5.163 global_cardinality

Constraint

Purpose

All solutions

	DESCRIPTION	LINKS	GRAPH	AUTOMATON
Origin	CHARME [298]			

Synonyms count, distribute, distribution, gcc, card_var_gcc, egcc, extended_global_cardinality.

global_cardinality(VARIABLES, VALUES)

Restrictions

required(VARIABLES, var)

required(VALUES, [val, noccurrence])

distinct(VALUES, val)

VALUES.noccurrence ≥ 0

VALUES.noccurrence ≤ |VARIABLES|

Each value VALUES[i].val (with $i \in [1, |VALUES|]$) should be taken by exactly VALUES[i].noccurrence variables of the VARIABLES collection.

Example $\left(\begin{array}{c} \langle 3,3,8,6 \rangle\,, \\ \sqrt{\begin{array}{c} \mathtt{val} - 3 & \mathtt{noccurrence} - 2, \\ \mathtt{val} - 5 & \mathtt{noccurrence} - 0, \\ \mathtt{val} - 6 & \mathtt{noccurrence} - 1 \end{array} \right)$

The global_cardinality constraint holds since values 3, 5 and 6 respectively occur 2, 0 and 1 times within the collection $\langle 3,3,8,6 \rangle$ and since no restriction was specified for value 8.

Figure 5.369 gives all solutions to the following non ground instance of the global_cardinality constraint: $V_1 \in [3,4], \ V_2 \in [2,3], \ V_3 \in [1,2], \ V_4 \in [2,4], \ V_5 \in [2,3], \ V_6 \in [1,2], \ O_1 \in [1,1], \ O_2 \in [2,3], \ O_3 \in [0,1], \ O_4 \in [2,3], \ global_cardinality(\langle V_1, V_2, V_3, V_4, V_5, V_6 \rangle, \langle 1 \ O_1, \ 2 \ O_2, \ 3 \ O_3, \ 4 \ O_4 \rangle).$

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\begin{tabular}{lll} \textbf{Typical} & | VARIABLES| > 1 \\ & \textbf{range}(VARIABLES.var) > 1 \\ & | VALUES| > 1 \\ & | VARIABLES| \geq | VALUES| \\ & \textbf{minval}(VARIABLES.var) = 0 \lor \textbf{in\_attr}(VARIABLES, var, VALUES, val) \\ \end{tabular}
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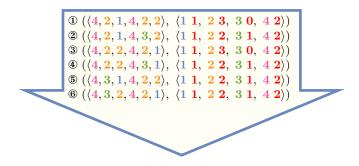


Figure 5.369: All solutions corresponding to the non ground example of the global_cardinality constraint of the **All solutions** slot

Symmetries

- Items of VARIABLES are permutable.
- Items of VALUES are permutable.
- An occurrence of a value of VARIABLES.var that does not belong to VALUES.val can be replaced by any other value that also does not belong to VALUES.val.
- All occurrences of two distinct values in VARIABLES.var or VALUES.val can be swapped; all occurrences of a value in VARIABLES.var or VALUES.val can be renamed to any unused value.

Arg. properties

- Functional dependency: VALUES.noccurrence determined by VARIABLES and VALUES.val.
- Contractible wrt. VALUES.

Usage

We show how to use the global_cardinality constraint in order to model the magic series problem [415, page 155] with a single global_cardinality constraint. A non-empty finite series $S = (s_0, s_1, \ldots, s_n)$ is magic if and only if there are s_i occurrences of i in S for each integer i ranging from 0 to n. This leads to the following model:

$$\begin{array}{c} \text{global_cardinality} \left(\begin{array}{c} \left\langle \begin{array}{c} \mathsf{var} - s_0, \mathsf{var} - s_1, \dots, \mathsf{var} - s_n \end{array} \right\rangle, \\ \mathsf{val} - 0 \quad \mathsf{noccurrence} - s_0, \\ \left\langle \begin{array}{c} \mathsf{val} - 1 \quad \mathsf{noccurrence} - s_1, \\ \vdots \\ \mathsf{val} - n \quad \mathsf{noccurrence} - s_n \end{array} \right) \end{array} \right)$$

