Accelerating Unstructured Mesh Applications using Custom Streaming Architectures

Kyrylo Tkachov

Imperial College London
kt208@doc.ic.ac.uk

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Supervisor: Prof. Paul Kelly
2nd marker: Dr. Tony Field

Unstructured meshes

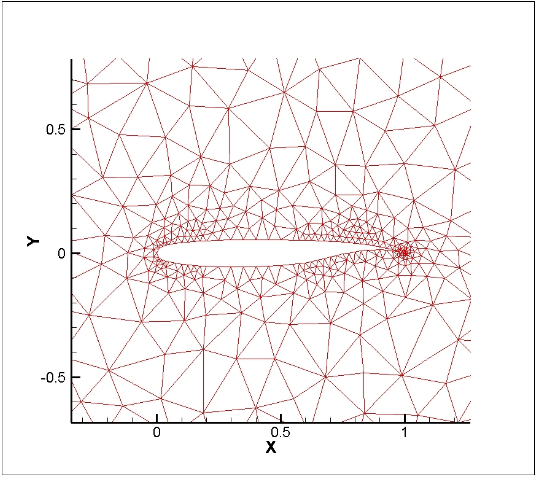


Image from Department of Environmental Engineering, University of Genoa

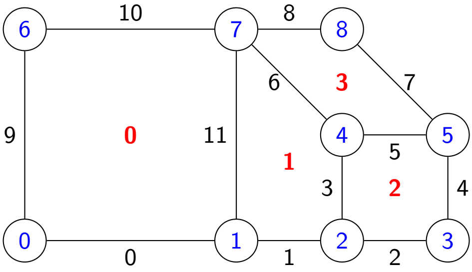
Airfoil: Indirection maps

Elements:

- Nodes
- Cells
- Edges

edge-to-node map = {0,1, 1,2, 2,3, 2,4, 3,5, 4,5, 4,7, 5,8, 7,8, 0,6, 6,7}

cell-to-node map = {0,9,10,11, 1,2,4,7, 2,3,4,5, 4,5,7,8}



Airfoil: Data sets

Data set name	Associated with	Type/Dimension
x	Nodes	$\mathbb{R} \times \mathbb{R}$
q	Cells	$\mathbb{R} \times \mathbb{R} \times \mathbb{R} \times \mathbb{R}$
q_old	Cells	$\mathbb{R} \times \mathbb{R} \times \mathbb{R} \times \mathbb{R}$
res	Cells	$\mathbb{R} \times \mathbb{R} \times \mathbb{R} \times \mathbb{R}$
adt	Cells	$\mathbb{R}$
bound	Edges	$\{0, 1\}$

Airfoil: Kernels

Kernel Name	Iterates over	Reads	Writes
save_soln	Cells	q	q_old
adt_calc	Cells	x, q	adt
res_calc	Edges	x, q, adt	res
bres_calc	(Boundary) Edges	x, q, adt, bound	res
update	Cells	q_old, adt, res	q, res

```
void res_calc(float *x1, float *x2, float *q1, float *q2,
                    float *adt1, float *adt2, float *res1, float *
                        res2) {
  float dx,dy,mu, ri, p1,vol1, p2,vol2, f;
  dx = x1[0] - x2[0];
  dy = x1[1] - x2[1];
  ri = 1.0f/q1[0];
  p1 = gm1*(q1[3]-0.5f*ri*(q1[1]*q1[1]+q1[2]*q1[2]));
  vol1 = ri*(q1[1]*dy - q1[2]*dx);
  ri = 1.0f/q2[0];
  p2 = gm1*(q2[3]-0.5f*ri*(q2[1]*q2[1]+q2[2]*q2[2]));
  vol2 = ri*(q2[1]*dy - q2[2]*dx);
  mu = 0.5f*((*adt1)+(*adt2))*eps;
  f = 0.5f*(vol1* q1[0] + vol2* q2[0] ) + mu*(q1[0]-q2[0]);
  res1[0] += f;
  res2[0] -= f;
  f = 0.5f*(vol1* q1[1] + p1*dy + vol2* q2[1] + p2*dy) + mu*(q1
      [1]-q2[1]);
  res1[1] += f;
  res2[1] -= f;
  f = 0.5f*(vol1* q1[2] - p1*dx + vol2* q2[2] - p2*dx) + mu*(q1
      [2]-q2[2]);
  res1[2] += f;
  res2[2] -= f;
  f = 0.5f*(vol1*(q1[3]+p1) + vol2*(q2[3]+p2) ) + mu*(q1[3]-q2[3])
      ;
  res1[3] += f;
  res2[3] -= f;
}
```

