Vincent Neiger

Laboratoire LIP6, Sorbonne Université

vincent.neiger@lip6.fr

polynomial matrices: approximation and interpolation, quasi-linear GCD

Algorithmes Efficaces en Calcul Formel
Master Parisien de Recherche en Informatique
16 December 2021

outline

- introduction
- shifted reduced forms
- fast algorithms
- next time

outline

- **introduction**
  - rational approximation and interpolation
  - the vector case
  - pol. matrices: reminders and motivation
- **shifted reduced forms**
- **fast algorithms**
- **next time**

introduction

⇓ earlier in the course ⇓

⇓ in this lecture ⇓

introduction

⇓ earlier in the course ⇓

- addition $f+g$, multiplication $f*g$
- division with remainder $f=qg+r$
- truncated inverse $f^{-1} \bmod X^d$
- extended GCD $uf+vg=\gcd(f,g)$
- multipoint eval. $f \mapsto f(\alpha_1),\dots,f(\alpha_d)$
- interpolation $f(\alpha_1),\dots,f(\alpha_d) \mapsto f$
- Padé approximation $f=\frac{p}{q} \bmod X^d$
- minpoly of linearly recurrent sequence

⇓ in this lecture ⇓

introduction

⇓ earlier in the course ⇓

$O(M(d))$

- addition $f + g$, multiplication $f * g$
- division with remainder $f = qg + r$
- truncated inverse $f^{-1} \bmod X^d$
- extended GCD $uf + vg = \gcd(f, g)$

$O(M(d) \log(d))$

- multipoint eval. $f \mapsto f(\alpha_1), \ldots, f(\alpha_d)$
- interpolation $f(\alpha_1), \ldots, f(\alpha_d) \mapsto f$
- Padé approximation $f = \frac{p}{q} \bmod X^d$
- minpoly of linearly recurrent sequence

⇓ in this lecture ⇓

introduction

⇓ earlier in the course ⇓

$O(M(d))$

- addition $f + g$, multiplication $f * g$
- division with remainder $f = qg + r$
- truncated inverse $f^{-1} \bmod X^d$
- extended GCD $uf + vg = \gcd(f, g)$

$O(M(d) \log(d))$

- multipoint eval. $f \mapsto f(\alpha_1), \ldots, f(\alpha_d)$
- interpolation $f(\alpha_1), \ldots, f(\alpha_d) \mapsto f$
- Padé approximation $f = \frac{p}{q} \bmod X^d$
- minpoly of linearly recurrent sequence

⇓ in this lecture ⇓

Padé approximation, sequence minpoly, extended GCD

$O(M(d) \log(d))$ operations in $\mathbb{K}$

matrix versions of these problems

$O(m^\omega M(d) \log(d))$ operations in $\mathbb{K}$

or a tiny bit more for matrix-GCD

rational approximation and interpolation

given power series $p(X)$ and $q(X)$ over $\mathbb{K}$ at precision d,
with $q(X)$ invertible,
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod X^d$ algo?? $O(??)$

introduction

rational approximation and interpolation

given power series $p(X)$ and $q(X)$ over $\mathbb{K}$ at precision d,
with $q(X)$ invertible,
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod X^d$

algo?? $O(??)$
inv+mul: $O(M(d))$

rational approximation and interpolation

given power series $p(X)$ and $q(X)$ over $\mathbb{K}$ at precision d,
with $q(X)$ invertible,
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod X^d$ algo?? $O(??)$
inv+mul: $O(M(d))$

given $M(X) \in \mathbb{K}[X]$ of degree $d > 0$,
given polynomials $p(X)$ and $q(X)$ over $\mathbb{K}$ of degree $< d$,
with $q(X)$ invertible modulo $M(X)$, what does that mean?
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod M(X)$ algo?? $O(??)$

introduction

rational approximation and interpolation

given power series $p(X)$ and $q(X)$ over $\mathbb{K}$ at precision d,
with $q(X)$ invertible,
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod X^d$ algo?? $O(??)$
inv+mul: $O(M(d))$

given $M(X) \in \mathbb{K}[X]$ of degree $d > 0$,
given polynomials $p(X)$ and $q(X)$ over $\mathbb{K}$ of degree $< d$,
with $q(X)$ invertible modulo $M(X)$, what does that mean?
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod M(X)$ algo?? $O(??)$
xgcd+mul+rem $O(M(d) \log(d))$

rational approximation and interpolation

given power series $p(X)$ and $q(X)$ over $\mathbb{K}$ at precision d,
with $q(X)$ invertible,
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod X^d$ algo?? $O(??)$

inv+mul: $O(M(d))$

given $M(X) \in \mathbb{K}[X]$ of degree $d > 0$,
given polynomials $p(X)$ and $q(X)$ over $\mathbb{K}$ of degree $< d$,
with $q(X)$ invertible modulo $M(X)$, what does that mean?
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod M(X)$ algo?? $O(??)$

xgcd+mul+rem $O(M(d)\log(d))$

given $M(X) = (X-\alpha_1)\cdots(X-\alpha_d) \in \mathbb{K}[X]$,
for pairwise distinct $\alpha_1, \ldots, \alpha_d \in \mathbb{K}$,
given polynomials $p(X)$ and $q(X)$ over $\mathbb{K}$ of degree $< d$,
with $q(X)$ invertible modulo $M(X)$, what does that mean?
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod M(X)$ algo?? $O(??)$

introduction

rational approximation and interpolation

given power series $p(X)$ and $q(X)$ over $\mathbb{K}$ at precision d,
with $q(X)$ invertible,
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod X^d$ algo?? $O(??)$
inv+mul: $O(M(d))$

given $M(X) \in \mathbb{K}[X]$ of degree $d > 0$,
given polynomials $p(X)$ and $q(X)$ over $\mathbb{K}$ of degree $< d$,
with $q(X)$ invertible modulo $M(X)$, what does that mean?
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod M(X)$ algo?? $O(??)$
xgcd+mul+rem $O(M(d)\log(d))$

given $M(X) = (X - \alpha_1)\cdots(X - \alpha_d) \in \mathbb{K}[X]$,
for pairwise distinct $\alpha_1, \ldots, \alpha_d \in \mathbb{K}$,
given polynomials $p(X)$ and $q(X)$ over $\mathbb{K}$ of degree $< d$,
with $q(X)$ invertible modulo $M(X)$, what does that mean?
$\rightarrow$ compute $\frac{p(X)}{q(X)} \bmod M(X)$ algo?? $O(??)$
eval+div+interp $O(M(d)\log(d))$

linearly recurrent sequences – reminder from October 21

Generating series of LRS and rational functions

Theorem

Given a monic polynomial P of degree d, a sequence $(a_n)_{n\in\mathbb{N}}$, and the series $A = \sum_{n\in\mathbb{N}} a_n X^n$, both following assertions are equivalent:

1. $(a_n)_{n\in\mathbb{N}}$ is an LRS with characteristic polynomial P;
2. there exists $N \in \mathbb{K}[X]$ of degree $< d$ such that $A = N/\operatorname{rec} P$ in $\mathbb{K}[[X]]$.

When these assertions hold, if moreover P is the minimal polynomial of $(a_n)_{n\in\mathbb{N}}$, then

$$d = \max\{1 + \deg N, \deg \operatorname{rec} P\} := m \quad \text{and} \quad \gcd(N, \operatorname{rec} P) = 1.$$

linearly recurrent sequences – reminder from October 21

Generating series of LRS and rational functions

Theorem

Given a monic polynomial P of degree d, a sequence $(a_n)_{n\in\mathbb{N}}$, and the series $A = \sum_{n\in\mathbb{N}} a_n X^n$, both following assertions are equivalent:

1. $(a_n)_{n\in\mathbb{N}}$ is an LRS with characteristic polynomial P;
2. there exists $N \in \mathbb{K}[X]$ of degree $< d$ such that $A = N/\operatorname{rec} P$ in $\mathbb{K}[[X]]$.

When these assertions hold, if moreover P is the minimal polynomial of $(a_n)_{n\in\mathbb{N}}$, then

$$d = \max\{1 + \deg N, \deg \operatorname{rec} P\} := m \quad \text{and} \quad \gcd(N, \operatorname{rec} P) = 1.$$

numerator N and charpoly P $\xrightarrow{\text{expand } \frac{N}{\operatorname{rev}(P)} \bmod X^\delta}$ first δ terms of the LRS $(a_n)_{n\in\mathbb{N}}$

linearly recurrent sequences – reminder from October 21

Generating series of LRS and rational functions

Theorem

Given a monic polynomial P of degree d, a sequence $(a_n)_{n\in\mathbb{N}}$, and the series $A = \sum_{n\in\mathbb{N}} a_n X^n$, both following assertions are equivalent:

1. $(a_n)_{n\in\mathbb{N}}$ is an LRS with characteristic polynomial P;
2. there exists $N \in \mathbb{K}[X]$ of degree $< d$ such that $A = N/\operatorname{rec} P$ in $\mathbb{K}[[X]]$.

When these assertions hold, if moreover P is the minimal polynomial of $(a_n)_{n\in\mathbb{N}}$, then

$$d = \max\{1 + \deg N, \deg \operatorname{rec} P\} := m \quad \text{and} \quad \gcd(N, \operatorname{rec} P) = 1.$$

expand $\frac{N}{\operatorname{rev}(P)} \bmod X^\delta$

numerator N and charpoly P $\longrightarrow$ first δ terms of the LRS $(a_n)_{n\in\mathbb{N}}$

reconstruct from $A(X) \bmod X^\delta$ $\rightsquigarrow$ **Padé approximation**

Padé approximation:

given power series $f(X)$ at precision d,
$\rightarrow$ compute $p(X), q(X)$ such that $f = \frac{p}{q} \bmod X^d$

Padé approximation:

given power series $f(X)$ at precision d,
$\rightarrow$ compute $p(X), q(X)$ such that $f = \frac{p}{q} \bmod X^d$

opinions on this algorithmic problem?

introduction

rational approximation and interpolation

Padé approximation:

given power series $f(X)$ at precision d,
given degree constraints $d_1, d_2 > 0$,
$\rightarrow$ compute polynomials $(p(X), q(X))$ of degrees $< (d_1, d_2)$
and such that $f = \frac{p}{q} \bmod X^d$

Padé approximation:

given power series $f(X)$ at precision d,
given degree constraints $d_1, d_2 > 0$,
$\rightarrow$ compute polynomials $(p(X), q(X))$ of degrees $< (d_1, d_2)$
and such that $f = \frac{p}{q} \bmod X^d$

Cauchy interpolation:

given $M(X) = (X - \alpha_1) \cdots (X - \alpha_d) \in \mathbb{K}[X]$,
for pairwise distinct $\alpha_1, \ldots, \alpha_d \in \mathbb{K}$,
given degree constraints $d_1, d_2 > 0$,
$\rightarrow$ compute polynomials $(p(X), q(X))$ of degrees $< (d_1, d_2)$
and such that $f = \frac{p}{q} \bmod M(X)$

introduction

rational approximation and interpolation

Padé approximation:

given power series $f(X)$ at precision d,
given degree constraints $d_1, d_2 > 0$,
$\rightarrow$ compute polynomials $(p(X), q(X))$ of degrees $< (d_1, d_2)$
and such that $f = \frac{p}{q} \bmod X^d$

Cauchy interpolation:

given $M(X) = (X - \alpha_1) \cdots (X - \alpha_d) \in \mathbb{K}[X]$,
for pairwise distinct $\alpha_1, \ldots, \alpha_d \in \mathbb{K}$,
given degree constraints $d_1, d_2 > 0$,
$\rightarrow$ compute polynomials $(p(X), q(X))$ of degrees $< (d_1, d_2)$
and such that $f = \frac{p}{q} \bmod M(X)$

- degree constraints specified by the context
- usual choices have $d_1 + d_2 \approx d$ and existence of a solution

introduction

approximation and structured linear system

$\mathbb{K} = \mathbb{F}_7$
$f = 2X^7 + 2X^6 + 5X^4 + 2X^2 + 4$
$d = 8$, $d_1 = 3$, $d_2 = 6$
$\rightarrow$ look for (p, q) of degree $< (3, 6)$ such that $f = \frac{p}{q} \bmod X^8$

$$\begin{bmatrix} q & p \end{bmatrix} \begin{bmatrix} f \\ -1 \end{bmatrix} = 0 \bmod X^8$$

introduction

approximation and structured linear system

$\mathbb{K} = \mathbb{F}_7$
$f = 2X^7 + 2X^6 + 5X^4 + 2X^2 + 4$
$d = 8,\ d_1 = 3,\ d_2 = 6$
$\rightarrow$ look for (p, q) of degree $< (3, 6)$ such that $f = \frac{p}{q} \bmod X^8$

$$\begin{bmatrix} q & p \end{bmatrix} \begin{bmatrix} f \\ -1 \end{bmatrix} = 0 \bmod X^8$$

$$\left[\begin{array}{cccccc|ccc} q_0 & q_1 & q_2 & q_3 & q_4 & 1 & p_0 & p_1 & p_2 \end{array}\right] \left[\begin{array}{cccccccc} 4 & 0 & 2 & 0 & 5 & 0 & 2 & 2 \\ & 4 & 0 & 2 & 0 & 5 & 0 & 2 \\ & & 4 & 0 & 2 & 0 & 5 & 0 \\ & & & 4 & 0 & 2 & 0 & 5 \\ & & & & 4 & 0 & 2 & 0 \\ & & & & & 4 & 0 & 2 \\ \hdashline 6 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ & 6 & 0 & 0 & 0 & 0 & 0 & 0 \\ & & 6 & 0 & 0 & 0 & 0 & 0 \end{array}\right] = 0$$

introduction

approximation and structured linear system

$\mathbb{K} = \mathbb{F}_7$
$f = 2X^7 + 2X^6 + 5X^4 + 2X^2 + 4$
$d = 8$, $d_1 = 3$, $d_2 = 6$
$\rightarrow$ look for (p, q) of degree $< (3, 6)$ such that $f = \frac{p}{q} \bmod X^8$

$$\begin{bmatrix} q & p \end{bmatrix} \begin{bmatrix} f \\ -1 \end{bmatrix} = 0 \bmod X^8$$

$$[q_0 \;\; q_1 \;\; q_2 \;\; q_3 \;\; q_4 \;\; 1 \mid p_0 \;\; p_1 \;\; p_2] \left[\begin{array}{cccccccc} 4 & 0 & 2 & 0 & 5 & 0 & 2 & 2 \\ & 4 & 0 & 2 & 0 & 5 & 0 & 2 \\ & & 4 & 0 & 2 & 0 & 5 & 0 \\ & & & 4 & 0 & 2 & 0 & 5 \\ & & & & 4 & 0 & 2 & 0 \\ & & & & & 4 & 0 & 2 \\ \hdashline 6 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ & 6 & 0 & 0 & 0 & 0 & 0 & 0 \\ & & 6 & 0 & 0 & 0 & 0 & 0 \end{array}\right] = 0$$

Sur la généralisation des fractions continues algébriques;

Par M. H. PADÉ,

Docteur ès Sciences mathématiques,
Professeur au lycée de Lille.

