

Slides 6.8.23

# Math

- Preorder is – Reflexive & Transitive
- If anti-symmetric then partial order
- Homomorphisms bestow a preorder
- Ensure we understand up/down set...

**Definition 3 (Up-/down-sets).** *Given a preorder  $(\mathcal{P}, \leq)$ , a set  $\mathcal{S} \subseteq \mathcal{P}$  is an up-set (or order filter) iff for all structures  $G$  and  $H$ , whenever  $G \in \mathcal{S}$  and  $G \leq H$ , then  $H \in \mathcal{S}$ .  $\mathcal{S}$  is a down-set (or order ideal) iff for all structures  $G$  and  $H$ , whenever  $H \in \mathcal{S}$  and  $G \leq H$ , then  $G \in \mathcal{S}$ .*

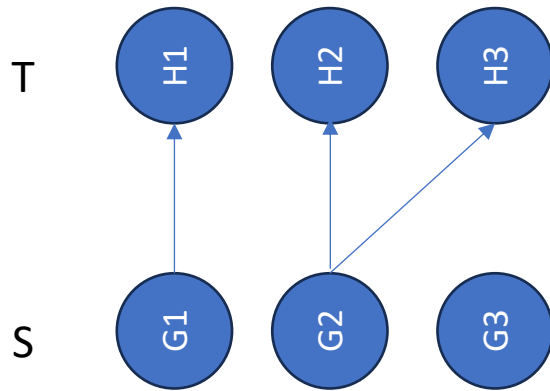
*The upward closure of a set  $\mathcal{S}$  is  $\phi(\mathcal{S}) = \{H \in \mathcal{P} \mid \exists G \in \mathcal{S} \wedge G \leq H\}$ . Similarly the downward closure of a set  $\mathcal{S}$  is  $\iota(\mathcal{S}) = \{G \in \mathcal{P} \mid \exists H \in \mathcal{S} \wedge G \leq H\}$ .*

The symbols  $\phi$  and  $\iota$  reflect the terminology of order filters and order ideals.

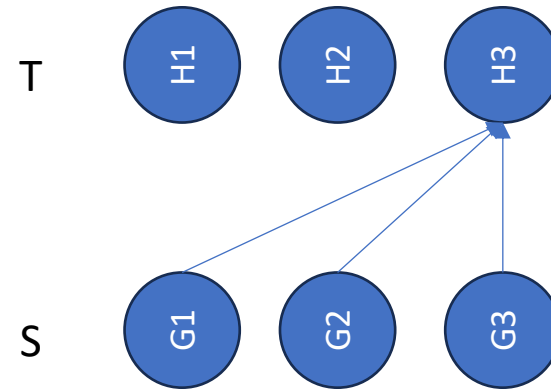
# Covers / supports

**Definition 10 (Supports, Covers).** Given two sets of graphs  $\mathcal{S}$  and  $\mathcal{T}$ , we say that  $\mathcal{S}$  supports  $\mathcal{T}$  iff for every  $H \in \mathcal{T}$ , there is some  $G \in \mathcal{S}$ , such that  $G \leq H$ . We say that  $\mathcal{T}$  covers  $\mathcal{S}$  iff for every  $G \in \mathcal{S}$  there is some  $H \in \mathcal{T}$  such that  $G \leq H$ .

$\mathcal{S}$  supports  $\mathcal{T}$



$\mathcal{T}$  covers  $\mathcal{S}$



# Copland preorder

$$c_1 \leq_{\mathcal{I}}^C c_2 \text{ iff } \{\llbracket c_2 \rrbracket_p^e\} \text{ covers } \{\llbracket c_1 \rrbracket_p^e\}$$

$$c_1 \leq_{\mathcal{F}}^C c_2 \text{ iff } \{\llbracket c_2 \rrbracket_p^e\} \text{ supports } \{\llbracket c_1 \rrbracket_p^e\}$$

$$c_1 \leq_{\mathcal{I}}^C c_2 \text{ iff } \llbracket c_1 \rrbracket_p^e \leq \llbracket c_2 \rrbracket_p^e$$

$$c_1 \leq_{\mathcal{F}}^C c_2 \text{ iff } \llbracket c_2 \rrbracket_p^e \leq \llbracket c_1 \rrbracket_p^e$$

6.2 slides

# Input to the problem

- Set of measurement operations
- Ranking of corruption events
- Dependency structure?

# Principles

1. Increasing volume of measurement operations may not confine the adversary
2. Stronger measurements mimic the dependency chain
  - If  $a$  depends on  $b$  which depends on  $c$  (all at some place  $p_1$ ) then  $ms(p_1, rtm, p_1, c) \rightarrow ms(p_1, c, p_1, b) \rightarrow ms(p_1, b, p_1, a)$  will be the strongest measurement
  - Same idea as *well-supported* measurement from confining paper
  - Hypothesis: if measurement chain is not well-supported, then it is “easy” for an adversary to corrupt

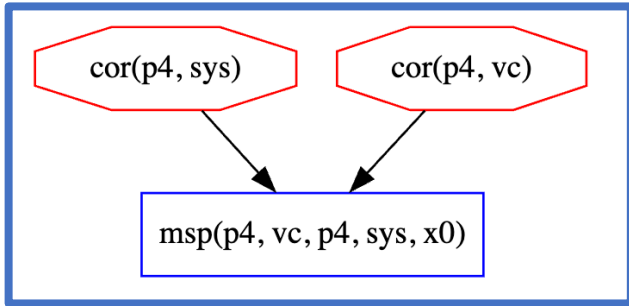
**Definition 6.** A measurement event  $e = ms(o_2, o_1)$  in execution  $E$  is well-supported iff either

- $o_2 = rtm$ , or
- for every  $o \in D^1(o_1)$ , there is a measurement event  $e' \prec_E e$  such that  $o$  is the target of  $e'$ .

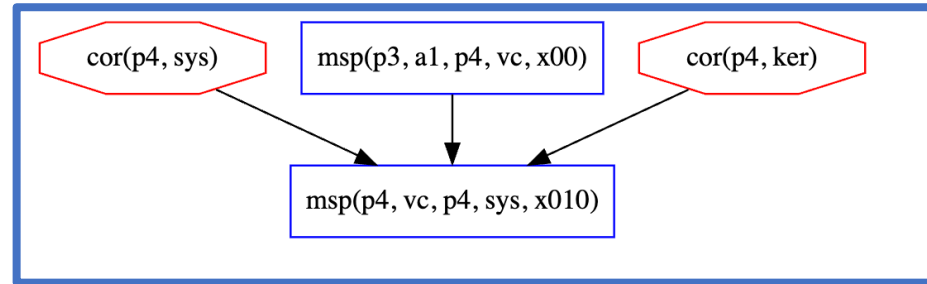
When  $e$  is well-supported, we call the set of  $e'$  from Condition ii above the support of  $e$ . An execution  $E$  measures bottom-up iff each measurement event  $e \in E$  is well-supported.

a
b
c
rtm

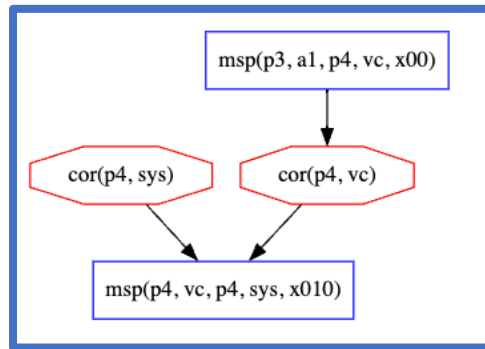
# Try different ordering



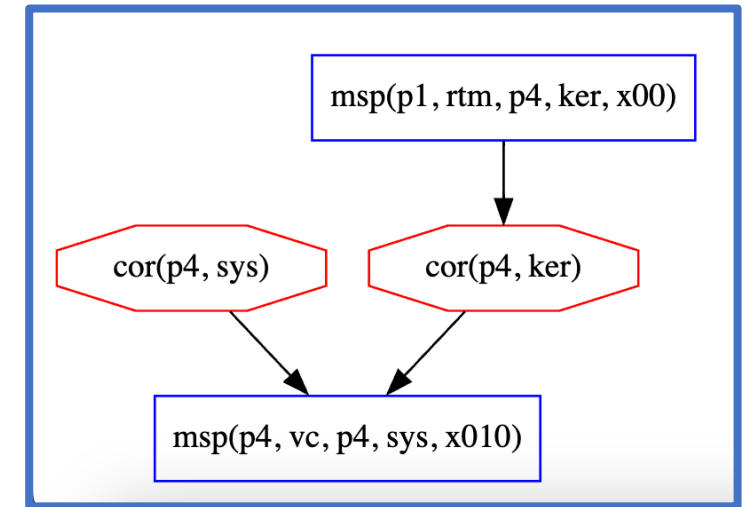
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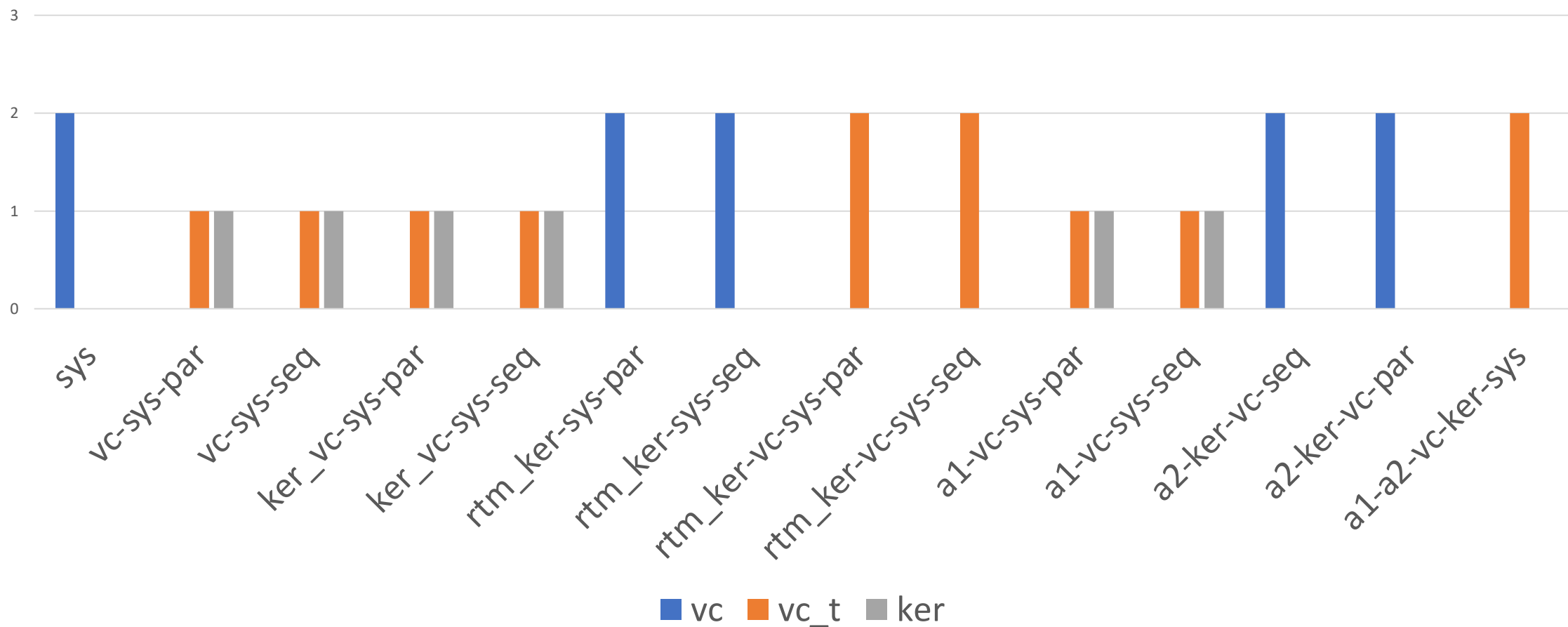
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Protocol Name	Protocol
sys	*target: @p4 [vc p4 sys]
rtm_ker-sys-par	*target: @p1 [rtm p4 ker +~+ @p4 vc p4 sys]
rtm_ker-sys-seq	*target: @p1 [rtm p4 ker +<+ @p4 vc p4 sys]
ker_vc-sys-par	*target: @p4 [(ker p4 vc) +~+ @p4 (vc p4 sys)]
ker_vc-sys-seq	*target: @p4 [(ker p4 vc) +<+ @p4 (vc p4 sys)]
rtm_ker-vc-sys-par	*target: @p1 [(rtm p4 ker) +~+ @p4 [(ker p4 vc) +~+ @p4 (vc p4 sys)]]
rtm_ker-vc-sys-seq	*target: @p1 [(rtm p4 ker) +<+ @p4 [(ker p4 vc) +<+ @p4 (vc p4 sys)]]
vc-sys-par	*target: @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]
vc-sys-seq	*target: @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]
a1-vc-sys-seq	*target: @p1 [(rtm p3 a1) +<+ @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]]
a1-vc-sys-par	*target: @p1 [(rtm p3 a1) +~+ @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]]
a2-ker-vc-seq	*target: @p1 [rtm p3 a2 +<+ @p3 [a2 p4 ker +<+ @p4 (vc p4 sys)]]
a2-ker-vc-par	*target: @p1 [(rtm p3 a2) +~+ @p3 [(a2 p4 ker) +~+ @p4 (vc p4 sys)]]
a1-a2-vc-ker-sys	*target: @p1 [(rtm p3 a1 +~+ rtm p3 a2) +<+ @p3 [(a1 p4 vc +~+ a2 p4 ker) +<+ @p4 (vc p4 sys1)]]

Measurement Operations with Minimum Cost Attack



5.26 slides

# Add protocols

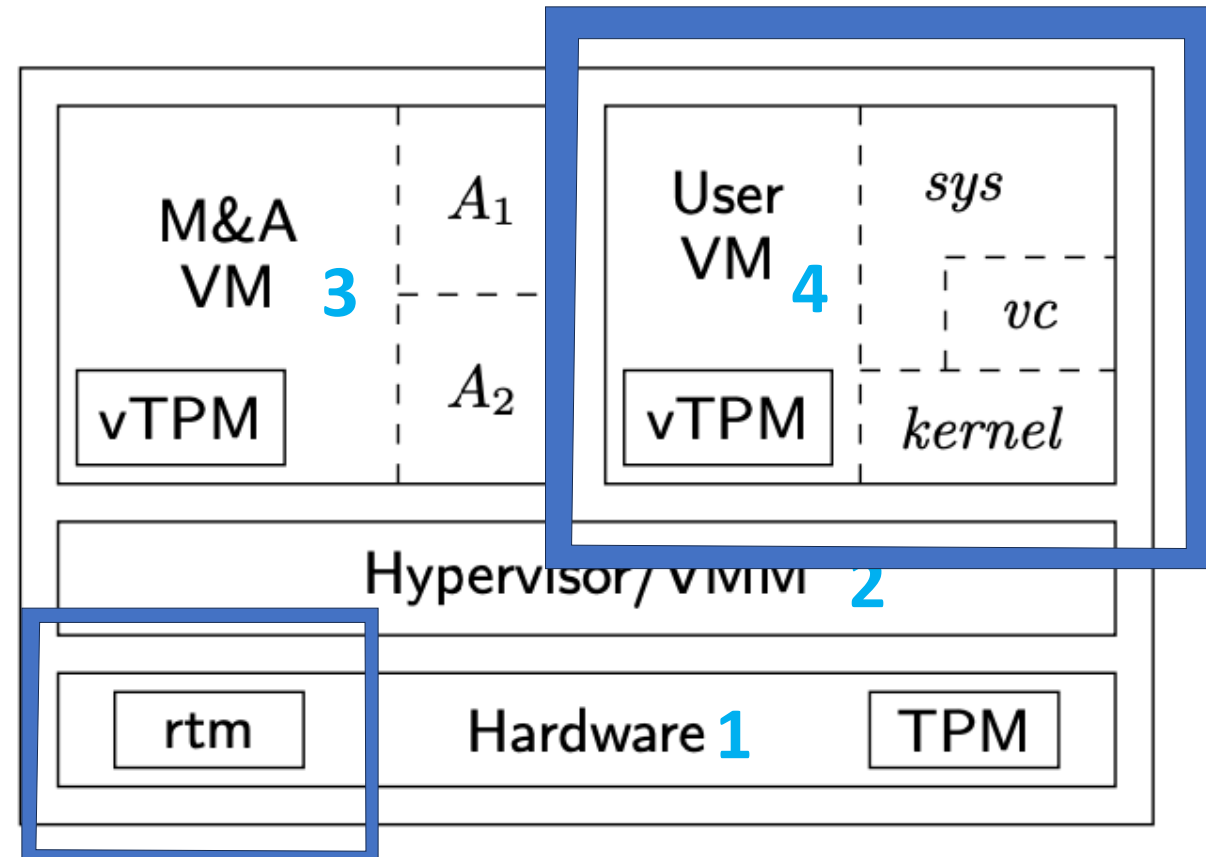
- \*target: @p4 [ker p4 vc] +~+ @p4 [vc p4 sys]
- \*target: @p4 [ker p4 vc] +<+ @p4 [vc p4 sys]
  
