

5.261 min\_width\_peak

	DESCRIPTION	LINKS	AUTOMATON
Origin	derived from <a href="#">peak</a>		
Constraint	min_width_peak(MIN_WIDTH, VARIABLES)		
Synonym	min_base_peak.		
Arguments	MIN_WIDTH : dvar VARIABLES : collection(var—dvar)		
Restrictions	MIN_WIDTH ≥ 0 MIN_WIDTH ≤  VARIABLES  − 2 <a href="#">required</a> (VARIABLES, var)		
Purpose	Given a sequence VARIABLES constraint MIN_WIDTH to be fixed to the width of the smallest peak, or to 0 if no peak exists.		
Example	<div>(5, ⟨4, 4, 2, 2, 3, 5, 5, 6, 3, 1, 1, 2, 2, 2, 2, 2, 1⟩)</div> <div>(5, ⟨4, 6, 7, 9, 8, 5, 4⟩)</div> <div>(0, ⟨4, 4, 2, 0, 0, 4, 5⟩)</div>		

The first min\_width\_peak constraint holds since the sequence 4 4 2 2 3 5 5 6 3 1 1 2 2 2 2 2 1 contains two peaks of respective width 5 and 6 (see Figure 5.564) and since its argument MIN\_WIDTH is fixed to the smallest value 5.

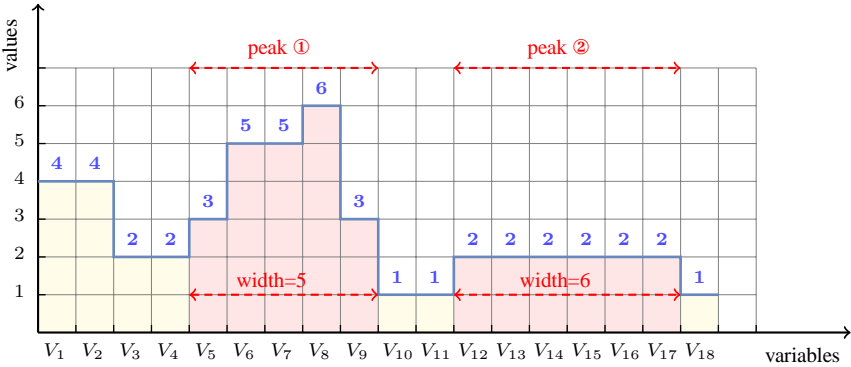


Figure 5.564: Illustration of the first example of the **Example** slot: a sequence of eighteen variables  $V_1, V_2, V_3, V_4, V_5, V_6, V_7, V_8, V_9, V_{10}, V_{11}, V_{12}, V_{13}, V_{14}, V_{15}, V_{16}, V_{17}, V_{18}$  respectively fixed to values 4, 4, 2, 2, 3, 5, 5, 6, 3, 1, 1, 2, 2, 2, 2, 2, 2, 1 and its two peaks of width 5 and 6.

Typical

```
MIN_WIDTH > 1
|VARIABLES| > 2
```

Symmetries

- Items of VARIABLES can be [reversed](#).
- One and the same constant can be [added](#) to the var attribute of all items of VARIABLES.

