5.92 counts

DESCRIPTION LINKS GRAPH AUTOMATON

Origin Derived from count.

Constraint counts(VALUES, VARIABLES, RELOP, LIMIT)

Arguments

VALUES : collection(val−int)
VARIABLES : collection(var−dvar)
RELOP : atom
LIMIT : dvar

Restrictions

required(VALUES, val)
distinct(VALUES, val)
required(VARIABLES, var)
RELOP $\in [=, \neq, <, \geq, >, \leq]$

Purpose Let N be the number of variables of the VARIABLES collection assigned to a value of the VALUES collection. Enforce condition N RELOP LIMIT to hold.

Example $(\langle 1, 3, 4, 9\rangle, \langle 4, 5, 5, 4, 1, 5\rangle, =, 3)$

Values 1, 3, 4 and 9 of the VALUES collection are assigned to 3 items of the VARIABLES $= \langle 4, 5, 5, 4, 1, 5\rangle$ collection. The counts constraint holds since this number is in fact equal (RELOP is set to $=$) to the last argument of the counts constraint.

Typical

$|$VALUES$| > 1$
$|$VARIABLES$| > 1$
range(VARIABLES.var) > 1
$|$VARIABLES$| > |$VALUES$|$
RELOP $\in [=, <, \geq, >, \leq]$
LIMIT > 0
LIMIT $< |$VARIABLES$|$

Symmetries

- Items of VALUES are permutable.
- Items of VARIABLES are permutable.
- An occurrence of a value of VARIABLES.var that belongs to VALUES.val (resp. does not belong to VALUES.val) can be replaced by any other value in VALUES.val (resp. not in VALUES.val).

Arg. properties

- Contractible wrt. VARIABLES when RELOP $\in [<, \leq]$.
- Extensible wrt. VARIABLES when RELOP $\in [\geq, >]$.
- Aggregate: VALUES(sunion), VARIABLES(union), RELOP(id), LIMIT(+) when RELOP $\in [<, \leq, \geq, >]$.

Usage Used in the **Constraint(s) on sets** slot for defining some constraints like `assign_and_counts`.

Reformulation The `count`(`VALUES`, `VARIABLES`, `RELOP` , `LIMIT`) constraint can be expressed in term of the conjunction `among`(N, `VARIABLES`, `VALUES`) $\wedge$ N `RELOP` `LIMIT`.

Systems `count` in **Gecode**.

Used in `assign_and_counts`.

See also **assignment dimension added:** `assign_and_counts` *(assignment dimension introduced)*.

common keyword: `among` *(value constraint,counting constraint)*.

specialisation: `count` *(*`variable` $\in$ `VALUES` *replaced by* `variable`=`VALUE`*)*.

Keywords **characteristic of a constraint:** automaton, automaton with counters.

constraint network structure: alpha-acyclic constraint network(2).

constraint type: value constraint, counting constraint.

filtering: arc-consistency.

final graph structure: acyclic, bipartite, no loop.

Arc input(s)	VARIABLES VALUES
Arc generator	$PRODUCT$ ↦ collection(variables, values)
Arc arity	2
Arc constraint(s)	variables.var = values.val
Graph property(ies)	NARC RELOP LIMIT
Graph class	• ACYCLIC • BIPARTITE • NO_LOOP

Graph model Because of the arc constraint variables.var = values.val and since each domain variable can take at most one value, NARC is the number of variables taking a value in the VALUES collection.

Parts (A) and (B) of Figure 5.223 respectively show the initial and final graph associated with the **Example** slot. Since we use the NARC graph property, the arcs of the final graph are stressed in bold.

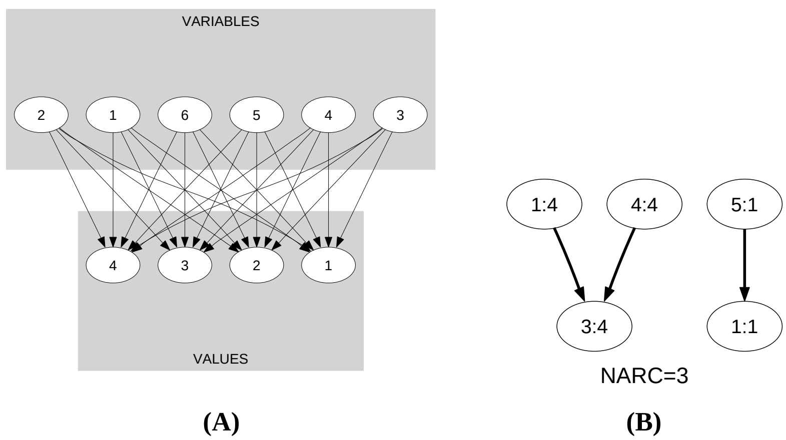

Figure 5.223: Initial and final graph of the counts constraint

Automaton Figure 5.224 depicts the automaton associated with the `counts` constraint. To each variable $\mathtt{VAR}_i$ of the collection `VARIABLES` corresponds a 0-1 signature variable S_i. The following signature constraint links $\mathtt{VAR}_i$ and S_i: $\mathtt{VAR}_i \in \mathtt{VALUES} \Leftrightarrow S_i$.

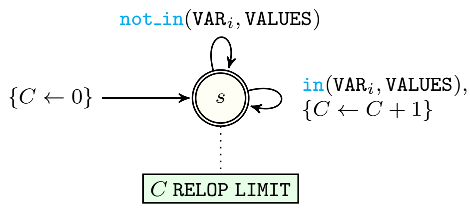

Figure 5.224: Automaton of the `counts` constraint

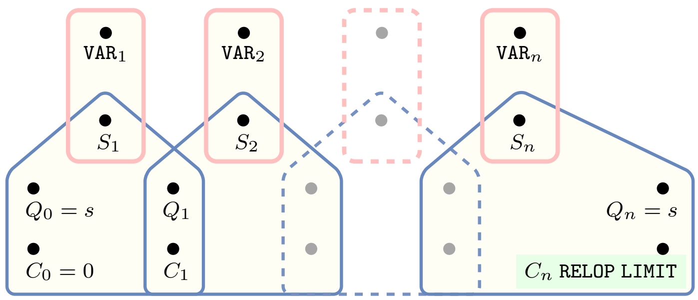

Figure 5.225: Hypergraph of the reformulation corresponding to the automaton (with one counter) of the `counts` constraint: since all states variables $Q_0, Q_1, \ldots, Q_n$ are fixed to the unique state s of the automaton, the transitions constraints share only the counter variable C and the constraint network is Berge-acyclic