



# JavaScript Static Analysis for Evolving Language Specifications

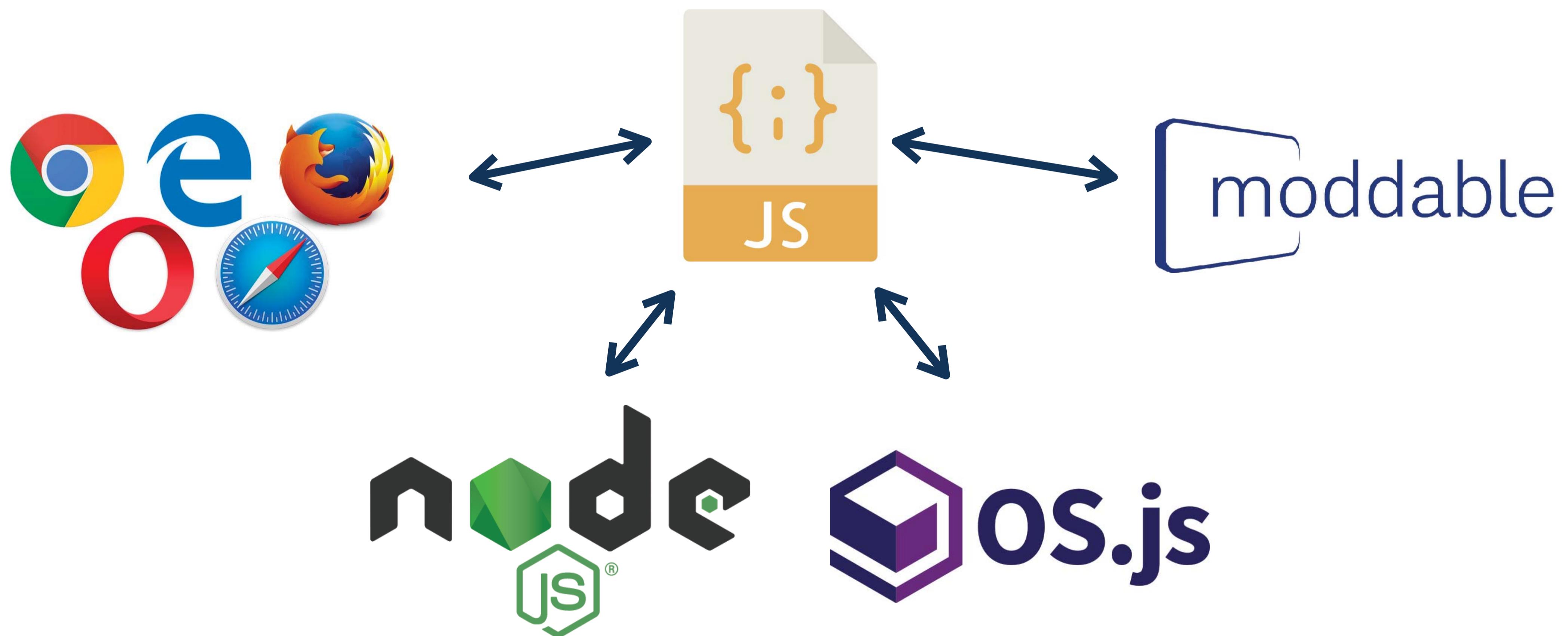
[SW재난연구센터] 겨울정기워크샵

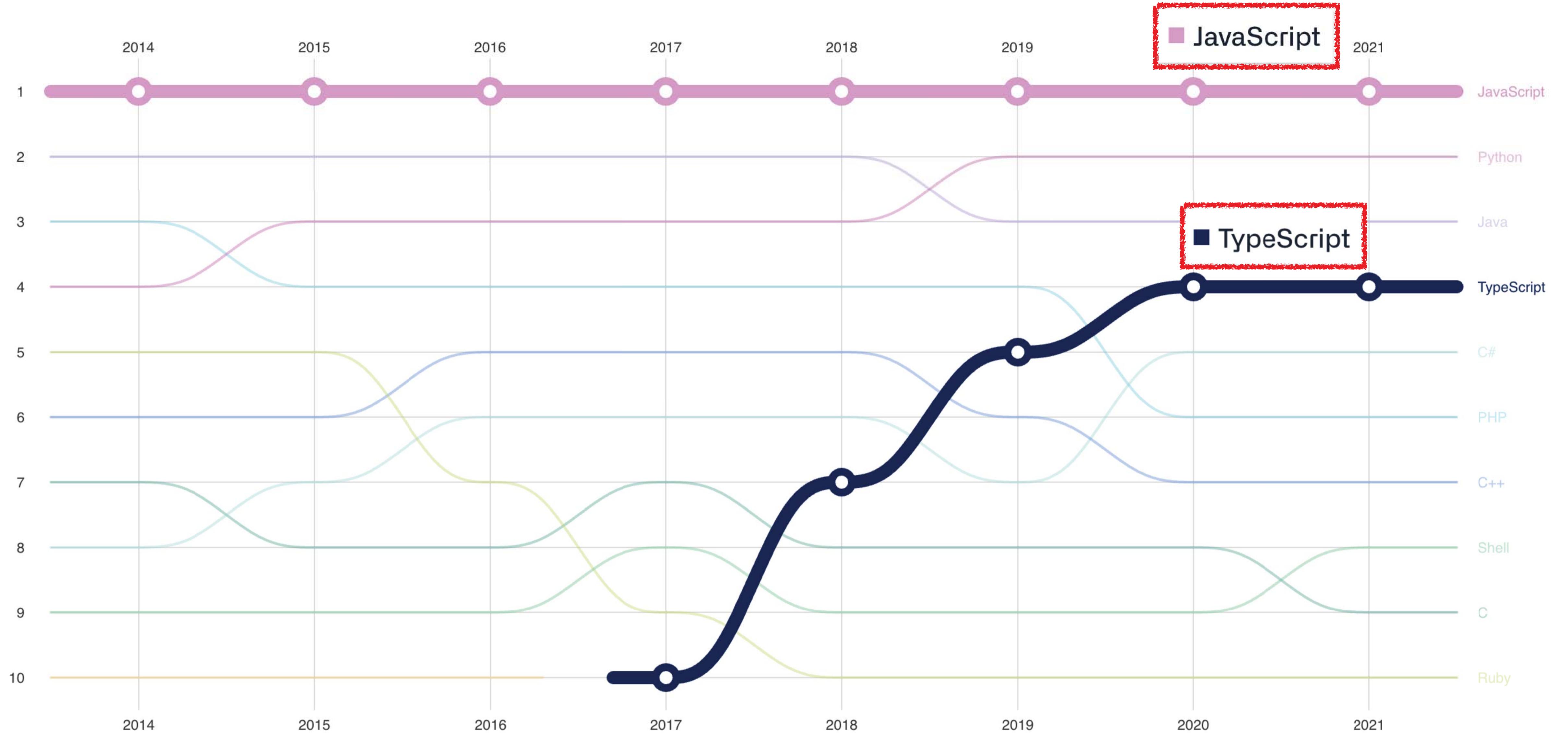
**Jihyeok Park**

PLRG @ KAIST

February 9, 2022

# JavaScript Is Everywhere





<https://octoverse.github.com/>

# JavaScript Complex Semantics

```
function f(x) { return x == !x; }
```

Always return **false**?

# JavaScript Complex Semantics

```
function f(x) { return x == !x; }
```

Always return **false**?

**NO!!**

```
f( []) -> [] == ![]
-> [] == false
-> +[] == +false
-> 0 == 0
-> true
```

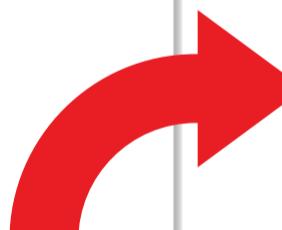
# ECMA-262: ECMAScript Specification



Semantics

Syntax

ArrayLiteral [Yield, Await] :  
[ Elision<sub>opt</sub> ]  
[ ElementList [?Yield, ?Await] ]  
[ ElementList [?Yield, ?Await] , Elision<sub>opt</sub> ]



## 13.2.5.2 Runtime Semantics: Evaluation

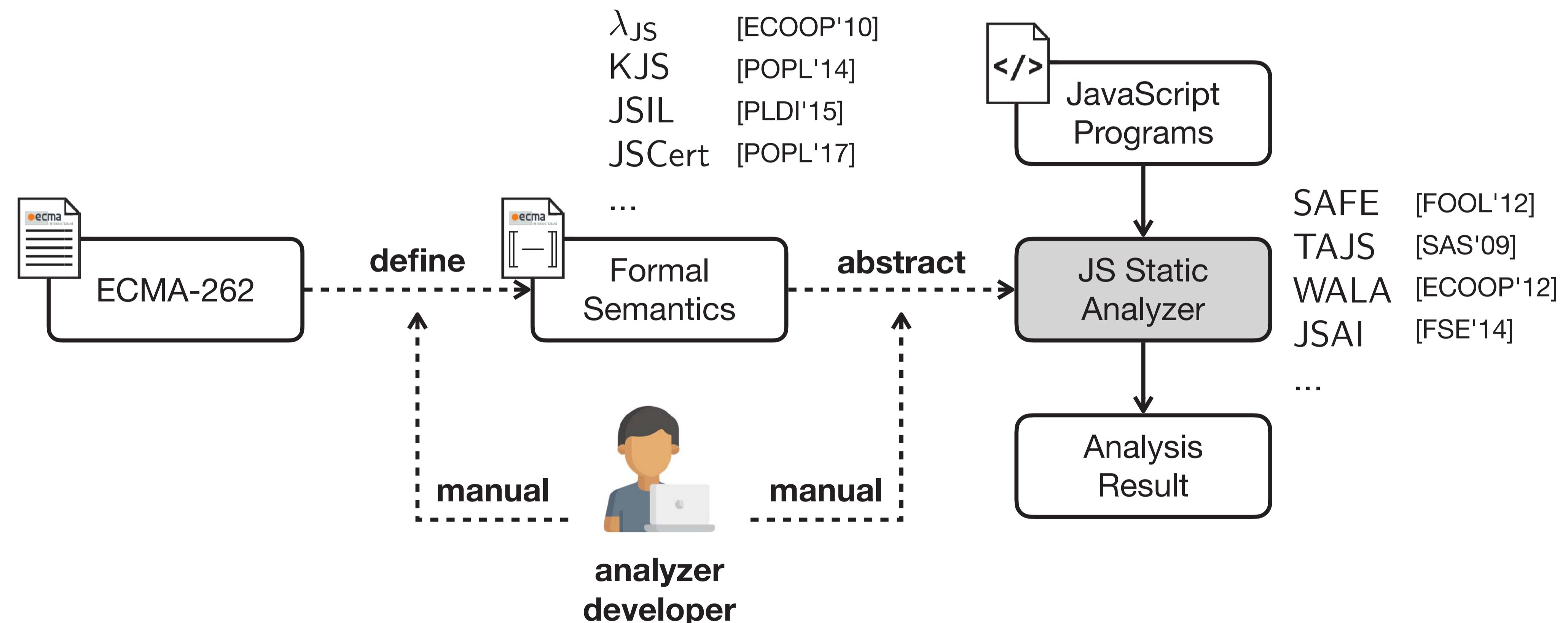
ArrayLiteral : [ ElementList , Elision<sub>opt</sub> ]

1. Let *array* be ! *ArrayCreate*(0).
2. Let *nextIndex* be the result of performing *ArrayAccumulation* for *ElementList* with arguments *array* and 0.
3. *ReturnIfAbrupt*(*nextIndex*).
4. If *Elision* is present, then
  - a. Let *len* be the result of performing *ArrayAccumulation* for *Elision* with arguments *array* and *nextIndex*.
  - b. *ReturnIfAbrupt*(*len*).
5. Return *array*.

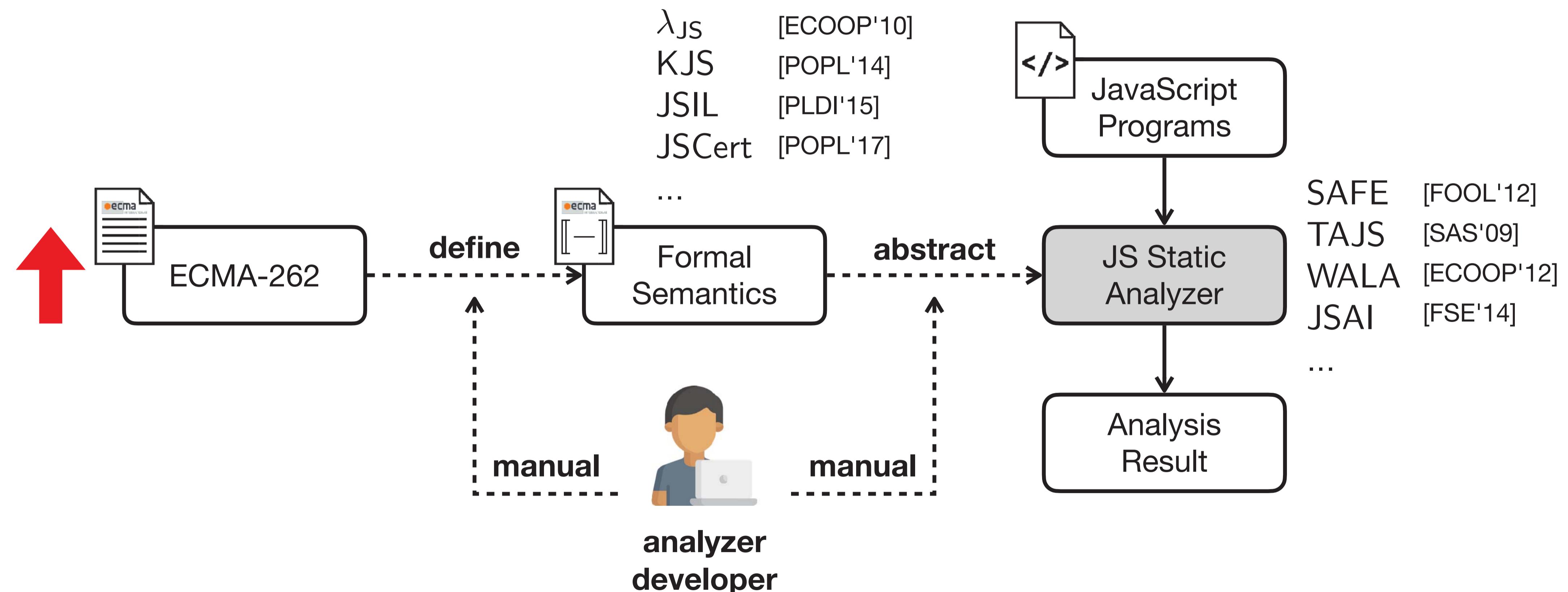
The production of *ArrayLiteral* in ES12

The Evaluation algorithm for  
the third alternative of *ArrayLiteral* in ES12

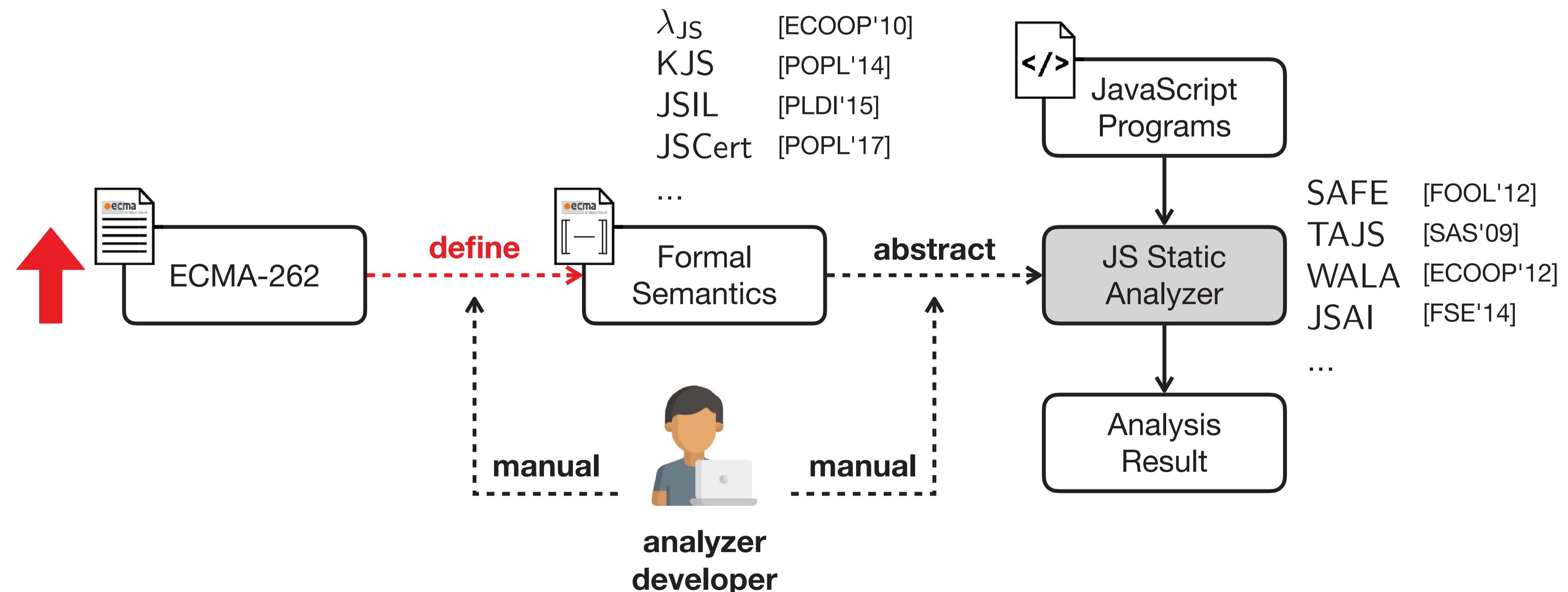
# Problem: Manual JavaScript Static Analyzer



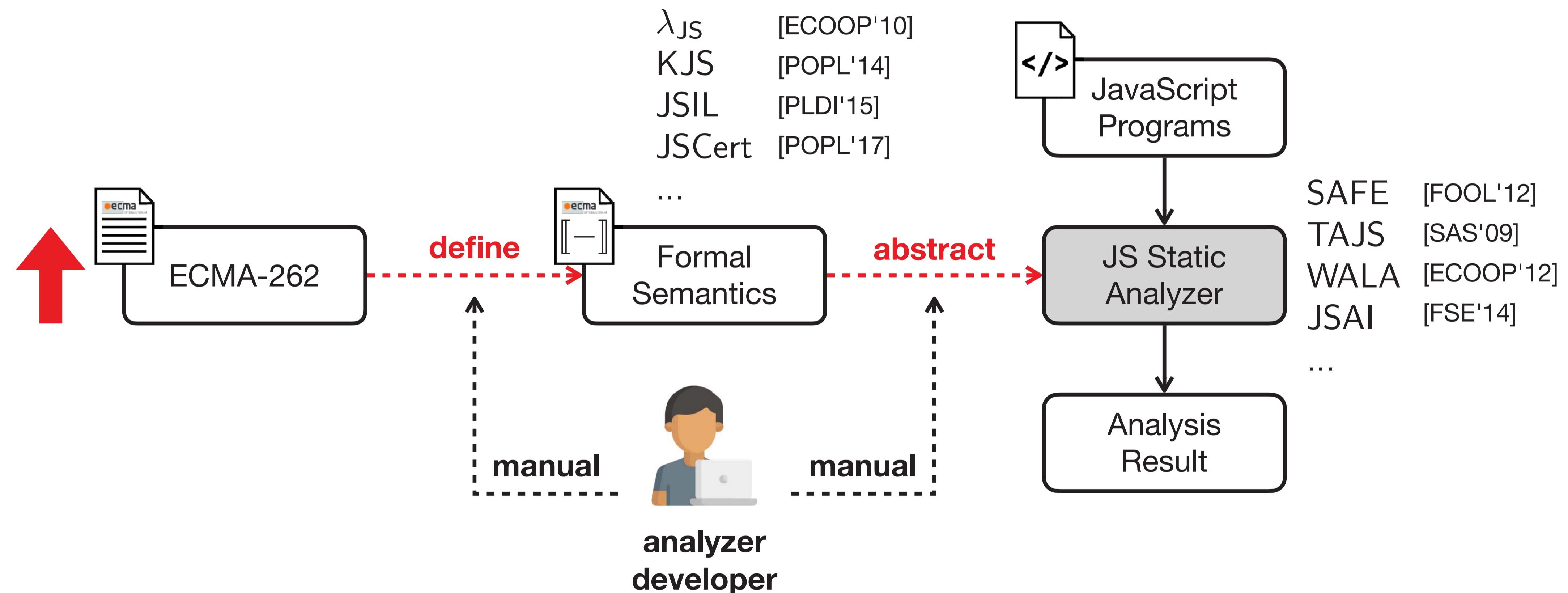
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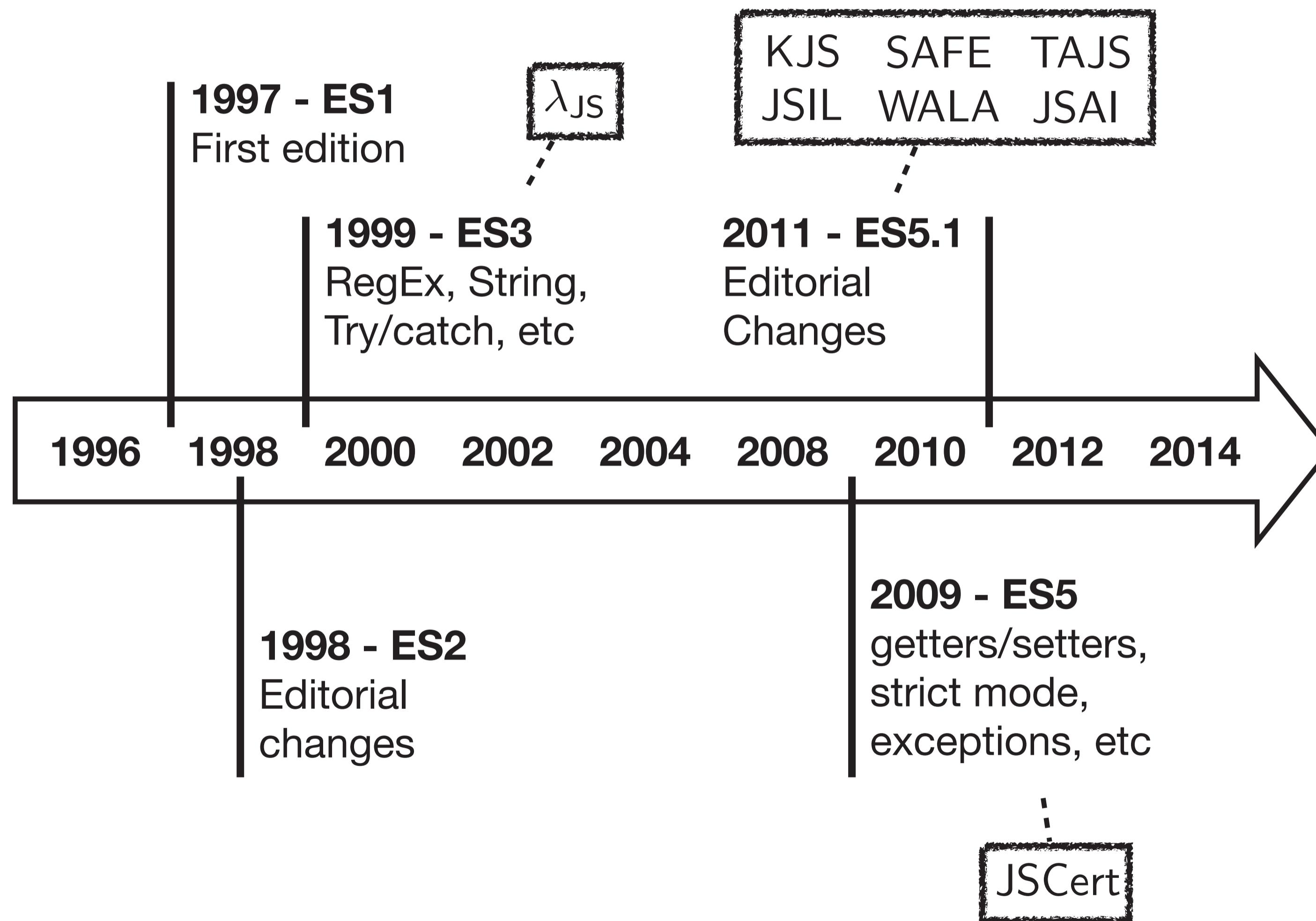
# Problem: Manual JavaScript Static Analyzer



