

## Quantum Error Correction

Sender  
Alice

Noisy channel

Receiver  
Bob

$$\begin{array}{r}
 \boxed{10011010} \\
 \text{---} \\
 10010\color{red}{0}010 \\
 \text{---} \\
 101100010 \\
 \text{---} \\
 10010\color{red}{1}1010
 \end{array}$$

consecutive bits  
are corrupted

Bit is sent ✓  
single bit error ✓  
multiple bit error ↗  
(Burst error)

$$\begin{array}{r}
 \boxed{1011001}0 \\
 \text{---} \\
 10100010
 \end{array}$$

parity bit  
no. of 1 = 4  
→ not acceptable

Logical Bits

$$\left\{ \begin{array}{l} 0_L = \underline{\underline{000}} \\ 1_L = \underline{\underline{111}} \end{array} \right.$$

Suppose you get : 101

then correct it to

111

## Errors in quantum communication

Decoherence

$$\begin{array}{c}
 |0\rangle + |1\rangle \\
 \hline
 \sqrt{2}
 \end{array}
 \rightarrow
 \begin{array}{c}
 |0\rangle - |1\rangle \\
 \hline
 \sqrt{2}
 \end{array}$$

Bit flips

continuous evolution of quantum states  
also introduce phase errors

$$|\psi\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}} ; \quad |\psi'\rangle = e^{\frac{i\alpha}{2}} \left( \frac{|0\rangle + |1\rangle}{\sqrt{2}} \right)$$

Say, a phase perturbation converts it to

$$|\psi'\rangle = \frac{|0\rangle + e^{\frac{i\alpha}{2}}|1\rangle}{\sqrt{2}}$$

A msmt of  $\sigma_z$  will still give you  
 $|0\rangle$  or  $|1\rangle$  with probability  $\frac{1}{2}$

But if you make a msmt of  $\sigma_x$   
it will project  $+1$  with probability

$$|\psi\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$|\psi'\rangle = \frac{|0\rangle + e^{i\alpha}|1\rangle}{\sqrt{2}} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ e^{i\alpha} \end{pmatrix} = \frac{1}{\sqrt{2}} e^{i\alpha/2} \begin{pmatrix} 1 \\ e^{i\alpha/2} \end{pmatrix}$$

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad x = \pm 1$$

$$x = +1, \quad |\underline{x=1}\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$P_{+1} = |+1\rangle \langle +1| = \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \cdot \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

$$P_{+1} |\psi'\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} \cos \alpha/2 \\ \sin \alpha/2 \end{pmatrix} e^{i\alpha/2} = \cancel{\cos \alpha/2} e^{i\alpha/2} \boxed{\frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}}$$

$$\cos^2 \frac{\alpha}{2}$$

Sender

Alice

$$|\psi\rangle = \underline{\alpha|0\rangle + \beta|1\rangle}$$

$$o_L = \underline{\underline{000}}$$

*logical*

$$\begin{aligned}
 |\psi\rangle &= (\alpha|0\rangle + \beta|1\rangle) \underline{|0\rangle} \underline{|0\rangle} \\
 &= (\alpha|00\rangle + \beta|11\rangle) |0\rangle \\
 |\Psi_L\rangle &= \alpha|000\rangle + \beta|111\rangle
 \end{aligned}$$

Alice is sending  $\Psi_L$

Say, noise in the channel flips a bit with probability  $p$  and leaving it unchanged is  $1-p$