- \*target: @p1 [(rtm p4 ker) +~+ @p4 [(ker p4 vc) +~+ @p4 (vc p4 sys)]]
- \*target: @p1 [(rtm p4 ker) +<+ @p4 [(ker p4 vc) +<+ @p4 (vc p4 sys)]]

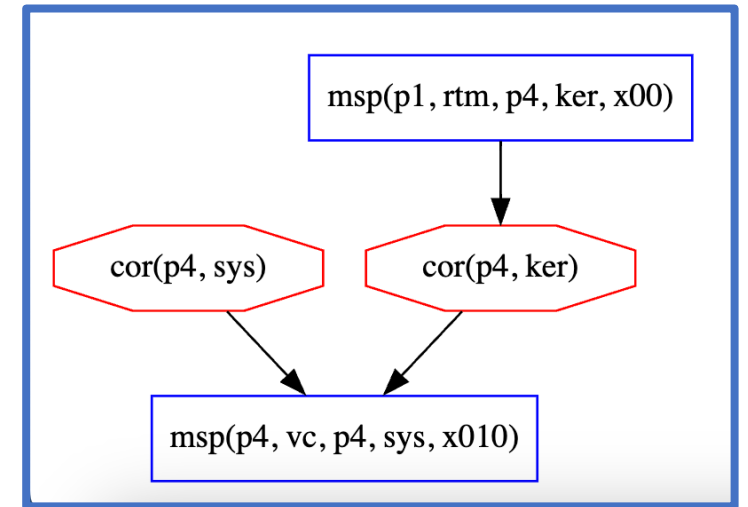
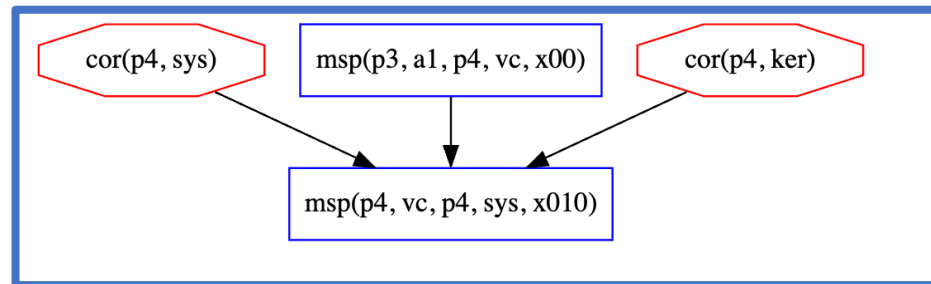
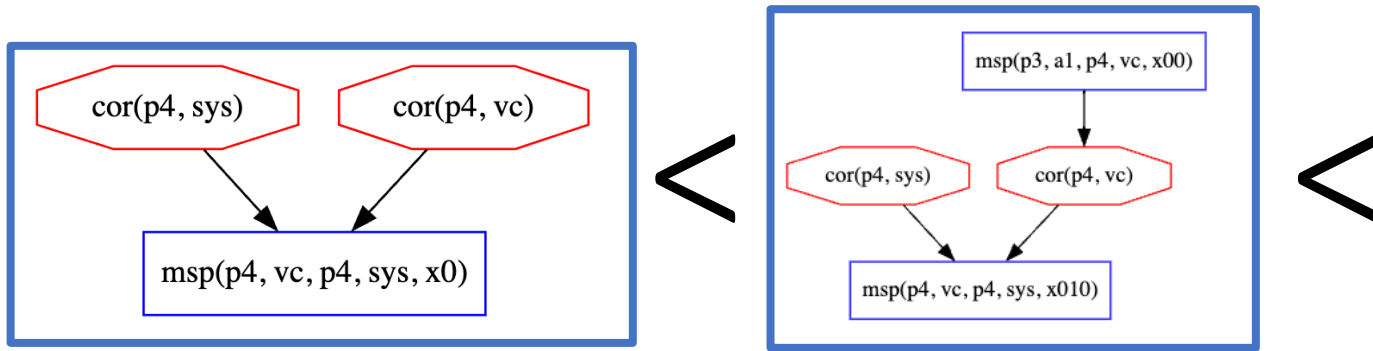
# Protocols

Protocol Name	Protocol
sys	*target: @p4 [vc p4 sys]
rtm_ker-sys-par	*target: @p1 [rtm p4 ker +~+ @p4 vc p4 sys]
rtm_ker-sys-seq	*target: @p1 [rtm p4 ker +<+ @p4 vc p4 sys]
ker_vc-sys-par	*target: @p4 [(ker p4 vc) +~+ @p4 (vc p4 sys)]
ker_vc-sys-seq	*target: @p4 [(ker p4 vc) +<+ @p4 (vc p4 sys)]
rtm_ker-vc-sys-par	*target: @p1 [(rtm p4 ker) +~+ @p4 [(ker p4 vc) +~+ @p4 (vc p4 sys)]]
rtm_ker-vc-sys-seq	*target: @p1 [(rtm p4 ker) +<+ @p4 [(ker p4 vc) +<+ @p4 (vc p4 sys)]]
vc-sys-par	*target: @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]
vc-sys-seq	*target: @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]
a1-vc-sys-seq	*target: @p1 [(rtm p3 a1) +<+ @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]]
a1-vc-sys-par	*target: @p1 [(rtm p3 a1) +~+ @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]]
a2-ker-vc-seq	*target: @p1 [rtm p3 a2 +<+ @p3 [a2 p4 ker +<+ @p4 (vc p4 sys)]]
a2-ker-vc-par	*target: @p1 [(rtm p3 a2) +~+ @p3 [(a2 p4 ker) +~+ @p4 (vc p4 sys)]]
a1-a2-vc-ker-sys	*target: @p1 [(rtm p3 a1 +~+ rtm p3 a2) +<+ @p3 [(a1 p4 vc +~+ a2 p4 ker) +<+ @p4 (vc p4 sys1)]]

# Say we have the architecture from “Confining the Adversary” Paper

- $ms(rtm, A1)$
- $ms(rtm, A2)$
- $ms(A1, vc)$
- $ms(A2, ker)$
- $msker(vc, sys)$





# Corruption event order

**$p4 < p1$**

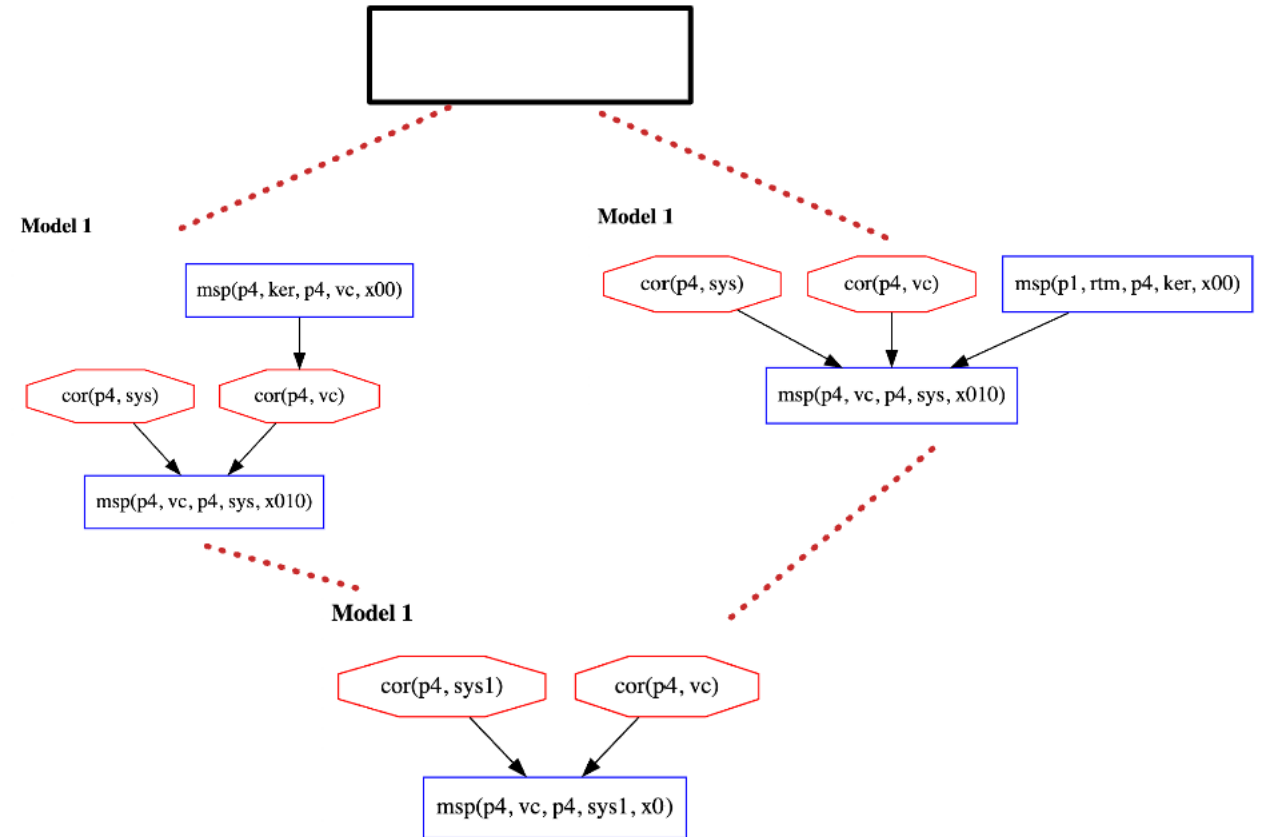
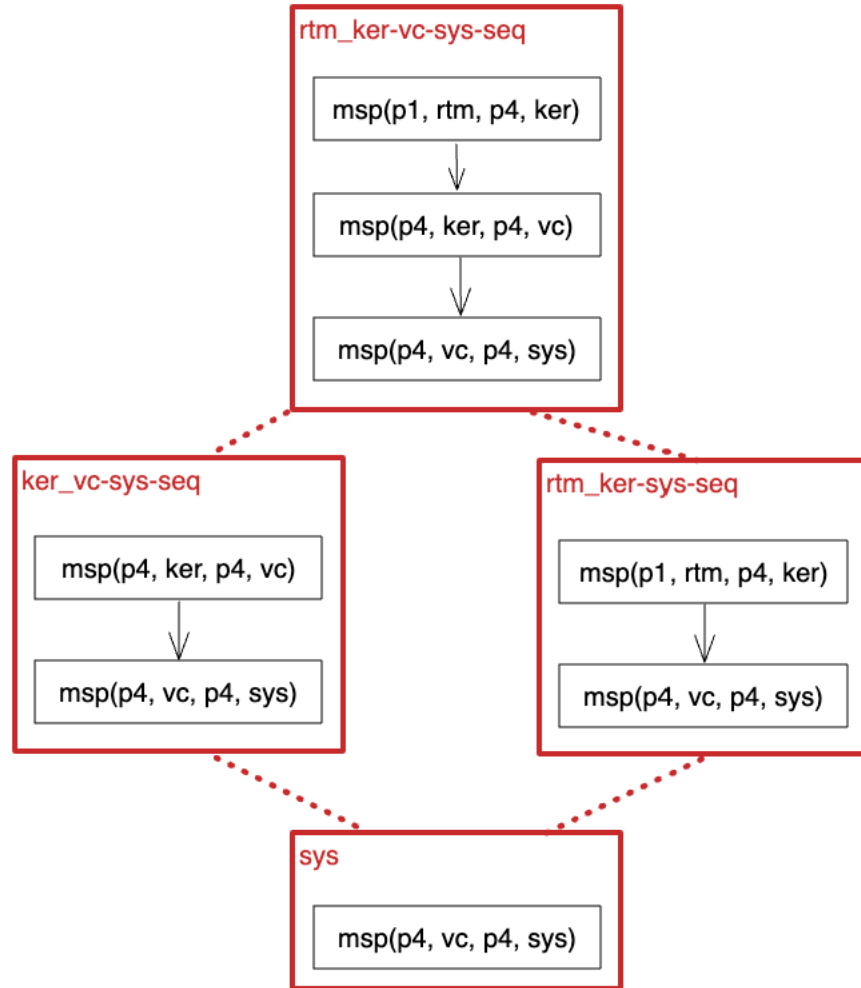
- To corrupt something at  $p4$  is lowest cost to adversary
- Corrupting at  $p1$  (rtm location) is the highest cost

**$sys < vc < vc\_t < ker < ker\_t$**

- Corrupting the system is easiest for the adversary
- Denote  $x\_t$  as a corruption event of  $x$  in a timely manner



- Possible lattice of protocols



- Ordering relation = subset relation

Slides (5/16)

# Goals of cost analysis

- \* Maximize the adversary's minimum cost
- Cost considerations
  - Weight of cost depends on component
  - Cor-rep-cor should be more expensive than just corrupt
  - Cost of repair
    - Maybe repair is not costly to an adversary
    - Maybe need to model partial repair (need to decompose system further for this)

# What I did for today (5/16)

1. Ran all protocols
2. Selected lowest cost/costs model
3. Took maximum of lowest costs to produce best (?) protocol

# Principles

- Increase cost after start event
- Any corruption of a deeper component is higher cost
- Add weight to cost event
  - Have some base cost to the corruption event
  - Add to cost more if its in a protected place

# Protocols

Protocol Name	Protocol
sys	*target: @p4 [vc p4 sys1]
vc-sys-par	*target: @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]
vc-sys-seq	*target: @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]
a1-vc-sys-seq	*target: @p1 [(rtm p3 a1) +<+ @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]]
a1-vc-sys-par	*target: @p1 [(rtm p3 a1) +~+ @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]]
a2-ker-vc-seq	*target: @p1 [rtm p3 a2 +<+ @p3 [a2 p4 ker +<+ @p4 (vc p4 sys)]]
a2-ker-vc-par	*target: @p1 [(rtm p3 a2) +~+ @p3 [(a2 p4 ker) +~+ @p4 (vc p4 sys)]]
a1-a2-vc-ker-sys	*target: @p1 [(rtm p3 a1 +~+ rtm p3 a2) +<+ @p3 [(a1 p4 vc +~+ a2 p4 ker) +<+ @p4 (vc p4 sys1)]]

# sys

Protocol: \*target: @p4 [vc p4 sys1]

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at our main measurement
% event. This is a measurement of sys
l(V) = msp(p4, M, p4, sys1, X)
=> corrupt_at(p4, sys1, V).

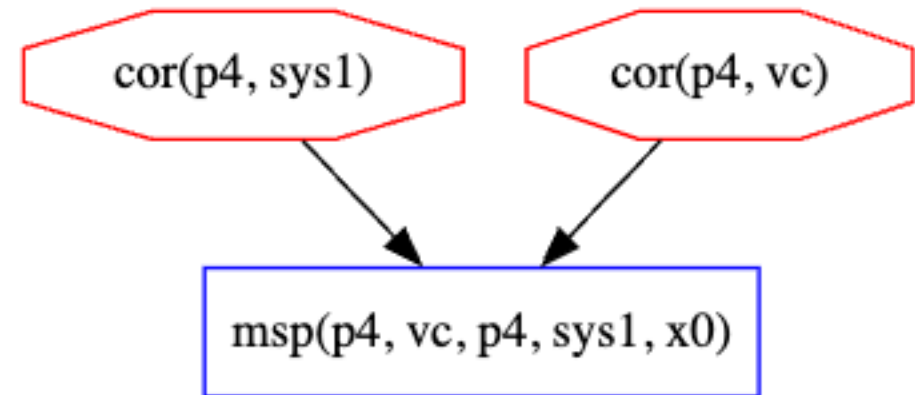
% Assume no dependencies

% No recent assumptions
% prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
% => false.

