
LABEL TRANSLATION (LINKING*)

➤ Introduction

➤ Illustration

Introduction

Labels are generated while translating the high level program constructs such as if-then-else, while-do and function calls into target machine code as an intermediate step in code generation. Labels are needed because the target address of a JUMP instruction may not be known at the time of generating code. However these labels must be replaced with target addresses in the final target code. Thus the compiler designer needs to write a program to replace all the labels in the program with the correct memory addresses eventually. This conversion is called Label translation.

The input to the label translator program is the machine code with labels (generated by `codeGen()`). The output is the target machine code without the labels (replaced with corresponding addresses).

The main idea behind label translation is to execute the following two steps.

1. The input (machine code with labels) is parsed and a table in which the labels are mapped with their corresponding addresses is created.
2. Parse the input again and replace all the labels with their corresponding addresses by looking into the table created above.

*Note: In real systems, the compiler generates an object file which contains labels for functions as well as variables. The object file is processed by a separate software module called the linker which converts the object file into an executable file. The advantage of this scheme is that a huge program could be divided into different stand alone source modules, each compiled separately and finally linked together. A change in one module does not require recompilation of other modules and only the linking phase needs to

be run again. In this project, we keep the linking task as a part of the compilation itself as ExpL does not allow a program to be compiled into several object modules.

In the ExpL project, variable addresses are resolved before reaching the Label translation phase and only symbolic references to functions are left to be replaced with addresses in this phase.

Illustration

The above procedure is demonstrated with the help of a program snippet.

```
1  decl
2      int fact(int n);
3  enddecl
4
5  int fact(int n)
6  {
7      decl
8          int f;
9      enddecl
10     begin
11         if(n<=1) then
12             f=1;
13         else
14             f=n*fact(
15         endif;
16         return f;
17     end
18 }
19
20 int main()
21 {
22     decl
23         int n,m,res;
24     enddecl
25     begin
26         read(n);
27         while( n >= 1 ) c
28             read(m);
29             res = fac
30             write(res
31             n = n-1;
```

```
32         endwhile;
33         return 0;
34     end
35 }
```

Factorial.c hosted with ❤ by [GitHub](#) [view raw](#)

The code generated after the codeGen() phase is shown in the figure 1 (It consists of labels).

Figure 1

1	2048 : 0
1	2049 : 2056
1	2050 : 0
1	2051 : 0
1	2052 : 0
1	2053 : 0
1	2054 : 0
1	2055 : 0
5	2056 : MOV SP,409
6	2058 : MOV BP,409
7	2060 : PUSH R0
8	2062 : CALL MAIN
9	2064 : INT 10
10	F0:
11	2066 : PUSH BP
12	2068 : MOV BP,SP
13	2070 : PUSH R0
14	2072 : MOV R1,BP
15	2074 : MOV R2,2
16	2076 : SUB R1,R2
17	2078 : MOV R2,1
18	2080 : SUB R1,R2
19	2082 : MOV R0,[R1
20	2084 : MOV R1,1
21	2086 : LE R0,R1
22	2088 : JZ R0,L0
23	2090 : MOV R0,1
24	2092 : MOV R2,BP
25	2094 : MOV R1,1
26	2096 : ADD R2,R1
27	2098 : MOV [R2],R
28	2100 : JMP L1
29	L0:
30	2102 : MOV R1,BP
31	2104 : MOV R2,2

32	2106 : SUB R1,R2
33	2108 : MOV R2,1
34	2110 : SUB R1,R2
35	2112 : MOV R0,[R1
36	2114 : PUSH R0
37	2116 : MOV R1,BP
38	2118 : MOV R2,2
39	2120 : SUB R1,R2
40	2122 : MOV R2,1
41	2124 : SUB R1,R2
42	2126 : MOV R0,[R1
43	2128 : MOV R1,1
44	2130 : SUB R0,R1
45	2132 : PUSH R0
46	2134 : PUSH R0
47	2136 : CALL F0
48	2138 : POP R0
49	2140 : POP R0
50	2142 : POP R0
51	2144 : MOV R1,3
52	2146 : MOV R2,SP
53	2148 : ADD R2,R1
54	2150 : MOV R1,[R2
55	2152 : MUL R0,R1
56	2154 : MOV R2,BP
57	2156 : MOV R1,1
58	2158 : ADD R2,R1
59	2160 : MOV [R2],R
60	L1:
61	2162 : MOV R1,BP
62	2164 : MOV R0,1
63	2166 : ADD R1,R0
64	2168 : MOV R0,[R1
65	2170 : MOV R1,BP
66	2172 : MOV R2,2
67	2174 : SUB R1,R2
68	2176 : MOV [R1],R
69	2178 : POP R0
70	2180 : MOV BP,[SP
71	2182 : POP R0
72	2184 : RET
73	MAIN:
74	2186 : PUSH BP
75	2188 : MOV BP,SP
76	2190 : PUSH R0
77	2192 : PUSH R0
78	2194 : PUSH R0