res_calc data requirements

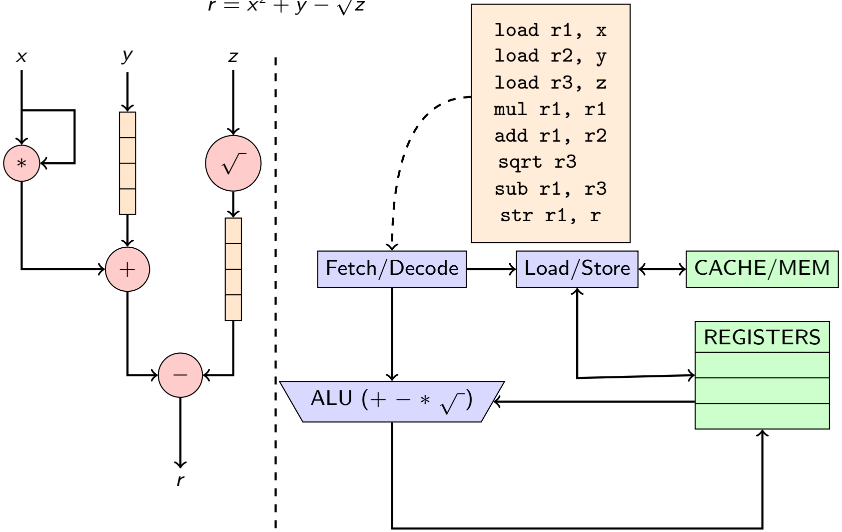
- Iterates over edges
- Processing each edge requires 2 cells, 2 nodes.
- Each edge **increments** two cells (+=).
- Most computationally intensive kernel in Airfoil.

Kernel application and double dereferencing

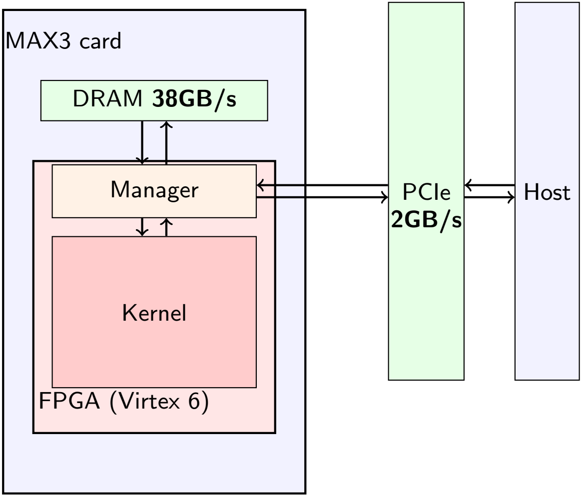
```
res_calc(
        &x[2*edge[2*i]],
        &x[2*edge[2*i+1]],
        &q[4*ecell[2*i]],
        &q[4*ecell[2*i+1]],
        &adt[ecell[2*i]],
        &adt[ecell[2*i+1]],
        &res[4*ecell[2*i]],
        &res[4*ecell[2*i+1]]
        );
```

Why custom streaming?

