Chapter 1

Library Manifest_NFM

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Require Import Lia.
Require Import Relations.
Require Import Logic.FunctionalExtensionality.
Require Import Lists.List.
Import LISTNOTATIONS.
Require Import String.
Require Import Cop. Copland.
Import Copland. Term.
Require Import Utils. Utilities.
```

 ${\it Manifest}$ defines a single attestation manger and its interconnections. Information includes:

- 1.1 asps: a list of ASPs (measurement operations the AM can preform)
- 1.2 M: can measure relation (other AMs the current AM knows of)
- 1.3 C: context relation (other AMs the current AM depends on)
- 1.4 Policy: local policy specific to the current AM.
- 1.5 Minimally includes privacy policy and may possibly include selection policy
- 1.6 Other information not necessary for reasoning includes:
- 1.7 key simulates public key
- 1.8 address simulates address information and
- 1.9 *tpm* simulates cruft necessary to initialize its TPM

```
Record Manifest := \{
asps : list \ ASP ;
M : list \ Plc ;
C : list \ Plc ;
Policy : ASP \rightarrow Plc \rightarrow \texttt{Prop} ;
\}.
Environment \ \text{is a set of AM's each defined by a } Manifest.
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- 1.10 The domain of an Environment provides names for each Manifest.
- 1.11 Names should be the hash of their public key, but this restriction
- 1.12 is not enforced here.

```
Definition Environment: Type := Plc \rightarrow (option\ Manifest).
Definition e_-empty : Environment := (fun_- \Rightarrow None).
Definition e\_update\ (m:Environment)\ (x:Plc)\ (v:(option\ Manifest)):=
     fun x' \Rightarrow \text{if } plc\_dec \ x \ x' \text{ then } v \text{ else } m \ x'.
Definition System := list Environment.
Definition hasASPe(k:Plc)(e:Environment)(a:ASP):Prop :=
match (e k) with
| None \Rightarrow False
| Some \ m \Rightarrow In \ a \ m.(asps)|
Fixpoint hasASPs(k:Plc)(s:System)(a:ASP):Prop :=
    {\tt match}\ s\ {\tt with}
     | | | \Rightarrow False
     |s1::s2 \Rightarrow (hasASPe \ k \ s1 \ a) \lor (hasASPs \ k \ s2 \ a)
Theorem hasASPe\_dec: \forall k \ e \ a, \{hasASPe \ k \ e \ a\} + \{\tilde{has}ASPe \ k \ e \ a\}.
Proof.
  intros k e a.
  unfold hasASPe.
  destruct (e \ k).
  + induction (asps m).
  ++ auto.
  ++ inverts IHl.
  +++ simpl. left. right. apply H.
  +++ simpl. assert (asp\_dec: \{a=a\theta\} + \{a\neq a\theta\}).
              { repeat decide equality. }
       inverts asp\_dec.
  ++++ left. auto.
  ++++ right. unfold not. intros. inverts H1; auto.
  + auto.
Defined.
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Theorem hasASPs\_dec: \forall k \ e \ a, \{hasASPs \ k \ e \ a\} + \{\tilde{has}ASPs \ k \ e \ a\}.
Proof.
  intros k e a.
  induction e.
  + simpl in *. right. unfold not. intros. apply H.
  + simpl in *. pose proof has ASPe\_dec \ k \ a0 \ a. inverts H.
  ++ left. left. apply H0.
  ++ inverts IHe.
  +++ left. right. apply H.
  +++ right. unfold not. intros. inverts H1; auto.
Defined.
Definition knowsOfe(k:Plc)(e:Environment)(p:Plc):Prop :=
match (e k) with
| None \Rightarrow False
| Some \ m \Rightarrow In \ p \ m.(M)
end.
Print System.
Print Environment.
   Determine if place k within the system s knows of p
Fixpoint knowsOfs(k:Plc)(s:System)(p:Plc):Prop :=
match s with
| | | \Rightarrow False
|s1::ss \Rightarrow (knowsOfe \ k \ s1 \ p) \lor (knowsOfs \ k \ ss \ p)
end.
Theorem knowsOfe\_dec: \forall k \ e \ p, \{(knowsOfe \ k \ e \ p)\} + \{\tilde{knowsOfe} \ k \ e \ p)\}.
Proof.
  intros k e p.
  unfold knowsOfe.
  destruct (e \ k); auto.
  + induction (M m).
  ++ auto.
  ++ assert (H: \{p=a\} + \{p \neq a\}). {repeat decide equality. }
      inversion H.
  +++ simpl. left. auto.
  +++ simpl. inverts IHl; auto. right. unfold not. intros. inverts H2; auto.
Defined.
Theorem knowsOfs\_dec: \forall k \ s \ p, \{(knowsOfs \ k \ s \ p)\} + \{\tilde{(knowsOfs \ k \ s \ p)}\}.
Proof.
     intros k \ s \ p.
     induction s; simpl in *.
    + right. unfold not. intros. inversion H.
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+ pose proof knowsOfe_dec k a p. inverts H.
    ++ left. left. apply H0.
    ++ inverts IHs.
    +++ left. right. apply H.
    +++ right. unfold not. intros. inversion H1; auto.
Defined.
   Determine if place k within the environment e depends on place p (the context relation)
Definition dependsOne (k:Plc)(e:Environment)(p:Plc):Prop :=
match (e k) with
None \Rightarrow False
| Some \ m \Rightarrow In \ p \ m.(C)
   Determine if place k within the system s depends on place p (the context relation)
Fixpoint dependsOns (k:Plc)(s:System)(p:Plc):Prop :=
match s with
| | | \Rightarrow False
|s1::ss \Rightarrow (dependsOne \ k \ s1 \ p) \lor (dependsOns \ k \ ss \ p)
Theorem dependsOne\_dec: \forall k \ e \ p, \{(dependsOne \ k \ e \ p)\} + \{ (dependsOne \ k \ e \ p) \}.
Proof.
  intros k e p.
  unfold dependsOne.
  destruct (e \ k).
  + induction (C m).
  ++ auto.
  ++ simpl. inversion IHl.
  +++ auto.
  ++++ assert (H': \{a=p\}+\{a\neq p\}). {repeat decide equality.} inversion H'.
  ++++ left. left. apply H0.
  ++++ right. unfold not. intros. inversion H1; auto.
  + auto.
Defined.
Theorem dependsOns\_dec: \forall k \ s \ p, \{dependsOns \ k \ s \ p\} + \{\tilde{\ } dependsOns \ k \ s \ p\}.
Proof.
  intros. induction s.
  + simpl. auto.
  + simpl. pose proof dependsOne_dec k a p. inversion IHs.
  ++ left. right. apply H0.
  ++ inversion {\it H.}
  +++ left. left. apply H1.
  +++ right. unfold not. intros. inversion H2; auto.
```

