

5.11 all\_min\_dist

	DESCRIPTION	LINKS	GRAPH
Origin	[343]		
Constraint	all_min_dist(MINDIST, VARIABLES)		
Synonyms	minimum_distance, inter_distance.		
Arguments	MINDIST : int VARIABLES : collection(var-dvar)		
Restrictions	MINDIST > 0 $  \text{VARIABLES}   < 2 \vee \text{MINDIST} < \text{range}(\text{VARIABLES.var})$ <code>required</code> (VARIABLES, var)		
Purpose	Enforce for each pair $(\text{var}_i, \text{var}_j)$ of distinct variables of the collection VARIABLES that $ \text{var}_i - \text{var}_j  \geq \text{MINDIST}$ .		
Example	<div>(2, &lt;5, 1, 9, 3&gt;)</div> <p>The all_min_dist constraint holds since the following expressions <math> 5 - 1 </math>, <math> 5 - 9 </math>, <math> 5 - 3 </math>, <math> 1 - 9 </math>, <math> 1 - 3 </math>, <math> 9 - 3 </math> are all greater than or equal to the first argument MINDIST = 2 of the all_min_dist constraint.</p>		
All solutions	<p>Figure 5.24 gives all solutions to the following non ground instance of the all_min_dist constraint: <math>V_1 \in [0, 5]</math>, <math>V_2 \in [3, 9]</math>, <math>V_3 \in [5, 7]</math>, <math>V_4 \in [2, 10]</math>, all_min_dist(3, &lt;V<sub>1</sub>, V<sub>2</sub>, V<sub>3</sub>, V<sub>4</sub>&gt;).</p> <div><div>① (3, &lt;0, 3, 6, 9&gt;) ② (3, &lt;0, 3, 6, 10&gt;) ③ (3, &lt;0, 3, 7, 10&gt;) ④ (3, &lt;0, 4, 7, 10&gt;) ⑤ (3, &lt;0, 9, 6, 3&gt;) ⑥ (3, &lt;1, 4, 7, 10&gt;)</div></div>		
Typical	MINDIST > 1 $  \text{VARIABLES}   > 1$		

Figure 5.24: All solutions corresponding to the non ground example of the all\_min\_dist constraint of the All solutions slot

**Symmetries**

- MINDIST can be [decreased](#) to any value  $\geq 1$ .
- Items of VARIABLES are [permutable](#).
- Two distinct values of VARIABLES.var can be [swapped](#).
- One and the same constant can be [added](#) to the var attribute of all items of VARIABLES.

**Arg. properties**

[Contractible](#) wrt. VARIABLES.

**Usage**

The `all_min_dist` constraint was initially created for handling frequency allocation problems. In [11] it is used for scheduling tasks that all have the same fixed duration in the context of [air traffic management](#) in the terminal radar control area of airports.

**Remark**

The `all_min_dist` constraint can be modelled as a set of tasks that should not overlap. For each variable `var` of the VARIABLES collection we create a task  $t$  where `var` and MINDIST respectively correspond to the origin and the duration of  $t$ .

Some solvers use in a pre-processing phase, while stating constraints of the form  $|X_i - X_j| \geq D_{ij}$  (where  $X_i$  and  $X_j$  are domain variables and  $D_{ij}$  is a constant), an algorithm for automatically extracting large cliques [88] from such inequalities in order to state `all_min_dist` constraints.

