VARIATIONAL QUANTUM LINEAR SOLVER

Introduction to Quantum Computing Kalamarakis Theodoros: 2018030022 February 8, 2023

GitHub Repository https://github.com/kthod/VQLS_algorithm.git

Ax = b where $A \in C^{N \times N}$ and $x, b \in C^N$

Existing classical algorithms solves the problem with polynomial scaling in N

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Exponential speedup can be achieved using the quantum algorithm **HHL** which:

- Determines a quantum state $|x\rangle$ that is proportional to $x:|x\rangle \sim x$
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• Demanding implementation

Instead a Variational Hybrid Quantum Classical Algorithm can be used

The Variational Quantum Linear Solver (VQLS)

• Can be performed in NISQ quantum computers

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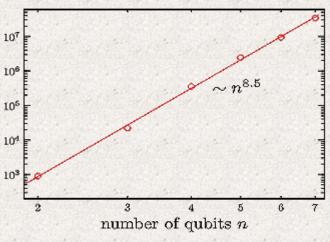
- Complexity analysis on heuristic algorithms like the VHQCA is difficult
- However complexity can be determined by performing numerical simulation
- The obtained relation for running time scaling with *n* is :

$$y \sim n^{8.5}$$

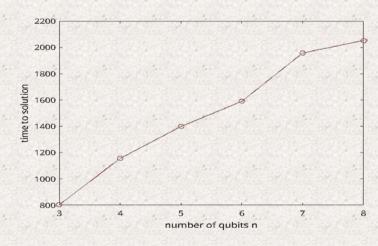
Since
$$N = 2^n \Leftrightarrow n = \log N$$

$$y \sim (log N)^{8.5}$$

Which is polylogarithmical in N



Time to solution scaling with n. Figure taken from original paper



Time to solution scaling with n. Measurements taken using our implementation of VQLS

Input

• A matrix A that, must be given as a linear combination of L unitaries A_1 , A_2 , ..., A_L such that

$$A = \sum_{l=1}^{L} c_l A_l$$

where c_l is a complex number

• An operator U that prepares a quantum state $|b\rangle$ that is proportional to the vector \mathbf{b} , such that

$$U|0\rangle = |b\rangle$$

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Goal

Determine a state $|x\rangle$ such that $A|x\rangle$ is proportional to $|b\rangle$, or equivalently:

$$|\psi\rangle := \frac{A|x\rangle}{\langle x|A^{\dagger}A|x\rangle} \approx |b\rangle$$

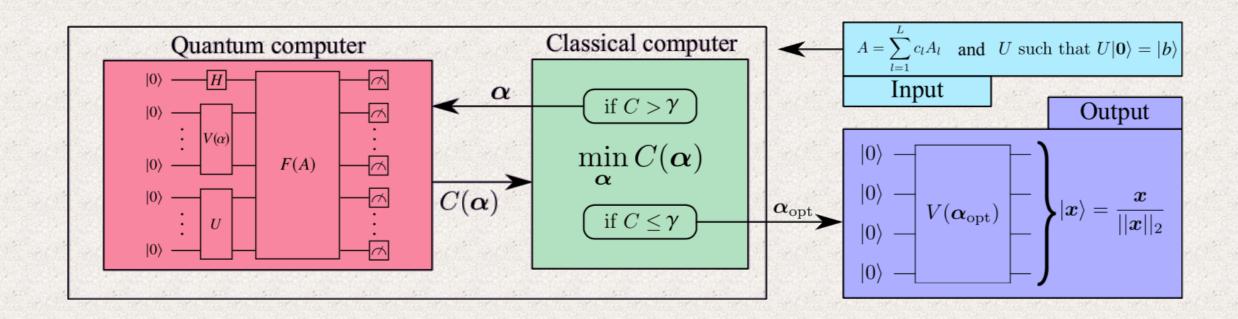
• $|x\rangle$ is approximated using a set of parameters $\alpha = (\alpha_1, \alpha_2, ...)$ such that

•
$$|x(\alpha)\rangle = V(\alpha)|\mathbf{0}\rangle$$

The operator $V(\alpha)$ is called **Ansatz** and is capable of generating any arbitrary state in Hilbert space

