

Introduction to Natural Language Processing I [Statistické metody zpracování přirozených jazyků I] (NPFL067)

<https://ufal.mff.cuni.cz/courses/npfl067>

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Intro to Stat. NLP I

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 - ÚFAL MFF UK, office: 420 / 422 MS, office hours: J. Hajic: Mon 10:00-11:00
 - preferred contact: {hajic,pecina}@ufal.mff.cuni.cz
- Room & time:
 - lecture: room S1, Tue 9:00-10:40 (only Oct, 4)
 - seminar [cvičení] room S5, Tue 12:20-14:00
 - Oct 4, 2021 – Jan 4, 2022, final written exam (probable) date: Jan 11, 2022
- Inverted course:
 - you watch the lectures at your leisure
 - Interactive part at seminars (in/person, online for those with travel difficulties)

Textbooks you need

- Manning, C. D., Schütze, H.:
 - *Foundations of Statistical Natural Language Processing*. The MIT Press. 1999. ISBN 0-262-13360-1. **[required]**
- Jurafsky, D., Martin, J.H.:
 - *Speech and Language Processing*. Prentice-Hall. 2000. ISBN 0-13-095069-6 and later editions. **[recommended]**.

Other reading

- Charniak, E:
 - *Statistical Language Learning*. The MIT Press. 1996. ISBN 0-262-53141-0.
- Cover, T. M., Thomas, J. A.:
 - *Elements of Information Theory*. Wiley. 1991. ISBN 0-471-06259-6.
- Jelinek, F.:
 - *Statistical Methods for Speech Recognition*. The MIT Press. 1998. ISBN 0-262-10066-5
- Proceedings of major conferences (ACL Anthology)
 - ACL (Assoc. of Computational Linguistics)
 - EACL/NAACL/IJCNLP (European/American/Asian Chapter of ACL)
 - EMNLP (Empirical Methods in NLP), CONLL
 - COLING (Intl. Committee of Computational Linguistics)
- arXiv, journals: Computational Linguistics, TACL

Course requirements

- Grade components: requirements & weights:
 - Homeworks (2): 33,3%
 - Final Exam: 33,3%
- Exam:
 - approx. 4-5 questions:
 - mostly explanatory answers (1/4 page or so),
 - algorithms
 - only a few multiple choice questions


Homeworks

- Homeworks:
 - Entropy, Language Modeling; Word Classes
- Organization
 - (little) paper-and-pencil exercises, lot of programming
 - turning-in mechanism: see the web
 - no plagiarism!
- Deadlines
 - See the course web (Jan. 31, 2021; Feb 28, 2021)
 - Late penalty: 5% of grade (0-100) per day (max. 50%)

Course segments

- Intro & Probability & Information Theory
 - The very basics: definitions, formulas, examples.
- Language Modeling
 - n-gram models, parameter estimation
 - smoothing (EM algorithm)
- Words and the Lexicon
 - word classes, mutual information, bit of lexicography
- Hidden Markov Models
 - background, algorithms, parameter estimation

NLP: The Main Issues

- Why is NLP difficult?
 - many “words”, many “phenomena” --> many “rules”
 - **OED: 400k words; Finnish lexicon (of forms): $\sim 2 \cdot 10^7$**
 - **sentences, clauses, phrases, constituents, coordination, negation, imperatives/questions, inflections, parts of speech, pronunciation, topic/focus, and much more!**
 - irregularity (exceptions, exceptions to the exceptions, ...)
 - **potato -> potato^{es} (tomato, hero,...); photo -> photo^s, and even: both mango -> mango^s or -> mango^{es}**
 - **Adjective / Noun order: new book, electrical engineering, general regulations, flower garden, garden flower, ...: but Governor General**
- 

Difficulties in NLP (cont.)

- ambiguity

- **books: NOUN or VERB?**

- you **need** many books vs. she **books** her flights online

- **No left turn weekdays 4-6 pm / except transit vehicles**
(*Charles Street at Cold Spring*)

- when may transit vehicles turn: Always? Never?

- **Thank you for not smoking, drinking, eating or playing radios without earphones.** (*MTA bus*)

- Thank you for not eating without earphones??

- or even: Thank you for ~~not~~ drinking without earphones!?

- **My neighbor's hat was taken by wind. He tried to catch it.**

- ...catch the wind or ...catch the hat ?

(Categorical) Rules or Statistics?

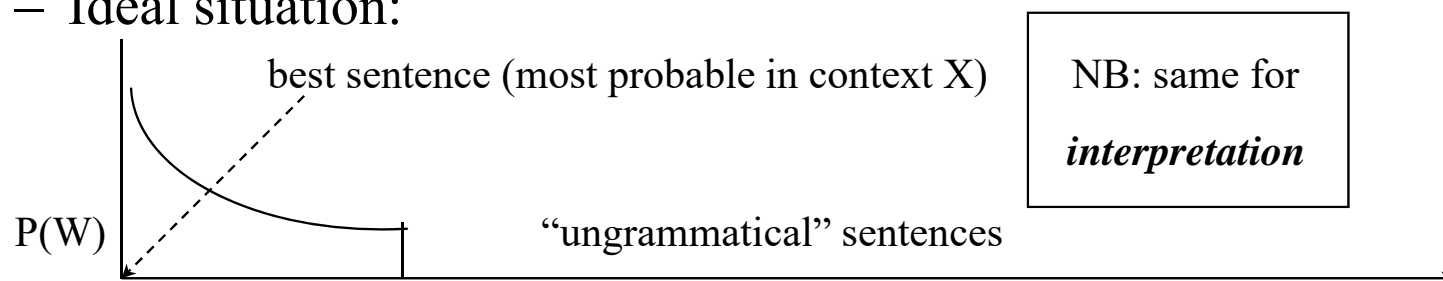
- Preferences:
 - clear cases: context clues: she books --> books is a verb
 - rule: if an ambiguous word (verb/nonverb) is preceded by a matching personal pronoun -> word is a verb
 - less clear cases: pronoun reference
 - she/he/it refers to the most recent noun or pronoun (?) (but maybe we can specify exceptions)
 - selectional:
 - catching hat >> catching wind (but why not?)
 - semantic:
 - never thank for drinking in a bus! (but what about the earphones?)

