Lecture 4 – Identifiers (1)

COSE212: Programming Languages

Jihyeok Park



2025 Fall





• ADT for Abstract Syntax of AE

```
enum Expr:
    case Num(number: BigInt)
    case Add(left: Expr, right: Expr)
    case Mul(left: Expr, right: Expr)
```

• Parser for Concrete Syntax of AE

```
lazy val expr: P[Expr] = ...
```

• Interpreter for Semantics of AE

```
def interp(expr: Expr): Value = ...
```





ADT for Abstract Syntax of AE

```
enum Expr:
   case Num(number: BigInt)
   case Add(left: Expr, right: Expr)
   case Mul(left: Expr, right: Expr)
```

• Parser for Concrete Syntax of AE

```
lazy val expr: P[Expr] = ...
```

• Interpreter for Semantics of AE

```
def interp(expr: Expr): Value = ...
```

• In this lecture, we will learn identifiers.

Contents



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE - AE with Variables

Concrete Syntax Abstract Syntax Examples

Contents



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE – AE with Variables Concrete Syntax

Abstract Syntax





An identifier is a name for a certain element in a program.

In Scala, there are diverse kinds of identifiers:

```
/* Scala */
// variable names
val x: Int = 42
// function and parameter names
def f(a: Int, b: Int): Int = a + b
// class and field names
case class Person(name: String, age: Int)
```



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
   val x: Int = a + b
   x + add(y, z)

add(x, b)
```

- A binding occurrence of an identifier denotes its definition site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



- A binding occurrence of an identifier denotes its definition site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
 val x: Int = a + b
 x + add(y, z)

add(x, b)
```

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
   val x: Int = a + b
   x + add(y, z)

add(x, b)
```

- A binding occurrence of an identifier denotes its definition site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



- A **binding occurrence** of an identifier denotes its **definition** site.
- A **scope** of an identifier denotes where the identifier is **usable**.
- A bound occurrence of an identifier denotes its lookup site.



A **bound identifier** is an identifier that is **defined** in a program.

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
   val x: Int = a + b
   x + add(y, z)

add(x, b)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
   val x: Int = a + b
   x + add(y, z)

add(x, b)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
    val x: Int = a + b
    x + add(y, z)

add(x, b)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A **bound occurrence** of an identifier denotes its **lookup** site.



A **bound identifier** is an identifier that is **defined** in a program.

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
   val x: Int = a + b
   x + add(y, z)

add(x, b)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A binding occurrence of an identifier denotes its definition site.
- A scope of an identifier denotes where the identifier is usable.
- A **bound occurrence** of an identifier denotes its **lookup** site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =
  val x: Int = a + b
  x + add(y, z)

add(x, b)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.



```
/* Scala */
val x: Int = 3

val y: Int = x + z

def add(a: Int, b: Int): Int =

val x Int = a + b

x + add(y, z)

add(x, b)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A **binding occurrence** of an identifier denotes its **definition** site.
- A scope of an identifier denotes where the identifier is usable.
- A bound occurrence of an identifier denotes its lookup site.

Free Identifiers



A **free identifier** is an identifier that is **not defined** in the current scope of the program.

Shadowing



```
/* Scala */ Shadowed
val x: Int = 3 Identifier

val y: Int = x + z

Shadowing
def add(a: Int, b: Int): Int =
val x: Int = a + b
x + add(y, z)

Shadowing
add(x, b) Identifier
```

Shadowing means that the innermost binding occurrence shadows the outer binding occurrences of the same name.

- A shadowing identifier is an identifier that shadows another
- A shadowed identifier is an identifier that is shadowed by another.

Shadowing



```
/* Scala */ Shadowed
val x: Int = 3 Identifier

val y: Int = x + z

Shadowing
def add(a: Int, b: Int): Int =
val x: Int = a + b
x + add(y, z)

Shadowing
add(x, b) Identifier
```

Shadowing means that the innermost binding occurrence shadows the outer binding occurrences of the same name.