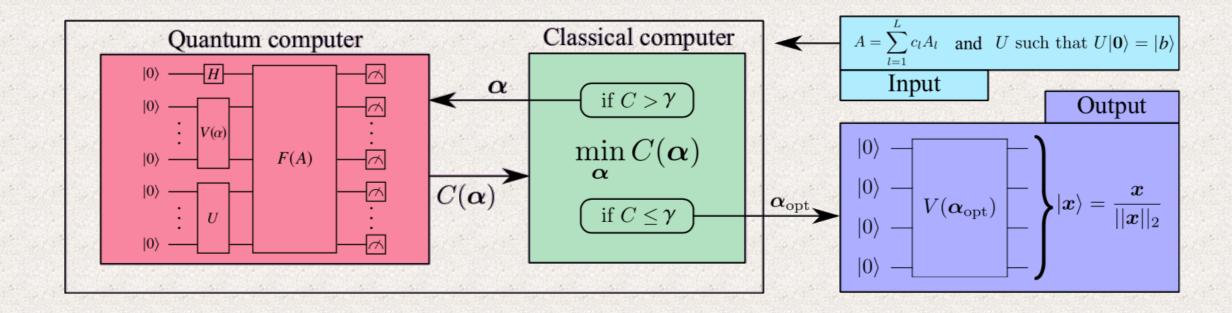
- The configuration of the parameters α is happening through a classical optimizer that aims to minimize a **cost function** $C(\alpha)$
- A quantum circuit evaluates the cost function $C(\alpha)$

Interaction between the Quantum circuit and the classical optimizer



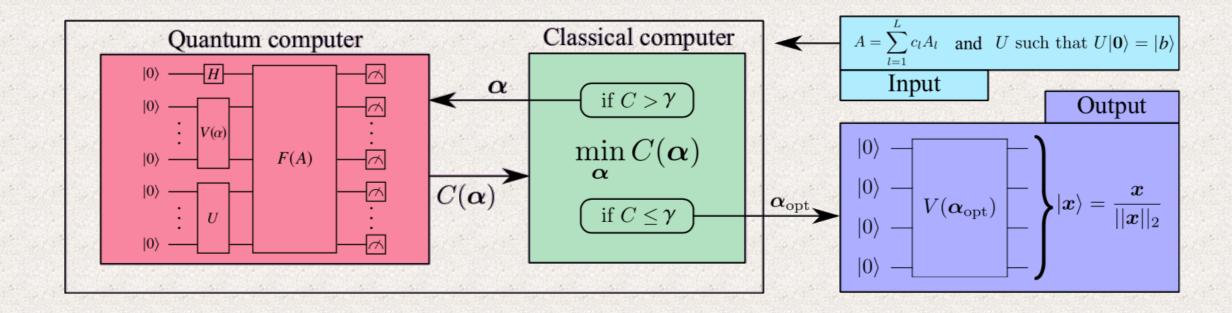
Interaction between the Quantum circuit and the classical optimizer

• The quantum circuit receives as input the parameters α and evaluates the cost function on this point $C(\alpha)$



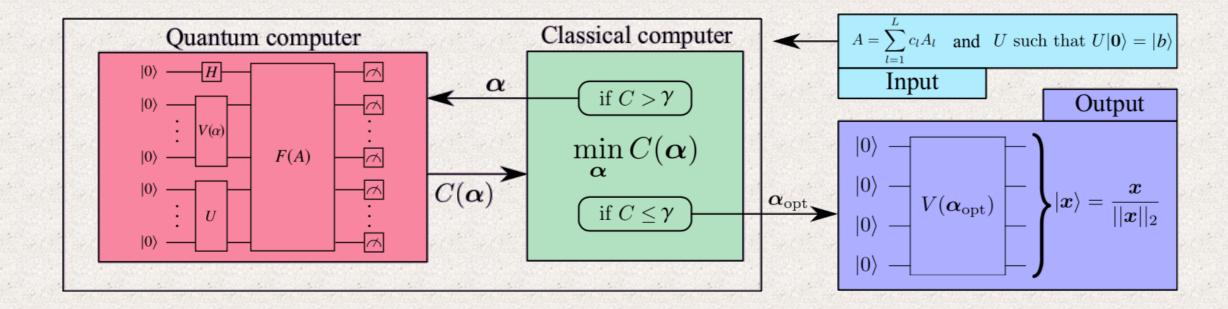
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- The quantum circuit receives as input the parameters α and evaluates the cost function on this point $C(\alpha)$
- The classical optimizer, uses the output value $C(\alpha)$, to determine a new set of parameters α' such that, $C(\alpha') < C(\alpha)$
- The quantum circuit is fed with the new set of parameters α' and the process repeats itself until the cost function reaches its global minimum: $C(\alpha_{opt}) = 0 \rightarrow |x(\alpha_{opt})\rangle = |x\rangle$



