

### 5.136 domain\_constraint

	DESCRIPTION	LINKS	GRAPH	AUTOMATON
Origin	[339]			
Constraint	domain_constraint(VAR, VALUES)			
Synonym	domain.			
Arguments	VAR : dvar VALUES : collection(var01-dvar, value-int)			
Restrictions	required(VALUES, [var01, value]) VALUES.var01 ≥ 0 VALUES.var01 ≤ 1 distinct(VALUES, value)			
Purpose	Make the link between a domain variable VAR and those 0-1 variables that are associated with each potential value of VAR: The 0-1 variable associated with the value that is taken by variable VAR is equal to 1, while the remaining 0-1 variables are all equal to 0.			
Example	$\left( 5, \left\langle \begin{array}{ll} \text{var01} - 0 & \text{value} - 9, \\ \text{var01} - 1 & \text{value} - 5, \\ \text{var01} - 0 & \text{value} - 2, \\ \text{var01} - 0 & \text{value} - 7 \end{array} \right\rangle \right)$			
	The domain_constraint holds since VAR = 5 is set to the value corresponding to the 0-1 variable set to 1, while the other 0-1 variables are all set to 0.			
Typical	VALUES  > 1			
Symmetry	Items of VALUES are <a href="#">permutable</a> .			
Usage	This constraint is used in order to make the link between a formulation using finite domain constraints and a formulation exploiting 0-1 variables.			
Reformulation	The domain_constraint(VAR, ⟨var01 - B <sub>1</sub> value - v <sub>1</sub> , var01 - B <sub>2</sub> value - v <sub>2</sub> , ..... var01 - B <sub> VALUES </sub> value - v <sub> VALUES </sub> ⟩) constraint can be expressed in term of the following reified constraint (VAR = v <sub>1</sub> ∧ B <sub>1</sub> = 1) ∨ (VAR = v <sub>2</sub> ∧ B <sub>2</sub> = 1) ∨ ... ∨ (VAR = v <sub> VALUES </sub> ∧ B <sub> VALUES </sub> = 1).			
Systems	<a href="#">domainChanneling</a> in <a href="#">Choco</a> , <a href="#">channel</a> in <a href="#">Gecode</a> , <a href="#">in</a> in <a href="#">SICStus</a> , <a href="#">in_set</a> in <a href="#">SICStus</a> .			

**See also**

**common keyword:** `link_set_to_booleans` (*channelling constraint*).

**related:** `roots`.

**Keywords**

**characteristic of a constraint:** `automaton`, `automaton without counters`,  
`reified automaton constraint`, `derived collection`.

**constraint network structure:** `centered cyclic(1)` `constraint network(1)`.

**constraint type:** `decomposition`.

**filtering:** `linear programming`, `arc-consistency`.

**modelling:** `channelling constraint`, `domain channel`, `Boolean channel`.

Derived Collection	$\text{col} \left( \begin{array}{l} \text{VALUE} - \text{collection}(\text{var01} - \text{int}, \text{value} - \text{dvar}), \\ [\text{item}(\text{var01} - 1, \text{value} - \text{VAR})] \end{array} \right)$
Arc input(s)	VALUE VALUES
Arc generator	<i>PRODUCT</i> $\mapsto$ collection(value, values)
Arc arity	2
Arc constraint(s)	value.value = values.value $\Leftrightarrow$ values.var01 = 1
Graph property(ies)	<i>NARC</i> =  VALUES

Graph model

The domain\_constraint constraint is modelled with the following bipartite graph:

- The first class of vertices corresponds to a single vertex containing the domain variable.
- The second class of vertices contains one vertex for each item of the collection VALUES.

*PRODUCT* is used in order to generate the arcs of the graph. In our context it takes a collection with a single item  $\langle \text{var01} - 1 \text{ value} - \text{VAR} \rangle$  and the collection VALUES.

