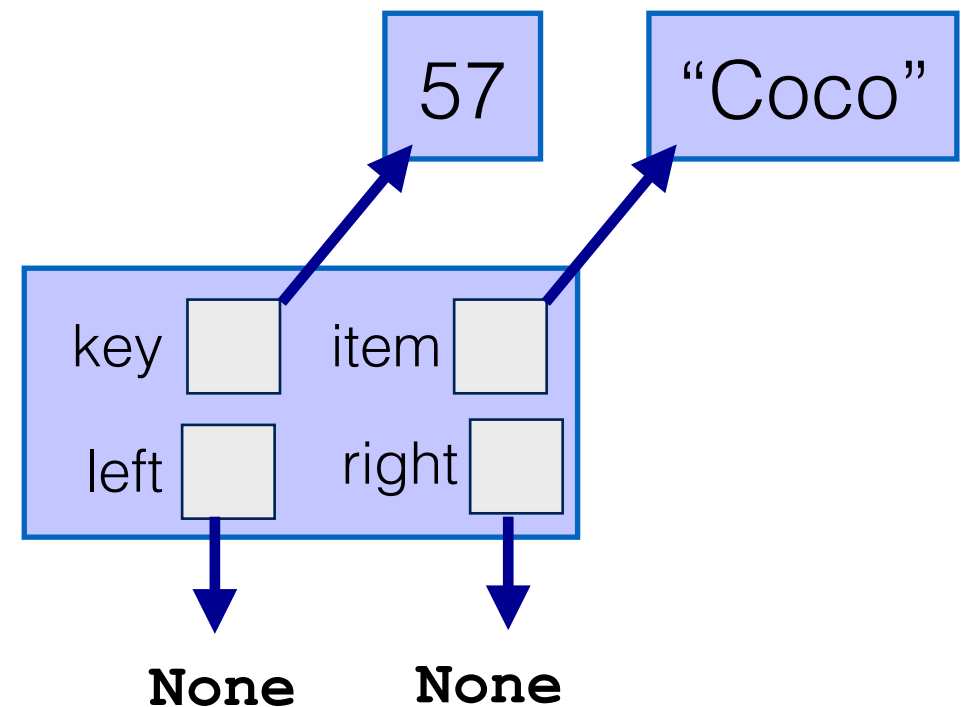
```
a_tree.insert(57, "Coco")
```

```
def insert(self, key, item):
    self.root = self._insert_aux(self.root, key, item)
```

```
def _insert_aux(self, current, key, item):
    if current is None: # base case: at the leaf
        current = BinarySearchTreeNode(key, item)
    elif key < current.key:
        self._insert_aux(current.left, key, item)
    elif key > current.key:
        self._insert_aux(current.right, key, item)
    else: # key == current.key
        raise ValueError("Duplicate Item")
```



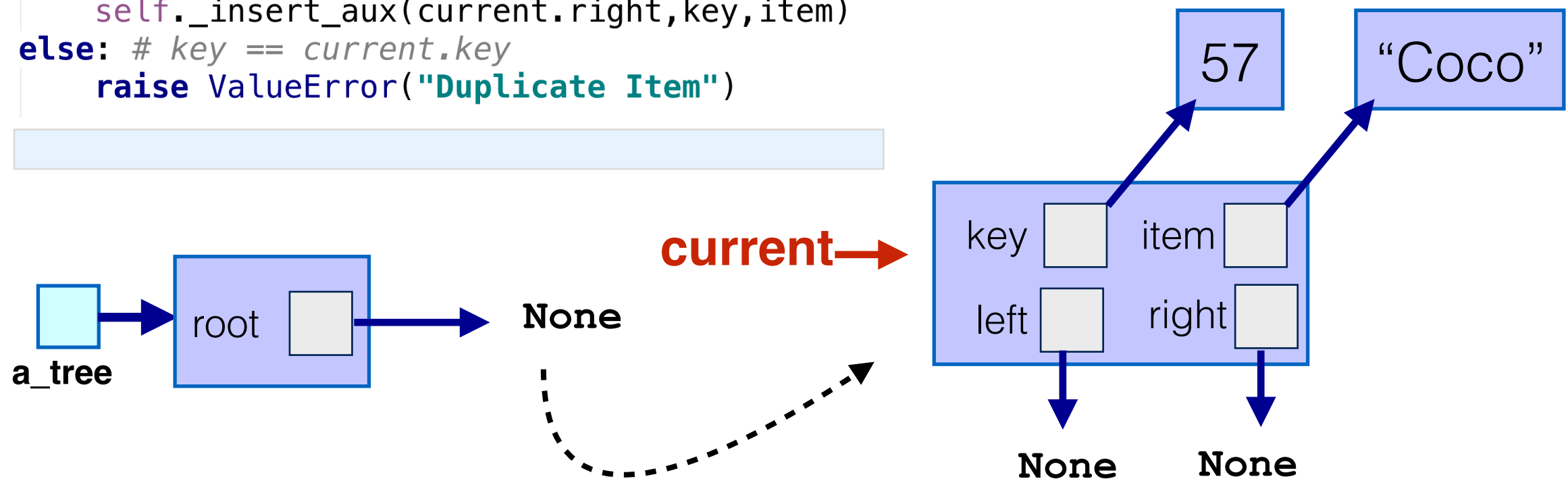
current →



`a_tree.insert(57, "Coco")`

```
def insert(self, key, item):
    self.root = self._insert_aux(self.root, key, item)
```

```
def _insert_aux(self, current, key, item):
    if current is None: # base case: at the leaf
        current = BinarySearchTreeNode(key, item)
    elif key < current.key:
        self._insert_aux(current.left, key, item)
    elif key > current.key:
        self._insert_aux(current.right, key, item)
    else: # key == current.key
        raise ValueError("Duplicate Item")
```



current needs to be returned!

```
def insert(self, key, item):  
    self.root = self._insert_aux(self.root, key, item)
```

```
def _insert_aux(self, current, key, item):  
    if current is None: # base case: at the leaf  
        current = BinarySearchTreeNode(key, item)  
    elif key < current.key:  
        current.left = self._insert_aux(current.left, key, item)  
    elif key > current.key:  
        current.right = self._insert_aux(current.right, key, item)  
    else: # key == current.key  
        raise ValueError("Duplicate Item")  
    return current
```

```
def insert(self, key, item):  
    self.root = self._insert_aux(self.root, key, item)
```

```
def _insert_aux(self, current, key, item):  
    if current is None: # base case: at the leaf  
        current = BinarySearchTreeNode(key, item)  
    elif key < current.key:  
        current.left = self._insert_aux(current.left, key, item)  
    elif key > current.key:  
        current.right = self._insert_aux(current.right, key, item)  
    else: # key == current.key  
        raise ValueError("Duplicate Item")  
    return current
```

__setitem__

```
def __setitem__(self, key, item):  
    self.root = self._setitem_aux_(self.root, key, item)  
  
def _setitem_aux_(self, current, key, item):  
    if current is None: # base case: at the leaf  
        current = BinarySearchTreeNode(key, item)  
    elif key < current.key:  
        current.left = self._setitem_aux_(current.left, key, item)  
    elif key > current.key:  
        current.right = self._setitem_aux_(current.right, key, item)  
    else: # key == current.key  
        current.item = item  
    return current
```

Search algorithm

- If we reach an empty node, item is not there...
return **False**.
- Else, if target key is equal to the current node's key,
return **True**
- Else, if target key is **less** than current node's key,
search the left sub-tree
- Else, if target key is **greater** than current node's
key, **search the right sub-tree**

search can be implemented by
__contains__
and
__getitem__


```
def __contains__(self, key):  
    return self._contains_aux(key, self.root)  
  
def _contains_aux(self, key, current_node):  
    if current_node is None: # base case  
        return False  
    elif key == current_node.key:  
        return True  
    elif key < current_node.key:  
        return self._contains_aux(key, current_node.left)  
    elif key > current_node.key:  
        return self._contains_aux(key, current_node.right)
```

Keys implement “rich comparison”

```
object.__lt__(self, other)
object.__le__(self, other)
object.__eq__(self, other)
object.__ne__(self, other)
object.__gt__(self, other)
object.__ge__(self, other)
```

These are the so-called “rich comparison” methods. The correspondence between operator symbols and method names is as follows: `x < y` calls `x.__lt__(y)`, `x <= y` calls `x.__le__(y)`, `x == y` calls `x.__eq__(y)`, `x != y` calls `x.__ne__(y)`, `x > y` calls `x.__gt__(y)`, and `x >= y` calls `x.__ge__(y)`.

<https://docs.python.org/3/reference/datamodel.html>

we want to get the item associated to a key...

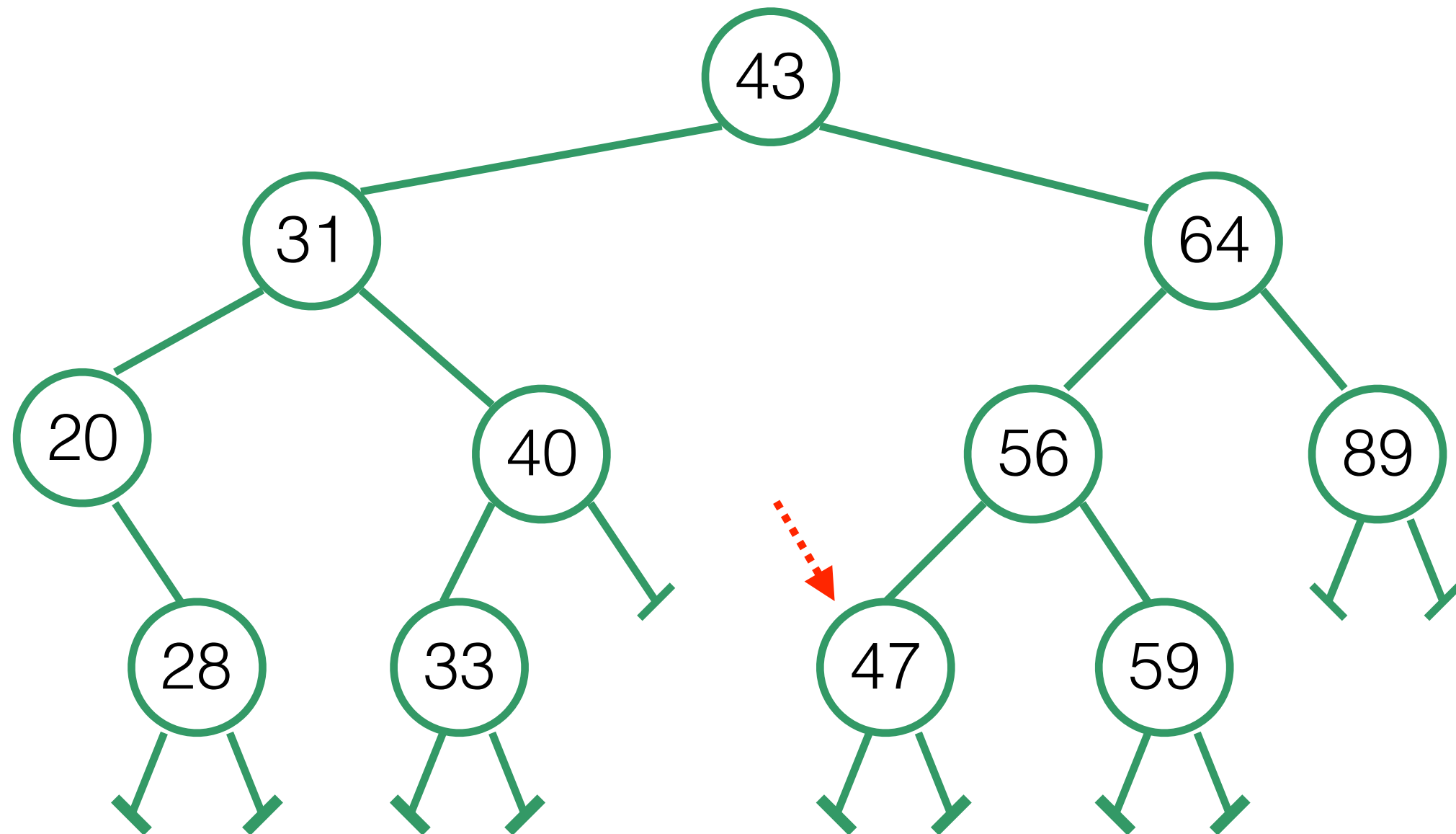
`__getitem__`

__getitem__

```
def __getitem__(self, key):  
    return self._getitem_aux(self.root, key)  
  
def _getitem_aux(self, current, key):  
    if current is None: # base case: empty  
        raise KeyError("Key not found")  
    elif key == current.key: # base case: found  
        return current.item  
    elif key < current.key:  
        return self._getitem_aux(current.left, key)  
    else: # key > current.key  
        return self._getitem_aux(current.right, key)
```

