

p	atomic proposition		
l	proof entry label		
n	index variable (subscript)		
$claim$	$::=$	claim	
	$judgment\ \rho$		
$judgment$	$::=$	judgment	
	$\bar{\phi} \vdash \phi$		
ϕ	$::=$	proposition	
	p	atomic	
	$\neg\phi$	negation	
	$\phi \wedge \phi'$	conjunction	
	$\phi \vee \phi'$	disjunction	
	$\phi \rightarrow \phi'$	implication	
	\perp	contradiction	
	(ϕ)	S	
$\bar{\phi}$	$::=$	list of propositions	
	$\phi_1, .., \phi_n$	M	
ρ	$::=$	proof	
	$entry_1 .. entry_n$		
	$()$	M	
	$entry\ \rho$	M	
$entry$	$::=$	proof entry	
	$derivation$	line	
	$[\rho]$	box	
	$invalid$		
	$last(\rho)$	M	last entry in proof
$derivation$	$::=$	derivation in proof	
	$l\ \phi\ reason$		
$reason$	$::=$		
	$assumption$		
	$justification$		
$justification$	$::=$	derivation justification	
	$premise$	premise	
	LEM	law of excluded middle	
	$copy\ l$	copying	
	$\wedge i\ l, l'$	conjunction introduction	
	$\wedge e_1\ l$	conjunction elimination	
	$\wedge e_2\ l$	conjunction elimination	
	$\vee i_1\ l$	disjunction introduction	
	$\vee i_2\ l$	disjunction introduction	
	$\rightarrow e\ l, l'$	implication elimination	
	$\neg e\ l, l'$	negation elimination	

		$\perp e\ l$	contradiction elimination
		$\neg\neg i\ l$	double negation introduction
		$\neg\neg e\ l$	double negation elimination
		MT l, l'	modus tollens
		$\rightarrow i\ l - l'$	implication introduction
		$\neg i\ l - l'$	negation introduction
		$\vee e\ l_1, l_2 - l_3, l_4 - l_5$	disjunction elimination
		PBC $l - l'$	proof by contradiction
<i>terminals</i>	::=		
		premise	
		LEM	
		copy	
		$\wedge i$	
		$\wedge e_1$	
		$\wedge e_2$	
		$\vee i_1$	
		$\vee i_2$	
		$\rightarrow e$	
		$\neg e$	
		$\perp e$	
		$\neg\neg i$	
		$\neg\neg e$	
		MT	
		$\rightarrow i$	
		$\neg i$	
		$\vee e$	
		PBC	
		assumption	
		\vdash	
		\neg	
		\wedge	
		\vee	
		\rightarrow	
		\perp	
		\in	
		\mapsto	
		last	
		invalid	
<i>dyadicprop</i>	::=		
		ϕ	M
		(ϕ, ϕ')	M
Γ	::=		admissible context
		\square	M
		$\Gamma[l \mapsto \phi]$	M

		$\Gamma[(l, l') \mapsto (\phi, \phi')]$	M	
<i>formula</i>	::=			
		<i>judgement</i>		judgement
		$\phi \in \bar{\phi}$	M	proposition in set of premises
		$\Gamma(l) = \phi$	M	map label to prop
		$\Gamma(l, l') = (\phi, \phi')$	M	map tuple of labels to tuple of props
		<i>entry</i> = <i>entry'</i>	M	entry equality
<i>validity</i>	::=			
		<i>claim</i>		valid claim
		$\Gamma, \bar{\phi} \vdash \rho$		valid proof
		$\Gamma, \bar{\phi} \vdash \textit{derivation}$		valid derivation
<i>judgement</i>	::=			
		<i>validity</i>		
<i>user_syntax</i>	::=			
		<i>p</i>		
		<i>l</i>		
		<i>n</i>		
		<i>claim</i>		
		<i>judgment</i>		
		ϕ		
		$\bar{\phi}$		
		ρ		
		<i>entry</i>		
		<i>derivation</i>		
		<i>reason</i>		
		<i>justification</i>		
		<i>terminals</i>		
		<i>dyadicprop</i>		
		Γ		
		<i>formula</i>		

claim valid claim

$$\frac{\text{last}(\rho) = l \ \phi \ \textit{justification} \quad \bar{\phi}, \bar{\phi} \vdash \rho}{\bar{\phi} \vdash \phi \ \rho} \quad \text{VC_CLAIM}$$

$\Gamma, \bar{\phi} \vdash \rho$ valid proof

$$\frac{}{\Gamma, \bar{\phi} \vdash ()} \quad \text{VP_EMPTY}$$

$$\frac{\Gamma, \bar{\phi} \vdash l \ \phi \ \textit{justification} \quad \Gamma[l \mapsto \phi], \bar{\phi} \vdash \rho}{\Gamma, \bar{\phi} \vdash l \ \phi \ \textit{justification} \ \rho} \quad \text{VP_DERIVATION}$$

$$\begin{array}{c}
\text{last } (l \ \phi \text{ assumption } \rho) = l' \ \phi' \text{ reason} \\
\Gamma[l \mapsto \phi], \bar{\phi} \vdash \rho \\
\Gamma[(l, l') \mapsto (\phi, \phi')], \bar{\phi} \vdash \rho' \\
\hline
\Gamma, \bar{\phi} \vdash [l \ \phi \text{ assumption } \rho] \rho'
\end{array} \quad \text{VP_BOX}$$