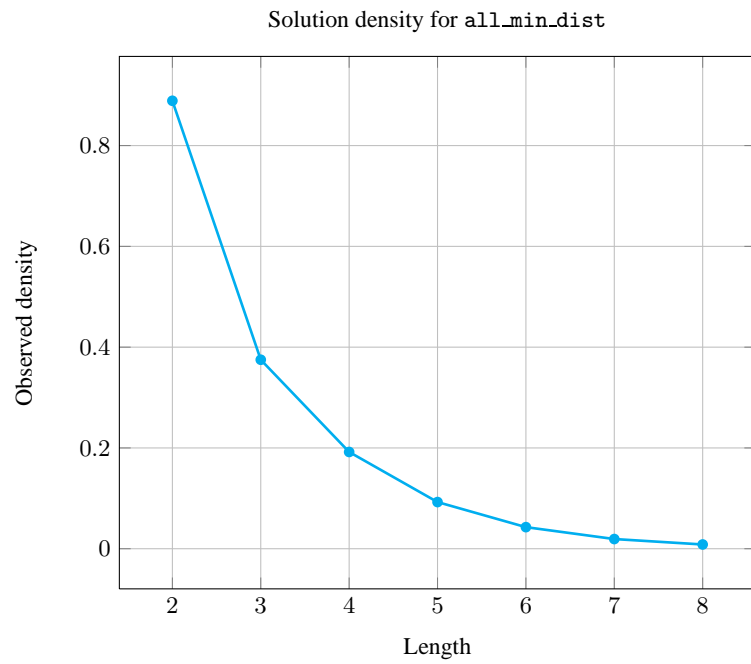
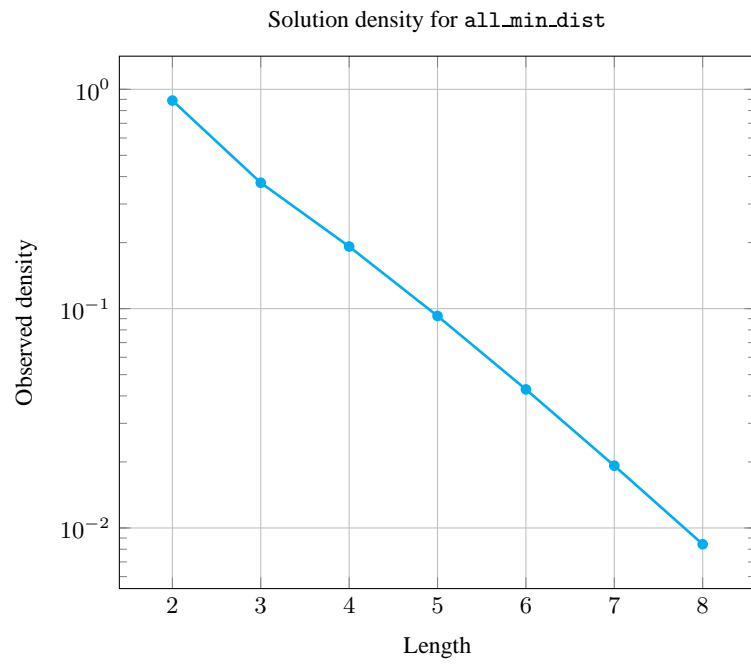
**Algorithm**

K. Artiouchine and P. Baptiste came up with a cubic time complexity algorithm achieving [bound-consistency](#) in [11, 12] based on the adaptation of a feasibility test algorithm from M.R. Garey *et al.* [185]. Later on, C.-G. Quimper *et al.*, proposed a quadratic algorithm achieving the same level of consistency in [332].

**Counting**

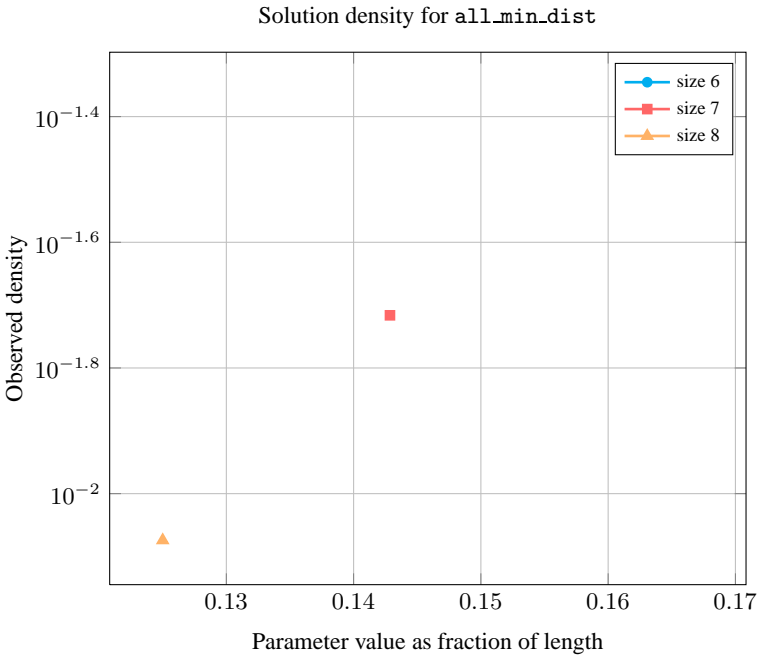
Length ( $n$ )	2	3	4	5	6	7	8
Solutions	8	24	120	720	5040	40320	362880