Capable of generating any arbitrary state on Hilbert space, including the entangled ones

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Consists of:

- Rotation single-qubit gates R_X , R_Y , R_Z for "exploring" the space. Each one of the rotation gates should correspond to one of the parameters $\alpha = (\alpha_1, \alpha_2, ...)$
- Entanglement two-qubit gates CNOT, CZ to provide entanglement

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Distinguishing two types of Ansatzes:

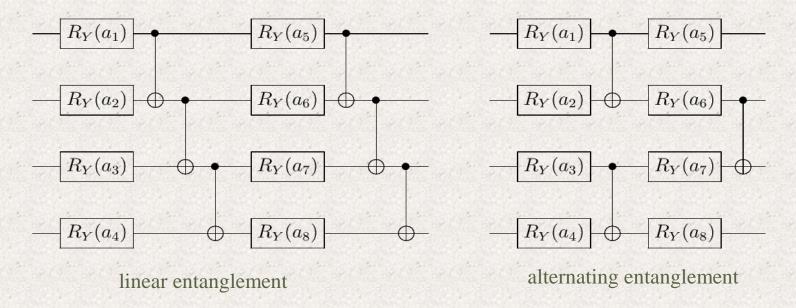
- The Fixed-Structure Anstatz
- The Variable-Structure Ansatz

We are only concerned with Fixed-Structure Anstatz

Fixed-Structure Anstaz

Placement of the gates remains the same throughout the execution of the algorithm

Two examples are the following:



Ansatz with alternating entanglement is more efficient in terms of effective parameters per two-qubit gate

Define $|\psi\rangle = |\psi(\alpha)\rangle := A|x(\alpha)\rangle$

We construct a cost function which takes its minimum value C=0 when $|\psi\rangle$ and $|b\rangle$ have the maximum overlap

The general form of a cost function in VQA is $C_G = \langle \psi | H_G | \psi \rangle$

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A suitable Hamiltonian H_G to take advantage of the overlap is

$$H_G = I - |b\rangle\langle b| = U (I - |0\rangle\langle 0|)U^{\dagger}.$$

Hence the cost function is

$$\hat{C}_G = \langle \psi | (\mathbf{I} - |b\rangle\langle b|) | \psi \rangle = \langle \psi | \psi \rangle - |\langle b | \psi \rangle|^2 \qquad \text{with } 0 \le \hat{C}_G \le \langle \psi | \psi \rangle$$

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A potential small $\langle \psi | \psi \rangle$ would be problematic, so we use instead a normalized version of the cost function

$$C_G = \frac{\hat{C}_G}{\langle \psi | \psi \rangle} = 1 - \frac{|\langle b | \psi \rangle|^2}{\langle \psi | \psi \rangle}$$

Using that $|b\rangle = U|\mathbf{0}\rangle$ and $|\psi\rangle = A|x(\boldsymbol{\alpha})\rangle = \sum_{l=1}^{L} c_l A_l V(\boldsymbol{\alpha})|\mathbf{0}\rangle$

$$C_{G} = 1 - \frac{\langle \psi | b \rangle \langle b | \psi \rangle}{\langle \psi | \psi \rangle} = 1 - \frac{\left(\sum_{l=1}^{L} c_{l}^{*} \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_{l}^{\dagger} U | \mathbf{0} \rangle\right) \left(\sum_{l=1}^{L} c_{l} \langle \mathbf{0} | U^{\dagger} A_{l} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle\right)}{\left(\sum_{l=1}^{L} c_{l}^{*} \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_{l}^{\dagger}\right) \left(\sum_{l=1}^{L} c_{l} A_{l} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle\right)} = 0$$

$$=1-\frac{\sum_{l=1}^{L}\sum_{l'=1}^{L}c_{l}^{*}c_{l}\langle\mathbf{0}|V^{\dagger}(\boldsymbol{\alpha})A_{l}^{\dagger}U|\mathbf{0}\rangle\langle\mathbf{0}|U^{\dagger}A_{l'}V(\boldsymbol{\alpha})|\mathbf{0}\rangle}{\sum_{l=1}^{L}\sum_{l'=1}^{L}c_{l}^{*}c_{l'}\langle\mathbf{0}|V^{\dagger}(\boldsymbol{\alpha})A_{l}^{\dagger}A_{l'}V(\boldsymbol{\alpha})|\mathbf{0}\rangle}$$

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The above is called Global Cost Function

