

Lecture 2 – Syntax and Semantics (1)

COSE212: Programming Languages

Jihyeok Park



2025 Fall

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1. Programming Languages

2. Syntax

- Concrete Syntax

- Abstract Syntax

- Concrete vs. Abstract Syntax

3. Operational Semantics

- Inference Rules

- Big-Step Operational (Natural) Semantics

- Small-Step Operational (Reduction) Semantics

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Definition (Programming Language)

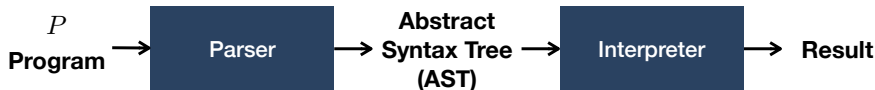
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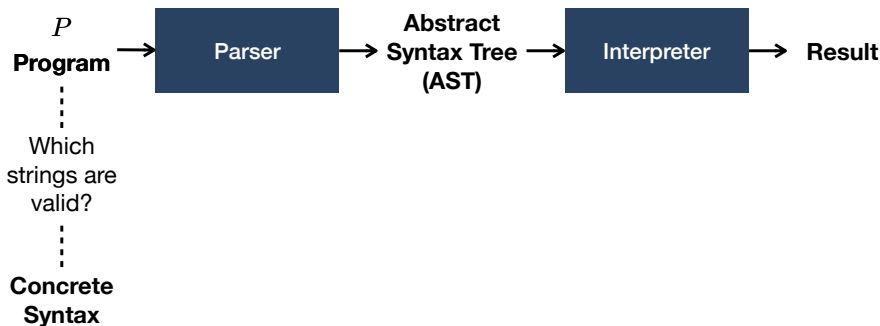
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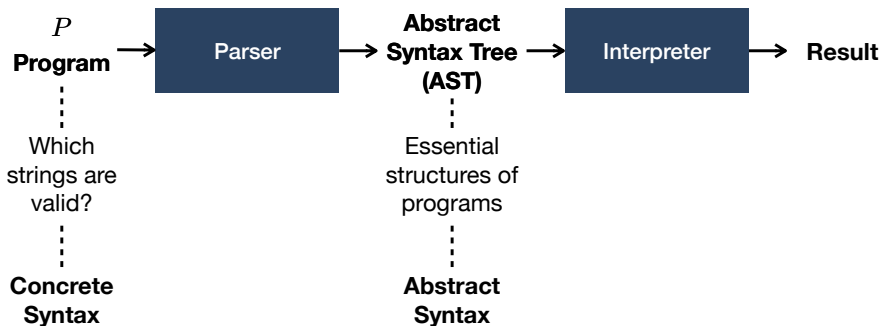
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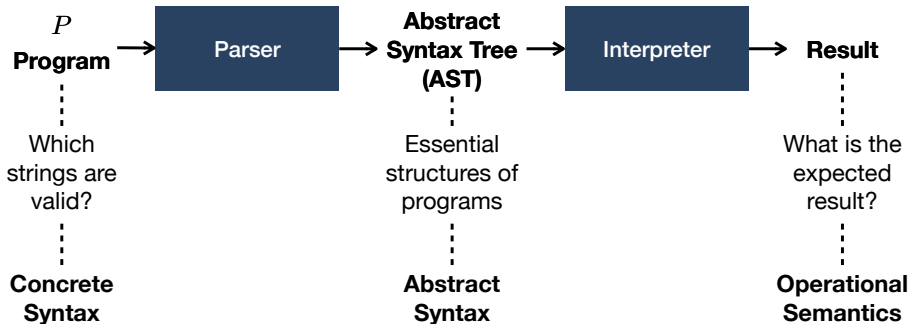
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For example, let's consider the arithmetic expressions (AE) supporting **addition** and **multiplication** of number (integer) values.

- $4 + 2$
- $1 * 24$
- $-42 + 4 * 10$
- $(1 + 2) * (2 + 3)$
- ...