Solutions

- Don't guess if you know:
 - **morphology (inflections)**
 - **lexicons (lists of words)**
 - **unambiguous names**
 - **perhaps some (really) fixed phrases**
 - **syntactic rules?**
- Use statistics (based on real-world data) for preferences (only?)
 - **No doubt: but this is the big question!**

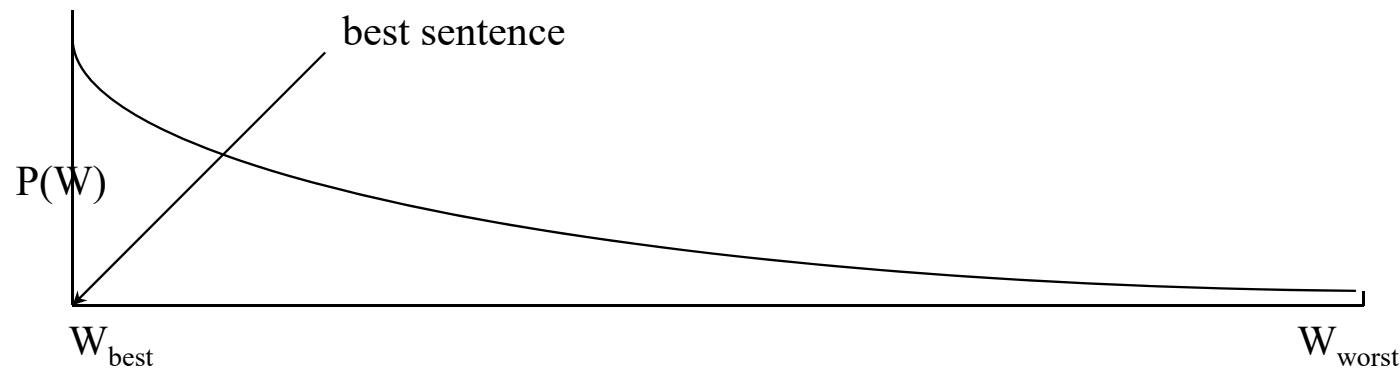
Statistical NLP

- Imagine:
 - Each sentence $W = \{ w_1, w_2, \dots, w_n \}$ gets a probability $P(W|X)$ in a context X (think of it in the intuitive sense for now)
 - For every possible context X , sort all the imaginable sentences W according to $P(W|X)$:
 - Ideal situation:



Real World Situation

- Unable to specify set of grammatical sentences today using fixed “categorical” rules (maybe never, cf. arguments in MS)
- Use statistical “model” based on **REAL WORLD DATA** and care about the best sentence only (disregarding the “grammaticality” issue)



Probability

Experiments & Sample Spaces

- Experiment, process, test, ...
- Set of possible basic outcomes: sample space Ω
 - coin toss ($\Omega = \{\text{head}, \text{tail}\}$), die ($\Omega = \{1..6\}$)
 - yes/no opinion poll, quality test (bad/good) ($\Omega = \{0,1\}$)
 - lottery ($|\Omega| \cong 10^7 \dots 10^{12}$)
 - # of traffic accidents somewhere per year ($\Omega = \mathbb{N}$)
 - spelling errors ($\Omega = Z^*$), where Z is an alphabet, and Z^* is a set of possible strings over such an alphabet
 - missing word ($|\Omega| \cong \text{vocabulary size}$)

Events

- Event A is a set of basic outcomes
- Usually $A \subset \Omega$, and all $A \in 2^\Omega$ (the event space)
 - Ω is then the certain event, \emptyset is the impossible event
- Example:
 - experiment: three times coin toss
 - $\Omega = \{\text{HHH}, \text{HHT}, \text{HTH}, \text{HTT}, \text{TTH}, \text{THT}, \text{TTH}, \text{TTT}\}$
 - count cases with exactly two tails: then
 - $A = \{\text{HTT}, \text{THT}, \text{TTH}\}$
 - all heads:
 - $A = \{\text{HHH}\}$

Probability

- Repeat experiment many times, record how many times a given event A occurred (“count” c_1).
- Do this whole series many times; remember all c_i s.
- Observation: if repeated really many times, the ratios of c_i/T_i (where T_i is the number of experiments run in the i -th series) are close to some (unknown but) **constant** value.
- Call this constant a **probability of A** . Notation: $p(A)$

Estimating probability

- Remember: ... close to an *unknown* constant.
- We can only estimate it:
 - from a single series (typical case, as mostly the outcome of a series is given to us and we cannot repeat the experiment), set

$$p(A) = c_1/T_1.$$

- otherwise, take the weighted average of all c_i/T_i (or, if the data allows, simply look at the set of series as if it is a single long series).
- This is the **best** estimate.

Example

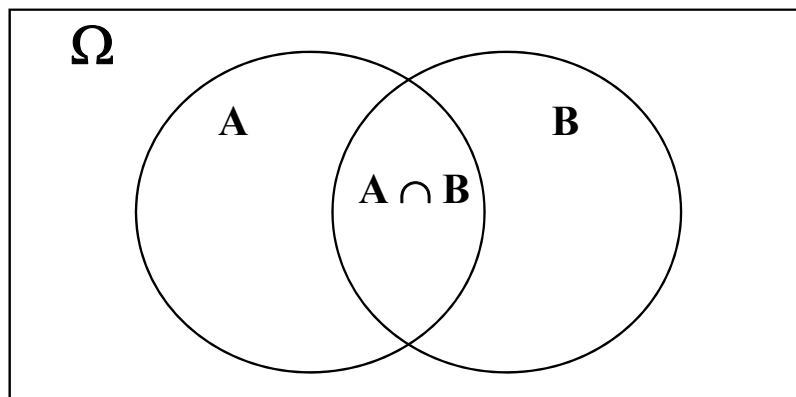
- Recall our example:
 - experiment: three times coin toss
 - $\Omega = \{\text{HHH}, \text{HHT}, \text{HTH}, \text{HTT}, \text{THH}, \text{THT}, \text{TTH}, \text{TTT}\}$
 - count cases with exactly two tails: $A = \{\text{HTT}, \text{THT}, \text{TTH}\}$
- Run an experiment 1000 times (i.e. 3000 tosses)
- Counted: 386 cases with two tails (HTT, THT, or TTH)
- estimate: $p(A) = 386 / 1000 = .386$
- Run again: 373, 399, 382, 355, 372, 406, 359
 - $p(A) = .379$ (weighted average) or simply $3032 / 8000$
- *Uniform* distribution assumption: $p(A) = 3/8 = .375$