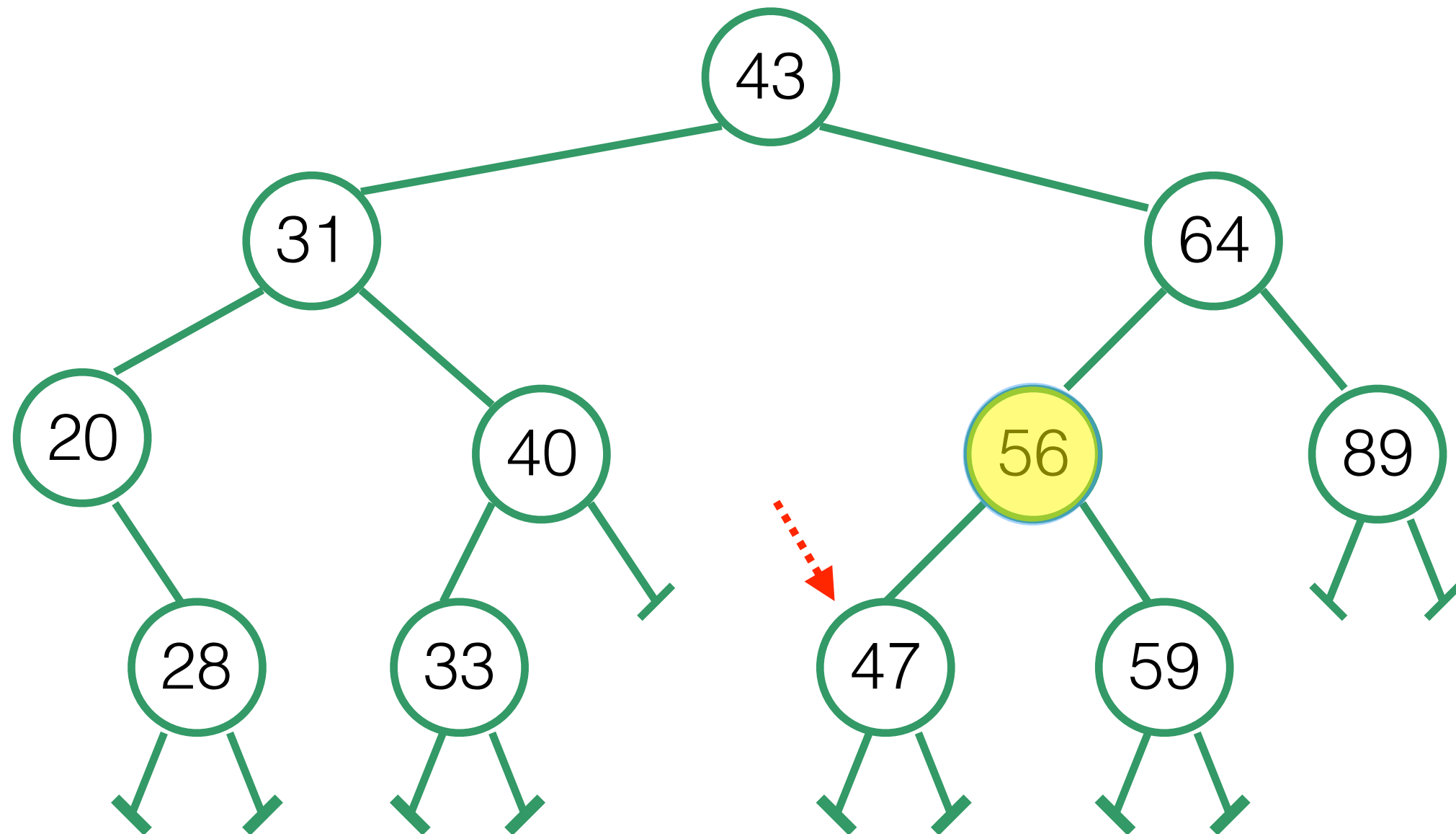
Delete

Delete 47



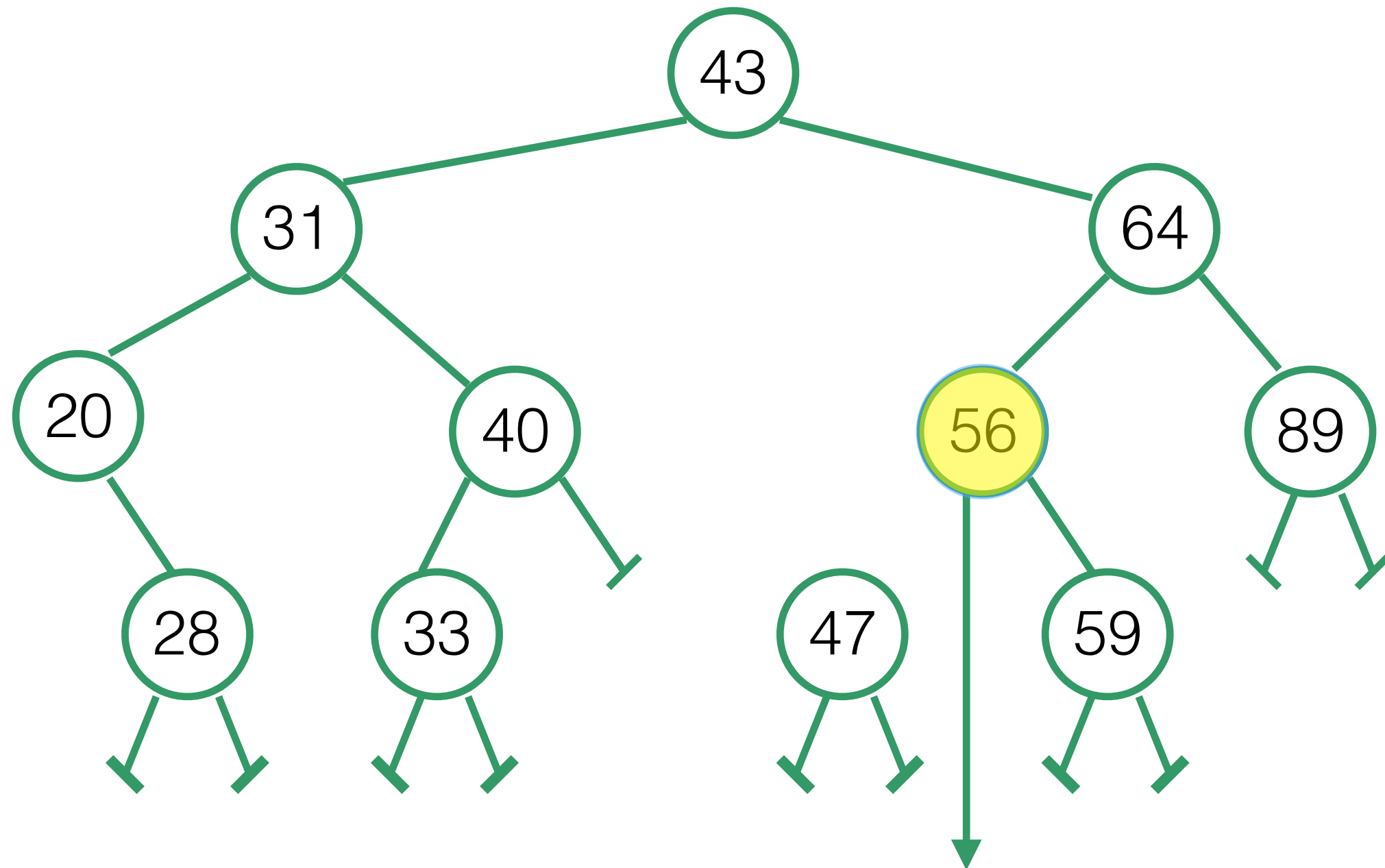
only showing **key**

Delete 47



only showing **key**

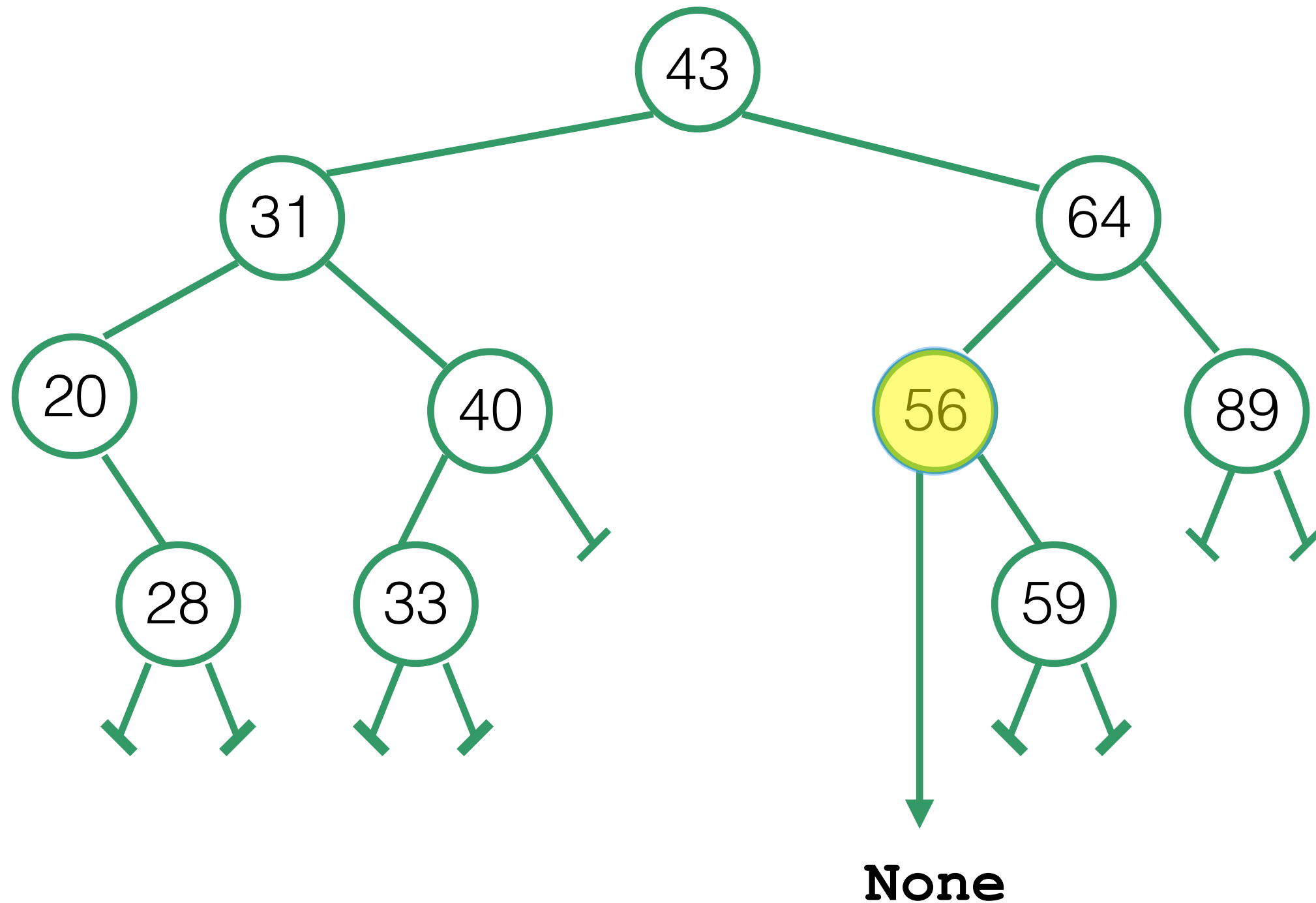
