

5.225 `lex_chain_greatereq`

	DESCRIPTION	LINKS	GRAPH
Origin	Derived from <code>lex_chain_lesseq</code>		
Constraint	<code>lex_chain_greatereq(VECTORS)</code>		
Usual name	<code>lex_chain</code>		
Type	VECTOR : <code>collection</code> (var-dvar)		
Argument	VECTORS : <code>collection</code> (vec - VECTOR)		
Restrictions	$ \text{VECTOR} \geq 1$ <code>required</code> (VECTOR, var) <code>required</code> (VECTORS, vec) <code>same_size</code> (VECTORS, vec)		
Purpose	<p>For each pair of consecutive vectors VECTOR_i and VECTOR_{i+1} of the VECTORS collection we have that VECTOR_i is lexicographically greater than or equal to VECTOR_{i+1}. Given two vectors, \vec{X} and \vec{Y} of n components, $\langle X_0, \dots, X_{n-1} \rangle$ and $\langle Y_0, \dots, Y_{n-1} \rangle$, \vec{X} is lexicographically greater than or equal to \vec{Y} if and only if $n = 0$ or $X_0 > Y_0$ or $X_0 = Y_0$ and $\langle X_1, \dots, X_{n-1} \rangle$ is lexicographically greater than or equal to $\langle Y_1, \dots, Y_{n-1} \rangle$.</p>		
Example	$((\text{vec} - \langle 5, 2, 6, 2 \rangle, \text{vec} - \langle 5, 2, 6, 2 \rangle, \text{vec} - \langle 5, 2, 3, 9 \rangle))$		
	<p>The <code>lex_chain_greatereq</code> constraint holds since:</p> <ul style="list-style-type: none"> The first vector $\langle 5, 2, 6, 2 \rangle$ of the VECTORS collection is lexicographically greater than or equal to the second vector $\langle 5, 2, 6, 2 \rangle$ of the VECTORS collection. The second vector $\langle 5, 2, 6, 2 \rangle$ of the VECTORS collection is lexicographically greater than or equal to the third vector $\langle 5, 2, 3, 9 \rangle$ of the VECTORS collection. 		
Typical	$ \text{VECTOR} > 1$ $ \text{VECTORS} > 1$		
Arg. properties	<ul style="list-style-type: none"> <code>Contractible</code> wrt. VECTORS. <code>Suffix-contractible</code> wrt. VECTORS.vec (remove items from same position). 		
Usage	<p>This constraint was motivated for breaking symmetry: more precisely when one wants to lexicographically order the consecutive columns of a matrix of decision variables. A further motivation is that using a set of lexicographic ordering constraints between two vectors does usually not allow to come up with a complete pruning.</p>		

Algorithm	<p>A filtering algorithm achieving arc-consistency for a chain of lexicographical ordering constraints is presented in [95].</p> <p>Six different ways of integrating a chain of lexicographical ordering constraints within non-overlapping constraints like diffn or geost and within their corresponding necessary condition like the cumulative constraint are shown in [3].</p>
See also	<p>common keyword: lex_between, lex_greater, lex_less, lex_lesseq (<i>lexicographic order</i>).</p> <p>implied by: lex_chain_greater (<i>non-strict order implied by strict order</i>).</p> <p>part of system of constraints: lex_greatereq.</p> <p>used in graph description: lex_greatereq.</p>
Keywords	<p>characteristic of a constraint: vector.</p> <p>constraint type: system of constraints, decomposition, order constraint.</p> <p>filtering: arc-consistency.</p> <p>heuristics: heuristics and lexicographical ordering.</p> <p>symmetry: symmetry, matrix symmetry, lexicographic order.</p>

Arc input(s)	VECTORS
Arc generator	$\text{PATH} \mapsto \text{collection}(\text{vectors1}, \text{vectors2})$
Arc arity	2
Arc constraint(s)	$\text{lex_lesseq}(\text{vectors1.vec}, \text{vectors2.vec})$
Graph property(ies)	$\text{NARC} = \text{VECTORS} - 1$

Graph model Parts (A) and (B) of Figure 5.498 respectively show the initial and final graph associated with the **Example** slot. Since we use the **NARC** graph property, the arcs of the final graph are stressed in bold. The **lex_chain_greatereq** constraint holds since all the arc constraints of the initial graph are satisfied.

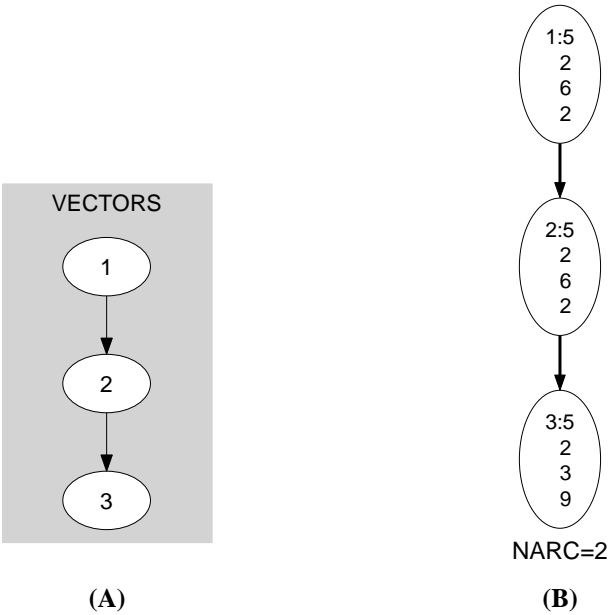


Figure 5.498: Initial and final graph of the **lex_chain_greatereq** constraint

Signature Since we use the *PATH* arc generator on the **VECTORS** collection the number of arcs of the initial graph is equal to $|\text{VECTORS}| - 1$. For this reason we can rewrite $\text{NARC} = |\text{VECTORS}| - 1$ to $\text{NARC} \geq |\text{VECTORS}| - 1$ and simplify $\overline{\text{NARC}}$ to $\overline{\text{NARC}}$.