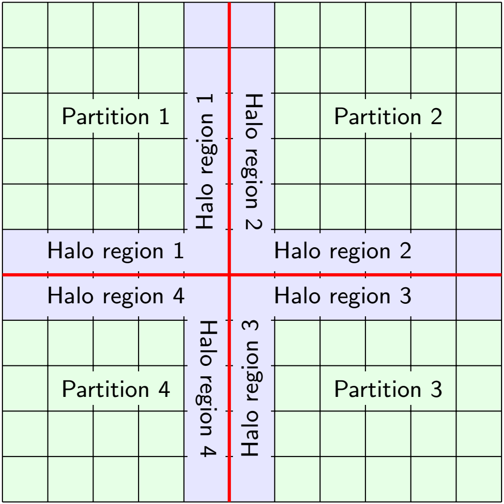
$r = x^2 + y - \sqrt{z}$



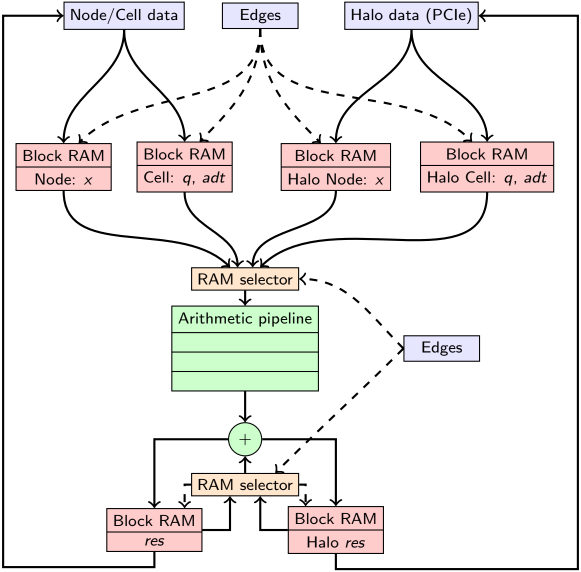
The hardware



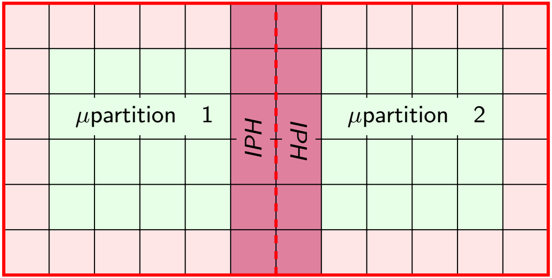
Partitioning and halos



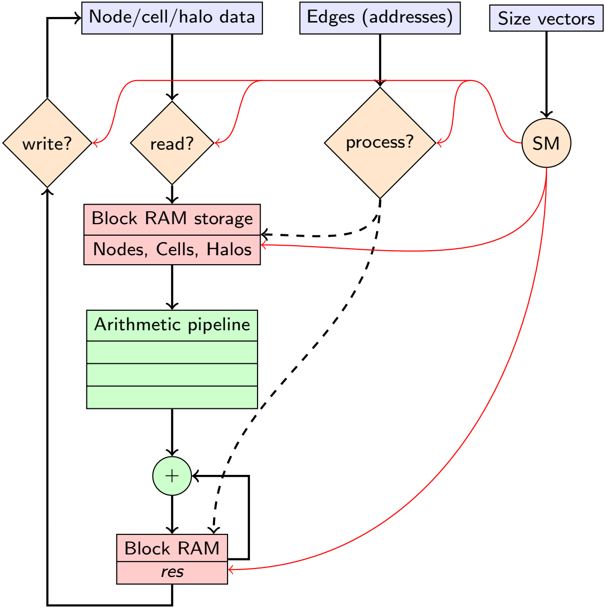
Architecture design: Accumulation and Halo Exchange