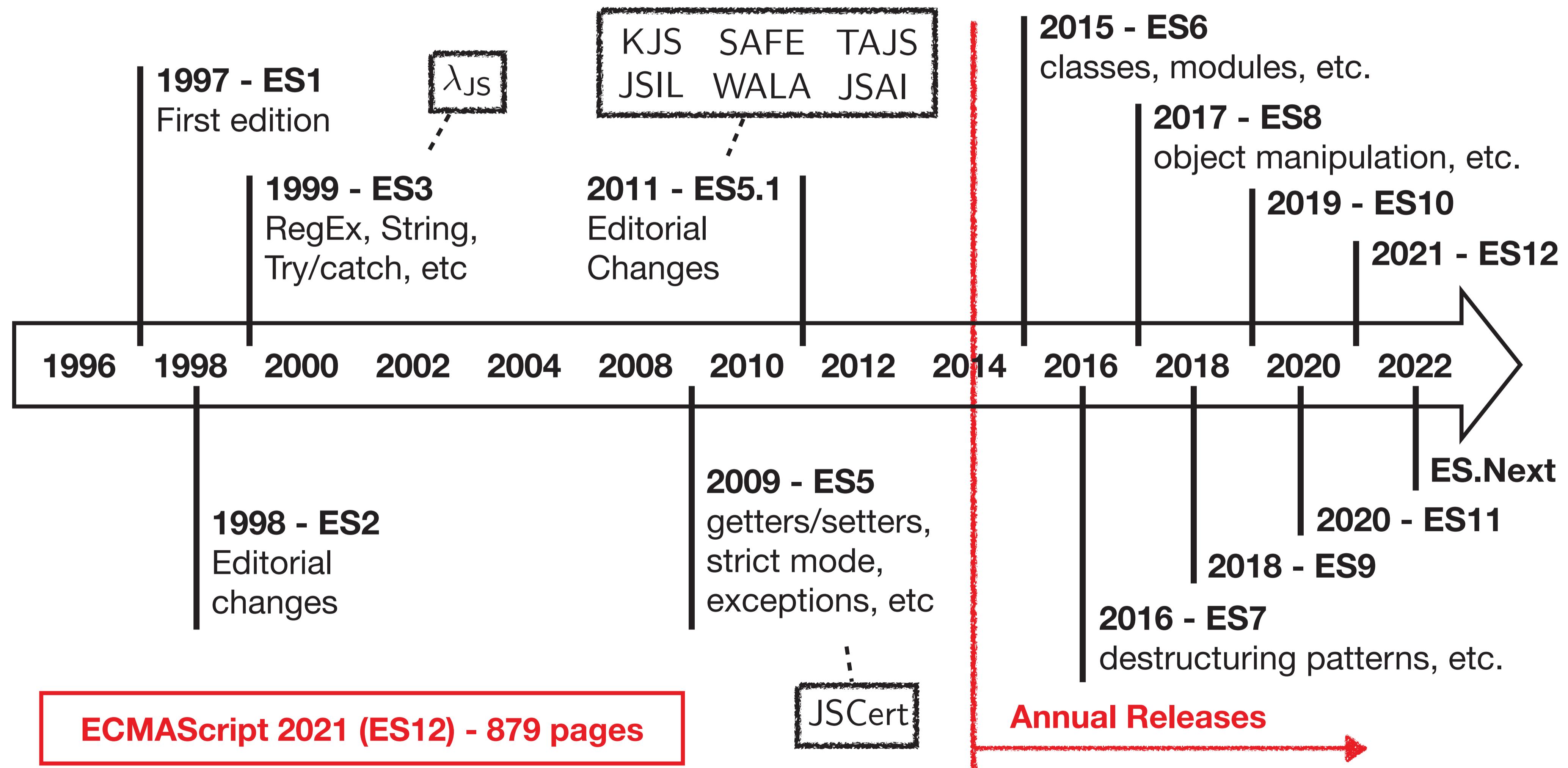
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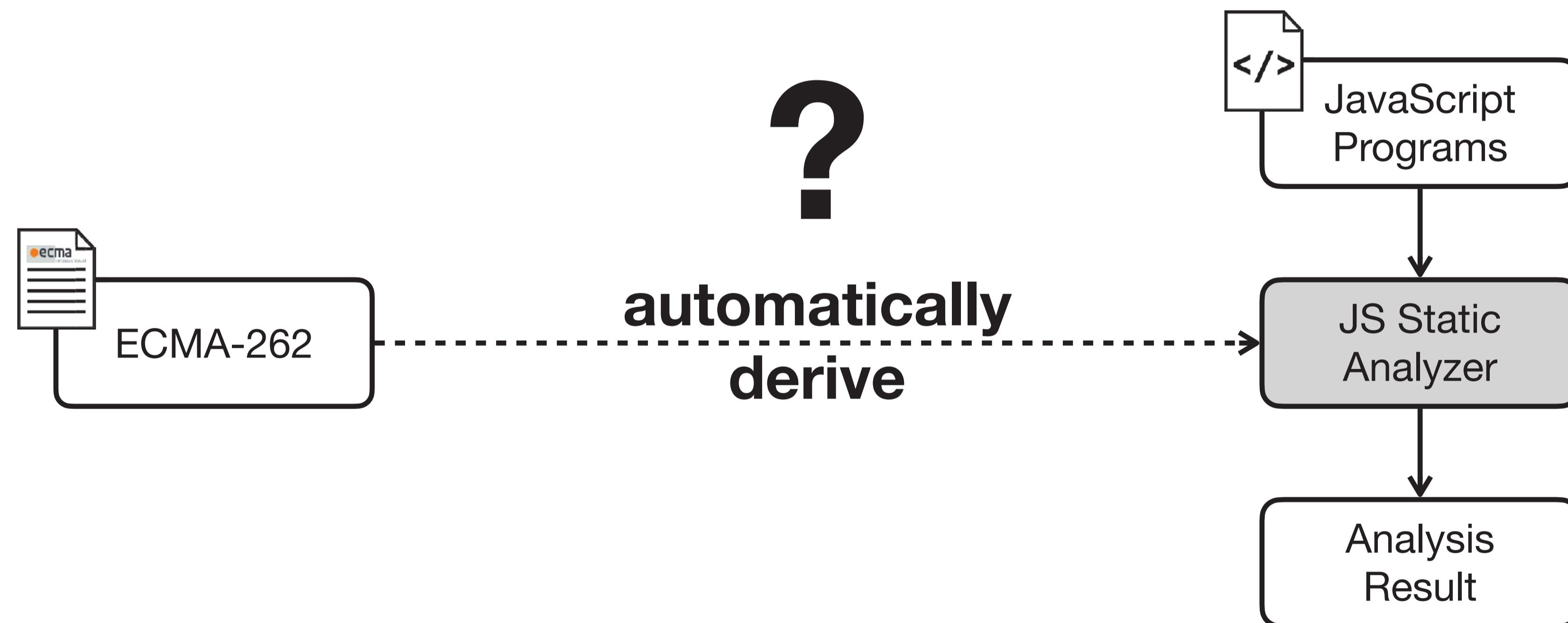
# Problem: Fast Evolving JavaScript



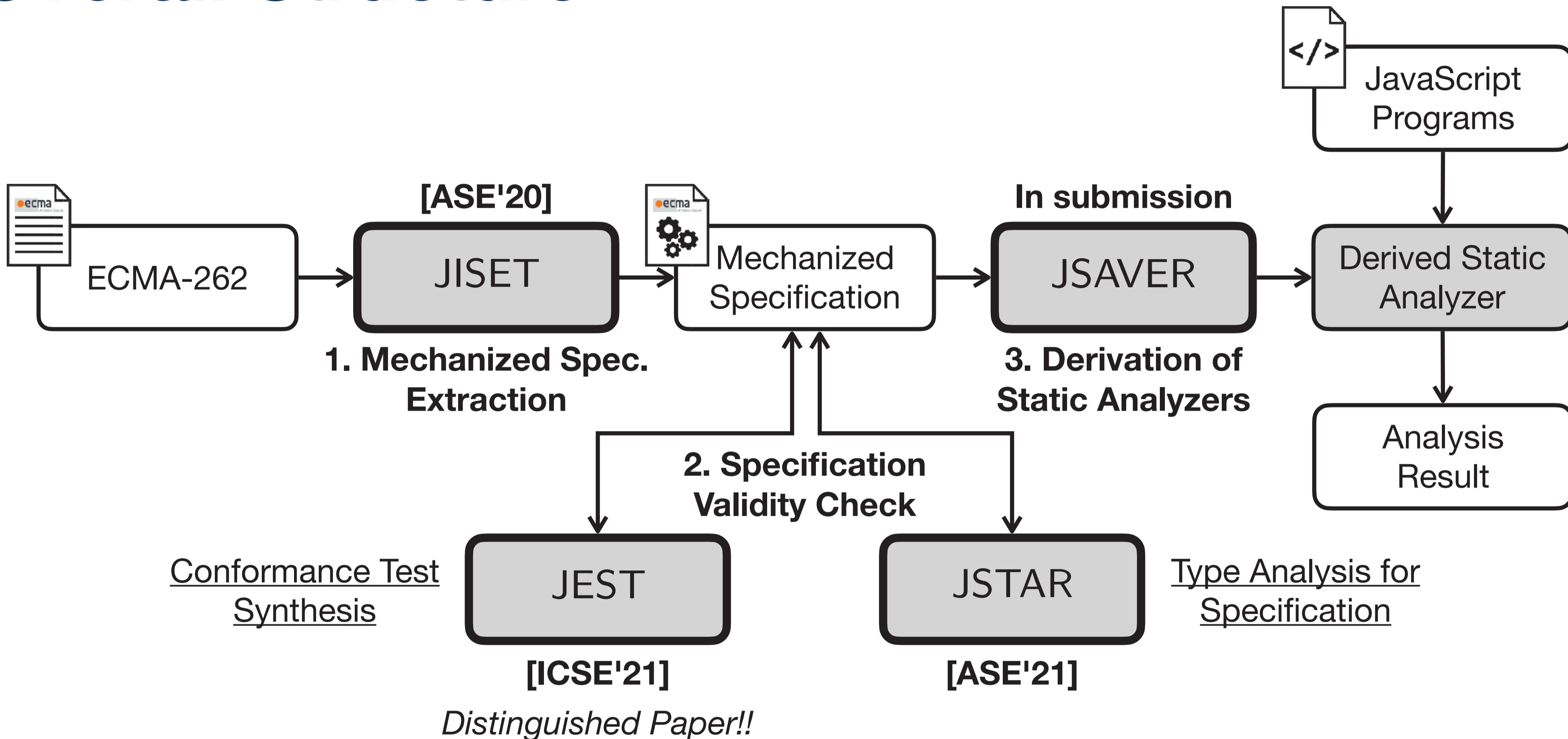
# Problem: Fast Evolving JavaScript



# Main Idea: Deriving Static Analyzer from Spec.

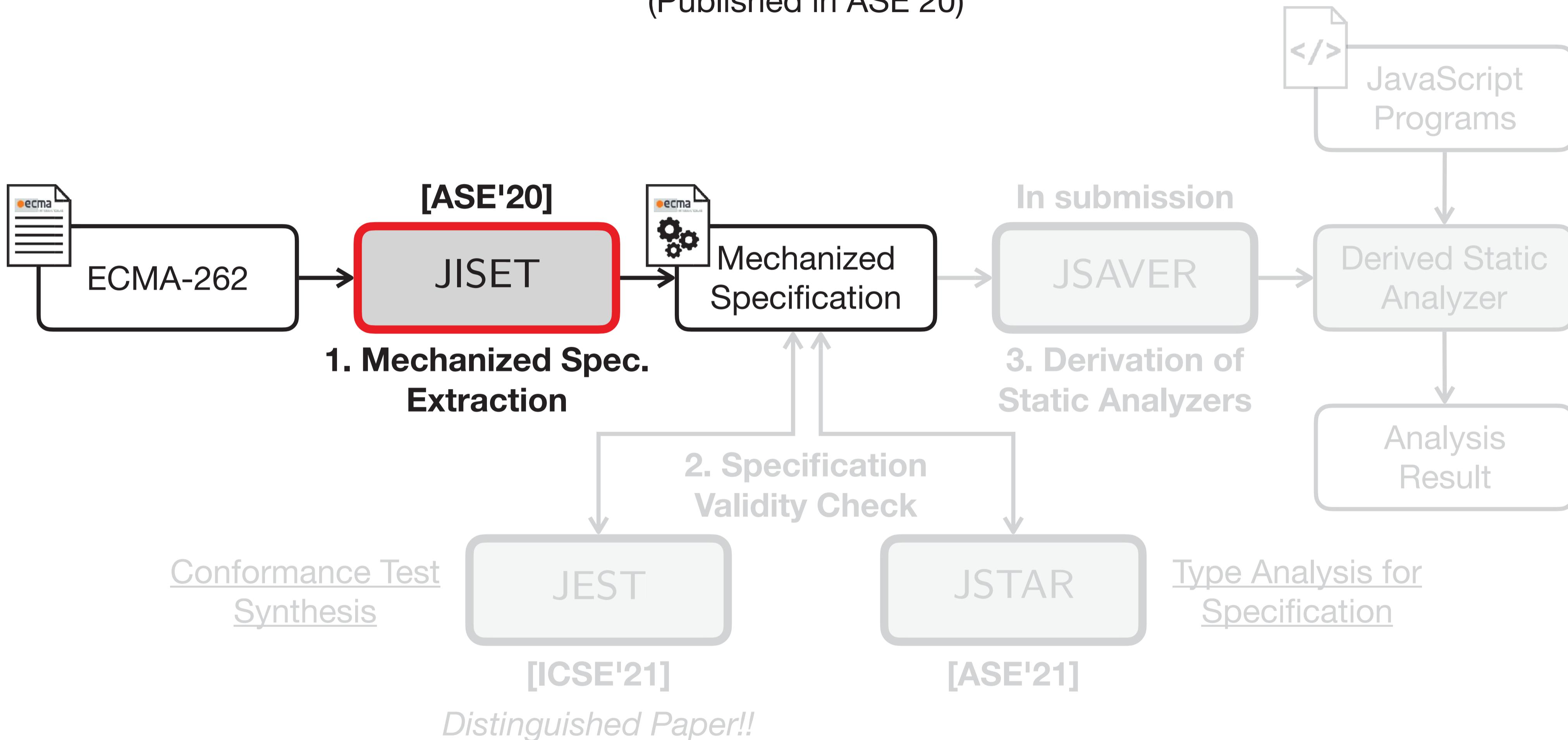


# Overall Structure



# JISET: JavaScript IR-based Semantics Extraction Toolchain

Jihyeok Park, Jihee Park, Seungmin An, and Sukyoung Ryu  
(Published in ASE'20)



# Motivation: Patterns in Writing Style of ECMA-262

## 13.2.5.2 Runtime Semantics: Evaluation

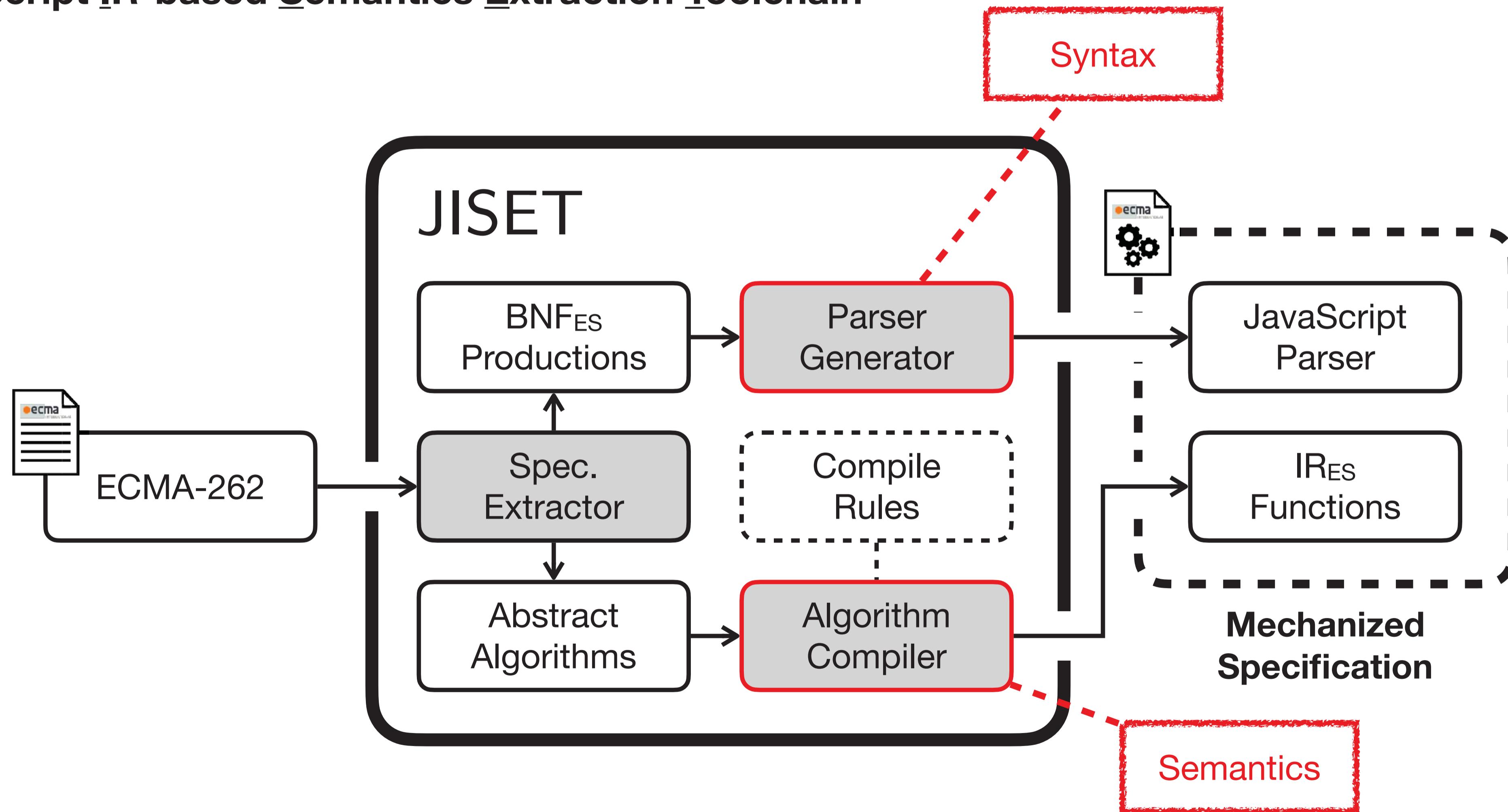
*ArrayLiteral* : [ *ElementList* , *Elision<sub>opt</sub>* ]

1. Let *array* be ! *ArrayCreate*(0).
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  - b. *ReturnIfAbrupt*(*len*).
5. Return *array*.