The state received by Bob via the noisy channel is  $|\Psi_2\rangle_L$

We will have the following possibilities

Bob receives

$$\alpha|000\rangle + \beta|111\rangle$$

$$\alpha|100\rangle + \beta|011\rangle$$

$$\alpha|010\rangle + \beta|101\rangle$$

$$\dots \sim \dots |1110\rangle$$

Probability of occurrence

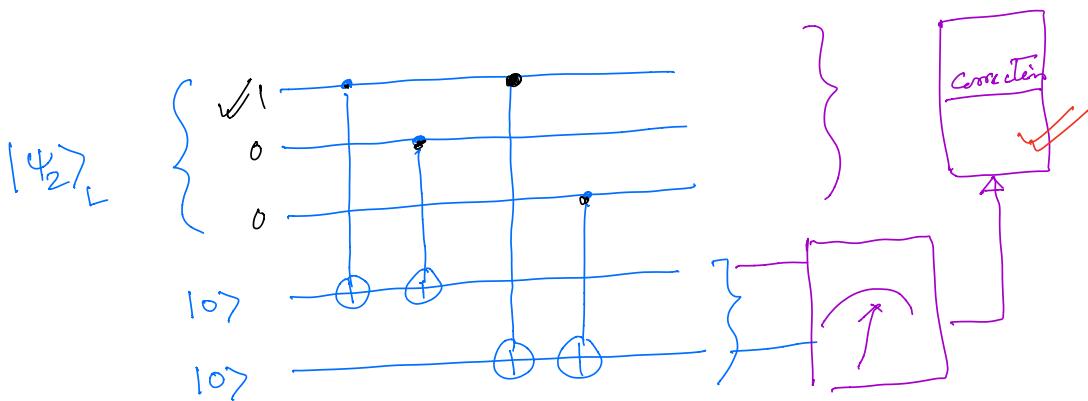
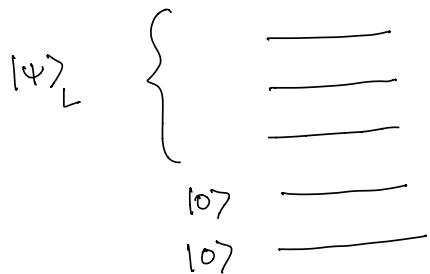
$$(1-p)^3$$

$$p(1-p)^2$$

$$p(1-p)^2$$

$$p(1-p)^2$$

$$\begin{aligned}
 & \alpha |001\rangle + \beta |111\rangle \quad / \quad \text{Total} = 3p(1-p)^2 \\
 & \alpha |110\rangle + \beta |001\rangle \quad \left. \right\} \quad p^2(1-p) \\
 & \alpha |101\rangle + \beta |010\rangle \quad \left. \right\} \quad \text{Total} = 3p^2(1-p) \\
 & \alpha |011\rangle + \beta |100\rangle \quad \left. \right\} \quad p^3 \\
 & \alpha |111\rangle + \beta |000\rangle \quad \left. \right\}
 \end{aligned}$$



$$(\alpha |000\rangle + \beta |111\rangle) |00\rangle = (\alpha |000\rangle + \beta |111\rangle) |00\rangle$$

$\Rightarrow$  ancilla bits are not affected at all.

$$|\Psi_2\rangle_L = (\underline{\alpha |100\rangle + \beta |011\rangle}) \underline{\underline{|00\rangle}}$$

=

Bob's action results in  $|4_3\rangle_L$

<u>State received</u>	<u>After Bob's gate operation</u>	<u>Probability</u>
$ 4_2\rangle_L$	$ 4_3\rangle_L$	$(1-p)^3$
$\alpha 000\rangle + \beta 111\rangle$	$(\alpha 000\rangle + \beta 111\rangle) 00\rangle$	$p(1-p)^2$
$\alpha 100\rangle + \beta 011\rangle$	$(\alpha 100\rangle + \beta 011\rangle) 11\rangle$	$p(1-p)^2$
$\cancel{\alpha 101\rangle + \beta 101\rangle}$	$(\alpha 101\rangle + \beta 101\rangle) 10\rangle$	$p(1-p)^2$
$\alpha 001\rangle + \beta 110\rangle$	$(\alpha 001\rangle + \beta 110\rangle) 01\rangle$	$p^2(1-p)$
$\alpha( 110\rangle + \beta 001\rangle)$	$(\alpha 110\rangle + \beta 001\rangle) 01\rangle$	$p^2(1-p)$
$\cancel{\alpha 101\rangle + \beta 010\rangle}$	$\cancel{ 10\rangle}$	$p^2(1-p)$
$\alpha 011\rangle + \beta 100\rangle$	$\cancel{ 11\rangle}$	$p^2(1-p)$
$\alpha 111\rangle + \beta 000\rangle$	$\cancel{ 00\rangle}$	$p^3$

msmt of ancilla may yield  
 $|00\rangle, |01\rangle, |10\rangle, |11\rangle$

Say Bob gets  $\cancel{|00\rangle}$

$\cancel{\alpha|100\rangle + \beta|111\rangle}$  with prob.  $(1-p)^3$  (without errors)