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What does parsing result of each AE look like? – (**abstract syntax**)

What is the evaluation result of each AE? – (**operational semantics**)

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We use the different notation for concrete and abstract syntax:

Description	Concrete Syntax	Abstract Syntax
Terminal	"a"	a
Nonterminal	<expr>	e
Optional	<expr>?	e [?]
Zero or more repetition	<expr>*	e [*]
One or more repetition	<expr>+	e ⁺

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For example, we can define a concrete syntax of integers as follows:

```
<digit> ::= "0" | "1" | "2" | "3" | "4" | "5" | "6" | "7" | "8" | "9"  
<number> ::= "-"? <digit>+
```

Let's define the **concrete syntax** of AE in BNF:

```
<expr> ::= <number>  
         | <expr> "+" <expr>  
         | <expr> "*" <expr>  
         | "(" <expr> ")"
```

It is the **surface-level** representation of programs with all the syntactic details to decide whether a given string is a valid AE or not.

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For example, $(1+2)*3$ is a valid AE:

$\langle \text{expr} \rangle \Rightarrow \langle \text{expr} \rangle * \langle \text{expr} \rangle$	$\Rightarrow (\langle \text{expr} \rangle) * \langle \text{expr} \rangle$
$\Rightarrow (\langle \text{expr} \rangle + \langle \text{expr} \rangle) * \langle \text{expr} \rangle$	$\Rightarrow (\langle \text{number} \rangle + \langle \text{expr} \rangle) * \langle \text{expr} \rangle$
$\Rightarrow (1 + \langle \text{expr} \rangle) * \langle \text{expr} \rangle$	$\Rightarrow (1 + \langle \text{number} \rangle) * \langle \text{expr} \rangle$
$\Rightarrow (1+2) * \langle \text{expr} \rangle$	$\Rightarrow (1+2) * \langle \text{number} \rangle$
$\Rightarrow (1+2) * 3$	

Let's define the **concrete syntax** of AE in BNF:

```
<expr> ::= <number>
         | <expr> "+" <expr>
         | <expr> "*" <expr>
         | "(" <expr> ")"
```

We need **associativity** and **precedence** rules to remove ambiguity:

- "+" and "*" are **left-associative**.

```
"1 + 2 + 3" == "(1 + 2) + 3"
"1 * 2 * 3" == "(1 * 2) * 3"
```

- "*" has higher **precedence** than "+".

```
"1 + 2 * 3" == "1 + (2 * 3)"
```

Let's define the **abstract syntax** of AE in BNF:

Numbers	$n \in \mathbb{Z}$	(BigInt)
Expressions	$e ::= n$	(Num)
	$e + e$	(Add)
	$e * e$	(Mul)

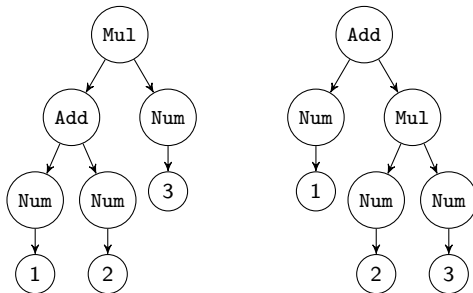
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The **abstract syntax trees (ASTs)** of "(1+2)*3" and "1+2*3":



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There might be **multiple** concrete syntax for the **same** abstract syntax:

```
<expr> ::= <number>
        | <expr> "+" <expr>
        | <expr> "*" <expr>
        | "(" <expr> ")"
```

```
<expr> ::= <number>
        | "(" "+" <expr> <expr> ")"
        | "(" "*" <expr> <expr> ")"
```

```
<expr> ::= <number>
        | "ADD[" <expr> ";" <expr> "]"
        | "MUL[" <expr> ";" <expr> "]"
```

$n \in \mathbb{Z}$	(BigInt)
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While **concrete syntax** is the **surface-level** representation of programs, **abstract syntax** captures the **essential structure** of programs.

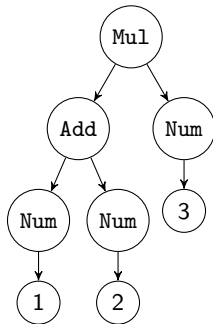
There might be **multiple** concrete syntax for the **same** abstract syntax:

$(1 + 2) * 3$

$(* (+ 1 2) 3)$

\Rightarrow

MUL[ADD[1; 2]; 3]



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Small-Step Operational (Reduction) Semantics

There exist diverse ways to define **semantics** of programming languages.

- **Axiomatic semantics** defines the meaning of a program by specifying the properties that hold after its execution.

$$\{x = n \wedge y = m\} \quad z = x + y \quad \{z = n + m\}$$

- **Denotational semantics** defines the meaning of a program by mapping it to a mathematical object that represents its meaning.

$$\llbracket e_1 + e_2 \rrbracket = \llbracket e_1 \rrbracket + \llbracket e_2 \rrbracket$$

- **Operational semantics** defines the meaning of a program by specifying how it executes on a machine.