$\Gamma, \bar{\phi} \vdash \text{derivation}$

valid derivation

$$\begin{array}{c}
\frac{\phi \in \bar{\phi}}{\Gamma, \bar{\phi} \vdash l \ \phi \text{ premise}} \quad \text{VD_PREMISE} \\
\\
\frac{}{\Gamma, \bar{\phi} \vdash l \ \phi \vee \neg \phi \text{ LEM}} \quad \text{VD_LEM} \\
\\
\frac{\Gamma(l') = \phi}{\Gamma, \bar{\phi} \vdash l \ \phi \text{ copy } l'} \quad \text{VD_COPY} \\
\\
\frac{\Gamma(l') = \perp}{\Gamma, \bar{\phi} \vdash l \ \phi \ \perp e \ l'} \quad \text{VD_CONTE} \\
\\
\frac{\Gamma(l_1) = \phi \quad \Gamma(l_2) = \phi'}{\Gamma, \bar{\phi} \vdash l \ \phi \wedge \phi' \ \wedge i \ l_1, l_2} \quad \text{VD_ANDI} \\
\\
\frac{\Gamma(l') = \phi \wedge \phi'}{\Gamma, \bar{\phi} \vdash l \ \phi \ \wedge e_1 \ l'} \quad \text{VD_ANDE1} \\
\\
\frac{\Gamma(l') = \phi \wedge \phi'}{\Gamma, \bar{\phi} \vdash l \ \phi' \ \wedge e_2 \ l'} \quad \text{VD_ANDE2} \\
\\
\frac{\Gamma(l') = \phi}{\Gamma, \bar{\phi} \vdash l \ \phi \vee \phi' \ \vee i_1 \ l'} \quad \text{VD_ORI1} \\
\\
\frac{\Gamma(l') = \phi'}{\Gamma, \bar{\phi} \vdash l \ \phi \vee \phi' \ \vee i_2 \ l'} \quad \text{VD_ORI2} \\
\\
\frac{\Gamma(l_1) = \phi' \quad \Gamma(l_2) = \phi' \rightarrow \phi}{\Gamma, \bar{\phi} \vdash l \ \phi \ \rightarrow e \ l_1, l_2} \quad \text{VD_IMPE} \\
\\
\frac{\Gamma(l') = \phi}{\Gamma, \bar{\phi} \vdash l \ \neg \neg \phi \ \neg i \ l'} \quad \text{VD_NEGNEGI} \\
\\
\frac{\Gamma(l') = \neg \neg \phi}{\Gamma, \bar{\phi} \vdash l \ \phi \ \neg \neg e \ l'} \quad \text{VD_NEGNEGE} \\
\\
\frac{\Gamma(l_1) = \phi \rightarrow \phi' \quad \Gamma(l_2) = \neg \phi'}{\Gamma, \bar{\phi} \vdash l \ \neg \phi \text{ MT } l_1, l_2} \quad \text{VD_MT} \\
\\
\frac{\Gamma(l_1) = \phi \quad \Gamma(l_2) = \neg \phi}{\Gamma, \bar{\phi} \vdash l \ \perp \ \neg e \ l_1, l_2} \quad \text{VD_NEGE} \\
\\
\frac{\Gamma(l_1, l_2) = (\phi, \phi')}{\Gamma, \bar{\phi} \vdash l \ \phi \rightarrow \phi' \ \rightarrow i \ l_1 - l_2} \quad \text{VD_IMPI} \\
\\
\frac{\Gamma(l_1, l_2) = (\phi, \perp)}{\Gamma, \bar{\phi} \vdash l \ \neg \phi \ \neg i \ l_1 - l_2} \quad \text{VD_NEGI}
\end{array}$$

$$\begin{array}{c}
\Gamma(l_1) = \phi \vee \phi' \\
\Gamma(l_2, l_3) = (\phi, \phi'') \\
\Gamma(l_4, l_5) = (\phi', \phi'') \\
\hline
\Gamma, \bar{\phi} \vdash l \phi'' \quad \vee e \ l_1, l_2 - l_3, l_4 - l_5 \quad \text{VD_ORE}
\end{array}$$

$$\begin{array}{c}
\Gamma(l_1, l_2) = (\neg\phi, \perp) \\
\hline
\Gamma, \bar{\phi} \vdash l \phi \quad \text{PBC } l_1 - l_2 \quad \text{VD_PBC}
\end{array}$$