Defined.

Is term t exectuable on the attestation manager named k in environment e? Are ASPs available at the right attestation managers and are necessary communications allowed? Fixpoint executable(t:Term)(k:Plc)(e:Environment):Prop :=match t with $asp \ a \Rightarrow hasASPe \ k \ e \ a$ att $p \ t \Rightarrow knowsOfe \ k \ e \ p \rightarrow executable \ t \ p \ e$ $lseq\ t1\ t2 \Rightarrow executable\ t1\ k\ e\ \land\ executable\ t2\ k\ e$ $bseq - t1 \ t2 \Rightarrow executable \ t1 \ k \ e \wedge executable \ t2 \ k \ e$ $\mid bpar \ _t1 \ t2 \Rightarrow executable \ t1 \ k \ e \land executable \ t2 \ k \ e$ end. Ltac $right_dest_contr \ H := right; unfold \ not; intros \ H; destruct \ H; contradiction.$ Theorem $executable_dec: \forall t \ k \ e, \{(executable \ t \ k \ e)\} + \{\tilde{\ }(executable \ t \ k \ e)\}.$ intros. generalize k. induction t; intros. + unfold executable. apply hasASPe_dec. + simpl. pose $proof\ knowsOfe_dec\ k0\ e\ p$. destruct H. ++ destruct (IHt p). +++ left; auto. +++ right. unfold not. intros; auto. ++ destruct (IHt p). +++ left; auto. +++ left. intros. congruence. + simpl. specialize IHt1 with k0. specialize IHt2 with k0. destruct IHt1,IHt2; try right_dest_contr H. ++ left. split; assumption. + simpl. specialize IHt1 with k0. specialize IHt2 with k0. destruct IHt1,IHt2; try $right_dest_contr\ H.$ ++ left. split; assumption. + simpl. specialize IHt1 with k0. specialize IHt2 with k0. destruct IHt1, IHt2; try $right_dest_contr\ H.$ ++ left. split; assumption. Defined.

Is term t executable on the attestation mnanager named k in

1.13 system s? Are ASPs available at the right attestation managers

1.14 and are necessary communications allowed?