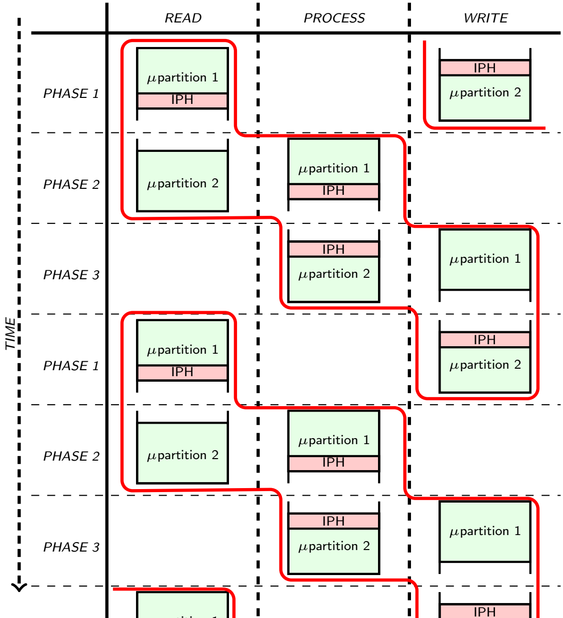
Two-level partitioning: edge processing and I/O interleaving



Architecture design: State machine



Accelerator phases and execution pattern



Performance Model

We can calculate:

- Time to stream micro-partition from DRAM:

$$\boxed{t_{DRAM} = \frac{Nonhalo\ node\ and\ cell\ data}{DRAM\ bandwidth}}$$

- Time to stream halo data for micro-partition from PCIe:

$$\boxed{t_{PCIe} = \frac{Halo\ node\ and\ cell\ data}{PCIe\ bandwidth}}$$

- Time to consume edge data during processing:

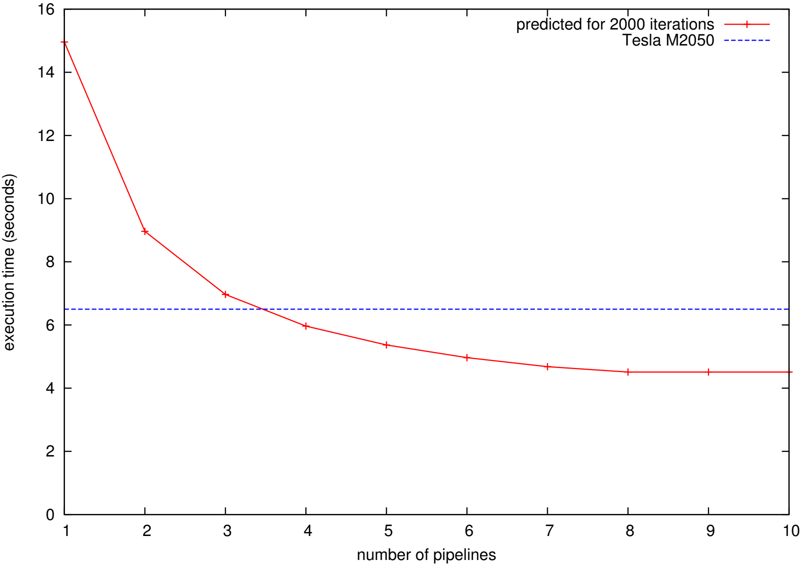
$$\boxed{t_{FPGA} = \frac{Number\ of\ edges}{frequency \times number\ of\ arithmetic\ pipelines}}$$