The Evaluation algorithm for  
the third alternative of *ArrayLiteral* in ES12

# JISET [ASE'20]

## JavaScript IR-based Semantics Extraction Toolchain



# JSET - Parser Generator (Syntax)

```
ArrayLiteral[Yield, Await] :  
  [ Elisionopt ]  
  [ ElementList[?Yield, ?Await] ]  
  [ ElementList[?Yield, ?Await] , Elisionopt ]
```

JavaScript Parser  
in Scala

Parsing Expression Grammar  
(+ Lookahead Parsing)

```
val ArrayLiteral: List[Boolean] => LAParser[T] = memo {  
  case List(Yield, Await) =>  
    "[" ~ opt(Elision) ~ "]" ^^ ArrayLiteral0 |  
    "[" ~ ElementList(Yield, Await) ~ "]" ^^ ArrayLiteral1 |  
    "[" ~ ElementList(Yield, Await) ~ ";" ~ opt(Elision) ~ "]" ^^ ArrayLiteral2  
}
```

(POPL'04) Bryan Ford, "Parsing Expression Grammars: A Recognition-based Syntactic Foundation"

# JSET - Algorithm Compiler (Semantics)

## 13.2.5.2 Runtime Semantics: Evaluation

*ArrayLiteral* : [ *ElementList* , *Elision*<sub>opt</sub> ]

1. Let *array* be ! *ArrayCreate*(0).
2. Let *nextIndex* be the result of performing *ArrayAccumulation* for *ElementList* with arguments *array* and 0.
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4. If *Elision* is present, then
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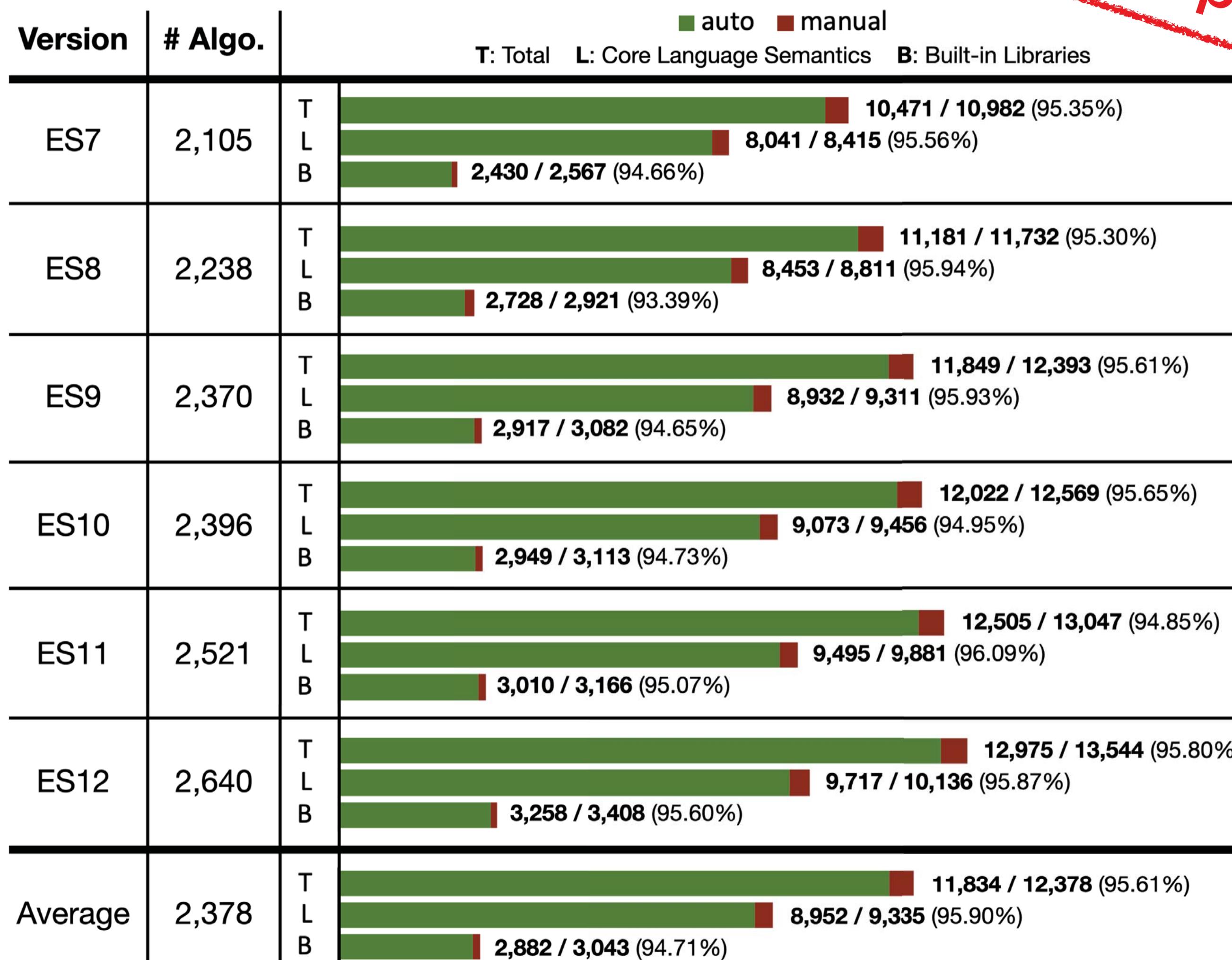
118 Compile Rules for  
Steps in Abstract Algorithms

```
syntax def ArrayLiteral[2].Evaluation(
    this, ElementList, Elision
) {
    let array = [! (ArrayCreate 0)]
    let nextIndex = (ElementList.ArrayAccumulation array 0)
    [? nextIndex]
    if (! (= Elision absent)) {
        let len = (Elision.ArrayAccumulation array nextIndex)
        [? len]
    }
    return array
}
```

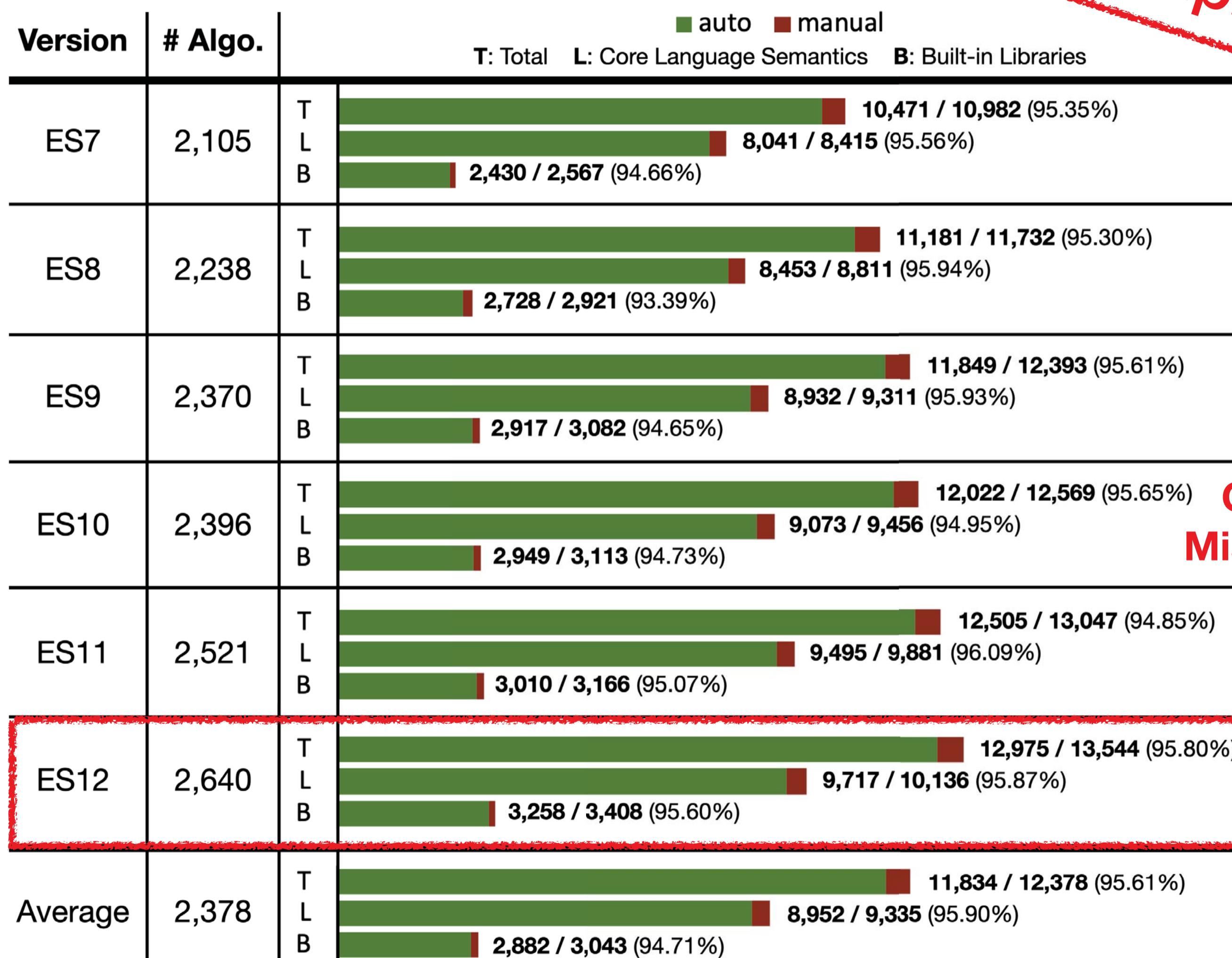
IRES  
Functions

# JISET - Evaluation

≈ 95%  
Compiled



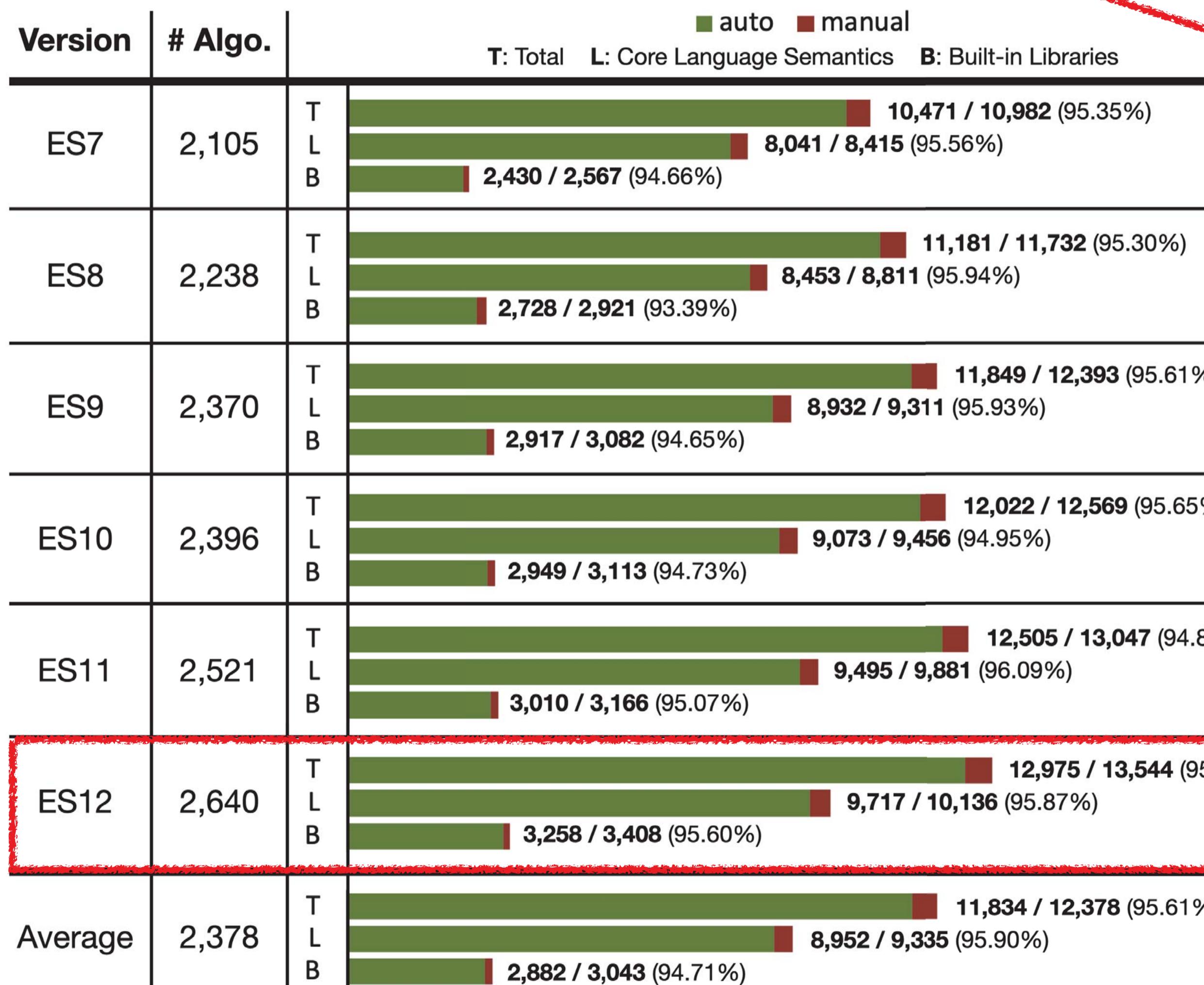
# JISET - Evaluation



$\approx 95\%$   
Compiled

Complete  
Missing Parts

# JISET - Evaluation



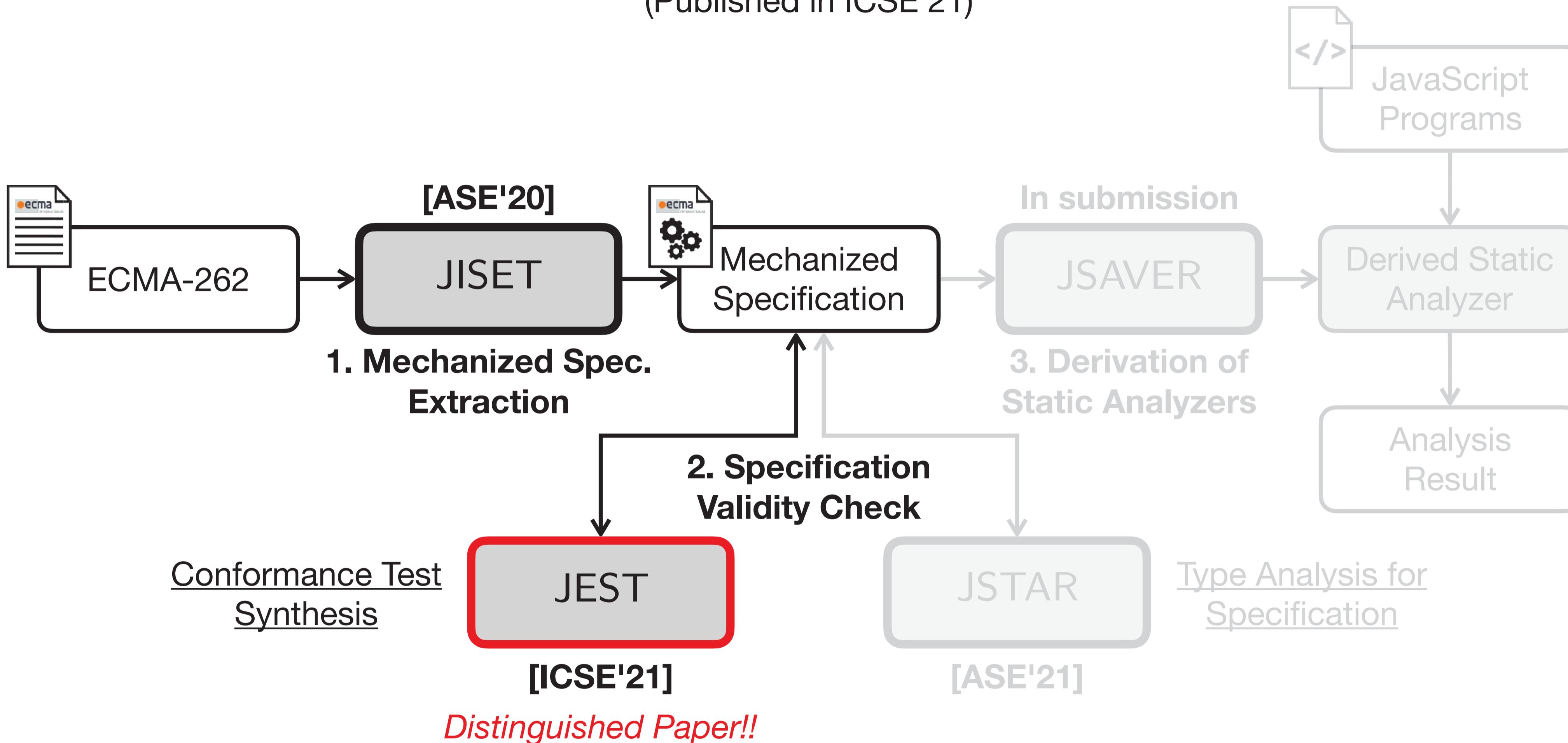
≈ 95%  
Compiled

Passed  
All Tests

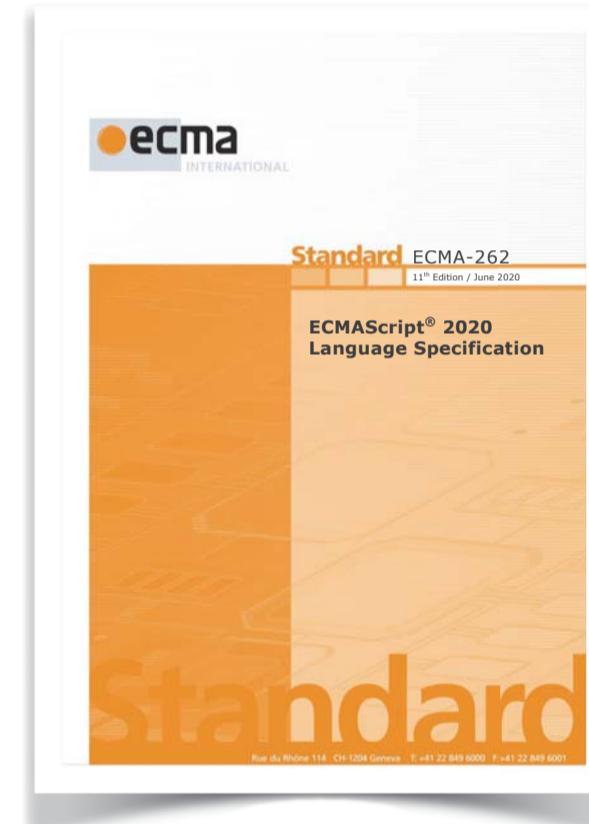
- **Test262**  
(Official Conformance Tests)
  - 18,556 applicable tests
- **Parsing tests**
  - Passed all 18,556 tests
- **Evaluation Tests**
  - Passed all 18,556 tests

# JEST: N+1-version Differential Testing of Both JavaScript Engines

Jihyeok Park, Seungmin An, Dongjun Youn, Gyeongwon Kim, and Sukyoung Ryu  
(Published in ICSE'21)



# JEST - Conformance with Engines



ECMA-262



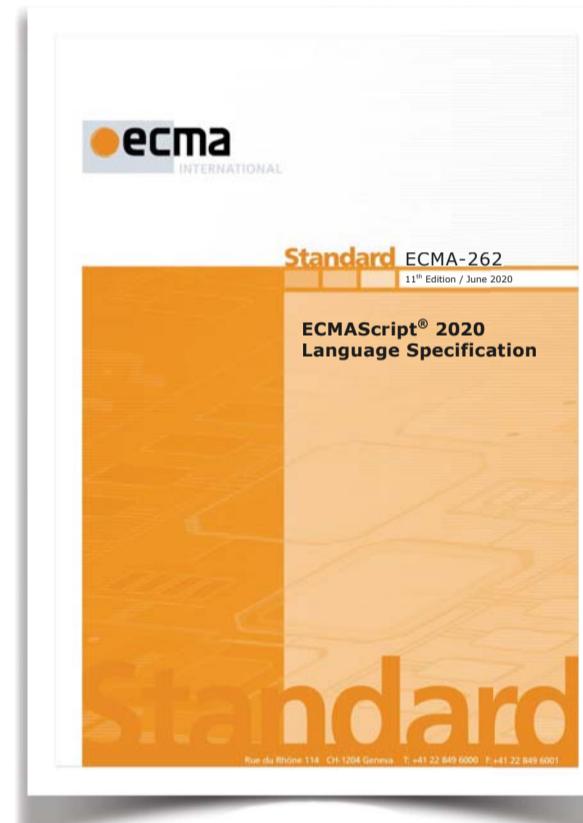
GraalVM™

QuickJS

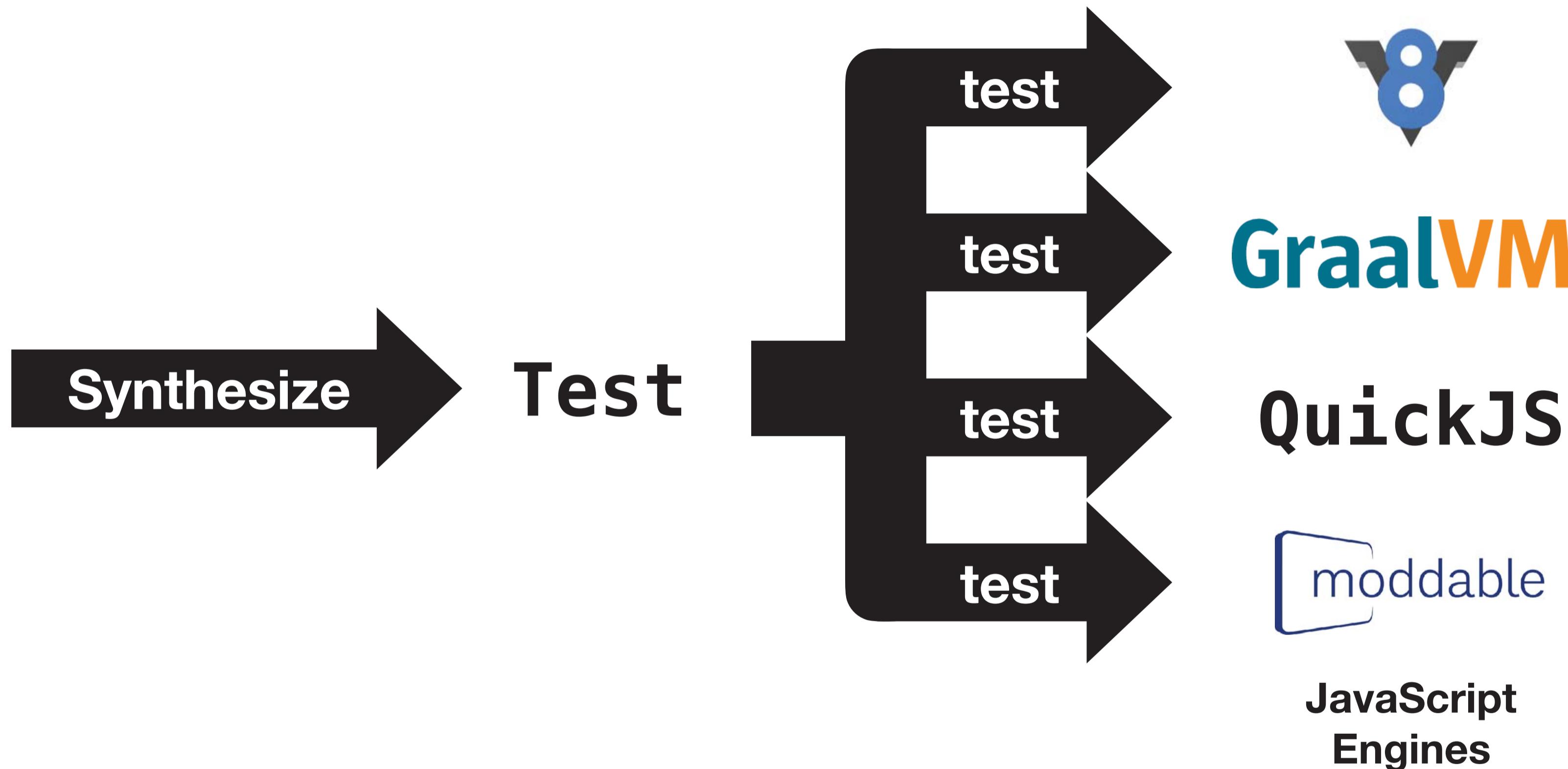


JavaScript  
Engines

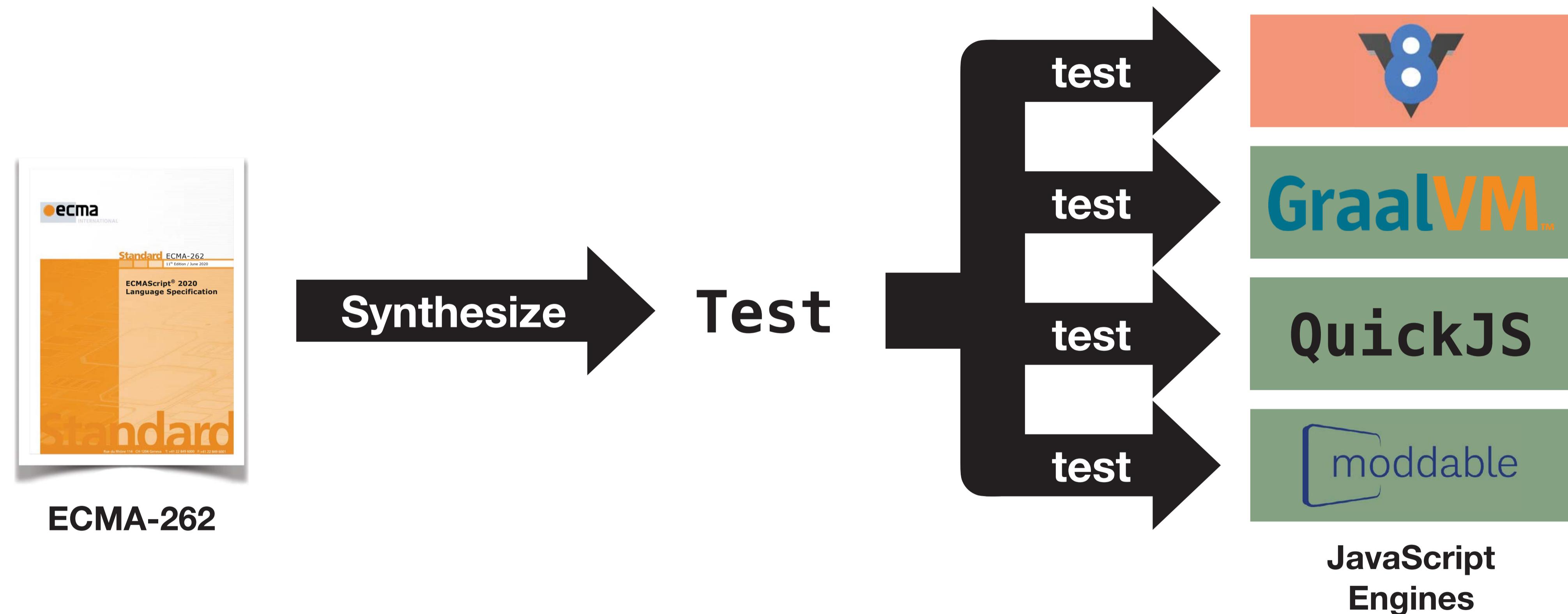
# JEST - N+1-version Differential Testing



