Time and Space Complexity

10 Dec 2021

Time and Space Complexity

- Efficiency
- Asymptotic notation
 - Big O
 - Big Ω
 - Big Θ
- What is time complexity?
- What is space complexity?
- Ascending order of Complexity
- Worst Accepted Algorithm
- Problems
- References

Efficiency

- How efficiency is the program you have written?
 - Time Complexity: How much time does it take program to complete?
 - Space Complexity: How much memory does this program use?
 - How do these complexities change as the amount of data changes?
 - E.g. From 1 to 10,000,000,000,000
 - What is the difference between the average case and worst case efficiency if any?

Asymptotic notation

Big O, Ω and Θ are formal notational methods for stating the growth of resource needs (efficiency and storage) of an algorithm.

Big O - Worst case

- Upper bound of an algorithm
- Rate of growth of an algorithm is less than or equal to a specific value

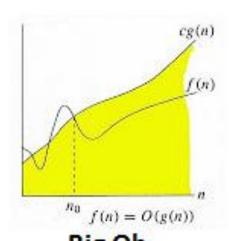
Big Ω Omega – Best case

- Lower bound of an algorithm
- Rate of growth is greater than or equal to a specified value

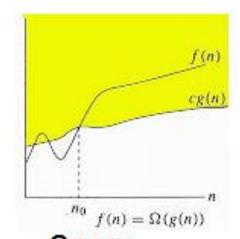
Big Θ Theta – Average case

- Tight bound of an algorithm
- Rate of growth is equal to a specified value

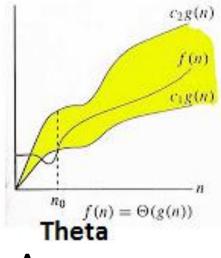
Asymptotic notation



Big Oh Worst case



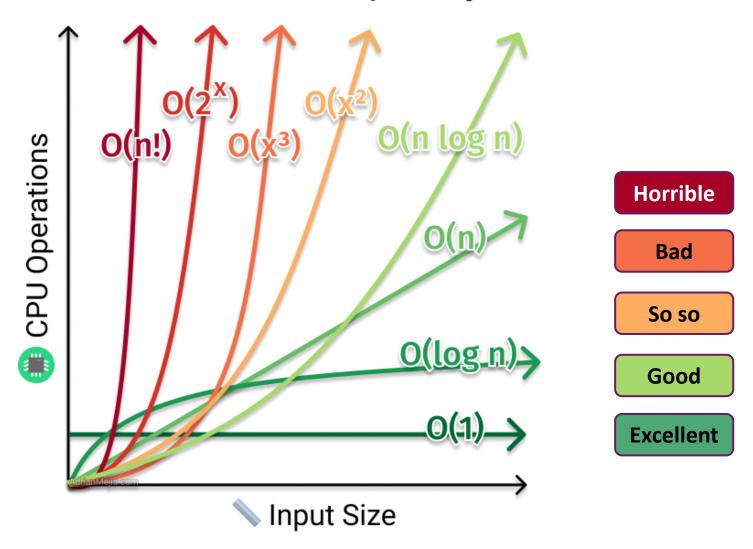
Omega Best case



Average case

Time Complexity

Time Complexity



Time Complexity

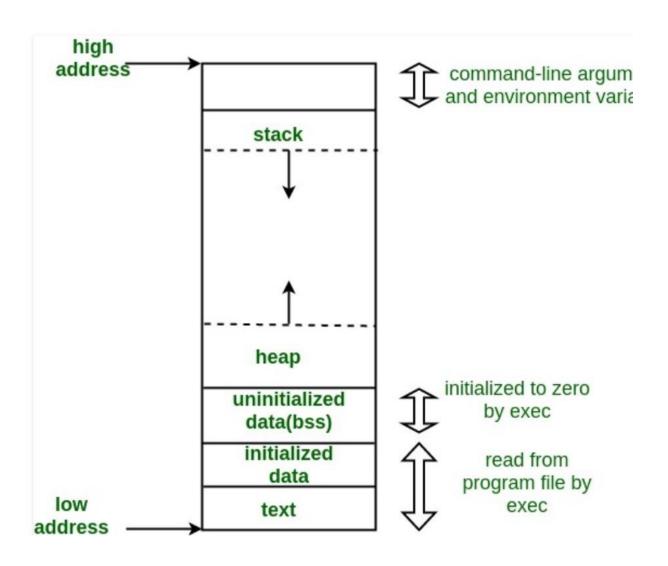
Assume N = 100,000 and processor speed is 1,000,000,000 operations per second

