Lecture 1 – Mathematical Preliminaries

COSE215: Theory of Computation

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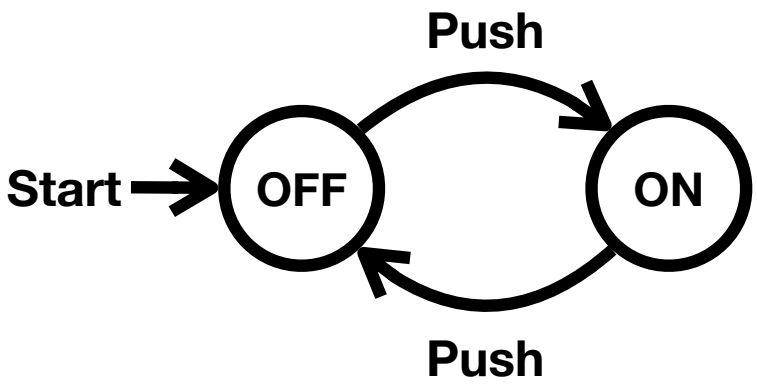
2024 Spring

Recall



Theorem

The current state is OFF *if and only if the button is pushed* **even** *times.*

- Is it possible to prove it?

Let's learn **mathematical background and notation.**

Contents

1. Mathematical Notations
    - Notations in Logics
    - Notations in Set Theory
2. Inductive Proofs
    - Inductions on Integers
    - Structural Inductions
    - Mutual Inductions
3. Notations in Languages
    - Symbols & Words
    - Languages

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Notations in Logics

Notation	Description
A, B	arbitrary **statements**.
$P(x)$	a **predicate** that involves a **variable** x.
$A \wedge B$	the **conjunction** of A and B. (i.e., "A and B").
$A \vee B$	the **disjunction** of A and B. (i.e., "A or B").
$\neg A$	the **negation** of A. (i.e., "not A").

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$$(\textbf{De Morgan's Laws}) = \begin{cases} \neg(A \wedge B) = \neg A \vee \neg B \\ \neg(A \vee B) = \neg A \wedge \neg B \end{cases}$$

$$(\textbf{Truth Table}) =$$

A	B	$\neg(A \wedge B)$	$\neg A \vee \neg B$
T	T	F	F
T	F	T	T
F	T	T	T
F	F	T	T

Notations in Logics

Notation	Description
$A \Rightarrow B$	the **implication** of A and B (i.e., "if A then B" or "A implies B") (i.e., $\neg A \vee B$).
$A \Leftrightarrow B$	A **if and only if (iff)** B (i.e., $A \Rightarrow B \wedge B \Rightarrow A$).
$\forall x \in X.\ P(x)$	the **universal quantifier** (i.e., "for all x in X, $P(x)$ holds").
$\exists x \in X.\ P(x)$	the **existential quantifier** (i.e., "there exists x in X such that $P(x)$ holds").

Notations in Set Theory

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- A **proper subset** X of a set Y is denoted by $X \subset Y$.

$$X \subset Y \iff X \subseteq Y \wedge X \neq Y$$

Notations in Set Theory

- The **union** of sets

$$X \cup Y = \{x \mid x \in X \vee x \in Y\}$$
$$\bigcup \mathcal{C} \quad = X_1 \cup X_2 \cup \cdots \cup X_n = \{x \mid \exists X \in \mathcal{C}.\ x \in X\}$$

where $\mathcal{C} = \{X_1, X_2, \cdots, X_n\}$.

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- The **difference** of sets

$$X \setminus Y = \{x \mid x \in X \wedge x \notin Y\}$$

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$$2^X = \mathcal{P}(X) = \{Y \mid Y \subseteq X\}$$

- The **Cartesian product** of sets X and Y is denoted by $X \times Y$.

$$X \times Y = \{(x, y) \mid x \in X \land y \in Y\}$$

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Inductions on Integers

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Definition (Inductions on Integers)

Let $P(n)$ be a predicate on integers, and if

- **(Basis Case)** $P(k)$ holds where k is an integer, and
- **(Induction Case)** for all integer $n \geq k$, $P(n) \Rightarrow P(n+1)$,

then $P(i)$ holds for all $i \geq k$.