Basic Properties

- Basic properties:
 - $p: 2^{\Omega} \rightarrow [0,1]$
 - $p(\Omega) = 1$
 - Disjoint events: $p(\cup A_i) = \sum_i p(A_i)$
- [NB: axiomatic definition of probability: take the above three conditions as axioms]
- Immediate consequences:
 - $p(\emptyset) = 0$, $p(\bar{A}) = 1 - p(A)$, $A \subseteq B \Rightarrow p(A) \leq p(B)$
 - $\sum_{a \in \Omega} p(a) = 1$

Joint and Conditional Probability

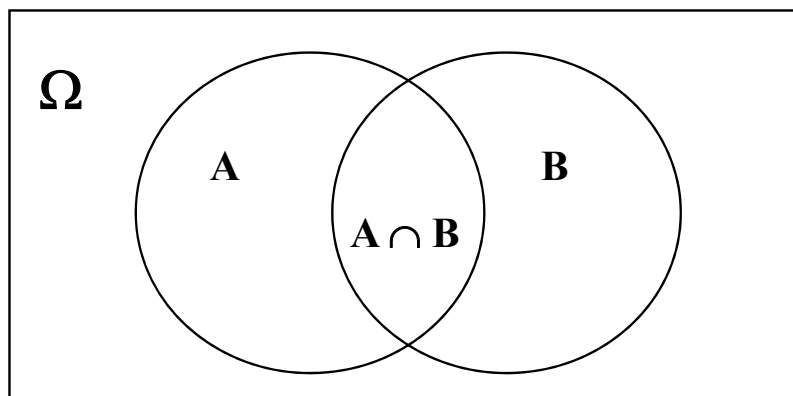
- $p(A,B) = p(A \cap B)$
- $p(A|B) = p(A,B) / p(B)$
 - Estimating from counts:
 - $p(A|B) = p(A,B) / p(B) = (c(A \cap B) / T) / (c(B) / T) = c(A \cap B) / c(B)$



Bayes Rule

- $p(A,B) = p(B,A)$ since $p(A \cap B) = p(B \cap A)$
 - therefore: $p(A|B) p(B) = p(B|A) p(A)$, and therefore

$$p(A|B) = p(B|A) p(A) / p(B)$$



Independence

- Can we compute $p(A,B)$ from $p(A)$ and $p(B)$?
- Recall from previous foil:

$$p(A|B) = p(B|A) p(A) / p(B)$$

$$p(A|B) p(B) = p(B|A) p(A)$$


$$p(A,B) = p(B|A) p(A)$$

... we're almost there: how $p(B|A)$ relates to $p(B)$?

– $p(B|A) = P(B)$ iff A and B are **independent**

- Example: two coin tosses, weather today and weather on March 4th 1789;
- Any two events for which $p(B|A) = P(B)$!

Chain Rule

$$p(A_1, A_2, A_3, A_4, \dots, A_n) =$$
$$p(A_1|A_2, A_3, A_4, \dots, A_n) \times p(A_2|A_3, A_4, \dots, A_n) \times$$
$$\times p(A_3|A_4, \dots, A_n) \times \dots p(A_{n-1}|A_n) \times p(A_n)$$


- this is a direct consequence of the Bayes rule.

The Golden Rule (of Classic Statistical NLP)

- Interested in an event A given B (when it is not easy or practical or desirable to estimate $p(A|B)$):
- take Bayes rule, max over all As:
- $\operatorname{argmax}_A p(A|B) = \operatorname{argmax}_A p(B|A) \cdot p(A) / p(B) =$

$\operatorname{argmax}_A p(B|A) p(A) \bullet$

- ... as $p(B)$ is constant when changing As

Random Variable

- is a function $X: \Omega \rightarrow Q$
 - in general: $Q = \mathbb{R}^n$, typically \mathbb{R}
 - easier to handle real numbers than real-world events
- random variable is *discrete* if Q is countable (i.e. also if finite)
- Example: *die*: natural “numbering” $[1,6]$, *coin*: $\{0,1\}$
- Probability distribution:
 - $p_X(x) = p(X=x) =_{\text{df}} p(A_x)$ where $A_x = \{a \in \Omega : X(a) = x\}$
 - often just $p(x)$ if it is clear from context what X is