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash e_1 + e_2 \Rightarrow n_1 + n_2}$$

- ...

In this course, we will focus on **operational semantics**,

In this course, we will focus on **operational semantics**, and there are two different representative styles:

- **Big-Step Operational (Natural) Semantics** defines the meaning of a program by specifying how it executes on a machine in **one big step**.

$$\frac{\dots}{\vdash e \Rightarrow n}$$

(The execution result of an expression e is n because of)

- **Small-Step Operational (Reduction) Semantics** defines the meaning of a program by specifying how it executes on a machine **step-by-step**.

$$e \rightarrow e' \rightarrow e'' \rightarrow \dots \rightarrow n$$

(An expression e is reduced to e' , then to e'' , and so on until n .)

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meaning that “if all the premises are true, then the conclusion is true”:

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meaning that “*if all the premises are true, then the conclusion is true*”:

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For example,

$$\frac{A \implies B \quad B \implies C}{A \implies C}$$

means that “*if A implies B, and B implies C, then A implies C*”.

(Syllogism – 삼단논법)

$$\boxed{\vdash e \Rightarrow n}$$

It means that “*the expression e evaluates to the number n* ”.

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Let's define the **big-step operational (natural) semantics** of AE:

$$\begin{array}{lcl}
 e ::= n & (\text{Num}) & \\
 | e + e & (\text{Add}) & \\
 | e * e & (\text{Mul}) &
 \end{array}
 \quad \Longrightarrow \quad
 \begin{array}{c}
 \text{NUM} \frac{}{\vdash n \Rightarrow n} \\
 \\
 \text{ADD} \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash e_1 + e_2 \Rightarrow n_1 + n_2} \\
 \\
 \text{MUL} \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash e_1 * e_2 \Rightarrow n_1 \times n_2}
 \end{array}$$

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Let's prove $\vdash (1 + 2) * 3 \Rightarrow 9$ by drawing a **derivation tree**:

$$\begin{array}{l} \text{NUM} \frac{}{\vdash n \Rightarrow n} \quad \text{ADD} \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash e_1 + e_2 \Rightarrow n_1 + n_2} \quad \text{MUL} \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash e_1 * e_2 \Rightarrow n_1 \times n_2} \end{array}$$

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Let's prove $\vdash 1 + (2 * 3) \Rightarrow 7$ by drawing a **derivation tree**:

$$\vdash 1 + (2 * 3) \Rightarrow$$

$$e_0 \rightarrow e_1$$

It means that “ e_0 is reduced to e_1 as the result of one-step evaluation”.

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Let's define the **small-step operational (reduction) semantics** of AE:

$$\begin{array}{lcl}
 e ::= & n & \text{(Num)} \\
 & | & e + e \quad \text{(Add)} \\
 & | & e * e \quad \text{(Mul)}
 \end{array}
 \implies
 \begin{array}{c}
 \frac{e_1 \rightarrow e'_1}{e_1 + e_2 \rightarrow e'_1 + e_2} \quad \frac{e_1 \rightarrow e'_1}{e_1 * e_2 \rightarrow e'_1 * e_2} \\
 \\
 \frac{e_2 \rightarrow e'_2}{n_1 + e_2 \rightarrow n_1 + e'_2} \quad \frac{e_2 \rightarrow e'_2}{n_1 * e_2 \rightarrow n_1 * e'_2} \\
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 \frac{}{n_1 + n_2 \rightarrow n_1 + n_2} \quad \frac{}{n_1 * n_2 \rightarrow n_1 \times n_2}
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Let's prove $(1 + 2) * 3 \rightarrow^* 9$ by showing a **reduction sequence**:

(Note that \rightarrow^* denotes the reflexive-transitive closure of \rightarrow .)

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Let's prove $1 + (2 * 3) \rightarrow^* 7$ by showing a **reduction sequence**:

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(See the language specification of AE.¹)

¹<https://github.com/ku-plrg-classroom/docs/blob/main/cose212/ae/ae-spec.pdf>

- Syntax and Semantics (2)

Jihyeok Park

jihyeok_park@korea.ac.kr

<https://plrg.korea.ac.kr>