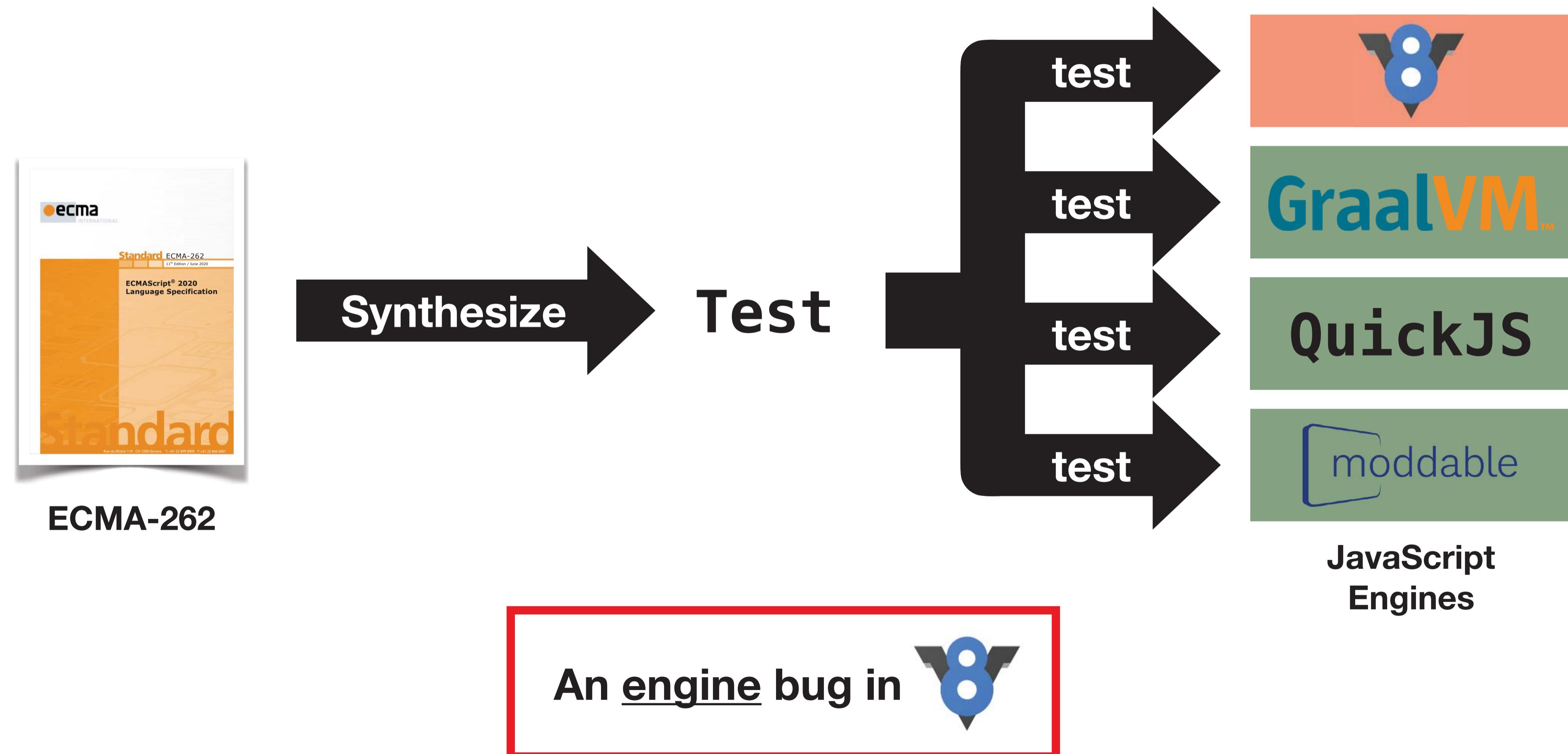
ECMA-262



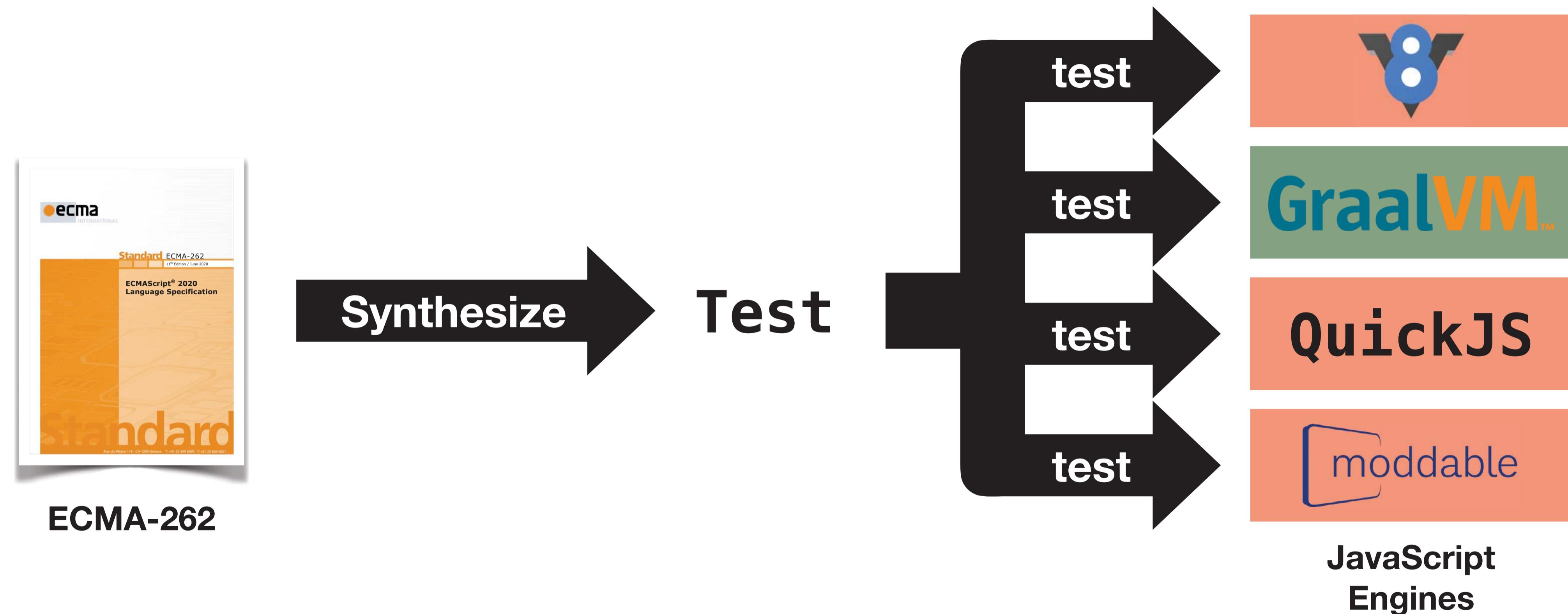
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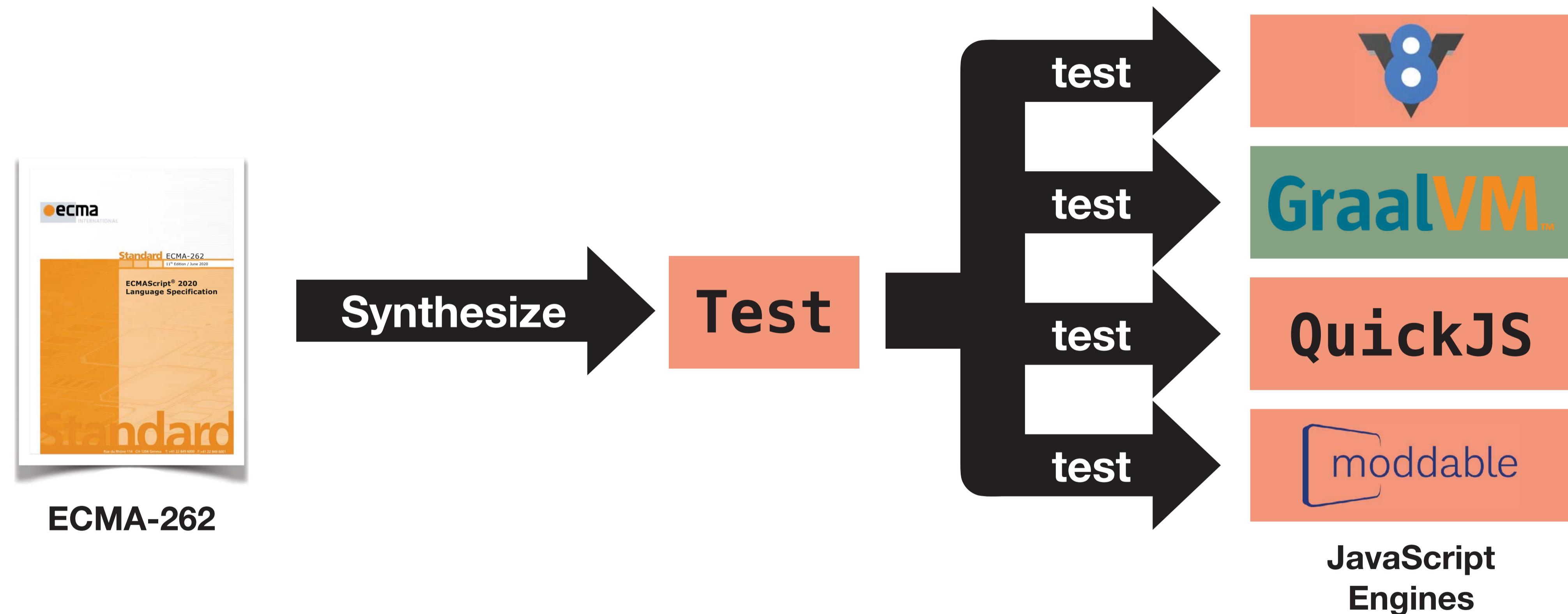
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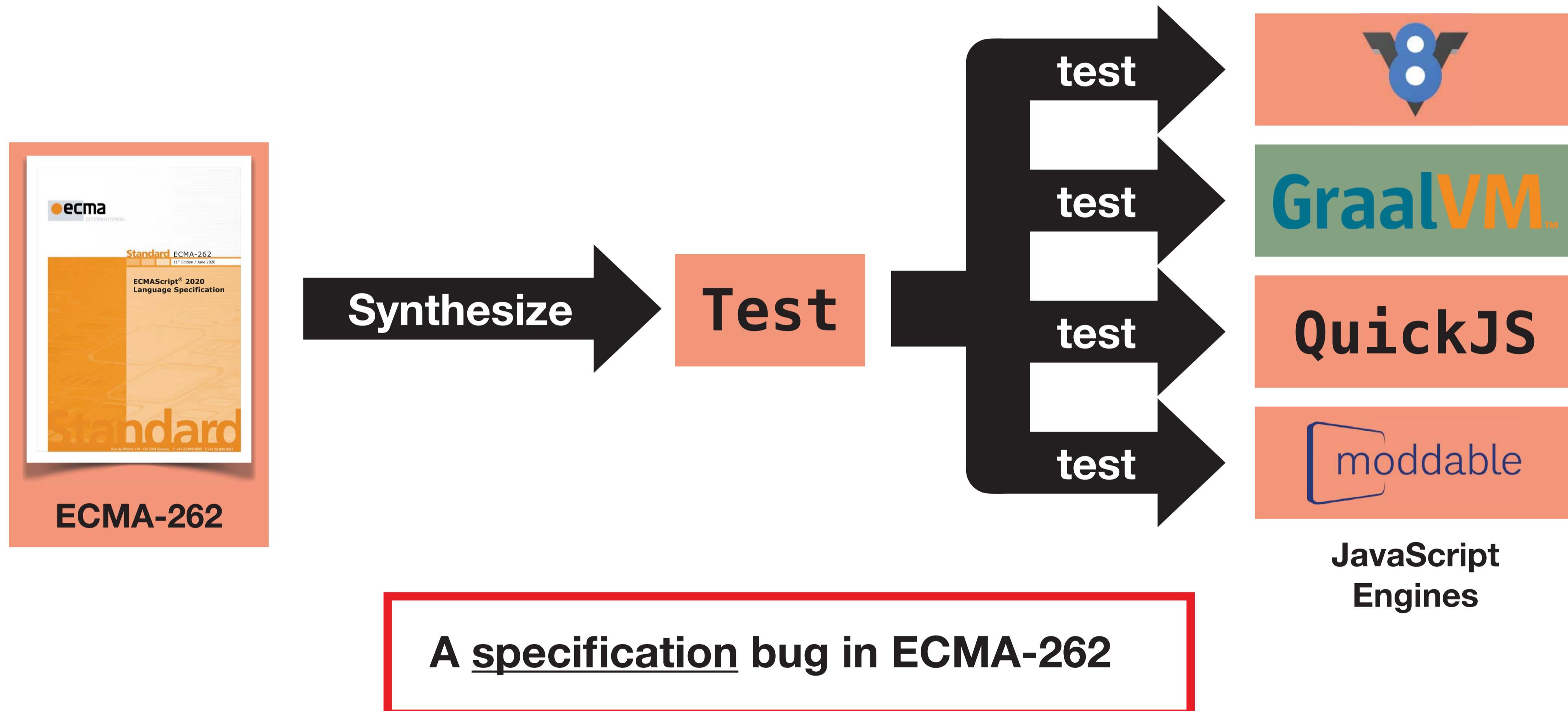
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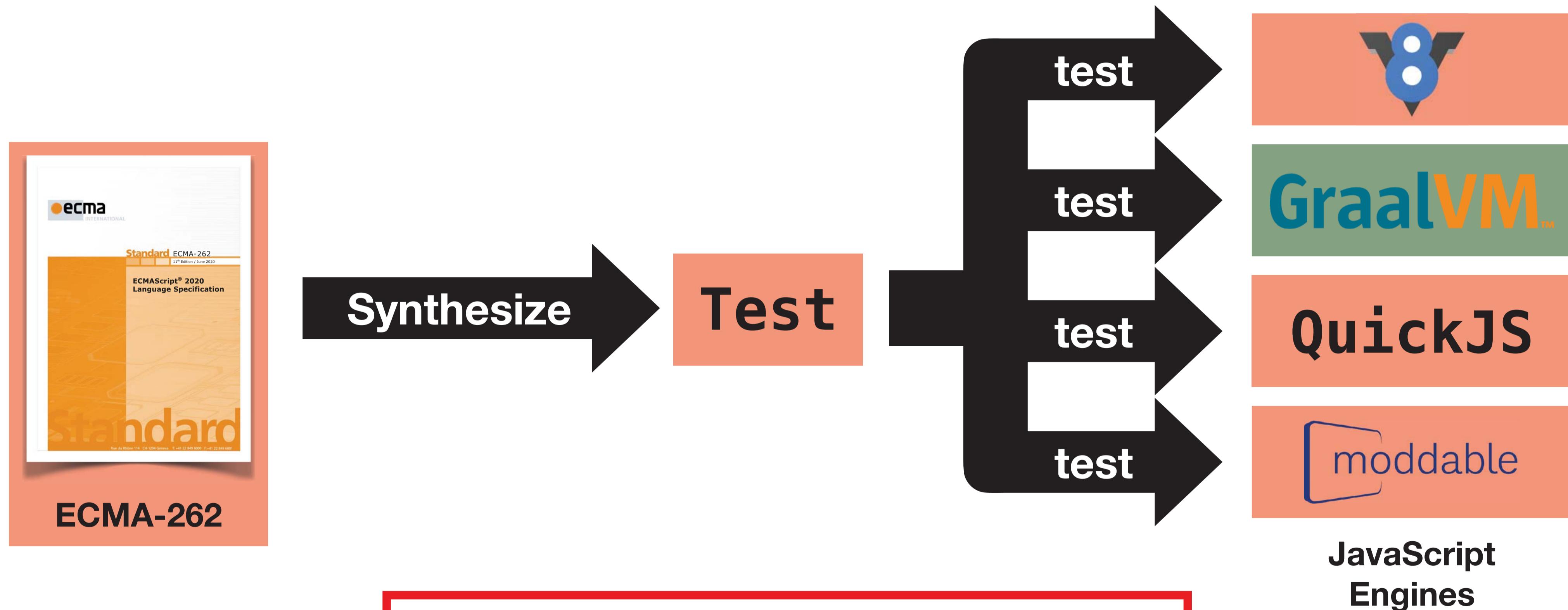
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# JEST - N+1-version Differential Testing



# JEST - N+1-version Differential Testing

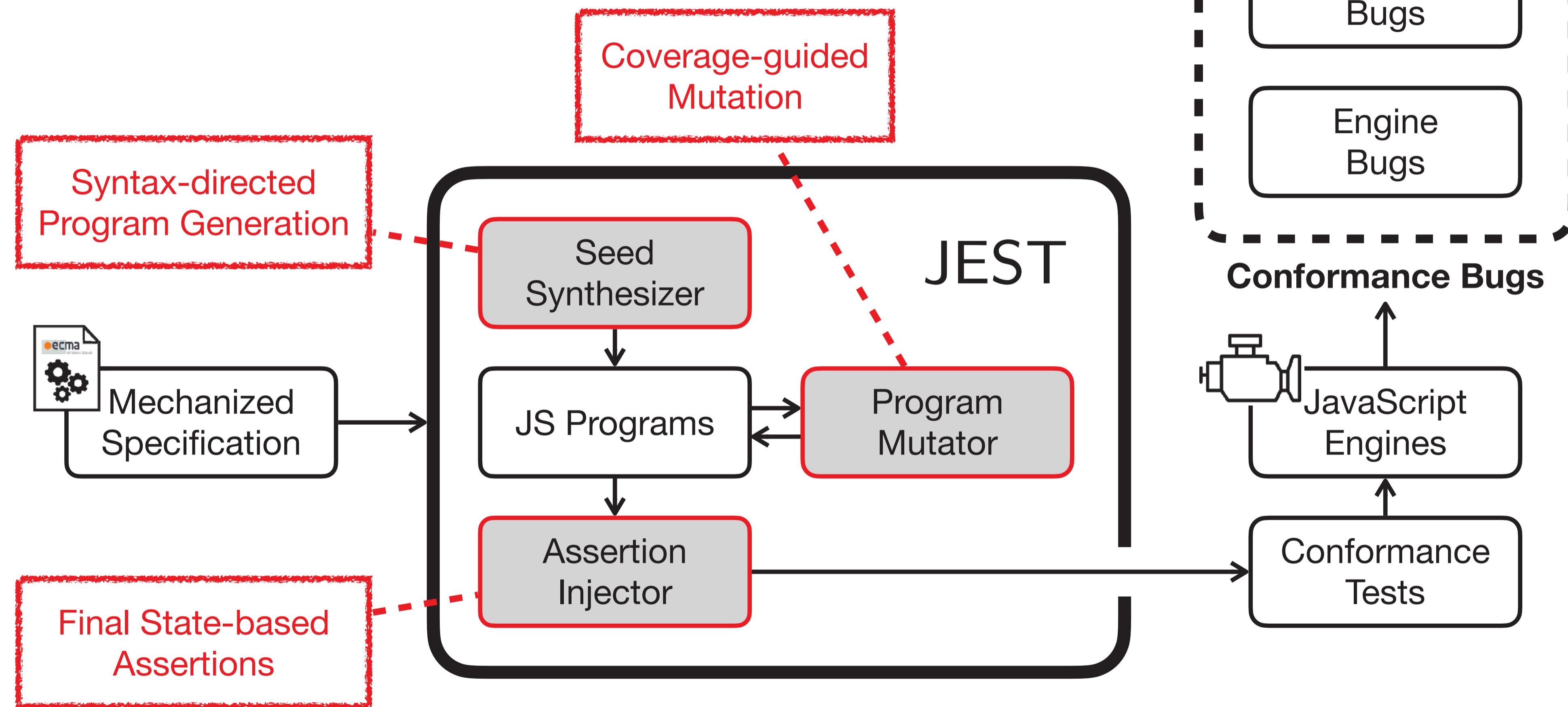


A specification bug in ECMA-262

An engine bug in **GraalVM™**

# JEST [ICSE'21]

## JavaScript Engines and Specification Tester



# JEST - Assertion Injector (7 Kinds)

```
var x = 1 + 2;
```

# JEST - Assertion Injector (7 Kinds)

```
var x = 1 + 2;
```

```
+ $assert.sameValue(x, 3);
```

# JEST - Assertion Injector (7 Kinds)

## 1. Exceptions (Exc)

```
+ // Throw
let x = 42;
function x() {};
```

## 2. Aborts (Abort)

```
+ // Abort
var x = 42; x++;
```

## 3. Variable Values (Var)

```
var x = 1 + 2;
+ $assert.sameValue(x, 3);
```

## 4. Object Values (Obj)

```
var x = {}, y = {}, z = { p: x, q: y };
+ $assert.sameValue(z.p, x);
+ $assert.sameValue(z.q, y);
```

# JEST - Assertion Injector (7 Kinds)

## 5. Object Properties (Desc)

```
var x = { p: 42 };
+ $verifyProperty(x, "p", {
+   value: 42.0, writable: true,
+   enumerable: true, configurable: true
+ });
```

## 6. Property Keys (Key)

```
var x = {[Symbol.match]: 0, p: 0, 3: 0, q: 0, 1: 0}
+ $assert.compareArray(
+   Reflect.ownKeys(x),
+   ["1", "3", "p", "q", Symbol.match]
+ );
```

## 7. Internal Methods and Slots (In)

```
function f() {}
+ $assert.sameValue(Object.getPrototypeOf(f),
+                   Function.prototype);
+ $assert.sameValue(Object.isExtensible(x), true);
+ $assert.callable(f);
+ $assert.constructable(f);
```

# JEST - Evaluation

44 Bugs  
in Engines

TABLE II: The number of engine bugs detected by JEST

Engines	Exc	Abort	Var	Obj	Desc	Key	In	Total
V8	0	0	0	0	0	2	0	2
GraalJS	6	0	0	0	2	8	0	16
QuickJS	3	0	1	0	0	2	0	6
Moddable XS	12	0	0	0	3	5	0	20
<b>Total</b>	21	0	1	0	5	17	0	44