[**1894**, Journal de mathématiques pures et appliquées]

INTRODUCTION.

M. Hermite s'est, dans un travail récemment paru (¹), occupé de la généralisation des fractions continues algébriques. La question est de déterminer les polynomes $X_1, X_2, \ldots, X_n$, de degrés $\mu_1, \mu_2, \ldots, \mu_n$, qui satisfont à l'équation

$$S_1 X_1 + S_2 X_2 + \ldots + S_n X_n = S\, x^{\mu_1+\mu_2+\ldots+\mu_n+n-1},$$

$S_1, S_2, \ldots, S_n$ étant des séries entières données, et S une série également entière. Ou plutôt, il s'agit d'obtenir un algorithme qui permette le calcul de proche en proche de ces systèmes de n polynomes, et qui soit analogue à l'algorithme par lequel le numérateur et le dénominateur d'une réduite d'une fraction continue se déduisent des numérateurs et dénominateurs des réduites précédentes. D'élégantes considé-

approximation and interpolation: the vector case

Hermite-Padé approximation

[Hermite 1893, Padé 1894]

input:

- polynomials $f_1, \ldots, f_m \in \mathbb{K}[X]$
- precision $d \in \mathbb{Z}_{>0}$
- degree bounds $d_1, \ldots, d_m \in \mathbb{Z}_{>0}$

output:

polynomials $p_1, \ldots, p_m \in \mathbb{K}[X]$ such that

- $p_1 f_1 + \cdots + p_m f_m = 0 \bmod X^d$
- $\mathrm{cdeg}([p_1 \cdots p_m]) < (d_1, \ldots, d_m)$

(Padé approximation: particular case $m = 2$ and $f_2 = -1$)

approximation and interpolation: the vector case

M-Padé approximation / vector rational interpolation

[Cauchy 1821, Mahler 1968]

input:

- polynomials $f_1, \ldots, f_m \in \mathbb{K}[X]$
- pairwise distinct points $\alpha_1, \ldots, \alpha_d \in \mathbb{K}$
- degree bounds $d_1, \ldots, d_m \in \mathbb{Z}_{>0}$

output:

polynomials $p_1, \ldots, p_m \in \mathbb{K}[X]$ such that

- $p_1(\alpha_i)f_1(\alpha_i) + \cdots + p_m(\alpha_i)f_m(\alpha_i) = 0$ for all $1 \leqslant i \leqslant d$
- $\mathrm{cdeg}([p_1 \cdots p_m]) < (d_1, \ldots, d_m)$

(rational interpolation: particular case $m = 2$ and $f_2 = -1$)

approximation and interpolation: the vector case

in this lecture: modular equation and fast algebraic algorithms

[van Barel-Bultheel 1992; Beckermann-Labahn 1994, 1997, 2000; Giorgi-Jeannerod-Villard 2003; Storjohann 2006; Zhou-Labahn 2012; Jeannerod-Neiger-Schost-Villard 2017, 2020]

input:

- polynomials $f_1, \ldots, f_m \in \mathbb{K}[X]$
- field elements $\alpha_1, \ldots, \alpha_d \in \mathbb{K}$ $\rightsquigarrow$ not necessarily distinct
- degree bounds $d_1, \ldots, d_m \in \mathbb{Z}_{>0}$ $\rightsquigarrow$ general "shift" $\mathbf{s} \in \mathbb{Z}^m$

output:

polynomials $p_1, \ldots, p_m \in \mathbb{K}[X]$ such that

- $p_1 f_1 + \cdots + p_m f_m = 0 \bmod \prod_{1 \leqslant i \leqslant d}(X - \alpha_i)$
- $\mathrm{cdeg}([p_1 \cdots p_m]) < (d_1, \ldots, d_m)$ $\rightsquigarrow$ minimal $\mathbf{s}$-row degree

(Hermite-Padé: $\alpha_1 = \cdots = \alpha_d = 0$; interpolation: pairwise distinct points)

interpolation and structured linear system

application of vector rational interpolation:
given pairwise distinct points $\{(\alpha_i, \beta_i), 1 \leqslant i \leqslant 8\}$
$= \{(24, 80), (31, 73), (15, 73), (32, 35), (83, 66), (27, 46), (20, 91), (59, 64)\}$,
compute a bivariate polynomial $p(X, Y) \in \mathbb{K}[X, Y]$
such that $p(\alpha_i, \beta_i) = 0$ for $1 \leqslant i \leqslant 8$

$$\left.\begin{array}{l} M(X) = (X-24)\cdots(X-59) \\ L(X) = \text{Lagrange interpolant} \end{array}\right\} \longrightarrow \text{solutions} = \text{ideal } \langle M(X), Y - L(X)\rangle$$

solutions of smaller X-degree: $p(X, Y) = p_0(X) + p_1(X)Y + p_2(X)Y^2$

$$p(X, L(X)) = \begin{bmatrix} p_0 & p_1 & p_2 \end{bmatrix} \begin{bmatrix} 1 \\ L \\ L^2 \end{bmatrix} = 0 \bmod M(X)$$

- instance of univariate rational vector interpolation
- with a structured input equation (powers of L mod M)

interpolation and structured linear system

application of vector rational interpolation:
given pairwise distinct points $\{(\alpha_i, \beta_i), 1 \leqslant i \leqslant 8\}$
$= \{(24, 80), (31, 73), (15, 73), (32, 35), (83, 66), (27, 46), (20, 91), (59, 64)\}$,
compute a bivariate polynomial $p(X, Y) \in \mathbb{K}[X, Y]$
such that $p(\alpha_i, \beta_i) = 0$ for $1 \leqslant i \leqslant 8$

add degree constraints: seek $p(X, Y)$ of the form
$p_{00} + p_{01}X + p_{02}X^2 + p_{03}X^3 + p_{04}X^4 + (p_{10} + p_{11}X + p_{12}X^2)Y + p_{20}Y^2$:

$$\left[\begin{array}{ccccc:ccc:c} p_{00} & p_{01} & p_{02} & p_{03} & p_{04} & p_{10} & p_{11} & p_{12} & p_{20} \end{array}\right] \left[\begin{array}{cccc} 1 & 1 & \cdots & 1 \\ \alpha_1 & \alpha_2 & \cdots & \alpha_8 \\ \alpha_1^2 & \alpha_2^2 & \cdots & \alpha_8^2 \\ \alpha_1^3 & \alpha_2^3 & \cdots & \alpha_8^3 \\ \alpha_1^4 & \alpha_2^4 & \cdots & \alpha_8^4 \\ \hdashline \beta_1 & \beta_2 & \cdots & \beta_8 \\ \alpha_1\beta_1 & \alpha_2\beta_2 & \cdots & \alpha_8\beta_8 \\ \alpha_1^2\beta_1 & \alpha_2^2\beta_2 & \cdots & \alpha_8^2\beta_8 \\ \hdashline \beta_1^2 & \beta_2^2 & \cdots & \beta_8^2 \end{array}\right] = 0$$

- $\mathbb{K}$-linear system
- two levels of structure

$$p(X, Y) = (2X^4 + 56X^3 + 42X^2 + 48X + 15) + (72X^2 + 12X + 30)Y + Y^2$$

polynomial matrices: reminder and motivation

why polynomial matrices here?

why polynomial matrices here?

omitting degree constraints, the set of solutions is
$\mathcal{S} = \{(p_1, \ldots, p_m) \in \mathbb{K}[X]^m \mid p_1 f_1 + \cdots + p_m f_m = 0 \bmod M\}$

recall $M(X) = \prod_{1 \leqslant i \leqslant d}(X - \alpha_i)$

polynomial matrices: reminder and motivation

why polynomial matrices here?

omitting degree constraints, the set of solutions is
$\mathcal{S} = \{(p_1, \ldots, p_m) \in \mathbb{K}[X]^m \mid p_1 f_1 + \cdots + p_m f_m = 0 \bmod M\}$

recall $M(X) = \prod_{1 \leqslant i \leqslant d} (X - \alpha_i)$

$\mathcal{S}$ is a "free $\mathbb{K}[X]$-module of rank m", meaning:

- stable under $\mathbb{K}[X]$-linear combinations
- admits a basis consisting of m elements
- basis = $\mathbb{K}[X]$-linear independence + generates all solutions

polynomial matrices: reminder and motivation

why polynomial matrices here?

omitting degree constraints, the set of solutions is
$\mathcal{S} = \{(p_1, \ldots, p_m) \in \mathbb{K}[X]^m \mid p_1 f_1 + \cdots + p_m f_m = 0 \bmod M\}$

recall $M(X) = \prod_{1 \leqslant i \leqslant d} (X - \alpha_i)$

$\mathcal{S}$ is a "free $\mathbb{K}[X]$-module of rank m", meaning:

- stable under $\mathbb{K}[X]$-linear combinations
- admits a basis consisting of m elements
- basis $=$ $\mathbb{K}[X]$-linear independence $+$ generates all solutions

- $\mathcal{S} \subset \mathbb{K}[X]^m \quad \Rightarrow \quad \mathcal{S}$ has rank $\leqslant m$
- $M(X)\mathbb{K}[X]^m \subset \mathcal{S} \quad \Rightarrow \quad \mathcal{S}$ has rank $\geqslant m$

remark: solutions are not considered modulo M
e.g. $(M, 0, \ldots, 0)$ is in $\mathcal{S}$ and may appear in a basis

polynomial matrices: reminder and motivation

why polynomial matrices here?

omitting degree constraints, the set of solutions is
$\mathcal{S} = \{(p_1, \ldots, p_m) \in \mathbb{K}[X]^m \mid p_1 f_1 + \cdots + p_m f_m = 0 \bmod M\}$

recall $M(X) = \prod_{1 \leqslant i \leqslant d}(X - \alpha_i)$

basis of solutions:

- square nonsingular matrix $\mathbf{P}$ in $\mathbb{K}[X]^{m \times m}$
- each row of $\mathbf{P}$ is a solution
- any solution is a $\mathbb{K}[X]$-combination $\mathbf{uP}, \mathbf{u} \in \mathbb{K}[X]^{1 \times m}$

i.e. $\mathcal{S}$ is the $\mathbb{K}[X]$-row space of $\mathbf{P}$

polynomial matrices: reminder and motivation

why polynomial matrices here?

omitting degree constraints, the set of solutions is
$\mathcal{S} = \{(p_1, \ldots, p_m) \in \mathbb{K}[X]^m \mid p_1 f_1 + \cdots + p_m f_m = 0 \bmod M\}$

recall $M(X) = \prod_{1 \leqslant i \leqslant d} (X - \alpha_i)$

basis of solutions:

- square nonsingular matrix $\mathbf{P}$ in $\mathbb{K}[X]^{m \times m}$
- each row of $\mathbf{P}$ is a solution
- any solution is a $\mathbb{K}[X]$-combination $\mathbf{uP}, \mathbf{u} \in \mathbb{K}[X]^{1 \times m}$

i.e. $\mathcal{S}$ is the $\mathbb{K}[X]$-row space of $\mathbf{P}$

prove: $\det(\mathbf{P})$ **is a divisor of** $M(X)^m$

polynomial matrices: reminder and motivation

why polynomial matrices here?

omitting degree constraints, the set of solutions is
$\mathcal{S} = \{(p_1, \ldots, p_m) \in \mathbb{K}[X]^m \mid p_1 f_1 + \cdots + p_m f_m = 0 \bmod M\}$

recall $M(X) = \prod_{1 \leqslant i \leqslant d}(X - \alpha_i)$

basis of solutions:

- square nonsingular matrix $\mathbf{P}$ in $\mathbb{K}[X]^{m \times m}$
- each row of $\mathbf{P}$ is a solution
- any solution is a $\mathbb{K}[X]$-combination $\mathbf{uP}$, $\mathbf{u} \in \mathbb{K}[X]^{1 \times m}$

i.e. $\mathcal{S}$ is the $\mathbb{K}[X]$-row space of $\mathbf{P}$

prove: $\det(\mathbf{P})$ **is a divisor of** $M(X)^m$

prove: any other basis is $\mathbf{UP}$ **for** $\mathbf{U} \in \mathbb{K}[X]^{m \times m}$ **with** $\det(\mathbf{U}) \in \mathbb{K} \setminus \{0\}$

polynomial matrices: reminder and motivation

why polynomial matrices here?

omitting degree constraints, the set of solutions is
$\mathcal{S} = \{(p_1, \ldots, p_m) \in \mathbb{K}[X]^m \mid p_1 f_1 + \cdots + p_m f_m = 0 \bmod M\}$

recall $M(X) = \prod_{1 \leqslant i \leqslant d}(X - \alpha_i)$

basis of solutions:

- square nonsingular matrix $\mathbf{P}$ in $\mathbb{K}[X]^{m \times m}$
- each row of $\mathbf{P}$ is a solution
- any solution is a $\mathbb{K}[X]$-combination $\mathbf{uP}$, $\mathbf{u} \in \mathbb{K}[X]^{1 \times m}$

i.e. $\mathcal{S}$ is the $\mathbb{K}[X]$-row space of $\mathbf{P}$

computing a basis of $\mathcal{S}$ with "minimal degrees"

- has many more applications than a single small-degree solution
- is in most cases the fastest known strategy anyway(!)