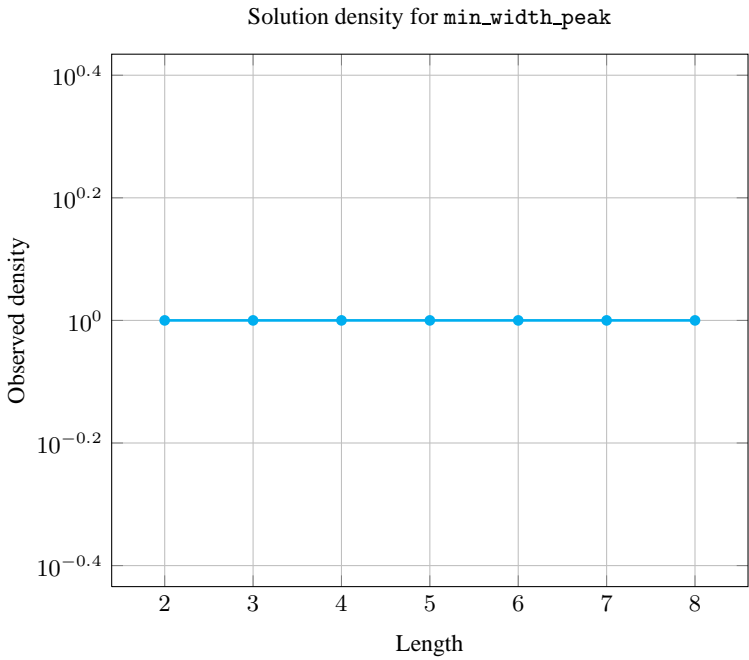
Arg. properties

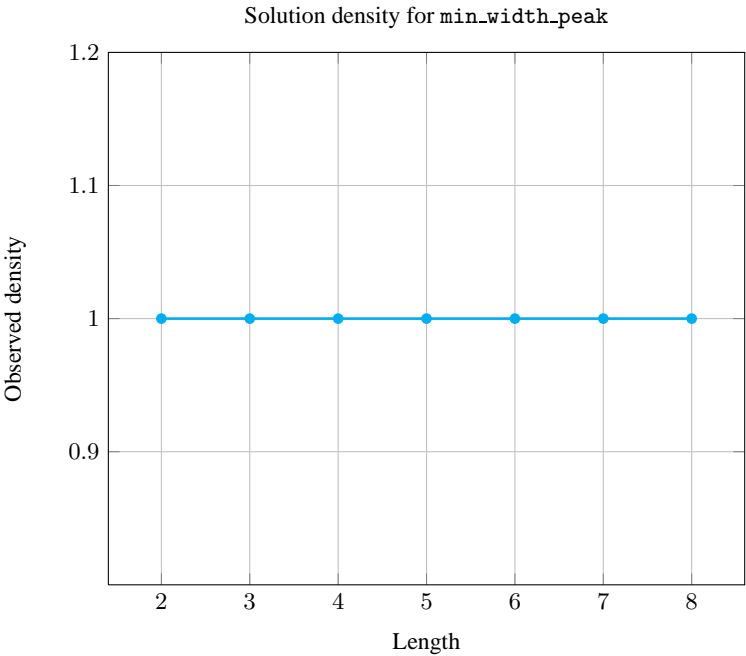
[Functional dependency](#): MIN\_WIDTH determined by VARIABLES.

Counting

Length ( <i>n</i> )	2	3	4	5	6	7	8
Solutions	9	64	625	7776	117649	2097152	43046721

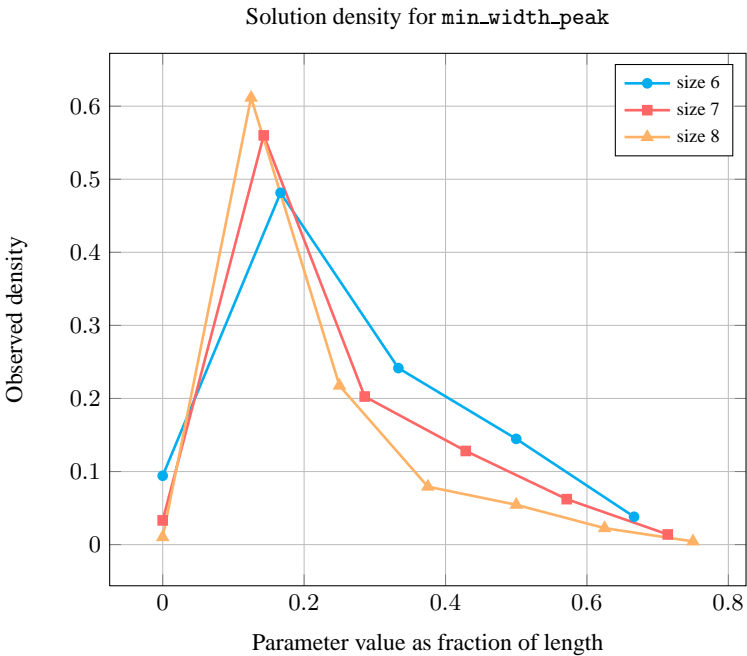
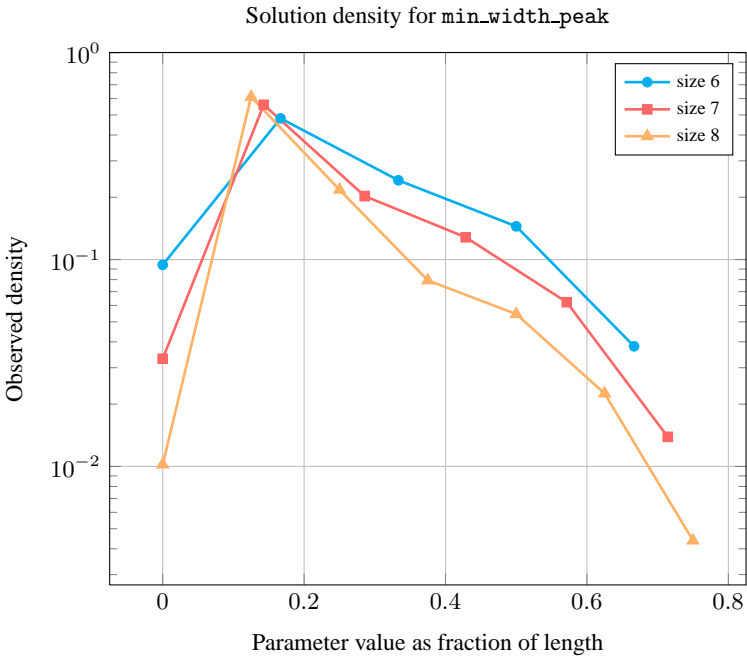
Number of solutions for min\_width\_peak: domains 0..*n*