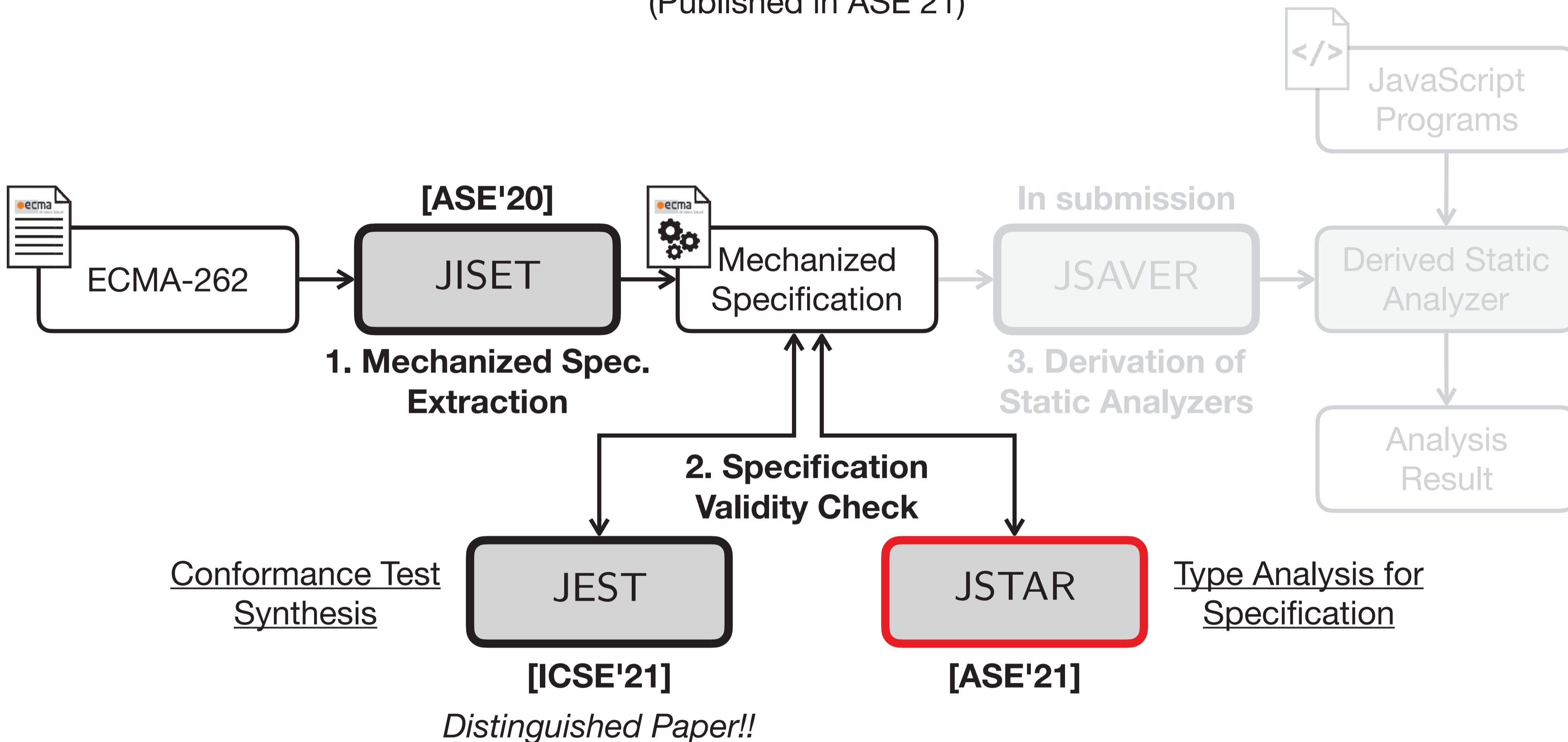
27 Bugs  
in Spec.

TABLE III: Specification bugs in ECMAScript 2020 (ES11) detected by JEST

Name	Feature	#	Assertion	Known	Created	Resolved	Existed
ES11-1	Function	12	Key	O	2019-02-07	2020-04-11	429 days
ES11-2	Function	8	Key	O	2015-06-01	2020-04-11	1,776 days
ES11-3	Loop	1	Exc	O	2017-10-17	2020-04-30	926 days
ES11-4	Expression	4	Abort	O	2019-09-27	2020-04-23	209 days
ES11-5	Expression	1	Exc	O	2015-06-01	2020-04-28	1,793 days
ES11-6	Object	1	Exc	X	2019-02-07	2020-11-05	637 days

# JSTAR: JavaScript Specification Type Analyzer using Refinement

Jihyeok Park, Seungmin An, Wonho Shin, Yusung Sim, and Sukyoung Ryu  
(Published in ASE'21)



# JSTAR - Types in Specification

## 20.3.2.28 Math.round ( $x$ )

1. Let  $n$  be ? ToNumber( $x$ ).
2. If  $n$  is an integral Number, return  $n$ .
3. If  $x < 0.5$  and  $x > 0$ , return +0.
4. If  $x < 0$  and  $x \geq -0.5$ , return -0.
- ...

<https://github.com/tc39/ecma262/tree/575149cf77aebcf3a129e165bd89e14caafc31c>

# JSTAR - Types in Specification

**20.3.2.28 Math.round ( $x$ )**  $x: (\text{String} \vee \text{Boolean} \vee \text{Number} \vee \text{Object} \vee \dots)$

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- ...
- Type Mismatch for numeric operator `>`

<https://github.com/tc39/ecma262/tree/575149cf77aebcf3a129e165bd89e14caafc31c>

# JSTAR - Types in Specification

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  4. If  $x < 0$  and  $x \geq -0.5$ , return  $-0$ .
- ...

Type Mismatch for  
numeric operator `>`

Math.round(true) = ???  
Math.round(false) = ???

<https://github.com/tc39/ecma262/tree/575149cf77aebcf3a129e165bd89e14caafc31c>

# JSTAR - Types in Specification

**20.3.2.28 Math.round ( $x$ )**  $x: (\text{String} \vee \text{Boolean} \vee \text{Number} \vee \text{Object} \vee \dots)$

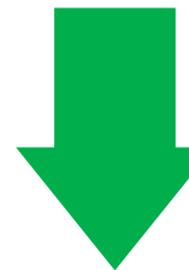
1. Let  $n$  be  $\text{?ToNumber}(x)$ .  $\text{ToNumber}(x): (\text{Number} \vee \text{Exception}) \wedge n: (\text{Number})$

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4. If  $x < 0$  and  $x \geq -0.5$ , return  $-0$ .

...

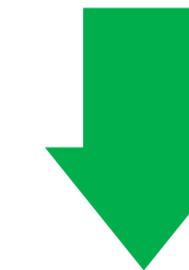


3. If  $n < 0.5$  and  $n > 0$ , return  $+0$ .

4. If  $n < 0$  and  $n \geq -0.5$ , return  $-0$ .

Type Mismatch for  
numeric operator `>`

Math.round(true) = ???  
Math.round(false) = ???

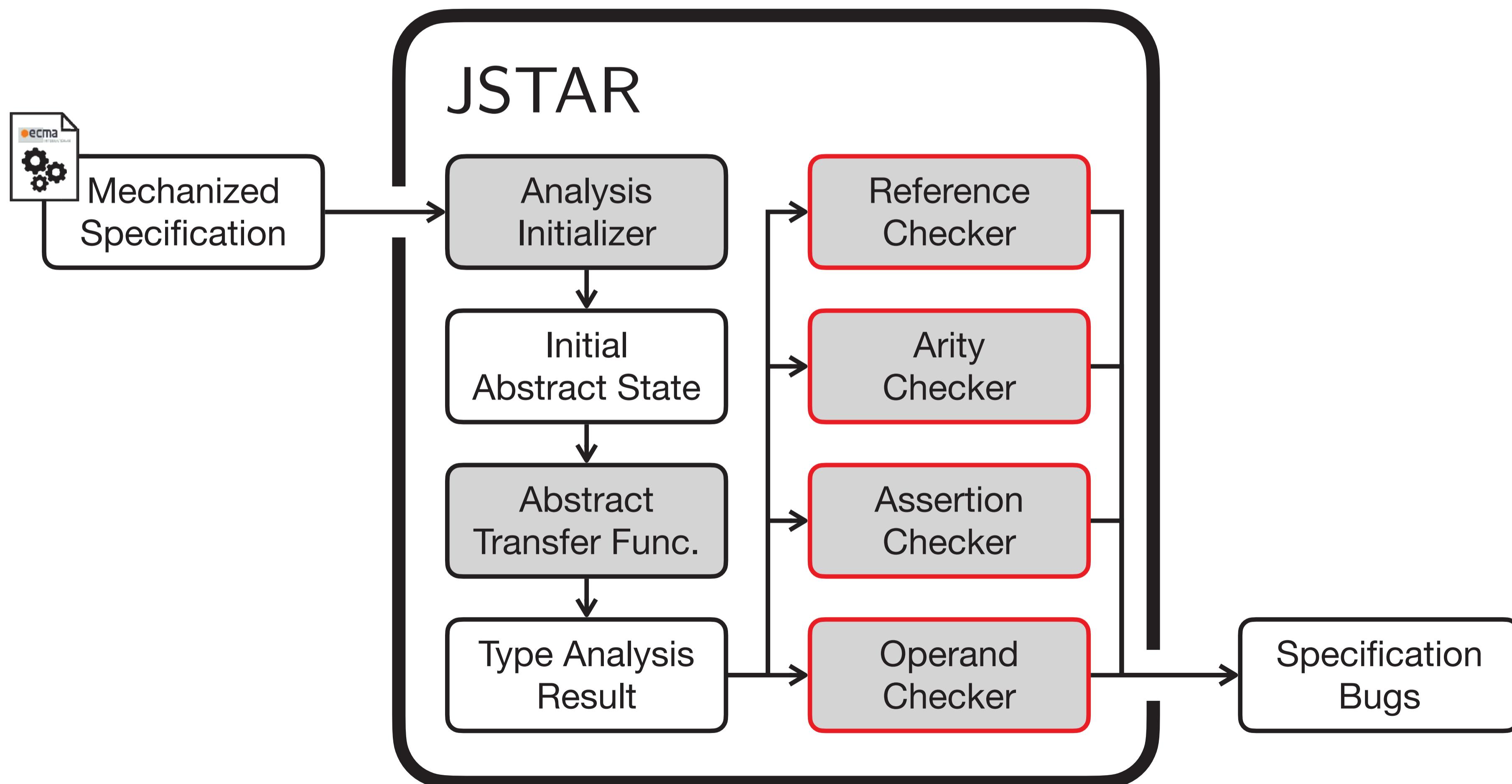


Math.round(true) = 1  
Math.round(false) = 0

<https://github.com/tc39/ecma262/tree/575149cf77aebcf3a129e165bd89e14caafc31c>

# JSTAR [ASE'21]

## JavaScript Specification Type Analyzer using Refinement



# JSTAR - Evaluation

- Type Analysis for 864 versions of ECMA-262

59.2%  
Precision

93 Bugs  
Detected

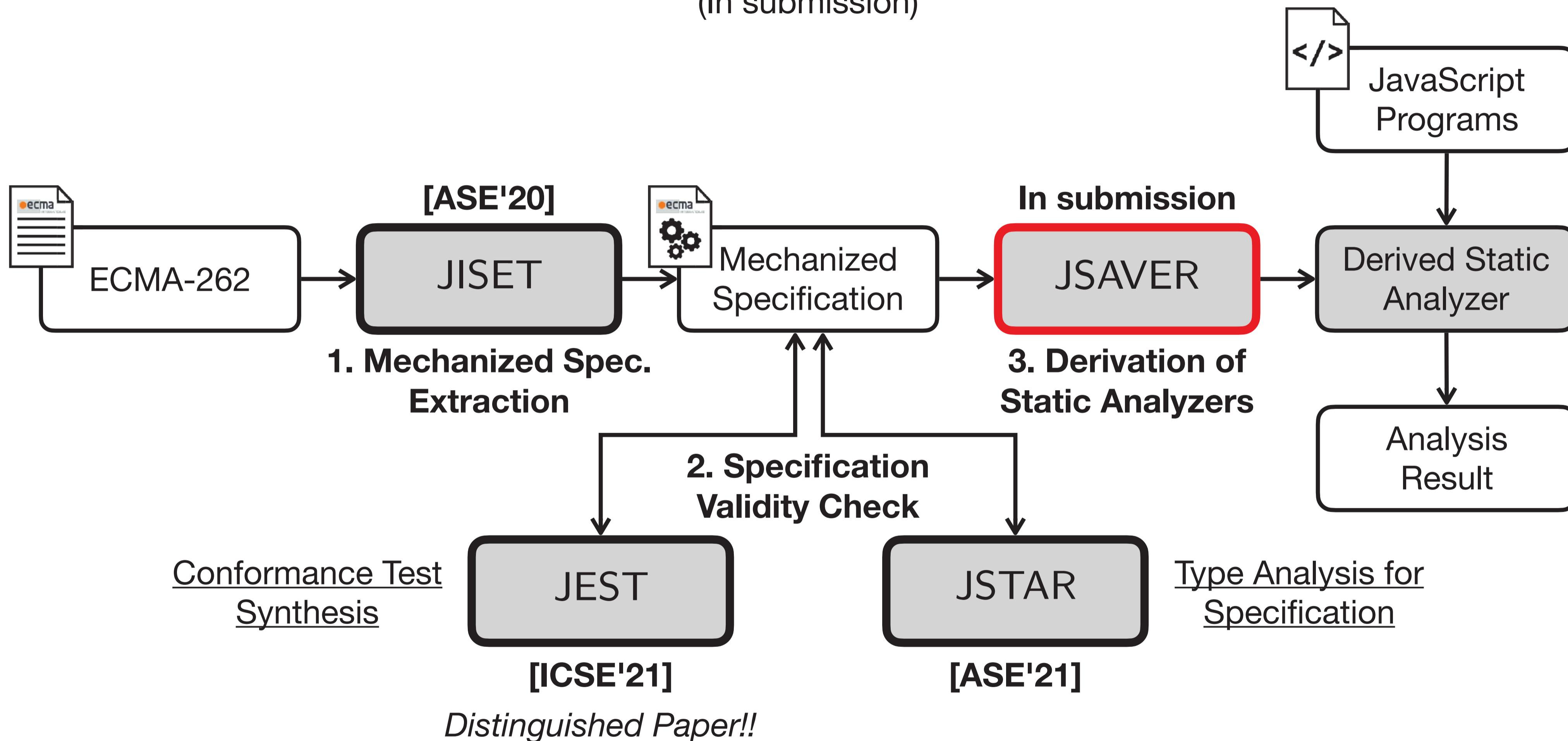
Checker	Bug Kind	Precision = (# True Bugs) / (# Detected Bugs)				
		no-refine		refine		Δ
Reference	UnknownVar	62 / 106	17 / 60	63 / 78	17 / 31	+1 / -28
	DuplicatedVar		45 / 46		46 / 47	
Arity	MissingParam	4 / 4	4 / 4	4 / 4	4 / 4	/ /
Assertion	Assertion	4 / 56	4 / 56	4 / 31	4 / 31	/ -25 / -25
Operand	NoNumber	22 / 113	2 / 65	22 / 44	2 / 6	/ -69 / -59
	Abrupt		20 / 48		20 / 38	
Total		92 / 279 (33.0%)		93 / 157 (59.2%)		+1 / -122 (+26.3%)

Name	Feature	#	Checker	Created	Life Span
ES12-1	Switch	3	Reference	2015-09-22	1,996 days
ES12-2	Try	3	Reference	2015-09-22	1,996 days
ES12-3	Arguments	1	Reference	2015-09-22	1,996 days
ES12-4	Array	2	Reference	2015-09-22	1,996 days
ES12-5	Async	1	Reference	2015-09-22	1,996 days
ES12-6	Class	1	Reference	2015-09-22	1,996 days
ES12-7	Branch	1	Reference	2015-09-22	1,996 days
ES12-8	Arguments	2	Operand	2015-12-16	1,910 days

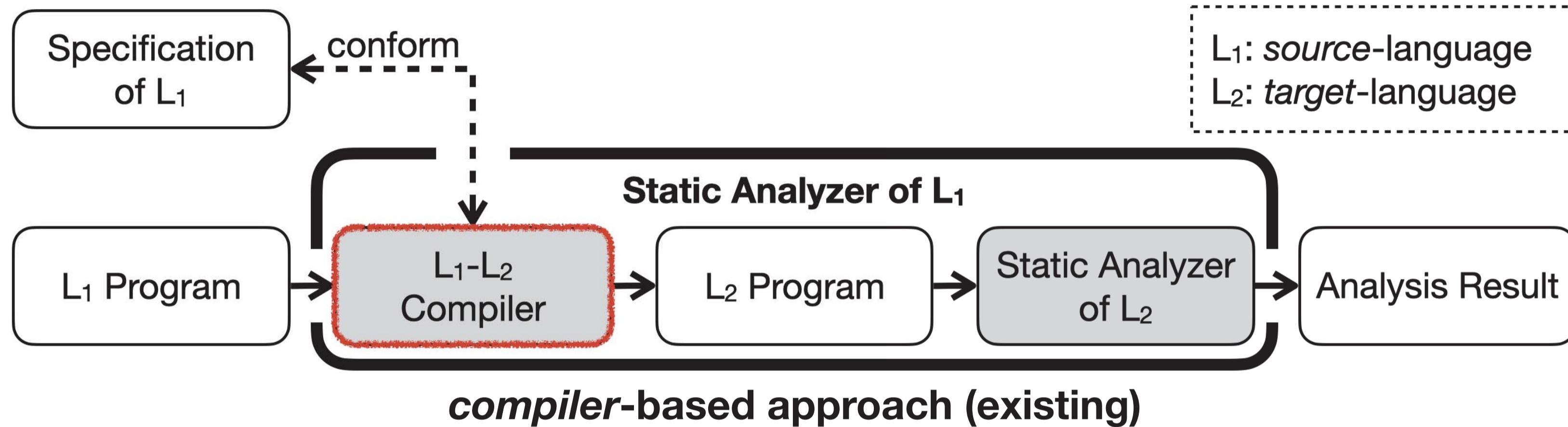
14 Bugs  
in ES12

# Automatically Deriving JavaScript Static Analyzers from Language Specifications

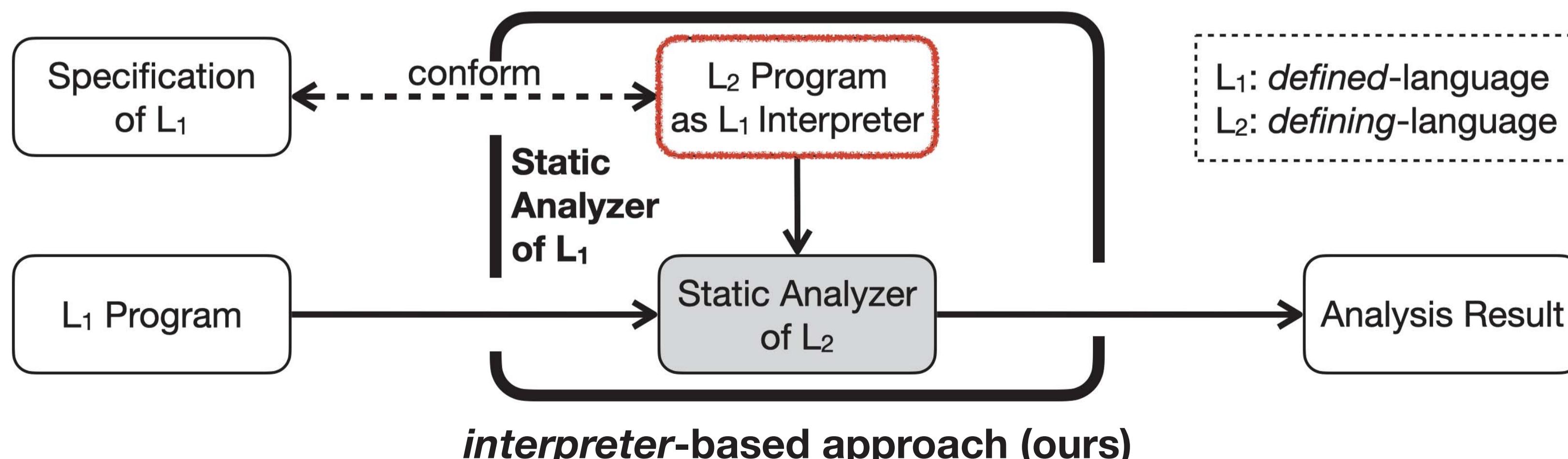
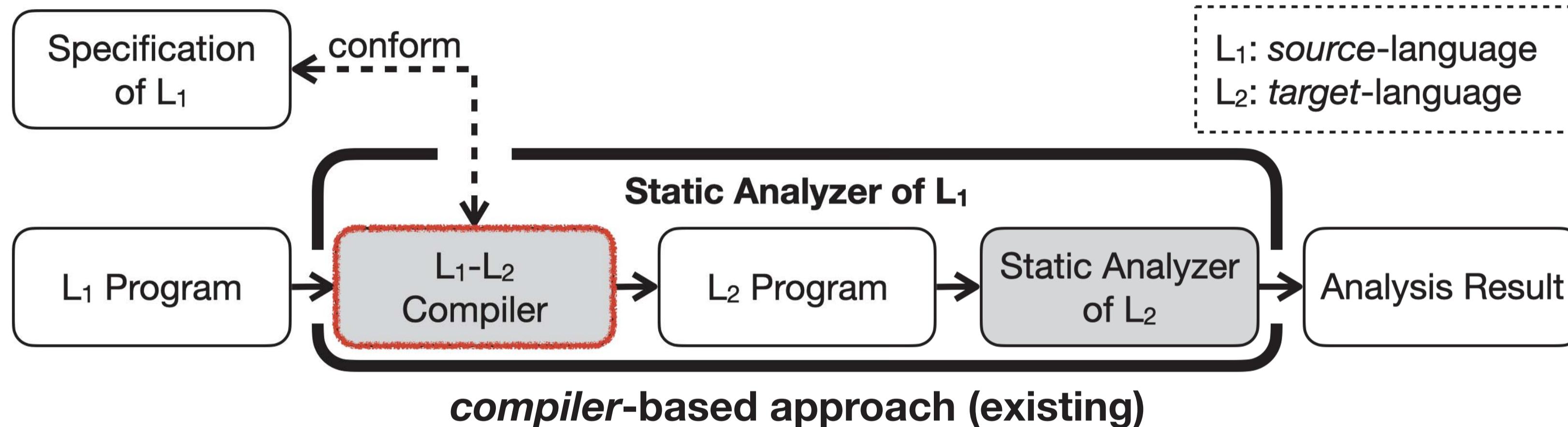
Jihyeok Park, Seungmin An, and Sukyoung Ryu  
(In submission)