Delete 47



None

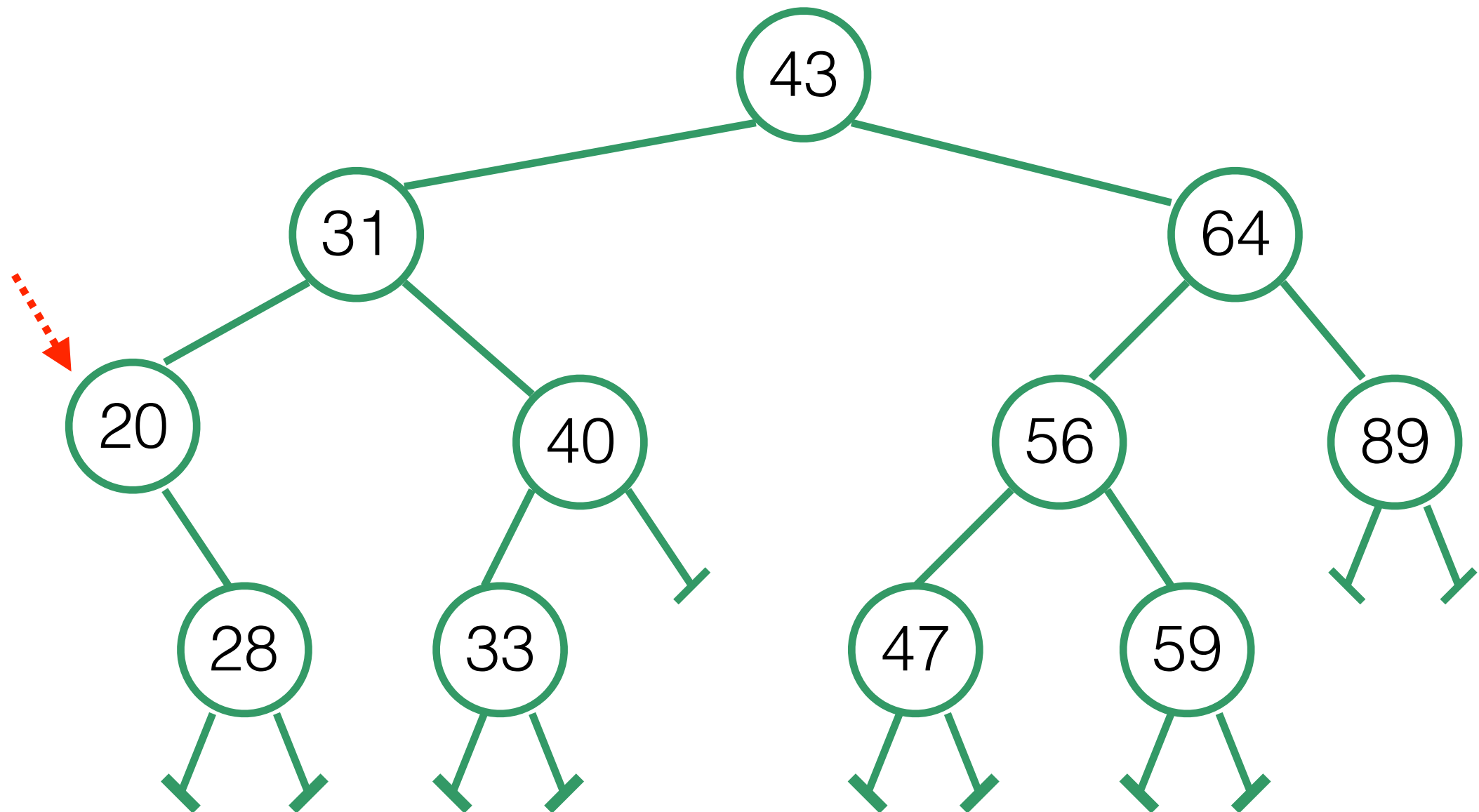
only showing **key**

Delete 47



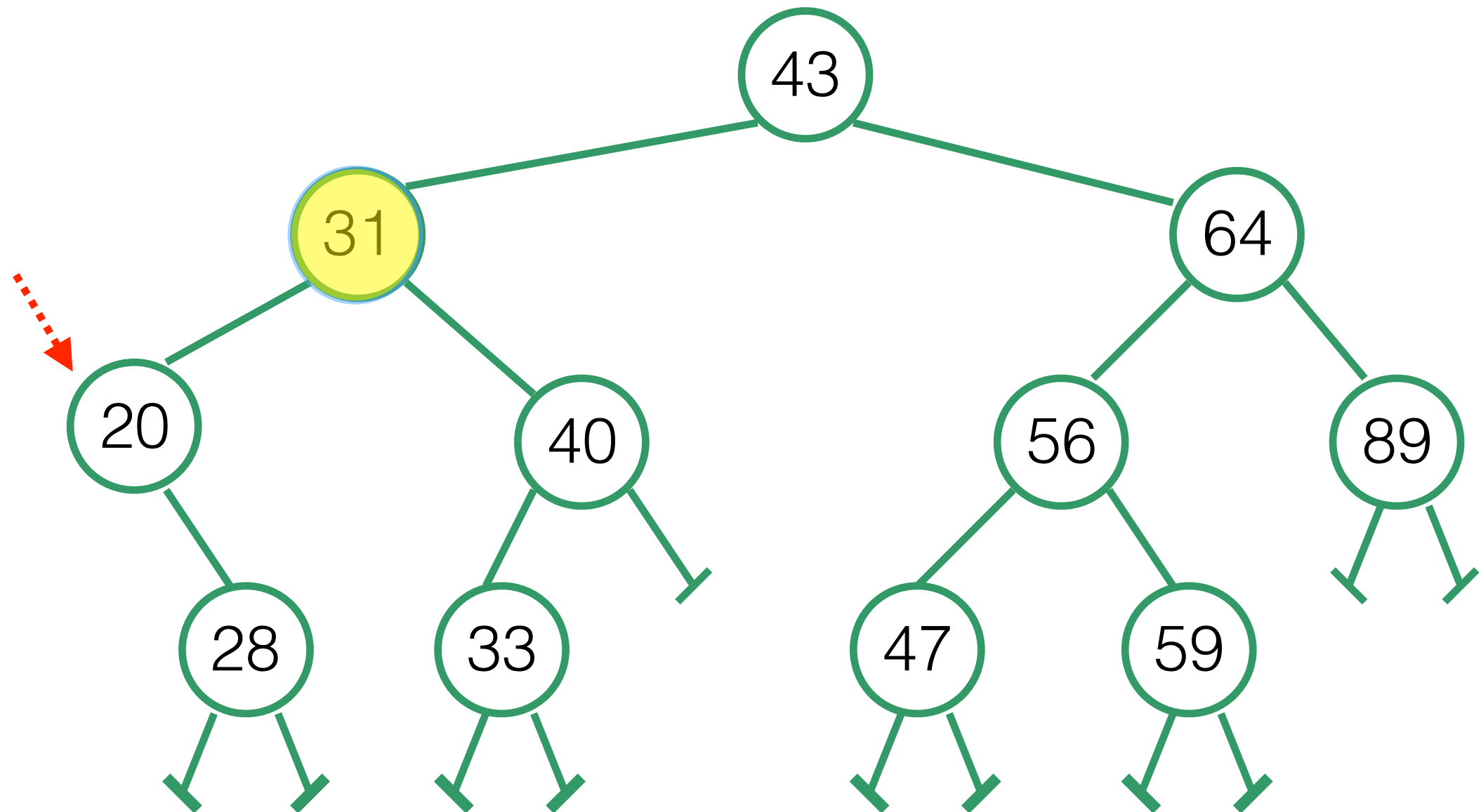
Node with no children:
Find parent - point to None