% No deep assumptions
l(V) = cor(p1, M) => false.

m4_include(`sys.gli')m4_dnl
m4_include(`sys_dist.gli')m4_dnl
m4_include(`thy.gli')m4_dnl
```

## Model 1



# vc-sys-par

- \*target: @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]

---

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at
% our main measurement event.
% This is a measurement of sys.
l(V) = msp(p4, M, p4, sys, X)
=> corrupt_at(p4, sys, V).

% Assume dependencies
depends(p4, C, p4, sys) => C = ker.
depends(p4, C, p4, vc) => C = ker.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.

% Assume no recent corruptions
%prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
%=> false.

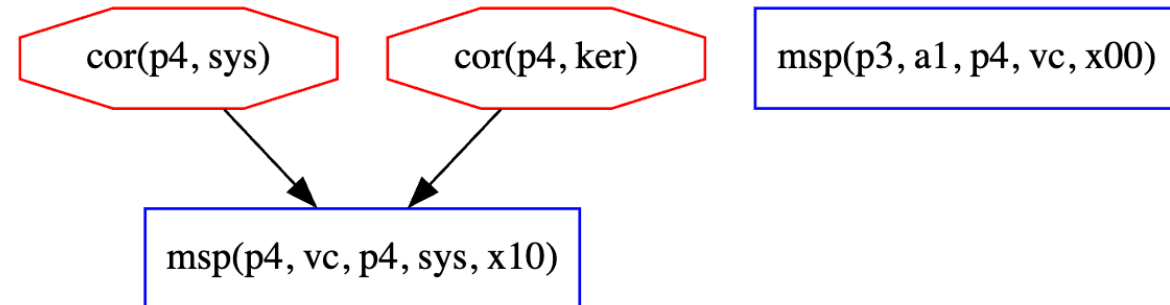
% Assume no deep corruptions
l(V) = cor(p1, M) => false.

m4_include(`vc-sys-par.gli')m4_dnl

m4_include(`vc-sys-par_dist.gli')m4_dnl

m4_include(`thy.gli')m4_dnl
```

## Model 2

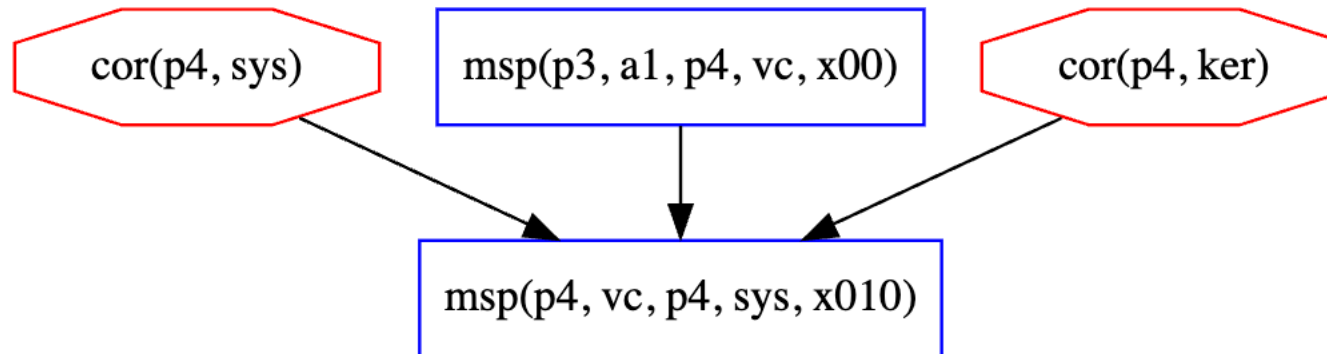




# vc-sys-seq

- \*target: @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]

## Model 2



# a1-vc-sys-par

Protocol:

\*target: @p1 [(rtm p3 a1) +~+ @p3 [(a1 p4 vc) +~+ @p4 (vc p4 sys)]]

```
% Assume adversary avoids detection at  
% our main measurement event.  
% This is a measurement of sys.  
l(V) = msp(p4, M, p4, sys, X)  
| => corrupt_at(p4, sys, V).
```

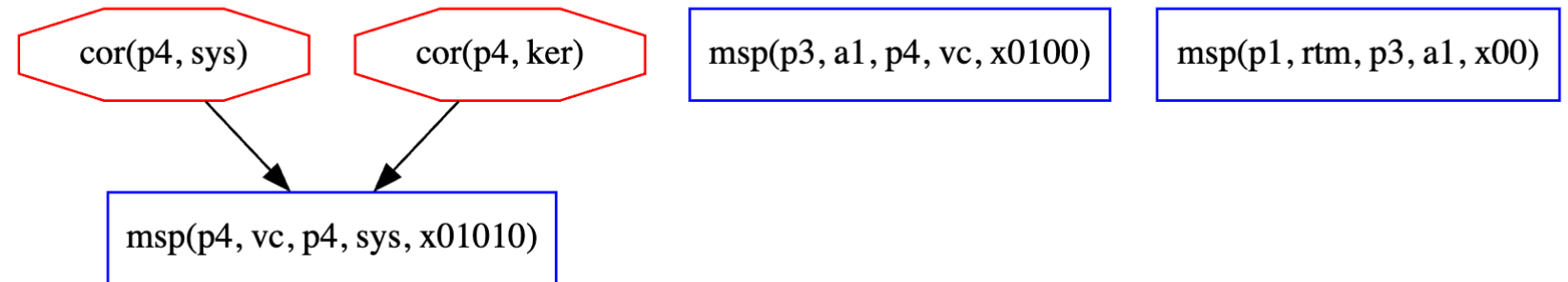
```
% system dependencies  
depends(p4, C, p4, sys) => C = ker.  
depends(p4, C, p4, vc) => C = ker.
```

```
% rtm has no dependencies  
depends(p1, C, p1, rtm) => false.
```

```
% Assume no recent corruptions  
% prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)  
% => false.
```

```
% Assume no deep corruptions  
l(V) = cor(p1, M) => false.
```

## Model 2

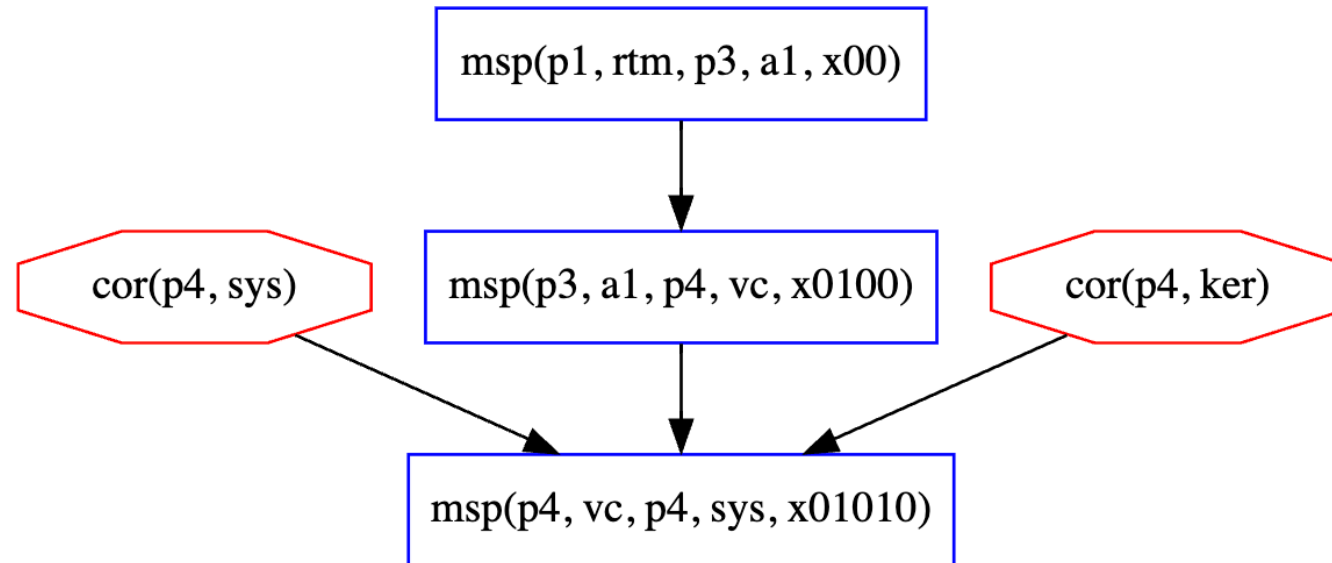


# A1-vc-sys-seq

Protocol:

\*target: @p1 [(rtm p3 a1) +<+ @p3 [(a1 p4 vc) +<+ @p4 (vc p4 sys)]]

## Model 2



# a2-ker-vc-par

- \*target: @p1 [(rtm p3 a2) +~+ @p3 [(a2 p4 ker) +~+ @p4 (ker p4 vc)]]

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at
% our main measurement event.
% This is a measurement of ker.
l(V) = msp(p4, M, p4, vc, X)
=> corrupt_at(p4, vc, V).

% dependencies
depends(p4, C, p4, sys) => C = ker.
depends(p4, C, p4, vc) => C = ker.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.

% Assume no recent corruptions
%prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
% => false.

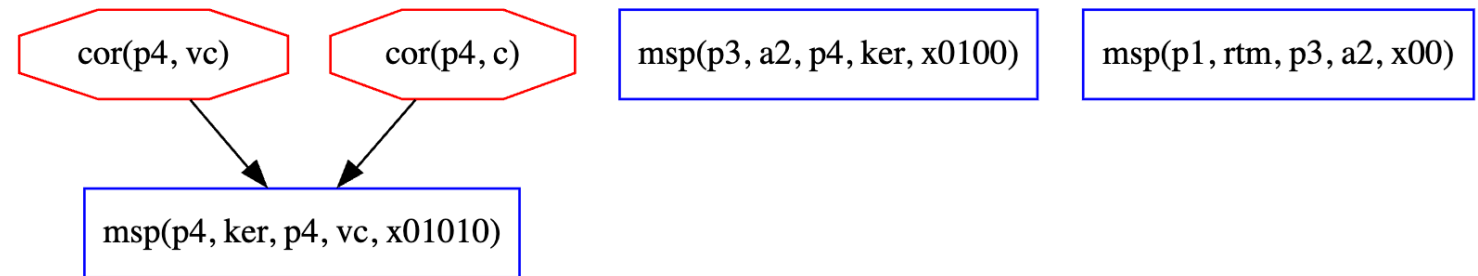
% Assume no deep corruptions
l(V) = cor(p1, M) => false.

m4_include(`a2-ker-vc-par.gli')m4_dnl

m4_include(`a2-ker-vc-par_dist.gli')m4_dnl

m4_include(`thy.gli')m4_dnl
```

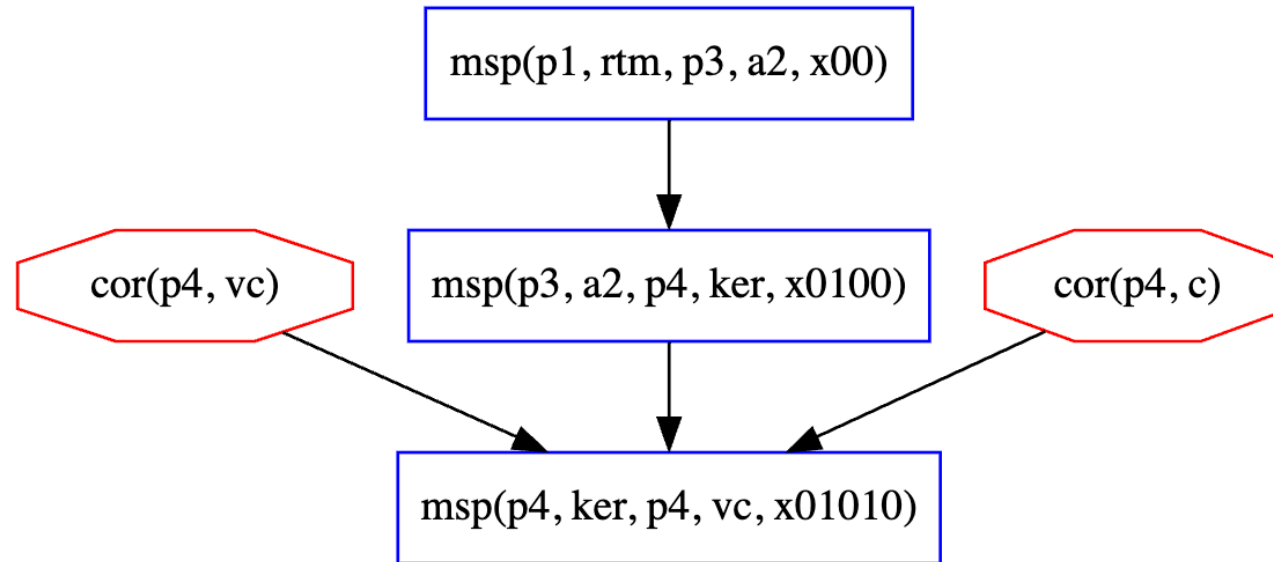
## Model 1



# a2-ker-vc-seq

- \*target: @p1 [rtm p3 a2 +<+ @p3 [a2 p4 ker +<+ @p4 (ker p4 vc)]]

## Model 1



# a1-a2-vc-ker-sys

- \*target: @p1 [(rtm p3 a1 +~+ rtm p3 a2) +<+ @p3 [(a1 p4 vc +~+ a2 p4 ker) +<+ @p4 (vc p4 sys1)]]

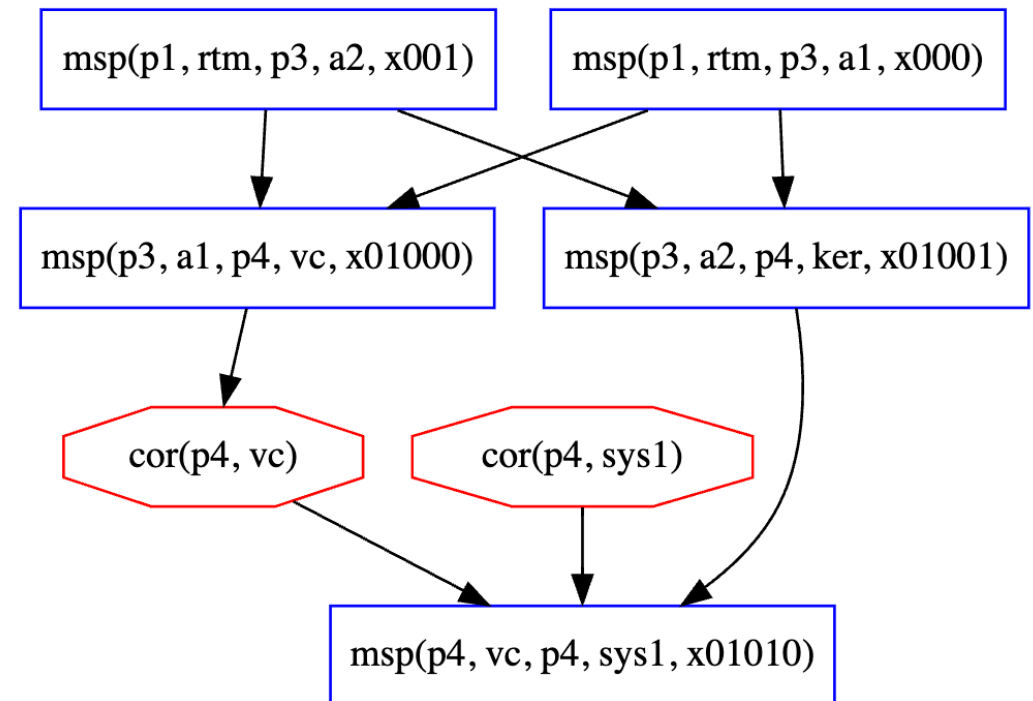
```
% dependencies
depends(p4, C, p4, sys1) => C = ker.
depends(p4, C, p4, vc) => C = ker.
% depends(p1, C, p4, ker) => C = rtm.
% depends(p1, C, p3, a1) => C = rtm.

% depends(P1, C, p3, a2) => P1 = p1 & C = rtm.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.