Fixpoint executables(t:Term)(k:Plc)(s:System):Prop :=

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match t with
    asp \ a \Rightarrow hasASPs \ k \ s \ a
    att p \ t \Rightarrow knowsOfs \ k \ s \ p \rightarrow executables \ t \ p \ s
   lseq\ t1\ t2 \Rightarrow executables\ t1\ k\ s \land executables\ t2\ k\ s
   bseq - t1 \ t2 \Rightarrow executables \ t1 \ k \ s \land executables \ t2 \ k \ s
   \mid bpar \perp t1 \mid t2 \Rightarrow executables \mid t1 \mid k \mid s \land executables \mid t2 \mid k \mid s
end.
Ltac prove\_exec :=
     match goal with
     |\vdash \{executables (asp\_)\_\_\} + \{\_\} \Rightarrow unfold \ executables; apply \ hasASPs\_dec
     |IHt1:\_,IHt2:\_\vdash \{executables\_?k?s\}+\{\_\} \Rightarrow \texttt{simpl}; \texttt{specialize}\ IHt1 \ \texttt{with}
k s; specialize IHt2 with k s; destruct IHt1, IHt2; try(left; split; assumption)
     end.
Theorem executables\_dec: \forall t \ k \ s, \{executables \ t \ k \ s\} + \{\tilde{e}executables \ t \ k \ s\}.
intros. generalize k s. induction t; intros; try prove_exec; try right_dest_contr H.
+ simpl. destruct (IHt p s\theta).
++ auto.
++ pose proof knowsOfs_dec k0 s0 p. destruct H.
+++ right. unfold not; intros. intuition.
+++ left. intros. congruence.
Defined.
Notation Rely := "Rely"\% string.
Notation Target := "Target" \% string.
Notation Appraise := "Appraise" \% string.
Notation aspc1 :=
  (ASPC ALL EXTD (asp_paramsC "asp1"%string ["x"%string;"y"%string] Target Tar-
qet)).
Notation aspc2 :=
  (ASPC ALL EXTD (asp_paramsC "asp2"%string ["x"%string] Target Target)).
Inductive rely\_Policy : ASP \rightarrow Plc \rightarrow \texttt{Prop} :=
| p\_aspc1 : \forall p, rely\_Policy aspc1 p.
Inductive tar\_Policy : ASP \rightarrow Plc \rightarrow \texttt{Prop} :=
p\_aspc2: tar\_Policy aspc2 Appraise
\mid p\_SIG : \forall p, tar\_Policy SIG p.
Inductive app\_Policy: ASP \rightarrow Plc \rightarrow \texttt{Prop}:=
\mid p\_HSH : \forall p, app\_Policy HSH p.
Definition e\theta := e_-empty.
Definition e_Rely :=
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 $e_update\ e_empty\ Rely\ (Some\ \{|\ asps:=[aspc1];\ M:=[Target]\ ;\ C:=[]\ ;\ Policy:=rely_Policy\ |\}).$

Definition $e_{-}Targ :=$

 $e_update\ e_empty\ Target\ (Some\ \{|\ asps:=[SIG;\ aspc2];\ M:=[Appraise]\ ;\ C:=[]\ ;\ Policy:=tar_Policy|\}).$

Definition $e_-App :=$

 $e_update\ e_empty\ Appraise\ (Some\ \{|\ asps:=[HSH]\ ;\ M:=[]\ ;\ C:=[Target]\ ;\ Policy:=app_Policy\ |\}).$

Definition $example_sys_1 := [e_Rely; e_Targ; e_App].$

Example ex1: hasASPe Rely e_Rely aspc1.

Proof. unfold hasASPe. simpl. left. reflexivity. Qed.

Example ex2: has ASPe Rely e_{-} Rely $CPY \rightarrow False$.

Proof. unfold has ASPe. simpl. intros. inverts H. inverts H0. auto. Qed.

Example ex3: hasASPs Rely (example_sys_1) aspc1.

Proof. unfold has ASPs. unfold has ASPe. simpl. left. left. reflexivity. Qed.

Example ex4: knowsOfs Rely example_sys_1 Target.

Proof.

unfold knowsOfs. simpl. left. unfold knowsOfe. simpl. auto.

Ned

Example ex5: $knowsOfe Rely e_App Appraise \rightarrow False$.

Proof.

unfold knowsOfe. simpl. intros. destruct H.

Qed.

Example ex5': knowsOfs Rely $example_sys_1$ $Appraise <math>\rightarrow$ False.

Proof.

unfold knowsOfs. simpl. unfold knowsOfe. simpl. intros. $inverts\ H$. $inverts\ H0$. $inverts\ H0$. apply $inverts\ H0$.

Example ex6: knowsOfs Rely example_sys_1 Target.

Proof.

unfold knowsOfs,knowsOfe. simpl. auto.

Qed.

Example ex7: knowsOfs Rely [e_Rely] Target.

Proof.

unfold *knowsOfs*, *knowsOfe*. simpl. auto.

Qed.

Example ex8: dependsOne Appraise e_App Target.

Proof.

unfold dependsOne. simpl. auto.

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Qed.
Example ex9: dependsOns Appraise example_sys_1 Target.
Proof.
unfold depends Ons. simpl. unfold depends One. simpl. auto.
Qed.
Ltac prove_exec' :=
     simpl; auto; match goal with
                    \mid \vdash hasASPe \_\_\_ \Rightarrow cbv; left; reflexivity
                    \mid \vdash knowsOfe \_\_\_ \Rightarrow unfold knowsOfe; simpl; left; reflexivity
                    | ⊢ _ ∧ _ ⇒ split; prove_exec'
                    |\vdash ?A \Rightarrow \mathsf{idtac}\ A
                    end.
Example ex10: (executable (asp SIG) Target e_Targ).
Proof. prove_exec'. Qed.
Example ex11: (executable (asp CPY) Target e\_App) \rightarrow False.
Proof.
  intros Hcontra; cbv in *; destruct Hcontra.
Qed.
Example ex12: (executable (lseq (asp SIG) (asp SIG)) Target e_Targ).
Proof. prove_exec'. Qed.
Example ex13: (executables (lseq (asp aspc1)
                                 (att Target
                                      (lseq (asp SIG)
                                      (asp\ SIG)))
                   Rely\ example\_sys\_1).
Proof.
  prove_exec'; cbv; auto.
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Qed.