The arc constraint between the variable VAR and one potential value  $v$  expresses the following:

- If the 0-1 variable associated with  $v$  is equal to 1, VAR is equal to  $v$ .
- Otherwise, if the 0-1 variable associated with  $v$  is equal to 0, VAR is not equal to  $v$ .

Since all arc constraints should hold the final graph contains exactly |VALUES| arcs.  
Parts (A) and (B) of Figure 5.308 respectively show the initial and final graph associated with the **Example** slot. Since we use the *NARC* graph property, the arcs of the final graph are stressed in bold.

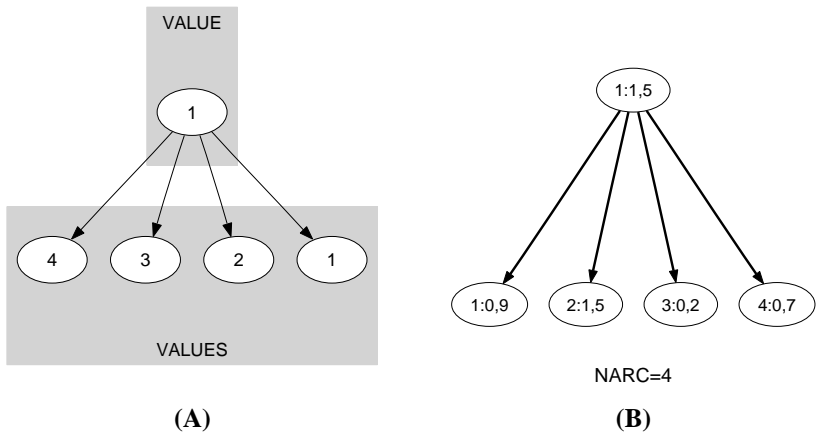


Figure 5.308: Initial and final graph of the domain\_constraint constraint

**Signature**

Since the number of arcs of the initial graph is equal to  $VALUES$  the maximum number of arcs of the final graph is also equal to  $VALUES$ . Therefore we can rewrite the graph property  $NARC = |VALUES|$  to  $NARC \geq |VALUES|$ . This leads to simplify  $NARC$  to  $NARC$ .

**Automaton**

Figure 5.309 depicts the automaton associated with the `domain_constraint` constraint. Let  $\text{VAR01}_i$  and  $\text{VALUE}_i$  respectively be the `var01` and the `value` attributes of the  $i^{\text{th}}$  item of the `VALUES` collection. To each triple  $(\text{VAR}, \text{VAR01}_i, \text{VALUE}_i)$  corresponds a 0-1 signature variable  $S_i$  as well as the following signature constraint:  $((\text{VAR} = \text{VALUE}_i) \Leftrightarrow \text{VAR01}_i) \Leftrightarrow S_i$ .

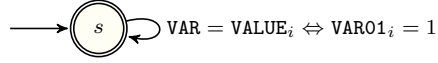
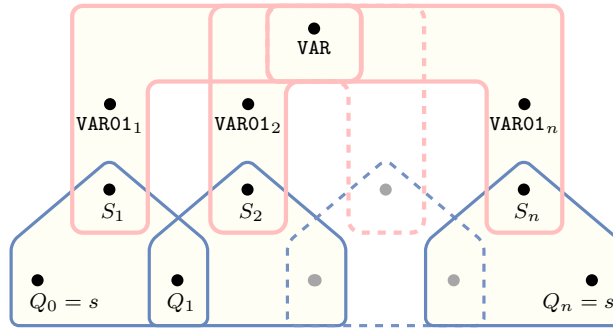
Figure 5.309: Automaton of the `domain_constraint` constraint

Figure 5.310: Hypergraph of the reformulation corresponding to the automaton of the `domain_constraint` constraint: since all states variables  $Q_0, Q_1, \dots, Q_n$  are fixed to the unique state  $s$  of the automaton, the transitions constraints involve only a single variable and the constraint network is Berge-acyclic

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