

Specifying Semantics of Programming Languages in K Framework

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faculty of mathematics and physics

Overview

K k-framework

The K Semantic Framework

- University of Illinois at Urbana-Champaign
- Scheme, Verilog, Python, C, ...

C c-semantics

Semantics of C in the K Framework

- An Executable Formal Semantics of C with Applications
- Chucky Ellison, Grigore Rosu
- POPL 2012



Weverca

- Web VERifiCAtion for PHP
- GAUK project



K Framework - Overview

- Executable
 - Interpreter, Model checker, Deductive reasoning
 - Static analyser
- Rewriting-style semantics
 - Uses Maude engine (developed since 80')
 - Use of OCAML (fast execution) and CoQ (verification) planned

<http://k-framework.org/>

K Framework - Example

Original language syntax	K Strictness	K Semantics
Pgm ::= var List{Id}; Stmt		$\langle k \rangle \text{var } x1:\text{ListItem}; s:\text{Stmt} \Rightarrow s \langle /k \rangle$ $\langle \text{state} \rangle x1 \mapsto 0 \langle / \text{state} \rangle$
Stmt ::= Id = AExp Stmt; Stmt if BExp then Stmt else Stmt while BExp do Stmt spawn Stmt	[strict(2)] [strict(1)]	$\langle k \rangle x:\text{Id} = i:\text{Int} \Rightarrow . _ \langle /k \rangle$ $\langle \text{state} \rangle _ x \mapsto i _ \langle / \text{state} \rangle$ $s1:\text{Stmt}; s2:\text{Stmt} \Rightarrow s1 \sim s2$ $\text{if true then } s1:\text{Stmt} \text{ else } _ \Rightarrow s1$ $\text{if false then } _ \text{ else } s2:\text{Stmt} \Rightarrow s2$ $\text{while } B:\text{BExp} \text{ do } S:\text{Stmt} \Rightarrow$ $\text{if } B \text{ then } (S ; \text{while } B \text{ do } S) \text{ else } .$ $\langle k \rangle \text{spawn } s:\text{Stmt} \Rightarrow . _ \langle /k \rangle . \Rightarrow \langle k \rangle s \langle /k \rangle$
BExp ::= Bool		
AExp ::= Int Id AExp + AExp	[strict]	$\langle k \rangle x:\text{Id} \Rightarrow i _ \langle /k \rangle$ $\langle \text{state} \rangle _ x \mapsto i _ \langle / \text{state} \rangle$ $i1:\text{Int} + i2:\text{Int} \Rightarrow i1 + \text{Int } i2$
KResult ::= Int Bool		

```
Configuration ≡  <k*>K</k*>
                  <state>Map{Id |-> Int}</state>
```



C Semantics

First complete formal semantics for C

Conforming

- Must accept all portable programs, but can also accept non portable programs

Freestanding

- All language features except complex numbers
- Subset of the standard library

Extensively tested

- E. g. against GCC torture tests (1093 test programs, 776 standard compliant. Of those 770 passed)
- Better results than Clang or GCC

<http://code.google.com/p/c-semantics/>

C Semantics – Features

Feature	GH	CCR	CR	Definition		BL	Le	ER
				No	Pa			
Bitfields	●	◐	○	○	◐	○	○	●
Enums	◐	●	○	○	●	○	○	●
Floats	○	○	○	○	◐	●	●	●
Struct/Union	●	●	●	◐	●	●	●	●
Struct as Value	○	○	○	●	○	○	○	●
Arithmetic	◐	●	●	○	●	●	●	●
Bitwise	○	●	○	○	●	●	●	●
Casts	◐	◐	○	◐	◐	●	●	●
Functions	●	●	◐	●	●	●	●	●
Exp. Side Effects	●	●	○	●	●	○	●	●
Variadic Funcs.	○	○	○	○	○	○	○	●
Eval. Strategies	○	◐	○	●	●	○	●	●
Concurrency	○	○	○	○	○	○	○	◐
Break/Continue	◐	●	◐	●	●	●	●	●
Goto	◐	○	○	○	●	○	●	●
Switch	◐	●	○	○	●	◐	◐	●
Longjmp	○	○	○	○	○	○	○	●
Malloc	○	○	○	○	○	○	○	●

●: Fully Described

◐: Partially Described

○: Not Described

GH denotes *Gurevich and Huggins* (1993),
 CCR is *Cook, Cohen, and Redmond* (1994),
 CR is *Cook and Subramanian* (1994),
 No is *Norrish* (1998),
 Pa is *Papaspyrou* (2001),
 BL is *Blazy and Leroy* (2009),
 Le is *Leroy* (unpublished, 2010), and
 ER is *Ellison and Roşu* (our work).