$\rightsquigarrow$ degree minimality ensured via shifted reduced forms

polynomial matrices: reminder and motivation

$$\mathbf{A} = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \\ 3X^3+X^2+5X+3 & 6X+5 & 2X+1 \end{bmatrix} \in \mathbb{K}[X]^{3\times 3}$$

3×3 matrix of degree 3
with entries in $\mathbb{K}[X] = \mathbb{F}_7[X]$

operations in $\mathbb{K}[X]^{m\times m}_{<d}$:

- combination of matrix and polynomial computations
- addition in $O(m^2 d)$, naive multiplication in $O(m^3 d^2)$
- some tools shared with $\mathbb{K}$-matrices, others specific to $\mathbb{K}[X]$-matrices

[Cantor-Kaltofen'91]

multiplication in $O(m^\omega d \log(d) \; + \; m^2 d \log(d) \log\log(d))$

$\in O(m^\omega \mathsf{M}(d)) \subset O\tilde{~}(m^\omega d)$

introduction

polynomial matrices: reminder and motivation

$$\mathbf{A} = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \\ 3X^3+X^2+5X+3 & 6X+5 & 2X+1 \end{bmatrix} \in \mathbb{K}[X]^{3\times 3}$$

3×3 matrix of degree 3
with entries in $\mathbb{K}[X] = \mathbb{F}_7[X]$

operations in $\mathbb{K}[X]^{m\times m}_{<d}$:

- combination of matrix and polynomial computations
- addition in $O(m^2 d)$, naive multiplication in $O(m^3 d^2)$
- some tools shared with $\mathbb{K}$-matrices, others specific to $\mathbb{K}[X]$-matrices

[Cantor-Kaltofen'91]

multiplication in $O(m^\omega d \log(d) \; + \; m^2 d \log(d) \log\log(d))$

$\in O(m^\omega M(d)) \subset O\tilde{}(m^\omega d)$

- Newton truncated inversion, matrix-QuoRem → fast $O\tilde{}(m^\omega d)$
- inversion and determinant via evaluation-interpolation → medium $O\tilde{}(m^{\omega+1} d)$
- vector rational approximation & interpolation → **???**

polynomial matrices: reminder and motivation

reductions of most problems to polynomial matrix multiplication

matrix $m \times m$ of degree d $\rightarrow$ $O\tilde{}(m^\omega d)$

of "average" degree $\frac{D}{m}$ $\rightarrow$ $O\tilde{}(m^\omega \frac{D}{m})$

classical matrix operations

- multiplication
- kernel, system solving
- rank, determinant
- inversion $O\tilde{}(m^3 d)$

univariate specific operations

- truncated inverse, QuoRem
- Hermite-Padé approximation
- vector rational interpolation
- syzygies / modular equations

transformation to **normal forms**

- triangularization: Hermite form
- row reduction: Popov form
- diagonalization: Smith form

reductions of most problems to polynomial matrix multiplication

matrix $m \times m$ of degree d $\quad\rightarrow\quad O\tilde{~}(m^\omega d)$

of "average" degree $\frac{D}{m}$ $\quad\rightarrow\quad O\tilde{~}(m^\omega \frac{D}{m})$

classical matrix operations

- multiplication
- kernel, system solving
- rank, determinant
- inversion $\quad O\tilde{~}(m^3 d)$

univariate specific operations

- truncated inverse, QuoRem
- Hermite-Padé approximation
- vector rational interpolation
- syzygies / modular equations

transformation to **normal forms**

- triangularization: Hermite form
- row reduction: Popov form
- diagonalization: Smith form

polynomial matrices: reminder and motivation

reductions of most problems to polynomial matrix multiplication

matrix $m \times m$ of degree d $\rightarrow$ $O\tilde{~}(m^\omega d)$

of "average" degree $\frac{D}{m}$ $\rightarrow$ $O\tilde{~}(m^\omega \frac{D}{m})$

classical matrix operations

- multiplication
- kernel, system solving
- rank, determinant
- inversion $O\tilde{~}(m^3 d)$

univariate specific operations

- truncated inverse, QuoRem
- Hermite-Padé approximation
- vector rational interpolation
- syzygies / modular equations

transformation to **normal forms**

- triangularization: Hermite form
- row reduction: Popov form
- diagonalization: Smith form

polynomial matrices: reminder and motivation

reductions of most problems to polynomial matrix multiplication

matrix $m \times m$ of degree d $\quad\rightarrow\quad$ $O\tilde{~}(m^\omega d)$

of "average" degree $\frac{D}{m}$ $\quad\rightarrow\quad$ $O\tilde{~}(m^\omega \frac{D}{m})$

classical matrix operations

- multiplication
- kernel, system solving
- rank, determinant
- inversion $\quad O\tilde{~}(m^3 d)$

univariate specific operations

- truncated inverse, QuoRem
- Hermite-Padé approximation
- vector rational interpolation
- syzygies / modular equations

transformation to **normal forms**

- triangularization: Hermite form
- row reduction: Popov form
- diagonalization: Smith form

outline

- **introduction**
  - ▸ rational approximation and interpolation
  - ▸ the vector case
  - ▸ pol. matrices: reminders and motivation
- **shifted reduced forms**
- **fast algorithms**
- **next time**

outline

- **introduction**
  - rational approximation and interpolation
  - the vector case
  - pol. matrices: reminders and motivation
- **shifted reduced forms**
  - reducedness: examples and properties
  - shifted forms and degree constraints
  - stability under multiplication
- **fast algorithms**
- **next time**

notation:

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with no zero row,
define $\mathbf{d} = (d_1, \ldots, d_m) = \mathrm{rdeg}(\mathbf{A})$

and $\mathbf{X^d} = \begin{bmatrix} X^{d_1} & & \\ & \ddots & \\ & & X^{d_m} \end{bmatrix} \in \mathbb{K}[X]^{m \times m}$

definition: (row-wise) leading matrix

the leading matrix of $\mathbf{A}$ is the unique matrix $\mathrm{lm}(\mathbf{A}) \in \mathbb{K}^{m \times n}$ such that $\mathbf{A} = \mathbf{X^d}\mathrm{lm}(\mathbf{A}) + \mathbf{R}$ with $\mathrm{rdeg}(\mathbf{R}) < \mathbf{d}$ entry-wise

equivalently, $\mathbf{X^{-d}A} = \mathrm{lm}(\mathbf{A}) +$ terms of strictly negative degree

notation:

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with no zero row,
define $\mathbf{d} = (d_1, \ldots, d_m) = \mathrm{rdeg}(\mathbf{A})$

and $\mathbf{X}^{\mathbf{d}} = \begin{bmatrix} X^{d_1} & & \\ & \ddots & \\ & & X^{d_m} \end{bmatrix} \in \mathbb{K}[X]^{m \times m}$

definition: (row-wise) leading matrix

the leading matrix of $\mathbf{A}$ is the unique matrix $\mathrm{lm}(\mathbf{A}) \in \mathbb{K}^{m \times n}$ such that $\mathbf{A} = \mathbf{X}^{\mathbf{d}}\mathrm{lm}(\mathbf{A}) + \mathbf{R}$ with $\mathrm{rdeg}(\mathbf{R}) < \mathbf{d}$ entry-wise

equivalently, $\mathbf{X}^{-\mathbf{d}}\mathbf{A} = \mathrm{lm}(\mathbf{A})$ + terms of strictly negative degree

definition: (row-wise) reduced matrix

$\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ is said to be reduced
if $\mathrm{lm}(\mathbf{A})$ has full row rank

shifted reduced forms

reducedness: examples and properties

consider the following matrices, with $\mathbb{K} = \mathbb{F}_7$:

$$\mathbf{A}_1 = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \end{bmatrix}$$

$$\mathbf{A}_2 = \begin{bmatrix} 3X+1 & 4X+3 & 5X+5 \\ 0 & 4X^2+6X & 5 \\ 4X^2+5X+2 & 5 & 6X^2+1 \end{bmatrix}$$

$\mathbf{A}_3 =$ transpose of $\mathbf{A}_1$

$\mathbf{A}_4 =$ transpose of $\mathbf{A}_2$

answer the following, for $i \in \{1, 2, 3, 4\}$:
1. what is $\mathrm{rdeg}(\mathbf{A}_i)$?
2. what is $\mathrm{lm}(\mathbf{A}_i)$?
3. is $\mathbf{A}_i$ reduced?

polynomial matrices in reduced form

reducedness: examples and properties

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with $m \leqslant n$,
the following are equivalent:

(i) $\mathbf{A}$ is reduced (i.e. $\mathsf{lm}(\mathbf{A})$ has full rank)

polynomial matrices in reduced form

reducedness: examples and properties

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with $m \leqslant n$,
the following are equivalent:

(i) $\mathbf{A}$ is reduced (i.e. $\mathrm{lm}(\mathbf{A})$ has full rank)

(ii) for any vector $\mathbf{u} = [\mathbf{u}_1 \;\; 1 \;\; \mathbf{u}_2] \in \mathbb{K}[X]^{1 \times m}$ with 1 at index i,
$\mathrm{rdeg}(\mathbf{uA}) \geqslant \mathrm{rdeg}(\mathbf{A}_{i,*})$

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with $m \leqslant n$,
the following are equivalent:

(i) $\mathbf{A}$ is reduced (i.e. $\mathrm{lm}(\mathbf{A})$ has full rank)

(ii) for any vector $\mathbf{u} = [\mathbf{u}_1 \;\; 1 \;\; \mathbf{u}_2] \in \mathbb{K}[X]^{1 \times m}$ with 1 at index i,
$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) \geqslant \mathrm{rdeg}(\mathbf{A}_{i,*})$

(iii) predictable degree: for any vector $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1 \times m}$,
$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \max_{1 \leqslant i \leqslant m}(\deg(u_i) + \mathrm{rdeg}(\mathbf{A}_{i,*}))$

polynomial matrices in reduced form

reducedness: examples and properties

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with $m \leqslant n$,
the following are equivalent:

(i) $\mathbf{A}$ is reduced (i.e. $\mathrm{lm}(\mathbf{A})$ has full rank)

(ii) for any vector $\mathbf{u} = [\mathbf{u}_1 \;\; 1 \;\; \mathbf{u}_2] \in \mathbb{K}[X]^{1 \times m}$ with 1 at index i,
$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) \geqslant \mathrm{rdeg}(\mathbf{A}_{i,*})$

(iii) predictable degree: for any vector $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1 \times m}$,
$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \max_{1 \leqslant i \leqslant m}(\deg(u_i) + \mathrm{rdeg}(\mathbf{A}_{i,*}))$

(iv) degree minimality: $\mathrm{rdeg}(\mathbf{A}) \preccurlyeq \mathrm{rdeg}(\mathbf{U}\mathbf{A})$ holds for any nonsingular matrix $\mathbf{U} \in \mathbb{K}[X]^{m \times m}$, where $\preccurlyeq$ sorts the tuples in nondecreasing order and then uses lexicographic comparison

polynomial matrices in reduced form

reducedness: examples and properties

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with $m \leqslant n$,
the following are equivalent:

(i) $\mathbf{A}$ is reduced (i.e. $\mathrm{lm}(\mathbf{A})$ has full rank)

(ii) for any vector $\mathbf{u} = [\mathbf{u}_1 \;\; 1 \;\; \mathbf{u}_2] \in \mathbb{K}[X]^{1 \times m}$ with 1 at index i,
$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) \geqslant \mathrm{rdeg}(\mathbf{A}_{i,*})$

(iii) predictable degree: for any vector $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1 \times m}$,
$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \max_{1 \leqslant i \leqslant m}(\deg(u_i) + \mathrm{rdeg}(\mathbf{A}_{i,*}))$

(iv) degree minimality: $\mathrm{rdeg}(\mathbf{A}) \preccurlyeq \mathrm{rdeg}(\mathbf{U}\mathbf{A})$ holds for any nonsingular matrix $\mathbf{U} \in \mathbb{K}[X]^{m \times m}$, where $\preccurlyeq$ sorts the tuples in nondecreasing order and then uses lexicographic comparison

(v) predictable determinantal degree: $\deg\det(\mathbf{A}) = |\mathrm{rdeg}(\mathbf{A})|$
(only when $m = n$)

shifted reduced forms

reducedness: examples and properties

recall the matrix, with $\mathbb{K} = \mathbb{F}_7$,

$$\mathbf{A} = \begin{bmatrix} 3X+1 & 4X+3 & 5X+5 \\ 0 & 4X^2+6X & 5 \\ 4X^2+5X+2 & 5 & 6X^2+1 \end{bmatrix}$$

1. what is $\deg\det(\mathbf{A})$?

2. what is $\mathrm{rdeg}([4X^2+1 \quad 2X \quad 4X+5]\,\mathbf{A})$?

3. is it possible to find a matrix
$$\mathbf{P} = \begin{bmatrix} p_{00} & p_{01} & p_{02} \\ p_{10} & p_{11} & p_{12} \end{bmatrix}$$
whose rank is 2, whose degree is 1, and which is a left-multiple of $\mathbf{A}$?

shifted reduced forms

reducedness: examples and properties

recall the matrix, with $\mathbb{K} = \mathbb{F}_7$,

$$\mathbf{A} = \begin{bmatrix} 3X+1 & 4X+3 & 5X+5 \\ 0 & 4X^2+6X & 5 \\ 4X^2+5X+2 & 5 & 6X^2+1 \end{bmatrix}$$

1. what is $\deg \det(\mathbf{A})$?

2. what is $\mathrm{rdeg}([4X^2+1 \quad 2X \quad 4X+5]\,\mathbf{A})$?