Delete 20



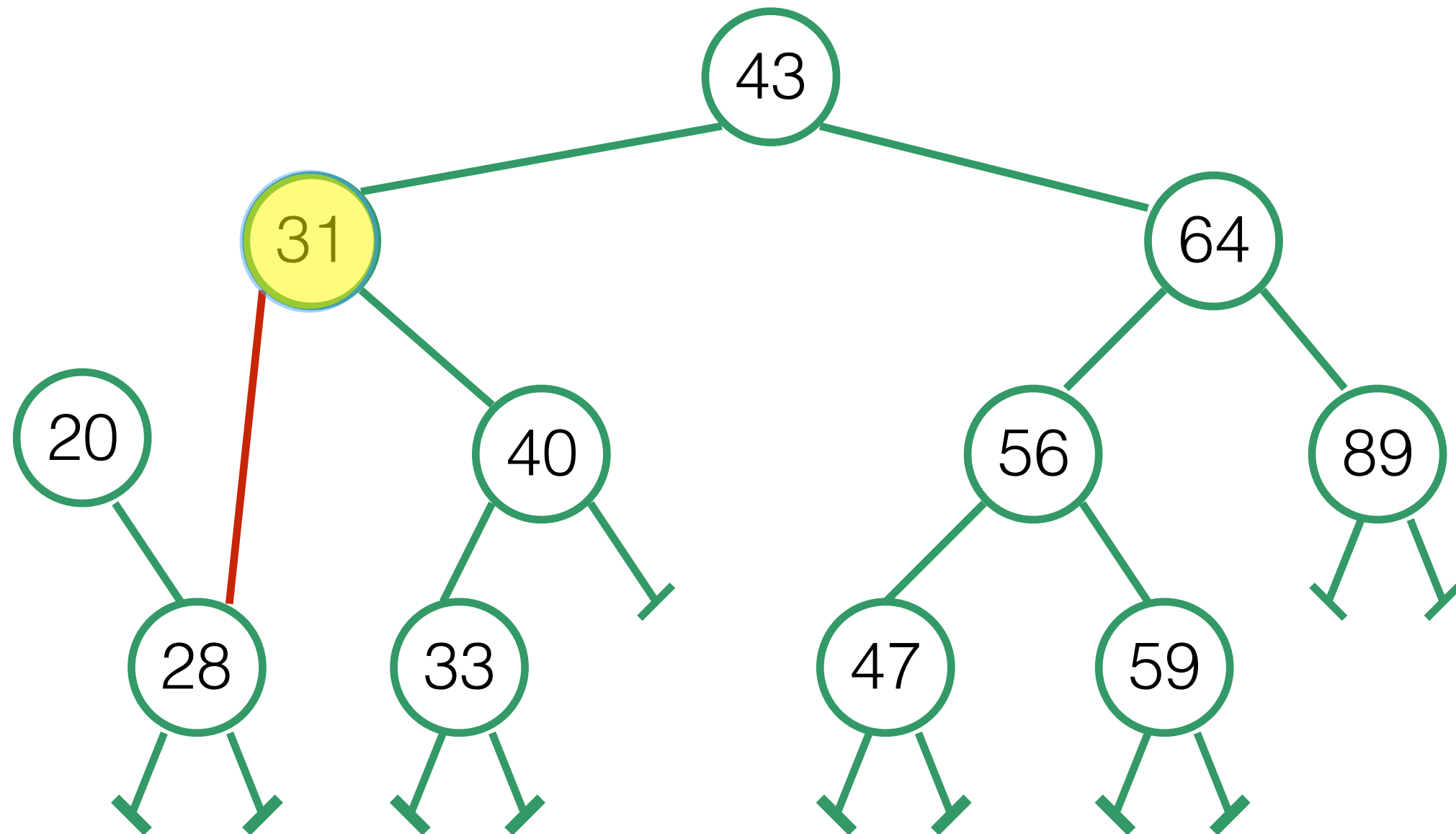
only showing **key**

Delete 20



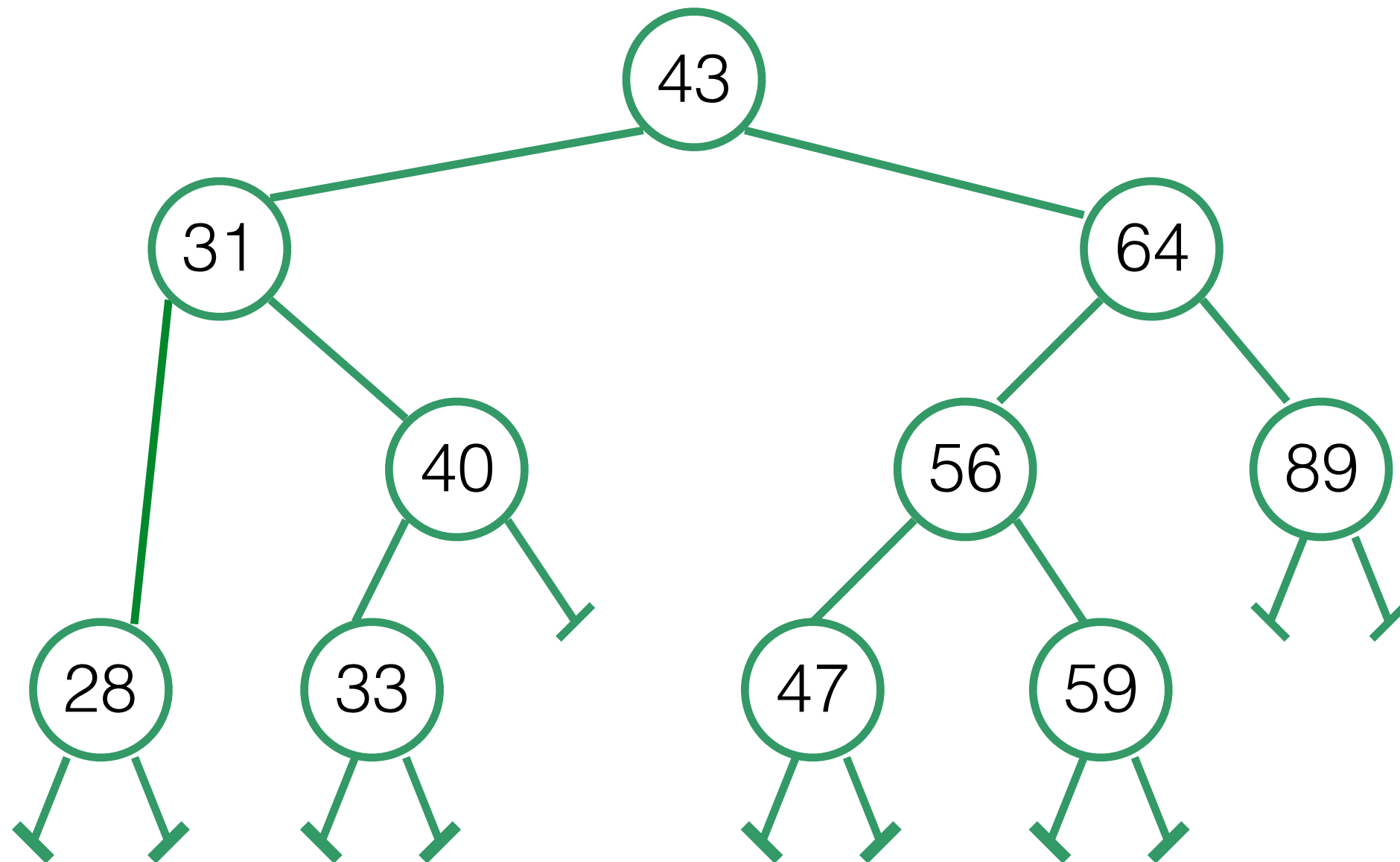
only showing **key**

Delete 20



only showing **key**

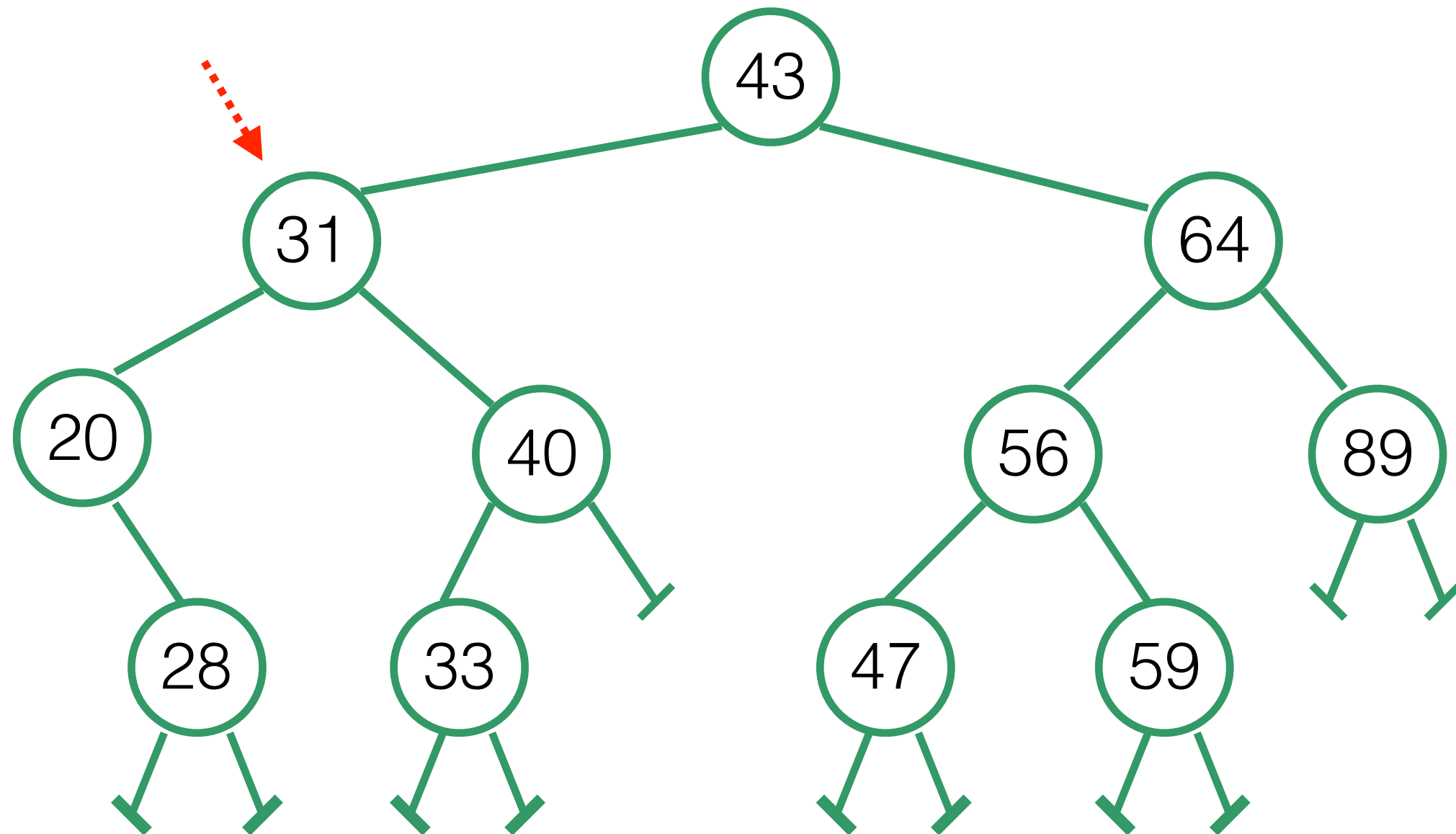
Delete 20



Node with one child:

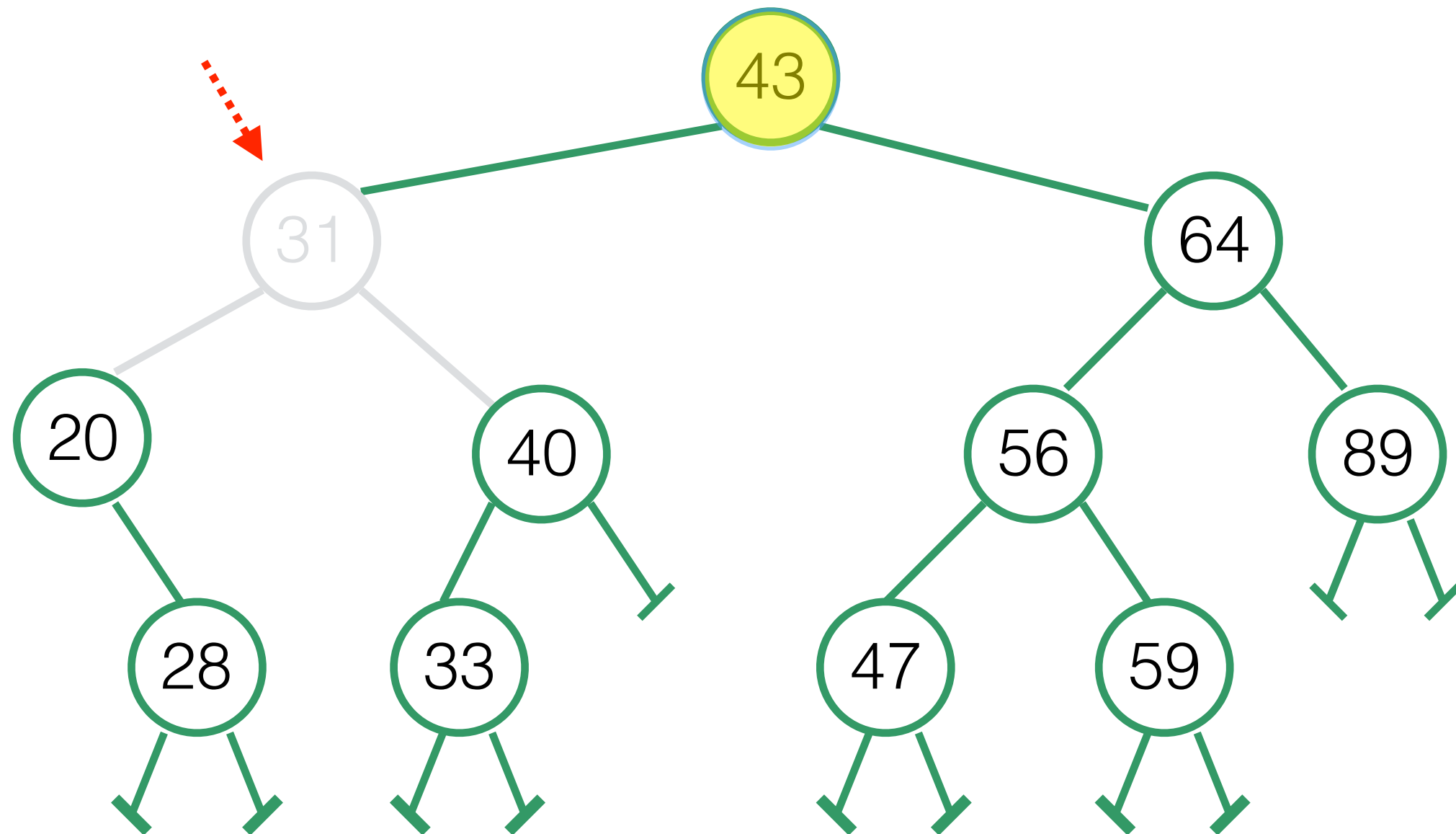
Find parent - point to child of deleted node