C Semantics – Implementation

- 75 cells
- 150 syntactic operators
- 5900 source lines of semantics
- 1200 rules
 - 80 rules for statements
 - 160 rules for expressions
 - 500 rules for declarations and types
 - 115 rules for standard library
 - ...
- 6 person-months



Existing C analysis/verification tools

- Lint/Purify/Coverity/Valgrind
- Blast
- Havoc
- Slam
- VCC
- ...

Based on approximative models of C

- Hard to argue for the soundness of the tools

C Semantics – Problems of Approximative models (1)

```
int main(void) {  
    int x = 0;  
    return (x = 1) + (x = 2);  
}
```

Undefined according to C standard

GCC4, MSVC returns 4

GCC3, ICC, Clang returns 3

Frama-C „proves“ it returns 4

C Semantics – Problems of Approximative models (2)

```
int main(void) {  
    "foo"[0] = 'x';  
    return "foo"[0];  
}
```

Undefined according to C standard

GCC4	doesn't compile
ICC, Clang	segmentation fault

Frama-C „proves“ it returns ,x‘

C Semantics – Problems of Approximative models (3)

```
int r;  
  
int f(int x) {  
    return (r = x);  
}  
  
int main(void) {  
    f(1) + f(2); return r;  
}
```

Defined (Could return 1 or 2)

GCC, ICC, MSVC, Clang returns 2

Both Frama-C and Havoc „prove“ it can only return 2

C Semantics – Built tools

Use an explicit and testable definition to build tools that conform to this semantics

Semantics-Based Analysis Tools

- Interpreter
- State-space explorer
- LTL Model-checker
- Debugger
- Program verifier

C Semantics – Interpreter

```
#include <string.h>
int main(void) {
    char dest[5], src[5] = "hello";
    strcpy(dest, src);
}
```

ERROR! KCC encountered an error
while executing this program.
Description: Reading outside the
bounds of an object.
File: buggy_strcpy.c
Function: strcpy
Line: 4

C Semantics – State-space Search (1)

```
int denominator = 5;
int setDenominator(int d) {
    return denominator = d;
}
int main(void) {
    return setDenominator(0) + (7 / denominator);
}
```

2 solutions found

Solution 1

Program got stuck

File: eval_order.c

Line: 8

Description: Division by 0.

Solution 2

Program completed successfully

Return value: 1

C Semantics – LTL-Based Model Checking

```
typedef enum {green, yellow, red} state;
state lightC = green; state lightW = red;
int changeC() {
    switch (lightC) {
        case(green): lightC = yellow; return 0;
        case(yellow): lightC = red; return 0;
        case(red):
            if (lightW == red) { lightC = green; } return 0;
    }
}
...
int main(void) { while(1) { changeC() + changeW(); } }
```

changeC, changeW => lightC=yel, lightW=red

changeW, changeC => lightC=red, lightW=yel

changeC, changeW => lightC=gre, lightW=red

changeC(); changeW();

```
#pragma __ltl safety: [] (lightC == red /\ lightW == red)
#pragma __ltl progressC: [] <> (lightC == green)
```

C Semantics – Deductive Verification

```
void listNode* reverse(struct listNode *x) {
    struct listNode *p;
    struct listNode *y;
    p = 0;
    while (x) {
        y = x->next;
        x->next = p;
        p = x;
        x = y;
    }
}
```

```
rule <k> $ => return p; </k>
    <heap>... list(x,A) => list(p,rev(A)) ...</heap>
```

