

5.295 open_alldifferent

	DESCRIPTION	LINKS	GRAPH
Origin	[427]		
Constraint	<code>open_alldifferent(S, VARIABLES)</code>		
Synonyms	<code>open_alldiff</code> , <code>open_alldistinct</code> , <code>open_distinct</code> .		
Arguments	<code>S</code> : <code>svar</code> <code>VARIABLES</code> : <code>collection</code> (<code>var—dvar</code>)		
Restrictions	$S \geq 1$ $S \leq \text{VARIABLES} $ <code>required</code> (<code>VARIABLES</code> , <code>var</code>)		
Purpose	Let \mathcal{V} be the variables of the collection <code>VARIABLES</code> for which the corresponding position belongs to the set <code>S</code> . Positions are numbered from 1. Enforce all variables of \mathcal{V} to take distinct values.		
Example	$(\{2, 3, 4\}, \langle 9, 1, 9, 3 \rangle)$		
	The <code>open_alldifferent</code> constraint holds since the last three (i.e., $S = \{2, 3, 4\}$) values of the collection $\langle 9, 1, 9, 3 \rangle$ are distinct.		
Typical	$ \text{VARIABLES} > 2$		
Symmetry	All occurrences of two distinct values of <code>VARIABLES.var</code> can be <code>swapped</code> ; all occurrences of a value of <code>VARIABLES.var</code> can be <code>renamed</code> to any unused value.		
Arg. properties	<code>Suffix-contractible</code> wrt. <code>VARIABLES</code> .		
Usage	<p>In their article [427], W.-J. van Hoes and J.-C. Régin motivate the <code>open_alldifferent</code> constraint by the following scheduling problem. Consider a set of activities (where each activity has a fixed duration 1 and a start variable) that can be processed on two factory lines such that all the activities that will be processed on a given line must be pairwise distinct. This can be modelled by using one <code>open_alldifferent</code> constraint for each line, involving all the start variables as well as a set variable whose final value specifies the set of activities assigned to that specific factory line.</p> <p>Note that this can also be directly modelled by a single <code>diffn</code> constraint. This is done by introducing an assignment variable for each activity. The initial domain of each assignment variable consists of two values that respectively correspond to the two factory lines.</p>		
Algorithm	A slight adaptation of the <code>flow</code> model that handles the original <code>global_cardinality</code> constraint [342] is described in [427]. The rightmost part of Figure 3.29 illustrates this flow model.		

See also

common keyword: [size_max_seq_alldifferent](#), [size_max_starting_seq_alldifferent](#) (*all different, disequality*).

generalisation: [open_global_cardinality](#) (control the number of occurrence of each active value¹³ with a counter variable), [open_global_cardinality_low_up](#) (control the number of occurrence of each active value with an interval).

hard version: [alldifferent](#).

used in graph description: [in_set](#).

Keywords

characteristic of a constraint: all different, disequality.

constraint arguments: constraint involving set variables.

constraint type: open constraint, soft constraint, value constraint.

filtering: flow.

¹³An *active value* corresponds to a value occuring at a position mentionned in the set S.

Arc input(s)	VARIABLES
Arc generator	<i>CLIQUE</i> \mapsto <i>collection</i> (variables1, variables2)
Arc arity	2
Arc constraint(s)	<ul style="list-style-type: none">• variables1.var = variables2.var• <i>in_set</i>(variables1.key, S)• <i>in_set</i>(variables2.key, S)
Graph property(ies)	<i>MAX_NSCC</i> \leq 1
Graph class	<i>ONE_SUCC</i>

Graph model

We generate a *clique* with an *equality* constraint between each pair of vertices (including a vertex and itself) and state that the size of the largest strongly connected component should not exceed one. Variables for which the corresponding position does not belong to the set S are removed from the final graph by the second and third conditions of the arc-constraint.

Parts (A) and (B) of Figure 5.630 respectively show the initial and final graph associated with the **Example** slot. Since we use the *MAX_NSCC* graph property we show one of the largest strongly connected components of the final graph. The *open_alldifferent* holds since all the strongly connected components have at most one vertex: a value is used at most once.

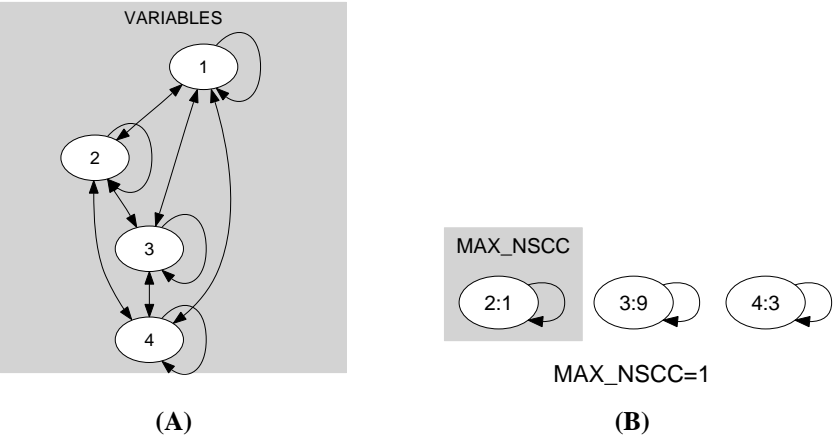


Figure 5.630: Initial and final graph of the *open_alldifferent* constraint

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