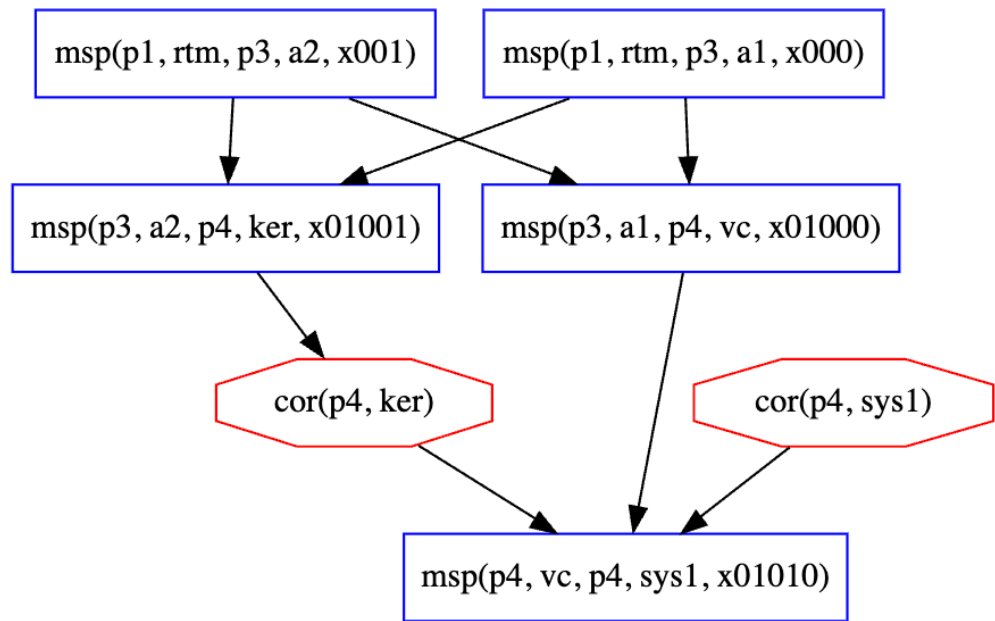
% Assume no recent corruptions
%prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
% => false.

% Assume no deep corruptions
l(V) = cor(p1, M) => false.
```

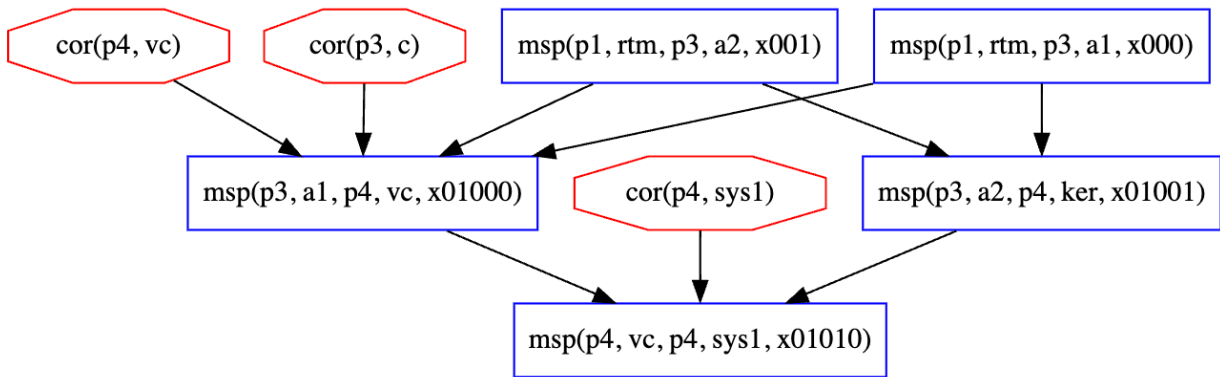
**Model 1**



**Model 2**



**Model 4**



# Same protocol... change theory

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at
% our main measurement event.
% This is a measurement of ker.
l(V) = msp(p4, M, p4, sys1, X)
=> corrupt_at(p4, sys1, V).

% dependencies
depends(p4, C, p4, sys1) => C = ker.
depends(p4, C, p4, vc) => C = ker.
depends(P1, C, p4, ker) => P1 = p1 & C = rtm.
depends(P1, C, p3, a1) => P1 = p1 & C = rtm.
depends(P1, C, p3, a2) => P1 = p1 & C = rtm.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.

% Assume no recent corruptions
%prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
% => false.

% Assume no deep corruptions
l(V) = cor(p1, M) => false.

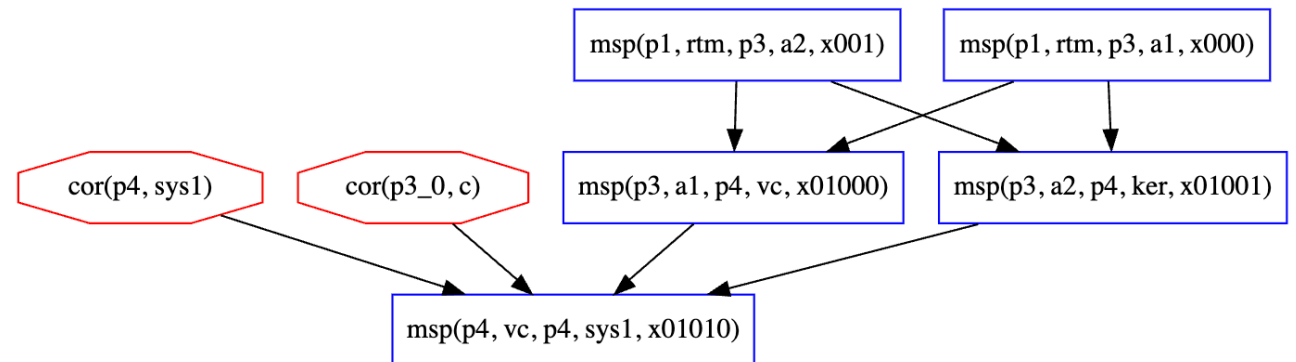
m4_include(`a1-a2-vc-ker-sys.gli')m4_dnl
m4_include(`a1-a2-vc-ker-sys_dist.gli')m4_dnl
m4_include(`thy.gli')m4_dnl
```

## % Rule 1

```
%l(V) = msp(P2, M, P1, T, X) & corrupt_at(P1, T, V)
% => corrupt_at(P2, M, V) | depends(P2, C, P2, M) & corrupt_at(P2, C, V).
%corruption events can be at different places
l(V) = msp(P2, M, P1, T, X) & corrupt_at(P1, T, V)
| => corrupt_at(P2, M, V) | depends(P3, C, P2, M) & corrupt_at(P3, C, V).
```

## Models

### Model 1

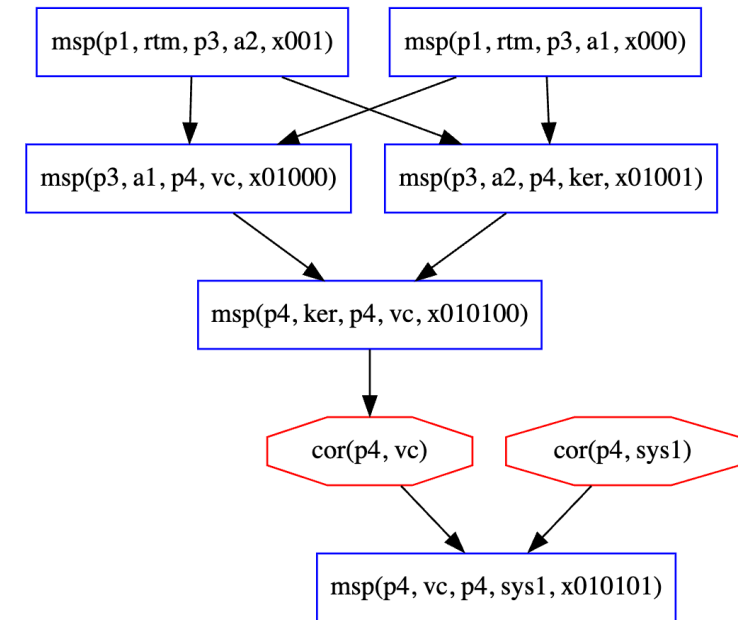




# Back to original theory... change measurement

- \*target: @p1 [( rtm p3 a1 +~+ rtm p3 a2) +<+  
@p3 [( a1 p4 vc +~+ a2 p4 ker ) +<+  
@p4 [(ker p4 vc) +<+ (vc p4 sys1)]]]

Model 1



# Ways to manipulate CHASE outcome

Things I tried but didn't have much impact...

1. Change measurement combination (parallel to sequence)
2. Add measurements
  - Not sure how to distinguish when something is a “useful” measurement
3. Changing dependencies

Slides (5/5)

# Assumptions

- Always assume deep theorem (remove recent theorem)

- $I(V) = \text{cor}(p1, M) \Rightarrow \text{false}.$

- Assumptions about system dependencies

- TPM is the root of trust... has no dependencies

- Virus checker depends on kernel (p4,ker)

- System depends on kernel (p4,ker)

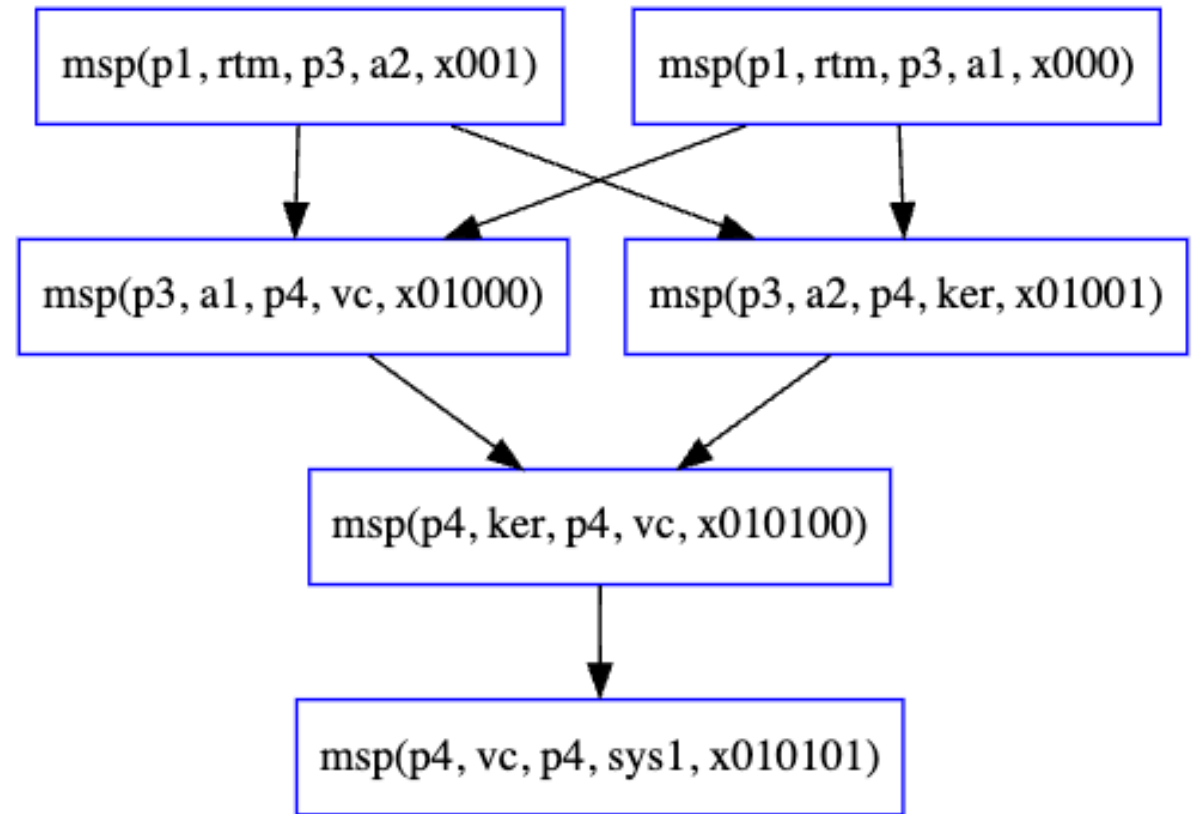
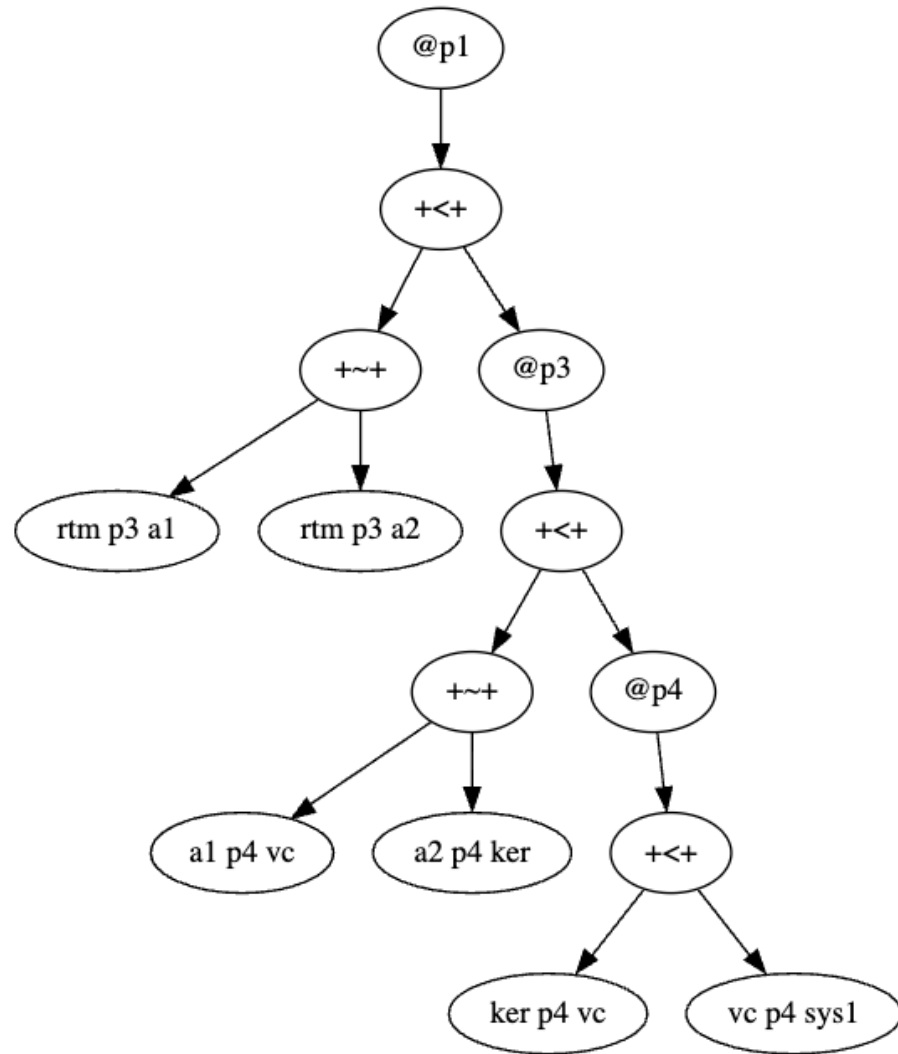
- Kernel depends on the hardware (p1,rtm)

- A1 depends on the hardware (p1,rtm)

- A2 depends on the hardware (p1,rtm)