# JSAVER - Meta-Level Static Analysis



# JSAVER - Meta-Level Static Analysis



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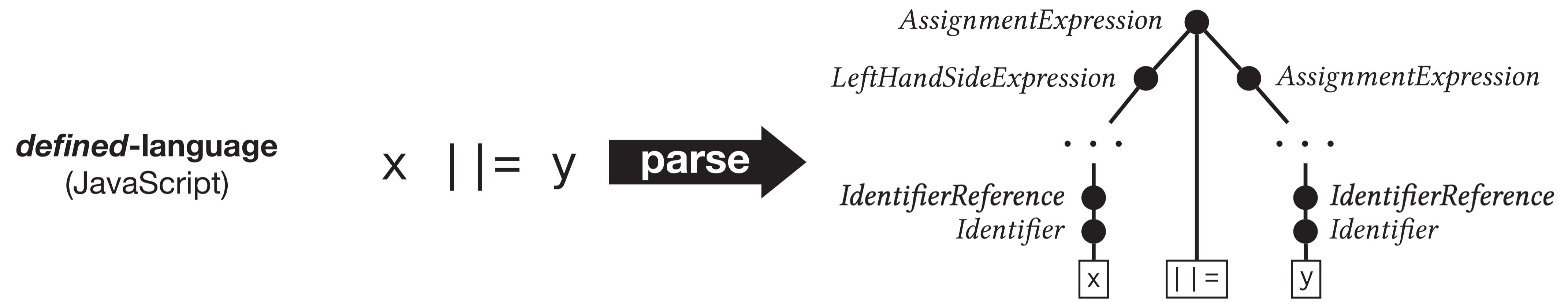
***defined-language***  
(JavaScript)

$x \mid |= y$



***defining-language***  
(IR<sub>ES</sub>)

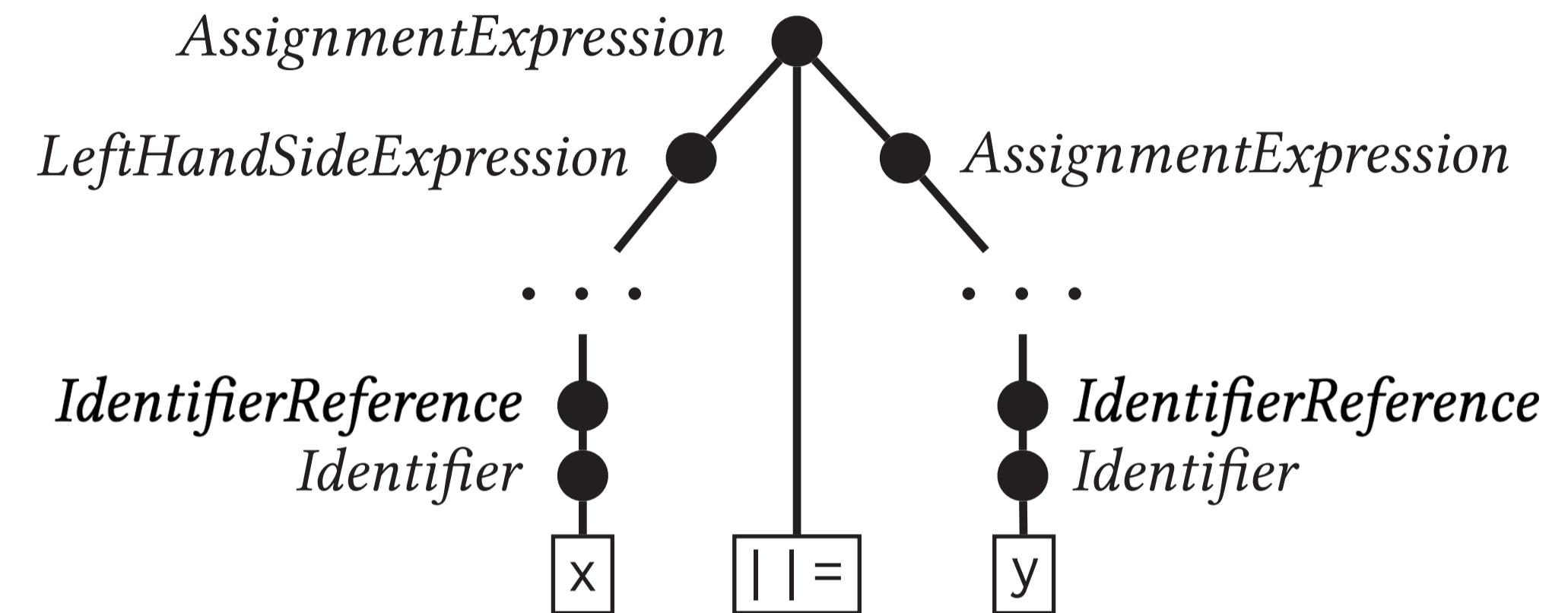
# JSAVER - Meta-Level Static Analysis



# JSAVER - Meta-Level Static Analysis

**defined-language**  
(JavaScript)

x | |= y **parse** →

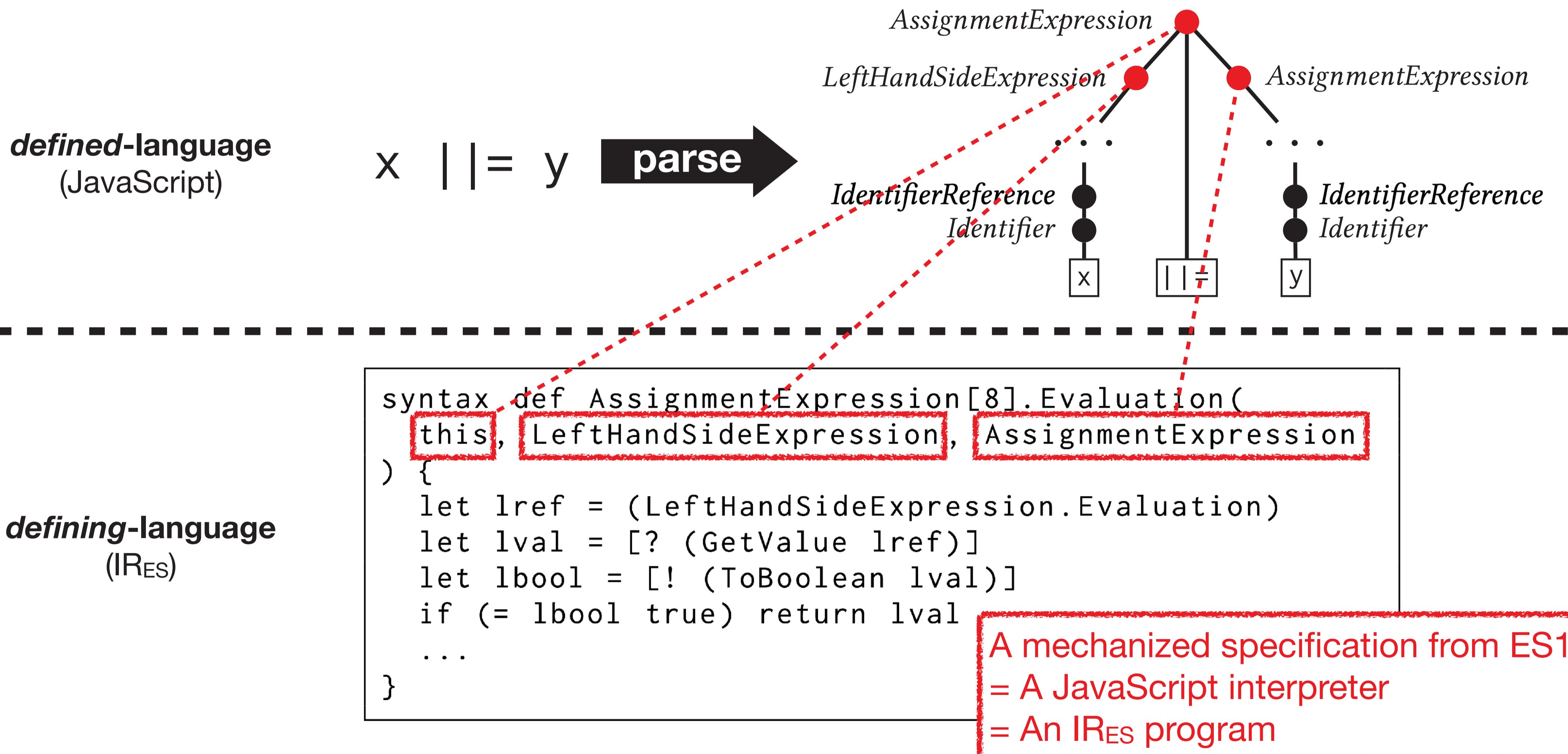


**defining-language**  
(IR<sub>ES</sub>)

```
syntax def AssignmentExpression[8].Evaluation(
  this, LeftHandSideExpression, AssignmentExpression
) {
  let lref = (LeftHandSideExpression.Evaluation)
  let lval = [? (GetValue lref)]
  let lbool = [! (ToBoolean lval)]
  if (= lbool true) return lval
  ...
}
```

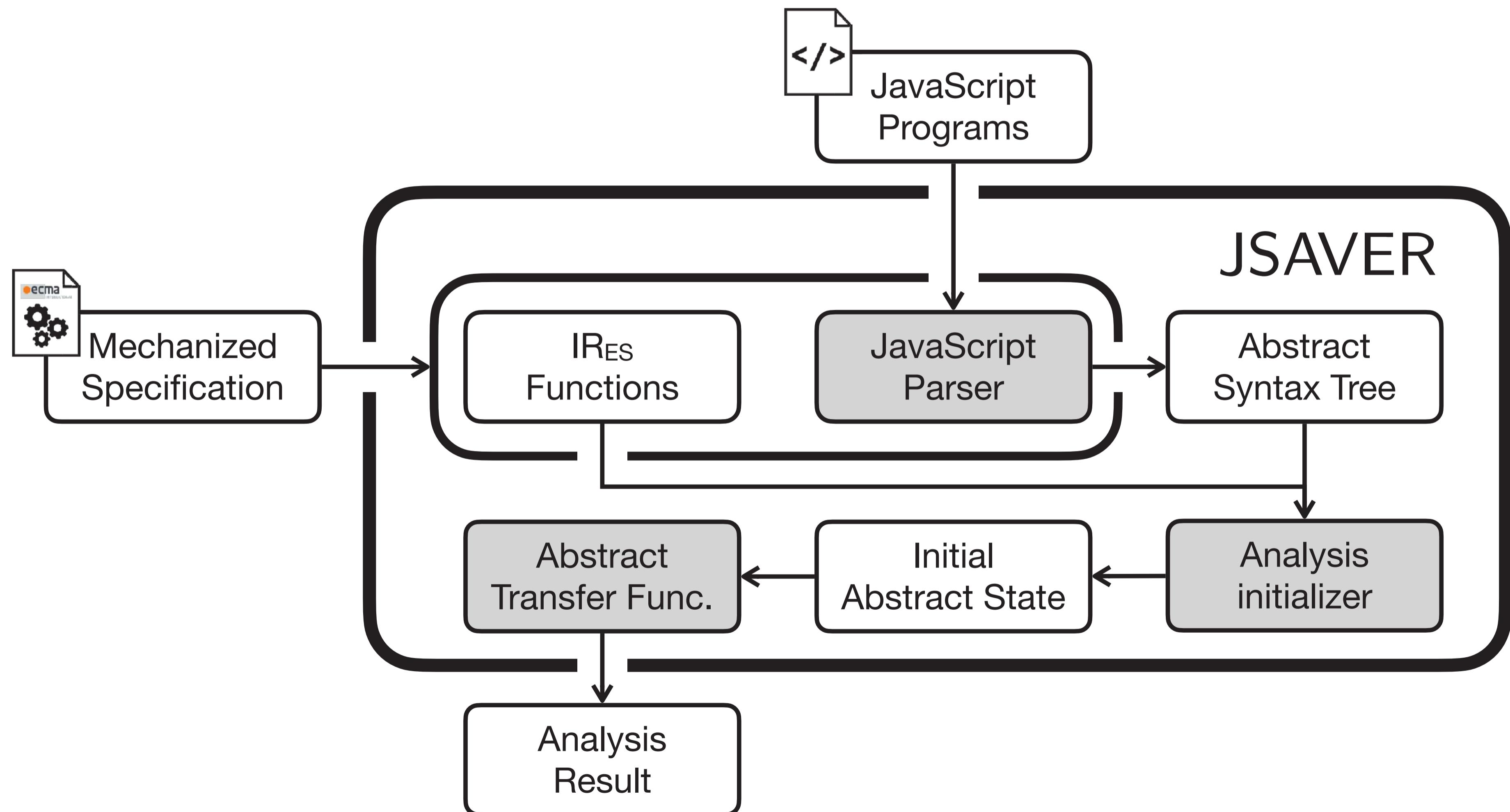
A mechanized specification from ES12  
= A JavaScript interpreter  
= An IR<sub>ES</sub> program

# JSAVER - Meta-Level Static Analysis

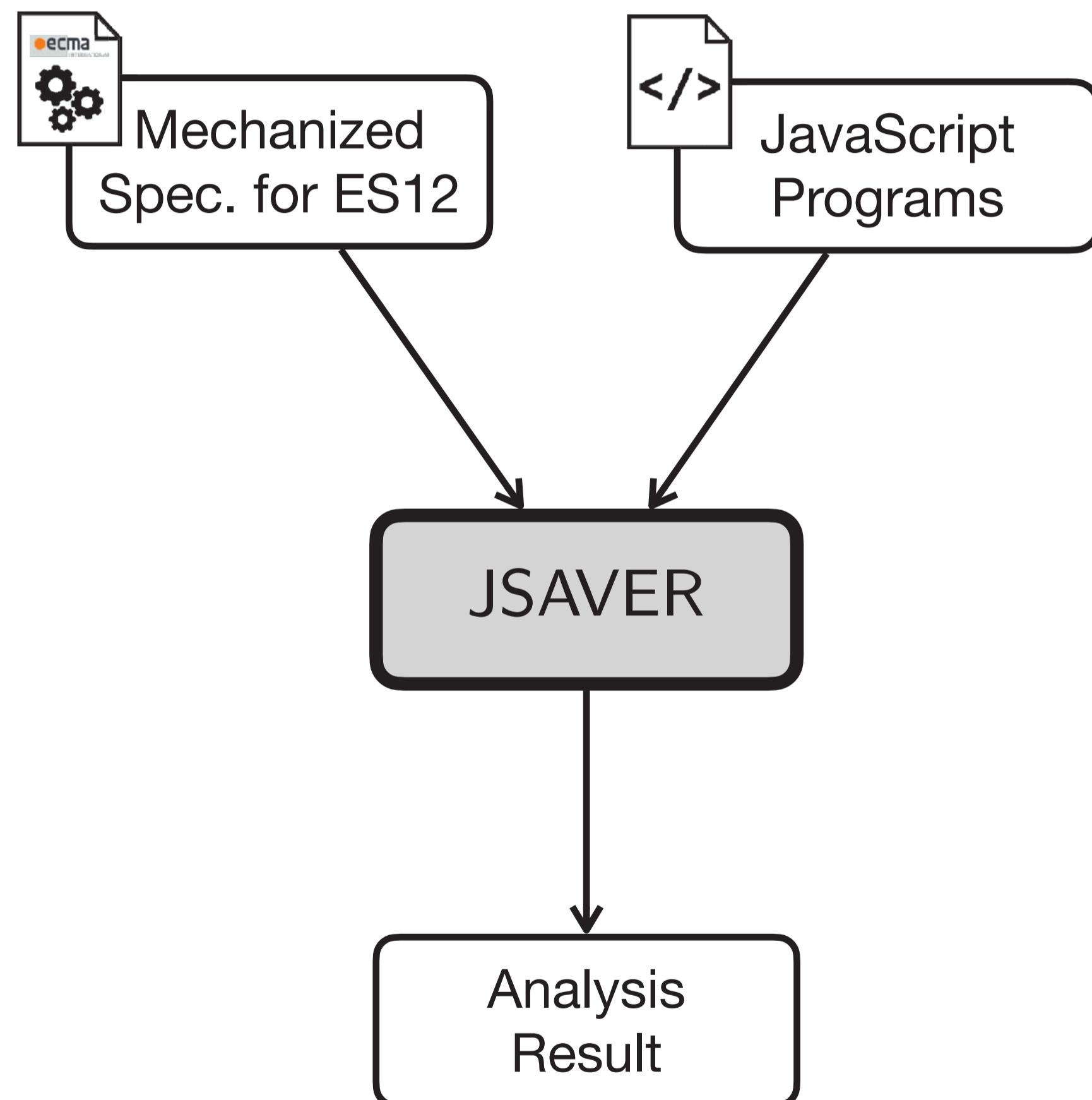


# JSAVER In submission

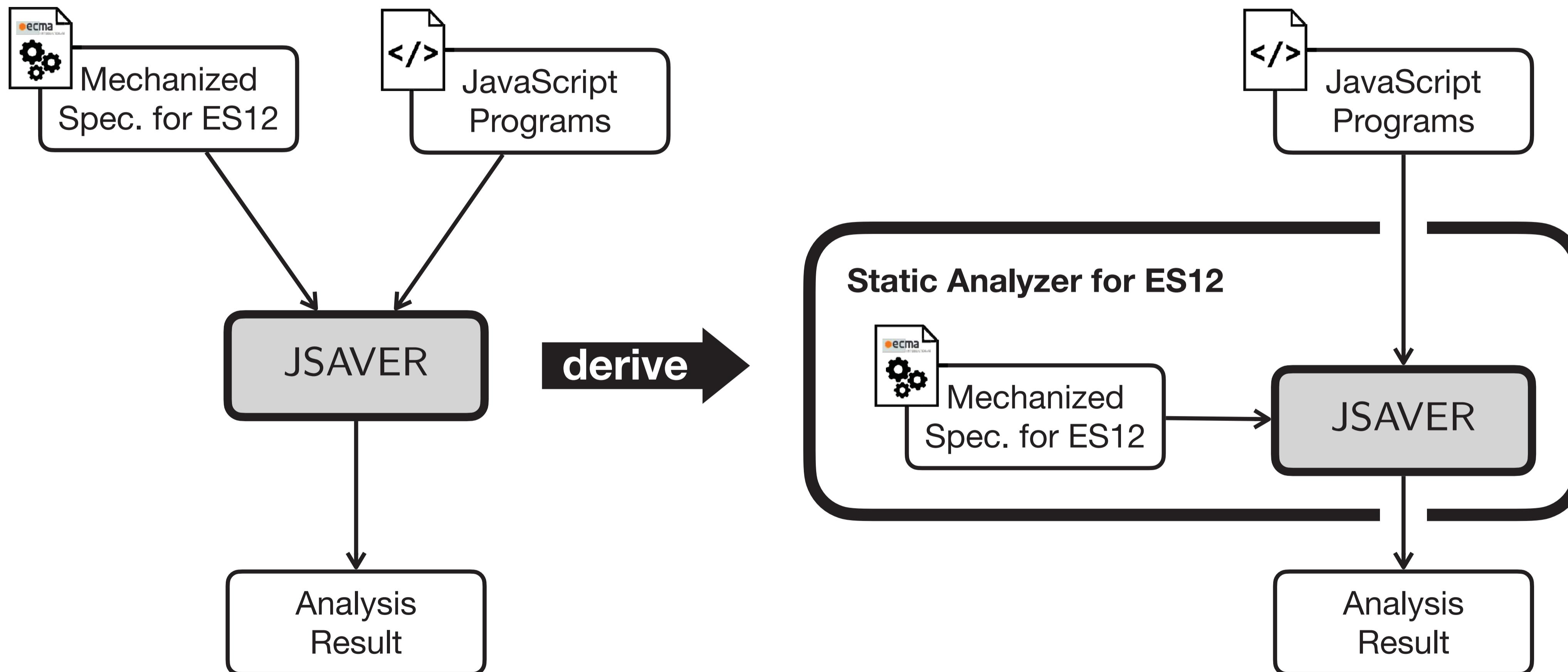
## JavaScript Static Analyzer via ECMAScript Representation