79	2196 : MOV R1, BP
80	2198 : MOV R0, 1
81	2200 : ADD R1, R0
82	2202 : PUSH R0
83	2204 : PUSH R1
84	2206 : MOV R0, "Re
85	2208 : PUSH R0
86	2210 : MOV R0, -1
87	2212 : PUSH R0
88	2214 : PUSH R1
89	2216 : PUSH R0
90	2218 : PUSH R0
91	2220 : CALL 0
92	2222 : POP R0
93	2224 : POP R0
94	2226 : POP R0
95	2228 : POP R0
96	2230 : POP R0
97	2232 : POP R0
98	2234 : POP R0
99	L2:
100	2236 : MOV R1, BP
101	2238 : MOV R0, 1
102	2240 : ADD R1, R0
103	2242 : MOV R0, [R1
104	2244 : MOV R1, 1
105	2246 : GE R0, R1
106	2248 : JZ R0, L3
107	2250 : MOV R1, BP
108	2252 : MOV R0, 2
109	2254 : ADD R1, R0
110	2256 : PUSH R0
111	2258 : PUSH R1
112	2260 : MOV R0, "Re
113	2262 : PUSH R0
114	2264 : MOV R0, -1
115	2266 : PUSH R0
116	2268 : PUSH R1
117	2270 : PUSH R0
118	2272 : PUSH R0
119	2274 : CALL 0
120	2276 : POP R0
121	2278 : POP R0
122	2280 : POP R0
123	2282 : POP R0
124	2284 : POP R0
125	2286 : POP R0

126	2288 : POP R0
127	2290 : MOV R1,BP
128	2292 : MOV R0,2
129	2294 : ADD R1,R0
130	2296 : MOV R0,[R1
131	2298 : PUSH R0
132	2300 : PUSH R0
133	2302 : CALL F0
134	2304 : POP R0
135	2306 : POP R0
136	2308 : MOV R0,2
137	2310 : MOV R1,SP
138	2312 : ADD R1,R0
139	2314 : MOV R0,[R1
140	2316 : MOV R2,BP
141	2318 : MOV R1,3
142	2320 : ADD R2,R1
143	2322 : MOV [R2],R
144	2324 : MOV R1,BP
145	2326 : MOV R0,3
146	2328 : ADD R1,R0
147	2330 : MOV R0,[R1
148	2332 : MOV [2042]
149	2334 : PUSH R0
150	2336 : MOV R0,"Wr
151	2338 : PUSH R0
152	2340 : MOV R0,-2
153	2342 : PUSH R0
154	2344 : MOV R0,204
155	2346 : PUSH R0
156	2348 : PUSH R0
157	2350 : PUSH R0
158	2352 : CALL 0
159	2354 : POP R0
160	2356 : POP R0
161	2358 : POP R0
162	2360 : POP R0
163	2362 : POP R0
164	2364 : POP R0
165	2366 : MOV R1,BP
166	2368 : MOV R0,1
167	2370 : ADD R1,R0
168	2372 : MOV R0,[R1
169	2374 : MOV R1,1
170	2376 : SUB R0,R1
171	2378 : MOV R2,BP
172	2380 : MOV R1,1

173	2382 : ADD R2,R1
174	2384 : MOV [R2],R
175	2386 : JMP L2
176	L3:
177	2388 : MOV R0,0
178	2390 : MOV R1,BP
179	2392 : MOV R2,2
180	2394 : SUB R1,R2
181	2396 : MOV [R1],R
182	2398 : POP R0
183	2400 : POP R0
184	2402 : POP R0
185	2404 : MOV BP,[SP
186	2406 : POP R0
187	2408 : RET