Expectation

Joint and Conditional Distributions

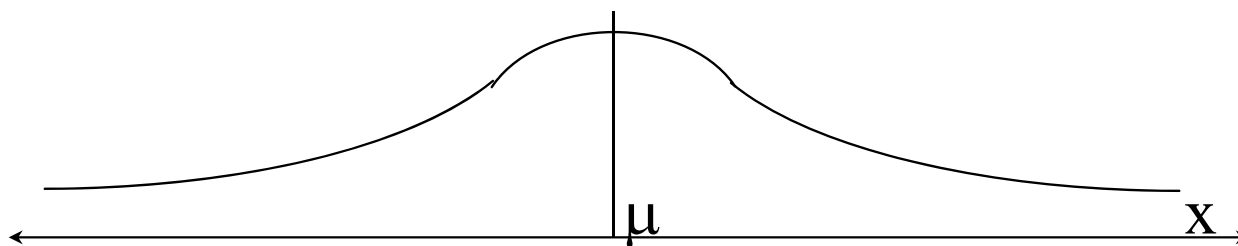
- is a mean of a random variable (weighted average)
 - $E(X) = \sum_{x \in X(\Omega)} x \cdot p_X(x)$
- Example: one six-sided die: 3.5, two dice (sum) 7
- Joint and Conditional distribution rules:
 - analogous to probability of events
- Bayes: $p_{X|Y}(x,y) \stackrel{\text{notation}}{=} p_{XY}(x|y) \stackrel{\text{even simpler notation}}{=} p(x|y) = p(y|x) \cdot p(x) / p(y)$
- Chain rule: $p(w,x,y,z) = p(z) \cdot p(y|z) \cdot p(x|y,z) \cdot p(w|x,y,z)$

Standard distributions

- Binomial (discrete)
 - outcome: 0 or 1 (thus: *binomial*)
 - make n trials
 - interested in the (probability of) number of successes r
- Must be careful: it's not uniform!
- $p_b(r|n) = \binom{n}{r} / 2^n$ (for equally likely outcome)
- $\binom{n}{r}$ counts how many possibilities there are for choosing r objects out of n ; $= n! / ((n-r)! r!)$

Continuous Distributions

- The normal distribution (“Gaussian”)
- $p_{\text{norm}}(x|\mu, \sigma) = e^{-(x-\mu)^2/(2\sigma^2)} / \sigma\sqrt{2\pi}$
- where:
 - μ is the mean (x-coordinate of the peak) (0)
 - σ is the standard deviation (1)



- other: hyperbolic, t

Essential Information Theory

The Notion of Entropy

- Entropy \sim “chaos”, fuzziness, opposite of order, ...
 - you know it:
 - it is much easier to create “mess” than to tidy things up...
- Comes from physics:
 - Entropy does not go down unless energy is applied
- Measure of uncertainty:
 - if low... low uncertainty; the higher the entropy, the higher uncertainty, but the higher “surprise” (information) we can get out of an experiment

The Formula

- Let $p_X(x)$ be a distribution of random variable X
- Basic outcomes (alphabet) Ω

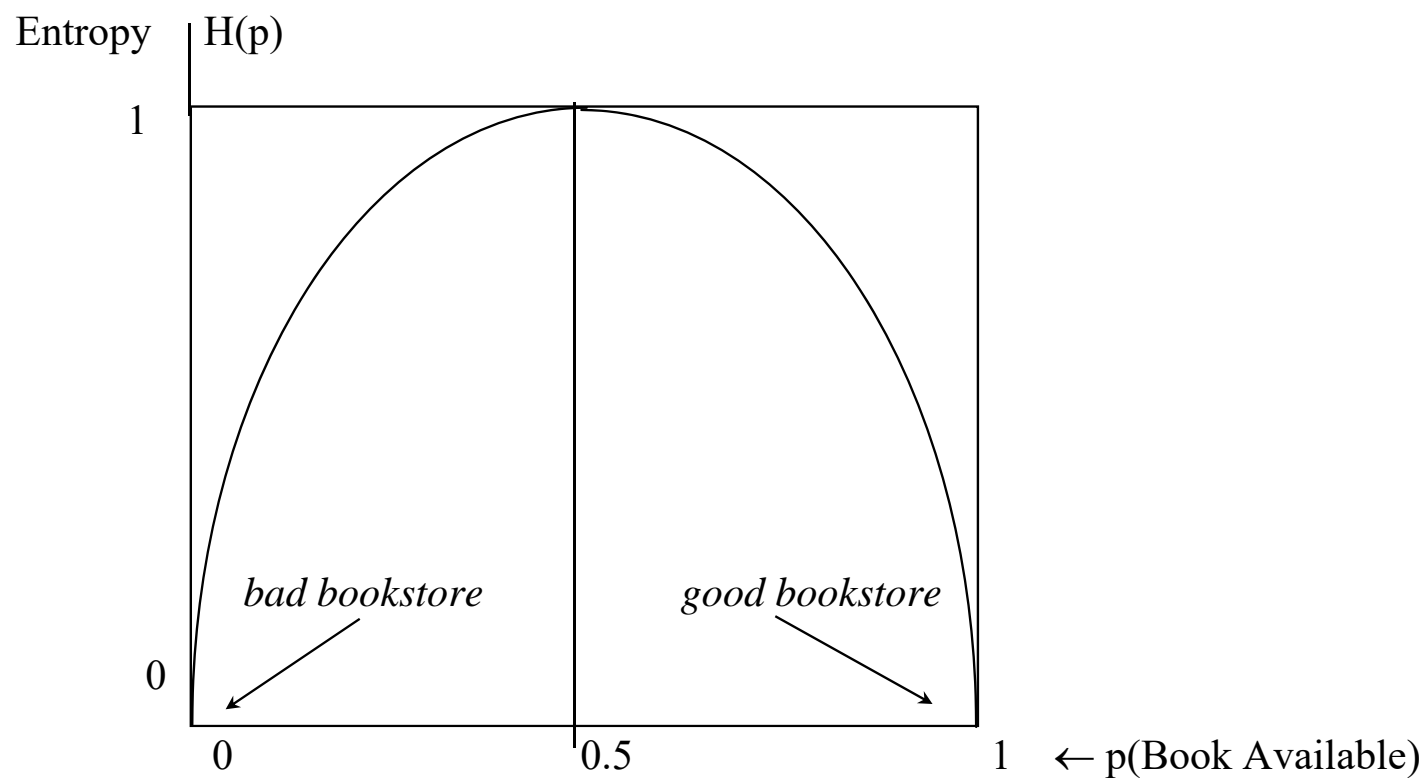
$$H(X) = - \sum_{x \in \Omega} p(x) \log_2 p(x) \quad !$$