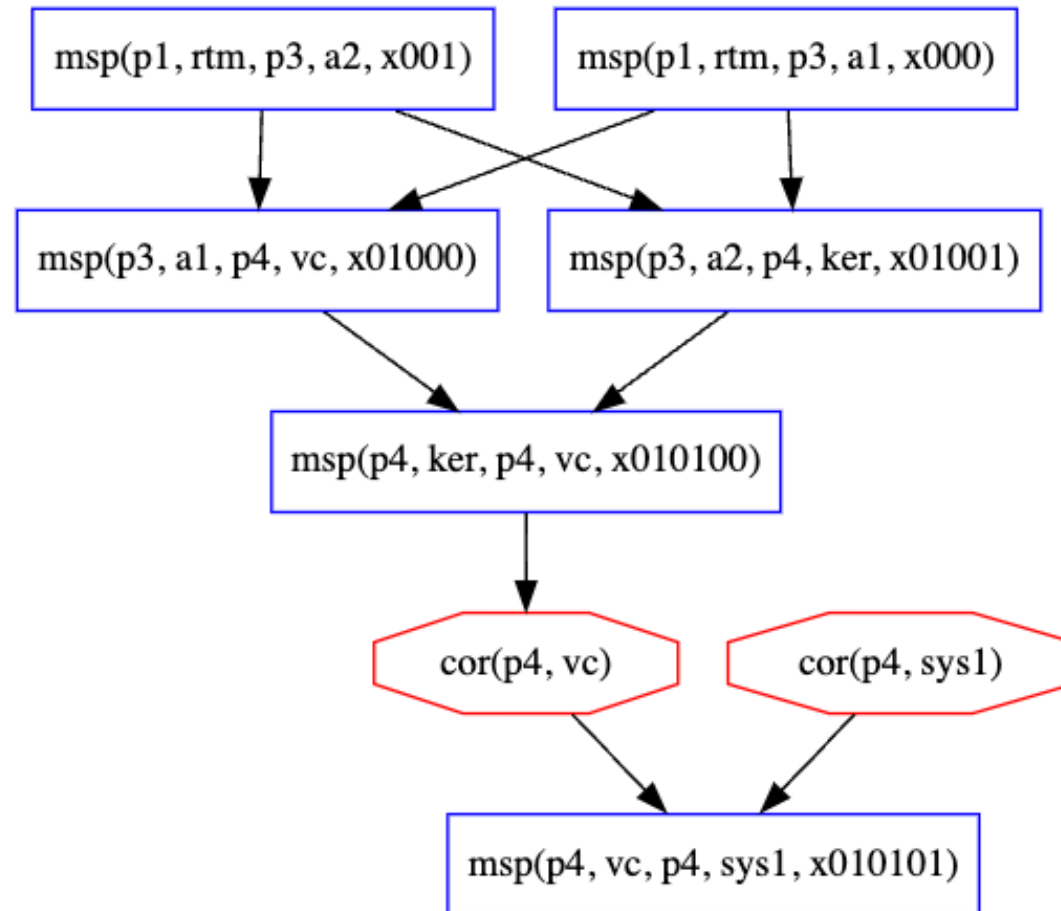
```
% dependencies
depends(p4, C, p4, sys1) => C = ker.
depends(p4, C, p4, vc) => C = ker.
depends(p1, C, p4, ker) => C = rtm.
depends(p1, C, p3, a1) => C = rtm.
depends(p1, C, p3, a2) => C = rtm.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.
```

## Abstract Syntax Tree

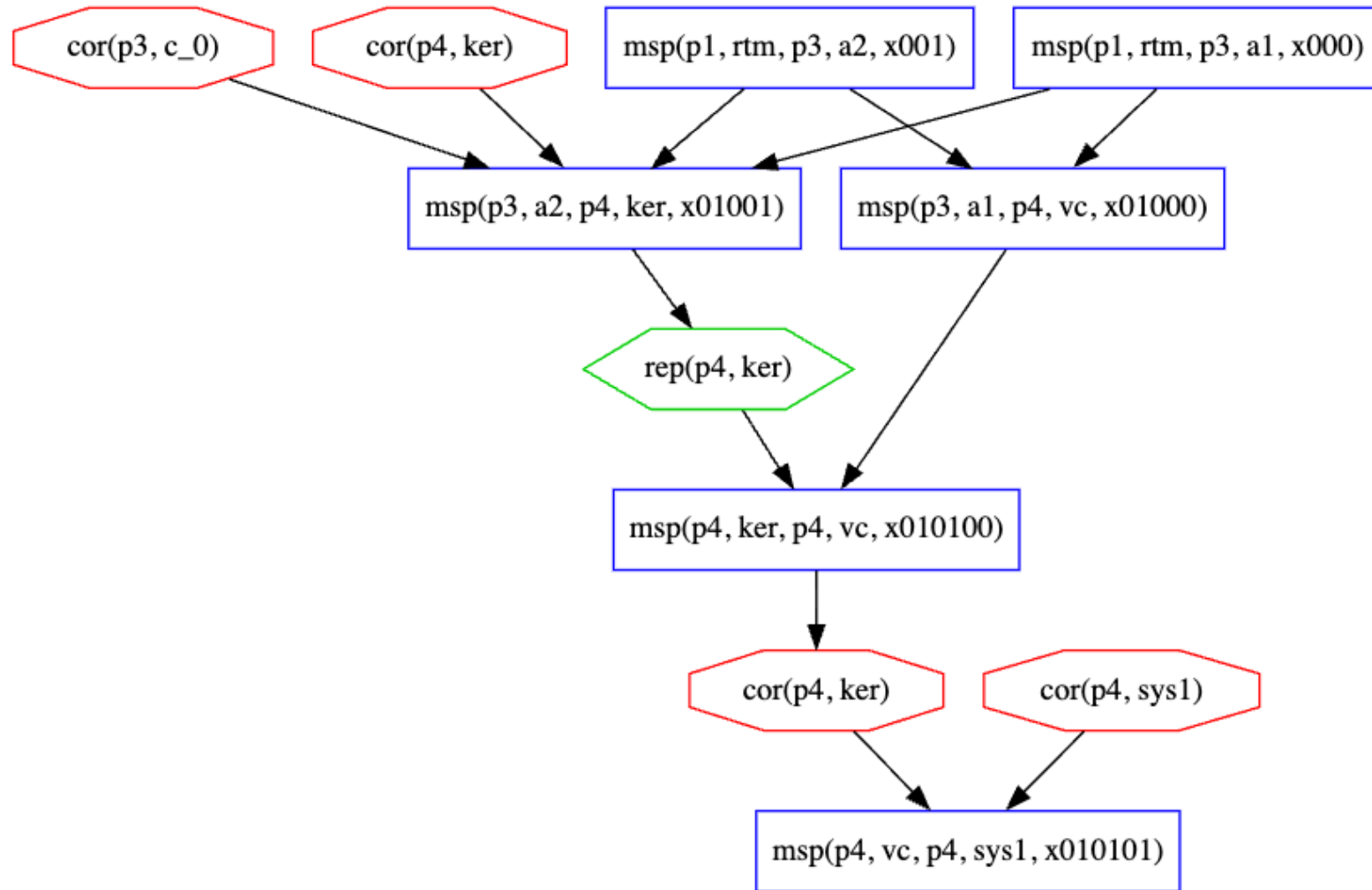


Assuming recent measurements may be corrupted there are 21 models...

**Model 1**



### Model 21



Order	Event	Cost	Present In	Details
low	cor(p4,sys1)	c1	all models	Always before the last measurement event
	cor(p4,c)	c2	4	happens before ms(ker,vc)
	cor(p4,vc)	c3	1 4 5 8 9 10 11 12 13 14 15 16 17 18(2) 19(2)	occurs after some attestation start event (between measurements) or before a measurement, sometimes happens twice (once and then after a repair)
	cor(p4, ker)	c4	2 3 5 6 7 10 11 12 13 14 15 16 17 20(2) 21(2)	occurs various places.. Before/after ms(a2,ker), before ms(vc,sys1)
	cor(p4,c_1)	c5	8	before ms(ker,vc)
	cor(p3,a1)	c6	8 10 14 15 18	before ms(a1,vc), always after the attestation begins... maybe this is most difficult because you have to consider time window for adversary
	cor(p3,c_3)	c7	9 11 16 17 19	before ms(a1,vc), no attestation start event... could be easiest for an adversary
	cor(p4,c_2)	c8	9	before ms(ker,vc), no attestation start event... could be easiest for adversary
	cor(p3,a2)	c9	6 12 14 16 20	between ms(rtm, a2) – ms(a2, ker)... close to root of trust. Difficult for an adversary
	cor(p3,c_4)	c10	13	before ms(a2,ker), no attestation start event... could be easiest for adversary
	cor(p3,c_5)	c11	15	before ms(a2,ker) no attestation start event... could be easiest for adversary
	cor(p3,c_6)	c12	17	before ms(a2,ker) , no attestation start event... could be easiest for adversary
	rep(p4,vc)	c13	18 19	between ms(a1,vc) – ms(ker,vc),
	rep(p4,ker)	c14	20 21	between ms(a2,ker) -- ms(ker, vc)
	cor(p3,c_0)	c15	7 21	before ms(a2,ker)



# Thoughts/Takeaways

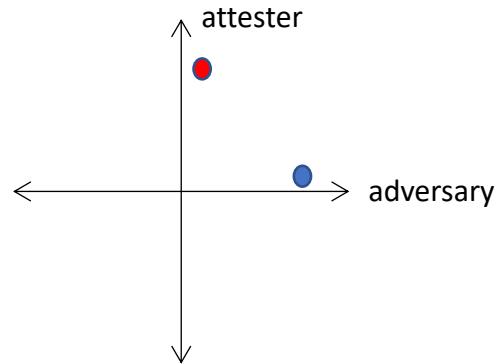
- Write some script to assign cost

Slides (3) 4.21.23

# Goals of cost analysis

Ultimate goal: guide selection of a protocol

- How:
  - Systematic variation of assumption
  - Assign abstract cost to each component that's corrupted
  - Define function to create order between cost and value
- Consider:
  - Cost to adversary
  - Cost to attester

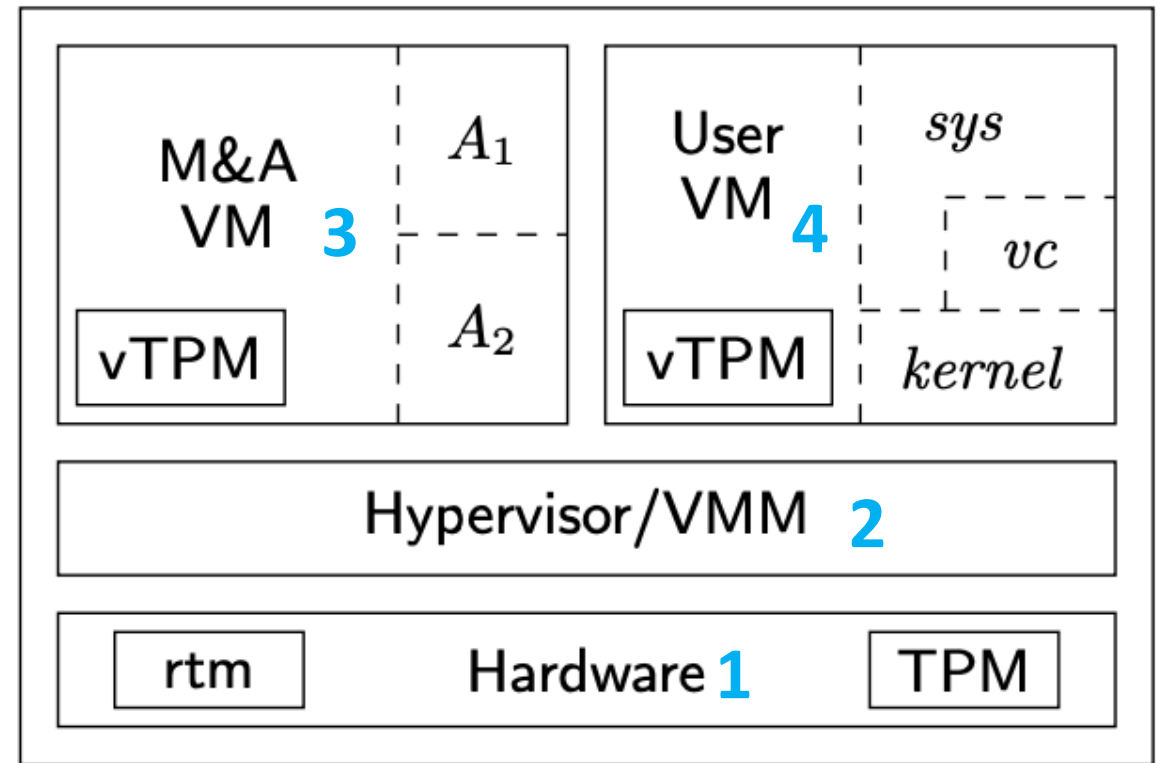


# What I did for today (5/5)

- Ran one complex protocol which considered all measurement operations in “Confining” system
- What I did not do...
  - Consider different ordering of the protocol
  - I think we have enough material to discuss with just one protocol

# Say we have the architecture from “Confining the Adversary” Paper

- $ms(rtm, A1)$
- $ms(rtm, A2)$
- $ms(A1, vc)$
- $ms(A2, ker)$
- $msker(vc, sys)$



# Assumptions

- Always assume deep theorem (remove recent theorem)
- Assumptions about system dependencies
  - TPM is the root of trust... has no dependencies
  - Virus checker depends on kernel (p4,ker)
  - System depends on kernel (p4,ker)
  - Kernel depends on the hardware (p1,rtm)
  - A1 depends on the hardware (p1,rtm)
  - A2 depends on the hardware (p1,rtm)

```
% dependencies
depends(p4, C, p4, sys1) => C = ker.
depends(p4, C, p4, vc) => C = ker.
depends(p1, C, p4, ker) => C = rtm.
depends(p1, C, p3, a1) => C = rtm.
depends(p1, C, p3, a2) => C = rtm.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.
```

# Principles

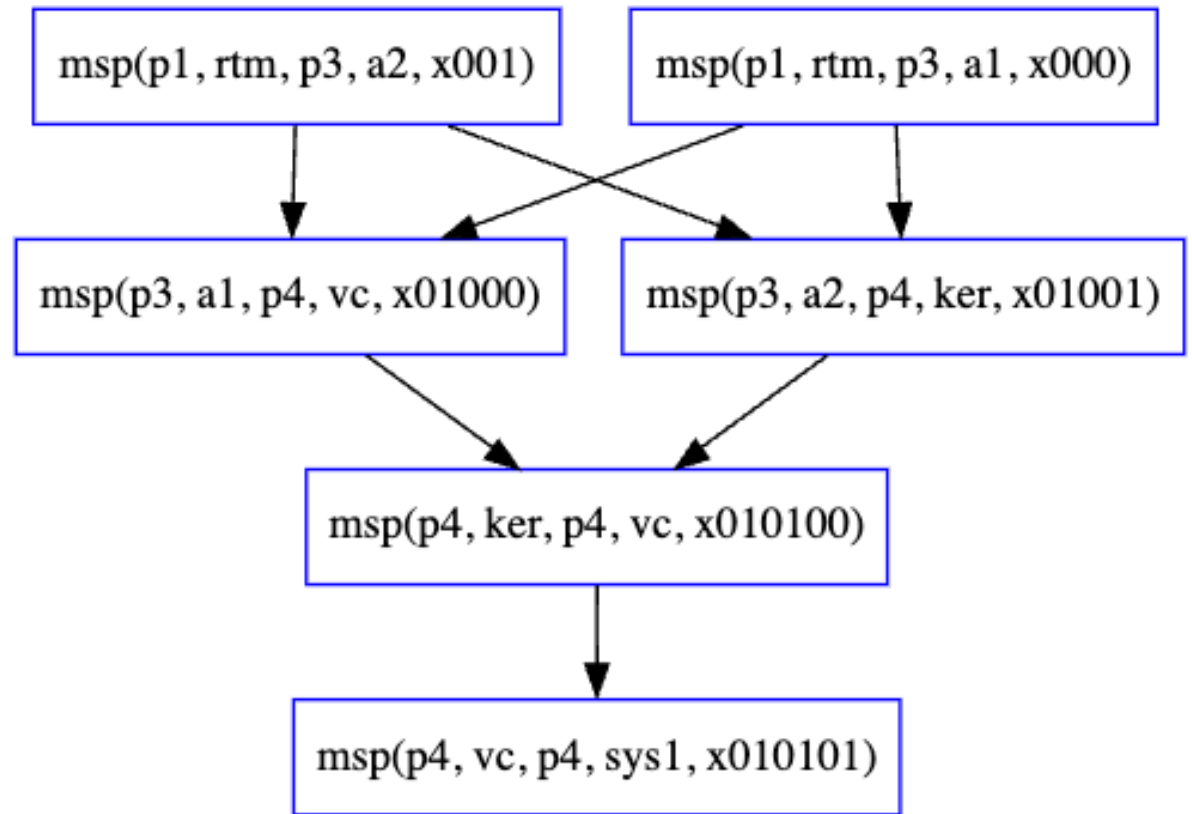
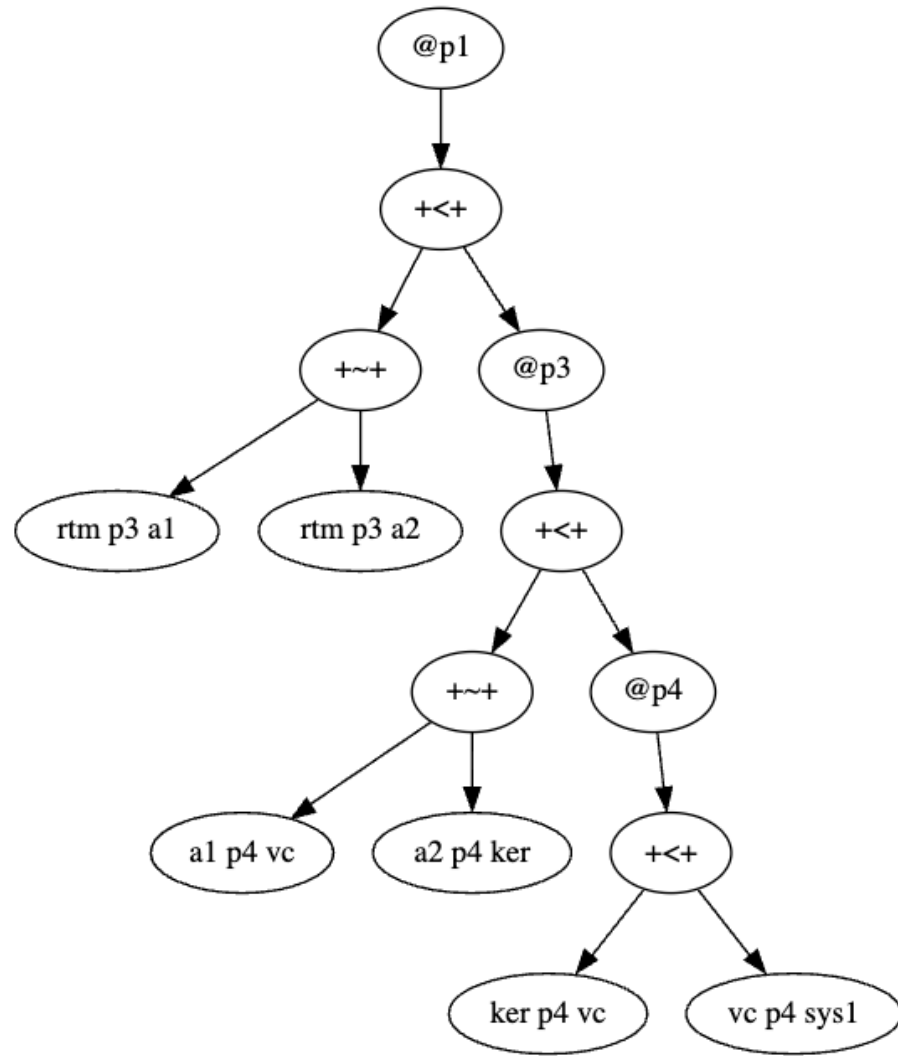
- Increase cost after start event
- Any corruption of a deeper component is higher cost
- Add weight to cost event
  - Have some base cost to the corruption event
  - Add to cost more if its in a protected place

