

Systems	<code>eq</code> in Choco , <code>neq</code> in Choco , <code>geq</code> in Choco , <code>gt</code> in Choco , <code>leq</code> in Choco , <code>lt</code> in Choco , <code>rel</code> in Gecode , <code>#<</code> in SICStus , <code>#=<</code> in SICStus , <code>#></code> in SICStus , <code>#>=</code> in SICStus , <code>#=</code> in SICStus , <code>#\=</code> in SICStus .
Used in	<code>arith_sliding</code> .
See also	<p>common keyword: <code>among</code>, <code>count</code> (<i>value constraint</i>).</p> <p>generalisation: <code>arith_or</code>(variable RELOP VALUE <i>replaced by</i> variable RELOP VALUE \vee variable RELOP VALUE).</p> <p>system of constraints: <code>arith_sliding</code>.</p>
Keywords	<p>characteristic of a constraint: <code>automaton</code>, <code>automaton without counters</code>, <code>reified automaton constraint</code>.</p> <p>constraint network structure: <code>Berge-acyclic constraint network</code>.</p> <p>constraint type: <code>decomposition</code>, <code>value constraint</code>.</p> <p>filtering: <code>arc-consistency</code>.</p> <p>modelling: <code>domain definition</code>.</p>
Cond. implications	<p><code>arith(VARIABLES, RELOP, VALUE)</code> with <code>RELOP</code> \in [<code><</code>] and <code>minval(VARIABLES.var)</code> ≥ 0 implies <code>range_ctr(VARIABLES, CTR, R)</code> when <code>CTR</code> \in [<code><</code>].</p>

Arc input(s)	VARIABLES
Arc generator	$\text{SELF} \mapsto \text{collection}(\text{variables})$
Arc arity	1
Arc constraint(s)	$\text{variables.var RELOP VALUE}$
Graph property(ies)	$\text{NARC} = \text{VARIABLES} $

Graph model Parts (A) and (B) of Figure 5.85 respectively show the initial and final graph associated with the **Example** slot. Since we use the **NARC** graph property, the loops of the final graph are stressed in bold.



Figure 5.85: Initial and final graph of the arith constraint

Automaton

Figure 5.86 depicts the automaton associated with the `arith` constraint. To each variable VAR_i of the collection `VARIABLES` corresponds a 0-1 signature variable S_i . The following signature constraint links VAR_i and S_i : $\text{VAR}_i \text{ RELOP VALUE} \Leftrightarrow S_i$. The automaton enforces for each variable VAR_i the condition $\text{VAR}_i \text{ RELOP VALUE}$.



Figure 5.86: Automaton of the `arith` constraint

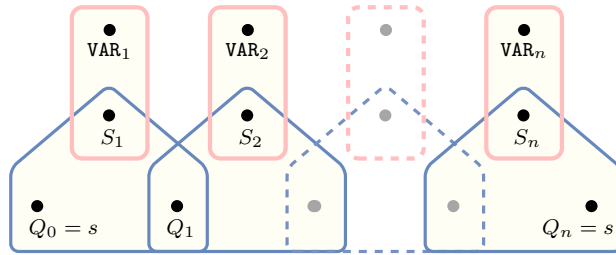


Figure 5.87: Hypergraph of the reformulation corresponding to the automaton of the `arith` constraint