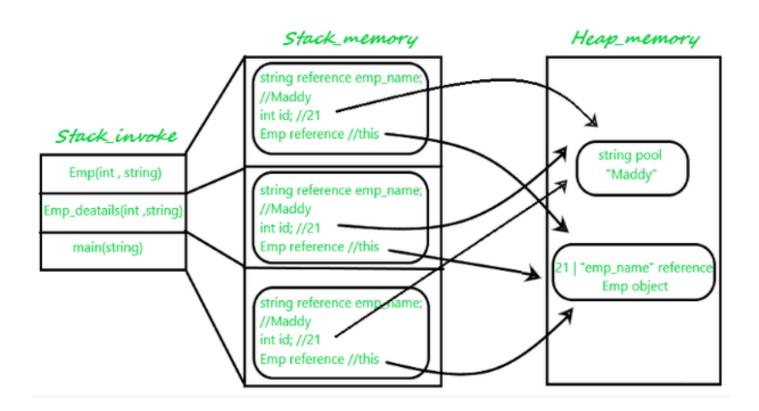
Function	Running Time
2 ^N	3.2 x 10 ^{30,086} years
N ⁴	3171 years
N^3	11.6 days
N^2	10 seconds
N N	0.032 seconds
N log N	0.0017 seconds
N	0.0001 seconds
N	3.2 x 10 ⁻⁷ seconds
log N	1.2 x 10 ⁻⁸ seconds

- The amount of memory used by the algorithm (including the input values to the algorithm) to execute and produce the result.
- Program instruction instruction space
 - amount of memory used to save the compiled version of instructions
- Environmental stack
 - One function calls another functions
 - Program stores the current variables to the system stack, while waiting for further execution

Data space

- Variables space used by variables and constants constant values, temporary values
- E.g. int , const, let , etc.





```
#include "holberton.h"
int main(void)
                                                        HEAP
        _puts_recursion("Holberton");
        return (0);
void _puts_recursion(char *s)
                                                          Stack
        if (*s == '\0')
               _putchar('\n');
       _putchar(*s);
        _puts_recursion(++s);
                                                          unitribiatized
                                                          initialited
                                                          +04+/(all
<.c" 16L, 199C written 16,3
                                    All
```

Ascending order of Complexity

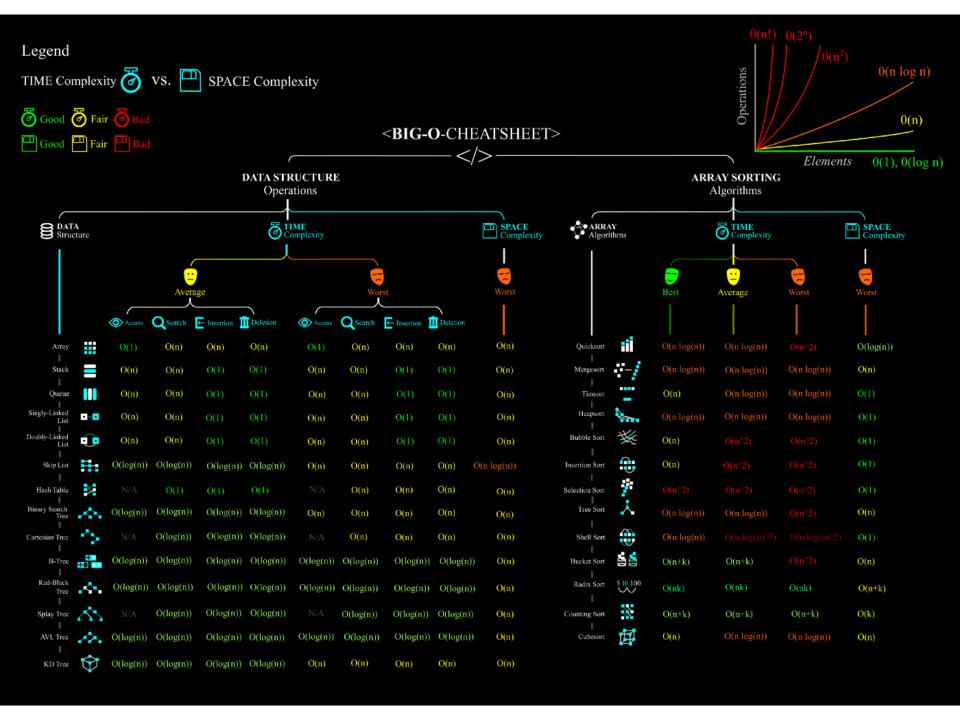
Function	Common Name
N!	factorial
2 ^N	Exponential
N ^d , d > 3	Polynomial
N^3	Cubic
N^2	Quadratic
N N	N Square root N
N log N	N log N
N	Linear
N	Root - n
log N	Logarithmic
1	Constant

Running time grows 'quickly' with more input.