# Protocols

Protocol Name	Protocol
sys	*target: @p4 [vc p4 sys1]
vc-sys-seq	*target: @p3 [a p4 vc] +<+ @p4 [vc p4 sys]
vc-sys-par	*target: @p3 [a p4 vc] +~+ @p4 [vc p4 sys]
a-vc-sys-seq	*target: @p1 [rtm p3 a] +<+ @p3 [a p4 vc] +<+ @p4 [vc p4 sys]
a-vc-sys-par	*target: @p1 [rtm p3 a +~+ @p3 [a p4 vc +~+ @p4 [vc p4 sys]]]
a1-a2-vc-ker-sys	*target: @p1 ( rtm p3 a1 +~+ rtm p3 a2) +<+ @p3 ( a1 p4 vc +~+ a2 p4 ker ) +<+ @p4 ((ker p4 vc) +<+ (vc p4 sys1 ))

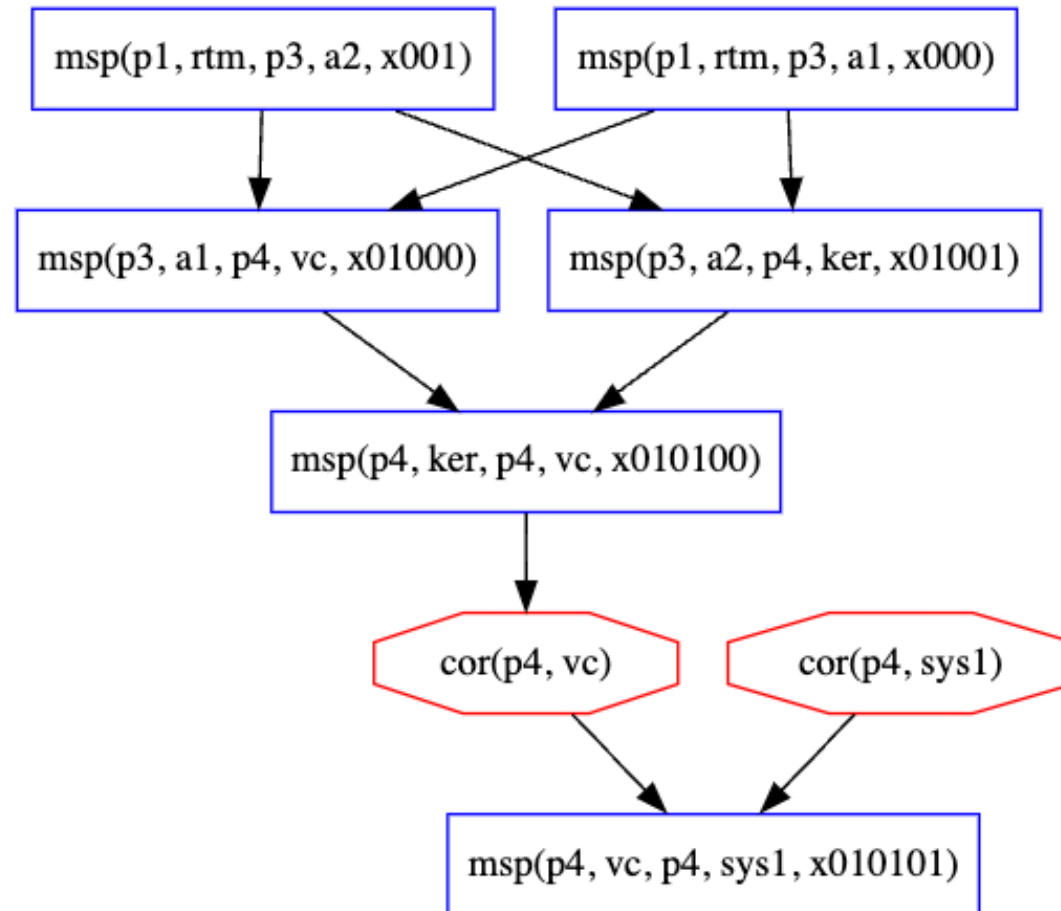


## Abstract Syntax Tree

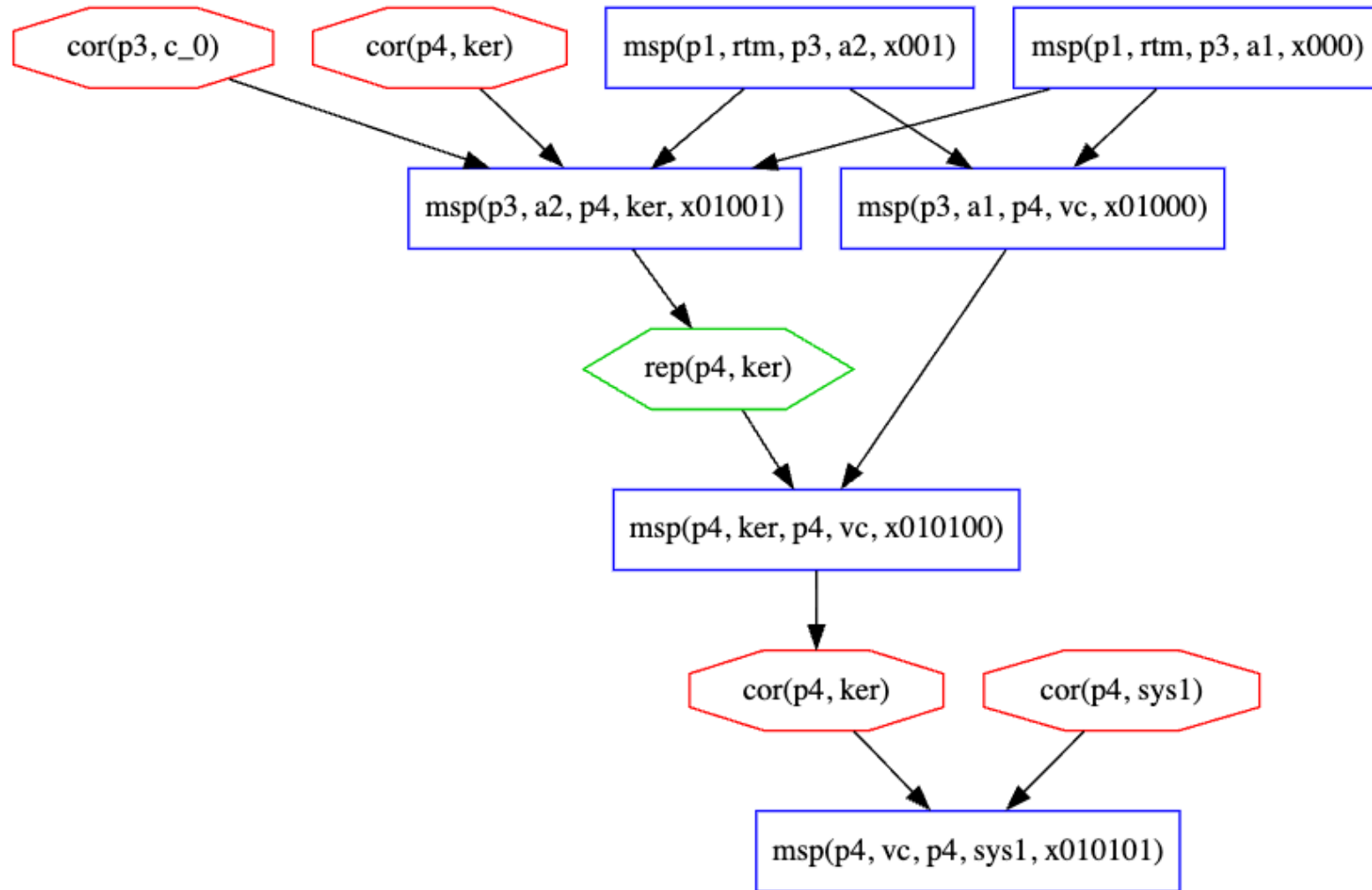


Assuming recent measurements may be corrupted there are 21 models...

**Model 1**



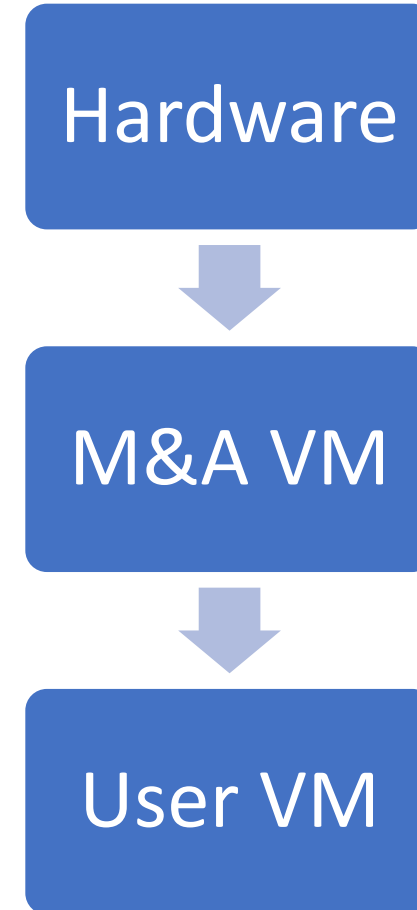
### Model 21



Order	Event	Cost	Present In	Details
low	cor(p4,sys1)	c1	all models	Always before the last measurement event
	cor(p4,c)	c2	4	happens before ms(ker,vc)
	cor(p4,vc)	c3	1 4 5 8 9 10 11 12 13 14 15 16 17 18(2) 19(2)	occurs after some attestation start event (between measurements) or before a measurement, sometimes happens twice (once and then after a repair)
	cor(p4, ker)	c4	2 3 5 6 7 10 11 12 13 14 15 16 17 20(2) 21(2)	occurs various places.. Before/after ms(a2,ker), before ms(vc,sys1)
	cor(p4,c_1)	c5	8	before ms(ker,vc)
	cor(p3,a1)	c6	8 10 14 15 18	before ms(a1,vc), always after the attestation begins... maybe this is most difficult because you have to consider time window for adversary
	cor(p3,c_3)	c7	9 11 16 17 19	before ms(a1,vc), no attestation start event... could be easiest for an adversary
	cor(p4,c_2)	c8	9	before ms(ker,vc), no attestation start event... could be easiest for adversary
	cor(p3,a2)	c9	6 12 14 16 20	between ms(rtm, a2) – ms(a2, ker)... close to root of trust. Difficult for an adversary
	cor(p3,c_4)	c10	13	before ms(a2,ker), no attestation start event... could be easiest for adversary
	cor(p3,c_5)	c11	15	before ms(a2,ker) no attestation start event... could be easiest for adversary
	cor(p3,c_6)	c12	17	before ms(a2,ker) , no attestation start event... could be easiest for adversary
	rep(p4,vc)	c13	18 19	between ms(a1,vc) – ms(ker,vc),
	rep(p4,ker)	c14	20 21	between ms(a2,ker) -- ms(ker, vc)
	cor(p3,c_0)	c15	7 21	before ms(a2,ker)

# Considering Cost to an Adversary

- Hardware = highest cost to adversary
- M&A = middle cost
- User VM = lowest cost to an adversary



# Considering Cost to an Attester

- Hardware = worst case for an adversary
- M&A = ??
- User VM = ??