$\cancel{\alpha|111\rangle + \beta|000\rangle}$  with prob.  $p^3$  (errors)

$\checkmark$  Bob takes no action ✓

Now, say Bob gets  $\cancel{|01\rangle}$

$\cancel{\alpha|001\rangle + \beta|110\rangle}$  with prob.  $p(1-p)^2$  (with one error)

$\cancel{\alpha|110\rangle + \beta|001\rangle}$  with prob.  $p^2(1-p)$  (with two errors)

Then Bob applies a  $\sigma_x$  on the third qubit

$$\begin{aligned}\sigma_x &\xrightarrow{\text{3rd}} (\underline{\alpha|001\rangle + \beta|110\rangle}) \longrightarrow \underline{\alpha|000\rangle + \beta|111\rangle} \\ \sigma_x &\xrightarrow{\text{3rd}} (\underline{\alpha|110\rangle + \beta|001\rangle}) \longrightarrow \underline{\alpha|111\rangle + \beta|000\rangle}\end{aligned}$$

✓ One is getting corrected fully while the other one is fully getting wrong!

Say Bob sets  $\underline{|10\rangle}$  as a msmt. of the ancilla bit:

$$(\underline{\alpha|010\rangle + \beta|101\rangle}) |10\rangle$$

Now, Bob applies  $\sigma_x$  on the 2nd qubit

$$\begin{aligned}\sigma_x &\xrightarrow{\text{2nd}} (\underline{\alpha|010\rangle + \beta|101\rangle}) |10\rangle \\ &\longrightarrow (\underline{\alpha|000\rangle + \beta|111\rangle}) |10\rangle \quad \checkmark \\ &\quad \text{with prob. } \underline{p(1-p)^2}\end{aligned}$$

$$\begin{aligned}\checkmark \underline{\alpha|101\rangle + \beta|010\rangle} &\xrightarrow{\sigma_x} \underline{\alpha|111\rangle + \beta|000\rangle} \\ &\quad \underline{p^2(1-p)}\end{aligned}$$

Finally

$|11\rangle$

$$\begin{aligned}(\underline{\alpha|100\rangle + \beta|011\rangle}) &\xrightarrow{\sigma_x} \underline{\alpha|000\rangle} \\ &\quad + \underline{\beta|111\rangle} \\ &\quad \text{with prob. } \underline{p(1-p)^2}\end{aligned}$$

Bob applies  $\sigma_x$  on the 1st qubit

$$\alpha|011\rangle + \beta|100\rangle \xrightarrow{\text{Error}} \alpha|111\rangle + \beta|000\rangle$$

with prob.  $\frac{p^2(1-p)}{s}$

Q. How many of these states will continue to have errors?

A. Probability of that happening is

$$3p^2(1-p) + p^3$$

$$= 3p^2 - 2p^3 < p$$

Suppose  $p = 0.01$

$$3 \times 10^{-4}$$

error is reduced by a factor of 300!

Shor's 9 qubit error correction code

Bit flip  $\alpha|0\rangle + \beta|1\rangle \rightarrow \underline{\underline{\alpha|1\rangle + \beta|0\rangle}}$

Phase flip  $\alpha|0\rangle + \beta|1\rangle \rightarrow \alpha|0\rangle - \beta|1\rangle$

$$\underline{|0\rangle + |1\rangle} \rightarrow \underline{\underline{|0\rangle - |1\rangle}}$$

$$\begin{matrix} \sqrt{2} \\ |+\rangle \end{matrix} \longrightarrow \begin{matrix} \sqrt{2} \\ |-\rangle \end{matrix}$$

Phase-flip errors

$$\begin{matrix} |0\rangle_L \\ |1\rangle_L \end{matrix} \equiv \begin{matrix} |000\rangle \\ |111\rangle \end{matrix}$$