Running time grows 'slowly' with more input.

Worst Accepted Algorithm

Length of Input (N)	Worst Accepted Algorithm
\leq [1011]	$O(N!), O(N^6)$
\leq [1518]	$O(2^N*N^2)$
\leq [1822]	$O(2^N*N)$
≤ 100	$O(N^4)$
≤ 400	$O(N^3)$
$\leq 2K$	$O(N^2 * log N)$
$\leq 10K$	$O(N^2)$
$\leq 1M$	O(N*logN)
$\leq 100M$	O(N), O(logN), O(1)



Data Structures Time and Space Complexity

Data Structures

Data Structure	Time Comple	Time Complexity							Space Complexity Worst
	Average				Worst				
	Indexing	Search	Insertion	Deletion	Indexing	Search	Insertion	Deletion	
Basic Array	0(1)	0(n)			0(1)	0(n)			0(n)
Dynamic Array	0(1)	(n)	O(n)	0(n)	0(1)	(n)	O(n)	O(n)	O(n)
Singly-Linked List	(n)	0(n)	0(1)	0(1)	(n)	0(n)	0(1)	0(1)	O(n)
Doubly-Linked List	(n)	0(n)	0(1)	0(1)	0(n)	0(n)	0(1)	0(1)	0(n)
Skip List	0(log(n))	0(log(n))	0(log(n))	O(log(n))	(n)	0(n)	0(n)	O(n)	O(n log(n))
Hash Table	•	0(1)	0(1)	0(1)		(n)	(n)	(n)	(O(n)
Binary Search Tree	0(log(n))	0(log(n))	0(log(n))	0(log(n))	(n)	0(n)	O(n)	(n)	0(n)
Cartresian Tree		0(log(n))	0(log(n))	O(log(n))		0(n)	0(n)	(n)	(O(n)
B-Tree	0(log(n))	O(log(n))	0(log(n))	0(log(n))	0(log(n))	O(log(n))	0(log(n))	O(log(n))	0(n)
Red-Black Tree	0(log(n))	0(log(n))	0(log(n))	0(log(n))	O(log(n))	0(log(n))	0(log(n))	0(log(n))	0(n)
Splay Tree		0(log(n))	0(log(n))	0(log(n))		0(log(n))	0(log(n))	0(log(n))	0(n)
AVL Tree	O(log(n))	0(log(n))	O(log(n))	0(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	0(n)

Array: Sorting-Time and Space Complexity

Algorithm		Space Complexity			
	Best	Average Worst		Worst	
Quicksort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	O(n^2)	O(log(n))	
Mergesort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	O(n log(n))	0(n)	
Timsort	$\Omega(n)$	$\Theta(n \log(n))$	O(n log(n))	0(n)	
Heapsort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	O(n log(n))	0(1)	
Bubble Sort	$\Omega(n)$	⊕ (n^2)	O(n^2)	0(1)	
Insertion Sort	$\Omega(n)$	⊕ (n^2)	O(n^2)	0(1)	
Selection Sort	$\Omega(n^2)$	⊕ (n^2)	O(n^2)	0(1)	
Tree Sort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	O(n^2)	0(n)	
Shell Sort	$\Omega(n \log(n))$	$\Theta(n(\log(n))^2)$	O(n(log(n))^2)	0(1)	
Bucket Sort	Ω (n+k)	⊖ (n+k)	O(n^2)	O(n)	
Radix Sort	Ω(nk)	⊖(nk)	O(nk)	O(n+k)	
Counting Sort	Ω (n+k)	⊕ (n+k)	0(n+k)	0(k)	
Cubesort	$\Omega(n)$	$\Theta(n \log(n))$	O(n log(n))	0(n)	