$\boxed{\text{Total time for each phase: } max(t_{DRAM}, t_{PCIe}, t_{FPGA})}$

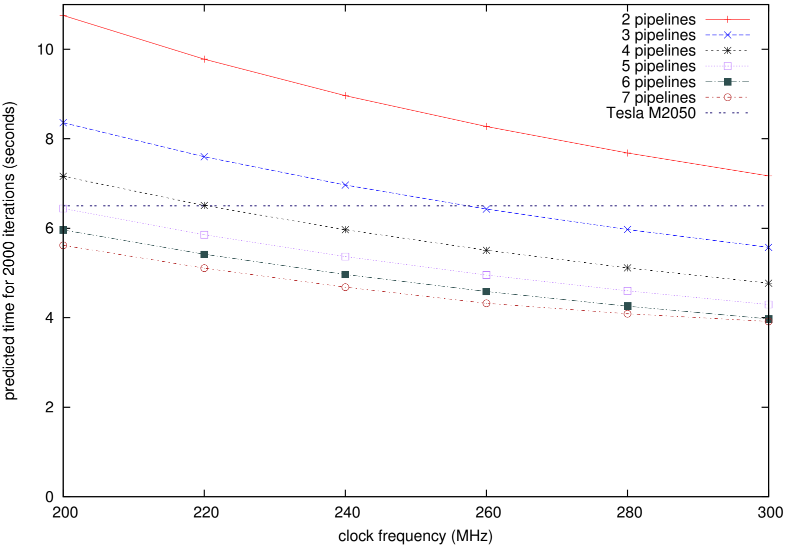
Design space exploration

- We defined a **family** of architectures.
- We can explore the design space using the performance model to find interesting ones.
- We can vary the problem and architecture parameters and predict the effect on performance.

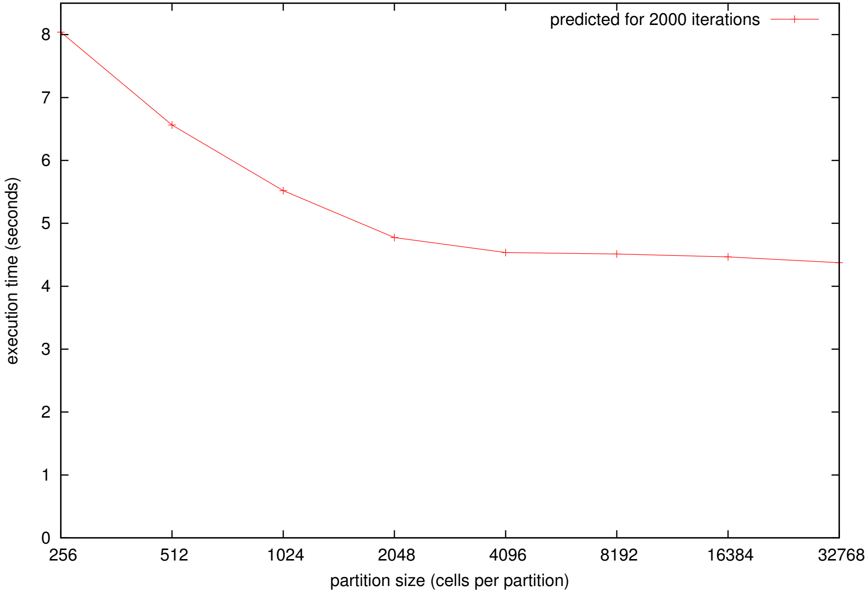
Execution time vs Number of pipelines



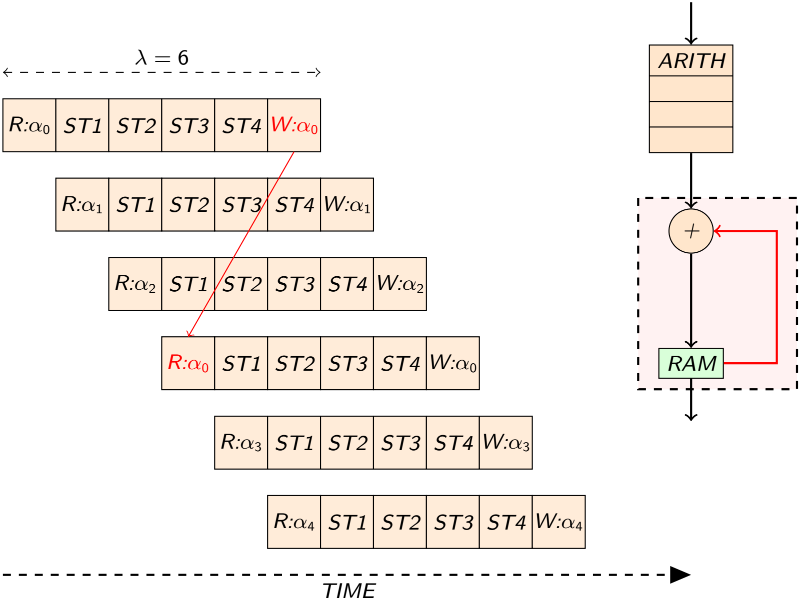
Execution time vs Number of pipelines and clock frequency