3. is it possible to find a matrix

$$\mathbf{P} = \begin{bmatrix} p_{00} & p_{01} & p_{02} \\ p_{10} & p_{11} & p_{12} \end{bmatrix}$$

whose rank is 2, whose degree is 1, and which is a left-multiple of $\mathbf{A}$?

find a row vector $\mathbf{u}$ of degree 1 such that $\mathbf{uA}$ has degree 2, where

$$\mathbf{A} = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \end{bmatrix}$$

keeping our problem in mind:

- input: f_i's and α_i's and degree constraints $d_1, \ldots, d_m \in \mathbb{Z}_{>0}$
- output: a solution $\mathbf{p}$ satisfying the constraints $\mathrm{cdeg}(\mathbf{p}) < (d_1, \ldots, d_m)$

obstacle:
computing a reduced basis of solutions ignores the constraints

exercice: suppose we have a reduced basis $\mathbf{P} \in \mathbb{K}[X]^{m \times m}$ of solutions

- think of particular constraints $(d_1, \ldots, d_m)$ that can be handled via $\mathbf{P}$
- give constraints $(d_1, \ldots, d_m)$ for which $\mathbf{P}$ is "typically" not satisfactory

shifted reduced forms

shifted forms and degree constraints

keeping our problem in mind:

- input: f_i's and α_i's and degree constraints $d_1, \ldots, d_m \in \mathbb{Z}_{>0}$
- output: a solution $\mathbf{p}$ satisfying the constraints $\mathrm{cdeg}(\mathbf{p}) < (d_1, \ldots, d_m)$

obstacle:
computing a reduced basis of solutions ignores the constraints

exercice: suppose we have a reduced basis $\mathbf{P} \in \mathbb{K}[X]^{m \times m}$ of solutions

- think of particular constraints $(d_1, \ldots, d_m)$ that can be handled via $\mathbf{P}$
- give constraints $(d_1, \ldots, d_m)$ for which $\mathbf{P}$ is "typically" not satisfactory

solution: compute $\mathbf{P}$ in **shifted** reduced form

$$\mathbf{A} = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \\ 3X^3+X^2+5X+3 & 6X+5 & 2X+1 \end{bmatrix}$$

using elementary row operations, transform $\mathbf{A}$ into...

Hermite form $\mathbf{H} = \begin{bmatrix} X^6+6X^4+X^3+X+4 & 0 & 0 \\ 5X^5+5X^4+6X^3+2X^2+6X+3 & X & 0 \\ 3X^4+5X^3+4X^2+6X+1 & 5 & 1 \end{bmatrix}$

Popov form $\mathbf{P} = \begin{bmatrix} X^3+5X^2+4X+1 & 2X+4 & 3X+5 \\ 1 & X^2+2X+3 & X+2 \\ 3X+2 & 4X & X^2 \end{bmatrix}$

shifted reduced forms

shifted forms and degree constraints

nonsingular $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$

elementary row transformations

Hermite form [Hermite, 1851]
- triangular
- column normalized

$$\begin{bmatrix} \mathbf{16} & & & \\ 15 & \mathbf{0} & & \\ 15 & & \mathbf{0} & \\ 15 & & & \mathbf{0} \end{bmatrix} \qquad \begin{bmatrix} \mathbf{4} & & & \\ 3 & \mathbf{7} & & \\ 1 & 5 & \mathbf{3} & \\ 3 & 6 & 1 & \mathbf{2} \end{bmatrix}$$

shifted reduced forms

shifted forms and degree constraints

nonsingular $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$

elementary row transformations

Hermite form [Hermite, 1851]
- triangular
- column normalized

Popov form [Popov, 1972]
- row reduced/distinct pivots
- column normalized

$$\begin{bmatrix} \mathbf{16} & & & \\ 15 & \mathbf{0} & & \\ 15 & & \mathbf{0} & \\ 15 & & & \mathbf{0} \end{bmatrix} \qquad \begin{bmatrix} \mathbf{4} & & & \\ 3 & \mathbf{7} & & \\ 1 & 5 & \mathbf{3} & \\ 3 & 6 & 1 & \mathbf{2} \end{bmatrix} \qquad \begin{bmatrix} \mathbf{4} & 3 & 3 & 3 \\ 3 & \mathbf{4} & 3 & 3 \\ 3 & 3 & \mathbf{4} & 3 \\ 3 & 3 & 3 & \mathbf{4} \end{bmatrix} \qquad \begin{bmatrix} \mathbf{7} & 0 & 1 & 5 \\ 0 & \mathbf{1} & & 0 \\ & & \mathbf{2} & \\ 6 & 0 & 1 & \mathbf{6} \end{bmatrix}$$

shifted reduced forms

shifted forms and degree constraints

nonsingular $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$

elementary row transformations

Hermite form [Hermite, 1851]

- triangular
- column normalized

Popov form [Popov, 1972]

- row reduced/distinct pivots
- column normalized

$$\begin{bmatrix} \mathbf{16} & & & \\ 15 & \mathbf{0} & & \\ 15 & & \mathbf{0} & \\ 15 & & & \mathbf{0} \end{bmatrix} \quad \begin{bmatrix} \mathbf{4} & & & \\ 3 & \mathbf{7} & & \\ 1 & 5 & \mathbf{3} & \\ 3 & 6 & 1 & \mathbf{2} \end{bmatrix} \qquad \begin{bmatrix} \mathbf{4} & 3 & 3 & 3 \\ 3 & \mathbf{4} & 3 & 3 \\ 3 & 3 & \mathbf{4} & 3 \\ 3 & 3 & 3 & \mathbf{4} \end{bmatrix} \quad \begin{bmatrix} \mathbf{7} & 0 & 1 & 5 \\ 0 & \mathbf{1} & & 0 \\ & & \mathbf{2} & \\ 6 & 0 & 1 & \mathbf{6} \end{bmatrix}$$

$\preccurlyeq_{\mathrm{pot}}$ reduced Gröbner basis $\preccurlyeq_{\mathrm{top}}$

$\mathbb{K}[X]$-module $\mathcal{S} \subset \mathbb{K}[X]^{1 \times m}$ of rank m

shifted reduced forms

shifted forms and degree constraints

nonsingular $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$

elementary row transformations

Hermite form [Hermite, 1851]
- triangular
- column normalized

Popov form [Popov, 1972]
- row reduced/distinct pivots
- column normalized

$$\begin{bmatrix} \mathbf{16} & & & \\ 15 & \mathbf{0} & & \\ 15 & & \mathbf{0} & \\ 15 & & & \mathbf{0} \end{bmatrix} \quad \begin{bmatrix} \mathbf{4} & & & \\ 3 & \mathbf{7} & & \\ 1 & 5 & \mathbf{3} & \\ 3 & 6 & 1 & \mathbf{2} \end{bmatrix} \qquad \begin{bmatrix} \mathbf{4} & 3 & 3 & 3 \\ 3 & \mathbf{4} & 3 & 3 \\ 3 & 3 & \mathbf{4} & 3 \\ 3 & 3 & 3 & \mathbf{4} \end{bmatrix} \quad \begin{bmatrix} \mathbf{7} & 0 & 1 & 5 \\ 0 & \mathbf{1} & & 0 \\ & & \mathbf{2} & \\ 6 & 0 & 1 & \mathbf{6} \end{bmatrix}$$

invariant: $D = \deg(\det(\mathbf{A})) = 4+7+3+2 = 7+1+2+6$

- average column degree is $\frac{D}{m}$
- size of object is $mD + m^2 = m^2(\frac{D}{m} + 1)$

shifted reduced forms

shifted forms and degree constraints

nonsingular $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$

elementary row transformations

Hermite form [Hermite, 1851]
- triangular
- column normalized

Popov form [Popov, 1972]
- row reduced/distinct pivots
- column normalized

$$\begin{bmatrix} \mathbf{16} & & & \\ 15 & \mathbf{0} & & \\ 15 & & \mathbf{0} & \\ 15 & & & \mathbf{0} \end{bmatrix} \quad \begin{bmatrix} \mathbf{4} & & & \\ 3 & \mathbf{7} & & \\ 1 & 5 & \mathbf{3} & \\ 3 & 6 & 1 & \mathbf{2} \end{bmatrix} \qquad \begin{bmatrix} \mathbf{4} & 3 & 3 & 3 \\ 3 & \mathbf{4} & 3 & 3 \\ 3 & 3 & \mathbf{4} & 3 \\ 3 & 3 & 3 & \mathbf{4} \end{bmatrix} \quad \begin{bmatrix} \mathbf{7} & 0 & 1 & 5 \\ 0 & \mathbf{1} & & 0 \\ & & \mathbf{2} & \\ 6 & 0 & 1 & \mathbf{6} \end{bmatrix}$$

[Beckermann-Labahn-Villard, 1999; Mulders-Storjohann, 2003]

shifted reduced form:
arbitrary degree constraints + **no** column normalization

$\approx$ minimal, non-reduced, $\prec$-Gröbner basis

shifted reduced forms

shift: integer tuple $\mathbf{s} = (s_1, \ldots, s_m)$ acting as column weights

$\rightarrow$ connects Popov and Hermite forms

$\mathbf{s} = (0,0,0,0)$ Popov	$\begin{bmatrix} \mathbf{4} & 3 & 3 & 3 \\ 3 & \mathbf{4} & 3 & 3 \\ 3 & 3 & \mathbf{4} & 3 \\ 3 & 3 & 3 & \mathbf{4} \end{bmatrix}$	$\begin{bmatrix} \mathbf{7} & 0 & 1 & 5 \\ 0 & \mathbf{1} & & 0 \\ & & \mathbf{2} & \\ 6 & 0 & 1 & \mathbf{6} \end{bmatrix}$
$\mathbf{s} = (0,2,4,6)$ $\mathbf{s}$-Popov	$\begin{bmatrix} \mathbf{7} & 4 & 2 & 0 \\ 6 & \mathbf{5} & 2 & 0 \\ 6 & 4 & \mathbf{3} & 0 \\ 6 & 4 & 2 & \mathbf{1} \end{bmatrix}$	$\begin{bmatrix} \mathbf{8} & 5 & 1 & \\ 7 & \mathbf{6} & 1 & \\ & & \mathbf{2} & \\ 0 & 1 & & \mathbf{0} \end{bmatrix}$
$\mathbf{s} = (0, D, 2D, 3D)$ Hermite	$\begin{bmatrix} \mathbf{16} & & & \\ 15 & \mathbf{0} & & \\ 15 & & \mathbf{0} & \\ 15 & & & \mathbf{0} \end{bmatrix}$	$\begin{bmatrix} \mathbf{4} & & & \\ 3 & \mathbf{7} & & \\ 1 & 5 & \mathbf{3} & \\ 3 & 6 & 1 & \mathbf{2} \end{bmatrix}$

- normal form, average column degree D/m
- shifted reduced form: same without normalization
- shifts arise naturally in algorithms (approximants, kernel, ...)

shifted row degree of a polynomial matrix
= the list of the maximum **shifted** degree in each of its rows

for $\mathbf{A} = (a_{i,j}) \in \mathbb{K}[X]^{m \times n}$, and $\mathbf{s} = (s_1, \ldots, s_n) \in \mathbb{Z}^n$,

$$\begin{aligned}\mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}) &= (\mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}_{1,*}), \ldots, \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}_{m,*})) \\ &= \left(\max_{1 \leqslant j \leqslant n} (\deg(\mathbf{A}_{1,j}) + s_j), \;\; \ldots, \;\; \max_{1 \leqslant j \leqslant n} (\deg(\mathbf{A}_{m,j}) + s_j) \right) \in \mathbb{Z}^m\end{aligned}$$

example: for the matrix $\mathbf{A} = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \end{bmatrix}$,
describe $\mathrm{rdeg}_{(0,0,0)}(\mathbf{A})$, $\mathrm{rdeg}_{(0,1,2)}(\mathbf{A})$, and $\mathrm{rdeg}_{(-1,-3,-2)}(\mathbf{A})$

shifted reduced forms

shifted forms and degree constraints

shifted row degree of a polynomial matrix
$=$ the list of the maximum **shifted** degree in each of its rows

for $\mathbf{A} = (a_{i,j}) \in \mathbb{K}[X]^{m \times n}$, and $\mathbf{s} = (s_1, \ldots, s_n) \in \mathbb{Z}^n$,

$$\mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}) = (\mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}_{1,*}), \ldots, \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}_{m,*}))$$
$$= \left(\max_{1 \leqslant j \leqslant n} (\deg(\mathbf{A}_{1,j}) + s_j), \;\; \ldots, \;\; \max_{1 \leqslant j \leqslant n} (\deg(\mathbf{A}_{m,j}) + s_j) \right) \in \mathbb{Z}^m$$

example: for the matrix $\mathbf{A} = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \end{bmatrix}$,
describe $\mathrm{rdeg}_{(0,0,0)}(\mathbf{A})$, $\mathrm{rdeg}_{(0,1,2)}(\mathbf{A})$, and $\mathrm{rdeg}_{(-1,-3,-2)}(\mathbf{A})$