• Can exhibit barren plateaus for a large number of qubits

To tackle this problem we divide the problem up into multiple single-qubit terms by replacing the term $|\mathbf{0}\rangle\langle\mathbf{0}|$ with $\frac{1}{n}\sum_{j=0}^{n-1}(|0_j\rangle\langle 0_j|\otimes I_{\bar{J}})$

The produced cost is function is called **Local Cost Function** $H_L = I - \frac{1}{n} \sum_{j=0}^{n-1} |0_j\rangle\langle 0_j| \otimes I_{\bar{J}}$

$$C_{L} = 1 - \frac{\sum_{l=1}^{L} \sum_{l'=1}^{L} c_{l}^{*} c_{l'} \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_{l}^{\dagger} U\left(\frac{1}{n} \sum_{j=0}^{n-1} \left(\left|0_{j}\right\rangle \langle 0_{j}\right| \otimes I_{\bar{j}}\right)\right) U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) |\mathbf{0}\rangle}{\sum_{l=1}^{L} \sum_{l'=1}^{L} c_{l}^{*} c_{l'} \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_{l}^{\dagger} A_{l'} V(\boldsymbol{\alpha}) |\mathbf{0}\rangle}$$

$$=1-\frac{\sum_{j=0}^{n-1}\sum_{l=1}^{L}\sum_{l'=1}^{L}c_{l}^{*}c_{l'}\langle\mathbf{0}|V^{\dagger}(\boldsymbol{\alpha})A_{l}^{\dagger}U\left(\left|0_{j}\right\rangle\langle0_{j}\right|\otimes I_{\bar{j}}\right)U^{\dagger}A_{l'}V(\boldsymbol{\alpha})|\mathbf{0}\rangle}{\sum_{l=1}^{L}\sum_{l'=1}^{L}c_{l}^{*}c_{l'}\langle\mathbf{0}|V^{\dagger}(\boldsymbol{\alpha})A_{l}^{\dagger}A_{l'}V(\boldsymbol{\alpha})|\mathbf{0}\rangle}$$

To prove that the local cost function suits our problem we use that

$$C_L \le C_G \le nC_L$$

Which implies that $C_G \to 0 \Leftrightarrow C_L \to 0$

Evaluation of the cost function using quantum circuit

• For Global Cost Function we have to evaluate the terms:

$$\beta_{ll'} = \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$$

and

$$\gamma_{ll'} = \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} U | \mathbf{0} \rangle \langle \mathbf{0} | U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$$

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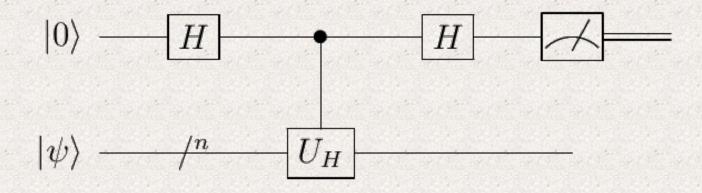
$$\delta_{ll'}^{(j)} = \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} U \left(\left| 0_j \right\rangle \left\langle 0_j \right| \otimes I_{\bar{j}} \right) U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$$

Using that $|0_j\rangle\langle 0_j|=\frac{1}{2}\;(I_j+Z_j)$, the last one can be rewritten as

$$\delta_{ll'}^{(j)} = \beta_{ll'} + \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} U \left(Z_j \otimes I_{\bar{j}} \right) U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$$

How can we evaluate the previous terms?

We employ a quantum circuit called **The Hadamard Test**

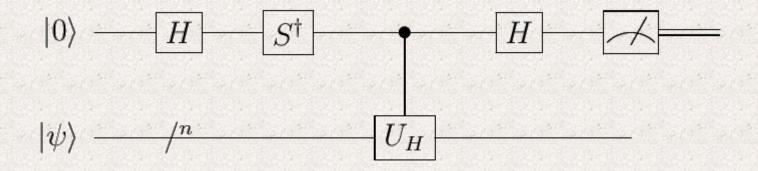


The circuit above can calculate the term $Re[\langle \psi | U_H | \psi \rangle]$

- The probability of measuring the first qubit (ancilla) to be 0 is $P(0) = \frac{1}{2}(1 + Re[\langle \psi | U_H | \psi \rangle])$
- The probability of measuring the ancilla qubit to be 1 is $P(1) = \frac{1}{2}(1 Re[\langle \psi | U_H | \psi \rangle])$

Hence
$$P(0) - P(1) = Re[\langle \psi | U_H | \psi \rangle]$$

For the Imaginary part we simple add a S^{\dagger} gate after the first Hadamard on the ancilla qubit