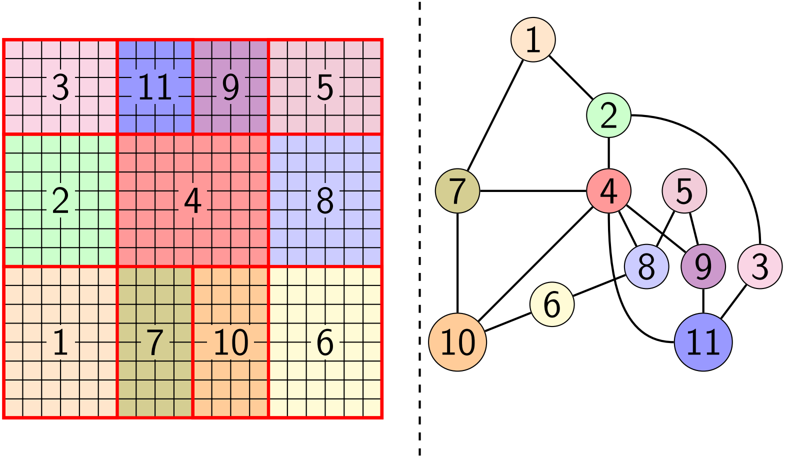
Execution time vs Partition size



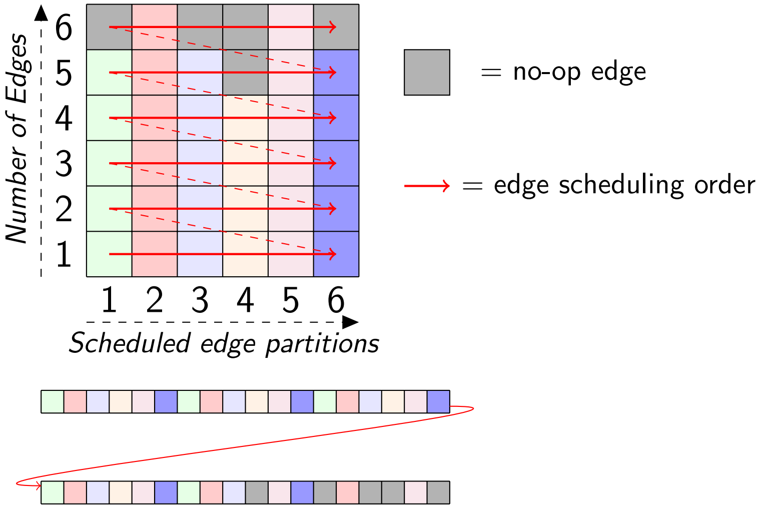
Implementation issues: Edge dependencies in the pipeline



Edge-partitions and adjacency graph scheduling



No-op edges



Number of Edges
1
2
3
4
5
6
1 2 3 4 5 6
Scheduled edge partitions
= no-op edge
= edge scheduling order

Complexity of edge scheduling

- Graph scheduling problem can be expressed as Hamiltonian path problem with extra adjacency constraint.
- **NP-hard!**.
- Best we can do is search through the schedule space.
- $O(n!)$ (n number of nodes in the adjacency graph).