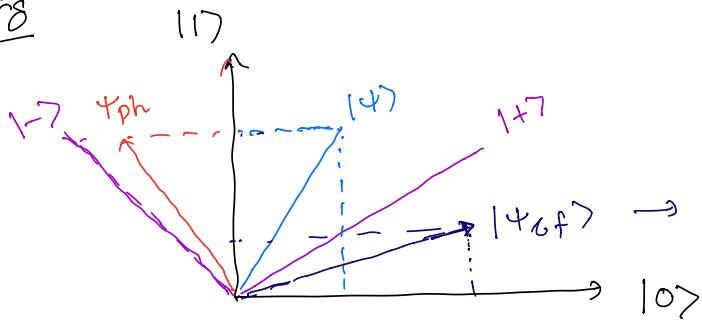
One can encode the triplet in a diagonal basis

Geometrical meaning — digression —

$$|\psi\rangle = a|0\rangle + b|1\rangle$$

$$\downarrow \quad \downarrow$$

$$|1\rangle \quad |0\rangle$$



This is the reflection of  $|\psi\rangle$  about the diagonal basis  $|+\rangle$

$$a|0\rangle + b|1\rangle$$

$$\rightarrow a|0\rangle - b|1\rangle$$

$$= -[-\underline{a}|0\rangle + \underline{b}|1\rangle]$$

$$\begin{matrix} a|000\rangle + b|111\rangle \\ |0\rangle_L \quad |1\rangle_L \end{matrix}$$

$$|\psi\rangle = a|+++> + b|--->$$

In the computational basis the state is:

$$\begin{aligned} |\psi\rangle &= \frac{a+b}{2\sqrt{2}} \left[ \underbrace{|000\rangle}_{\checkmark} + \underbrace{|011\rangle}_{\checkmark} + \underbrace{|101\rangle}_{\cancel{\checkmark}} + \underbrace{|110\rangle}_{\checkmark} \right] \\ &\quad + \frac{a-b}{2\sqrt{2}} \left[ \underbrace{|001\rangle}_{\checkmark} + \underbrace{|010\rangle}_{\checkmark} + \underbrace{|100\rangle}_{\checkmark} + \underbrace{|111\rangle}_{\cancel{\checkmark}} \right] \end{aligned}$$

Q. what happens if we measure any of the qubits in a computational basis?

Say, you measure qubit no. 1

The state collapses either to:

$$\frac{a+b}{\sqrt{2}} \left( |000\rangle + |011\rangle \right) + \frac{a-b}{\sqrt{2}} \left( |001\rangle + |010\rangle \right) \rightarrow (c)$$

or

$$\frac{a+b}{\sqrt{2}} \left( |100\rangle + |110\rangle \right) + \frac{a-b}{\sqrt{2}} \left( |101\rangle + |111\rangle \right) \rightarrow (cc)$$

Say, the state collapsed to (c)

↙  $\frac{a+b}{\sqrt{2}} \left( \cancel{|000\rangle} + \cancel{|011\rangle} \right) + \frac{a-b}{\sqrt{2}} \left( \cancel{|001\rangle} + \cancel{|010\rangle} \right)$

Introduce two ancilla and measure  
the ancilla after CNOT operations

$$|0\rangle \rightarrow \frac{|+\rangle + |-\rangle}{\sqrt{2}}$$

$$|1\rangle \rightarrow \frac{|+\rangle - |-\rangle}{\sqrt{2}}$$

In the diagonal basis the collapsed  
state is

$$\left( a |+++ \rangle + b |--- \rangle \right) + \cancel{\left( a |++- \rangle + b |+-+ \rangle \right)} \times$$

⇒ The state is either without error  
or with a single bit flip error

Shor's qubit error code

bit flip  $\times \begin{pmatrix} a \\ b \end{pmatrix} \rightarrow \begin{pmatrix} b \\ a \end{pmatrix}$

...  $|a\rangle \rightarrow |a\rangle$

phase flip  $\rightarrow (-\alpha)$

phase and bit flip error :  $\gamma(\alpha, \beta) \rightarrow (-\alpha, -\beta)$

