

5.107 cyclic_change

| | DESCRIPTION | LINKS | GRAPH | AUTOMATON |
|------------------------|--|-------|-------|-----------|
| Origin | Derived from change . | | | |
| Constraint | <code>cyclic_change(NCHANGE, CYCLE_LENGTH, VARIABLES, CTR)</code> | | | |
| Arguments | NCHANGE : dvar CYCLE_LENGTH : int VARIABLES : collection (var — dvar) CTR : atom | | | |
| Restrictions | NCHANGE ≥ 0 NCHANGE $< \text{VARIABLES} $ CYCLE_LENGTH > 0 required (VARIABLES, var) VARIABLES. var ≥ 0 VARIABLES. var $< \text{CYCLE_LENGTH}$ CTR $\in [=, \neq, <, \geq, >, \leq]$ | | | |
| Purpose | NCHANGE is the number of times that constraint $((X + 1) \bmod \text{CYCLE_LENGTH}) \text{ CTR } Y$ holds; X and Y correspond to consecutive variables of the collection VARIABLES. | | | |
| Example | <div>(2, 4, ⟨3, 0, 2, 3, 1⟩, ≠)</div> <p>Since CTR is set to \neq and since CYCLE_LENGTH is set to 4, a change between two consecutive items X and Y of the VARIABLES collection corresponds to the fact that the condition $((X + 1) \bmod 4) \neq Y$ holds. Consequently, the <code>cyclic_change</code> constraint holds since we have the two following changes (i.e., NCHANGE = 2) within ⟨3, 0, 2, 3, 1⟩:</p> <ul style="list-style-type: none"> • A first change between the consecutive values 0 and 2, • A second change between the consecutive values 3 and 1. <p>However, the sequence 3 0 does not correspond to a change since $(3 + 1) \bmod 4$ is equal to 0.</p> | | | |
| Typical | NCHANGE > 0 VARIABLES > 1 range (VARIABLES. var) > 1 CTR $\in [\neq]$ | | | |
| Symmetry | Items of VARIABLES can be shifted . | | | |
| Arg. properties | Functional dependency: NCHANGE determined by CYCLE_LENGTH, VARIABLES and CTR. | | | |

Usage

This constraint may be used for personnel [cyclic](#) timetabling problems where each person has to work according to cycles. In this context each variable of the VARIABLES collection corresponds to the type of work a person performs on a specific day. Because of some perturbation (e.g., illness, unavailability, variation of the workload) it is in practice not reasonable to ask for perfect [cyclic](#) solutions. One alternative is to use the [cyclic_change](#) constraint and to ask for solutions where one tries to minimise the number of cycle breaks (i.e., the variable NCHANGE).

See also

common keyword: [change](#), [cyclic_change_joker](#) (*number of changes*).

implies: [cyclic_change_joker](#).

Keywords

characteristic of a constraint: [cyclic](#), [automaton](#), [automaton with counters](#).

constraint arguments: pure functional dependency.

constraint network structure: sliding [cyclic\(1\)](#) constraint [network\(2\)](#).

constraint type: timetabling constraint.

final graph structure: [acyclic](#), [bipartite](#), [no loop](#).

modelling: number of changes, functional dependency.

| | |
|---------------------|---|
| Arc input(s) | VARIABLES |
| Arc generator | <i>PATH</i> \mapsto collection(variables1, variables2) |
| Arc arity | 2 |
| Arc constraint(s) | (variables1.var + 1) mod CYCLE_LENGTH CTR variables2.var |
| Graph property(ies) | NARC = NCHANGE |
| Graph class | <ul style="list-style-type: none">• ACYCLIC• BIPARTITE• NO_LOOP |

Graph model Parts (A) and (B) of Figure 5.252 respectively show the initial and final graph associated with the **Example** slot. Since we use the **NARC** graph property, the arcs of the final graph are stressed in bold.

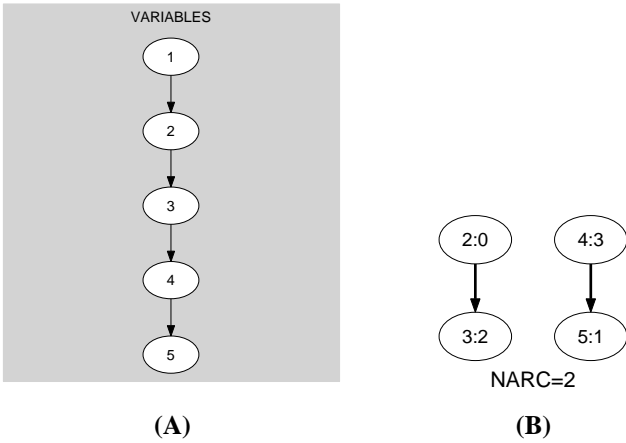
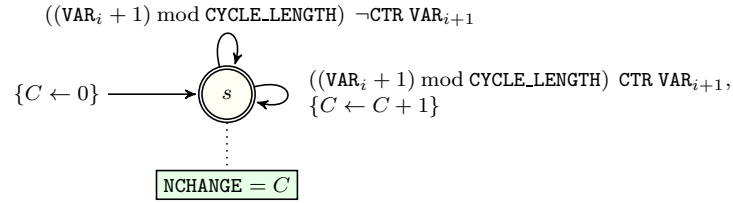
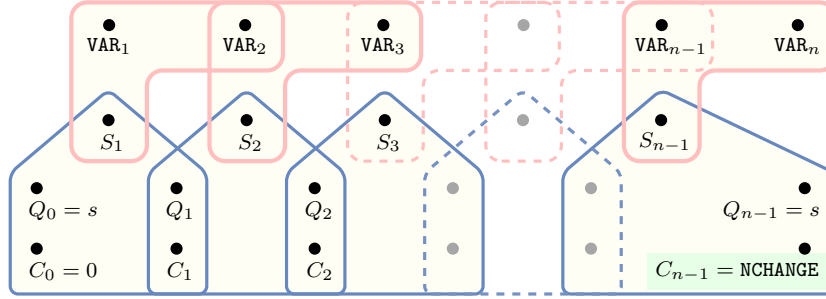


Figure 5.252: Initial and final graph of the cyclic_change constraint

Automaton

Figure 5.253 depicts the automaton associated with the `cyclic_change` constraint. To each pair of consecutive variables $(\text{VAR}_i, \text{VAR}_{i+1})$ of the collection `VARIABLES` corresponds a 0-1 signature variable S_i . The following signature constraint links VAR_i , VAR_{i+1} and S_i :

Figure 5.253: Automaton of the `cyclic_change` constraintFigure 5.254: Hypergraph of the reformulation corresponding to the automaton of the `cyclic_change` constraint