Horrible Bad Fair Good Excellent

```
int a = 0, b = 0;
for (i = 0; i < N; i++) {
    a = a + rand();
}
for (j = 0; j < M; j++) {
    b = b + rand();
}</pre>
```

- 1. O(N * M) time, O(1) space
- 2. O(N + M) time, O(N + M) space
- 3. O(N + M) time, O(1) space
- 4. O(N * M) time, O(N + M) space

```
int a = 0, b = 0;
for (i = 0; i < N; i++) {
    a = a + rand();
}
for (j = 0; j < M; j++) {
    b = b + rand();
}</pre>
```

- 1. O(N * M) time, O(1) space
- 2. O(N + M) time, O(N + M) space
- 3. O(N + M) time, O(1) space
- 4. O(N * M) time, O(N + M) space

```
int a = 0;
for (i = 0; i < N; i++) {
    for (j = N; j > i; j--) {
        a = a + i + j;
    }
}
```

- 1. O(N)
- 2. O(N*log(N))
- 3. O(N * Sqrt(N))
- 4. O(N*N)

```
int a = 0;
for (i = 0; i < N; i++) {
    for (j = N; j > i; j--) {
        a = a + i + j;
    }
}
```

- 1. O(N)
- 2. O(N*log(N))
- 3. O(N * Sqrt(N))
- 4. O(N*N)

```
int i, j, k = 0;
for (i = n / 2; i <= n; i++) {
    for (j = 2; j <= n; j = j * 2) {
        k = k + n / 2;
    }
}</pre>
```

- 1. O(n)
- 2. O(nLogn)
- 3. O(n^2)
- 4. O(n^2Logn)

```
int i, j, k = 0;
for (i = n / 2; i <= n; i++) {
    for (j = 2; j <= n; j = j * 2) {
        k = k + n / 2;
    }
}</pre>
```

- 1. If n = 32
- 2. Outer loop
 - 1. i = 32/2 = 16, i will be incremented by 1 (n)
- 3. Inner loop
 - 1. $J = 2, j \le n; j = 2 * 2 = 4 (log n)$
 - 2. K = k + n/2

- 1. O(n)
- 2. O(nLogn)
- 3. $O(n^2)$
- 4. O(n^2Logn)

```
int a = 0, i = N;
while (i > 0) {
    a += i;
    i /= 2;
}
```

- 1. O(N)
- 2. O(Sqrt(N))
- 3. O(N/2)
- 4. O(log N)

```
int a = 0, i = N;
while (i > 0) {
    a += i;
    i /= 2;
}
```

- 1. O(N)
- 2. O(Sqrt(N))
- 3. O(N/2)
- 4. O(log N)

1. If
$$a = 0$$
, $i = 32$

- 2. i > 0
- 3. a= 32
- 4. i = 32/2 = 16
- 1. i > 0
- 2. a= 48
- 3. i = 16/2 = 8
- 1. i > 0
- 2. a= 56
- 3. i = 8/2 = 4

1.
$$i > 0$$

$$2. a = 60$$

3.
$$i = 4/2 = 2$$

- 1. i > 0
- 2. a= 62
- 3. i = 2/2 = 1

1.
$$i > 0$$

3.
$$i = 1/2 = 0$$

```
int a = 0, i = N;
while (i > 0) {
    a += i;
    i /= 2;
}
```

- 1. O(N)
- 2. O(Sqrt(N))
- 3. O(N/2)
- 4. O(log N)

912. Sort an Array

Given an array of integers nums, sort the array in ascending order.

Example 1:

Input: nums = [5,2,3,1]
Output: [1,2,3,5]

Example 2:

Input: nums = [5,1,1,2,0,0]
Output: [0,0,1,1,2,5]

Constraints:

- 1 <= nums.length <= 5 * 10⁴
- $-5 * 10^4 \le nums[i] \le 5 * 10^4$

Problem - 5 Leetcode

(3) Sort an Array - LeetCode

Problem - 5 Leetcode

```
// Bubble sort
// In i-th pass of Bubble Sort (ascending order),
//last (i-1) elements are already sorted
// i-th largest element is placed at (N-i)-th position
class Solution {
  public int[] sortArray(int[] nums) {
     for(int i = 0; i < nums.length; i++){
      for(int j = 0; j < nums.length-1; j++){
         if (nums[j] > nums[j+1]){
           int temp = nums[j];
           nums[j] = nums[j+1];
           nums[i+1] = temp;
    return nums;
```

```
// Selection Sort
// It divides the array into two parts:
     -- sorted (left) and unsorted (right) subarray.
// It repeatedly selects the next smallest element.
class Solution {
  public int[] sortArray(int[] nums) {
    for(int i = 0; i < nums.length; i++){
       int minIndex = i;
      for(int j = i + 1; j < nums.length; j++){
         if (nums[j] < nums[minIndex]){</pre>
           int temp = nums[minIndex];
           nums[minIndex] = nums[j];
           nums[j] = temp;
    return nums;
```