Shor's code

We will encode

$$\checkmark |0\rangle \text{ as } |0\rangle_L = |+++ \rangle \quad \checkmark$$

$$|1\rangle \text{ as } |1\rangle_L = |--- \rangle \quad \checkmark$$

Our goal is to send

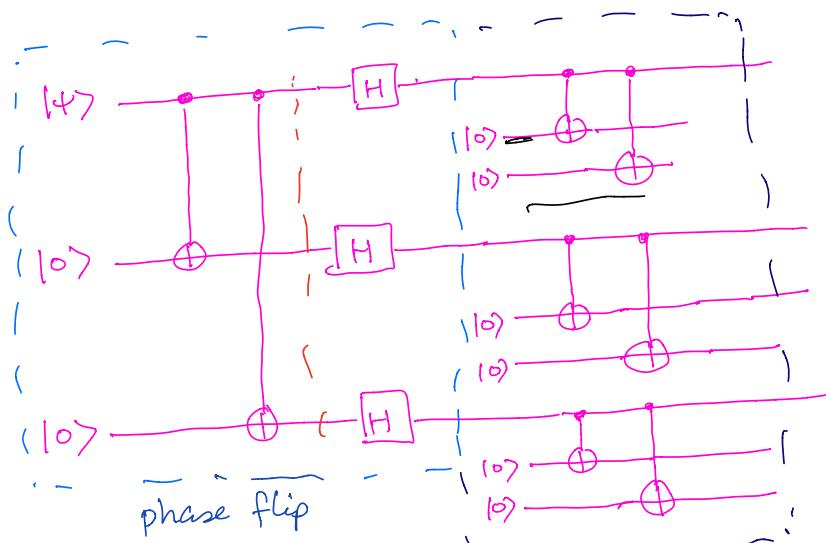
~~$a|0\rangle + b|1\rangle$~~  through a quantum channel

phase-flip encoding

$$(a|0\rangle + b|1\rangle)|0\rangle|0\rangle \xrightarrow{\text{CNOT}} a|000\rangle + b|111\rangle$$

↓  
Apply Hadamard gates on all 3 qubits

$$\begin{aligned} & \frac{a}{2\sqrt{2}} [(|0\rangle + |1\rangle)(|0\rangle + |1\rangle)(|0\rangle + |1\rangle)] \\ & + \frac{b}{2\sqrt{2}} [(|0\rangle - |1\rangle)(|0\rangle - |1\rangle)(|0\rangle - |1\rangle)] \end{aligned}$$



coding

Bit-flip  
coding

$$\frac{a}{2\sqrt{2}} \left[ (\underline{|0\rangle + |1\rangle})|00\rangle \quad (\underline{|0\rangle + |1\rangle})|00\rangle \quad (\underline{|0\rangle + |1\rangle})|00\rangle \right] \\ + \frac{b}{2\sqrt{2}} \left[ (\underline{|0\rangle - |1\rangle})|00\rangle \quad (\underline{|0\rangle - |1\rangle})|00\rangle \quad (\underline{|0\rangle - |1\rangle})|00\rangle \right]$$

CNOT operations

$$\frac{a}{2\sqrt{2}} \left[ (\underline{|000\rangle + |111\rangle}) (\underline{|000\rangle + |111\rangle}) (\underline{|000\rangle + |111\rangle}) \right] \\ + \frac{b}{2\sqrt{2}} \left[ (\underline{|000\rangle - |111\rangle}) (\underline{|000\rangle - |111\rangle}) (\underline{|000\rangle - |111\rangle}) \right]$$

Use short hand notation:

$$(+ \rightarrow \frac{|000\rangle + |111\rangle}{\sqrt{2}}) \quad \checkmark \\ (- \rightarrow \frac{|000\rangle - |111\rangle}{\sqrt{2}}) \quad \checkmark$$

=  $a |+++\rangle + b |---\rangle$

Assume that only one qubit is affected after going through the noisy channel at most

Say  $\beta$  is the probability that a single qubit is affected.