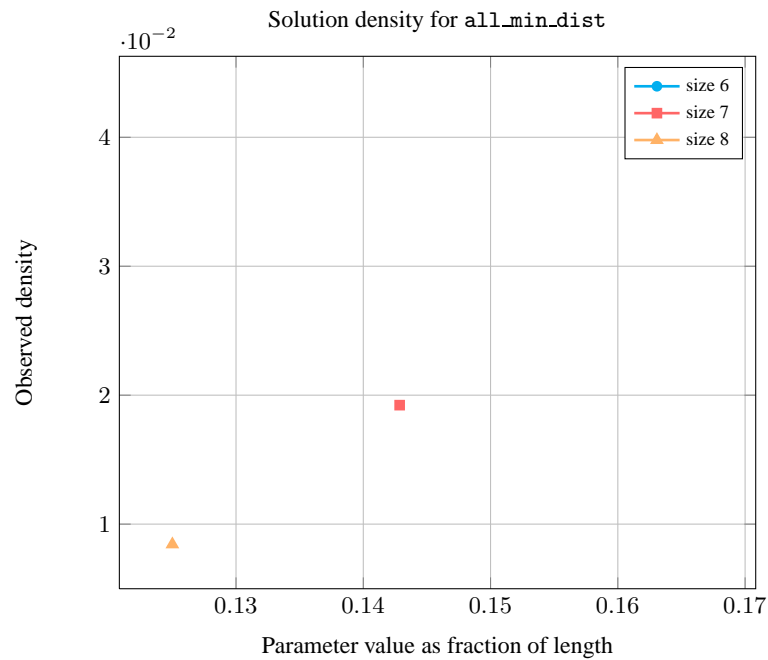
Number of solutions for `all_min_dist`: domains  $0..n$



Length ( $n$ )		2	3	4	5	6	7	8
Total		8	24	120	720	5040	40320	362880
Parameter value	1	6	24	120	720	5040	40320	362880
	2	2	-	-	-	-	-	-

Solution count for all\_min\_dist: domains 0.. $n$



**See also**

**generalisation:** `diffn`(line segment, of same length, replaced by *orthotope*),  
**disjunctive**(line segment, of same length, replaced by line segment),  
**multi\_inter\_distance**(LIMIT parameter introduced to specify capacity  $\geq 1$ ).

**implies:** `alldifferent_interval`, `soft_alldifferent_var`.

**related:** `distance`.

**specialisation:** `alldifferent`(line segment, of same length, replaced by variable).

**Keywords**

**application area:** frequency allocation problem, air traffic management.

**characteristic of a constraint:** sort based reformulation.

**constraint type:** value constraint, decomposition, scheduling constraint.

**filtering:** bound-consistency.

**final graph structure:** acyclic.

**problems:** maximum clique.

**Cond. implications**

`all_min_dist`(MINDIST, VARIABLES)

**implies** `soft_all_equal_max_var`(N, VARIABLES)

when  $N \geq |\text{VARIABLES}| - 1$ .

<b>Arc input(s)</b>	VARIABLES
<b>Arc generator</b>	$CLIQUE(<) \mapsto \text{collection}(\text{variables1}, \text{variables2})$
<b>Arc arity</b>	2
<b>Arc constraint(s)</b>	$\text{abs}(\text{variables1.var} - \text{variables2.var}) \geq \text{MINDIST}$
<b>Graph property(ies)</b>	$NARC =  \text{VARIABLES}  * ( \text{VARIABLES}  - 1) / 2$
<b>Graph class</b>	<ul style="list-style-type: none"> <li>• ACYCLIC</li> <li>• NO_LOOP</li> </ul>

**Graph model**

We generate a *clique* with a minimum distance constraint between each pair of distinct vertices and state that the number of arcs of the final graph should be equal to the number of arcs of the initial graph.

Parts (A) and (B) of Figure 5.25 respectively show the initial and final graph associated with the **Example** slot. The `all_min_dist` constraint holds since all the arcs of the initial graph belong to the final graph: all the minimum distance constraints are satisfied.

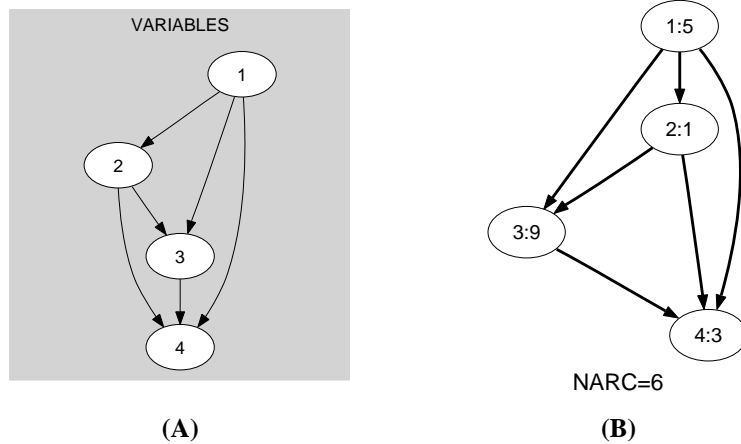


Figure 5.25: Initial and final graph of the `all_min_dist` constraint