The circuit above can calculate the term $Im[\langle \psi | U_H | \psi \rangle]$

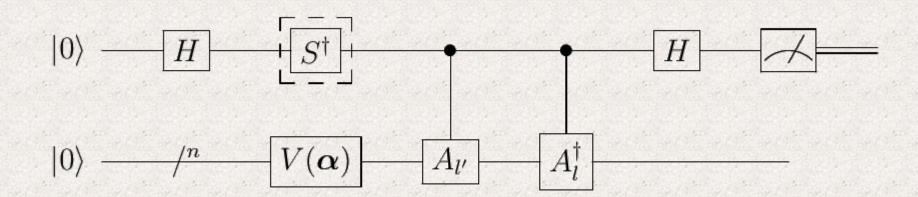
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Hence
$$P(0) - P(1) = Im[\langle \psi | U_H | \psi \rangle]$$

How can we apply the hadamard test for our purpose?

To calculate the terms $\beta_{ll'} = \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$ we set:

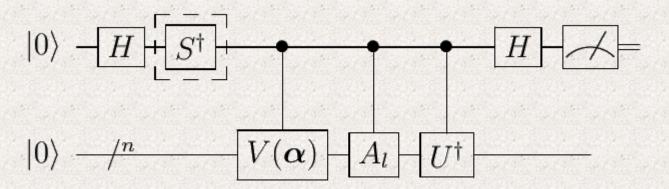
$$|\psi\rangle = V(\alpha)|\mathbf{0}\rangle$$
 and $U_H = A_l^{\dagger}A_{l'}$



To calculate the terms $\gamma_{ll'} = \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} U | \mathbf{0} \rangle \langle \mathbf{0} | U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$

$$\langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_{l}^{\dagger} U | \mathbf{0} \rangle = \left(\langle \mathbf{0} | U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle \right)^{\dagger} = \left(\langle \mathbf{0} | U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle \right)^{*}$$

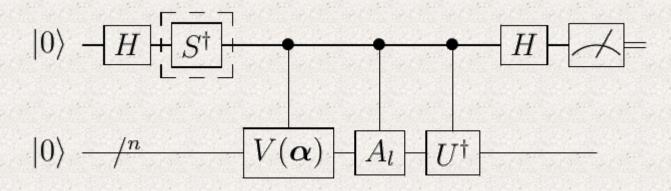
Thus by setting $|\psi\rangle = |\mathbf{0}\rangle$ and $U_H = U^{\dagger}A_lV(\alpha)$ we can calculate $\gamma_{ll'}$ by using the following circuit twice



To calculate the terms $\gamma_{ll'} = \langle \mathbf{0} | V^{\dagger}(\alpha) A_l^{\dagger} U | \mathbf{0} \rangle \langle \mathbf{0} | U^{\dagger} A_{l'} V(\alpha) | \mathbf{0} \rangle$:

$$\langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_{l}^{\dagger} U | \mathbf{0} \rangle = \left(\langle \mathbf{0} | U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle \right)^{\dagger} = \left(\langle \mathbf{0} | U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle \right)^{*}$$

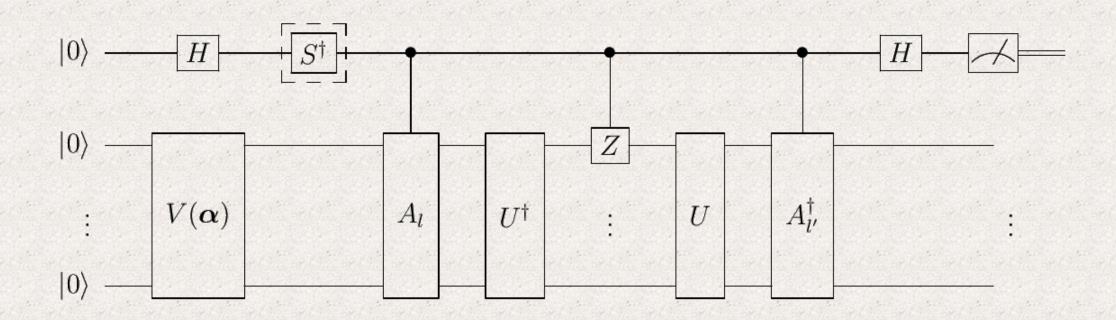
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• An alternative way to calculate these term is by using an different subroutine called **The Hadamard-Overlap Test**