Remark

This is a generalised form of the original global_cardinality constraint: in the original global_cardinality constraint [342], one specifies for each value its minimum and maximum number of occurrences (i.e., see global_cardinality_low_up). Here we give for each value v a domain variable that indicates how many time value v is actually used. By setting the minimum and maximum values of this variable to the appropriate constants we can express the same thing as in the original global_cardinality constraint. However, as shown in the *magic series* problem, we can also use this variable in other constraints. By reduction from 3-SAT, Claude-Guy Quimper shows in [331] that it is NP-hard to achieve arc-consistency for the count variables.

A last difference with the original global_cardinality constraint comes from the fact that there is no constraint on the values that are not explicitly mentioned in the VALUES collection. In the original global_cardinality these values could not be assigned to the variables of the VARIABLES collection. However allowing values that are not mentioned in VALUES to be assigned to variables of VARIABLES can potentially avoid mentioning a huge number of unconstrained values in the VALUES collection, and as a side effect, prevent possibly generating a dense graph (i.e., see DFS-bottleneck) for the corresponding underlying flow model).

Within [83] the global_cardinality constraint is called distribution. Within [350] the global_cardinality constraint is called card_var_gcc. Within [70] the global_cardinality constraint is called egcc or rgcc. This later case corresponds to the fact that some variables are duplicated within the VARIABLES collection.

The global_cardinality constraint can be seen as a system (i.e., a conjunction) of among constraints.

When all count variables (i.e., the variables VALUES[i].noccurrence with $i \in [1, |VALUES|])$ do not occur in any other constraints of the problem, it may be operationally more efficient to replace the global_cardinality constraint by a global_cardinality_low_up constraint where each count variable VALUES[i].noccurrence is replaced by the corresponding interval [VALUES[i].noccurrence, VALUES[i].noccurrence]. This stands for two reasons:



- First, by using a global_cardinality_low_up constraint rather than a global_cardinality constraint, we avoid the filtering algorithm related to the count variables.
- Second, unlike the global_cardinality constraint where we need to fix all its variables to get entailment, the global_cardinality_low_up constraint can be entailed before all its variables get fixed. As a result, this potentially avoid unnecessary calls to its filtering algorithm.

When all values that can be assigned to the variables of the VARIABLES collection occur in the val attribute of the VALUES collection, two implicit *necessary* conditions⁷ inferred by double counting with the global_cardinality constraint are depicted by the following expressions:



$$|\mathtt{VARIABLES}| = \sum_{i=1}^{|\mathtt{VALUES}|} \mathtt{VALUES}[i].\mathtt{noccurrence}$$

$$\sum_{i=1}^{|\mathtt{VARIABLES}|} \mathtt{VARIABLES}[i].\mathtt{var} = \sum_{i=1}^{|\mathtt{VALUES}|} \mathtt{VALUES}[i].\mathtt{val} \cdot \mathtt{VALUES}[i].\mathtt{noccurrence}$$

Within [317, pages 50–51] the previous condition where terms involving identical variables are grouped together (i.e., rule 5 of MALICE [316]) is mentioned as a crucial deduction rule for the autoref problem.

⁶ Of course one could also, while generating a flow model, detect all unconstrained values in order to generate a single vertex in the flow model for the set of unconstrained values.

⁷Note that such necessary conditions can be derived by assigning an integer weight to each value [385], e.g. 1 for the first condition, the value itself for the second condition.

W.-J. van Hoeve *et al.* present two soft versions of the global_cardinality constraint in [424].

In MiniZinc (http://www.minizinc.org/) there is also a distribute constraint where the val attribute is not necessarily initially fixed and where a same value may occur more than once. Their is also a global_cardinality_closed constraint where all variables must be assigned a value from the val attribute.

Algorithm

A flow algorithm that handles the original global_cardinality constraint is described in [342]. The two approaches that were used to design bound-consistency algorithms for alldifferent were generalised for the global_cardinality constraint. The algorithm in [334] identifies Hall intervals and the one in [233] exploits convexity to achieve a fast implementation of the flow-based arc-consistency algorithm. The later algorithm can also compute bound-consistency for the count variables [234, 231]. An improved algorithm for achieving arc-consistency is described in [333].