- Unit: bits (\log_{10} : nats)
- Notation: $H(X) = H_p(X) = H(p) = H_X(p) = H(p_X)$

Using the Formula: Example

- Toss a fair coin: $\Omega = \{\text{head}, \text{tail}\}$
 - $p(\text{head}) = .5, p(\text{tail}) = .5$
 - $\mathbf{H}(\mathbf{p}) = -0.5 \log_2(0.5) + (-0.5 \log_2(0.5)) = 2 \times ((-0.5) \times (-1)) = 2 \times 0.5 = 1$
- Take fair, 32-sided die: $p(x) = 1 / 32$ for every side x
 - $\mathbf{H}(\mathbf{p}) = -\sum_{i=1..32} p(x_i) \log_2 p(x_i) = -32 (p(x_1) \log_2 p(x_1))$
(since for all i $p(x_i) = p(x_1) = 1/32$)
 $= -32 \times ((1/32) \times (-5)) = \mathbf{5}$ (*now you see why it's called **bits**?*)
- Unfair coin:
 - $p(\text{head}) = .2 \dots \mathbf{H}(\mathbf{p}) = .722$; $p(\text{head}) = .01 \dots \mathbf{H}(\mathbf{p}) = .081$

Example: Book Availability



The Limits

- When $H(p) = 0$?
 - if a result of an experiment is *known* ahead of time:
 - necessarily:
$$\exists x \in \Omega; p(x) = 1 \ \& \ \forall y \in \Omega; y \neq x \Rightarrow p(y) = 0$$
- Upper bound?
 - none in general
 - for $|\Omega| = n$: $H(p) \leq \log_2 n$
 - **nothing can be more uncertain than the uniform distribution**

Entropy and Expectation

- Recall:

$$- E(X) = \sum_{x \in X(\Omega)} p_X(x) \times x$$

- Then:

$$E(\log_2(1/p_X(x))) = \sum_{x \in X(\Omega)} p_X(x) \log_2(1/p_X(x)) =$$

$$= - \sum_{x \in X(\Omega)} p_X(x) \log_2 p_X(x) =$$

$$= H(p_X) =_{\text{notation}} H(p)$$

Perplexity: motivation

- Recall:
 - 2 equiprobable outcomes: $H(p) = 1$ bit
 - 32 equiprobable outcomes: $H(p) = 5$ bits
 - 4.3 billion equiprobable outcomes: $H(p) \approx 32$ bits
- What if the outcomes are not equiprobable?
 - 32 outcomes, 2 equiprobable at .5, rest impossible:
 - **$H(p) = 1$ bit**
 - Any measure for comparing the entropy (i.e. uncertainty/difficulty of prediction) (also) for random variables with different number of outcomes?

Perplexity

- Perplexity:
 - $G(p) = 2^{H(p)}$
- ... so we are back at 32 (for 32 eqp. outcomes), 2 for fair coins, etc.
- it is easier to imagine:
 - NLP example: vocabulary size of a vocabulary with uniform distribution, which is equally hard to predict
- the “wilder” (biased) distribution, the better:
 - lower entropy, lower perplexity

Joint Entropy and Conditional Entropy

- Two random variables: X (space Ω), Y (Ψ)
- Joint entropy:
 - no big deal: $((X,Y)$ considered a single event):

$$H(X,Y) = - \sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 p(x,y)$$

- Conditional entropy:

$$H(Y|X) = - \sum_{x \in \Omega} \sum_{y \in \Psi} \underline{p(x,y)} \log_2 p(y|x)$$

recall that $H(X) = E(\log_2(1/p_X(x)))$

(weighted “average”, and weights are not conditional)

Conditional Entropy (Using the Calculus)

- other definition:

$$H(Y|X) = \sum_{x \in \Omega} p(x) H(Y|X=x) =$$

for $H(Y|X=x)$, we can use the
variable definition ($x \sim \text{constant}$)

single-

$$= \sum_{x \in \Omega} p(x) \left(- \sum_{y \in \Psi} p(y|x) \log_2 p(y|x) \right) =$$

$$= - \sum_{x \in \Omega} \sum_{y \in \Psi} p(y|x) p(x) \log_2 p(y|x) =$$

$$= - \sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 p(y|x)$$

Properties of Entropy I

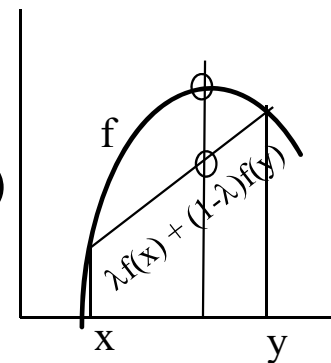
- Entropy is non-negative:
 - $H(X) \geq 0$
 - proof: (recall: $H(X) = - \sum_{x \in \Omega} p(x) \log_2 p(x)$)
 - $\log(p(x))$ is negative or zero for $x \leq 1$,
 - $p(x)$ is non-negative; their product $p(x)\log(p(x))$ is thus negative;
 - sum of negative numbers is negative;
 - and $-f$ is positive for negative f
- Chain rule:
 - $H(X,Y) = H(Y|X) + H(X)$, as well as
 - $H(X,Y) = H(X|Y) + H(Y)$ (since $H(Y,X) = H(X,Y)$)

Properties of Entropy II

- Conditional Entropy is better (than unconditional):
 - $H(Y|X) \leq H(Y)$
- $H(X,Y) \leq H(X) + H(Y)$ (follows from the previous (in)equalities)
 - **equality iff X,Y independent**
 - **[recall: X,Y independent iff $p(X,Y) = p(X)p(Y)$]**
- $H(p)$ is concave (remember the book availability graph?)
 - concave function \underline{f} over an interval (a,b) :

$$\forall x,y \in (a,b), \forall \lambda \in [0,1]:$$

$$f(\lambda x + (1-\lambda)y) \geq \lambda f(x) + (1-\lambda)f(y)$$
 - **function \underline{f} is convex if $-\underline{f}$ is concave**
- [for proofs and generalizations, see Cover/Thomas]



