5.123 disjoint

DESCRIPTION LINKS GRAPH AUTOMATON

Origin Derived from alldifferent.

Constraint disjoint(VARIABLES1, VARIABLES2)

Restrictions required(VARIABLES1, var)
required(VARIABLES2, var)

Each variable of the collection VARIABLES1 should take a value that is distinct from all the values assigned to the variables of the collection VARIABLES2.

Example $(\langle 1, 9, 1, 5 \rangle, \langle 2, 7, 7, 0, 6, 8 \rangle)$

Purpose

In this example, values 1,5,9 are used by the variables of VARIABLES1 and values 0,2,6,7,8 by the variables of VARIABLES2. Since there is no intersection between the two previous sets of values the disjoint constraint holds.

All solutions Figure 5.292 gives all solutions to the following non ground instance of the disjoint constraint: $U_1 \in [0,2], U_2 \in [1,2], U_3 \in [1,2], V_1 \in [0,1], V_2 \in [1,2],$ disjoint($\langle U_1, U_2, U_3 \rangle, \langle V_1, V_2 \rangle$).

 $(\langle 0, 2, 2 \rangle, \langle 1, 1 \rangle)$ $(\langle 1, 1, 1 \rangle, \langle 0, 2 \rangle)$ $(\langle 2, 2, 2 \rangle, \langle 0, 1 \rangle)$ $(\langle 2, 2, 2 \rangle, \langle 1, 1 \rangle)$

Figure 5.292: All solutions corresponding to the non ground example of the disjoint constraint of the **All solutions** slot

Typical |VARIABLES1| > 1|VARIABLES2| > 1 20000315

Symmetries

- Arguments are permutable w.r.t. permutation (VARIABLES1, VARIABLES2).
- Items of VARIABLES1 are permutable.
- Items of VARIABLES2 are permutable.
- An occurrence of a value of VARIABLES1.var can be replaced by any value of VARIABLES1.var.
- An occurrence of a value of VARIABLES2.var can be replaced by any value of VARIABLES2.var.
- All occurrences of two distinct values in VARIABLES1.var or VARIABLES2.var
 can be swapped; all occurrences of a value in VARIABLES1.var or
 VARIABLES2.var can be renamed to any unused value.

Arg. properties

- Contractible wrt. VARIABLES1.
- Contractible wrt. VARIABLES2.

Remark

Despite the fact that this is not an uncommon constraint, it can not be modelled in a compact way neither with a *disequality* constraint (i.e., two given variables have to take distinct values) nor with the alldifferent constraint. The disjoint constraint can bee seen as a special case of the common(NCOMMON1, NCOMMON2, VARIABLES1, VARIABLES2) constraint where NCOMMON1 and NCOMMON2 are both set to 0.

MiniZinc (http://www.minizinc.org/) has a disjoint constraint between two set variables rather than between two collections of variables.

Algorithm

Let us note:

- n₁ the minimum number of distinct values taken by the variables of the collection VARIABLES1.
- n₂ the minimum number of distinct values taken by the variables of the collection VARIABLES2.
- n₁₂ the maximum number of distinct values taken by the union of the variables of VARIABLES1 and VARIABLES2.

One invariant to maintain for the disjoint constraint is $n_1 + n_2 \le n_{12}$. A lower bound of n_1 and n_2 can be obtained by using the algorithms provided in [27, 40]. An exact upper bound of n_{12} can be computed by using a bipartite matching algorithm.

Used in

k_disjoint.

See also

generalisation: disjoint_tasks (variable replaced by task).

 $implies: \verb|alldifferent_on_intersection|, \verb|lex_different|.|$

system of constraints: k_disjoint.

Keywords

 $\textbf{characteristic of a constraint:} \ disequality, \ automaton, \ automaton \ with \ array \ of \ counters.$

filtering: bipartite matching.

constraint type: value constraint.

modelling: empty intersection.

Arc input(s) VARIABLES1 VARIABLES2

Arc generator $PRODUCT \mapsto collection(variables1, variables2)$

Arc arity 2

Arc constraint(s) variables1.var = variables2.var

Graph property(ies) NARC= 0

Graph model

PRODUCT is used in order to generate the arcs of the graph between all variables of VARIABLES1 and all variables of VARIABLES2. Since we use the graph property **NARC** = 0 the final graph will be empty. Figure 5.293 shows the initial graph associated with the **Example** slot. Since we use the **NARC** = 0 graph property the final graph is empty.

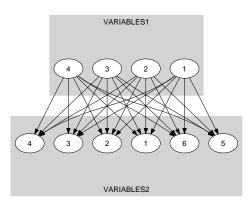


Figure 5.293: Initial graph of the disjoint constraint (the final graph is empty)

Signature

Since 0 is the smallest number of arcs of the final graph we can rewrite NARC = 0 to $NARC \le 0$. This leads to simplify NARC to NARC.

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Automaton

Figure 5.294 depicts the automaton associated with the disjoint constraint. To each variable VAR1 $_i$ of the collection VARIABLES1 corresponds a signature variable S_i that is equal to 0. To each variable VAR2 $_i$ of the collection VARIABLES2 corresponds a signature variable $S_{i+|\text{VARIABLES1}|}$ that is equal to 1.

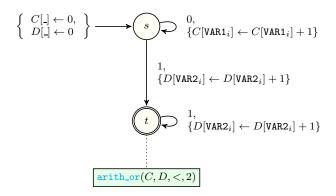


Figure 5.294: Automaton of the $\mathtt{disjoint}(\mathtt{VARIABLES1},\mathtt{VARIABLES2})$ constraint, where state s handles variables of the collection $\mathtt{VARIABLES1}$ and state t handles variables of the collection $\mathtt{VARIABLES2}$