

ALICE O² data model proposal

Mikolaj Krzewicki Sylvain Chapeland Roberto Divia
Matthias Richter David Rohr

for the CWG4 data model group.

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Abstract

This note presents a proposal for the ALICE O² data model. A base data layout and metadata format that allow for efficient resource use are proposed. Application of the data model to online/offline data processing and quality control is discussed.

1 Introduction

The ALICE online-offline (O²) computing system [1, 2] is a computing facility and a software framework designed for the processing of the ALICE data in the upcoming LHC Run 3. The design aims at high data throughput and parallelism using a multiprocess model. It does not, however, exclude the use of multithreading and other forms of concurrent processing inside of individual processes.

The data exchange between processes running within the O² system (called O² devices) is taken care of by the ALICE-FAIR (Alfa) framework [3]. Since this is the only communication mechanism foreseen for data exchange, it effectively serves the role of an API between the devices. The Alfa framework provides data transport and synchronisation primitives via the FairMQ message queue library. FairMQ messages consist of raw memory buffers which are asynchronously queued and atomically delivered.

The online data processed by the O² system consists of a set of data buffers originating from the detector hardware (raw data) and the processing devices (derived data). The data fragments are logically grouped into (sub-) time frames. A (sub-) time frame contains raw data associated to a period of data taking (typically several tens of ms, as dictated by the heartbeat trigger [4]) and/or the results of the processing of these data. In addition any data that might be necessary to describe and qualify the data set can be added to the logical group.

In the unified online-offline software model also derived data is handled within the same software framework. O² devices dedicated to quality control (QC) and physics analysis tasks should use the same set of interfaces as pure online components. The requirements for the data used in these tasks tend to be different from the online components: high level abstractions and ease of use (of e.g. ROOT[5] objects) is sometimes preferred to high performance low-level data structures. Transparent support for high level data structures is part of the proposed data model.

A single time frame data volume is expected to be of the order of tens of gigabytes. The data model facilitates communication approaches that minimize resource use, i.e. avoid unnecessary copies of data.

In order to assure consistent navigation within a time frame, each data fragment is described by a small metadata block containing the information about the content type of the payload, its origin and serialization strategy. In addition, the metadata block can be extended by the processing devices with additional information without the need to modify the payload.

2 Vectored IO

Vectored IO (also referred to as scatter/gather IO) is an important feature when dealing with multiple data buffers as it allows, in principle, to avoid the cost associated with serializing data into a single IO buffer. Vectored IO is provided in FairMQ in the form of multi-part messages. A multi-part message consists of multiple independent buffers and is delivered atomically as a whole while preserving the ordering of the buffers.

The multi-part approach, in addition to minimizing the resource strain associated with IO buffer construction, also by construction reduces the need for (re-)synchronisation and event building; data fragments once associated to a single time frame remain that way throughout the entire chain regardless of the networking topology. Another benefit is that additional data parts can be attached or removed by other processing devices without copy overhead at any point of the processing chain.

3 O² Message structure

The O² message consists of a sequence of metadata-payload pairs contained within a multi-part message. Each payload is described by metadata contained in a separate message part. Since FairMQ preserves the ordering of the parts, the natural choice is to precede each data part with the associated metadata part in the message, as illustrated in figure 1.

Storing the metadata in separate buffers offers several advantages:

- Since the metadata is separated from the payload already at transport level, efficient navigation is possible as only the (small) metadata parts need to be inspected.
- The size of the metadata buffer is not fixed enabling a scheme with flexible metadata content.
- The content of the data buffers, once produced by the hardware or a processing device, is immutable to other devices. Since the metadata is encapsulated in a separate buffer, it becomes possible to add additional information to the metadata with minimal cost and without modifying the payload downstream from the data producer.

4 Metadata format

The O² metadata consists of a contiguous buffer containing a sequence of headers (header stack). The byte representation of a header consists of a user defined body following a BaseHeader struct containing fields needed to:

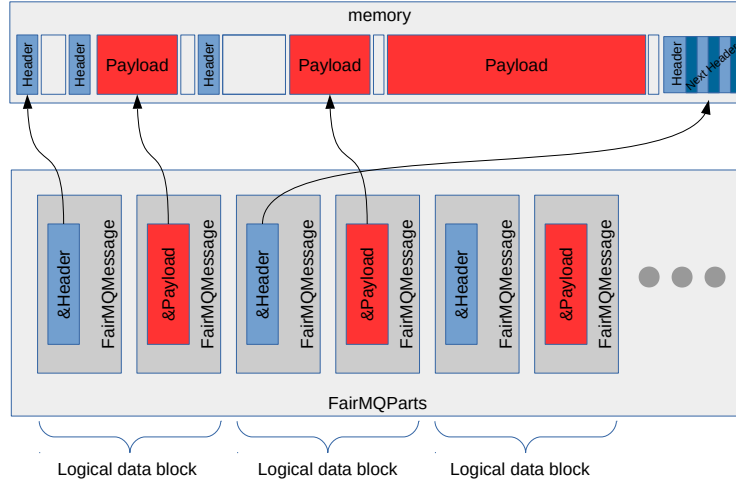


Figure 1: The O² message structure. Logical data blocks consisting of payload and metadata parts are contained in a multi-part FairMQ message (bottom). The buffers associated to the message parts do not have to be continuous in memory (top).

- 87 • Verify that the following data (header body) belongs to an O² header.
- 88 • Define the size of the entire header.
- 89 • Flag whether another header follows this one in the stack.
- 90 • Verify the version of the header.
- 91 • Signal the type of the header.
- 92 • Signal the packing/serialization scheme of the metadata body carried
- 93 by the header.