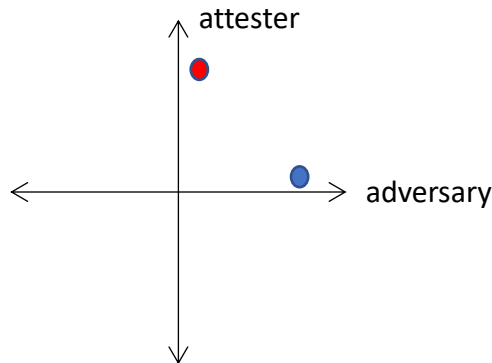
# Cost

Cost	Reasoning
high	<ul style="list-style-type: none"><li>• corruption events that occur between measurement events are difficult... Thus, high cost</li><li>• corruption events closer to root of trust are difficult. Thus, high cost.</li><li>• corruption then repair then corruption requires a lot of work from adversary. This is a high cost.</li></ul>
medium	<ul style="list-style-type: none"><li>• corruption event at M&amp;A domain is medium as it is in the middle of the architecture</li></ul>
low	<ul style="list-style-type: none"><li>• corruption before last measurement is probably the easiest thing for an adversary therefore the lowest cost.</li></ul>

Slides 4.14.23

# Goals of cost analysis

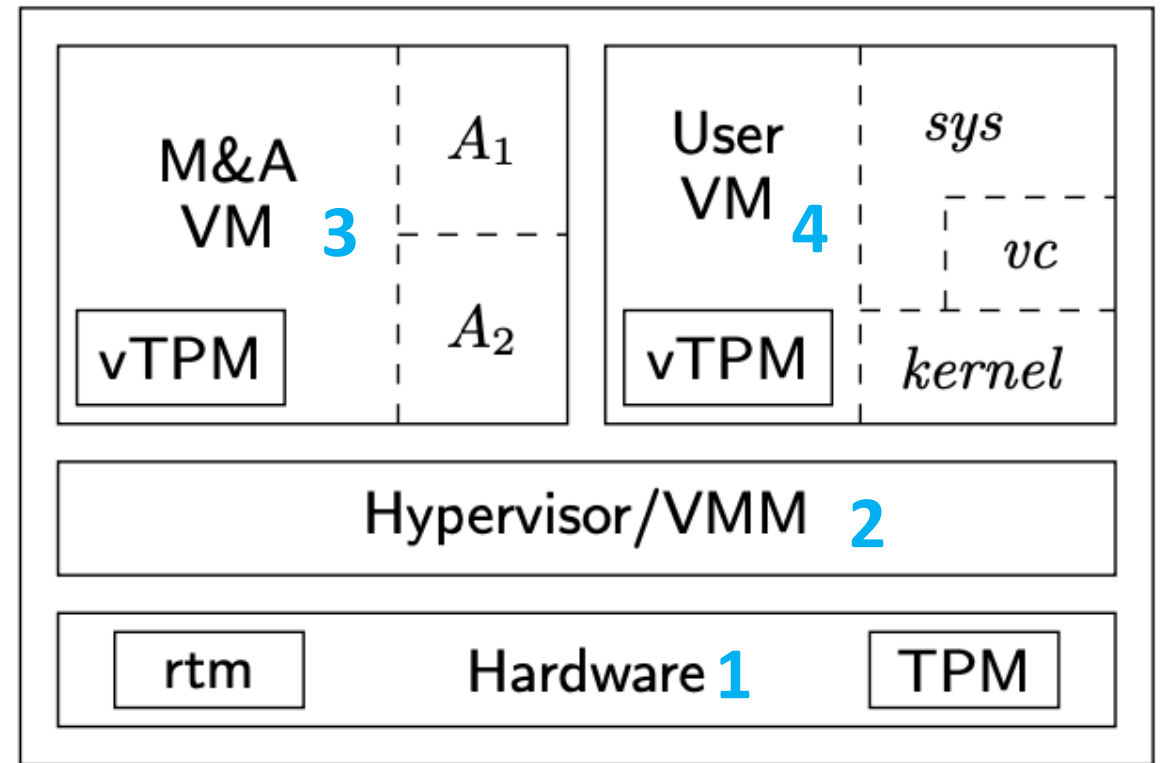
- Ultimate goal: guide selection of a protocol
- How:
  - Systematic variation of assumption
  - Assign abstract cost to each component that's corrupted
- Consider:
  - Cost to adversary
  - Cost to attester





# Say we have the architecture from “Confining the Adversary” Paper

- $ms(rtm, A1)$
- ~~$ms(rtm, A2)$~~
- $ms(A1, vc)$
- ~~$ms(A2, ker)$~~
- $msker(vc, sys)$



# Control Variables

```
% Assume sys depends on kernel
% if sys1 or vc depend on anything, that thing is the root of trust
depends(p1, C, p4, sys) => C = rtm.
depends(p1, C, p4, vc) => C = rtm.
depends(p1, C, p3, a) => C = rtm.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.
```

- Assumptions

- Always assume recent/deep
- Make no assumptions about system dependencies except...
  - TPM is the root of trust... has no dependencies
  - Virus checker and system depend on the hardware (p1,rtm)
  - A1 depends on the hardware (p1,rtm)

```
% Assume no deep corruptions
l(V) = cor(p1, M) => false.
```

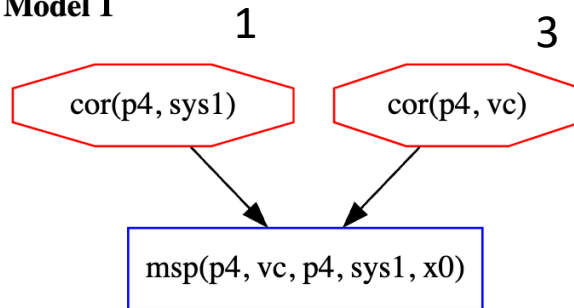
Side note: I changed all theory files to the original... allows for corruption only at the same place

- If I made it allow for corruption at different places... CHASE seemed to introduce corruption events with odd labels

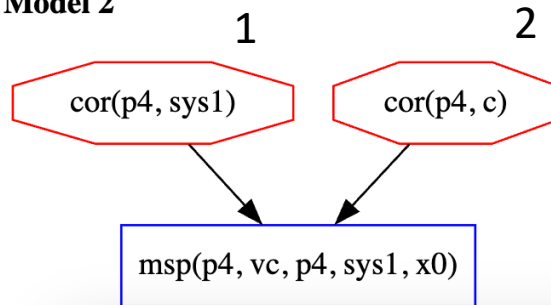
# First protocol.... Just measure *sys* using *vc*

## Models

### Model 1



### Model 2



Event	Cost
$\text{cor}(p4, \text{sys}1)$	$c1$
$\text{cor}(p4, \text{vc})$	$c3$
$\text{cor}(p4, c)$	$c2$
MODEL 1 COST	$c1 + c3$
MODEL 2 COST	$c1 + c2$

# Measure vc and sys in parallel

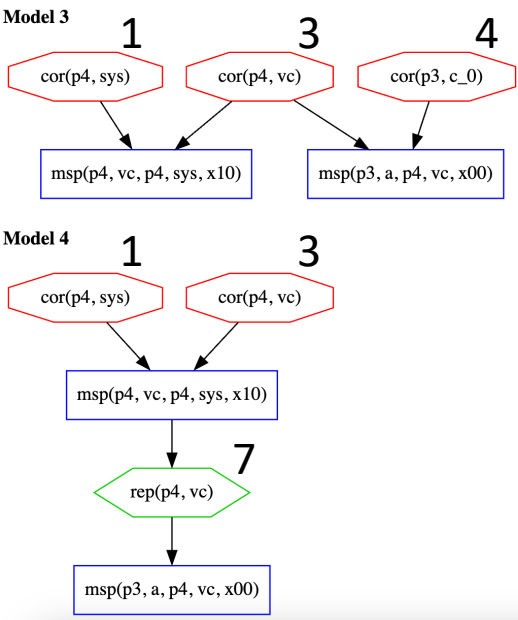
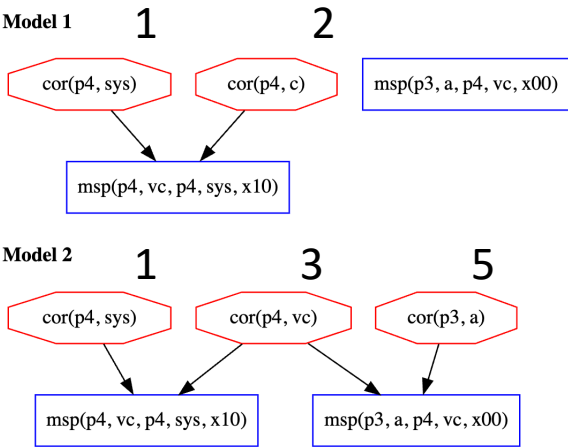
- Protocol
  - \*target: @p3 [a p4 vc]  
+~+ @p4 [vc p4 sys]

```
% Assume dependencies
% if sys1 or vc depend on anything, that thing is the root of trust
depends(p1, C, p4, sys) => C = rtm.
depends(p1, C, p4, vc) => C = rtm.
depends(p1, C, p3, a) => C = rtm.
% rtm has no dependencies
depends(p1, C, p1, rtm) => false.

% Assume no recent corruptions
prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
=> false.

% Assume no deep corruptions
l(V) = cor(p1, M) => false.
```

Models



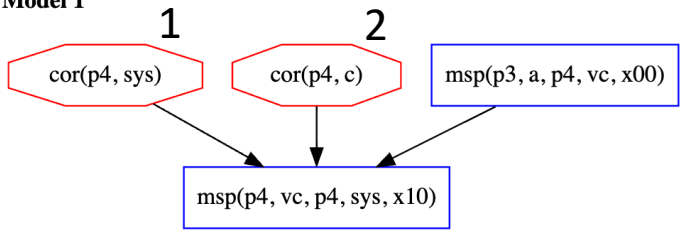
Model	Total cost
1	c1 + c2
2	c1 + c3 + c5
3	c1 + c3 + c4
4	c1 + c3 + c7

# Measure vc and sys in sequence

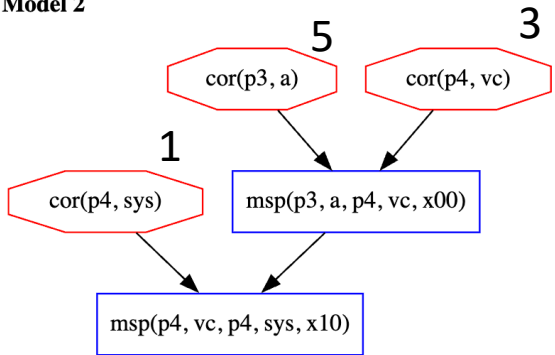
- Protocol
  - \*target: @p3 [a p4 vc]  
+<+ @p4 [vc p4 sys]

Models

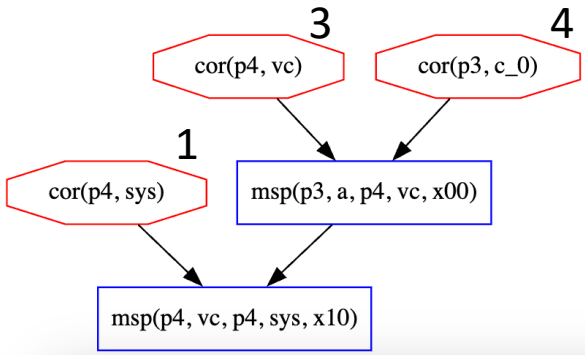
Model 1



Model 2



Model 3



Model	Total cost
1	c1 + c2
2	c1 + c3 + c4
3	c1 + c5 + c3

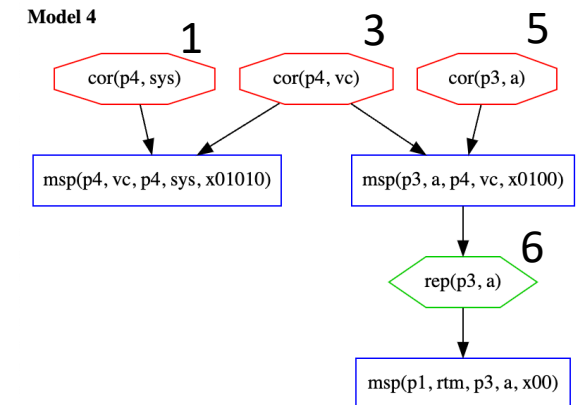
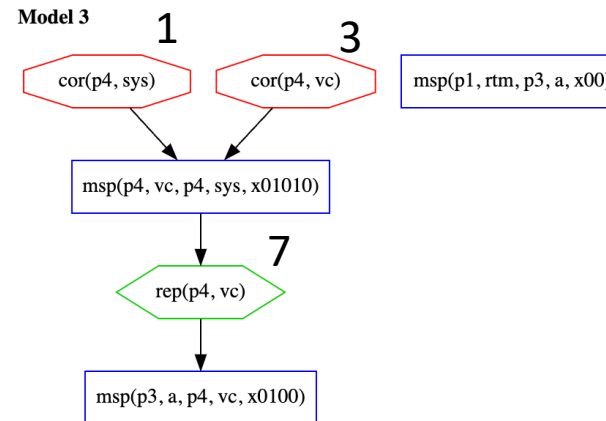
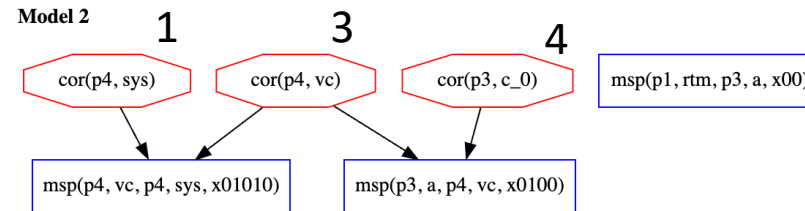
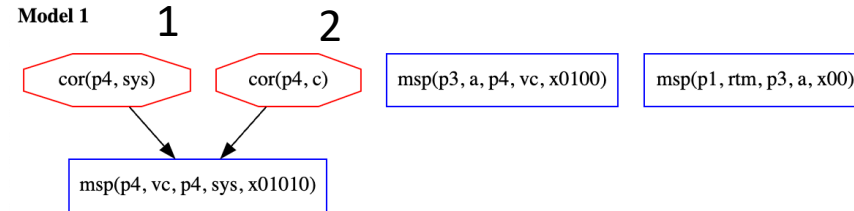
If you add deep thm about p3 model 2 and 3 are removed

# Measure $a$ then $vc$ then $sys$ in parallel

- Protocol
  - \*target: @p1 [rtm p3  $a$   
 $+ \sim +$  @p3 [a p4  $vc$   
 $+ \sim +$  @p4 [vc p4  $sys$ ]]]]

Model	Total cost
1	$c1 + c2$
2	$c1 + c2 + c4$
3	$c1 + c3 + c7$
4	$c1 + c3 + c5 + c6$

## Models



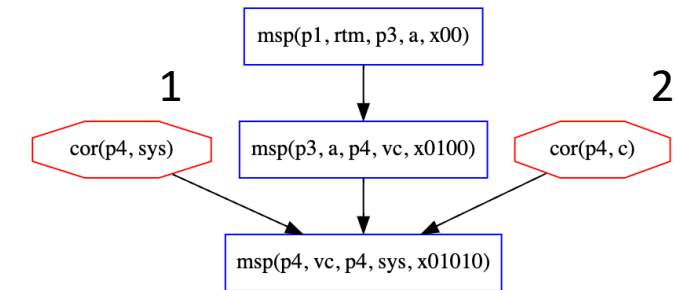
# Measure $a$ then $vc$ then $sys$ in sequence

- Protocol
  - \*target: @p1 [rtm p3  $a$   
   +<+ @p3 [ $a$  p4  $vc$   
   +<+ @p4 [ $vc$  p4  $sys$ ]]]]

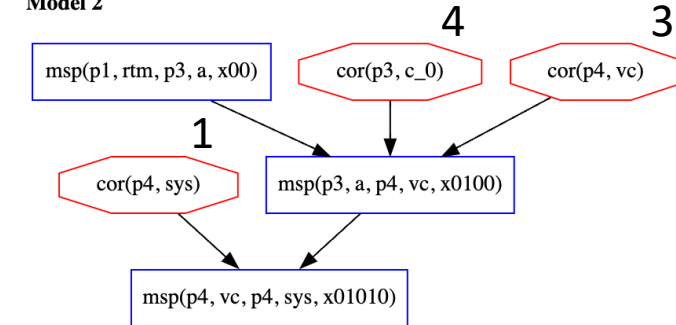
Model	Total cost
1	$c1 + c2$
2	$c1 + c3 + c4$

## Models

### Model 1



### Model 2



# All together

label	protocol	total cost
sys	*target: @p4 [vc p4 sys1]	$(c1 + c3)$ OR $(c1 + c2)$
vc-sys-par	*target: @p3 [a p4 vc] +~+ @p4 [vc p4 sys]	$(c1 + c2)$ OR $(c1 + c3 + c5)$ OR $(c1 + c3 + c4)$ OR $(c1 + c3 + c7)$
vc-sys-seq	*target: @p3 [a p4 vc] +<+ @p4 [vc p4 sys]	$(c1 + c2)$ OR $(c1 + c3 + c4)$ OR $(c1 + c5 + c3)$
a-vc-sys-par	*target: @p1 [rtm p3 $\alpha$ +~+ @p3 [a p4 vc +~+ @p4 [vc p4 sys]]]	$(c1 + c2)$ OR $(c1 + c2 + c4)$ OR $(c1 + c3 + c7)$ OR $(c1 + c3 + c5 + c6)$
a-vc-sys-seq	*target: @p1 [rtm p3 $\alpha$ +<+ @p3 [a p4 vc +<+ @p4 [vc p4 sys]]]	$(c1 + c2)$ OR $(c1 + c3 + c4)$