- $\mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}) = \mathrm{rdeg}(\mathbf{A}\mathbf{X}^{\mathbf{s}})$
- $\mathrm{rdeg}_{\mathbf{s}}(\mathbf{A})$ only depends on the degrees in $\mathbf{A}$
- $\mathrm{rdeg}_{\mathbf{s}+(c,\ldots,c)}(\mathbf{A}) = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}) + c$

notation:

let $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ with no zero row, and $\mathbf{s} \in \mathbb{Z}^n$,
define $\mathbf{d} = (d_1, \ldots, d_m) = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A})$

and $$\mathbf{X}^{\mathbf{d}} = \begin{bmatrix} X^{d_1} & & \\ & \ddots & \\ & & X^{d_m} \end{bmatrix} \in \mathbb{K}[X, X^{-1}]^{m \times m}$$

definition: s-leading matrix / s-reduced matrix

assuming $\mathbf{s} \geqslant 0$,

- the $\mathbf{s}$-leading matrix of $\mathbf{A}$ is $\mathrm{lm}_{\mathbf{s}}(\mathbf{A}) = \mathrm{lm}(\mathbf{A}\mathbf{X}^{\mathbf{s}}) \in \mathbb{K}^{m \times n}$
- $\mathbf{A} \in \mathbb{K}[X]^{m \times n}$ is reduced if $\mathrm{lm}_{\mathbf{s}}(\mathbf{A})$ has full row rank

shifted reduced forms

shifted forms and degree constraints

notation:

let $\mathbf{A} \in \mathbb{K}[X]^{m\times n}$ with no zero row, and $\mathbf{s} \in \mathbb{Z}^n$,
define $\mathbf{d} = (d_1, \ldots, d_m) = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A})$

and $$\mathbf{X}^{\mathbf{d}} = \begin{bmatrix} X^{d_1} & & \\ & \ddots & \\ & & X^{d_m} \end{bmatrix} \in \mathbb{K}[X, X^{-1}]^{m\times m}$$

definition: s-leading matrix / s-reduced matrix

assuming $\mathbf{s} \geqslant 0$,

- the $\mathbf{s}$-leading matrix of $\mathbf{A}$ is $\mathrm{lm}_{\mathbf{s}}(\mathbf{A}) = \mathrm{lm}(\mathbf{A}\mathbf{X}^{\mathbf{s}}) \in \mathbb{K}^{m\times n}$
- $\mathbf{A} \in \mathbb{K}[X]^{m\times n}$ is reduced if $\mathrm{lm}_{\mathbf{s}}(\mathbf{A})$ has full row rank

- these notions are invariant under $\mathbf{s} \to \mathbf{s} + (c, \ldots, c)$
- they coincide with the non-shifted case when $\mathbf{s} = (0, \ldots, 0)$
- $\mathbf{X}^{-\mathbf{d}}\mathbf{A}\mathbf{X}^{\mathbf{s}} = \mathrm{lm}_{\mathbf{s}}(\mathbf{A}) +$ terms of strictly negative degree

exercise: for each of the matrices below, and each shift s,
1. give the s-leading matrix
2. deduce whether the matrix is s-reduced

$$\mathbf{A} = \begin{bmatrix} 3X+4 & X^3+4X+1 & 4X^2+3 \\ 5 & 5X^2+3X+1 & 5X+3 \\ 3X^3+X^2+5X+3 & 6X+5 & 2X+1 \end{bmatrix}$$

$$\mathbf{H} = \begin{bmatrix} X^6+6X^4+X^3+X+4 & 0 & 0 \\ 5X^5+5X^4+6X^3+2X^2+6X+3 & X & 0 \\ 3X^4+5X^3+4X^2+6X+1 & 5 & 1 \end{bmatrix}$$

$$\mathbf{P} = \begin{bmatrix} X^3+5X^2+4X+1 & 2X+4 & 3X+5 \\ 1 & X^2+2X+3 & X+2 \\ 3X+2 & 4X & X^2 \end{bmatrix}$$

$$\mathbf{s} = (0,0,0),\ \mathbf{s} = (0,5,6),\ \mathbf{s} = (-3,-2,-2)$$

shifted reduced forms

shifted forms and degree constraints

the characterizations generalize to the $\mathbf{s}$-shifted case,
using $\mathbf{s}$-row degrees and $\mathbf{s}$-leading matrices where appropriate

(proofs: direct reductions, with: $\mathbf{A}$ is $\mathbf{s}$-reduced $\Leftrightarrow$ $\mathbf{A}\mathbf{X}^{\mathbf{s}}$ is reduced)

for example recall the predictable degree property:

$\mathbf{A}$ is reduced if and only if for any $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1\times m}$,
$$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \max_{1\leqslant i\leqslant m}(\deg(u_i) + \mathrm{rdeg}(\mathbf{A}_{i,*}))$$

shifted reduced forms

shifted forms and degree constraints

the characterizations generalize to the $\mathbf{s}$-shifted case,
using $\mathbf{s}$-row degrees and $\mathbf{s}$-leading matrices where appropriate

(proofs: direct reductions, with: $\mathbf{A}$ is $\mathbf{s}$-reduced $\Leftrightarrow$ $\mathbf{A}\mathbf{X}^{\mathbf{s}}$ is reduced)

for example recall the predictable degree property:

$\mathbf{A}$ is reduced if and only if for any $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1\times m}$,
$$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \max_{1\leqslant i\leqslant m}(\deg(u_i) + \mathrm{rdeg}(\mathbf{A}_{i,*}))$$

- this means $\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \mathrm{rdeg}_{\mathbf{t}}(\mathbf{u})$ where $\mathbf{t} = \mathrm{rdeg}(\mathbf{A})$
- i.e. $\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \mathrm{rdeg}(\mathbf{u}\mathbf{X}^{\mathrm{rdeg}(\mathbf{A})})$, “no surprising cancellation”
- proof: let $\delta = \mathrm{rdeg}_{\mathbf{t}}(\mathbf{u})$, our goal is to show $\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \delta$

terms of $X^{-\delta}\mathbf{u}\mathbf{A}$ have degree $\leqslant 0$,
and $X^{-\delta}\mathbf{u}\mathbf{A} = (X^{-\delta}\mathbf{u}\mathbf{X}^{\mathbf{t}})(\mathbf{X}^{-\mathbf{t}}\mathbf{A})$;
the term of degree 0 is $\mathrm{lm}_{\mathbf{t}}(\mathbf{u})\mathrm{lm}(\mathbf{A})$,
it is nonzero since $\mathrm{lm}(\mathbf{A})$ has full rank and $\mathrm{lm}_{\mathbf{t}}(\mathbf{u}) \neq 0$
(the case $\mathbf{u} = \mathbf{0}$ is trivial)

shifted reduced forms

shifted forms and degree constraints

the characterizations generalize to the s-shifted case,
using s-row degrees and s-leading matrices where appropriate

(proofs: direct reductions, with: $\mathbf{A}$ is s-reduced $\Leftrightarrow$ $\mathbf{A}\mathbf{X}^{\mathbf{s}}$ is reduced)

for example recall the predictable degree property:

$\mathbf{A}$ is reduced if and only if for any $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1\times m}$,
$$\mathrm{rdeg}(\mathbf{u}\mathbf{A}) = \max_{1\leqslant i\leqslant m}(\deg(u_i) + \mathrm{rdeg}(\mathbf{A}_{i,*}))$$

$\mathbf{A}$ is s-reduced if and only if for any $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1\times m}$,
$$\mathrm{rdeg}_{\mathbf{s}}(\mathbf{u}\mathbf{A}) = \max_{1\leqslant i\leqslant m}(\deg(u_i) + \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}_{i,*}))$$

this means $\mathrm{rdeg}_{\mathbf{s}}(\mathbf{u}\mathbf{A}) = \mathrm{rdeg}_{\mathbf{t}}(\mathbf{u})$, where $\mathbf{t} = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A})$

the characterizations generalize to the s-shifted case,
using s-row degrees and s-leading matrices where appropriate

(proofs: direct reductions, with: $\mathbf{A}$ is s-reduced $\Leftrightarrow$ $\mathbf{A}\mathbf{X}^{s}$ is reduced)

for example recall the predictable degree property:

$\mathbf{A}$ is reduced if and only if for any $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1\times m}$,
$$\mathrm{rdeg}(\mathbf{uA}) = \max_{1\leqslant i\leqslant m}(\deg(u_i) + \mathrm{rdeg}(\mathbf{A}_{i,*}))$$

$\mathbf{A}$ is s-reduced if and only if for any $\mathbf{u} = [u_1 \cdots u_m] \in \mathbb{K}[X]^{1\times m}$,
$$\mathrm{rdeg}_{\mathbf{s}}(\mathbf{uA}) = \max_{1\leqslant i\leqslant m}(\deg(u_i) + \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A}_{i,*}))$$
this means $\mathrm{rdeg}_{\mathbf{s}}(\mathbf{uA}) = \mathrm{rdeg}_{\mathbf{t}}(\mathbf{u})$, where $\mathbf{t} = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{A})$

- s-reduced forms provide vectors of minimal s-degree in the module
- satisfying degree constraints $(d_1, \ldots, d_m)$ $\Rightarrow$ taking $\mathbf{s} = (-d_1, \ldots, -d_m)$
- indeed $\mathrm{cdeg}([p_1 \;\; \cdots \;\; p_m]) < (d_1, \ldots, d_m)$

if and only if $\mathrm{rdeg}_{(-d_1,\ldots,-d_m)}([p_1 \;\; \cdots \;\; p_m]) < 0$

shifted reduced forms

stability under multiplication

algorithms based on polynomial matrix multiplication

[iterative: van Barel-Bultheel 1991, Beckermann-Labahn 2000]
[divide and conquer: Beckermann-Labahn 1994, Giorgi-Jeannerod-Villard 2003]

- compute a first basis $\mathbf{P}_1$ for a subproblem
- update the input instance to get the second subproblem
- compute a second basis $\mathbf{P}_2$ for this second subproblem
- the output basis of solutions is $\mathbf{P}_2\mathbf{P}_1$

we want $\mathbf{P}_2\mathbf{P}_1$ to be reduced:

1. is it implied by "$\mathbf{P}_1$ reduced and $\mathbf{P}_2$ reduced"?
2. any idea of how to fix this?

shifted reduced forms

stability under multiplication

algorithms based on polynomial matrix multiplication

[iterative: van Barel-Bultheel 1991, Beckermann-Labahn 2000]
[divide and conquer: Beckermann-Labahn 1994, Giorgi-Jeannerod-Villard 2003]

- compute a first basis $\mathbf{P}_1$ for a subproblem
- update the input instance to get the second subproblem
- compute a second basis $\mathbf{P}_2$ for this second subproblem
- the output basis of solutions is $\mathbf{P}_2\mathbf{P}_1$

we want $\mathbf{P}_2\mathbf{P}_1$ to be reduced:
1. is it implied by "$\mathbf{P}_1$ reduced and $\mathbf{P}_2$ reduced"?
2. any idea of how to fix this?

we want $\mathbf{P}_2\mathbf{P}_1$ to be reduced
theorem: implied by "$\mathbf{P}_1$ is reduced and $\mathbf{P}_2$ is $\mathbf{t}$-reduced"
where $\mathbf{t} = \mathrm{rdeg}(\mathbf{P}_1)$

stability under multiplication

algorithms based on polynomial matrix multiplication

[iterative: van Barel-Bultheel 1991, Beckermann-Labahn 2000]
[divide and conquer: Beckermann-Labahn 1994, Giorgi-Jeannerod-Villard 2003]

- compute a first basis $\mathbf{P}_1$ for a subproblem
- update the input instance to get the second subproblem
- compute a second basis $\mathbf{P}_2$ for this second subproblem
- the output basis of solutions is $\mathbf{P}_2\mathbf{P}_1$

we want $\mathbf{P}_2\mathbf{P}_1$ to be reduced:
1. is it implied by "$\mathbf{P}_1$ reduced and $\mathbf{P}_2$ reduced"?
2. any idea of how to fix this?

we want $\mathbf{P}_2\mathbf{P}_1$ to be $\boldsymbol{s}$-reduced
theorem: implied by "$\mathbf{P}_1$ is $\boldsymbol{s}$-reduced and $\mathbf{P}_2$ is $\mathbf{t}$-reduced"
where $\mathbf{t} = \mathrm{rdeg}_{\boldsymbol{s}}(\mathbf{P}_1)$

shifted reduced forms

stability under multiplication

let $\mathcal{M} \subseteq \mathcal{M}_1$ be two $\mathbb{K}[X]$-submodules of $\mathbb{K}[X]^m$ of rank m,
let $\mathbf{P}_1 \in \mathbb{K}[X]^{m \times m}$ be a basis of $\mathcal{M}_1$,
let $\mathbf{s} \in \mathbb{Z}^m$ and $\mathbf{t} = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{P}_1)$,

- the rank of the module $\mathcal{M}_2 = \{\boldsymbol{\lambda} \in \mathbb{K}[X]^{1 \times m} \mid \boldsymbol{\lambda}\mathbf{P}_1 \in \mathcal{M}\}$ is m and for any basis $\mathbf{P}_2 \in \mathbb{K}[X]^{m \times m}$ of $\mathcal{M}_2$, the product $\mathbf{P}_2\mathbf{P}_1$ is a basis of $\mathcal{M}$
- if $\mathbf{P}_1$ is $\mathbf{s}$-reduced and $\mathbf{P}_2$ is $\mathbf{t}$-reduced, then $\mathbf{P}_2\mathbf{P}_1$ is $\mathbf{s}$-reduced

shifted reduced forms

stability under multiplication

let $\mathcal{M} \subseteq \mathcal{M}_1$ be two $\mathbb{K}[X]$-submodules of $\mathbb{K}[X]^m$ of rank m,
let $\mathbf{P}_1 \in \mathbb{K}[X]^{m \times m}$ be a basis of $\mathcal{M}_1$,
let $s \in \mathbb{Z}^m$ and $t = \mathrm{rdeg}_s(\mathbf{P}_1)$,

- the rank of the module $\mathcal{M}_2 = \{\boldsymbol{\lambda} \in \mathbb{K}[X]^{1 \times m} \mid \boldsymbol{\lambda}\mathbf{P}_1 \in \mathcal{M}\}$ is m and for any basis $\mathbf{P}_2 \in \mathbb{K}[X]^{m \times m}$ of $\mathcal{M}_2$, the product $\mathbf{P}_2\mathbf{P}_1$ is a basis of $\mathcal{M}$
- if $\mathbf{P}_1$ is s-reduced and $\mathbf{P}_2$ is t-reduced, then $\mathbf{P}_2\mathbf{P}_1$ is s-reduced