Systems

globalCardinality in Choco, count in Gecode, gcc in JaCoP, global_cardinality in MiniZinc, global_cardinality in SICStus.

See also

common keyword: count, max_nvalue, min_nvalue (value constraint, counting constraint),
nvalue (counting constraint),

open_global_cardinality_low_up (assignment, counting constraint).

cost variant: global_cardinality_with_costs(cost associated with each
variable,value pair).

implied by: global_cardinality_with_costs(forget about cost),
same_and_global_cardinality(conjoin same and global_cardinality).

part of system of constraints: among.

related: roots, sliding_card_skip0 (counting constraint of a set of values on maximal sequences).

shift of concept: global_cardinality_no_loop (assignment of a variable to its position is ignored), ordered_global_cardinality (restrictions are done on nested sets of values, all starting from first value), symmetric_cardinality, symmetric_gcc.

soft variant: open_global_cardinality (a set variable defines the set of variables that are actually considered).

system of constraints: colored_matrix (one global_cardinality constraint for each row and each column of a matrix of variables).

uses in its reformulation: tree_range, tree_resource.

Keywords

application area: assignment.

characteristic of a constraint: core, automaton, automaton with array of counters. **complexity:** 3-SAT.

Cond. implications

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constraint arguments: pure functional dependency.
constraint type: value constraint, counting constraint, system of constraints.
filtering: Hall interval, bound-consistency, flow, duplicated variables, DFS-bottleneck.
modelling: functional dependency.
modelling exercises: magic series.
puzzles: magic series, autoref.
• global_cardinality(VARIABLES, VALUES)
   with minval(VARIABLES.var) = 0
 implies and(VAR, VARIABLES)
   when VAR = 0.
• global_cardinality(VARIABLES, VALUES)
   with maxval(VARIABLES.var) = 1
 implies or(VAR, VARIABLES)
  when VAR = 1.
• global_cardinality(VARIABLES, VALUES)
   with minval(VARIABLES.var) > 0
 implies min_size_full_zero_stretch(MINSIZE, VARIABLES)
  when MINSIZE = |VARIABLES|.
• global_cardinality(VARIABLES, VALUES)
   with maxval(VARIABLES.var) < 0
 implies min_size_full_zero_stretch(MINSIZE, VARIABLES)
   when MINSIZE = |VARIABLES|.
• global_cardinality(VARIABLES, VALUES)
   with range(VALUES.val) =nval(VALUES.val)
   and minval(VALUES.val) \left\left\) minval(VARIABLES.var)
  and maxval(VALUES.val) ≥maxval(VARIABLES.var)
 implies among_diff_O(NVAR, VARIABLES).
• global_cardinality(VARIABLES, VALUES)
   with range(VALUES.val) =nval(VALUES.val)
   and minval(VALUES.val) ≤minval(VARIABLES.var)
   and maxval(VALUES.val) ≥maxval(VARIABLES.var)
 implies atmost_nvalue(NVAL, VARIABLES).
• global_cardinality(VARIABLES, VALUES)
   with range(VALUES.noccurrence) = 1
   and range(VALUES.val) = nval(VALUES.val)
   and minval(VALUES.val) =minval(VARIABLES.var)
   and maxval(VALUES.val) =maxval(VARIABLES.var)
 implies balance(BALANCE, VARIABLES)
   when BALANCE = 0.
• global_cardinality(VARIABLES, VALUES)
   with range(VALUES.val) = nval(VALUES.val)
  and minval(VALUES.val) ≤minval(VARIABLES.var)
   and maxval(VALUES.val) ≥maxval(VARIABLES.var)
 implies max_n(MAX, RANK, VARIABLES).
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```
• global_cardinality(VARIABLES, VALUES)
  with range(VALUES.val) =nval(VALUES.val)
  and minval(VALUES.val) \leftlef{\text{sinval}}(VARIABLES.var)
  and maxval(VALUES.val) >maxval(VARIABLES.var)
 implies max_nvalue(MAX, VARIABLES).
• global_cardinality(VARIABLES, VALUES)