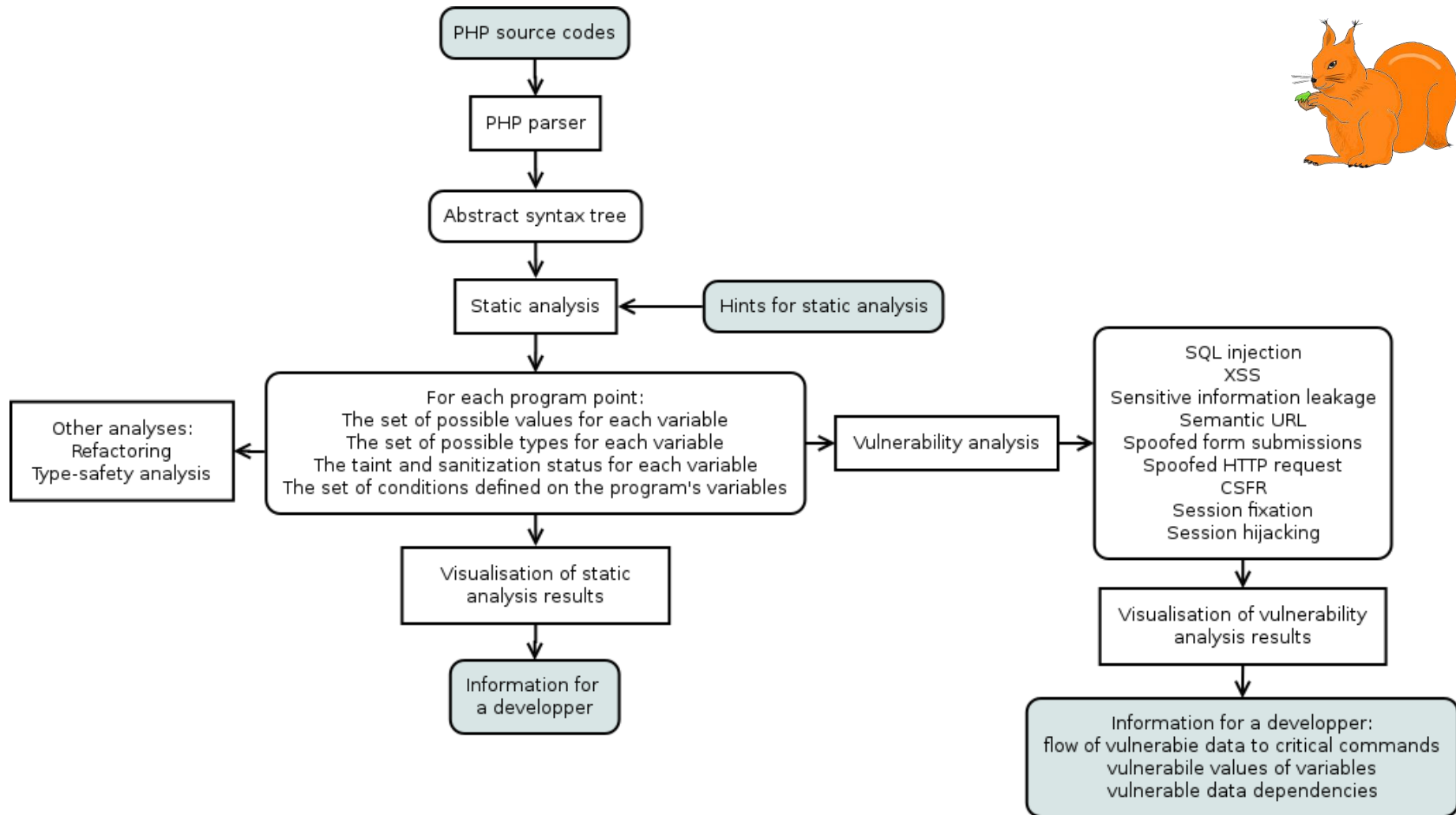
\$ is the body of the function



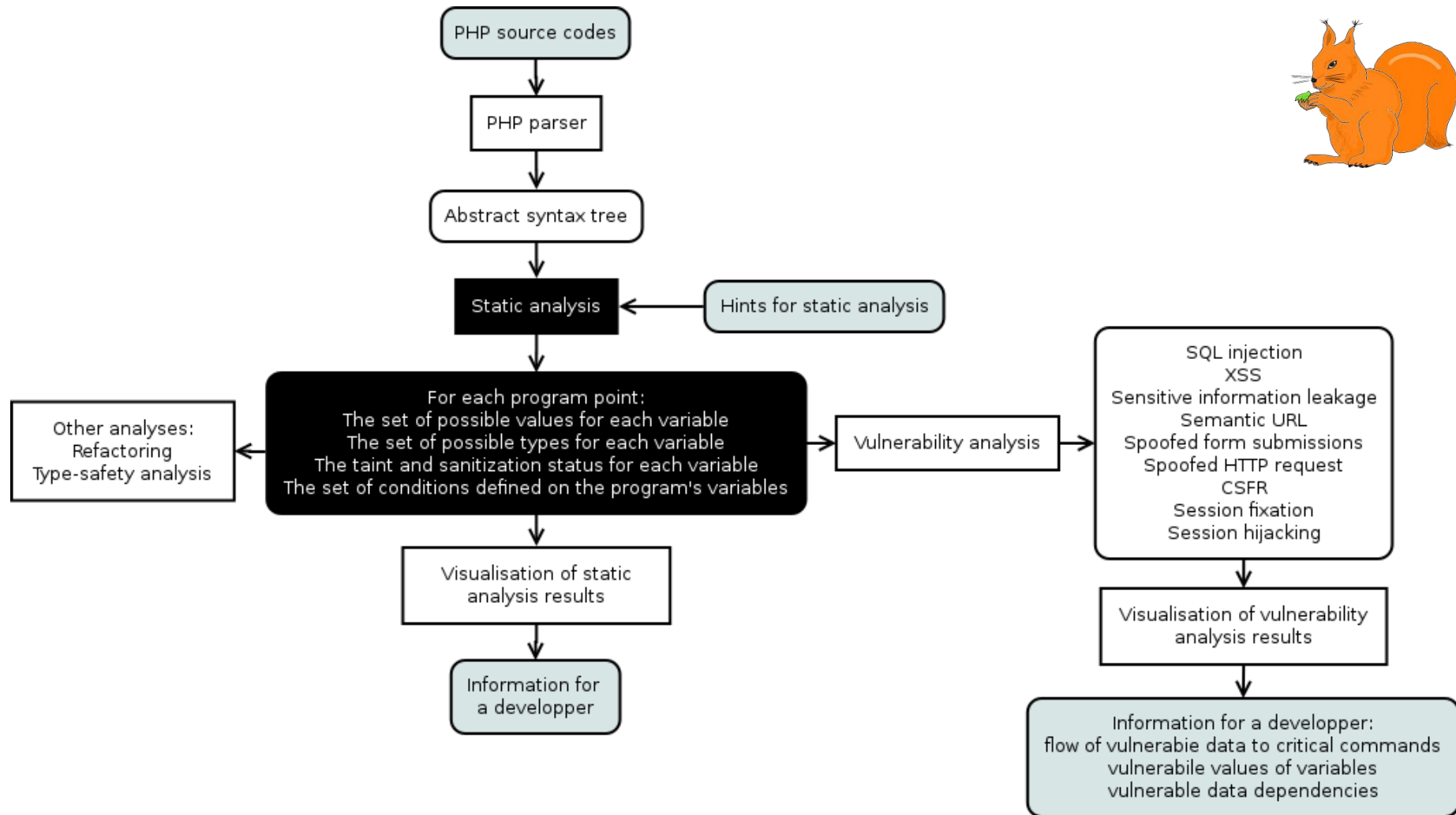
Weverca

- Security Analysis of PHP applications
 - Flow of sensitive data into critical commands

Weverca – Overview



Weverca – Work in Progress



Summary

K k-framework

The K Semantic Framework

- Defining semantics of programming languages

C c-semantics

Semantics of C in the K Framework

- Semantics-based analysis tools



Weverca

- Static analyzer in K