Delete 31



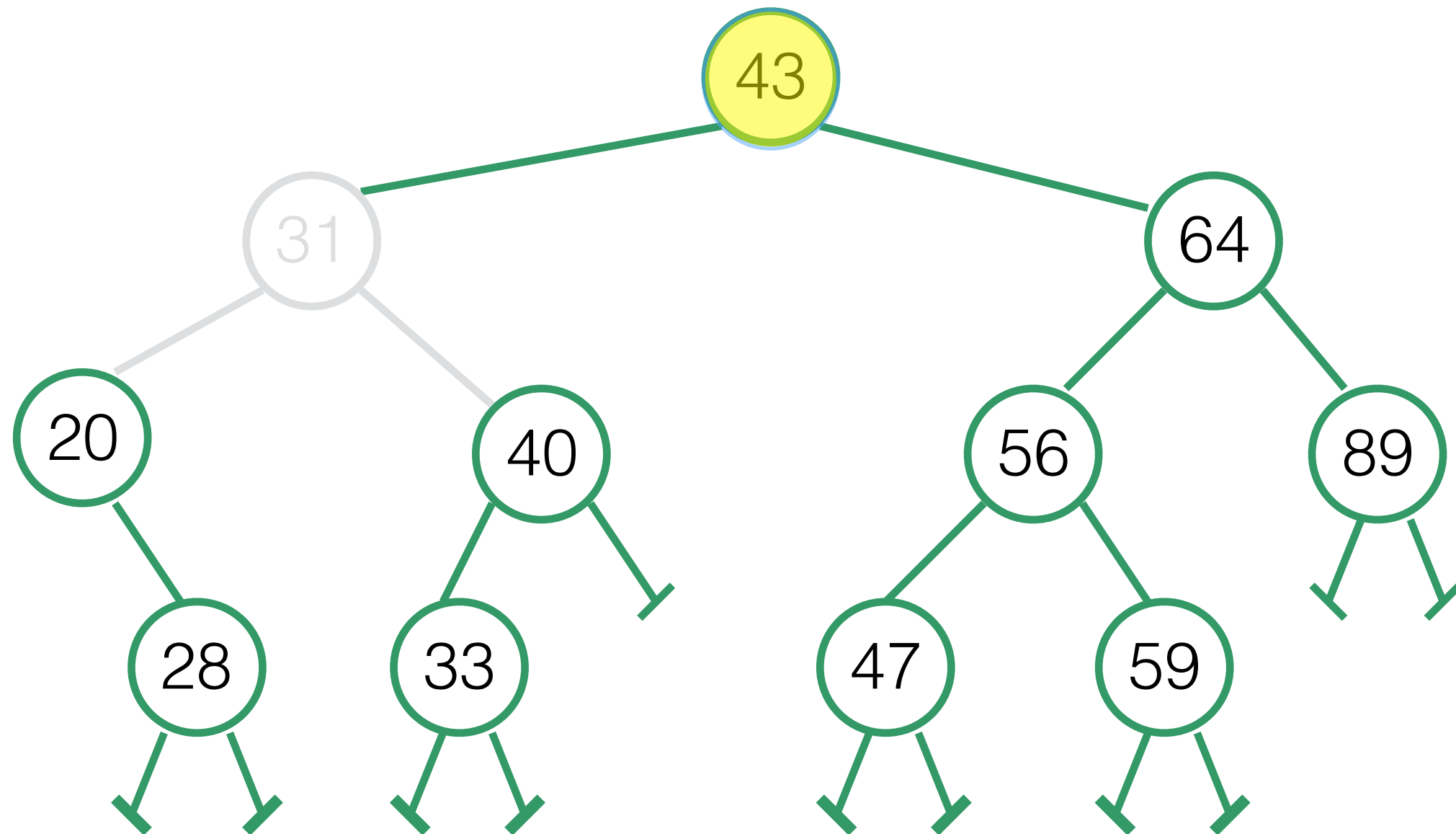
only showing **key**

Delete 31

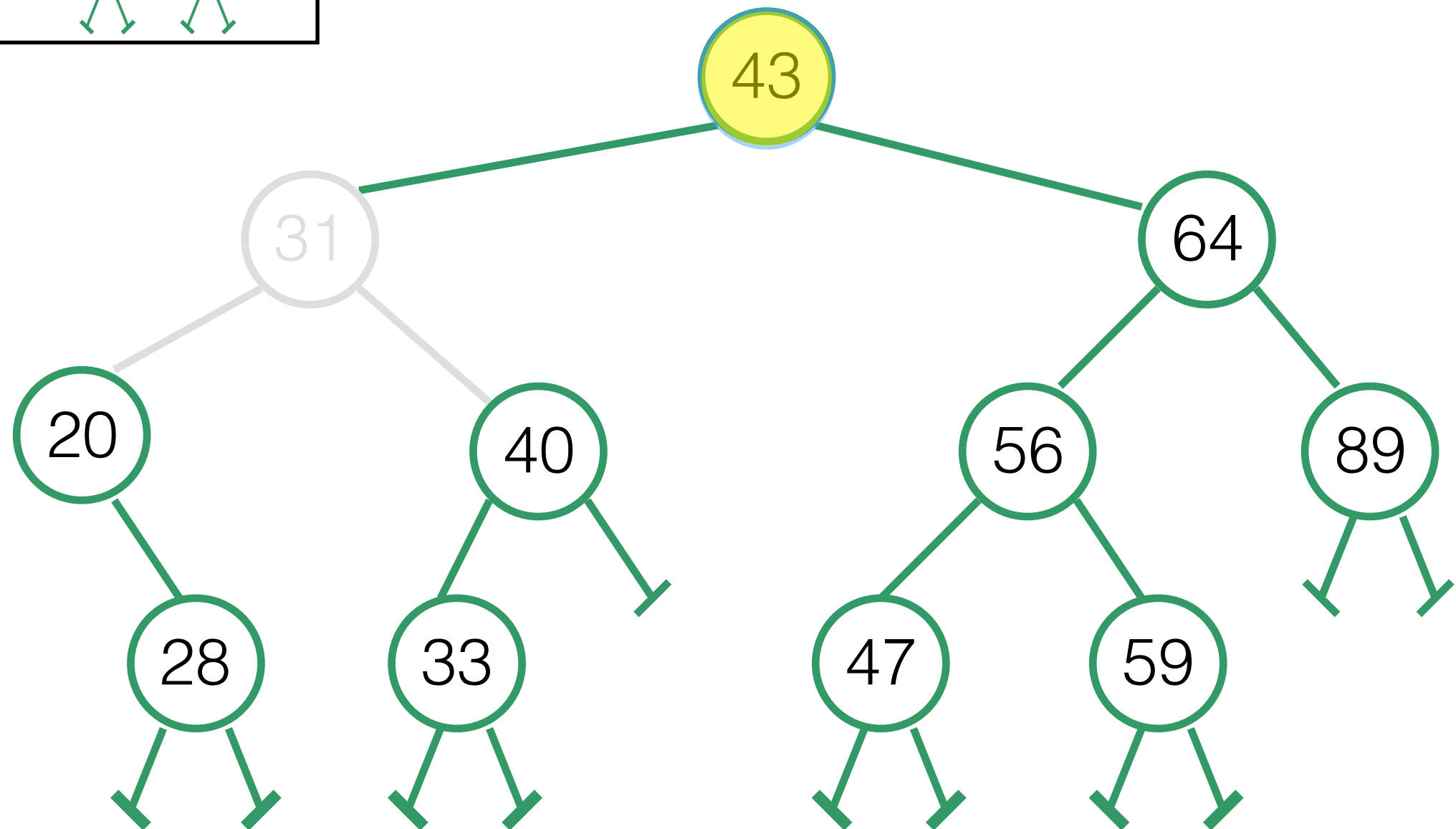
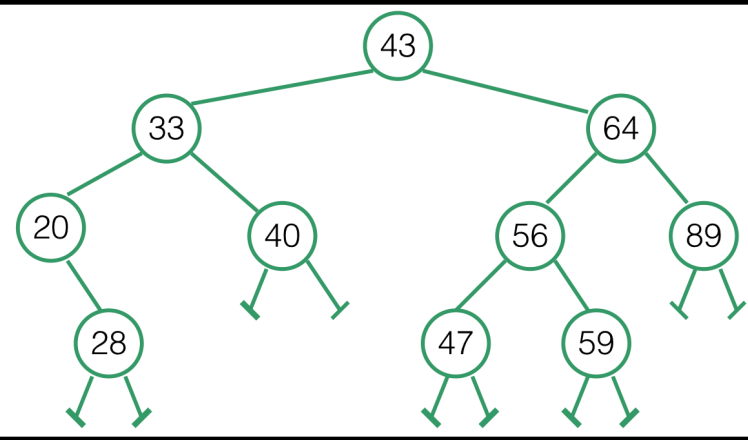