  with range(VALUES.val) = nval(VALUES.val)
  and minval(VALUES.val) ≤minval(VARIABLES.var)
  and maxval(VALUES.val) ≥maxval(VARIABLES.var)
 implies min_n(MIN, RANK, VARIABLES).
• global_cardinality(VARIABLES, VALUES)
  with range(VALUES.val) =nval(VALUES.val)
  and minval(VALUES.val) <minval(VARIABLES.var)
  and maxval(VALUES.val) ≥maxval(VARIABLES.var)
 implies min_nvalue(MIN, VARIABLES).
• global_cardinality(VARIABLES, VALUES)
  with range(VALUES.val) =nval(VALUES.val)
  and minval(VALUES.val) ≤minval(VARIABLES.var)
  and maxval(VALUES.val) ≥maxval(VARIABLES.var)
 implies range_ctr(VARIABLES, CTR, R).
```

For all items of VALUES:

Arc input(s) VARIABLES

Arc generator SELF→collection(variables)

Arc arity 1

Arc constraint(s) variables.var = VALUES.val

Graph property(ies) NVERTEX = VALUES.noccurrence

Graph model

Since we want to express one unary constraint for each value we use the "For all items of VALUES" iterator. Part (A) of Figure 5.370 shows the initial graphs associated with each value 3, 5 and 6 of the VALUES collection of the **Example** slot. Part (B) of Figure 5.370 shows the two corresponding final graphs respectively associated with values 3 and 6 that are both assigned to the variables of the VARIABLES collection (since value 5 is not assigned to any variable of the VARIABLES collection the final graph associated with value 5 is empty). Since we use the **NVERTEX** graph property, the vertices of the final graphs are stressed in bold.

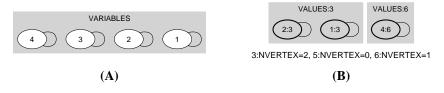


Figure 5.370: Initial and final graph of the global_cardinality constraint

Automaton

Figure 5.371 depicts the automaton associated with the global_cardinality constraint. To each item of the collection VARIABLES corresponds a signature variable S_i that is equal to 0. To each item of the collection VALUES corresponds a signature variable $S_{i+|\text{VARIABLES}|}$ that is equal to 1.

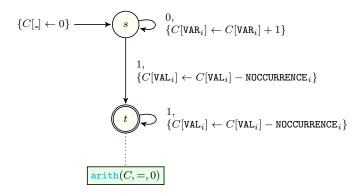


Figure 5.371: Automaton of the global_cardinality constraint

Quiz

EXERCISE 1 (checking whether a ground instance holds or not)

- A. Does the constraint global_cardinality($\langle 2, 4, 2, 2, 1 \rangle$, $\langle 0 \ 0, 1 \ 1, 2 \ 3, 3 \ 0, 4 \ 1 \rangle$) hold?
- **B.** Does the constraint global_cardinality($\langle 0,0,1,1\rangle,\langle 0\ 2,1\ 2,2\ 1,3\ 0,4\ 0\rangle$) hold?
- C. Does the constraint global_cardinality($\langle 2,3,4,5\rangle,\langle 0\ 0,1\ 0,2\ 1,3\ 1,4\ 1\rangle$) hold?

EXERCISE 2 (finding all solutions)

Give all the solutions to the constraint:

$$\left\{ \begin{array}{ll} V_1 \in [1,2], & V_2 \in [1,2], & V_3 \in [1,2], \\ V_4 \in [2,3], & V_5 \in [3,3], \\ O_1 \in [1,2], & O_2 \in [2,3], & O_3 \in [0,1], \\ \\ \text{global_cardinality} \left(\begin{array}{ll} \langle V_1, V_2, V_3, V_4, V_5 \rangle, \\ \langle \text{val} - 1 & \text{occurrence} - O_1, \\ \text{val} - 2 & \text{occurrence} - O_2, \\ \text{val} - 3 & \text{occurrence} - O_3 \end{array} \right) \right). \\ \end{array} \right.$$

$$^a \text{Hint: focus on the variables of the first argument (since the counting variables)}$$

^aHint: go back to the definition of global_cardinality.

^aHint: focus on the variables of the first argument (since the counting variables of the second argument are functionally determined by the first argument), and enumerate solutions in lexicographic order.

EXERCISE 3 (identifying infeasible values)^a

Identify all variable-value pairs (V_i, val) (respectively (O_i, val)) (with $i \in [1, 5]$), such that the following constraint has no solution when variable V_i (respectively O_i) is assigned value val:

```
\left\{ \begin{array}{ll} V_1 \in [2,3], & V_2 \in [1,5], & V_3 \in [3,4], \\ V_4 \in [1,3], & V_5 \in [1,4], \\ O_1 \in [1,4], & O_2 \in [0,1], & O_3 \in [0,1], \\ O_4 \in [1,5], & O_5 \in [1,4], \\ \end{array} \right. \\ \left\{ \begin{array}{ll} \left\langle V_1, V_2, V_3, V_4, V_5 \right\rangle, \\ \left\langle \text{val} - 1 & \text{occurrence} - O_1, \\ \text{val} - 2 & \text{occurrence} - O_2, \\ \text{val} - 3 & \text{occurrence} - O_3, \\ \text{val} - 4 & \text{occurrence} - O_4, \\ \text{val} - 5 & \text{occurrence} - O_5 \end{array} \right. \right\}.
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^aHint: first restrict the occurrence variables O_1,O_2,\ldots,O_5 , second restrict the decision variables V_1,V_2,\ldots,V_5 , third check that all remaining values occur in at least one solution.

EXERCISE 4 (modelling a nurse assignment problem)^a

Given a 24 hour period, you must schedule a pool of six nurses Bea, Lea, Leo, Lio, Lili and Tom to at least two and at most three morning shifts, to at least two and at most three afternoon shifts, to at least one night shift, while the other nurses are off-duty. In addition, due to past work, we have the following extra requirements:

- Since on the previous 24 hour period Bea, Lea and Leo were working in the afternoon shift they cannot be assigned to the night shift
- Leo, Lio and Lili have to work since they already took all their days off.
- Bea and Tom have to work together since Bea supervises Tom.

Provide a model of this problem that uses the global_cardinality constraint.

- A. Provide a solution that satisfies all the constraints, i.e., for each nurse give his/her assignment (morning, afternoon, night, off-duty).
- **B.** Identify the decision variables and the values of the problem, i.e., how do we model the fact that nurse $x \in \{\text{Bea}, \text{Lea}, \text{Leo}, \text{Lio}, \text{Lili}, \text{Tom}\}$ is assigned shift $y \in \{\text{morning}, \text{afternoon}, \text{night}, \text{off-duty}\}$?
- C. Using a bipartite graph, draw the relations between the variables and the values identified in the previous question and display the solution you came up with in the first question.
- **D.** Provide a model of the problem that uses a single global_cardinality constraint.
 - Explain how the minimum/maximum capacity constraints (i.e., at least/at most) are modelled.
 - Explain how each extra requirement is modelled in your solution

^aHint: focus on what is a variable and what is a value in your model, and how to model the capacity constraints with global_cardinality.

SOLUTION TO EXERCISE 1

- **A.** Yes, since within (2, 4, 2, 2, 1), values 0, 1, 2, 3 and 4 are respectively used zero, one, three, zero, and one times.
- **B.** No, since within (0,0,1,1), value 2 is not used one time.
- C. Yes, since within (2,3,4,5), value (0,1,2,3) and (4,3) are respectively used zero, zero, one, one, and one times. The presence of a 5 in the solution does not matter since value 5 is not mentioned in the values of the second argument of the global_cardinality constraint.