“Coding” Interpretation of Entropy

- The least (average) number of bits needed to encode a message (string, sequence, series,...) (each element having being a result of a random process with some distribution p): $= H(p)$
- Remember various compressing algorithms?
 - they do well on data with repeating (= easily predictable = low entropy) patterns
 - their results though have high entropy \Rightarrow compressing compressed data does nothing

Coding: Example

- How many bits do we need for ISO Latin 1?
 - \Rightarrow the trivial answer: 8
- Experience: some chars are more common, some (very) rare:
 - ...so what if we use more bits for the rare, and less bits for the frequent? [be careful: want to decode (easily)!]
 - suppose: $p('a') = 0.3$, $p('b') = 0.3$, $p('c') = 0.3$, the rest: $p(x) \cong .0004$
 - code: 'a' ~ 00 , 'b' ~ 01 , 'c' ~ 10 , rest: $11b_1b_2b_3b_4b_5b_6b_7b_8$
 - code acbbécbaac: $0010010111000011111001000010$

$$\begin{array}{cccccccccccccccc} a & c & b & b & & & \acute{e} & & & c & b & a & a & c \end{array}$$
 - number of bits used: 28 (vs. 80 using “naive” coding)
- code length $\sim 1 / \text{probability}$; conditional prob OK!

Entropy of a Language

- Imagine that we produce the next letter using

$$p(l_{n+1}|l_1, \dots, l_n),$$

where l_1, \dots, l_n is the sequence of *all* the letters which had been uttered so far (i.e. n is really big!); let's call l_1, \dots, l_n the *history* h (h_{n+1}), and all histories H :

- Then compute its entropy:
 - $\sum_{h \in H} \sum_{l \in A} p(l, h) \log_2 p(l|h)$
- Not very practical, isn't it?

Kullback-Leibler Distance (Relative Entropy)

- Remember:
 - long series of experiments... c_i/T_i oscillates around some number... we can only estimate it... to get a distribution \underline{q} .
- So we get a distribution \underline{q} ; (sample space Ω , r.v. X)
the true distribution is, however, \underline{p} . (same Ω , X)
 \Rightarrow how big error are we making?
- $D(\underline{p}||\underline{q})$ (the Kullback-Leibler distance):

$$D(\underline{p}||\underline{q}) = \sum_{x \in \Omega} \underline{p}(x) \log_2 (p(x)/q(x)) = E_{\underline{p}} \log_2 (p(x)/q(x))$$

Comments on Relative Entropy

- Conventions:
 - $0 \log 0 = 0$
 - $p \log (p/0) = \infty$ (for $p > 0$)
- Distance? (less “misleading”: Divergence)
 - not quite:
 - **not symmetric: $D(p||q) \neq D(q||p)$**
 - **does not satisfy the triangle inequality**
 - but useful to look at it that way
- $H(p) + D(p||q)$: bits needed for encoding p if q is used

Mutual Information (MI)

in terms of relative entropy

- Random variables X, Y ; $p_{X \cap Y}(x,y)$, $p_X(x)$, $p_Y(y)$
- Mutual information (between two random variables X, Y):

$$I(X, Y) = D(p(x, y) \parallel p(x)p(y))$$

- $I(X, Y)$ measures how much (our knowledge of) Y contributes (on average) to easing the prediction of X
- or, how $p(x, y)$ deviates from (independent) $p(x)p(y)$

Mutual Information: the Formula

- Rewrite the definition: [recall: $D(r||s) = \sum_{v \in \Omega} r(v) \log_2 (r(v)/s(v))$;
substitute $r(v) = p(x,y)$, $s(v) = p(x)p(y)$; $\langle v \rangle \sim \langle x,y \rangle$]

$$\begin{aligned} I(X,Y) &= D(p(x,y) || p(x)p(y)) = \\ &= \sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 (p(x,y)/p(x)p(y)) \end{aligned} \quad !$$

- Measured in bits (what else? :-)

From Mutual Information to Entropy

- by how many bits the knowledge of Y lowers the entropy $H(X)$:

$$\begin{aligned}
 I(X,Y) &= \sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 \left(\frac{p(x,y)}{p(y)p(x)} \right) = \\
 &\quad \dots \text{use } p(x,y)/p(y) = p(x|y) \\
 &= \sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 \left(\frac{p(x|y)}{p(x)} \right) = \\
 &\quad \dots \text{use } \log(a/b) = \log a - \log b \text{ (} a \sim p(x|y), b \sim p(x) \text{), distribute sums} \\
 &= \underbrace{\sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 p(x|y)}_{\dots \text{use def. of } H(X|Y) \text{ (left term)}} - \underbrace{\sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 p(x)}_{\dots \text{use } \sum_{y \in \Psi} p(x,y) = p(x) \text{ (right term)}} = \\
 &= -H(X|Y) + \left(- \sum_{x \in \Omega} p(x) \log_2 p(x) \right) = \\
 &\quad \dots \text{use def. of } H(X) \text{ (right term), swap terms} \\
 &= H(X) - H(X|Y) \quad \dots \text{by symmetry, } = H(Y) - H(Y|X)
 \end{aligned}$$

Properties of MI vs. Entropy

- $I(X,Y) = H(X) - \underline{H(X|Y)}$ = number of bits the knowledge of Y lowers the entropy of X
 $= H(Y) - H(Y|X)$ (prev. foil, symmetry)

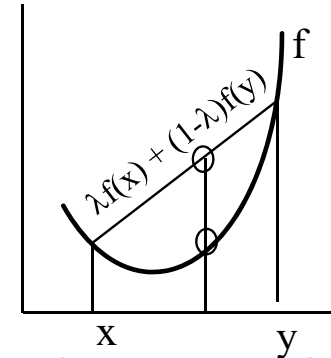
Recall: $H(X,Y) = H(X|Y) + H(Y) \Rightarrow -H(X|Y) = H(Y) - H(X,Y) \Rightarrow$

- $I(X,Y) = H(X) + \underline{H(Y) - H(X,Y)}$
- $I(X,X) = H(X)$ (since $H(X|X) = 0$)
- $I(X,Y) = I(Y,X)$ (just for completeness)
- $I(X,Y) \geq 0$... let's prove that now (as promised).