only showing **key**

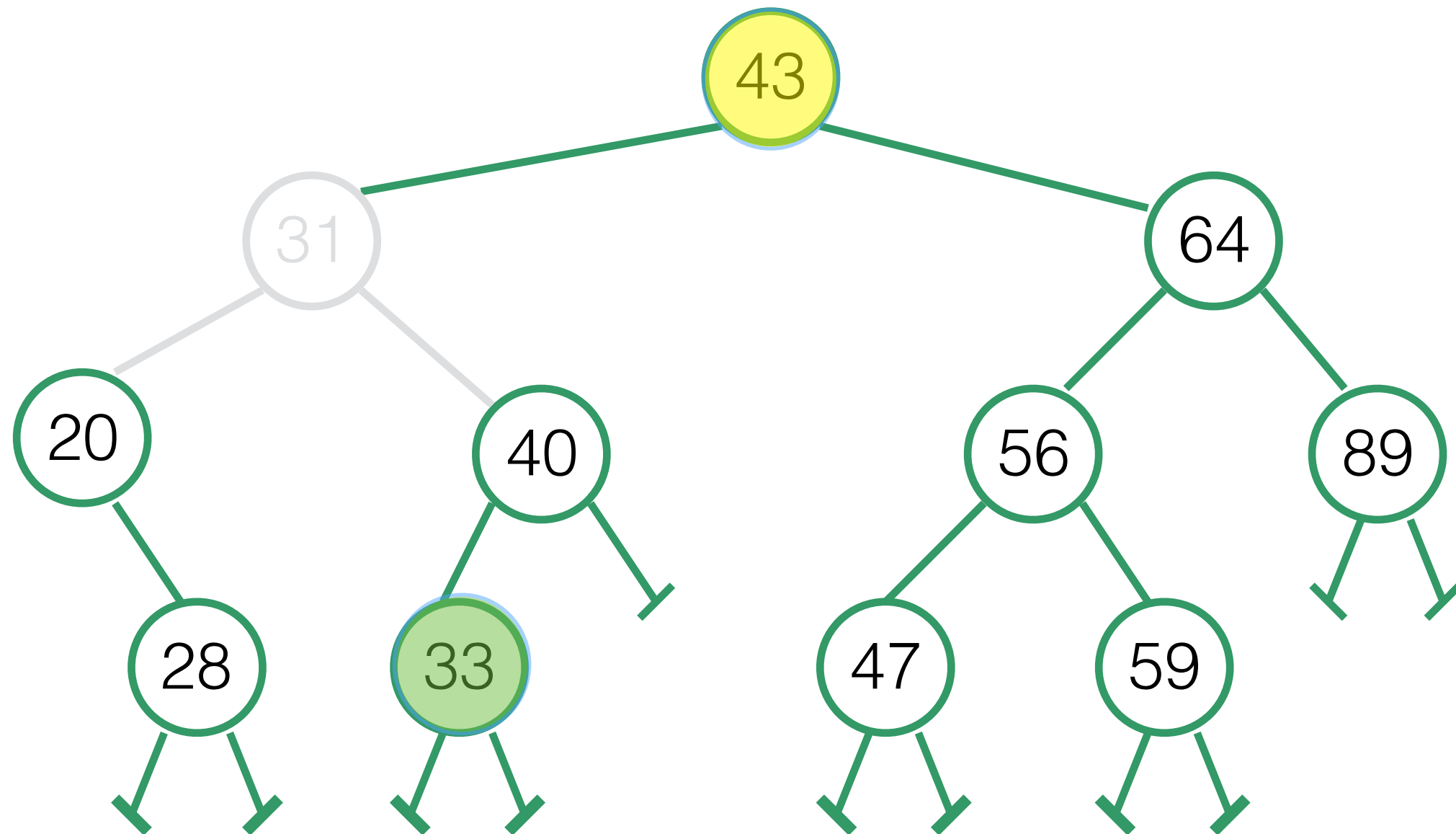
Delete 31



Delete 31

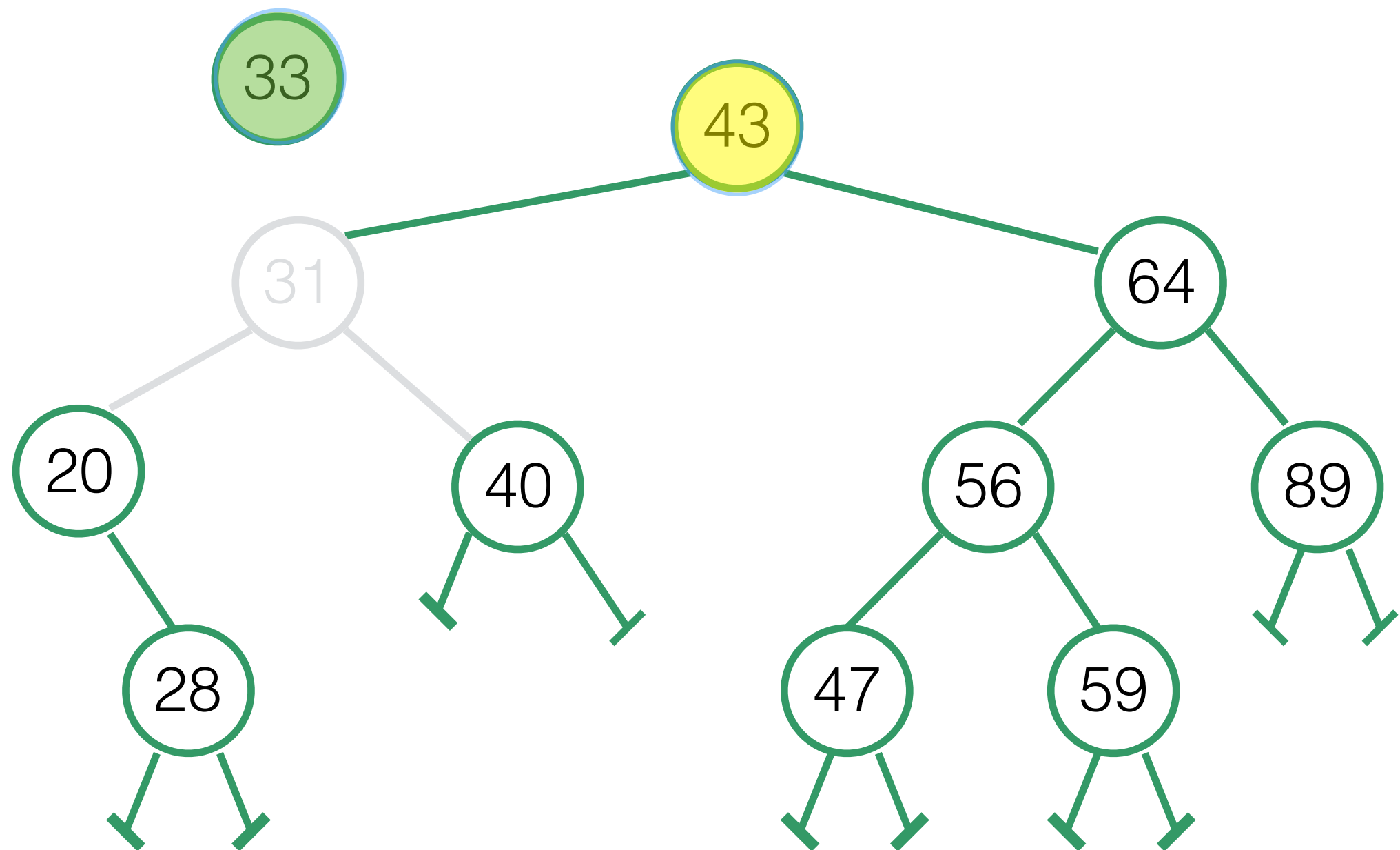


Delete 31

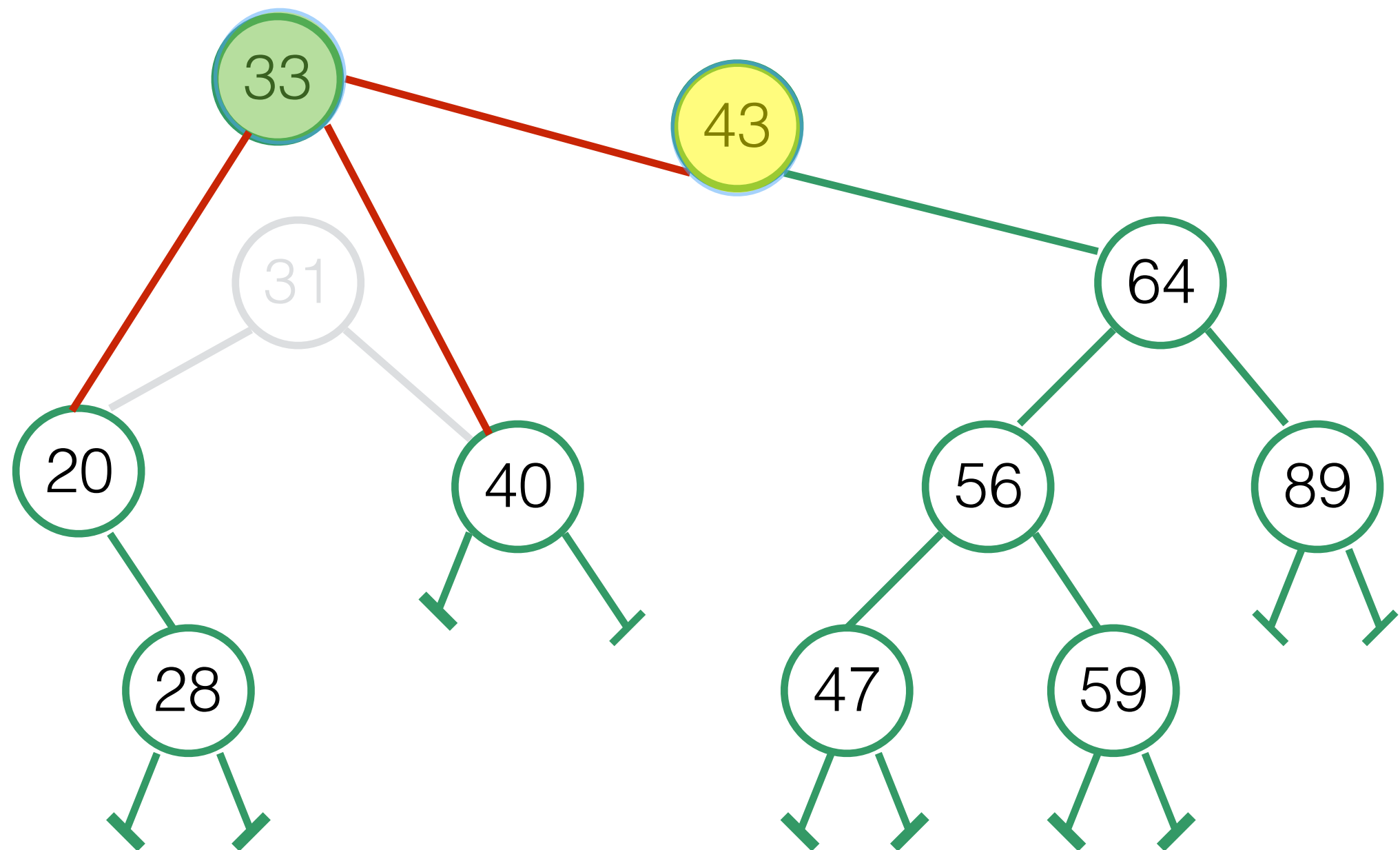


Successor of a node: node with **next** larger key.

