Hill Climbing Algorithm

1. Generate initial feasible solution s.
2. Find s′ ∈ $N(s)$ such that $z(s') = \min\{z(s^+)\ \forall s^+ \in N(s)\}$.
3. While $z(s^+) < z(s)$:
4. | s = s’
5. | Find s′ ∈ $N(s)$ such that $z(s') = \min\{z(s^+)\ \forall s^+ \in N(s)\}$.
6. Return s

Simulated Annealing Algorithm

1. Generate initial feasible solution s.
2. Initialize $s^* = s$ and T.
3. Generate s′ ∈ $N(s)$.
4. While not stopping condition:
5. | If z$(s^+) < z(s)$:
6. | | s = s’
7. | | If z$(S) < z(s^*)$:
8. | | | $s^* = s$
9. | else:
10. | | s = s’ with probability $p = e^{\frac{-(z(s')-z(s))}{k\,T}}$
11. | Decrease T.
12. | Generate s′ ∈ $N(s)$.
13. Return s^*.

Tabu Search Algorithm

1. Generate initial feasible solution s.
2. Initialize $s^* = s$ and $TL = \emptyset$.
3. Find s′ ∈ $N(s)$ such that $z(s') = \min\{z(s^+)\ \forall s^+ \in N(s), s^+ \in TL\}$.
4. While not stopping condition:
5. $TL = TL\ \cup \{s\}$
6. s = s’
7. If $z(s) < z(s^*)$:
8. $s^* = s$
9. Return s^*.