# JSAVER - Static Analyzer Derivation

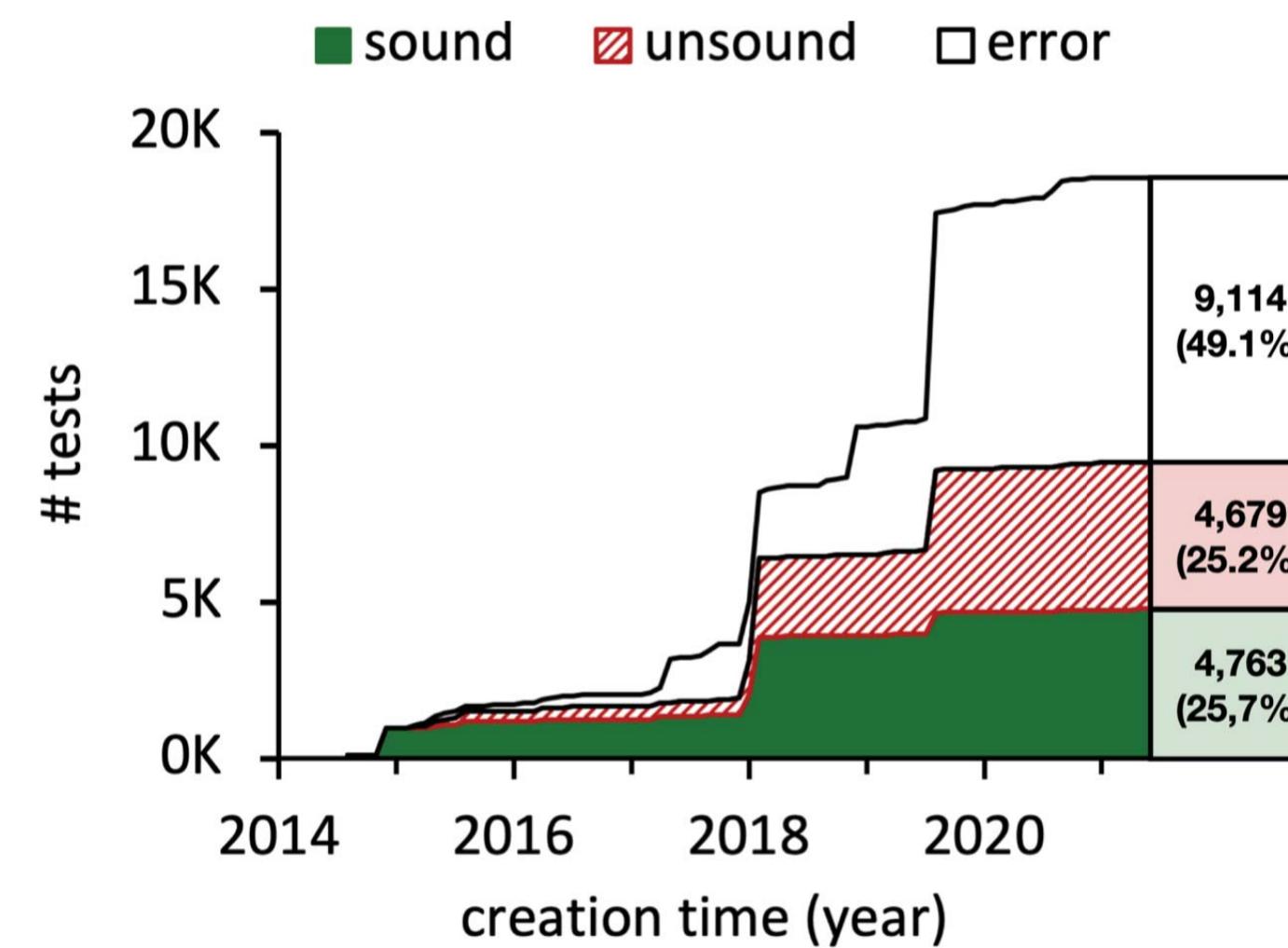


# JSAVER - Static Analyzer Derivation

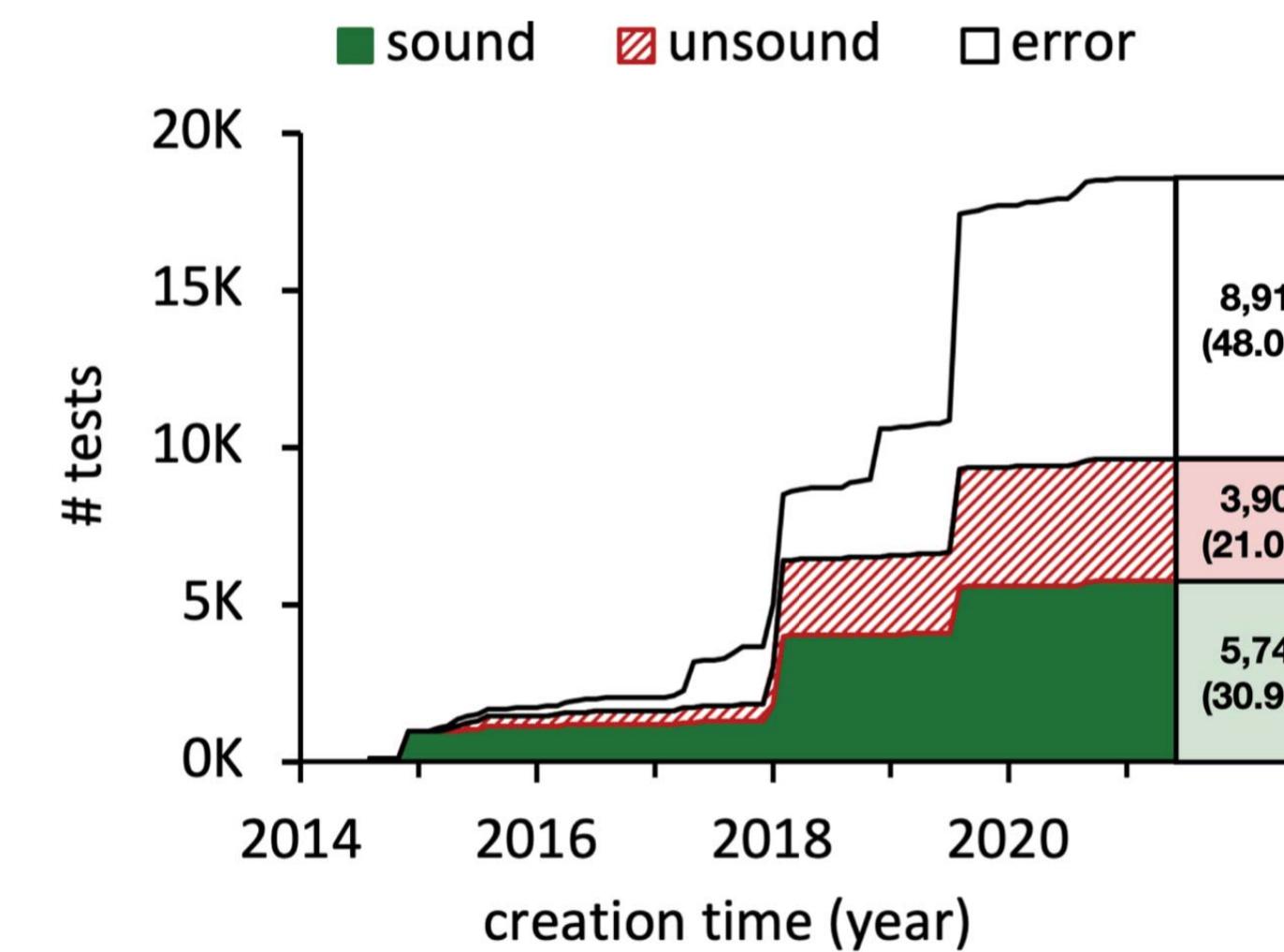


# JSAVER - Evaluation

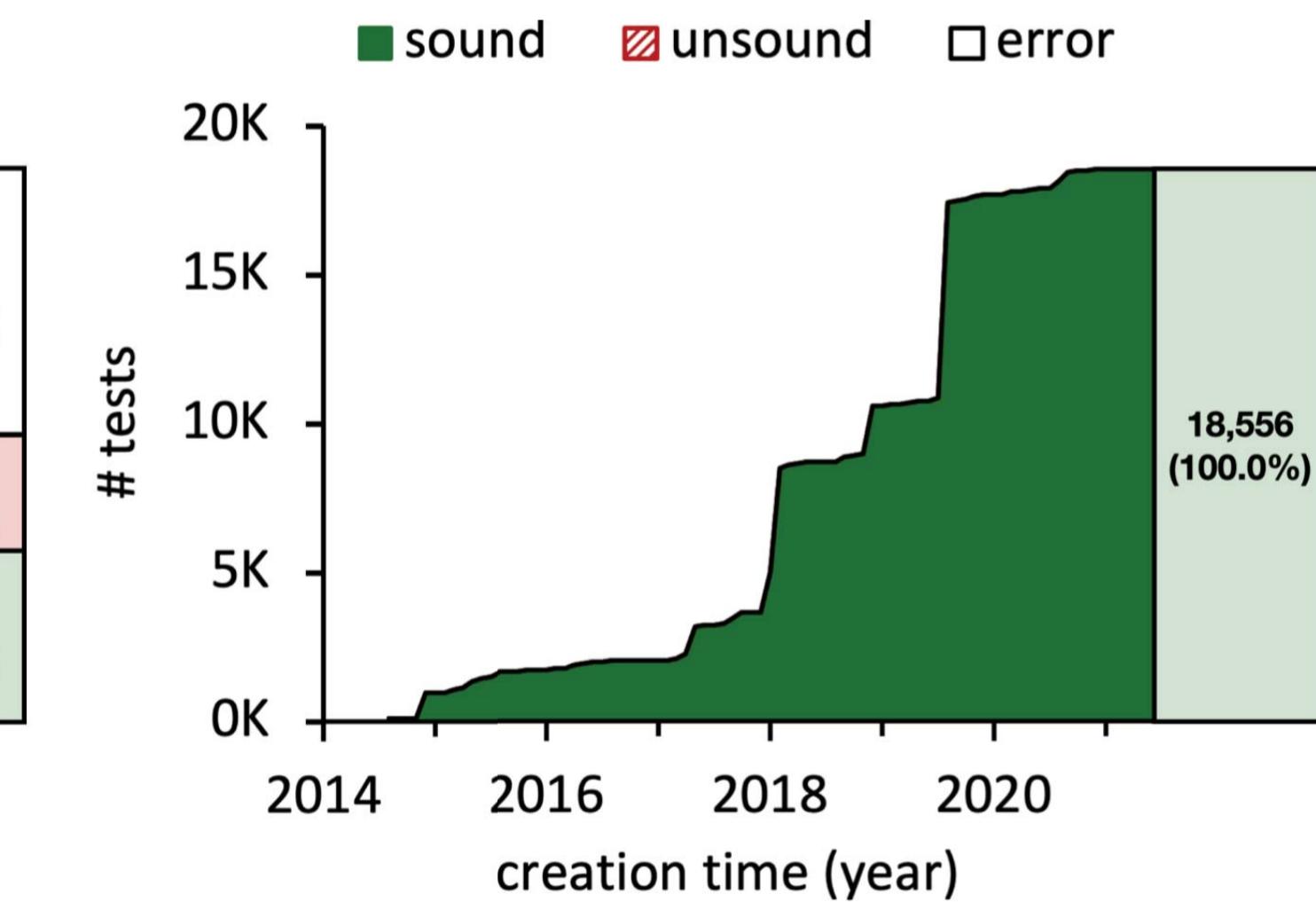
- **Soundness / Precision / Performance**
  - 18,556 applicable tests in Test262
  - 3,903 tests analyzable by all the three analyzers



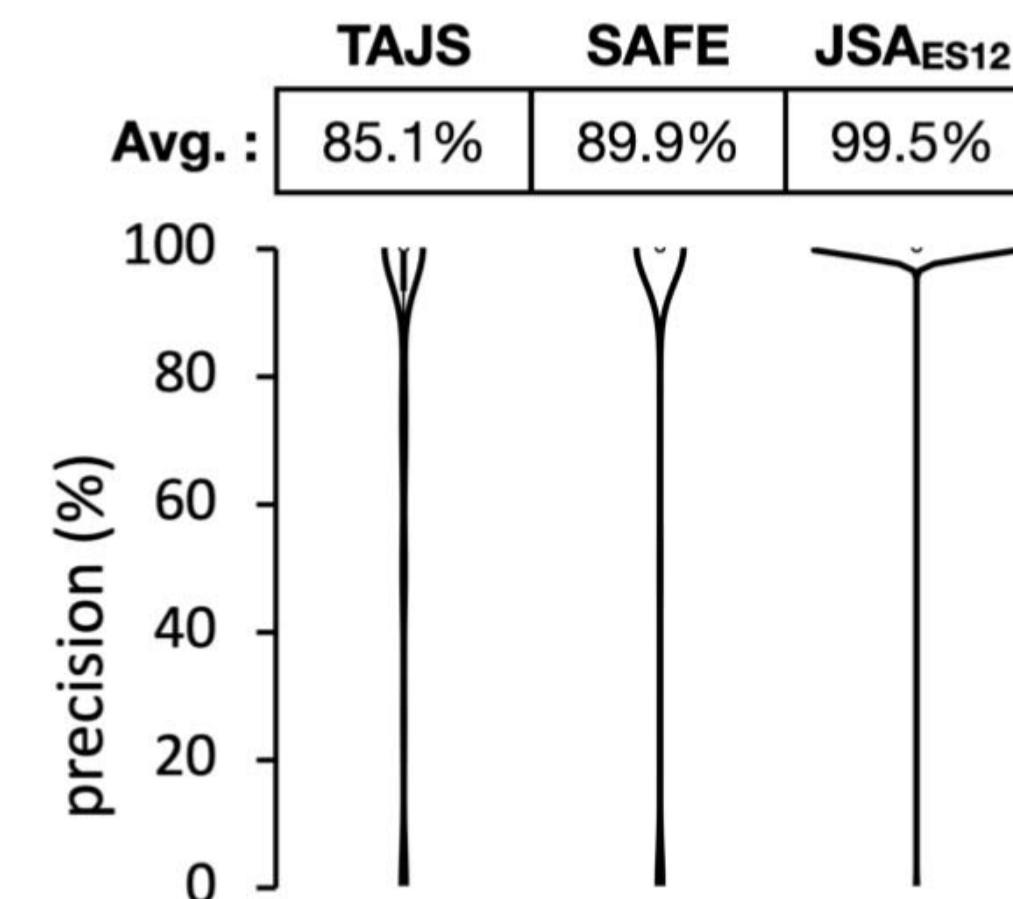
(a) Analysis results of TAJS



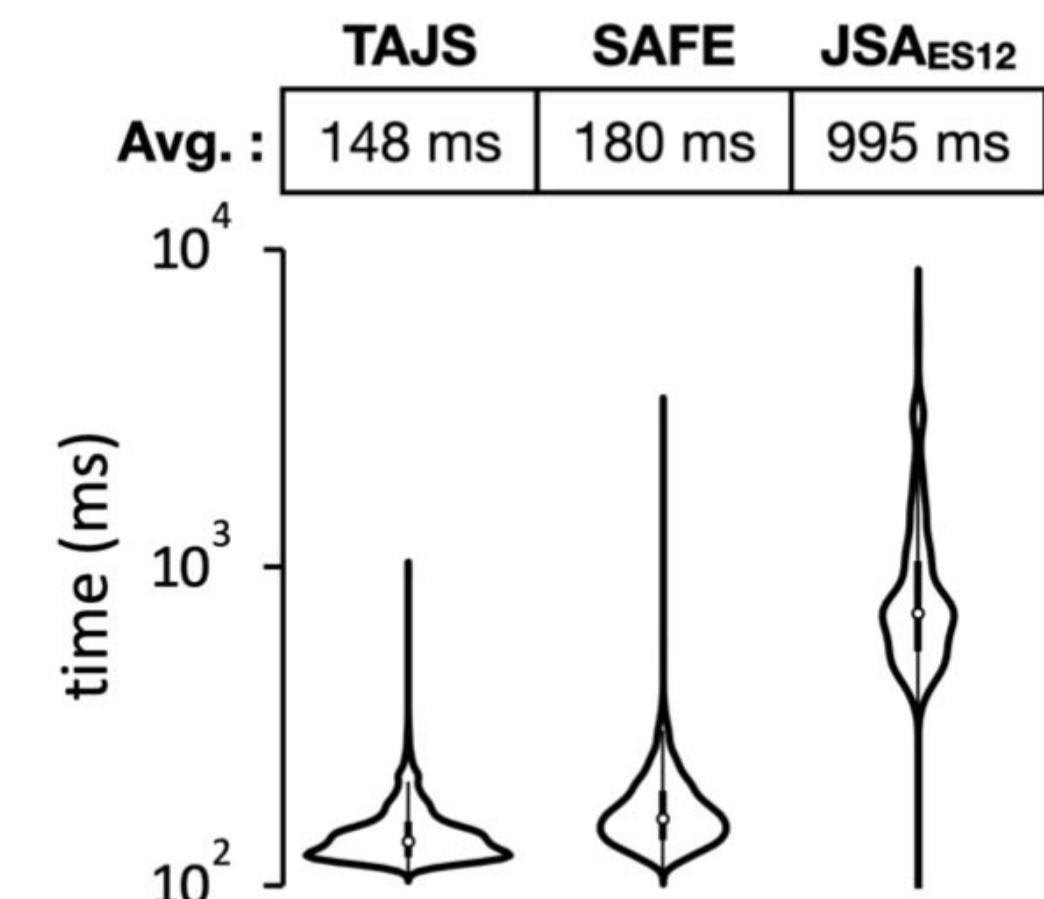
(b) Analysis results of SAFE



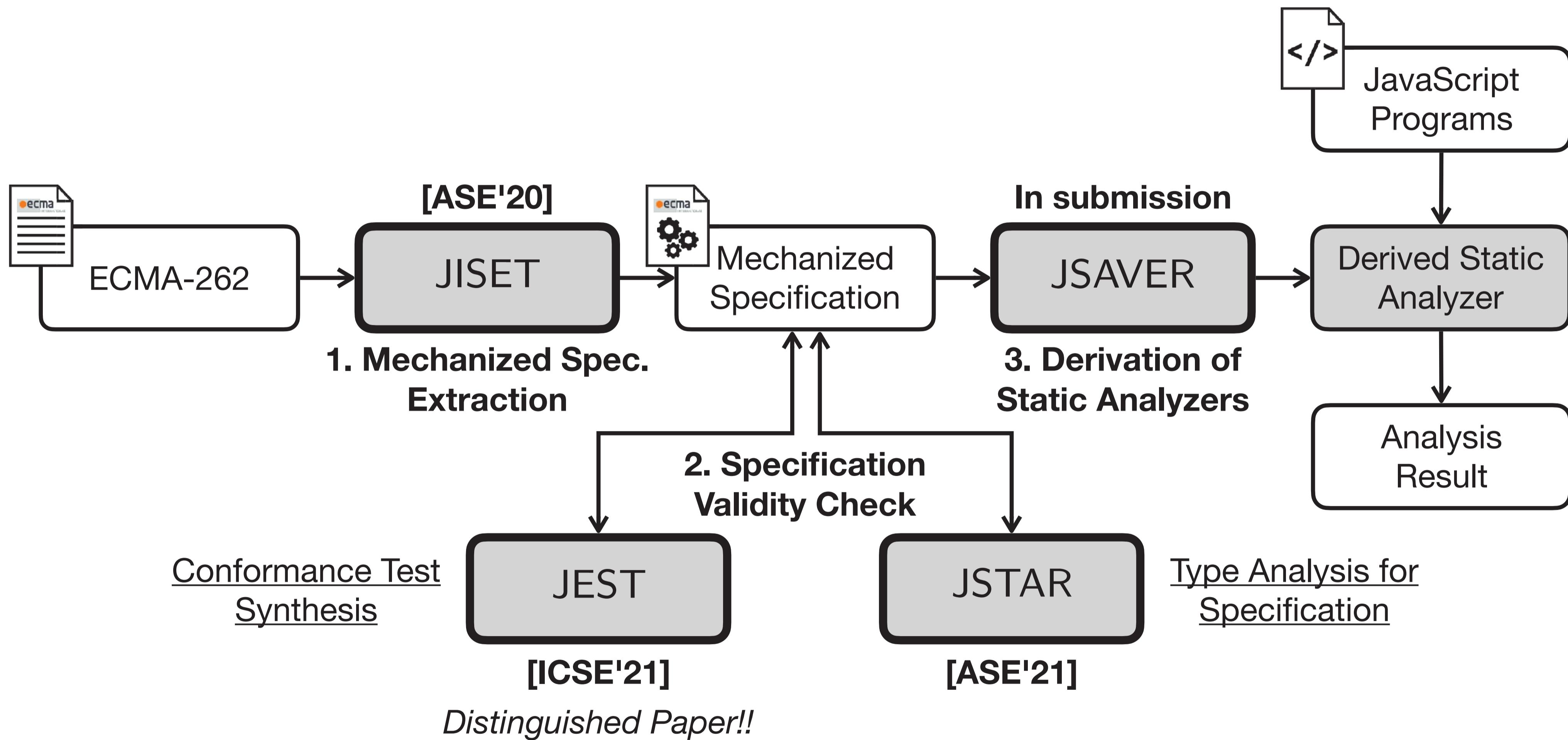
(c) Analysis results of JSA<sub>ES12</sub>



(a) The analysis precision



(b) The analysis performance



ESMeta - es-meta/esmeta: ECMAScript Metalanguage for Generation of Language-based Tools

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jhnaldo Removed unnecessary leg... 38 minutes ago 233

.github/workfl... Update ci.yml 15 days ago

project More project setting 20 days ago

src Removed unnecessary legacy files 38 minutes ago

tests Removed unnecessary legacy files 38 minutes ago

.gitignore Added tests for grammar Stringifi... 15 days ago

.jvmopts Added .jvmopts 14 days ago

.scalafmt.conf Reformatted code with modified s... 2 days ago

LICENSE sbt project setting 21 days ago

ECMAScript Metalanguage for Generation of Language-based Tools

Readme

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3 watching

0 forks

Releases

No releases published

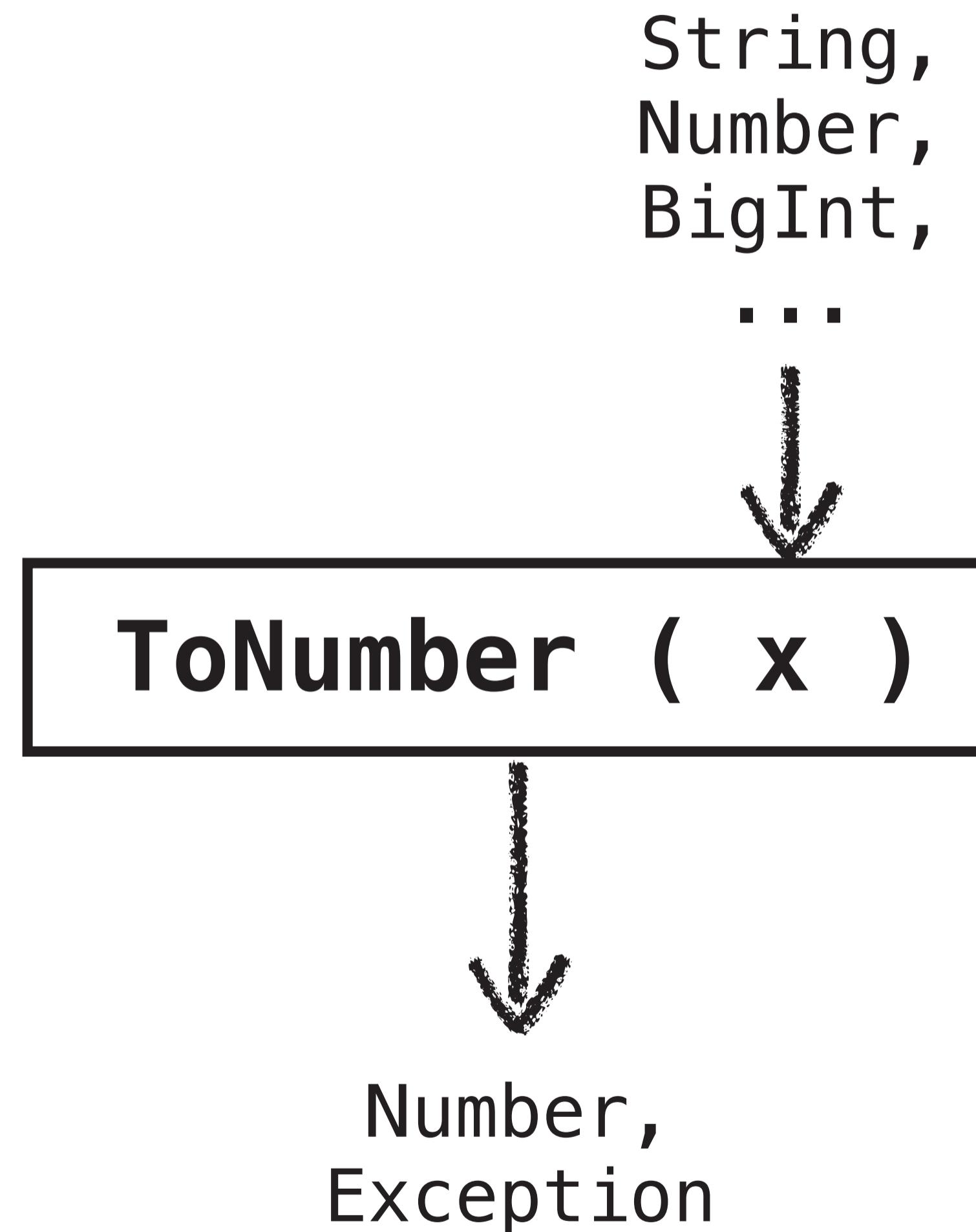
Create a new release

File	Action	Time Ago
.github/workfl...	Update ci.yml	15 days ago
project	More project setting	20 days ago
src	Removed unnecessary legacy files	38 minutes ago
tests	Removed unnecessary legacy files	38 minutes ago
.gitignore	Added tests for grammar Stringifi...	15 days ago
.jvmopts	Added .jvmopts	14 days ago
.scalafmt.conf	Reformatted code with modified s...	2 days ago
LICENSE	sbt project setting	21 days ago