Assembly Code view raw
Before Label Translation
hosted with ❤ by GitHub

Figure 2

1	2048 : 0
1	2049 : 2056
1	2050 : 0
1	2051 : 0
1	2052 : 0
1	2053 : 0
1	2054 : 0
1	2055 : 0
5	2056 : MOV SP,409
6	2058 : MOV BP,409
7	2060 : PUSH R0
8	2062 : CALL 2186
9	2064 : INT 10
10	2066 : PUSH BP
11	2068 : MOV BP,SP
12	2070 : PUSH R0
13	2072 : MOV R1,BP
14	2074 : MOV R2,2
15	2076 : SUB R1,R2
16	2078 : MOV R2,1
17	2080 : SUB R1,R2
18	2082 : MOV R0,[R1
19	2084 : MOV R1,1
20	2086 : LE R0,R1
21	2088 : JZ R0,2102
22	2090 : MOV R0,1

23	2092 : MOV R2,BP
24	2094 : MOV R1,1
25	2096 : ADD R2,R1
26	2098 : MOV [R2],R
27	2100 : JMP 2162
28	2102 : MOV R1,BP
29	2104 : MOV R2,2
30	2106 : SUB R1,R2
31	2108 : MOV R2,1
32	2110 : SUB R1,R2
33	2112 : MOV R0,[R1
34	2114 : PUSH R0
35	2116 : MOV R1,BP
36	2118 : MOV R2,2
37	2120 : SUB R1,R2
38	2122 : MOV R2,1
39	2124 : SUB R1,R2
40	2126 : MOV R0,[R1
41	2128 : MOV R1,1
42	2130 : SUB R0,R1
43	2132 : PUSH R0
44	2134 : PUSH R0
45	2136 : CALL 2066
46	2138 : POP R0
47	2140 : POP R0
48	2142 : POP R0
49	2144 : MOV R1,3
50	2146 : MOV R2,SP
51	2148 : ADD R2,R1
52	2150 : MOV R1,[R2
53	2152 : MUL R0,R1
54	2154 : MOV R2,BP
55	2156 : MOV R1,1
56	2158 : ADD R2,R1
57	2160 : MOV [R2],R
58	2162 : MOV R1,BP
59	2164 : MOV R0,1
60	2166 : ADD R1,R0
61	2168 : MOV R0,[R1
62	2170 : MOV R1,BP
63	2172 : MOV R2,2
64	2174 : SUB R1,R2
65	2176 : MOV [R1],R
66	2178 : POP R0
67	2180 : MOV BP,[SP
68	2182 : POP R0
69	2184 : RET

70	2186 : PUSH BP
71	2188 : MOV BP,SP
72	2190 : PUSH R0
73	2192 : PUSH R0
74	2194 : PUSH R0
75	2196 : MOV R1,BP
76	2198 : MOV R0,1
77	2200 : ADD R1,R0
78	2202 : PUSH R0
79	2204 : PUSH R1
80	2206 : MOV R0,"Re
81	2208 : PUSH R0
82	2210 : MOV R0,-1
83	2212 : PUSH R0
84	2214 : PUSH R1
85	2216 : PUSH R0
86	2218 : PUSH R0
87	2220 : CALL 0
88	2222 : POP R0
89	2224 : POP R0
90	2226 : POP R0
91	2228 : POP R0
92	2230 : POP R0
93	2232 : POP R0
94	2234 : POP R0
95	2236 : MOV R1,BP
96	2238 : MOV R0,1
97	2240 : ADD R1,R0
98	2242 : MOV R0,[R1
99	2244 : MOV R1,1
100	2246 : GE R0,R1
101	2248 : JZ R0,2388
102	2250 : MOV R1,BP
103	2252 : MOV R0,2
104	2254 : ADD R1,R0
105	2256 : PUSH R0
106	2258 : PUSH R1
107	2260 : MOV R0,"Re
108	2262 : PUSH R0
109	2264 : MOV R0,-1
110	2266 : PUSH R0
111	2268 : PUSH R1
112	2270 : PUSH R0
113	2272 : PUSH R0
114	2274 : CALL 0
115	2276 : POP R0
116	2278 : POP R0