- A shadowing identifier is an identifier that shadows another
- A **shadowed identifier** is an identifier that is shadowed by another.

Note that shadowing is **NOT** a mutation.

Contents



1. Identifiers

Bound Identifiers
Free Identifiers
Shadowing

2. VAE - AE with Variables

Concrete Syntax Abstract Syntax Examples





Now, we want to extend AE into VAE with variables:

```
/* VAE */
val x = 1 + 2; // x = 1 + 2 = 3
val y = x + 3; // y = x + 3 = 3 + 3 = 6
y + 4 // 6 + 4 = 10
```





Now, we want to extend AE into VAE with variables:

```
/* VAE */
val x = 1 + 2; // x = 1 + 2 = 3
val y = x + 3; // y = x + 3 = 3 + 3 = 6
y + 4 // 6 + 4 = 10
```

First, we define the **concrete syntax** of **identifiers** used in VAE:



Now, we want to extend AE into VAE with variables:

```
/* VAE */
val x = 1 + 2; // x = 1 + 2 = 3
val y = x + 3; // y = x + 3 = 3 + 3 = 6
y + 4 // 6 + 4 = 10
```

First, we define the **concrete syntax** of **identifiers** used in VAE:

For example, the following are valid identifiers:

x y get_name getName add42

Concrete Syntax



Then, let's define the **concrete syntax** of VAE in BNF:

Concrete Syntax



Then, let's define the **concrete syntax** of VAE in BNF:

Note that each variable definition creates a **new scope**.





Then, let's define the **concrete syntax** of VAE in BNF:

Note that each variable definition creates a **new scope**. For example:

```
/* VAE */
val x = 1 + 2;
val y = x + 3;
y + 4
```





Then, let's define the **concrete syntax** of VAE in BNF:

Note that each variable definition creates a **new scope**. For example:

```
/* VAE */
val x = 1 + 2;
val y = x + 3;
y + 4
```

means

```
/* VAE */
val x = 1 + 2;
{ // scope of x
 val y = x + 3;
 { // scope of y
 y + 4
 }
}
```

Abstract Syntax



Let's define the **abstract syntax** of VAE in BNF:





Let's define the **abstract syntax** of VAE in BNF:

We can define an **ADT** for the abstract syntax of VAE in Scala:

```
enum Expr:
    case Num(number: BigInt)
    case Add(left: Expr, right: Expr)
    case Mul(left: Expr, right: Expr)
    // variable definition
    case Val(name: String, init: Expr, body: Expr)
    // variable lookup
    case Id(name: String)
```





```
enum Expr:
   case Num(number: BigInt)
   case Add(left: Expr, right: Expr)
   case Mul(left: Expr, right: Expr)
   case Val(name: String, init: Expr, body: Expr)
   case Id(name: String)
```

Parser implementation is given and you don't need to implement it.

You can freely use Expr to parse VAE programs as follows:

```
Expr("val x = 1; x + 2")
// Val("x", Num(1), Add(Id("x"), Num(2)))

Expr("val a = 1; val b = 2; a + b")
// Val("a", Num(1), Val("b", Num(2), Add(Id("a"), Id("b"))))
```





For each VAE program, please draw:

- an arrow from each bound occurrence to its binding occurrence.
- a dotted arrow from each shadowing variable to its shadowed one.
- an X mark on each free variable.

```
/* VAE */
val x = 1; x
```

```
/* VAE */
val x = x + 1;
val y = x * 2;
val x = y + x;
x * z
```

```
/* VAE */
val x = 1;
val y = {
  val x = 2 * x;
  { val y = x; y } + { val y = 3; y }
};
x + y
```

Summary



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE - AE with Variables

Concrete Syntax Abstract Syntax Examples

Next Lecture



• Identifiers (2)

Jihyeok Park
 jihyeok_park@korea.ac.kr
https://plrg.korea.ac.kr