Length ( <i>n</i> )		2	3	4	5	6	7	8
Total		9	64	625	7776	117649	2097152	43046721
Parameter value	0	9	50	295	1792	11088	69498	439791
	1	-	14	230	3205	56637	1174398	26327058
	2	-	-	100	2100	28420	424928	9363060
	3	-	-	-	679	17024	268722	3413256
	4	-	-	-	-	4480	130452	2345982
	5	-	-	-	-	-	29154	968946
	6	-	-	-	-	-	-	188628

Solution count for min\_width\_peak: domains 0..*n*



See also [common keyword: peak \(sequence\)](#).

Keywords [characteristic of a constraint:](#) [automaton](#), [automaton with counters](#),

automaton with same input symbol.

**combinatorial object:** sequence.

**constraint arguments:** reverse of a constraint, pure functional dependency.

**filtering:** glue matrix.

**modelling:** functional dependency.

**Automaton**

Figure 5.565 depicts the automaton associated with the `min_width_peak` constraint. To each pair of consecutive variables  $(\text{VAR}_i, \text{VAR}_{i+1})$  of the collection `VARIABLES` corresponds a signature variable  $S_i$ . The following signature constraint links  $\text{VAR}_i$ ,  $\text{VAR}_{i+1}$  and  $S_i$ :  $(\text{VAR}_i < \text{VAR}_{i+1} \Leftrightarrow S_i = 0) \wedge (\text{VAR}_i = \text{VAR}_{i+1} \Leftrightarrow S_i = 1) \wedge (\text{VAR}_i > \text{VAR}_{i+1} \Leftrightarrow S_i = 2)$ .

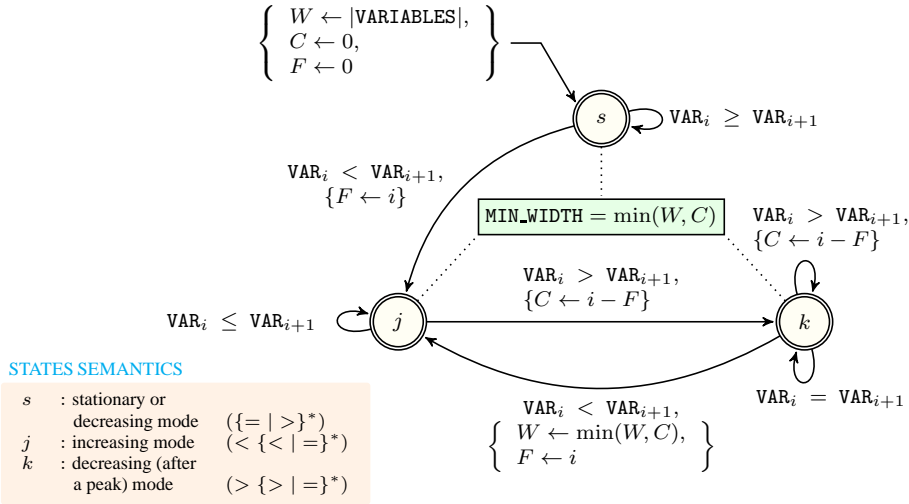


Figure 5.565: Automaton of the `min_width_peak` constraint: the start of the first potential peak is discovered while triggering the transition from  $s$  to  $j$ , the top of a peak is discovered while triggering the transition from  $j$  to  $k$ , the end of a peak and the start of the next potential peak are discovered while triggering the transition from  $k$  to  $j$ ; the counters  $W$ ,  $C$  and  $F$  respectively stand for *min\_width*, *current* and *first*.

Blue matrix where  $\vec{W}$ ,  $\vec{C}$ ,  $\vec{F}$  and  $\overleftarrow{W}$ ,  $\overleftarrow{C}$ ,  $\overleftarrow{F}$  resp. represent the counters values  $W$ ,  $C$ ,  $F$  at the end of a prefix and at the end of the corresponding reverse suffix that partitions the sequence VARIABLES;  $\overrightarrow{\text{MIN\_WIDTH}}$  (resp.  $\overleftarrow{\text{MIN\_WIDTH}}$ ) stands for  $\min(\vec{W}, \vec{C})$  (resp.  $\min(\overleftarrow{W}, \overleftarrow{C})$ ).