Let $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$ denote the adjugate of $\mathbf{P}_1$. Then, we have $\mathbf{A}\mathbf{P}_1 = \det(\mathbf{P}_1)\mathbf{I}_m$. Thus, $\mathbf{p}\mathbf{A}\mathbf{P}_1 = \det(\mathbf{P}_1)\mathbf{p} \in \mathcal{M}$ for all $\mathbf{p} \in \mathcal{M}$, and therefore $\mathcal{M}\mathbf{A} \subseteq \mathcal{M}_2$. Now, the nonsingularity of $\mathbf{A}$ ensures that $\mathcal{M}\mathbf{A}$ has rank m; this implies that $\mathcal{M}_2$ has rank m as well (see e.g. [Dummit-Foote 2004, Sec. 12.1, Thm. 4]). The matrix $\mathbf{P}_2\mathbf{P}_1$ is nonsingular since $\det(\mathbf{P}_2\mathbf{P}_1) \neq 0$. Now let $\mathbf{p} \in \mathcal{M}$; we want to prove that $\mathbf{p}$ is a $\mathbb{K}[X]$-linear combination of the rows of $\mathbf{P}_2\mathbf{P}_1$. First, $\mathbf{p} \in \mathcal{M}_1$, so there exists $\boldsymbol{\lambda} \in \mathbb{K}[X]^{1 \times m}$ such that $\mathbf{p} = \boldsymbol{\lambda}\mathbf{P}_1$. But then $\boldsymbol{\lambda} \in \mathcal{M}_2$, and thus there exists $\boldsymbol{\mu} \in \mathbb{K}[X]^{1 \times m}$ such that $\boldsymbol{\lambda} = \boldsymbol{\mu}\mathbf{P}_2$. This yields the combination $\mathbf{p} = \boldsymbol{\mu}\mathbf{P}_2\mathbf{P}_1$.

shifted reduced forms

stability under multiplication

let $\mathcal{M} \subseteq \mathcal{M}_1$ be two $\mathbb{K}[X]$-submodules of $\mathbb{K}[X]^m$ of rank m,
let $\mathbf{P}_1 \in \mathbb{K}[X]^{m \times m}$ be a basis of $\mathcal{M}_1$,
let $\mathbf{s} \in \mathbb{Z}^m$ and $\mathbf{t} = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{P}_1)$,

- the rank of the module $\mathcal{M}_2 = \{\boldsymbol{\lambda} \in \mathbb{K}[X]^{1 \times m} \mid \boldsymbol{\lambda}\mathbf{P}_1 \in \mathcal{M}\}$ is m and for any basis $\mathbf{P}_2 \in \mathbb{K}[X]^{m \times m}$ of $\mathcal{M}_2$, the product $\mathbf{P}_2\mathbf{P}_1$ is a basis of $\mathcal{M}$
- if $\mathbf{P}_1$ is $\mathbf{s}$-reduced and $\mathbf{P}_2$ is $\mathbf{t}$-reduced, then $\mathbf{P}_2\mathbf{P}_1$ is $\mathbf{s}$-reduced

Let $\mathbf{d} = \mathrm{rdeg}_{\mathbf{t}}(\mathbf{P}_2)$; we have $\mathbf{d} = \mathrm{rdeg}_{\mathbf{s}}(\mathbf{P}_2\mathbf{P}_1)$ by the predictable degree property. Using $\mathbf{X}^{-\mathbf{d}}\mathbf{P}_2\mathbf{P}_1\mathbf{X}^{\mathbf{s}} = \mathbf{X}^{-\mathbf{d}}\mathbf{P}_2\mathbf{X}^{\mathbf{t}}\mathbf{X}^{-\mathbf{t}}\mathbf{P}_1\mathbf{X}^{\mathbf{s}}$, we obtain that $\mathrm{lm}_{\mathbf{s}}(\mathbf{P}_2\mathbf{P}_1) = \mathrm{lm}_{\mathbf{t}}(\mathbf{P}_2)\mathrm{lm}_{\mathbf{s}}(\mathbf{P}_1)$. By assumption, $\mathrm{lm}_{\mathbf{t}}(\mathbf{P}_2)$ and $\mathrm{lm}_{\mathbf{s}}(\mathbf{P}_1)$ are invertible, and therefore $\mathrm{lm}_{\mathbf{s}}(\mathbf{P}_2\mathbf{P}_1)$ is invertible as well; thus $\mathbf{P}_2\mathbf{P}_1$ is $\mathbf{s}$-reduced.

outline

- **introduction**
  - rational approximation and interpolation
  - the vector case
  - pol. matrices: reminders and motivation
- **shifted reduced forms**
  - reducedness: examples and properties
  - shifted forms and degree constraints
  - stability under multiplication
- **fast algorithms**
- **next time**

outline

- **introduction**
  - rational approximation and interpolation
  - the vector case
  - pol. matrices: reminders and motivation
- **shifted reduced forms**
  - reducedness: examples and properties
  - shifted forms and degree constraints
  - stability under multiplication
- **fast algorithms**
  - iterative algorithm and output size
  - base case: modulus of degree 1
  - recursion: residual and basis multiplication
- **next time**

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

input: vector $\mathbf{F} = \begin{bmatrix} f_1 \\ \vdots \\ f_m \end{bmatrix}$, points $\alpha_1, \ldots, \alpha_d \in \mathbb{K}$, shift $\mathbf{s} = (s_1, \ldots, s_m) \in \mathbb{Z}^m$

1. $\mathbf{P} = \begin{bmatrix} -\mathbf{p}_1- \\ \vdots \\ -\mathbf{p}_m- \end{bmatrix}$ = identity matrix in $\mathbb{K}[X]^{m \times m}$

2. for i from 1 to d:

a. evaluate updated vector $\begin{bmatrix} (\mathbf{p}_1 \cdot \mathbf{F})(\alpha_i) \\ \vdots \\ (\mathbf{p}_m \cdot \mathbf{F})(\alpha_i) \end{bmatrix} = (\mathbf{P} \cdot \mathbf{F})(\alpha_i)$

b. choose pivot π with smallest s_π such that $(\mathbf{p}_\pi \cdot \mathbf{F})(\alpha_i) \neq 0$
update pivot shift $s_\pi = s_\pi + 1$

c. eliminate: `/* after this,` $\forall j \neq \pi$, $(\mathbf{p}_j \cdot \mathbf{F})(\alpha_i) = 0$ `*/`

for $j \neq \pi$ do $\mathbf{p}_j \leftarrow \mathbf{p}_j - \dfrac{(\mathbf{p}_j \cdot \mathbf{F})(\alpha_i)}{(\mathbf{p}_\pi \cdot \mathbf{F})(\alpha_i)} \mathbf{p}_\pi$; $\qquad \mathbf{p}_\pi \leftarrow (X - \alpha_i)\mathbf{p}_\pi$

after i iterations: $\mathbf{P}$ is an $\mathbf{s}$-reduced basis of solutions for $(\alpha_1, \ldots, \alpha_i)$

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^\mathsf{T}$

iteration: $i = 1$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[0 \quad 2 \quad 4 \quad 6]$

basis $\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$

values $\begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 80 & 73 & 73 & 35 & 66 & 46 & 91 & 64 \\ 95 & 91 & 91 & 61 & 88 & 79 & 36 & 22 \\ 34 & 47 & 47 & 1 & 85 & 45 & 75 & 50 \end{bmatrix}$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 1$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[0 \quad 2 \quad 4 \quad 6]$

basis
$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 80 & 73 & 73 & 35 & 66 & 46 & 91 & 64 \\ 95 & 91 & 91 & 61 & 88 & 79 & 36 & 22 \\ 34 & 47 & 47 & 1 & 85 & 45 & 75 & 50 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 1$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[0 \quad 2 \quad 4 \quad 6]$

basis $\begin{bmatrix} 1 & 0 & 0 & 0 \\ 17 & 1 & 0 & 0 \\ 2 & 0 & 1 & 0 \\ 63 & 0 & 0 & 1 \end{bmatrix}$

values $\begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 90 & 90 & 52 & 83 & 63 & 11 & 81 \\ 0 & 93 & 93 & 63 & 90 & 81 & 38 & 24 \\ 0 & 13 & 13 & 64 & 51 & 11 & 41 & 16 \end{bmatrix}$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $\quad m = 4$ $\quad \mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 1$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[1 \quad 2 \quad 4 \quad 6]$

basis $\begin{bmatrix} X + 73 & 0 & 0 & 0 \\ 17 & 1 & 0 & 0 \\ 2 & 0 & 1 & 0 \\ 63 & 0 & 0 & 1 \end{bmatrix}$

values $\begin{bmatrix} 0 & 7 & 88 & 8 & 59 & 3 & 93 & 35 \\ 0 & 90 & 90 & 52 & 83 & 63 & 11 & 81 \\ 0 & 93 & 93 & 63 & 90 & 81 & 38 & 24 \\ 0 & 13 & 13 & 64 & 51 & 11 & 41 & 16 \end{bmatrix}$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $\quad m = 4$ $\quad \mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 2$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[1 \;\; 2 \;\; 4 \;\; 6]$

basis
$$\begin{bmatrix} X+73 & 0 & 0 & 0 \\ 17 & 1 & 0 & 0 \\ 2 & 0 & 1 & 0 \\ 63 & 0 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 0 & 7 & 88 & 8 & 59 & 3 & 93 & 35 \\ 0 & 90 & 90 & 52 & 83 & 63 & 11 & 81 \\ 0 & 93 & 93 & 63 & 90 & 81 & 38 & 24 \\ 0 & 13 & 13 & 64 & 51 & 11 & 41 & 16 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 2$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[1 \quad 2 \quad 4 \quad 6]$

basis
$$\begin{bmatrix} X + 73 & 0 & 0 & 0 \\ X + 90 & 1 & 0 & 0 \\ 56X + 16 & 0 & 1 & 0 \\ 12X + 66 & 0 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 0 & 7 & 88 & 8 & 59 & 3 & 93 & 35 \\ 0 & 0 & 81 & 60 & 45 & 66 & 7 & 19 \\ 0 & 0 & 74 & 26 & 96 & 55 & 8 & 44 \\ 0 & 0 & 2 & 63 & 80 & 47 & 90 & 48 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $\quad m = 4$ $\quad \mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 2$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[2 \quad 2 \quad 4 \quad 6]$

basis
$$\begin{bmatrix} X^2 + 42X + 65 & 0 & 0 & 0 \\ X + 90 & 1 & 0 & 0 \\ 56X + 16 & 0 & 1 & 0 \\ 12X + 66 & 0 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 0 & 0 & 47 & 8 & 61 & 85 & 44 & 10 \\ 0 & 0 & 81 & 60 & 45 & 66 & 7 & 19 \\ 0 & 0 & 74 & 26 & 96 & 55 & 8 & 44 \\ 0 & 0 & 2 & 63 & 80 & 47 & 90 & 48 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $\quad m = 4$ $\quad \mathbf{s} = (0, 2, 4, 6)$, $\quad$ base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 3$ $\qquad$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $\quad [2 \;\; 2 \;\; 4 \;\; 6]$

basis

$$\begin{bmatrix} X^2 + 42X + 65 & 0 & 0 & 0 \\ X + 90 & 1 & 0 & 0 \\ 56X + 16 & 0 & 1 & 0 \\ 12X + 66 & 0 & 0 & 1 \end{bmatrix}$$

values

$$\begin{bmatrix} 0 & 0 & 47 & 8 & 61 & 85 & 44 & 10 \\ 0 & 0 & 81 & 60 & 45 & 66 & 7 & 19 \\ 0 & 0 & 74 & 26 & 96 & 55 & 8 & 44 \\ 0 & 0 & 2 & 63 & 80 & 47 & 90 & 48 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 3$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[3 \quad 2 \quad 4 \quad 6]$

basis
$$\begin{bmatrix} X^3 + 27X^2 + 17X + 92 & 0 & 0 & 0 \\ 54X^2 + 38X + 11 & 1 & 0 & 0 \\ 17X^2 + 91X + 54 & 0 & 1 & 0 \\ 66X^2 + 68X + 88 & 0 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 0 & 0 & 0 & 39 & 74 & 50 & 26 & 52 \\ 0 & 0 & 0 & 7 & 41 & 0 & 55 & 74 \\ 0 & 0 & 0 & 65 & 66 & 45 & 77 & 20 \\ 0 & 0 & 0 & 9 & 32 & 31 & 84 & 29 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 4$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[3 \;\; 2 \;\; 4 \;\; 6]$

basis
$$\begin{bmatrix} X^3 + 27X^2 + 17X + 92 & 0 & 0 & 0 \\ 54X^2 + 38X + 11 & 1 & 0 & 0 \\ 17X^2 + 91X + 54 & 0 & 1 & 0 \\ 66X^2 + 68X + 88 & 0 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 0 & 0 & 0 & 39 & 74 & 50 & 26 & 52 \\ 0 & 0 & 0 & 7 & 41 & 0 & 55 & 74 \\ 0 & 0 & 0 & 65 & 66 & 45 & 77 & 20 \\ 0 & 0 & 0 & 9 & 32 & 31 & 84 & 29 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 4$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[3 \quad 3 \quad 4 \quad 6]$

basis $\begin{bmatrix} X^3 + 31X^2 + 27X + 3 & 36 & 0 & 0 \\ 54X^3 + 56X^2 + 56X + 36 & X + 65 & 0 & 0 \\ 56X^2 + 43X + 35 & 60 & 1 & 0 \\ 52X^2 + 33X + 60 & 68 & 0 & 1 \end{bmatrix}$

values $\begin{bmatrix} 0 & 0 & 0 & 0 & 95 & 50 & 66 & 0 \\ 0 & 0 & 0 & 0 & 54 & 0 & 19 & 58 \\ 0 & 0 & 0 & 0 & 4 & 45 & 79 & 95 \\ 0 & 0 & 0 & 0 & 7 & 31 & 41 & 17 \end{bmatrix}$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 5$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[4 \quad 3 \quad 4 \quad 6]$

basis $\begin{bmatrix} X^4 + 45X^3 + 73X^2 + 90X + 42 & 36X + 19 & 0 & 0 \\ 81X^3 + 20X^2 + 9X + 20 & X + 67 & 0 & 0 \\ 2X^3 + 21X^2 + 41 & 35 & 1 & 0 \\ 52X^3 + 15X^2 + 79X + 22 & 0 & 0 & 1 \end{bmatrix}$