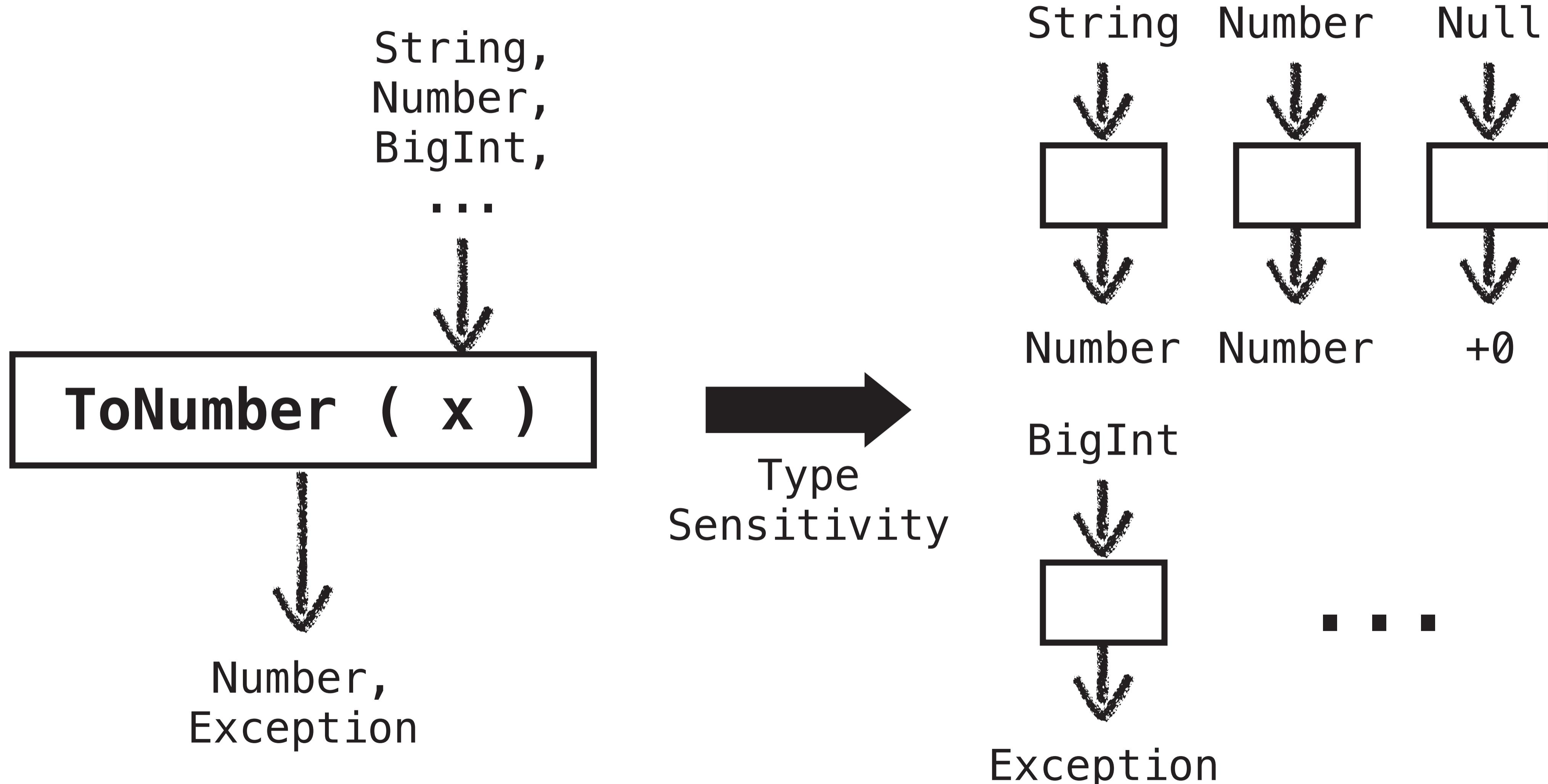
<https://github.com/es-meta/esmeta>

# Backup Slides

# JSTAR - Precision $\uparrow$ - 1) Type Sensitivity



# JSTAR - Precision $\uparrow - 1$ ) Type Sensitivity



# JSTAR - Precision $\uparrow$ - 2) Type Refinement

$$\text{refine}(!e, b)(\sigma^\#) = \text{refine}(e, \neg b)(\sigma^\#)$$

$$\text{refine}(e_0 \sqcup e_1, b)(\sigma^\#) = \begin{cases} \sigma_0^\# \sqcup \sigma_1^\# & \text{if } b \\ \sigma_0^\# \sqcap \sigma_1^\# & \text{if } \neg b \end{cases}$$

$$\text{refine}(e_0 \& e_1, b)(\sigma^\#) = \begin{cases} \sigma_0^\# \sqcap \sigma_1^\# & \text{if } b \\ \sigma_0^\# \sqcup \sigma_1^\# & \text{if } \neg b \end{cases}$$

$$\text{refine}(x.\text{Type} == c_{\text{normal}}, \#t)(\sigma^\#) = \sigma^\#[x \mapsto \tau_x^\# \sqcap \text{normal}(\mathbb{T})]$$

$$\text{refine}(x.\text{Type} == c_{\text{normal}}, \#f)(\sigma^\#) = \sigma^\#[x \mapsto \tau_x^\# \sqcap \{\text{abrupt}\}]$$

$$\text{refine}(x == e, \#t)(\sigma^\#) = \sigma^\#[x \mapsto \tau_x^\# \sqcap \tau_e^\#]$$

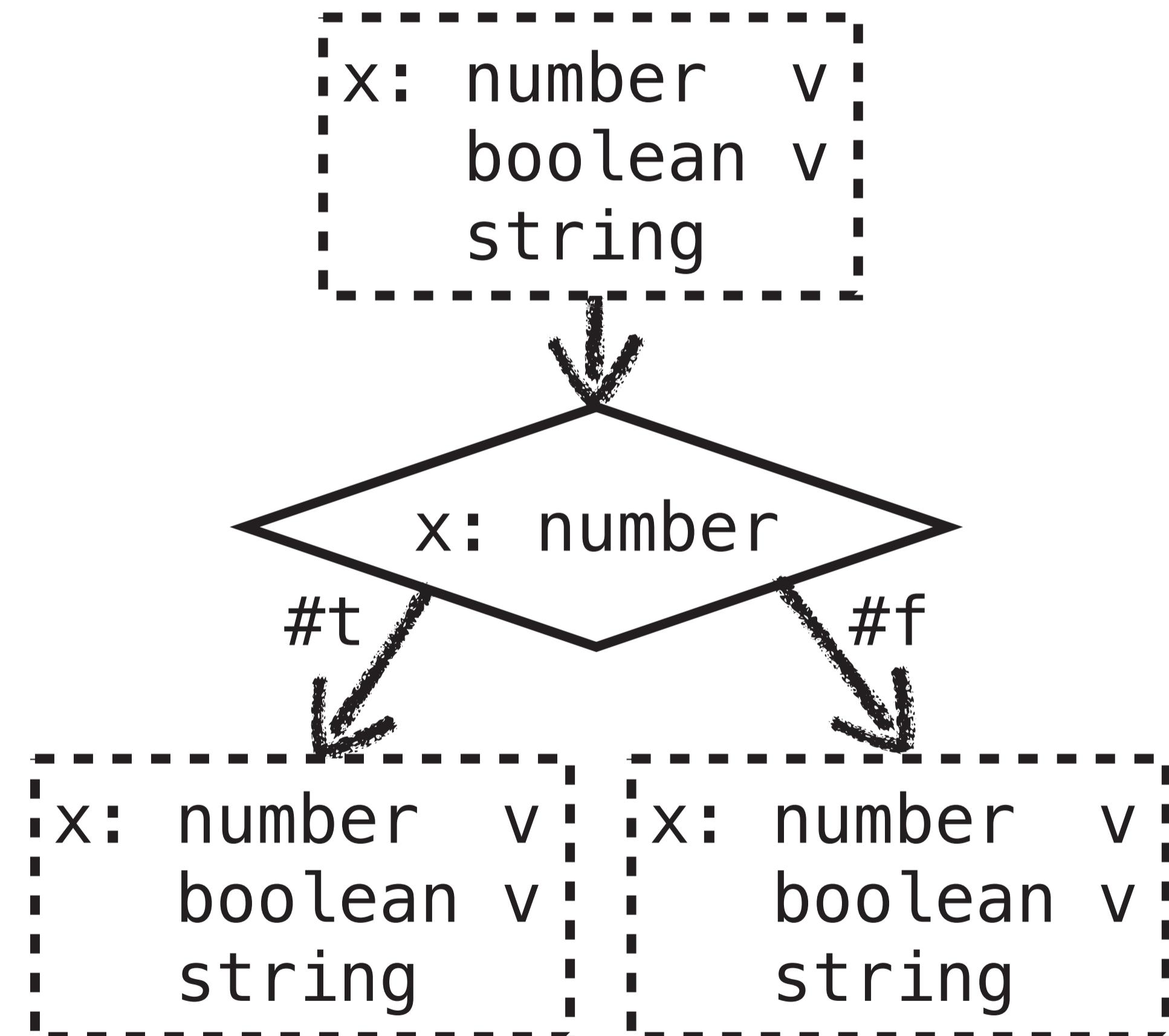
$$\text{refine}(x == e, \#f)(\sigma^\#) = \sigma^\#[x \mapsto \tau_x^\# \setminus [\tau_e^\#]]$$

$$\text{refine}(x : \tau, \#t)(\sigma^\#) = \sigma^\#[x \mapsto \tau_x^\# \sqcap \{\tau\}]$$

$$\text{refine}(x : \tau, \#f)(\sigma^\#) = \sigma^\#[x \mapsto \tau_x^\# \setminus \{\tau' \mid \tau' <: \tau\}]$$

$$\text{refine}(e, b)(\sigma^\#) = \sigma^\#$$

where  $\sigma_j^\# = \text{refine}(e_j, b)(\sigma^\#)$  for  $j = 0, 1$ ,  $\tau_e^\# = \llbracket e \rrbracket_e^\#(\sigma^\#)$ , and  $[\tau^\#]$  returns  $\{\tau\}$  if  $\tau^\#$  denotes a singleton type  $\tau$ , or returns  $\emptyset$ , otherwise.



# JSTAR - Precision $\uparrow$ - 2) Type Refinement

$$\text{refine}(!e, b)(\sigma^\#) = \text{refine}(e, \neg b)(\sigma^\#)$$

$$\text{refine}(e_0 \sqcup e_1, b)(\sigma^\#) = \begin{cases} \sigma_0^\# \sqcup \sigma_1^\# & \text{if } b \\ \sigma_0^\# \sqcap \sigma_1^\# & \text{if } \neg b \end{cases}$$

$$\text{refine}(e_0 \& e_1, b)(\sigma^\#) = \begin{cases} \sigma_0^\# \sqcap \sigma_1^\# & \text{if } b \\ \sigma_0^\# \sqcup \sigma_1^\# & \text{if } \neg b \end{cases}$$

$$\text{refine}(x.\text{Type} == c_{\text{normal}}, \#t)(\sigma^\#) = \sigma^\# [x \mapsto \tau_x^\# \sqcap \text{normal}(\mathbb{T})]$$

$$\text{refine}(x.\text{Type} == c_{\text{normal}}, \#f)(\sigma^\#) = \sigma^\# [x \mapsto \tau_x^\# \sqcap \{\text{abrupt}\}]$$

$$\text{refine}(x == e, \#t)(\sigma^\#) = \sigma^\# [x \mapsto \tau_x^\# \sqcap \tau_e^\#]$$

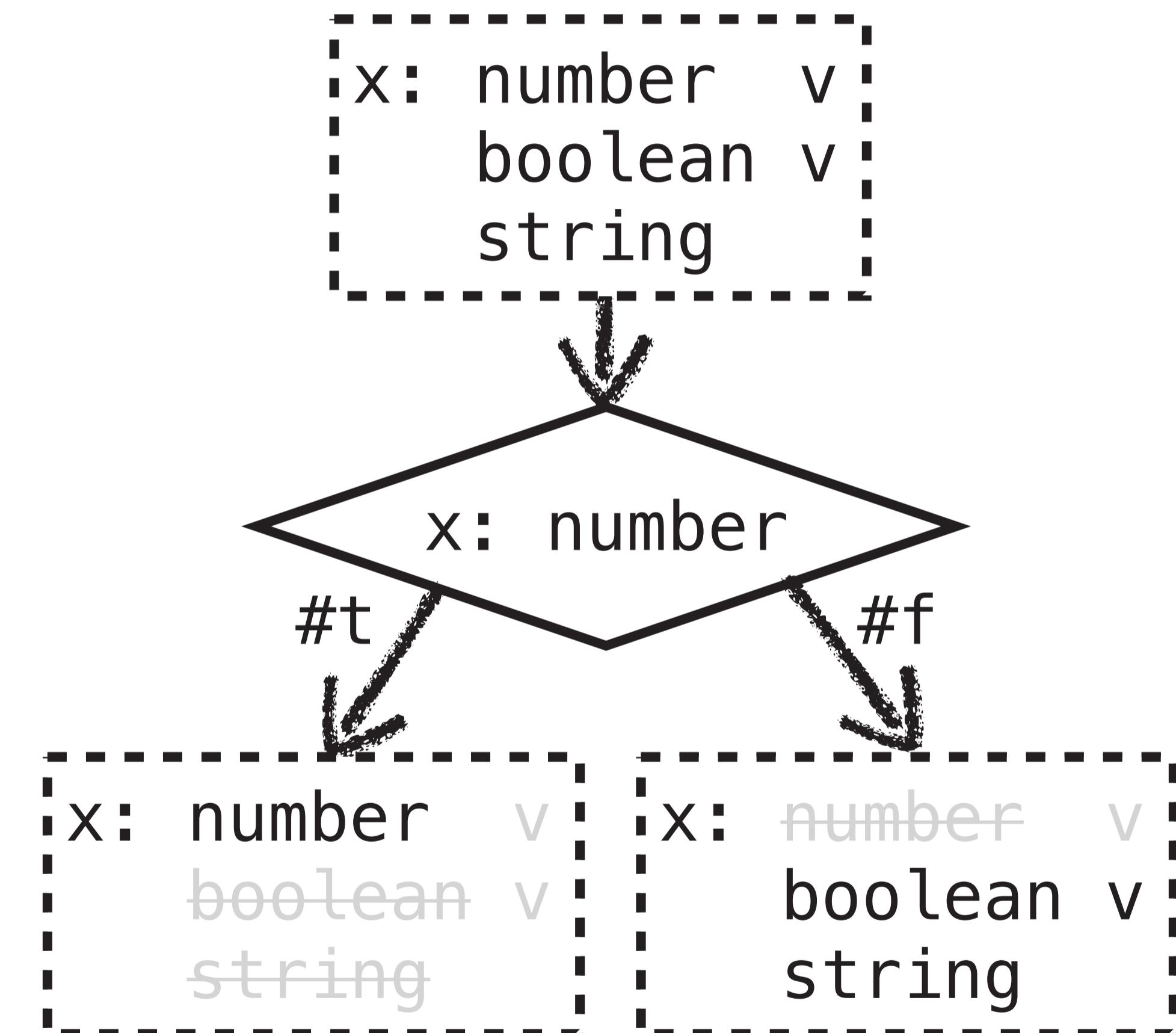
$$\text{refine}(x == e, \#f)(\sigma^\#) = \sigma^\# [x \mapsto \tau_x^\# \setminus [\tau_e^\#]]$$

$$\text{refine}(x : \tau, \#t)(\sigma^\#) = \sigma^\# [x \mapsto \tau_x^\# \sqcap \{\tau\}]$$

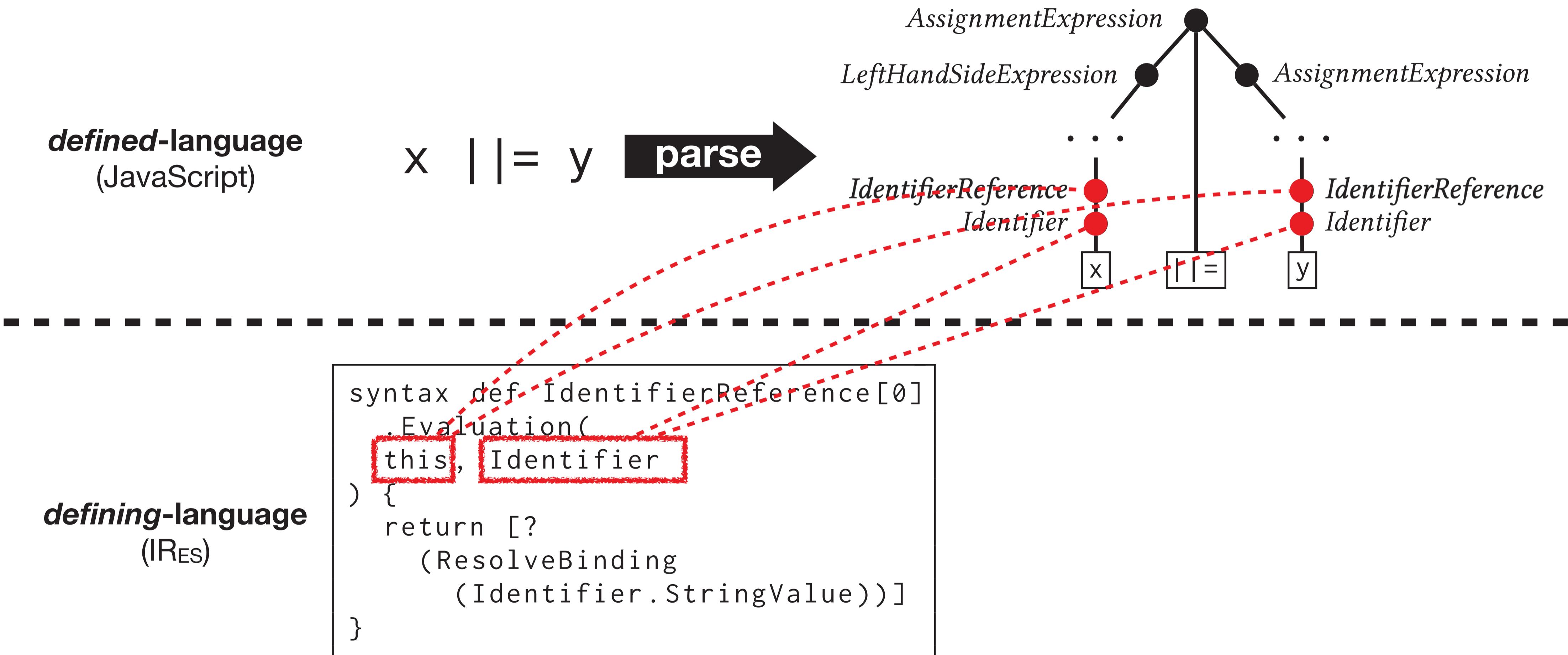
$$\text{refine}(x : \tau, \#f)(\sigma^\#) = \sigma^\# [x \mapsto \tau_x^\# \setminus \{\tau' \mid \tau' <: \tau\}]$$

$$\text{refine}(e, b)(\sigma^\#) = \sigma^\#$$

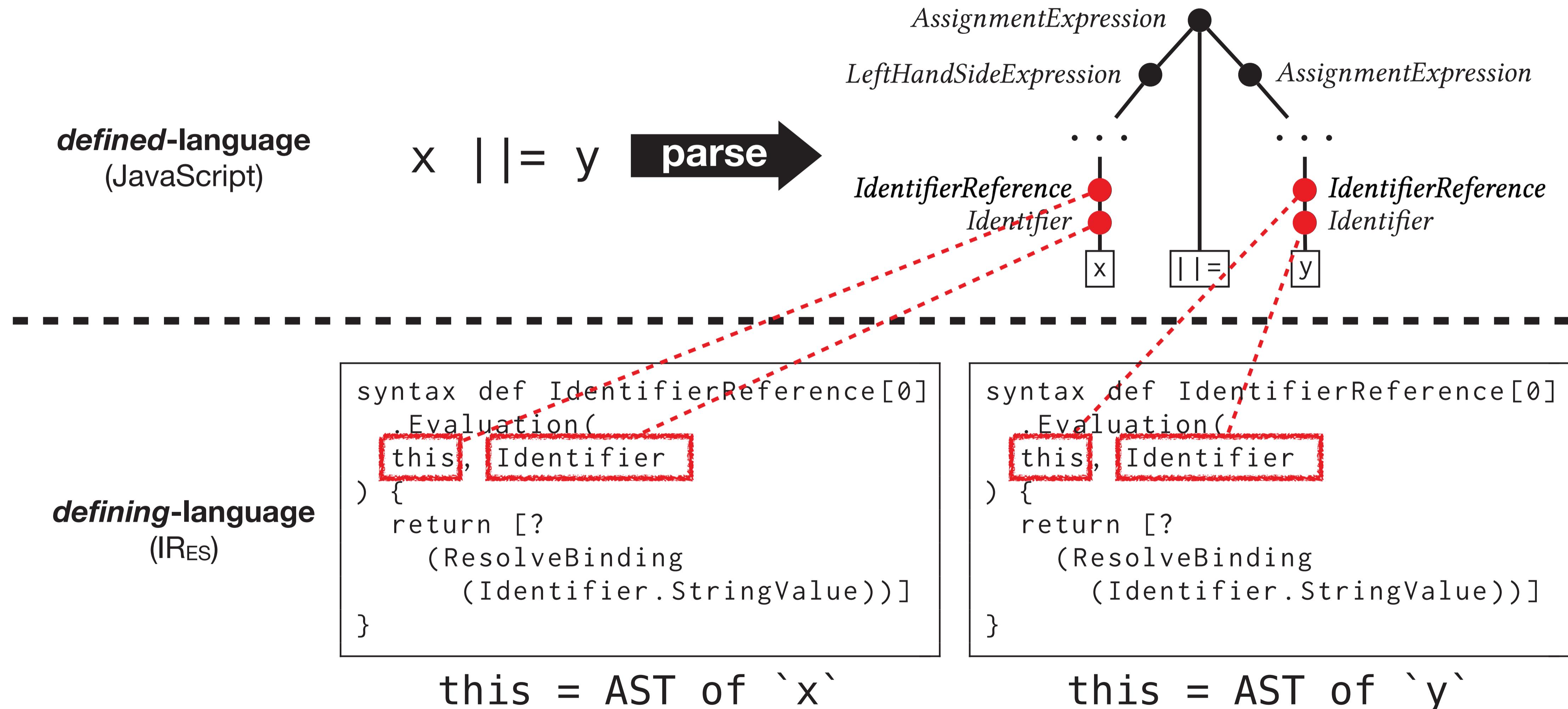
where  $\sigma_j^\# = \text{refine}(e_j, b)(\sigma^\#)$  for  $j = 0, 1$ ,  $\tau_e^\# = \llbracket e \rrbracket_e^\#(\sigma^\#)$ , and  $[\tau^\#]$  returns  $\{\tau\}$  if  $\tau^\#$  denotes a singleton type  $\tau$ , or returns  $\emptyset$ , otherwise.



# JSAVER - AST Sensitivity



# JSAVER - AST Sensitivity



# JSAVER - AST Sensitivity

defined-language (JavaScript)	defining-language (IR <sub>ES</sub> )
flow-sensitivity	$\delta^{\text{js-flow}}(t_{\perp}) = \{\sigma = (\_, \_, \bar{c}, \_) \in \mathbb{S} \mid \text{ast}(\bar{c}) = t_{\perp}\}$
k-callsite sensitivity	$\delta^{\text{js-}k\text{-cfa}}([t_1, \dots, t_n]) = \{\sigma = (\_, \_, \bar{c}, \_) \in \mathbb{S} \mid n \leq k \wedge (n = k \vee \text{js-ctxt}^{n+1}(\bar{c}) = \perp) \wedge \forall 1 \leq i \leq n. \text{ast} \circ \text{js-ctxt}^i(\bar{c}) = t_i\}$