Jensen's Inequality

- Recall: f is convex on interval (a,b) iff $\forall x,y \in (a,b), \forall \lambda \in [0,1]$:

$$f(\lambda x + (1-\lambda)y) \leq \lambda f(x) + (1-\lambda)f(y)$$



- J.I.: for distribution $p(x)$, r.v. X on Ω , and convex f ,

$$f\left(\sum_{x \in \Omega} p(x) x\right) \leq \sum_{x \in \Omega} p(x) f(x)$$

- Proof (idea): by induction on the number of basic outcomes;
- start with $|\Omega| = 2$ by:
 - $p(x_1)f(x_1) + p(x_2)f(x_2) \geq f(p(x_1)x_1 + p(x_2)x_2)$ (\Leftarrow def. of convexity)
 - for the induction step ($|\Omega| = k \rightarrow k+1$), just use the induction hypothesis and def. of convexity (again).

Information Inequality

$$D(p||q) \geq 0 \quad !$$

- Proof:

$$\begin{aligned} 0 &= -\log 1 = -\log \sum_{x \in \Omega} q(x) = -\log \sum_{x \in \Omega} (q(x)/p(x))p(x) \leq \\ &\quad \dots \text{apply Jensen's inequality here (} \underline{-\log} \text{ is convex)} \dots \\ &\leq \sum_{x \in \Omega} p(x) (-\log(q(x)/p(x))) = \sum_{x \in \Omega} p(x) \log(p(x)/q(x)) = \\ &\quad = D(p||q) \end{aligned}$$

Other (In)Equalities and Facts

- Log sum inequality: for $r_i, s_i \geq 0$

$$\sum_{i=1..n} (r_i \log(r_i/s_i)) \leq (\sum_{i=1..n} r_i) \log(\sum_{i=1..n} r_i / \sum_{i=1..n} s_i)$$

- $D(p||q)$ is convex [in p, q] (\Leftarrow log sum inequality)
- $H(p_X) \leq \log_2 |\Omega|$, where Ω is the sample space of p_X

Proof: uniform $u(x)$, same sample space Ω : $\sum p(x) \log u(x) = -\log_2 |\Omega|$;

$$\log_2 |\Omega| - H(X) = -\sum p(x) \log u(x) + \sum p(x) \log p(x) = D(p||u) \geq 0$$

- $H(p)$ is concave [in p]:

Proof: from $H(X) = \log_2 |\Omega| - D(p||u)$, $D(p||u)$ convex $\Rightarrow H(x)$ concave

Cross-Entropy

- Typical case: we've got series of observations
 $T = \{t_1, t_2, t_3, t_4, \dots, t_n\}$ (numbers, words, ...; $t_i \in \Omega$);
estimate (simple):
 $\forall y \in \Omega: \tilde{p}(y) = c(y) / |T|$, def. $c(y) = |\{t \in T; t = y\}|$
- ...but the true p is unknown; every sample is too small!
- Natural question: how well do we do using \tilde{p} [instead of p]?
- Idea: simulate actual p by using a different T'
(or rather: by using different observation we simulate the insufficiency of T vs. some other data (“random” difference))

Cross Entropy: The Formula

- $H_{p'}(\tilde{p}) = H(p') + D(p' \| \tilde{p})$

$$H_{p'}(\tilde{p}) = - \sum_{x \in \Omega} p'(x) \log_2 \tilde{p}(x) \quad \bullet$$

- p' is certainly not the true p , but we can consider it the “real world” distribution against which we test \tilde{p}
- note on notation (confusing...): $p/p' \leftrightarrow \tilde{p}$, also $H_T(p)$
- (Cross)Perplexity: $G_{p'}(p) = G_T(p) = 2^{H_{p'}(\tilde{p})}$

Conditional Cross Entropy

- So far: “unconditional” distribution(s) $p(x)$, $p'(x)$...
- In practice: virtually always conditioning on context
- Interested in: sample space Ψ , r.v. Y , $y \in \Psi$;
context: sample space Ω , r.v. X , $x \in \Omega$;;
“our” distribution $p(y|x)$, test against $p'(y,x)$,
which is taken from some independent data:

$$H_{p'}(p) = - \sum_{y \in \Psi, x \in \Omega} p'(y,x) \log_2 p(y|x)$$

Sample Space vs. Data

- In practice, it is often inconvenient to sum over the sample space(s) Ψ, Ω (especially for cross entropy!)
- Use the following formula:

$$H_p(p) = - \sum_{y \in \Psi, x \in \Omega} p'(y,x) \log_2 p(y|x) = - 1/|T'| \sum_{i=1..|T'|} \log_2 p(y_i|x_i) \quad !$$

- This is in fact the normalized log probability of the “test” data:

$$H_p(p) = - 1/|T'| \log_2 \prod_{i=1..|T'|} p(y_i|x_i)$$

Computation Example

- $\Omega = \{a, b, \dots, z\}$, prob. distribution (assumed/estimated from data):
 $p(a) = .25, p(b) = .5, p(\alpha) = 1/64$ for $\alpha \in \{c..r\}$, $= 0$ for the rest: s,t,u,v,w,x,y,z
- Data (test): barb $p'(a) = p'(r) = .25, p'(b) = .5$
- Sum over Ω :

α	a	b	c	d	e	f	g	...	p	q	r	s	t	...	z
$-p'(\alpha)\log_2 p(\alpha)$.5	+ .5	+ 0	+ 0	+ 0	+ 0	+ 0	+ 0	+ 0	+ 0	+ 1.5	+ 0	+ 0	+ 0	+ 0 = <u>2.5</u>