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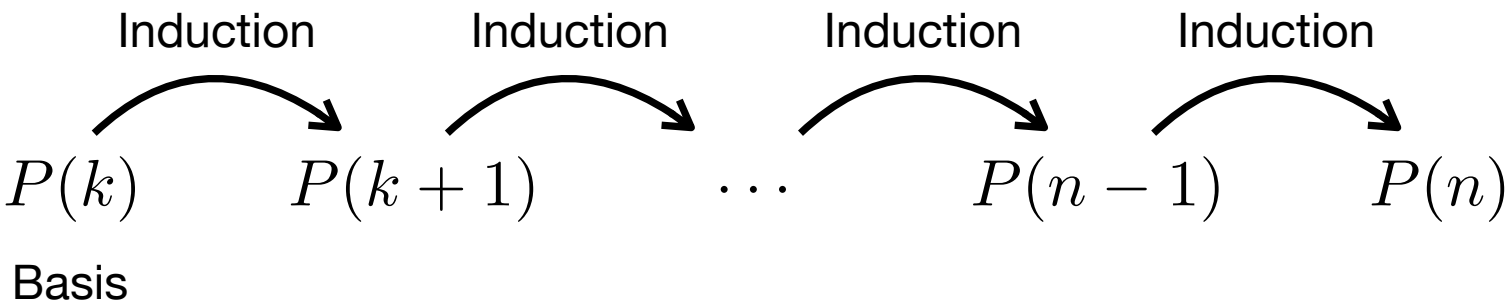

Inductions on Integers – Examples

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Example

Prove that $\forall n \geq 0.\ \sum_{i=0}^{n} i = \frac{n(n+1)}{2}$.

Proof)

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- (Basis Case): $0 = 0(0+1)/2$ $\square$

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- (Basis Case): $0 = 0(0+1)/2$ $\square$
- (Induction Case): Assume that it holds for n (I.H.). Then,

$$\begin{aligned}\sum_{i=0}^{n+1} i &= (n+1) + \sum_{i=0}^{n} i \\ &= (n+1) + \frac{n(n+1)}{2} \qquad (\because I.H.) \\ &= \frac{(n+1)(n+2)}{2} \quad \square\end{aligned}$$

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$$\begin{aligned}\sum_{i=0}^{n+1} i^2 &= (n+1)^2 + \sum_{i=0}^{n} i^2 \\ &= (n+1)^2 + \frac{n(n+1)(2n+1)}{6} \qquad (\because I.H.) \\ &= \frac{(n+1)(n+2)(2(n+1)+1)}{6}\end{aligned}$$ □

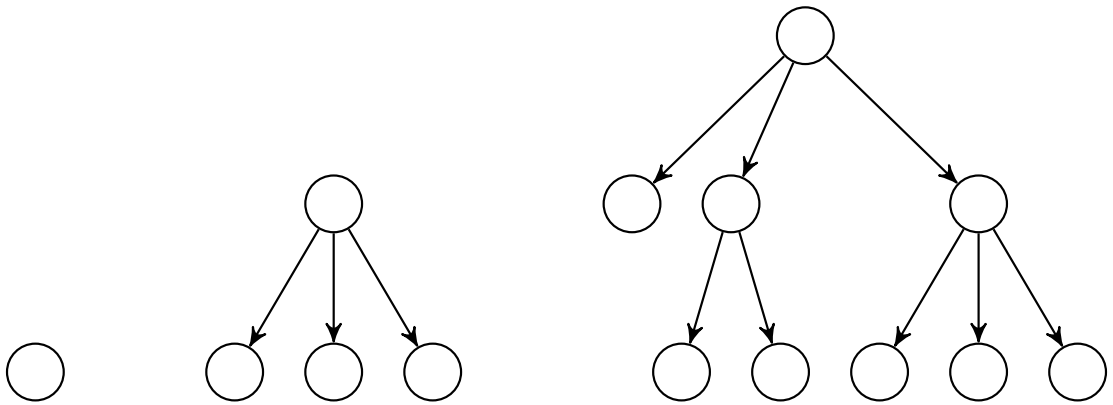
Structural Inductions – Inductive Definitions

In CS, we often define somethings as **inductively-defined sets**.
For example, we can define **trees** as follows:

Example (Inductive Definition of Trees)

A **tree** is defined as follows:

- **(Basis Case)** A single **node** N is a tree.
- **(Induction Case)** If $T_1, \cdots, T_n$ are trees, then a graph defined with a new root node N and edges from N to $T_1, \cdots, T_n$ is a tree.



Structural Inductions – Inductive Definitions

Another example is a set of **arithmetic expressions**:

Example (Inductive Definition of Arithmetic Expressions)

An **arithmetic expression** is defined as follows:

- **(Basis Case)** A **number** or a **variable** is an arithmetic expression.
- **(Induction Case)** If E and F are arithmetic expressions, then so are E+F, E*F, and (E).

```
42          x          x + y
42 * x      (x)        (x * y) * z
(2 + x) * y  x * (x * y)  ((((x))))
```

Structural Inductions

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Definition (Structural Inductions)

Let $P(x)$ be a predicate on a **inductively-defined set** X, and if

- **(Basis Case)** $P(b_1), \cdots, P(b_k)$ hold for all basis cases $b_1, \cdots, b_k$.
- **(Induction Case)** for all $x \in X$,

$$P(x_1) \wedge \cdots \wedge P(x_n) \Rightarrow P(x)$$

where $x_1, \cdots, x_n$ are the **sub-structures** of x.

then $P(x)$ holds for all $x \in X$.