Then the probability that no qubit is affected is  $(1-\beta)^9$

$$\beta \leq 0.01 \\ (1-\beta)^9 = 1 - 9\beta + 36\beta^2 + \dots$$

The prob. that any 9 qubit is subjected to error and remaining 8 unaffected:

$$9p(1-p)^8 = 9p(1 - 8p + \dots) \\ = \cancel{9p} - \cancel{72p^2} + \dots$$

Thus, probability of zero or 1 error

$$(1 - 9p + 36p^2) + (\cancel{9p} - \cancel{72p^2}) \\ = 1 - 36p^2$$

$\Rightarrow$  Probability of having more than one error is  $\underline{\underline{36p^2}}$   
Because  $p$  is small  $36p^2 \ll p$

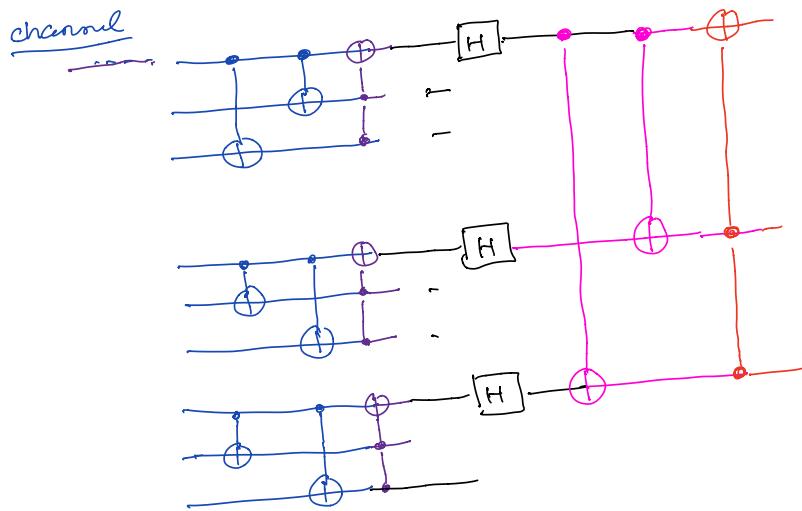
### Decoding

At the end of coding we obtain:

$$a \left[ \frac{|000\rangle + |111\rangle}{\sqrt{2}}, \frac{|000\rangle + |111\rangle}{\sqrt{2}}, \frac{|000\rangle + |111\rangle}{\sqrt{2}} \right] \\ + b \left[ \frac{|000\rangle - |111\rangle}{\sqrt{2}}, \frac{|000\rangle - |111\rangle}{\sqrt{2}}, \frac{|000\rangle - |111\rangle}{\sqrt{2}} \right]$$

Let us assume that the 1st qubit is flipped,  $\Leftrightarrow$  error

$$a \left[ \frac{|110\rangle - |011\rangle}{\sqrt{2}}, \frac{|000\rangle + |111\rangle}{\sqrt{2}}, \frac{|000\rangle + |111\rangle}{\sqrt{2}} \right] \\ + b \left[ \frac{|110\rangle + |011\rangle}{\sqrt{2}}, \frac{|000\rangle - |111\rangle}{\sqrt{2}}, \frac{|000\rangle - |111\rangle}{\sqrt{2}} \right]$$



CNOT operation

$$a \left[ \frac{|111\rangle - |000\rangle}{\sqrt{2}}, \frac{|000\rangle + |111\rangle}{\sqrt{2}}, \frac{|100\rangle + |100\rangle}{\sqrt{2}} \right] \\ + b \left[ \frac{|111\rangle + |000\rangle}{\sqrt{2}}, \frac{|000\rangle - |111\rangle}{\sqrt{2}}, \frac{|100\rangle - |100\rangle}{\sqrt{2}} \right]$$