SOLUTION TO EXERCISE 2

the six solutions

```
 \begin{array}{c} \langle v_1, v_2, v_3, v_4, v_5 \rangle & \langle 1 \ O_1, 2 \ O_2, 3 \ O_3 \rangle \\ \hline 0 & (\langle 1, 1, 2, 2, 3 \rangle, \ \langle 1 \ 2, 2 \ 2, 3 \ 1 \rangle) \\ \hline 0 & (\langle 1, 2, 1, 2, 3 \rangle, \ \langle 1 \ 2, 2 \ 2, 3 \ 1 \rangle) \\ \hline 0 & (\langle 1, 2, 2, 2, 3 \rangle, \ \langle 1 \ 1, 2 \ 3, 3 \ 1 \rangle) \\ \hline 0 & (\langle 2, 1, 1, 2, 3 \rangle, \ \langle 1 \ 2, 2 \ 2, 3 \ 1 \rangle) \\ \hline 0 & (\langle 2, 1, 2, 2, 3 \rangle, \ \langle 1 \ 1, 2 \ 3, 3 \ 1 \rangle) \\ \hline 0 & (\langle 2, 2, 1, 2, 3 \rangle, \ \langle 1 \ 1, 2 \ 3, 3 \ 1 \rangle) \\ \hline \end{array}
```

SOLUTION TO EXERCISE 3

As suggested by the hint we go through the following steps:

- **A.** [RESTRICTING THE OCCURRENCE VARIABLES O_1, O_2, \ldots, O_5]
 - (a) [PRUNING WRT THE MAXIMUM NUMBER OF OCCURRENCES OF EACH VALUE]
 Since values 1, 2, 3, 4 and 5 can respectively be assigned to at most 3, 4, 5, 3 and 1 decision variables (e.g., value 1 can only be assigned to V₂, V₄ and V₅) we have O₁ ≤ min(3, 4), O₂ ≤ min(4, 1), O₃ ≤ min(5, 1), O₄ ≤ min(3, 5), and O₅ ≤ min(1, 4).
 - (b) [PRUNING WRT $\sum_{i=1}^5 O_i = 5$ AND THE DOMAIN OF V_1] Since we have five decision variables the sum of the occurrence variables is equal to five (i.e., $O_1 + O_2 + O_3 + O_4 + O_5 = 5$). Since values 2 or 3 have to be assigned to the decision variable V_1 we have $O_2 + O_3 \geq 1$. It follows that $O_1 + O_4 + O_5 \leq 4$. Since $O_1 \in [1,3]$, $O_4 \in [1,3]$ and $O_5 = 1$ we get $O_1 + O_4 \leq 3$ and consequently $O_1 \leq 2$ and $O_4 \leq 2$.
- **B.** [RESTRICTING THE DECISION VARIABLES V_1, V_2, \ldots, V_5] At the end of step **A** we obtain $O_1 \in [1, 2]$, $O_2 \in [0, 1]$, $O_3 \in [0, 1]$, $O_4 \in [1, 2]$, and $O_5 \in [1, 1]$. Since $O_5 = 1$ and since V_2 is the only decision variable that can be assigned value 5 we have $V_2 = 5$. Consequently $V_1 \in [2, 3]$, $V_2 \in [5, 5]$, $V_3 \in [3, 4]$, $V_4 \in [1, 3]$, and $V_5 \in [1, 4]$.
- C. [CHECKING FOR A SUPPORT]

To show that no value can be removed from the domain of the decision and occurrence variables we show that every value that is still in the domain of a variable is part of a solution.

- (a) A solution with $O_1 = 2$ is $V_1 = 2$, $V_2 = 5$, $V_3 = 4$, $V_4 = 1$, $V_5 = 1$ and $O_1 = 2$, $O_2 = 1$, $O_3 = 0$, $O_4 = 1$, $O_5 = 1$.
- (b) A solution with $O_4=2$ is $V_1=2$, $V_2=5$, $V_3=4$, $V_4=1$, $V_5=4$ and $O_1=1$, $O_2=1$, $O_3=0$, $O_4=2$, $O_5=1$.
- (c) We now assume that $O_1 = O_2 = O_3 = O_4 = O_5 = 1$, i.e., all decision variables must be distinct. Without loos of generality we ignore V_2 , which is fixed to 5. We provide a set of solutions where V_1 , V_3 , V_4 and V_5 can respectively be assigned to all the values of their domains:

i.
$$V_1 = 2$$
, $V_3 = 3$, $V_4 = 1$, $V_5 = 4$,

ii.
$$V_1 = 2$$
, $V_3 = 4$, $V_4 = 3$, $V_5 = 1$,

iii.
$$V_1 = 2$$
, $V_3 = 4$, $V_4 = 1$, $V_5 = 3$,

iv.
$$V_1 = 3$$
, $V_3 = 4$, $V_4 = 1$, $V_5 = 2$,

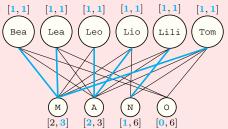
$$V_1 = 3$$
, $V_3 = 4$, $V_4 = 2$, $V_5 = 1$.

SOLUTION TO EXERCISE 4

A. A feasible solution is the following assignment

```
Bea: morning Lea: morning Leo: afternoon Lio: afternoon Lili: night Tom: morning since:
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- the number of morning shifts is between 2 and 3,
- the number of afternoon shifts is between 2 and 3,
- the number of night shifts is at least 1,
- Bea, Lea and Leo are not assigned to a night shift,
- Leo, Lio and Lili work,
- Bea and Tom are both assigned the same shift.
- **B.** To each nurse corresponds a variable whose initial domain is set to the types of shifts that nurse can actually perform (i.e., each shift type is encoded by a unique integer value).
- C. The next figure provides a graphical representation of the assignment problem. To each nurse and to each shift type corresponds a vertex. There is an edge between a given nurse and a given shift type if and only if that nurse can perform that shift type. The solution given to question A is displayed with thick blue lines. The interval on top or below each vertex indicates the minimum and maximum number of edges that can reach the corresponding vertex in any solution; values in blue correspond to the number of edges of the displayed solution.



D. We get the following model

$$\left\{ \begin{array}{ll} \mathsf{M} = 1, & \mathsf{A} = 2, & \mathsf{N} = 3, & \mathsf{O} = 4, \\ \mathsf{Bea} \in [\mathsf{M}, \mathsf{O}], & \mathsf{Lea} \in [\mathsf{M}, \mathsf{O}], & \mathsf{Leo} \in [\mathsf{M}, \mathsf{O}], \\ \mathsf{Lio} \in [\mathsf{M}, \mathsf{O}], & \mathsf{Lili} \in [\mathsf{M}, \mathsf{O}], & \mathsf{Tom} \in [\mathsf{M}, \mathsf{O}], \\ O_{\mathsf{M}} \in [2, 3], & O_{\mathsf{A}} \in [2, 3], & O_{\mathsf{N}} \in [1, 6], & O_{\mathsf{O}} \in [0, 6], \\ \mathsf{Bea} \neq \mathsf{N}, & \mathsf{Lea} \neq \mathsf{N}, & \mathsf{Leo} \neq \mathsf{N}, \\ \mathsf{Leo} \neq \mathsf{O}, & \mathsf{Lio} \neq \mathsf{O}, & \mathsf{Lili} \neq \mathsf{O}, \\ \mathsf{Bea} = \mathsf{Tom}, & \left(\begin{array}{c} \langle \mathsf{Bea}, \mathsf{Lea}, \mathsf{Leo}, \mathsf{Lio}, \mathsf{Lili}, \mathsf{Tom} \rangle, \\ \mathsf{val} - \mathsf{M} & \mathsf{occurrence} - O_{\mathsf{M}}, \\ \mathsf{val} - \mathsf{N} & \mathsf{occurrence} - O_{\mathsf{N}}, \\ \mathsf{val} - \mathsf{O} & \mathsf{occurrence} - O_{\mathsf{O}}, \end{array} \right)$$

where:

- line 1 declares the integer value of each shift type,
- lines 2, 3 and 4 declare the nurse and occurrence variables.
- line 5 enforces Bea, Lea and Leo not to work on a night shift,
- line 6 imposes Leo, Lio and Lili to work,
- line 7 constrains Bea and Tom to work on the same shift,
- line 8 restricts each shift type to occur within a given range.