$P(x_1), \cdots, P(x_n)$ are called **induction hypotheses**.

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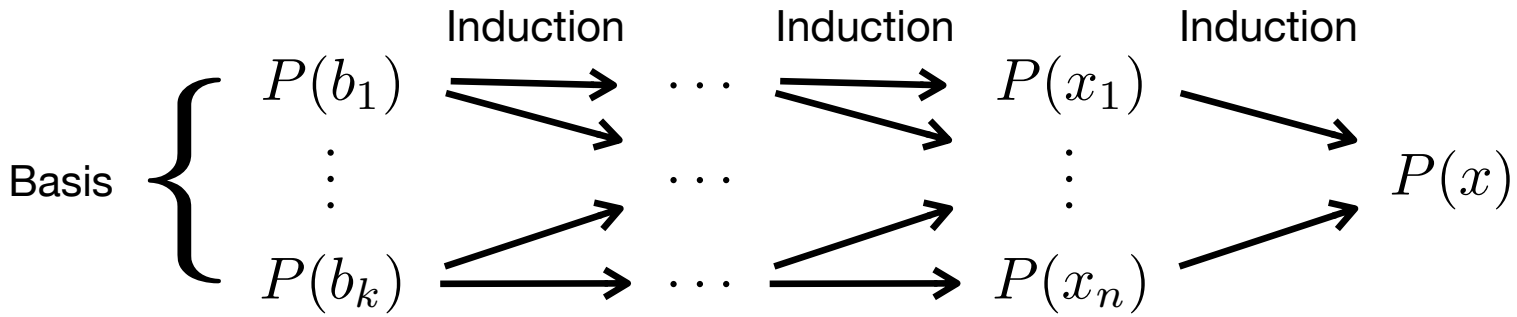

Structural Inductions – Examples

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Prove that for all tree T, the number of nodes in T is equal to the number of edges in T plus one.

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Prove that for all tree T, the number of nodes in T is equal to the number of edges in T plus one.

Proof) Let $N(T)$ be the number of node and $E(T)$ be the number of edges in T. Let's prove $\forall T.\ N(T) = E(T) + 1$.

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- (Basis Case): $N(T) = 1$ and $E(T) = 0$. □

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- (Basis Case): $N(T) = 1$ and $E(T) = 0$. □
- (Induction Case): Assume that it holds for $T_1, \cdots, T_n$ (I.H.). Then,

$$
\begin{aligned}
N(T) &= 1 + \sum_{i=1}^{n} N(T_i) \\
&= 1 + \sum_{i=1}^{n} (E(T_i) + 1) \qquad (\because I.H.) \\
&= 1 + n + \sum_{i=1}^{n} E(T_i) \\
&= 1 + E(T) \quad \square
\end{aligned}
$$

Structural Inductions – Examples

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Example

Prove that for all arithmetic expression E, the number of left parentheses in E is equal to the number of right parentheses in E.

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Prove that for all arithmetic expression E, the number of left parentheses in E is equal to the number of right parentheses in E.

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- (Induction Case): Assume that it holds for E and F (I.H.). Then,

$$
\begin{aligned}
L(E{+}F) &= L(E) + L(F) = R(E) + R(F) \qquad (\because \mathit{I.H.}) \\
&= R(E{+}F) \quad \square \\
L(E{*}F) &= L(E) + L(F) = R(E) + R(F) \qquad (\because \mathit{I.H.}) \\
&= R(E{*}F) \quad \square \\
L((E)) &= L(E) + 1 = R(E) + 1 \qquad (\because \mathit{I.H.}) \\
&= R((E)) \quad \square
\end{aligned}
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Mutual Inductions

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Let $P(x)$ and $Q(x)$ are predicates on integers, and if