Problem - 5 Leetcode

```
// Insertion Sort
// Compare current element temp to its predececessor
// If key < , compare it to the elements before
// Move the greater elements one position up
class Solution {
  public int[] sortArray(int[] nums) {
     for(int i = 1; i < nums.length; i++){
       int temp = nums[i];
       int j = i - 1;
       // Move elements of nums, that are greater than temp to
       // one position ahead of their current position
       while(j \ge 0 \&\& nums[j] > temp)
          nums[i+1] = nums[i];
          j--;
       nums[j+1] = temp;
     return nums;
```

Problem - 5 Leetcode

```
// Insertion Sort
// Compare current element temp to its predececessor
// If key < , compare it to the elements before
// Move the greater elements one position up
class Solution {
  public int[] sortArray(int[] nums) {
     for(int i = 1; i < nums.length; i++){
       int temp = nums[i];
       int j = i - 1;
       // Move elements of nums, that are greater than temp to
       // one position ahead of their current position
       while(j \ge 0 \&\& nums[j] > temp)
          nums[i+1] = nums[i];
          j--;
       nums[j+1] = temp;
     return nums;
```

```
class Solution {
    public int[] sortArray(int[] nums) {
        mergeSortRecursive(nums, 0, nums.length -1);
        return nums;
    private static void mergeSortRecursive(int[] nums, int low, int high){
        if (high - low + 1 <= 1) return;
        if (mid = low + (high - low)/2);
        mergeSortRecursive(nums, low, mid);
        mergeSortRecursive(nums, mid+1, high);
        merge(nums, low, mid, high);
    private static void merge(int[] nums, int low, int mid, int high){
        int[] temp = new int[high - low+1];
        int i = low;
        int j = mid+1;
        int tempIndex = 0;
        while(i <= mid && j <= high){</pre>
            if(nums[i] < nums[j]){</pre>
                temp[tempIndex++] = nums[i++];
            }else{
                temp[tempIndex++] = nums[j++];
        while(i <= mid){</pre>
            temp[tempIndex++] = nums[i++];
        }
        while(j <= high){</pre>
            temp[tempIndex++] = nums[j++];
        for(int x = low; x <= high; x++){
            nums[x++] = temp[x-low];
        }
```

Problem - 5 Leetcode

Array: Sorting-Time and Space Complexity

Algorithm		Space Complexity		
	Best	Average	Worst	Worst
Quicksort	$\Omega(n \log(n))$	$\Theta(n \log(n))$ $O(n^2)$		O(log(n))
Mergesort	Ω(n log(n))	$\Theta(n \log(n))$	O(n log(n))	O(n)
Timsort	Ω(n)	$\Theta(n \log(n))$	O(n log(n))	0(n)
Heapsort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	O(n log(n))	0(1)
Bubble Sort	Ω(n)	⊕(n^2)	O(n^2)	0(1)
Insertion Sort	Ω(n)	⊕(n^2)	O(n^2)	0(1)
Selection Sort	Ω(n^2)	⊖(n^2)	O(n^2)	0(1)
Tree Sort	Ω(n log(n))	$\Theta(n \log(n))$	O(n^2)	0(n)
Shell Sort	$\Omega(n \log(n))$	$\Theta(n(\log(n))^2)$	O(n(log(n))^2)	0(1)
Bucket Sort	Ω(n+k)	⊕ (n+k)	O(n^2)	0(n)
Radix Sort	Ω(nk)	<mark>⊕(nk)</mark>	O(nk)	0(n+k)
Counting Sort	Ω (n+k)	⊕ (n+k)	O(n+k)	0(k)
Cubesort	$\Omega(n)$	$\Theta(n \log(n))$	O(n log(n))	0(n)

Horrible Bad Fair Good Excellent

References

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- Practice Questions on Time Complexity Analysis, https://www.geeksforgeeks.org/practice-questions-time-complexity-analysis/