94 The header stack should contain at least the standard DataHeader
 95 struct describing the basic payload properties common to all payloads.
 96 The DataHeader representation starts with the BaseHeader followed by
 97 bytes representing:

- 98 • The functional data description uniquely determining the data type
- 99 contained in the payload.
- 100 • The payload serialization method (e.g. ROOT, FlatBuffers, none)
- 101 complementing the data description field.
- 102 • The origin of the data to identify the producer (e.g. detector system
- 103 or a software subsystem).
- 104 • A data type dependent 64 bit specification. This can be used by the
- 105 detectors to store e.g. the fine grained equipment ID like the link
- 106 number for raw data or cluster finder instance for clusters¹.

¹Based on HLT experience: most data types use some kind of fine grained ID. This field has been added here to avoid the overhead of a full header for what in most cases would be just one (64 bit) field.

107 • The payload buffer size².

108 Other headers can be defined similarly to DataHeader and included
109 in the header stack. Examples include trigger information for triggered
110 detector data, object name for ROOT objects used in quality control etc.

111 4.1 Header memory layout

112 The memory layout of the DataHeader struct is illustrated by the following
113 definition:

```
114       struct BaseHeader  
115       {  
116           uint32_t    magicString;  
117           uint32_t    headerSize;  
118           uint32_t    flags;  
119           uint32_t    headerVersion;  
120           uint64_t    headerDescription;  
121           uint64_t    headerSerialization;  
122       };  
123  
124       struct DataHeader : public BaseHeader  
125       {  
126           uint64_t    dataDescription[2];  
127           uint32_t    dataOrigin;  
128           uint32_t    reserved;  
129           uint64_t    payloadSerializationMethod;  
130           uint64_t    subSpecification;  
131           uint64_t    payloadSize;  
132       };
```

133 4.2 Header interface

134 Data members representing (in principle) arbitrary integers (headerSize,
135 payloadSize, subSpecification) or bitfields (flags) are directly accessible
136 and settable as integers.

137 Access to other members is protected by strongly typed interfaces. Ini-
138 tialization can only be performed using predefined constants, consistency
139 is then assured by the type system at compile time.

140 Access to header information contained in a metadata buffer is also
141 implemented in a type safe way.

142 4.2.1 Example

143 The user may check if a given buffer contains the desired header informa-
144 tion in a type safe way. If the desired header is of type DataHeader, a call
145 to:

```
146       DataHeader* header = Header::get<DataHeader>(buffer);
```

147 will yield a valid pointer if a header of type DataHeader is part of the
148 header stack inside the buffer pointed to by the buffer pointer.

²This is not strictly necessary online as the transport framework keeps track of the buffer sizes. Keeping this information in the header is useful for persistent storage and debugging purposes.

5 Data formats

The O^2 data model does not impose any limitations on the data types exchanged between devices. The only constraint from the data transport layer is that the transported payload must be contiguous in memory. For higher level data types it usually means that they need to be serialized which penalizes performance (to a varying degree).

6 Data in memory

The recommendation is to use flat POD data types where performance and/or memory usage is critical. Flat data means naturally contiguous data that does not require a serialization step and does not contain on any process specific run time dependent information like virtual function table pointers or pointer/reference members.

Outside of the critical path where only relatively a low volume of data needs to be transported between devices, serialization schemes possibly impose acceptable overhead. The data model supports serialization schemes and facilitates transparent handling of serialized data at the user code level.

6.1 Persistent storage

The in-memory data representation of header and payload buffers contains enough information to be stored on-disk directly as a sequence of buffers. The data represented by POD data are, however, not suitable for long term storage and need to be transformed to a portable and extensible format.

In the modular O^2 design only a dedicated device would handle persistent storage making the translation steps from and to persistent storage format transparent to other devices which only need to deal with the in-memory format.

7 Interfaces

The base functionality of a device running in the O^2 system is contained in the FairMQDevice class. This is a generic class providing all the necessary logic to support the O^2 system and data model, like control abstractions, a state machine and multi-part messaging. FairMQ support for the multi-part messaging is, however, unaware of the O^2 specific message layout. In order to enforce a consistent handling of the metadata the necessary interface is implemented on top of FairMQDevice as O2device.

Each device in the O^2 system should inherit from the O2device class in order to be able to use the data model. Interfaces are provided that insure consistency in the handling of the data and associated metadata.

7.1 Examples

To construct a message, the AddMessage(...) family of methods should be used to insure correct association of metadata to payloads, e.g. in the simplest case:

```
bool AddMessage( FairMQParts& message ,
```

```

192         AliceO2::Header::Block&& headerStack ,
193         FairMQMessagePtr payloadPart );

```

194 where in a single call the header stack to payload association is made and
195 data is appended to the message.

196 After the message is constructed a call to the standard FairMQ send
197 method suffices:

```

198 int64_t Send( const FairMQParts& parts ,
199              const std::string& channel );

```

200 This queues the entire message for atomic delivery via a specified data
201 channel.

202 For decoding the message on the receiving end functionality is provided
203 to access the metadata and it's associated payload in a consistent way.
204 The user implements a function that takes the metadata and the payload
205 buffers as input, and uses that function via e.g. a call to:

```

206 template <typename T>
207 bool ForEach(
208     O2message& parts ,
209     bool (T::*memberFunction)(
210         const byte* headerBuffer , size_t headerBufferSize ,
211         const byte* dataBuffer , size_t dataBufferSize
212     )
213 );

```

214 Inside the user provided function the metadata and the payload are de-
215 coded and processed.

216 References

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