117	2280 : POP R0
118	2282 : POP R0
119	2284 : POP R0
120	2286 : POP R0
121	2288 : POP R0
122	2290 : MOV R1,BP
123	2292 : MOV R0,2
124	2294 : ADD R1,R0
125	2296 : MOV R0,[R1
126	2298 : PUSH R0
127	2300 : PUSH R0
128	2302 : CALL 2066
129	2304 : POP R0
130	2306 : POP R0
131	2308 : MOV R0,2
132	2310 : MOV R1,SP
133	2312 : ADD R1,R0
134	2314 : MOV R0,[R1
135	2316 : MOV R2,BP
136	2318 : MOV R1,3
137	2320 : ADD R2,R1
138	2322 : MOV [R2],R
139	2324 : MOV R1,BP
140	2326 : MOV R0,3
141	2328 : ADD R1,R0
142	2330 : MOV R0,[R1
143	2332 : MOV [2042]
144	2334 : PUSH R0
145	2336 : MOV R0,"Wr
146	2338 : PUSH R0
147	2340 : MOV R0,-2
148	2342 : PUSH R0
149	2344 : MOV R0,204
150	2346 : PUSH R0
151	2348 : PUSH R0
152	2350 : PUSH R0
153	2352 : CALL 0
154	2354 : POP R0
155	2356 : POP R0
156	2358 : POP R0
157	2360 : POP R0
158	2362 : POP R0
159	2364 : POP R0
160	2366 : MOV R1,BP
161	2368 : MOV R0,1
162	2370 : ADD R1,R0
163	2372 : MOV R0,[R1

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164	2374 : MOV R1,1
165	2376 : SUB R0,R1
166	2378 : MOV R2,BP
167	2380 : MOV R1,1
168	2382 : ADD R2,R1
169	2384 : MOV [R2],R
170	2386 : JMP 2236
171	2388 : MOV R0,0
172	2390 : MOV R1,BP
173	2392 : MOV R2,2
174	2394 : SUB R1,R2
175	2396 : MOV [R1],R
176	2398 : POP R0
177	2400 : POP R0
178	2402 : POP R0
179	2404 : MOV BP,[SP
180	2406 : POP R0
181	2408 : RET

Assembly Code view raw
after Label Translation
hosted with ❤ by GitHub

In the machine code generated after codegen section, labels occur in two different ways.

1. Label declarations in which labels are followed by semicolon (:).
Examples : F0: is a label for the instruction at address 2066, L0: for 2102, L1: for 2162 and MAIN: for 2186 as in figure1.
2. Instructions which contain labels in them are JMP, JZ, JNZ and CALL.
Examples : JMP L1 at address 2100, JZ R0,L3 at address 2248, CALL MAIN at address 2062 etc., as in figure1.

Now, we will understand the label translation procedure with the help of an example label F0 from the figure1.

First F0: need to be recognized and the memory address of the label occurrence need to be stored in a table called Label-Address Table. In this case, the address corresponding to F0: is 2066. We need to parse the entire machine code to identify all labels and store all (label, address) pairs in the Label-Address table. In this pass, we remove the label F0: as well as other labels from the program. The next task is to replace labels occurring in instructions with the addresses of the labels from the label-address table. Referring to the figure1, there are two instructions that use the label F0: CALL F0 occurring at address 2136 and CALL F0 occurring at address 2302. We will

replace the labels in these instructions with the address of F0 (2066) by looking up the label-address table. Thus, after translation, both these instructions will translate to CALL 2066 as shown in figure2.

The label translator program need to parse the entire machine code two times. In the first parse we identify all the label declarations and the memory addresses in the label-address table. The table constructed after the first parse of the above code is shown below.

LABEL-ADDRESS Table

LABEL	ADDRESS
F0	2066
L0	2102
L1	2162
MAIN	2186
L2	2236
L3	2388

In the second parse, we remove the label declarations and replace the labels in instructions like JUMP, CALL etc. with the corresponding memory address from the label-address table.

The target code after the label translation is shown in the figure 2.

Implementing Label Translation

The above two pass translation process can be implemented using a single Lex program. The following functions provided by Lex can be useful while implementing the parser :

1. `yylless(k)` : Returns all but the first `n` characters of the current token back to the input stream, where they will be rescanned when the scanner looks for the next match. `yytext` and `yyleng` are adjusted appropriately (e.g., `yyleng` will now be equal to `n`).

For example, on the input "foobar" the following will write out "foobarbar":

```
%%
```

```
foobar ECHO; yyless(3);
```

```
[a-z]+ ECHO;
```

An argument of 0 to `yyless` will cause the entire current input string to be scanned again. Unless you've changed how the scanner will subsequently process its input, this will result in an endless loop. Note that `yyless` is a macro and can only be used in the flex input file, not from other source files.

2. `yytext+k` : ignores the first k characters in `yytext`.