To calculate the terms $\delta_{ll'}^{(j)} = \beta_{ll'} + \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} U \left(Z_j \otimes I_{\bar{J}} \right) U^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$:

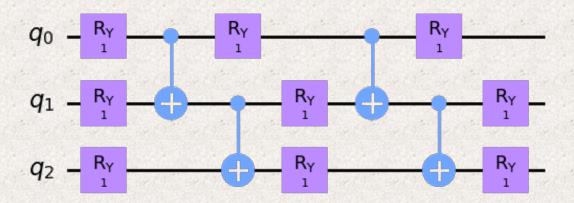
We set $|\psi\rangle = V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$ and $U_H = A_l^{\dagger} U \left(Z_j \otimes I_{\bar{J}} \right) U^{\dagger} A_{l'}$



Using Qiskit we attempt to implement the VQLS algorithm and solve an example of a 3-qubits linear equations system Ax = b with

$$\mathbf{A} = 0.4H_2 + 0.3Z_1 + 0.3X_3$$
 and $|b\rangle = H_1 \otimes H_2 \otimes H_3 |\mathbf{0}\rangle$

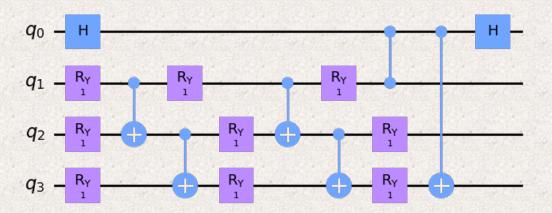
Linear entanglement ansatz was used with 9 rotation gates R_Y (therefore 9 parameters) and 4 *CNOTs* gates



Since A and $|b\rangle$ are real, only rotation around the y-axis is required

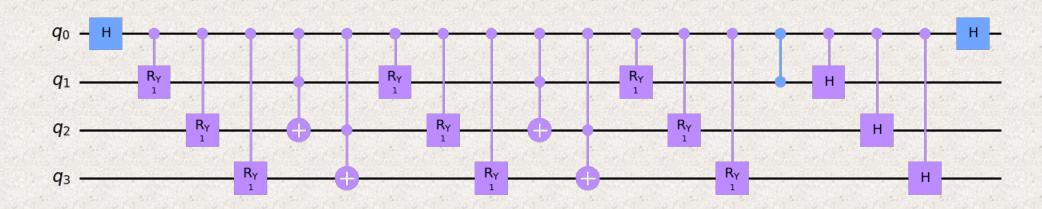
Global cost function was used:

The Hadamard Test for evaluating the terms $\beta_{ll'} = \langle \mathbf{0} | V^{\dagger}(\boldsymbol{\alpha}) A_l^{\dagger} A_{l'} V(\boldsymbol{\alpha}) | \mathbf{0} \rangle$



With
$$A_l = Z_1$$
 and $A_{l'} = X_3$

... and for the terms $\gamma_{ll'} = \langle \mathbf{0} | V^{\dagger}(\alpha) A_l^{\dagger} U | \mathbf{0} \rangle \langle \mathbf{0} | U^{\dagger} A_{l'} V(\alpha) | \mathbf{0} \rangle$



With
$$A_l = Z_1$$

Our cost function is:

$$C_G = 1 - \frac{\sum_{l=1}^{L} \sum_{l'=1}^{L} c_l^* c_l \gamma_{ll'}}{\sum_{l=1}^{L} \sum_{l'=1}^{L} c_l^* c_{l'} \beta_{ll'}}$$

We minimize it using preferably a gradient free optimizer, like COBYLA, that is less vulnerable to barren plateaus

The output of the algorithm is:

- fun is the final value of the cost function which is close to 0, as desired
- x is optimal point α_{opt} which corresponds to the solution $|x\rangle = V(\alpha_{opt})|0\rangle$

To confirm that the obtained solution is correct we calculate the inner product

 $\langle xA|b\rangle$

with the expectation of being close to 1. The output is:

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$$\langle xA|b\rangle$$

with the expectation of being close to 1. The output is:

$$\langle xA|b \rangle = (0.9998736206641057-0j)$$

... It works !!!

Referances

VQLS ALgorithm

https://arxiv.org/pdf/1909.05820.pdf

https://qiskit.org/textbook/ch-paper-implementations/vqls.html

https://pennylane.ai/qml/demos/tutorial_vqls.html

Anstaz

https://arxiv.org/pdf/2111.13730.pdf

Cost Function

https://qiskit.org/textbook/ch-paper-implementations/vqls.html