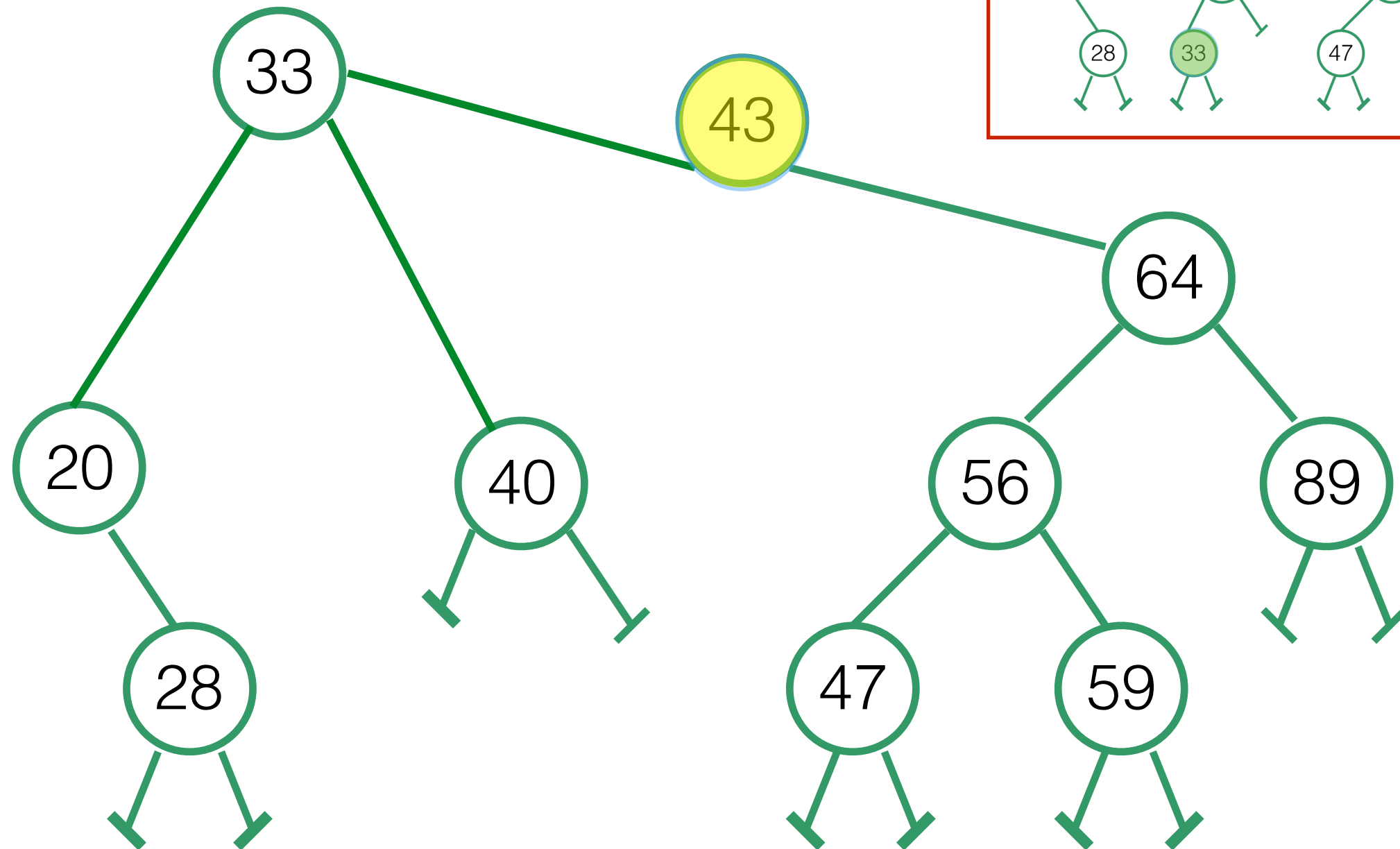
Delete 31



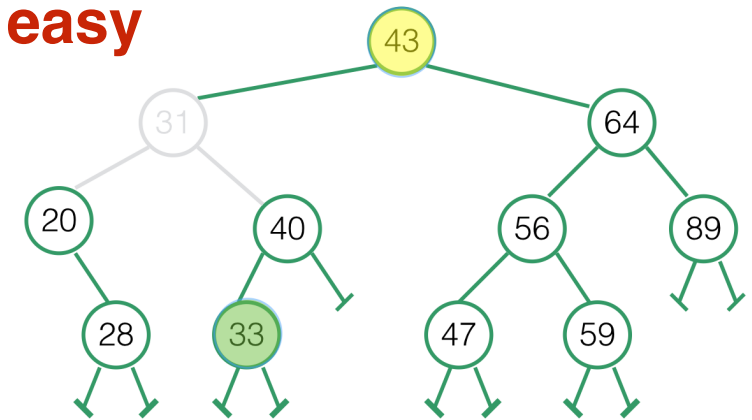
Delete 31



Delete 31



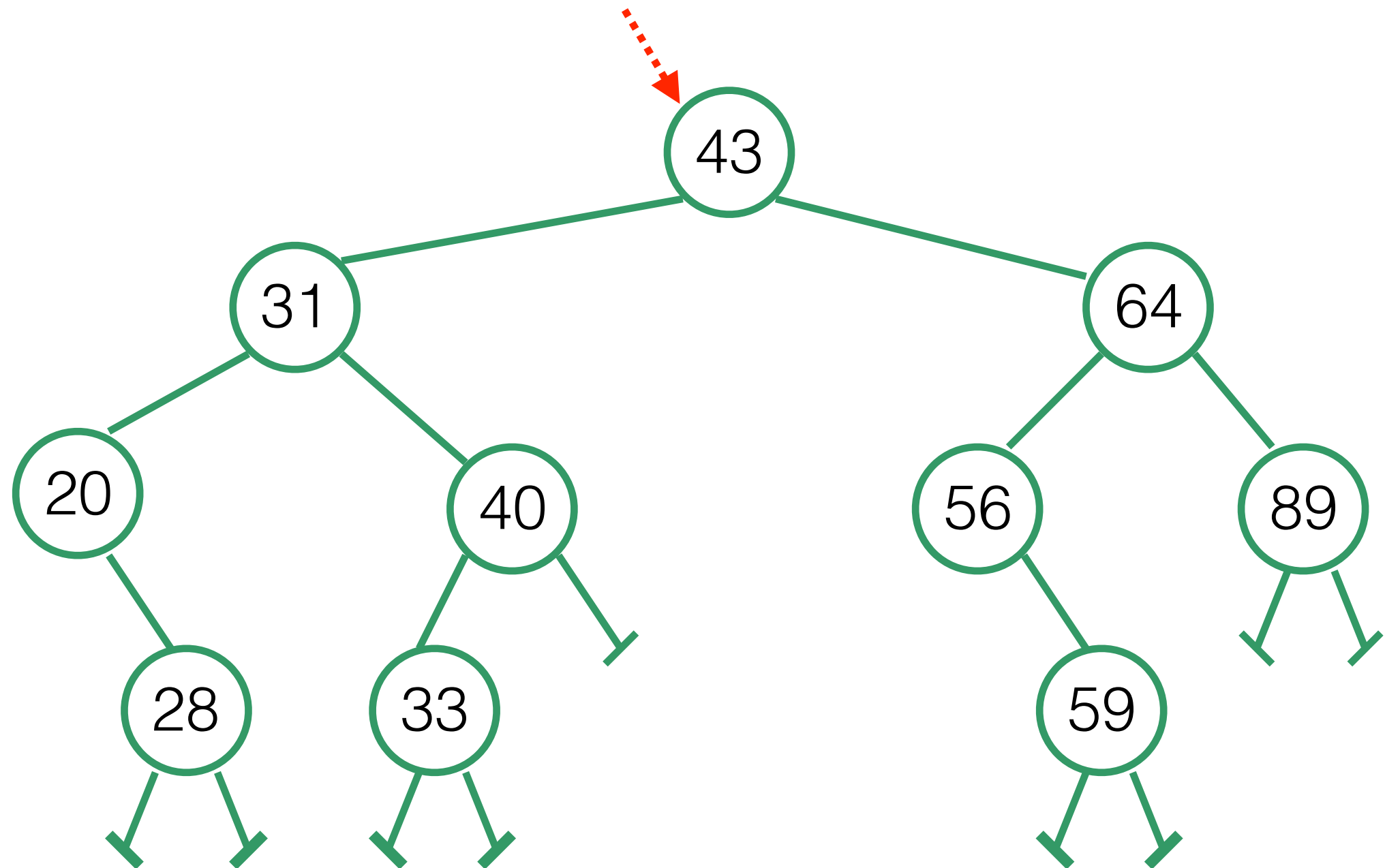
easy



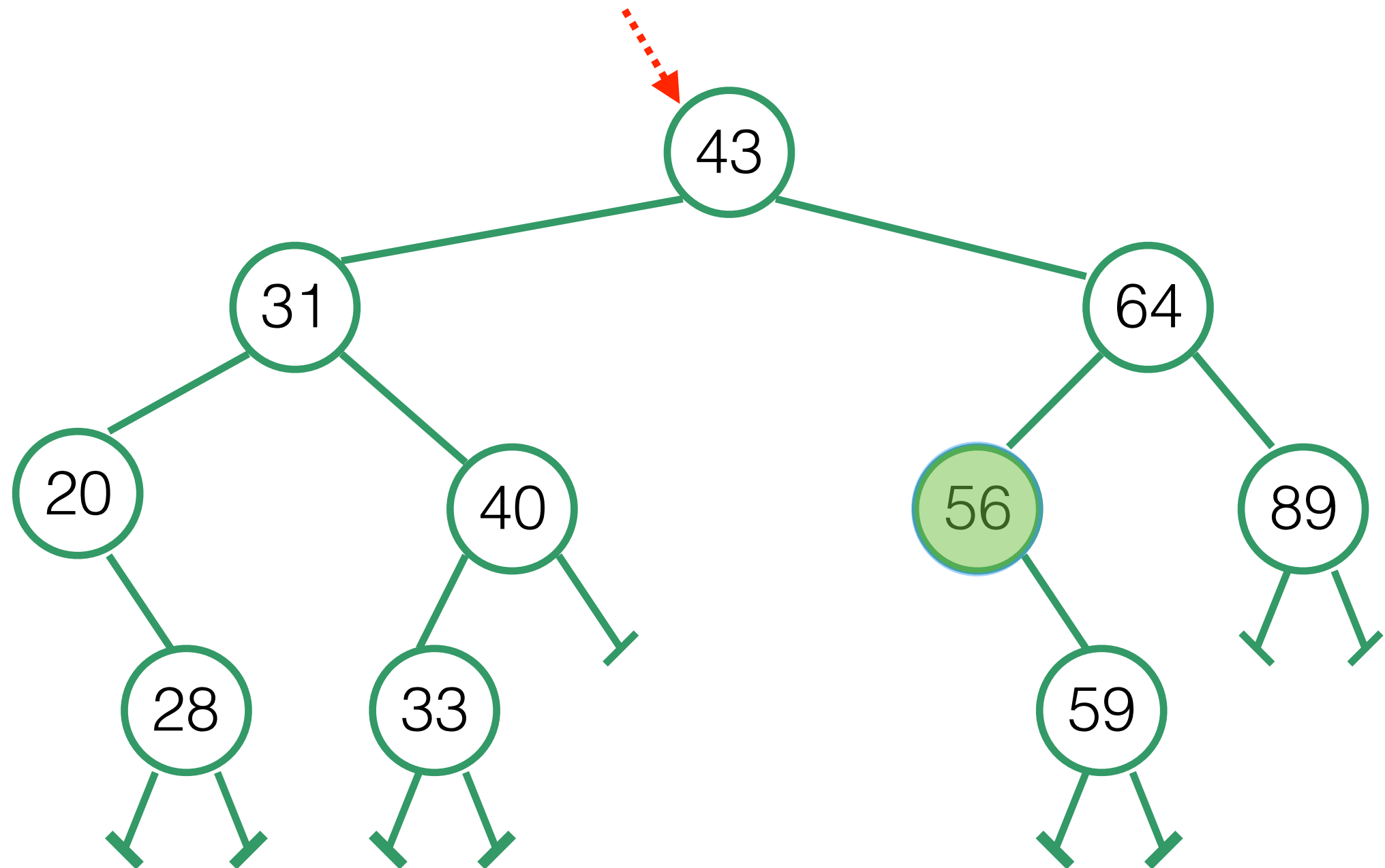
Node with two children:

