## Simple but longterm card-terminal authorization protocol based on one time passwords - sketch of protocol

## **Prerequisites**

- Each card has a unique ID (*Card<sub>ID</sub>*) and stores its current state (*ST*), which is simultaneously a symmetric key used for secure communication with the terminal.
- Terminal stores a mapping from card IDs to pair of their current and previous state. We assume that for a given card, the initial state of the card and the corresponding current state that terminal holds are the same (initial previous state stored by the terminal is taken at random).

## **Definitions**

```
• I\mathcal{D} - card IDs space (\{0,1\}^{32})
```

```
• \mathcal{R} - challenges space (\{0,1\}^{64})
```

```
• \mathcal{K} - key space (\{0,1\}^{128} \times \{0,1\}^{256})
```

- Enc encryption (AES in CFB mode with 256b key)
- Dec decryption (AES in CFB mode with 256b key)
- ACRT acceptable card response time (exact value to be defined)
- *time*() function that returns current timestamp
- $f: I\mathcal{D} \to \mathcal{K} \times \mathcal{K}$  mapping from card IDs to pair of states (previous and the current one).

Authentication protocol (simple pre-shared key challenge-response authentication): If the protocol is executed successfully, terminal opens the door to the secure location.

## **ASN.1 Documentation**

```
CardProtocol DEFINITIONS ::= BEGIN
    CardHello ::= SEQUENCE {
    cardId BIT STRING
    }

RandomChallenge ::= SEQUENCE challenge BIT STRING
    StageOne ::= SEQUENCE oldState BIT STRING
    StageTwo ::= SEQUENCE newState BIT STRING
    END
```

Terminal $(f)$	Transmission	Card $(Card_{ID}, ST)$
1.		
	$\leftarrow Card_{ID}$	
2. Take $r \in \mathcal{R}$ uniformly at random.		
Let $t := time()$		
	$\rightarrow r$	- ()
3.		$m_1 := Enc_{ST}(r)$
4 I at 4	$\leftarrow m_1$	
4. Let $t' := time()$ . Check if $t' - t < ACRT$ (If not, show error mes-		
sage about card response being too long and		
abort.)		
abort.)		
Let $(k_{prev}, k_{curr}) := f(Card_{ID})$ and check		
if (1) $Dec_{k_{prev}}(m_1) = r$ or (2) $Dec_{k_{curr}}(m_1) = r$		
(If (1) is fulfilled, show warning that terminal		
and card got desynchronized. If none is ful-		
filled, show error message about card being in		
an incorrect state, suggesting that it may have		
been cloned and abort).		
Let $k_{good} \in \{k_{prev}, k_{curr}\}$ be the one that ful-		
filled one of the equalities.		
Take $k' \in \mathcal{K} \setminus \{k_{good}\}$ uniformly at random		
and update $f$ so that $f(Card_{ID}) = (k_{good}, k')$ .		
$m_2 := Enc_{k_{good}}(k')$		
m2. Zinckgood (N)	$\rightarrow m_2$	
5.		$ST := Dec_{ST}(m_2)$