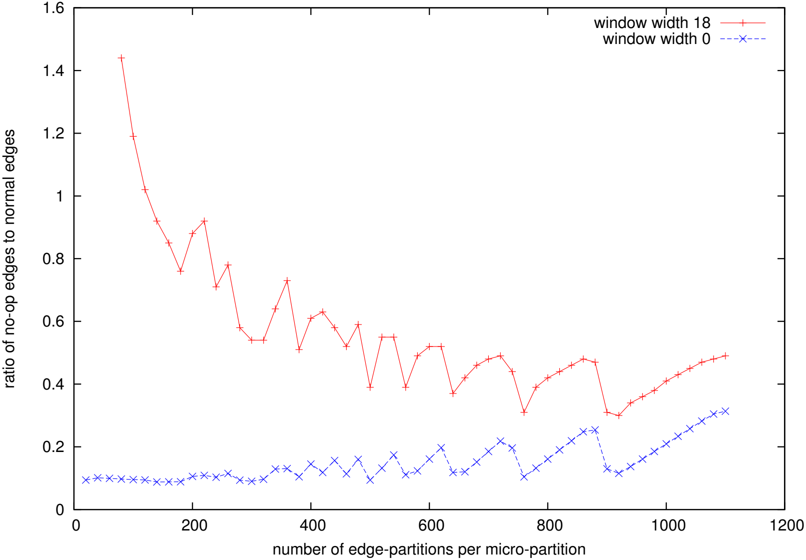
Graph colouring and no-op edge-partitions

- Can group together partitions with same colour.
- To produce schedule with window-width λ add λ no-op edge-partitions after each colour group. Add $\lambda \times c$ no-op partitions (c-number of colours used to colour graph).
- Optimal colouring still **NP**-complete, but we can efficiently find sub-optimal but adequate colouring.
- Greedy graph colouring algorithm. Assign lowest colour not assigned to neighbours of node. Worst-case time complexity is $O(n^3)$.

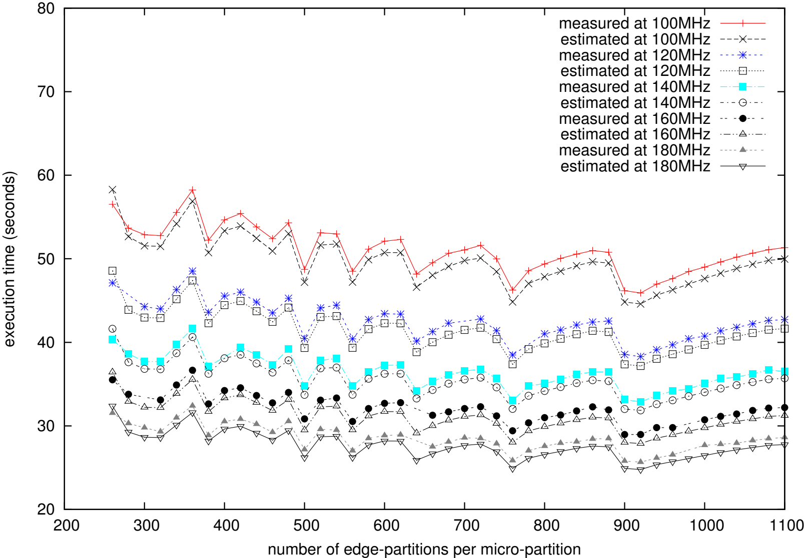
A note on correctness

- Sample implementation gives wrong arithmetic results.
- Through simulations and debugging tracked down to result commiting part of accelerator design.
- *NaN*s from no-op edges commited to result RAMs.
- Kernel consumes and produces correct amount of data in the correct order. Processes correct number of edges.
- *Can still trust the performance results.*

Evaluation: No-op edges and METIS



Evaluation: Performance model validation, various frequencies



Conclusions

- Performance model is validated!
- We can accurately predict the performance of a design space of architectures.
- Simple memory hierarchy provides high predictability. No cache-misses, no non-deterministic thread scheduling by OS.
- Unstructured memory accesses transformed into highly predictable and easily modelled streaming model!
- We showed that an interesting speedup can be achieved.
- Performance rivaling 448-core GPU implementation with only 4-5 pipelines running at a fraction of the clock frequency!