Find parent and successor - successor is the new parent of the (orphan) children

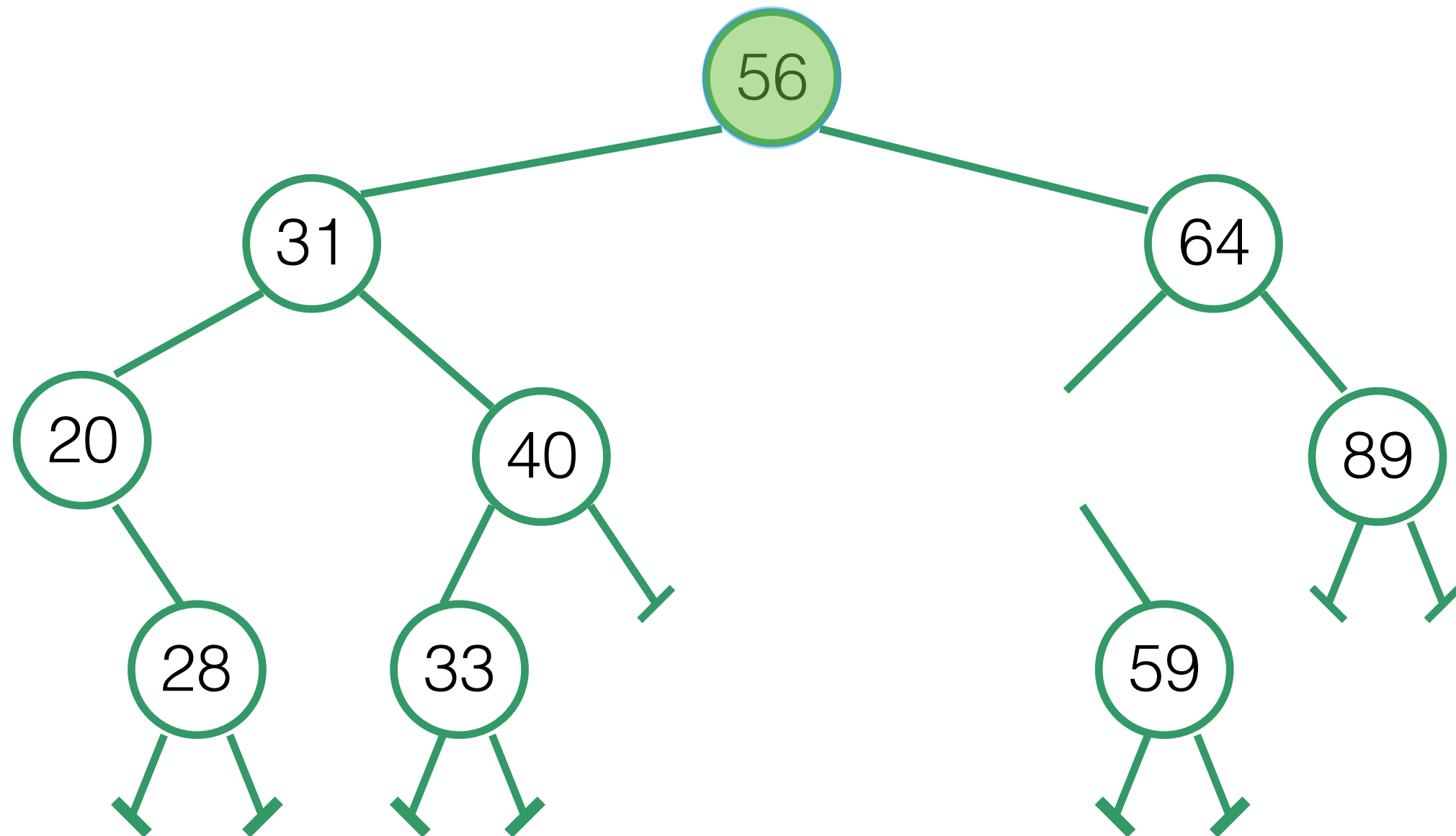
Delete 43



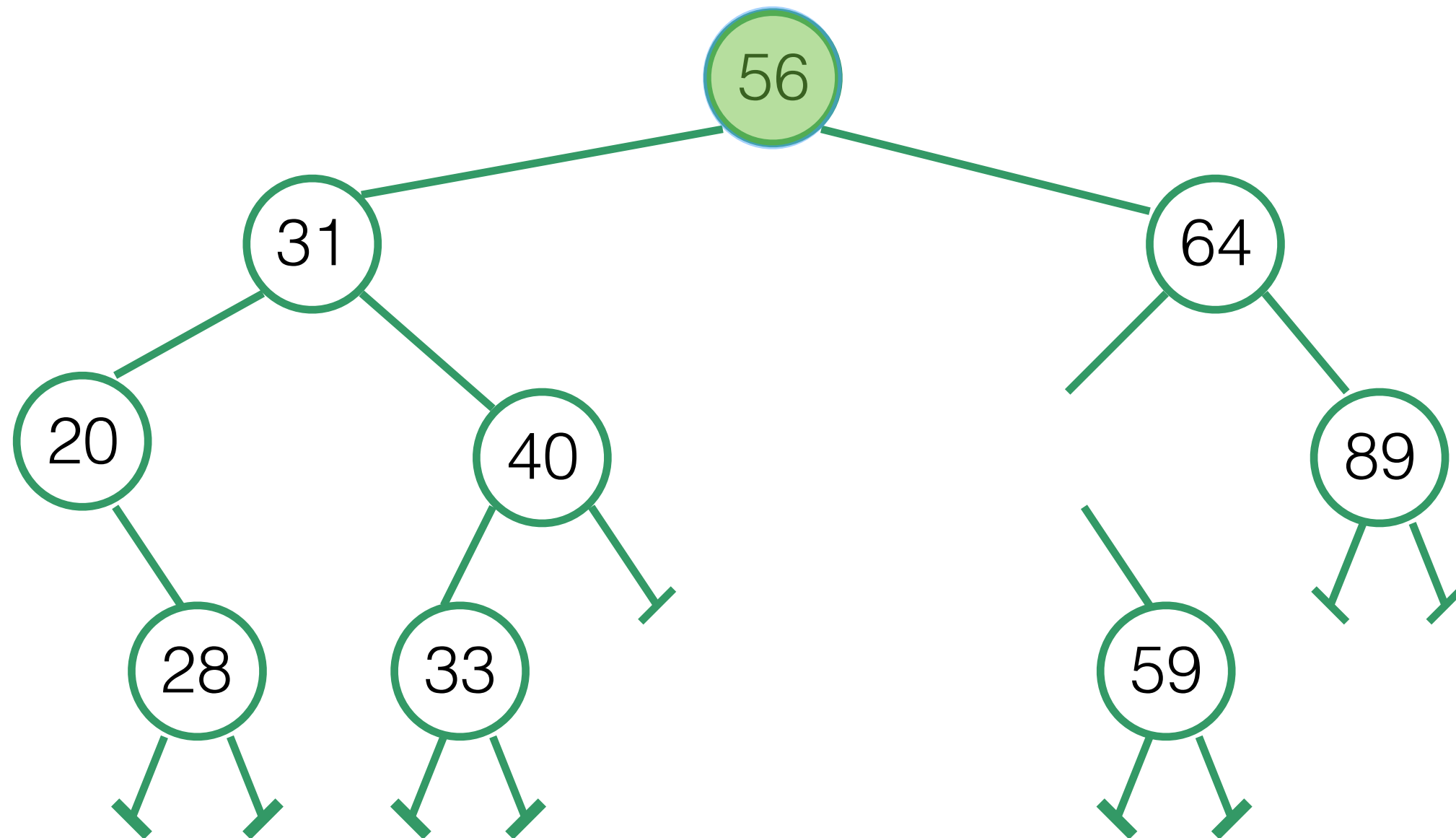
Delete 43



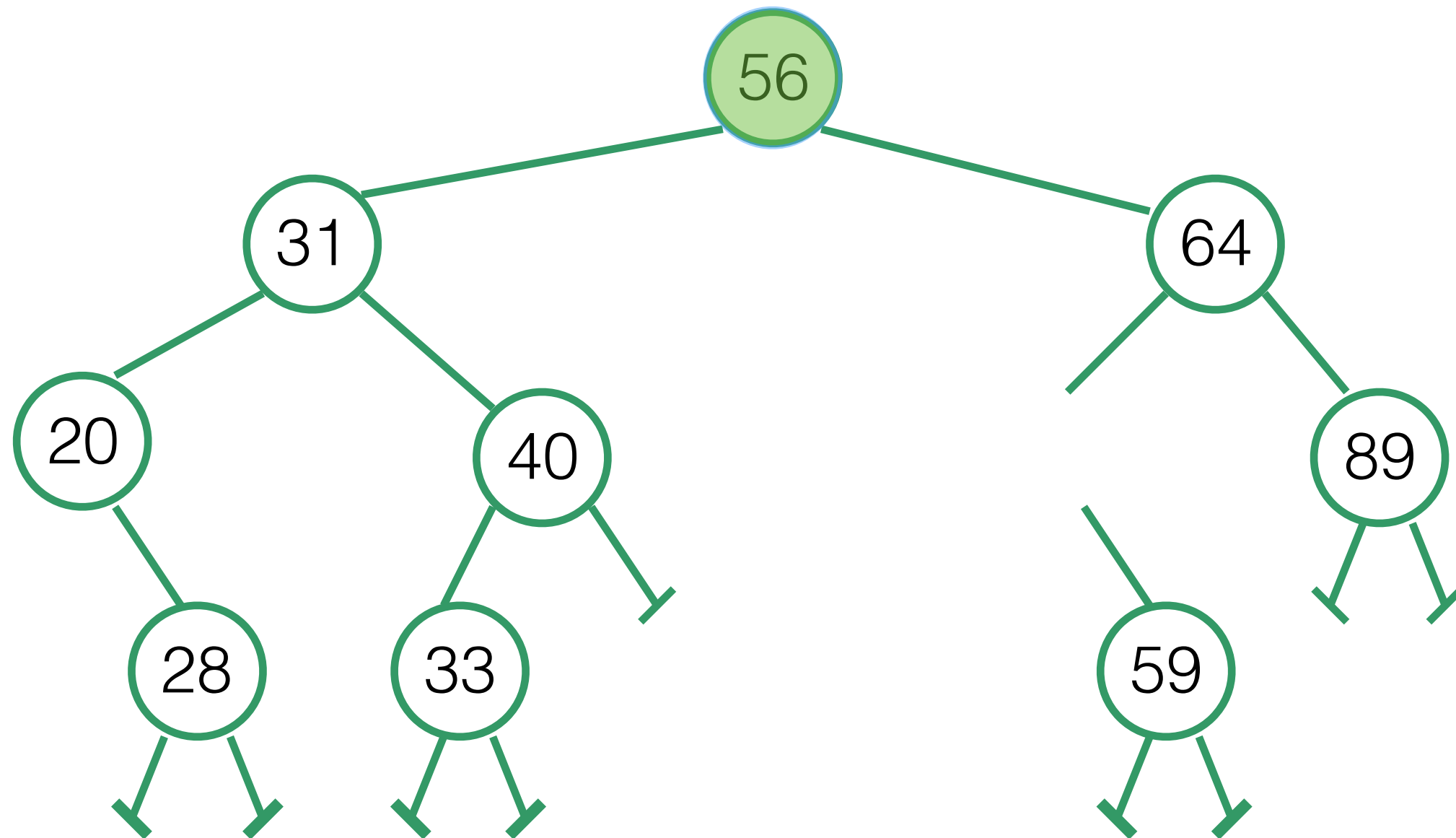
Delete 43



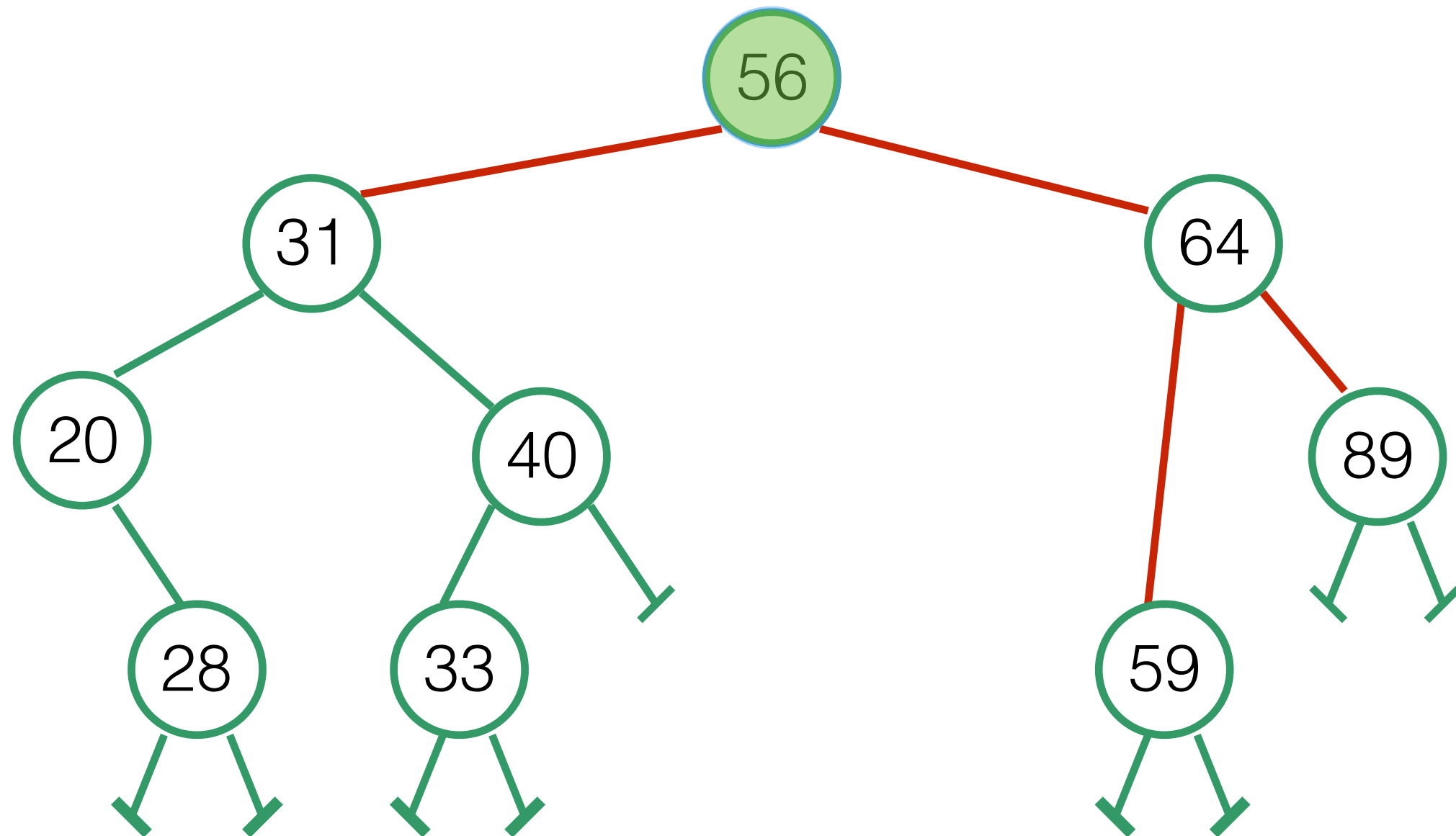
Delete 43



Delete 43



Delete 43



Delete

Input: key of element to delete.

Idea: Find key and successor...

- Try to find the key...
 - ➔ **If it is a leaf?** Set parent's reference to None
 - ➔ **It has one child?** Parent's reference set to child ("bypass").
 - ➔ **It has two children?** Find **successor**. Successor takes position of deleted node. **If successor leaves an orphan child, it should be linked to the successor's parent.**

Summary

- Binary search trees: search, insertion and deletion