values $\begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 13 & 13 & 0 \\ 0 & 0 & 0 & 0 & 0 & 89 & 55 & 58 \\ 0 & 0 & 0 & 0 & 0 & 48 & 17 & 95 \\ 0 & 0 & 0 & 0 & 0 & 12 & 78 & 17 \end{bmatrix}$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $m = 4$ $\mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 6$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $[4 \quad 4 \quad 4 \quad 6]$

basis
$$\begin{bmatrix} X^4 + 19X^3 + 57X^2 + 44X + 26 & 74X + 43 & 0 & 0 \\ 81X^4 + 64X^3 + 51X^2 + 68X + 42 & X^2 + 40X + 34 & 0 & 0 \\ 3X^3 + 44X^2 + 54X + 64 & 6X + 49 & 1 & 0 \\ 28X^3 + 45X^2 + 44X + 52 & 50X + 52 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 66 & 70 \\ 0 & 0 & 0 & 0 & 0 & 0 & 3 & 13 \\ 0 & 0 & 0 & 0 & 0 & 0 & 56 & 55 \\ 0 & 0 & 0 & 0 & 0 & 0 & 15 & 7 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $\quad m = 4$ $\quad \mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 7$ **point:** 24, 31, 15, 32, 83, 27, 20, 59

shift $\quad [5 \quad 4 \quad 4 \quad 6]$

basis

$$\begin{bmatrix} X^5 + 96X^4 + 65X^3 + 68X^2 + 19X + 62 & 74X^2 + 18X + 13 & 0 & 0 \\ 6X^4 + 94X^3 + 44X^2 + 66X + 32 & X^2 + 19X + 10 & 0 & 0 \\ 55X^4 + 78X^3 + 75X^2 + 49X + 39 & 2X + 86 & 1 & 0 \\ 13X^4 + 81X^3 + 10X^2 + 34X + 2 & 42X + 29 & 0 & 1 \end{bmatrix}$$

values

$$\begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 14 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 25 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 44 \end{bmatrix}$$

fast algorithms

iterative algorithm [van Barel-Bultheel / Beckermann-Labahn]

parameters: $d = 8$ $\quad m = 4$ $\quad \mathbf{s} = (0, 2, 4, 6)$, base field $\mathbb{F}_{97}$

input: $(24, 31, 15, 32, 83, 27, 20, 59)$ and $\mathbf{F} = [1 \;\; L \;\; L^2 \;\; L^3]^T$

iteration: $i = 8$ **point:** $24, 31, 15, 32, 83, 27, 20, 59$

shift $[5 \quad 5 \quad 4 \quad 6]$

basis
$$\begin{bmatrix} X^5 + 12X^4 + 10X^3 + 34X^2 + 65X + 2 & 60X^2 + 43X + 67 & 0 & 0 \\ 6X^5 + 31X^4 + 27X^3 + 89X^2 + 18X + 52 & X^3 + 57X^2 + 53X + 89 & 0 & 0 \\ 2X^4 + 56X^3 + 42X^2 + 48X + 15 & 72X^2 + 12X + 30 & 1 & 0 \\ 40X^4 + 19X^3 + 14X^2 + 40X + 49 & 53X^2 + 79X + 74 & 0 & 1 \end{bmatrix}$$

values
$$\begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

fast algorithms

iterative algorithm – complexity aspects

- input size: $md + d$ elements from $\mathbb{K}$
  - md coefficients of $\mathbf{F}$, assumed reduced modulo $M(X)$
  - d points $\alpha_1, \ldots, \alpha_d$
- output size: $\leqslant m^2(d+1)$ elements from $\mathbb{K}$
  - $m \times m$ matrix $\mathbf{P}$ of degree at most i at step i

is this output size bound tight?

fast algorithms

iterative algorithm – complexity aspects

- input size: $md + d$ elements from $\mathbb{K}$
  - . md coefficients of $\mathbf{F}$, assumed reduced modulo $M(X)$
  - . d points $\alpha_1, \ldots, \alpha_d$
- output size: $\leqslant m^2(d+1)$ elements from $\mathbb{K}$
  - . $m \times m$ matrix $\mathbf{P}$ of degree at most i at step i

is this output size bound tight?

- one can prove $\deg(\det(\mathbf{P})) \leqslant d$
  - . $\mathbf{P}$ is a basis of $\mathcal{S}$, which is the kernel of $\mathbb{K}[X]^m \to \mathbb{K}[X]/\langle M(X)\rangle, \mathbf{p} \mapsto \mathbf{pF}$
  - . $\mathbb{K}[X]^m/\mathcal{S}$ has $\mathbb{K}$-dimension at most $\dim_{\mathbb{K}}(\mathbb{K}[X]/\langle M(X)\rangle) = d$
- normalized bases have average column degree $\leqslant d$, and size $\leqslant m(d+1)$
- yet the bound $\Theta(m^2(d+1))$ is tight for this algorithm
  - . normalizing at each step is feasible for the iterative version
  - . but is much harder to incorporate in fast divide and conquer versions

example instance of Hermite-Padé approximation
where the output size is in $\Omega(m^2 d)$

parameters: $\mathbb{K} = \mathbf{F}_{97}$, $m = 4$, $d = 128$, $\mathbf{s} = (0, \dots, 0)$

choose random polynomial $R(X)$ of degree < 128

$$\mathbf{F} = \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ f_4 \end{bmatrix} = \begin{bmatrix} R \\ R + XR \\ XR + X^2 R \\ X^2 R + X^3 R \end{bmatrix}$$

- solution $\mathbf{p}$ means $\mathbf{pF} = 0 \bmod X^{128}$
- $\mathbf{F}$ has small vectors in its left kernel

$\Rightarrow$ reduced approximant basis has unbalanced row degrees $(1, 1, 1, 125)$

- will help to build an example with output size $\Omega(m^2 d)$

running the iterative algorithm:

i	1
s	$(0, 0, 0, 0)$
f_1	R
f_2	$R + XR$
f_3	$XR + X^2R$
f_4	$X^2R + X^3R$

P

running the iterative algorithm:

i	1	2
$\mathbf{s}$	$(0, 0, 0, 0)$	$(1, 0, 0, 0)$
f_1	R	XR
f_2	$R + XR$	XR
f_3	$XR + X^2R$	$XR + X^2R$
f_4	$X^2R + X^3R$	$X^2R + X^3R$
$\mathbf{P}$	$\begin{bmatrix} 1 & & & \\ 0 & 0 & & \\ & & 0 & \\ & & & 0 \end{bmatrix}$	

running the iterative algorithm:

i	1	2	3
$\mathbf{s}$	$(0,0,0,0)$	$(1,0,0,0)$	$(1,1,0,0)$
f_1	R	XR	0
f_2	$R+XR$	XR	X^2R
f_3	$XR+X^2R$	$XR+X^2R$	X^2R
f_4	X^2R+X^3R	X^2R+X^3R	X^2R+X^3R
$\mathbf{P}$	$\begin{bmatrix} 1 & & & \\ 0 & 0 & & \\ & & 0 & \\ & & & 0 \end{bmatrix}$	$\begin{bmatrix} 1 & 0 & & \\ 1 & 1 & & \\ 0 & 0 & 0 & \\ & & & 0 \end{bmatrix}$	

running the iterative algorithm:

i	1	2	3	4
$\mathbf{s}$	$(0,0,0,0)$	$(1,0,0,0)$	$(1,1,0,0)$	$(1,1,1,0)$
f_1	R	XR	0	0
f_2	$R+XR$	XR	X^2R	0
f_3	$XR+X^2R$	$XR+X^2R$	X^2R	X^3R
f_4	X^2R+X^3R	X^2R+X^3R	X^2R+X^3R	X^3R
$\mathbf{P}$	$\begin{bmatrix} 1 & & & \\ 0 & 0 & & \\ & & 0 & \\ & & & 0 \end{bmatrix}$	$\begin{bmatrix} 1 & 0 & & \\ 1 & 1 & & \\ 0 & 0 & 0 & \\ & & & 0 \end{bmatrix}$	$\begin{bmatrix} 1 & 0 & & \\ 1 & 1 & 0 & \\ 1 & 1 & 1 & \\ 0 & 0 & 0 & 0 \end{bmatrix}$	

running the iterative algorithm:

i	1	2	3	4	$\cdots$
$\mathbf{s}$	$(0,0,0,0)$	$(1,0,0,0)$	$(1,1,0,0)$	$(1,1,1,0)$	$\cdots$
f_1	R	XR	0	0	0
f_2	$R+XR$	XR	X^2R	0	0
f_3	$XR+X^2R$	$XR+X^2R$	X^2R	X^3R	0
f_4	X^2R+X^3R	X^2R+X^3R	X^2R+X^3R	X^3R	X^4R
$\mathbf{P}$	$\begin{bmatrix}1 & & & \\ 0 & 0 & & \\ & & 0 & \\ & & & 0\end{bmatrix}$	$\begin{bmatrix}1 & 0 & & \\ 1 & 1 & & \\ 0 & 0 & 0 & \\ & & & 0\end{bmatrix}$	$\begin{bmatrix}1 & 0 & & \\ 1 & 1 & 0 & \\ 1 & 1 & 1 & \\ 0 & 0 & 0 & 0\end{bmatrix}$	$\begin{bmatrix}1 & 0 & & \\ 1 & 1 & 0 & \\ 1 & 1 & 1 & 0\\ 1 & 1 & 1 & 1\end{bmatrix}$	$\cdots$

running the iterative algorithm:

i	1	2	3	4	$\cdots$
$\mathbf{s}$	$(0,0,0,0)$	$(1,0,0,0)$	$(1,1,0,0)$	$(1,1,1,0)$	$\cdots$
f_1	R	XR	0	0	0
f_2	$R+XR$	XR	X^2R	0	0
f_3	$XR+X^2R$	$XR+X^2R$	X^2R	X^3R	0
f_4	X^2R+X^3R	X^2R+X^3R	X^2R+X^3R	X^3R	X^4R
$\mathbf{P}$	$\begin{bmatrix} 1 & & & \\ 0 & 0 & & \\ & & 0 & \\ & & & 0 \end{bmatrix}$	$\begin{bmatrix} 1 & 0 & & \\ 1 & 1 & & \\ 0 & 0 & 0 & \\ & & & 0 \end{bmatrix}$	$\begin{bmatrix} 1 & 0 & & \\ 1 & 1 & 0 & \\ 1 & 1 & 1 & \\ 0 & 0 & 0 & 0 \end{bmatrix}$	$\begin{bmatrix} 1 & 0 & & \\ 1 & 1 & 0 & \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 \end{bmatrix}$	$\cdots$

degrees and "pivots" in final basis $\mathbf{P}$: $\begin{bmatrix} 1 & 0 & & \\ 1 & 1 & 0 & \\ 1 & 1 & 1 & 0 \\ 125 & 125 & 125 & 125 \end{bmatrix}$

fast algorithms

iterative algorithm – complexity aspects

parameters: $m = 8, \quad d = 128, \quad \mathbf{s} = (0, 0, 0, 0, d, d, d, d)$

input $\mathbf{F}$: same f_1, f_2, f_3, f_4 / random f_5, f_6, f_7, f_8

$i = 4$

$$\begin{bmatrix}
1 & 0 & & & & & & \\
1 & 1 & 0 & & & & & \\
1 & 1 & 1 & 0 & & & & \\
1 & 1 & 1 & 1 & & & & \\
0 & 0 & 0 & 0 & 0 & & & \\
0 & 0 & 0 & 0 & & 0 & & \\
0 & 0 & 0 & 0 & & & 0 & \\
0 & 0 & 0 & 0 & & & & 0
\end{bmatrix}$$

fast algorithms

iterative algorithm – complexity aspects

parameters: $m = 8$, $d = 128$, $\mathbf{s} = (0, 0, 0, 0, d, d, d, d)$

input $\mathbf{F}$: same f_1, f_2, f_3, f_4 / random f_5, f_6, f_7, f_8

$i = 4$

$$\begin{bmatrix} 1 & 0 & & & & & & \\ 1 & 1 & 0 & & & & & \\ 1 & 1 & 1 & 0 & & & & \\ 1 & 1 & 1 & 1 & & & & \\ 0 & 0 & 0 & 0 & 0 & & & \\ 0 & 0 & 0 & 0 & & 0 & & \\ 0 & 0 & 0 & 0 & & & 0 & \\ 0 & 0 & 0 & 0 & & & & 0 \end{bmatrix}$$

$i = 128$

$$\begin{bmatrix} 1 & 0 & & & & & & \\ 1 & 1 & 0 & & & & & \\ 1 & 1 & 1 & 0 & & & & \\ 125 & 125 & 125 & 125 & & & & \\ 124 & 124 & 124 & 124 & 0 & & & \\ 124 & 124 & 124 & 124 & & 0 & & \\ 124 & 124 & 124 & 124 & & & 0 & \\ 124 & 124 & 124 & 124 & & & & 0 \end{bmatrix}$$

fast algorithms

iterative algorithm – complexity aspects

parameters: $m = 8, \quad d = 128, \quad \mathbf{s} = (0,0,0,0,d,d,d,d)$

input $\mathbf{F}$: same f_1, f_2, f_3, f_4 / random f_5, f_6, f_7, f_8

$i = 4$

$$\begin{bmatrix} 1 & 0 & & & & & & \\ 1 & 1 & 0 & & & & & \\ 1 & 1 & 1 & 0 & & & & \\ 1 & 1 & 1 & 1 & & & & \\ 0 & 0 & 0 & 0 & 0 & & & \\ 0 & 0 & 0 & 0 & & 0 & & \\ 0 & 0 & 0 & 0 & & & 0 & \\ 0 & 0 & 0 & 0 & & & & 0 \end{bmatrix}$$

$i = 128$

$$\begin{bmatrix} 1 & 0 & & & & & & \\ 1 & 1 & 0 & & & & & \\ 1 & 1 & 1 & 0 & & & & \\ 125 & 125 & 125 & 125 & & & & \\ 124 & 124 & 124 & 124 & 0 & & & \\ 124 & 124 & 124 & 124 & & 0 & & \\ 124 & 124 & 124 & 124 & & & 0 & \\ 124 & 124 & 124 & 124 & & & & 0 \end{bmatrix}$$