	$s (\{>   =\}^*)$	$j (\{<   =\}^*)$	$k (\{>   =\}^*)$
$s (\{>   =\}^*)$			
$j (\{<   =\}^*)$		$\min \left( \begin{array}{c} \vec{W}, \\ n - \vec{F} - \overleftarrow{F}, \\ \overleftarrow{W} \end{array} \right)$	$\min \left( \begin{array}{c} \overrightarrow{\text{MIN\_WIDTH}}, \\ n - \vec{F} - \overleftarrow{F}, \\ \overleftarrow{\text{MIN\_WIDTH}} \end{array} \right)$
$k (\{>   =\}^*)$		$\min \left( \begin{array}{c} \overrightarrow{\text{MIN\_WIDTH}}, \\ n - \vec{F} - \overleftarrow{F}, \\ \overleftarrow{\text{MIN\_WIDTH}} \end{array} \right)$	$\min \left( \begin{array}{c} \overrightarrow{\text{MIN\_WIDTH}}, \\ \overrightarrow{\text{MIN\_WIDTH}} \end{array} \right)$

Figure 5.566: Glue matrix associated with the automaton of the `min_width_peak` constraint, where  $n$  stands for  $|\text{VARIABLES}|$

`min_width_peak(MIN_WIDTH = 5, VARIABLES = (4, 6, 7, 9, 8, 5, 4))`

	4 6 7 9					9 8 5 4				
	< < <					< < <				
$i$	0	1	2	3	⋮	3	2	1	0	$i$
$\vec{Q}_i$	$s$	$j$	$j$	$j$	⋮	$j$	$j$	$j$	$s$	$\overleftarrow{Q}_i$
$\vec{W}_i$	7	7	7	7	⋮	7	7	7	7	$\overleftarrow{W}_i$
$\vec{C}_i$	0	0	0	0	⋮	0	0	0	0	$\overleftarrow{C}_i$
$\vec{F}_i$	0	1	1	1	⋮	1	1	1	0	$\overleftarrow{F}_i$
$\overrightarrow{\text{MIN\_WIDTH}}_i$	0	0	0	0	⋮	0	0	0	0	$\overleftarrow{\text{MIN\_WIDTH}}_i$

`min_width_peak`  $\left( \begin{array}{c} \overrightarrow{\text{MIN\_WIDTH}}_3 = 0, \\ \langle 4, 6, 7, 9 \rangle \end{array} \right)$  `min_width_peak`  $\left( \begin{array}{c} \overleftarrow{\text{MIN\_WIDTH}}_3 = 0, \\ \langle 4, 5, 8, 9 \rangle \end{array} \right)$

glue matrix entry associated with the state pair  $(j, j)$ :

$$\text{MIN\_WIDTH} = \min(\overrightarrow{W}_3, |\text{VARIABLES}| - \overleftarrow{F}_3 - \overleftarrow{F}_3, \overleftarrow{W}_3) = \min(7, 7 - 1 - 1, 7) = 5$$

Figure 5.567: Illustrating the use of the state pair  $(j, j)$  of the glue matrix for linking `MIN_WIDTH` with the counters variables obtained after reading the prefix 4, 6, 7, 9 and corresponding suffix 9, 8, 5, 4 of the sequence 4, 6, 7, 9, 8, 5, 4; note that the suffix 9, 8, 5, 4 (in pink) is proceed in reverse order; the left (resp. right) table shows the initialisation (for  $i = 0$ ) and the evolution (for  $i > 0$ ) of the state of the automaton and its counters  $W$ ,  $C$  and  $F$  upon reading the prefix 4, 6, 7, 9 (resp. the reverse suffix 4, 5, 8, 9).

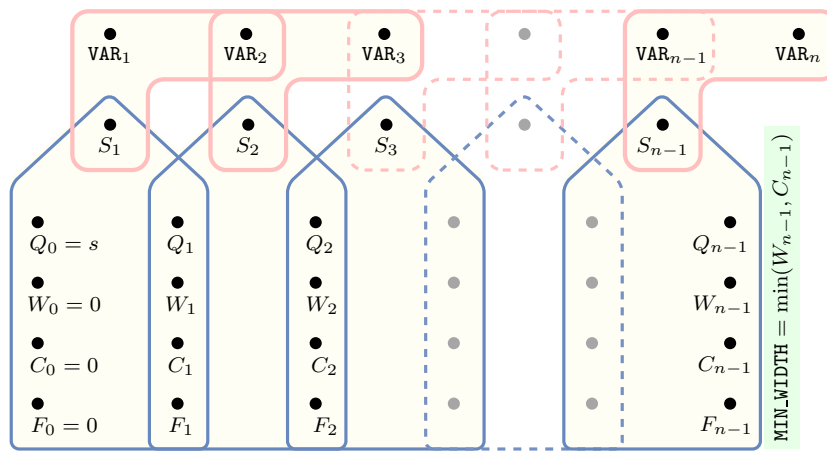


Figure 5.568: Hypergraph of the reformulation corresponding to the automaton of the `min_width_peak` constraint