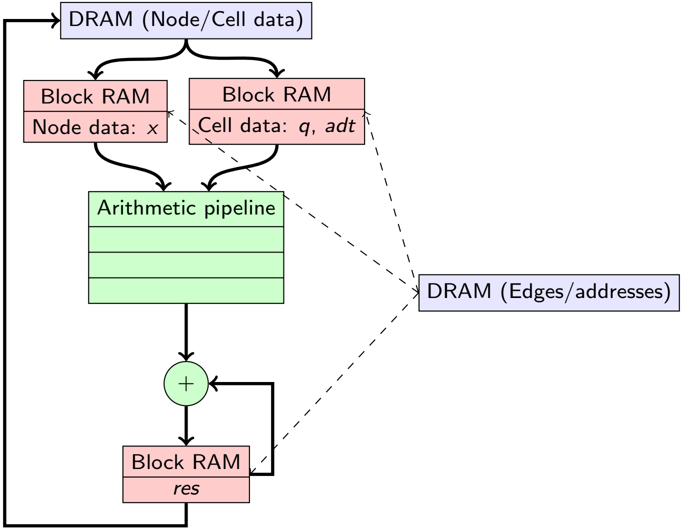
Further work

- Accelerate other kernels: different element iteration requires different data layout!
- Compilation system: plug in architecture parameters and generate host code and accelerator.
- Data formatting: reduce padding, increase bandwidth utilisation.
- Build multi-pipe designs: model predicts they offer the most performance benefits.
- And more!

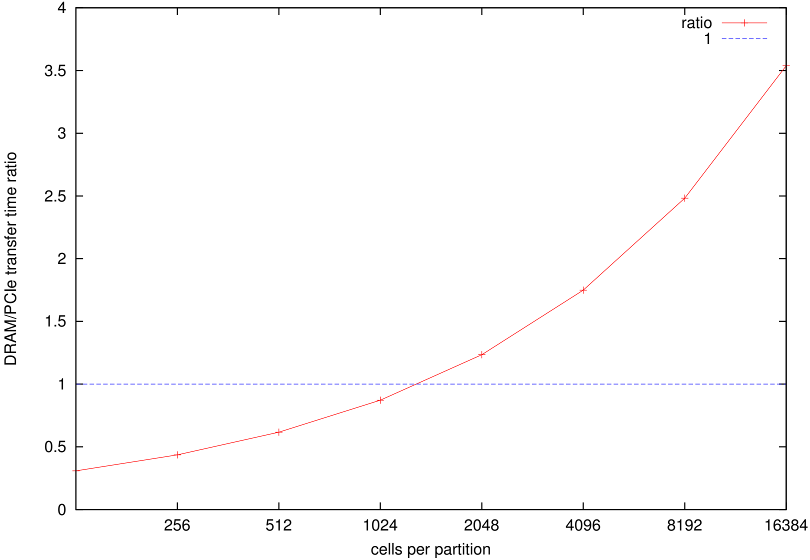
Thank you!

Questions

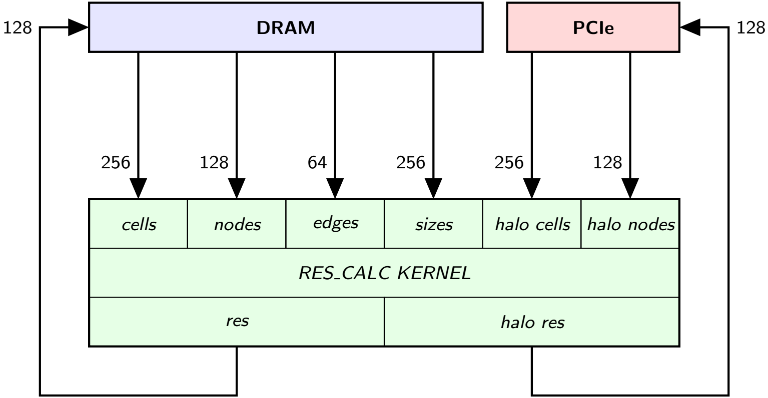
Architecture design: 1st iteration



DRAM/PCIe transfer ratio vs Partition size



Implementation issues: FPGA accelerator, manager configuration



Two-port limitation on RAMs