- $1/4$ of the entries have degree $\approx d$: size $\Theta(m^2 d)$
- opinions on the complexity of iterative algorithm?
- opinions on a “reasonable” target cost for fast algorithms?

base case: modulus of degree 1

modular vector equation

input:

- vector $\mathbf{F} = [f_1 \quad \cdots \quad f_m]^\mathsf{T} \in \mathbb{K}[X]^{m \times 1}$ of degree $< d$
- field elements $(\alpha_1, \ldots, \alpha_d) \in \mathbb{K}^d$
- shift $\boldsymbol{s} = (s_1, \ldots, s_m) \in \mathbb{Z}^m$

output:

matrix $\mathbf{P} \in \mathbb{K}[X]^{m \times m}$ such that

- $\mathbf{PF} = 0 \bmod \prod_{1 \leqslant i \leqslant d} (X - \alpha_i)$
- $\mathbf{P}$ generates all vectors $\mathbf{p}$ such that $\mathbf{pF} = 0 \bmod \prod_{1 \leqslant i \leqslant d} (X - \alpha_i)$
- $\mathbf{P}$ is $\boldsymbol{s}$-reduced

modular vector reconstruction: base case

input:

- vector $\mathbf{F} = [\mathsf{f}_1 \ \ \cdots \ \ \mathsf{f}_\mathsf{m}]^\mathsf{T} \in \mathbb{K}[X]^{\mathsf{m}\times 1}$ of degree < 1
- field element $\alpha \in \mathbb{K}$
- shift $\mathbf{s} = (s_1, \ldots, s_\mathsf{m}) \in \mathbb{Z}^\mathsf{m}$

output:

matrix $\mathbf{P} \in \mathbb{K}[X]^{\mathsf{m}\times\mathsf{m}}$ such that

- $\mathbf{PF} = 0 \bmod (X - \alpha)$
- $\mathbf{P}$ generates all vectors $\mathbf{p}$ such that $\mathbf{pF} = 0 \bmod (X - \alpha)$
- $\mathbf{P}$ is $\mathbf{s}$-reduced

modular vector reconstruction: base case

input:

- vector $\mathbf{F} = [f_1 \ \cdots \ f_m]^\mathsf{T} \in \mathbb{K}[X]^{m\times 1}$ of degree < 1 $\qquad \mathbf{F} \in \mathbb{K}^{m\times 1}$
- field element $\alpha \in \mathbb{K}$
- shift $\mathbf{s} = (s_1, \ldots, s_m) \in \mathbb{Z}^m$

output:

matrix $\mathbf{P} \in \mathbb{K}[X]^{m\times m}$ such that

- $\mathbf{PF} = 0 \bmod (X - \alpha)$ $\qquad (\mathbf{PF})(\alpha) = 0$
- $\mathbf{P}$ generates all vectors $\mathbf{p}$ such that $\mathbf{pF} = 0 \bmod (X - \alpha)$
- $\mathbf{P}$ is $\mathbf{s}$-reduced

modular vector reconstruction: base case

iterative algorithm: $$\mathbf{P} = \begin{bmatrix} \mathbf{I}_{\pi-1} & \boldsymbol{\lambda}_1 & \mathbf{0} \\ \mathbf{0} & X-\alpha & \mathbf{0} \\ \mathbf{0} & \boldsymbol{\lambda}_2 & \mathbf{I}_{m-\pi} \end{bmatrix}$$

where

- π minimizes s_π among indices such that $(\mathbf{p}_\pi \mathbf{F})(\alpha_i) \neq 0$
- the vectors $\boldsymbol{\lambda}_1 \in \mathbb{K}^{(\pi-1)\times 1}$ and $\boldsymbol{\lambda}_2 \in \mathbb{K}^{(m-\pi)\times 1}$ are constant

fast algorithms

base case: modulus of degree 1

modular vector reconstruction: base case

iterative algorithm: $$\mathbf{P} = \begin{bmatrix} \mathbf{I}_{\pi-1} & \boldsymbol{\lambda}_1 & \mathbf{0} \\ \mathbf{0} & X-\alpha & \mathbf{0} \\ \mathbf{0} & \boldsymbol{\lambda}_2 & \mathbf{I}_{m-\pi} \end{bmatrix}$$

where

- π minimizes s_π among indices such that $(\mathbf{p}_\pi \mathbf{F})(\alpha_i) \neq 0$
- the vectors $\boldsymbol{\lambda}_1 \in \mathbb{K}^{(\pi-1)\times 1}$ and $\boldsymbol{\lambda}_2 \in \mathbb{K}^{(m-\pi)\times 1}$ are constant

algorithm:

- $\mathbf{P}$ = identity matrix in $\mathbb{K}[X]^{m\times m}$
- for i from 1 to d:

  a. from the evaluation $\mathbf{F}(\alpha_i)$, find $\mathbf{P}_i$ as above

  b. update shift $s_\pi \leftarrow s_\pi + 1$

  c. update $\mathbf{P} \leftarrow \mathbf{P}_i\mathbf{P}$ as well as $\mathbf{F} \leftarrow \mathbf{P}_i\mathbf{F} \bmod \prod_{i+1\leqslant j\leqslant d}(X-\alpha_j)$

called residual vector

modular vector reconstruction: base case

iterative algorithm: $$\mathbf{P} = \begin{bmatrix} \mathbf{I}_{\pi-1} & \boldsymbol{\lambda}_1 & \mathbf{0} \\ \mathbf{0} & X-\alpha & \mathbf{0} \\ \mathbf{0} & \boldsymbol{\lambda}_2 & \mathbf{I}_{m-\pi} \end{bmatrix}$$

where

- π minimizes s_π among indices such that $(\mathbf{p}_\pi \mathbf{F})(\alpha_i) \neq 0$
- the vectors $\boldsymbol{\lambda}_1 \in \mathbb{K}^{(\pi-1)\times 1}$ and $\boldsymbol{\lambda}_2 \in \mathbb{K}^{(m-\pi)\times 1}$ are constant

complexity $O(m^2 d^2)$:

- iteration with d steps
- each step: evaluation of $\mathbf{F}$ + multiplications $\mathbf{P}_i\mathbf{F}$ and $\mathbf{P}_i\mathbf{P}$
- at any stage $\mathbf{F}$ has degree $< d$ and size $m \times 1$
- at any stage $\mathbf{P}$ has degree $\leqslant d$ and size $m \times m$

normalizing at each step + refined analysis yields $O(md^2)$

fast algorithms

base case: modulus of degree 1

modular vector reconstruction: base case

iterative algorithm: $$\mathbf{P} = \begin{bmatrix} \mathbf{I}_{\pi-1} & \boldsymbol{\lambda}_1 & \mathbf{0} \\ \mathbf{0} & X - \alpha & \mathbf{0} \\ \mathbf{0} & \boldsymbol{\lambda}_2 & \mathbf{I}_{m-\pi} \end{bmatrix}$$

where

- π minimizes s_π among indices such that $(\mathbf{p}_\pi \mathbf{F})(\alpha_i) \neq 0$
- the vectors $\boldsymbol{\lambda}_1 \in \mathbb{K}^{(\pi-1)\times 1}$ and $\boldsymbol{\lambda}_2 \in \mathbb{K}^{(m-\pi)\times 1}$ are constant

correctness:

- the main task is to prove the base case with $\mathbf{P}_i$
- then, direct consequence of the "basis multiplication theorem"

divide and conquer algorithm:

input: $\mathbf{F}, (\alpha_1, \ldots, \alpha_d), \mathbf{s}$ | output: $\mathbf{P}$

- if $d = 1$, use the base case algorithm to find $\mathbf{P}$ and return
- otherwise:

a. $M_1 \leftarrow (X - x_1) \cdots (X - x_{\lfloor d/2 \rfloor})$; $M_2 \leftarrow (X - x_{\lceil d/2 \rceil} \cdots (X - x_d)$

b. $\mathbf{P}_1 \leftarrow$ call the algorithm on $\mathbf{F}$ rem $M_1, (\alpha_1, \ldots, \alpha_{\lfloor d/2 \rfloor}), \mathbf{s}$

c. updated shift: $\mathbf{t} \leftarrow \mathrm{rdeg}_{\mathbf{s}}(\mathbf{P}_1)$

d. residual: $\mathbf{G} \leftarrow \frac{1}{M_1}\mathbf{P}_1\mathbf{F}$

e. $\mathbf{P}_2 \leftarrow$ call the algorithm on $\mathbf{G}$ rem $M_2, (\alpha_{\lceil d/2 \rceil}, \ldots, \alpha_d), \mathbf{t}$

f. return the product $\mathbf{P}_2\mathbf{P}_1$

fast algorithms

recursion: residual and basis multiplication

divide and conquer algorithm:

input: $\mathbf{F}, (\alpha_1, \ldots, \alpha_d), \mathbf{s}$ | output: $\mathbf{P}$

- if $d = 1$, use the base case algorithm to find $\mathbf{P}$ and return
- otherwise:

  a. $M_1 \leftarrow (X - x_1) \cdots (X - x_{\lfloor d/2 \rfloor})$; $M_2 \leftarrow (X - x_{\lceil d/2 \rceil} \cdots (X - x_d)$

  b. $\mathbf{P}_1 \leftarrow$ call the algorithm on $\mathbf{F}$ rem $M_1, (\alpha_1, \ldots, \alpha_{\lfloor d/2 \rfloor}), \mathbf{s}$

  c. updated shift: $\mathbf{t} \leftarrow \mathrm{rdeg}_{\mathbf{s}}(\mathbf{P}_1)$

  d. residual: $\mathbf{G} \leftarrow \frac{1}{M_1}\mathbf{P}_1\mathbf{F}$

  e. $\mathbf{P}_2 \leftarrow$ call the algorithm on $\mathbf{G}$ rem $M_2, (\alpha_{\lceil d/2 \rceil}, \ldots, \alpha_d), \mathbf{t}$

  f. return the product $\mathbf{P}_2\mathbf{P}_1$

complexity $O(m^\omega M(d) \log(d))$:

- d calls to the base case amounts to $O(m^2 d)$
- most expensive step in the recursion is the product $\mathbf{P}_2\mathbf{P}_1$
- hence: $\mathcal{C}(m, d) = \mathcal{C}(m, \lfloor d/2 \rfloor) + \mathcal{C}(m, \lceil d/2 \rceil) + O(m^\omega M(d))$

fast algorithms

recursion: residual and basis multiplication

divide and conquer algorithm:

input: $\mathbf{F}, (\alpha_1, \ldots, \alpha_d), \mathbf{s}$ | output: $\mathbf{P}$

- if $d = 1$, use the base case algorithm to find $\mathbf{P}$ and return
- otherwise:

  a. $M_1 \leftarrow (X - x_1) \cdots (X - x_{\lfloor d/2 \rfloor})$; $M_2 \leftarrow (X - x_{\lceil d/2 \rceil} \cdots (X - x_d)$

  b. $\mathbf{P}_1 \leftarrow$ call the algorithm on $\mathbf{F}$ rem $M_1, (\alpha_1, \ldots, \alpha_{\lfloor d/2 \rfloor}), \mathbf{s}$

  c. updated shift: $\mathbf{t} \leftarrow \mathrm{rdeg}_{\mathbf{s}}(\mathbf{P}_1)$

  d. residual: $\mathbf{G} \leftarrow \frac{1}{M_1}\mathbf{P}_1\mathbf{F}$

  e. $\mathbf{P}_2 \leftarrow$ call the algorithm on $\mathbf{G}$ rem $M_2, (\alpha_{\lceil d/2 \rceil}, \ldots, \alpha_d), \mathbf{t}$

  f. return the product $\mathbf{P}_2\mathbf{P}_1$

correctness:

- correctness of base case
- then, direct consequence of the "basis multiplication theorem"
- about the residual: $\{\mathbf{p} \mid \mathbf{p}\mathbf{P}_1\mathbf{F} = 0 \bmod M\} = \{\mathbf{p} \mid \mathbf{p}\mathbf{G} = 0 \bmod M_2\}$

state of the art:

- the above algorithm is in [Beckermann-Labahn 1994] (for Hermite-Padé)
- it also works for $\mathbf{F} \in \mathbb{K}[X]^{m\times n}$ with $n > 1$
- [Giorgi-Jeannerod-Villard 2003] obtained the complexity $O(m^{\omega} \mathsf{M}(d) \log(d))$ for the case $n \geqslant 1$
- today, $\tilde{O}(m^{\omega} \frac{nd}{m})$ has been reached for the general case [Storjohann 2006] [Zhou-Labahn 2012] [Jeannerod-Neiger-Villard 2020]

outline

- **introduction**
  - rational approximation and interpolation
  - the vector case
  - pol. matrices: reminders and motivation
- **shifted reduced forms**
  - reducedness: examples and properties
  - shifted forms and degree constraints
  - stability under multiplication
- **fast algorithms**
  - iterative algorithm and output size
  - base case: modulus of degree 1
  - recursion: residual and basis multiplication
- **next time**

outline

introduction

- rational approximation and interpolation
- the vector case
- pol. matrices: reminders and motivation

shifted reduced forms

- reducedness: examples and properties
- shifted forms and degree constraints
- stability under multiplication

fast algorithms

- iterative algorithm and output size
- base case: modulus of degree 1
- recursion: residual and basis multiplication

next time

- fast extended gcd
- block Wiedemann techniques
- minimal kernel bases and row bases

summary

introduction

- rational approximation and interpolation
- the vector case
- pol. matrices: reminders and motivation

shifted reduced forms

- reducedness: examples and properties
- shifted forms and degree constraints
- stability under multiplication

fast algorithms

- iterative algorithm and output size
- base case: modulus of degree 1
- recursion: residual and basis multiplication

next time

- fast extended gcd
- block Wiedemann techniques
- minimal kernel bases and row bases