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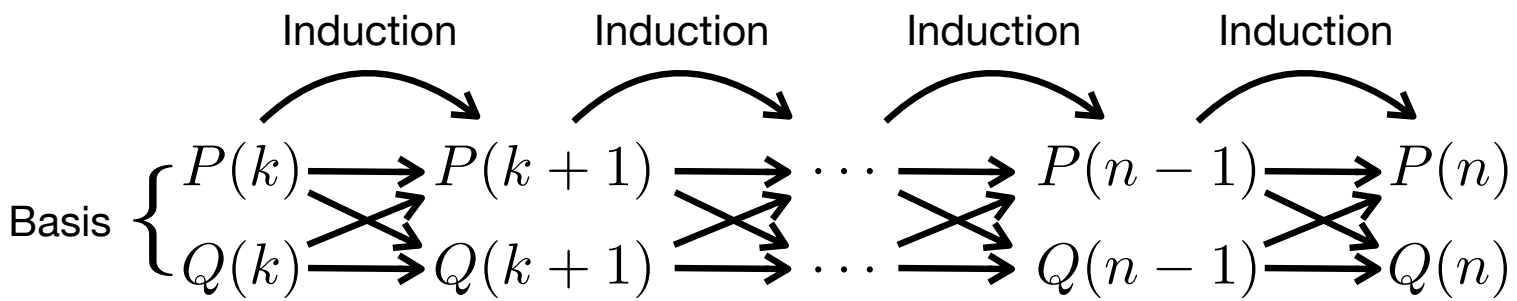

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Theorem

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The current state is OFF *if and only if the button is pushed* **even** *times, and the current state is* ON *if and only if the button is pushed* **odd** *times.*

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Proof) Let $S(i)$ be the current state after i times of pushing. Let's prove

$$\forall i \geq 0.\ S(i) = \text{OFF} \iff i \equiv 0 \pmod 2 \qquad (P)$$

$$\forall i \geq 0.\ S(i) = \text{ON} \iff i \equiv 1 \pmod 2 \qquad (Q)$$

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- (Basis Case): Known facts: $S(0) = \text{OFF}$ and $0 \equiv 0 \pmod 2$

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- (Basis Case): Known facts: $S(0) = \text{OFF}$ and $0 \equiv 0 \pmod 2$
  - $(P, \Rightarrow)$: $0 \equiv 0 \pmod 2 \implies S(0) = \text{OFF} \Rightarrow 0 \equiv 0 \pmod 2$

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- (Basis Case): Known facts: $S(0) = \text{OFF}$ and $0 \equiv 0 \pmod 2$
  - $(P, \Rightarrow)$: $0 \equiv 0 \pmod 2 \implies S(0) = \text{OFF} \Rightarrow 0 \equiv 0 \pmod 2$
  - $(P, \Leftarrow)$: $S(0) = \text{OFF} \implies S(0) = \text{OFF} \Leftarrow 0 \equiv 0 \pmod 2$

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Proof) Let $S(i)$ be the current state after i times of pushing. Let's prove

$$\forall i \geq 0.\ S(i) = \text{OFF} \iff i \equiv 0 \pmod 2 \qquad (P)$$
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- (Basis Case): Known facts: $S(0) = \text{OFF}$ and $0 \equiv 0 \pmod 2$
  - $(P, \Rightarrow)$: $0 \equiv 0 \pmod 2 \implies S(0) = \text{OFF} \Rightarrow 0 \equiv 0 \pmod 2$
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Mutual Inductions – Examples

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1. Mathematical Notations
    - Notations in Logics
    - Notations in Set Theory

2. Inductive Proofs
    - Inductions on Integers
    - Structural Inductions
    - Mutual Inductions

3. Notations in Languages
    - Symbols & Words
    - Languages

Symbols & Words

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  - $\Sigma = \{a \mid a \text{ is an Unicode character}\}$ – Unicode characters.

$$\epsilon, 안녕하세요, こんにちは, ★■▲⊕, \cdots \in \Sigma^*$$

Symbols & Words

Notation	Description
ϵ	the **empty word**.
$w_1 w_2$	the **concatenation** of w_1 and w_2. (w_1 is a **prefix** of $w_1 w_2$ and w_2 is a **suffix** of $w_1 w_2$)
w^R	the **reverse** of w.
$\lvert w \rvert$	the **length** of w.
Σ^k	the set of all words of length k.
Σ^*	the set of all words (the **Kleene star**). (i.e., $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \cdots = \bigcup_{k \geq 0} \Sigma^k$)
Σ^+	the set of all words except ϵ (the **Kleene plus**). (i.e., $\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \cdots = \bigcup_{k \geq 1} \Sigma^k$)

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- $L = \{10, 11, 101, 111, 1011, \cdots\}$ – ???

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- The **union**, **intersection**, and **difference** of languages:

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- The **concatenation** of languages:

$$L_1 L_2 = \{w_1 w_2 \mid w_1 \in L_1 \wedge w_2 \in L_2\}$$

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Next Lecture

- Basic Introduction of Scala

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