CCNOT operation

$$a \left[ \frac{|011\rangle - |111\rangle}{\sqrt{2}}, \frac{|100\rangle + |100\rangle}{\sqrt{2}}, \frac{|100\rangle + |100\rangle}{\sqrt{2}} \right] \\ + b \left[ \frac{|011\rangle + |111\rangle}{\sqrt{2}}, \frac{|100\rangle - |100\rangle}{\sqrt{2}}, \frac{|100\rangle - |100\rangle}{\sqrt{2}} \right]$$

$$= a \left[ \frac{|10\rangle - |11\rangle}{\sqrt{2}}, |11\rangle, \frac{|10\rangle + |11\rangle}{\sqrt{2}}, |100\rangle, \frac{|10\rangle + |11\rangle}{\sqrt{2}}, |100\rangle \right] \\ + b \left[ \frac{|10\rangle + |11\rangle}{\sqrt{2}}, |11\rangle, \frac{|10\rangle - |11\rangle}{\sqrt{2}}, |100\rangle, \frac{|10\rangle - |11\rangle}{\sqrt{2}}, |100\rangle \right]$$

1st, 4th, 7th Qubit undergo Hadamard operation

$$a \left[ |1\rangle_1, |0\rangle_4, |0\rangle_7 \right] + b \left[ |0\rangle_1, |1\rangle_4, |1\rangle_7 \right]$$

(we are not worried about 2nd, 3rd, 5th, 6th, 8th, 9th)

$$\xrightarrow{\text{CNOT}} \alpha \left[ \begin{smallmatrix} |1\rangle_1 & |+\rangle_4 & |+\rangle_7 \\ \equiv & & \end{smallmatrix} \right] + \beta \left[ \begin{smallmatrix} |0\rangle_1 & |1\rangle_4 & |-\rangle_7 \\ \equiv & & \end{smallmatrix} \right]$$

$$\xrightarrow{\text{CCNOT}} \left[ \begin{smallmatrix} \alpha|0\rangle_1 + \beta|1\rangle_1 \\ \equiv & \end{smallmatrix} \right] |+\rangle_4 |+\rangle_7$$

$\xrightarrow{\quad}$

$\alpha|0\rangle + \beta|1\rangle$

↑

You got the original qubit

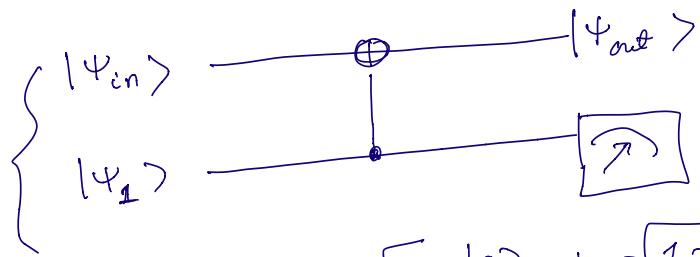
Problem

In a three qubit code, assume that the probability of a single bit flip error is 0.05. If the ancilla bit is measured to be  $|00\rangle$  find the probability that the measured state is error free.

$$\begin{array}{ll} \cancel{\alpha|000\rangle + \beta|111\rangle} & \text{with prob. } (1-p)^3 \text{ (without errors)} \\ \cancel{\alpha|111\rangle + \beta|000\rangle} & \text{with prob. } p^3 \text{ (errors)} \end{array}$$

$$\begin{aligned} \text{required probability} &= \frac{(1-p)^3}{(1-p)^3 + p^3} \\ &= 0.999854 \end{aligned}$$

Problem



$$|\Psi_1\rangle = \sqrt{p} |0\rangle + \sqrt{1-p} |1\rangle$$

what kind of process (bit flip, phase flip or both) is depicted and with what probability?

$$|\Psi_{in}\rangle = a |0\rangle + b |1\rangle$$

$$\text{Input} = \sqrt{p} (a |00\rangle + b |10\rangle) + \sqrt{1-p} (a |01\rangle + b |11\rangle)$$

$$\text{Output} = \sqrt{p} (a |00\rangle + b |10\rangle) + \sqrt{1-p} (a \cancel{|1\rangle} + b \cancel{|0\rangle})$$

Say a measurement is made on the 2nd qubit.

If the 2nd qubit measurement turns out to be |1\rangle, then the input gets flipped in the output, otherwise input does not change.

Thus, prob. of bit flip = Prob. (second qubit=|1\rangle)

$$\begin{aligned}
 &= (\sqrt{1-p} a)^2 + (\sqrt{1-p} b)^2 \\
 &= (1-p)(a^2 + b^2) \\
 &= \underline{\underline{1-p}}
 \end{aligned}$$