## 5.283 not\_all\_equal

DESCRIPTION LINKS GRAPH AUTOMATON

Origin CHIP

Constraint not\_all\_equal(VARIABLES)

Argument VARIABLES : collection(var-dvar)

Restrictions required(VARIABLES, var)

 $|{\tt VARIABLES}| > 1$ 

**Purpose** The variables of the collection VARIABLES should take more than a single value.

Example  $(\langle 3, 1, 3, 3, 3 \rangle)$ 

The not\_all\_equal constraint holds since the collection (3,1,3,3,3) involves more than one value (i.e., values 1 and 3).

Typical |VARIABLES| > 2nval(VARIABLES.var) > 2

Symmetries • Items of VARIABLES are permutable.

All occurrences of two distinct values of VARIABLES.var can be swapped; all
occurrences of a value of VARIABLES.var can be renamed to any unused value.

Arg. properties Extensible wrt. VARIABLES.

If the intersection of the domains of the variables of the VARIABLES collection is empty the  $not\_all\_equal$  constraint is entailed. Otherwise, when only a single variable V remains not fixed, remove the unique value (unique since the constraint is not entailed) taken by the other variables from the domain of V.

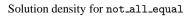
**Reformulation** The not\_all\_equal(VARIABLES) constraint can be expressed as atleast\_nvalue(2, VARIABLES).

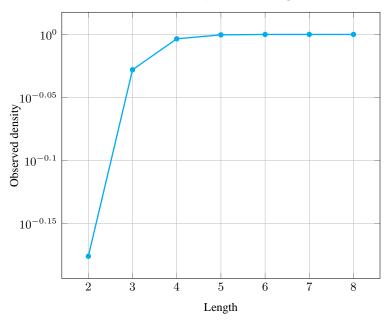
**Counting** 

Algorithm

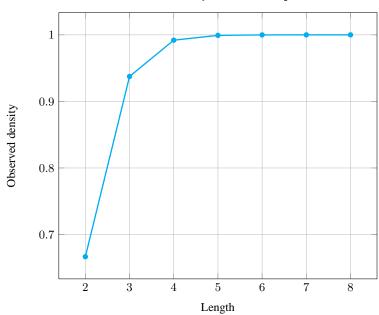
Length (n)	2	3	4	5	6	7	8
Solutions	6	60	620	7770	117642	2097144	43046712

Number of solutions for not\_all\_equal: domains 0..n





## Solution density for not\_all\_equal



Systems

rel in Gecode.

See also

**generalisation:** nvalue (introduce a variable for counting the number of distinct values).

implied by: alldifferent.

negation: all\_equal.

**specialisation:** neq (when go down to two variables).

used in reformulation: atleast\_nvalue.

Keywords

characteristic of a constraint: disequality, automaton, automaton without counters,

reified automaton constraint.

**constraint network structure:** sliding cyclic(1) constraint network(1).

constraint type: value constraint.

**filtering:** arc-consistency.

final graph structure: equivalence.

Arc input(s)	VARIABLES			
Arc generator	$CLIQUE \mapsto \texttt{collection}(\texttt{variables1}, \texttt{variables2})$			
Arc arity	2			
Arc constraint(s)	${\tt variables1.var} = {\tt variables2.var}$			
Graph property(ies)	NSCC>1			

## Graph model

Parts (A) and (B) of Figure 5.608 respectively show the initial and final graph associated with the **Example** slot. Since we use the **NSCC** graph property we show the different strongly connected components of the final graph. Each strongly connected component corresponds to a value that is assigned to some variables of the VARIABLES collection. The not\_all\_equal holds since the final graph contains more than one strongly connected component.

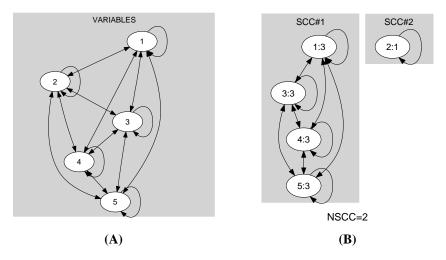


Figure 5.608: Initial and final graph of the not\_all\_equal constraint

Automaton

Figure 5.609 depicts the automaton associated with the not\_all\_equal constraint. To each pair of consecutive variables (VAR $_i$ , VAR $_{i+1}$ ) of the collection VARIABLES corresponds a signature variable  $S_i$ . The following signature constraint links VAR $_i$ , VAR $_{i+1}$  and  $S_i$ : VAR $_i =$ VAR $_{i+1} \Leftrightarrow S_i$ .

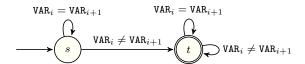


Figure 5.609: Automaton of the not\_all\_equal constraint

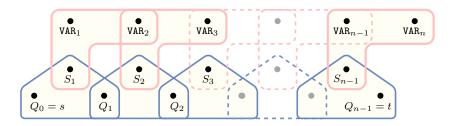


Figure 5.610: Hypergraph of the reformulation corresponding to the automaton of the not\_all\_equal constraint