# Event with label and cost

Event	Label	Cost	Present In
cor(p4,sys)	1	c1	sys(1,2),vc-sys-par(1,2,3,4), vc-sys-seq(1,2,3), a-vc-sys-par(1,2,3,4), a-vc-sys-seq(1,2)
cor(p4,c)	2	c2	sys(2), vc-sys-par(2), vc-sys-seq(1), a-vc-sys-par(1,2), a-vc-sys-seq(1)
cor(p4,vc)	3	c3	sys(1), vc-sys-par(2,3,4), vc-sys-seq(2,3), a-vc-sys-par(3,4), a-vc-sys-seq(2)
cor(p3, c_0)	4	c4	vc-sys-par(3), vc-sys-seq(2), a-vc-sys-par(2), a-vc-sys-seq(2)
cor(p3,a)	5	c5	vc-sys-par(2), vc-sys-seq(3), a-vc-sys-par(4)
rep(p3,a)	6	c6	a-vc-sys-par(4)
rep(p4,vc)	7	c7	vc-sys-par(4), a-vc-sys-par(3)

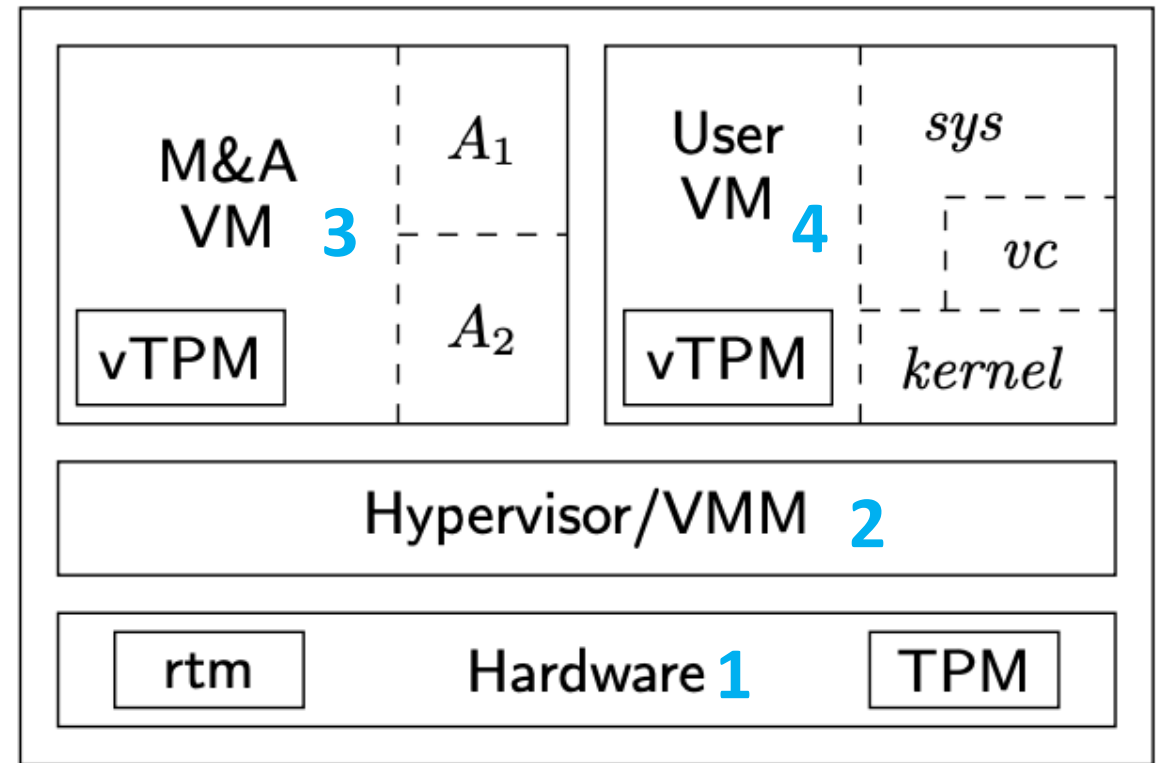
Slides 4.7.23

# Goals of cost analysis

- Ultimate goal: guide selection of a protocol
- How:
  - systematic variation of assumption
  - assigning cost to each component that's corrupted
    - Assign low (or high?) values to difficult actions
    - Realize set of protocols, one with minimum (maximum) cost
    - Cost may reflect ordering

# Say we have the architecture from “Confining the Adversary” Paper

- $ms(rtm, A1)$
- $ms(rtm, A2)$
- $ms(A1, vc)$
- $ms(A2, ker)$
- $msker(vc, sys)$



# First protocol.... Just measure *sys* using *vc*

- Protocol:

- @4 [vc 4 sys1]

msp(p4, vc, p4, sys1, x0)

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at our main measurement
% event. This is a measurement of sys
l(V) = msp(p4, M, p4, sys1, X)
=> corrupt_at(p4, sys1, V).

% Assume no dependencies
depends(p4, C, p4, sys1) => false.

% No recent assumptions
% No deep assumptions

m4_include(`sys.gli')m4_dnl
m4_include(`sys_dist.gli')m4_dnl
m4_include(`thy.gli')m4_dnl
```

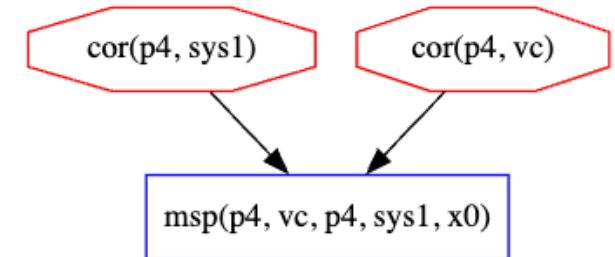
- Cost?

- Potentially 2?

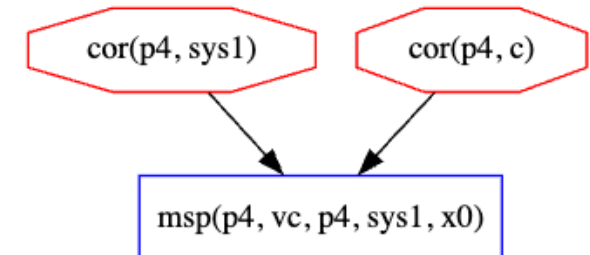
- 2 places where corruptions could occur

## Models

### Model 1



### Model 2



# Measure *vc* and *sys* in parallel

- Protocol
  - \*target: @p3 [a p4 vc]  
+~+ @p4 [vc p4 sys]
- Cost?
  - Two corruption events and a repair event...
  - What should be the cost of a repair?

msp(p4, vc, p4, sys, x10)

msp(p3, a, p4, vc, x00)

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at
% our main measurement event.
% This is a measurement of sys.
l(V) = msp(p4, M, p4, sys, X)
=> corrupt_at(p4, sys, V).

% Assume sys depends on kernel
depends(p4, C, p4, sys) => false.
depends(p4, C, p4, vc) => false.
depends(p3, C, p3, a) => false.

% Assume no recent corruptions
prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
=> false.

% Assume no deep corruptions
l(V) = cor(p3, M) => false.

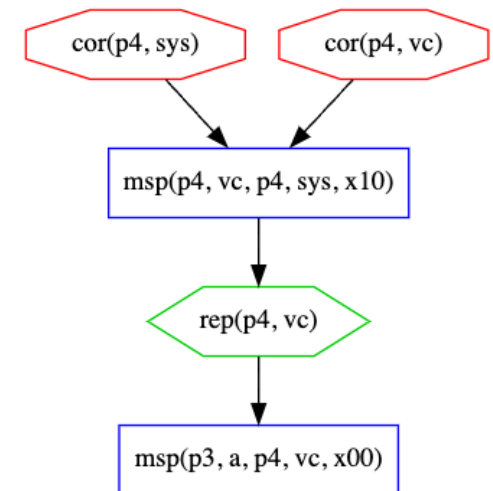
m4_include(`vc-sys.gli')m4_dnl

m4_include(`vc-sys_dist.gli')m4_dnl

m4_include(`thy.gli')m4_dnl
```

## Models

### Model 1



# Measure vc and sys in sequence

- Protocol
  - \*target: @p3 [a p4 vc]  
+<+ @p4 [vc p4 sys]
- Analysis
  - No models if recent or deep assumption... this is expected

msp(p3, a, p4, vc, x00)



msp(p4, vc, p4, sys, x10)

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at
% our main measurement event.
% This is a measurement of sys.
l(V) = msp(p4, M, p4, sys, X)
=> corrupt_at(p4, sys, V).

% Assume sys depends on kernel
% depends(p3, C, p3, a) => C = p1.
depends(p4, C, p4, sys) => false.
depends(p4, C, p4, vc) => false.
depends(p3, C, p3, a) => false.

% Assume no recent corruptions
prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
=> false.

% Assume no deep corruptions
l(V) = cor(p3, M) => false.

m4_include(`vc-sys-seq.gli')m4_dnl

m4_include(`vc-sys-seq_dist.gli')m4_dnl

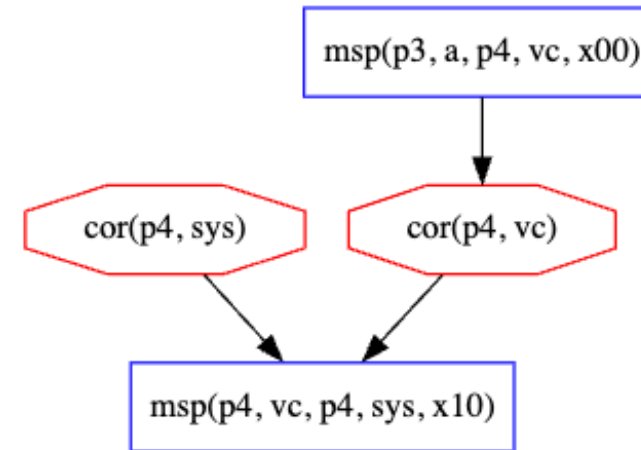
m4_include(`thy.gli')m4_dnl
```

Same protocol....  
No recent or deep  
assumption

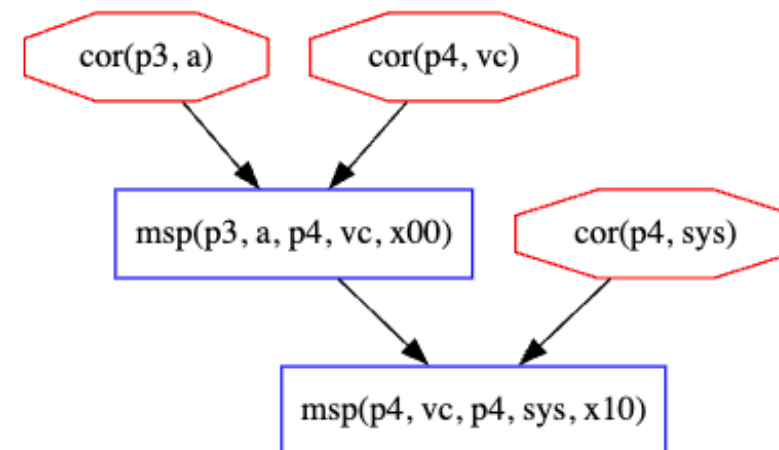
- This makes me think... What is the cost of including the recent/deep theorem?

## Models

### Model 1



### Model 2





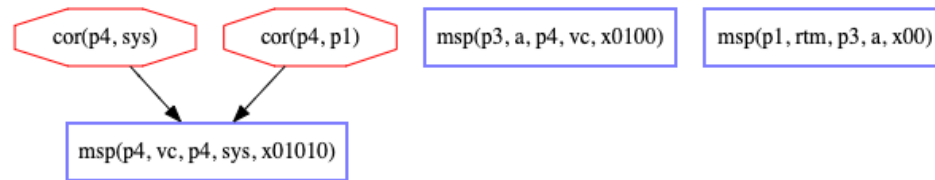
# Measure $a$ then $vc$ then $sys$ in parallel

- Protocol

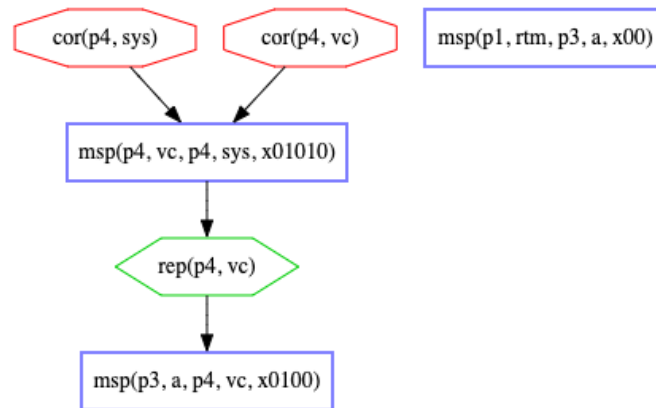
- \*target: @p1 [rtm p3  $a$   
 $+ \sim +$  @p3 [ $a$  p4  $vc$   
 $+ \sim +$  @p4 [ $vc$  p4  $sys$ ]]]]

## Models

### Model 1



### Model 2



msp(p3, a, p4, vc, x0100)

msp(p1, rtm, p3, a, x00)

msp(p4, vc, p4, sys, x01010)

## Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]
```

```
% Assume adversary avoids detection at
% our main measurement event.
% This is a measurement of sys.
l(V) = msp(p4, M, p4, sys, X)
=> corrupt_at(p4, sys, V).
```

```
% system dependencies
depends(p3, C, p3, a) => C = p1.
depends(p1, C, p1, rtm) => false.
depends(p4, C, p4, sys) => C = p1.
depends(p4, C, p4, vc) => C = p1.
```

```
% Assume no recent corruptions
prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
=> false.
```

```
% Assume no deep corruptions
l(V) = cor(p3, M) => false.
```

```
m4_include(`a-vc-sys-par.gli')m4_dnl
```

```
m4_include(`a-vc-sys-par_dist.gli')m4_dnl
```

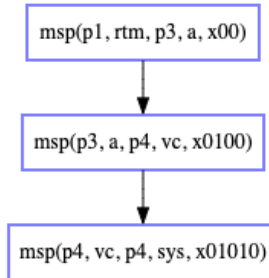
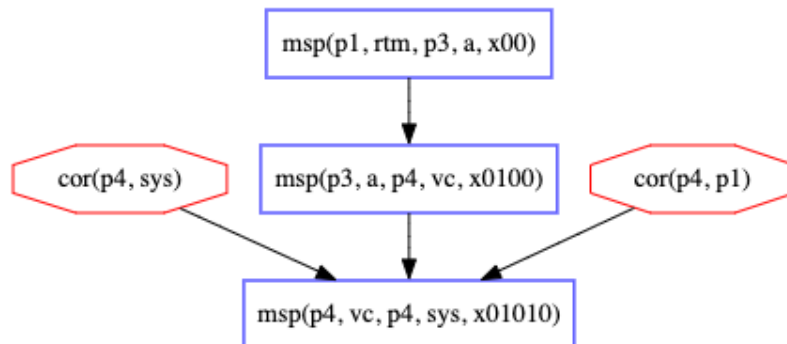
```
m4_include(`thy.gli')m4_dnl
```

# Measure $a$ then $vc$ then $sys$ in sequence

- Protocol
  - \*target: @p1 [rtm p3  $a$   
+<+ @p3 [ $a$  p4  $vc$   
+<+ @p4 [ $vc$  p4  $sys$ ]]]]

## Models

### Model 1



### Problem Configuration

```
[ bound = 500, limit = 5000, input_order ]

% Assume adversary avoids detection at
% our main measurement event.
% This is a measurement of sys.
l(V) = msp(p4, M, p4, sys, X)
=> corrupt_at(p4, sys, V).

% Assume sys depends on kernel
depends(p3, C, p3, a) => C = p1.
depends(p1, C, p1, rtm) => false.
depends(p4, C, p4, sys) => C = p1.
depends(p4, C, p4, vc) => C = p1.

% Assume no recent corruptions
prec(V, V1) & l(V1) = cor(P,C) & ms_evt(V)
=> false.

% Assume no deep corruptions
l(V) = cor(p3, M) => false.

m4_include(`a-vc-sys-seq.gli')m4_dnl

m4_include(`a-vc-sys-seq_dist.gli')m4_dnl

m4_include(`thy.gli')m4_dnl
```

# Thoughts/Takeaways

- Cost of adding an assumption?
- Cost of adding a dependencies?
- Cost of applying recent/deep theorem?
  - Should we consider this a standard assumption?
- Cost of a corruption/repair event?
  - Maybe turn protocol execution into a tree... then could look at depth of corruption/repair event and that could be the event's cost. Sum all costs together and that is the total cost.

## Models

### Model 1