- Sum over data:

i / s _i	1/b	2/a	3/r	4/b		1/ T'			
-log ₂ p(s _i)	1	+	2	+	6	+	1	= 10	(1/4) × 10 = 2.5

Cross Entropy: Some Observations

- $H(p)$?? $<, =, >$?? $H_{p'}(p)$: ALL!
- Previous example:
 $[p(a) = .25, p(b) = .5, p(\alpha) = 1/64 \text{ for } \alpha \in \{c..r\}, = 0 \text{ for the rest: } s,t,u,v,w,x,y,z]$

$$H(p) = 2.5 \text{ bits} = H(p') \text{ (barb)}$$
- Other data: probable: $(1/8)(6+6+6+1+2+1+6+6) = 4.25$

$$H(p) < 4.25 \text{ bits} = H(p') \text{ (probable)}$$
- And finally: abba: $(1/4)(2+1+1+2) = 1.5$

$$H(p) > 1.5 \text{ bits} = H(p') \text{ (abba)}$$
- But what about: baby $-p'(\vec{y})\log_2 p(\vec{y}) = -.25\log_2 0 = \infty$ (??)

Cross Entropy: Usage

- Comparing data??
 - NO! (we believe that we test on real data!)
- Rather: comparing distributions (vs. real data)
- Have (got) 2 distributions: p and q (on some Ω , X)
 - which is better?
 - better: has lower cross-entropy (perplexity) on real data S
- “Real” data: S
- $H_S(p) = -1/|S| \sum_{i=1..|S|} \log_2 p(y_i|x_i)$?? $H_S(q) = -1/|S| \sum_{i=1..|S|} \log_2 q(y_i|x_i)$

Comparing Distributions

Test data S: probable

- $p(.)$ from prev. example:

$$H_S(p) = 4.25$$

$p(a) = .25, p(b) = .5, p(\alpha) = 1/64$ for $\alpha \in \{c..r\}, = 0$ for the rest: s,t,u,v,w,x,y,z

- $q(.|.)$ (conditional; defined by a table):

$q(. .)\rightarrow$ ↓	a	b	e	l	o	p	r	other
a	0	.5	0	0	0	.125	0	0
b	1	0	0	0	1	.125	0	0
e	0	0	0	1	0	.125	0	0
l	0	.5	0	0	0	.125	0	0
o	0	0	0	0	0	.125	1	0
p	0	0	0	0	0	.125	0	1
r	0	0	0	0	0	.125	0	0
other	0	0	1	0	0	.125	0	0

ex.: $q(o|r) = 1$

$q(r|p) = .125$

$(1/8) (\log(p|oth.) + \log(r|p) + \log(o|r) + \log(b|o) + \log(a|b) + \log(b|a) + \log(l|b) + \log(e|l))$

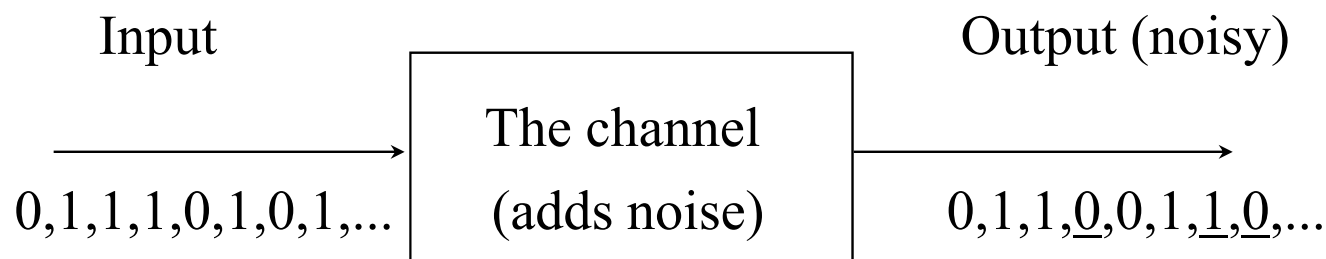
$(1/8) (0 + 3 + 0 + 0 + 1 + 0 + 1 + 0)$

$$H_S(q) = .625$$

Language Modeling (and the Noisy Channel)

The Noisy Channel

- Prototypical case:



- Model: probability of error (noise):
- Example: $p(0|1) = .3$ $p(1|1) = .7$ $p(1|0) = .4$ $p(0|0) = .6$
- The Task:
known: the noisy output; want to know: the input (decoding)

Noisy Channel Applications

- OCR
 - straightforward: text → print (adds noise), scan → image
- Handwriting recognition
 - text → neurons, muscles (“noise”), scan/digitize → image
- Speech recognition (dictation, commands, etc.)
 - text → conversion to acoustic signal (“noise”) → acoustic waves
- Machine Translation
 - text in target language → translation (“noise”) → source language
- Also: Part of Speech Tagging
 - sequence of tags → selection of word forms → text

Noisy Channel: The Golden Rule of ...

OCR, ASR, HR, MT, ...



- Recall:

$$p(A|B) = p(B|A) p(A) / p(B) \quad (\text{Bayes formula})$$

$$A_{\text{best}} = \operatorname{argmax}_A p(B|A) p(A) \quad (\text{The Golden Rule})$$

- $p(B|A)$: the acoustic/image/translation/lexical model
 - application-specific name
 - will explore later
- $p(A)$: *the language model*

The Perfect Language Model

- Sequence of word forms [forget about tagging for the moment]
- Notation: $A \sim W = (w_1, w_2, w_3, \dots, w_d)$
- The big (modeling) question:

$$p(W) = ?$$

- Well, we know (Bayes/chain rule \rightarrow):

$$\begin{aligned} p(W) &= p(w_1, w_2, w_3, \dots, w_d) = \\ &= p(w_1) \times p(w_2|w_1) \times p(w_3|w_1, w_2) \times \dots \times p(w_d|w_1, w_2, \dots, w_{d-1}) \end{aligned}$$

- Not practical (even short $W \rightarrow$ too many parameters)

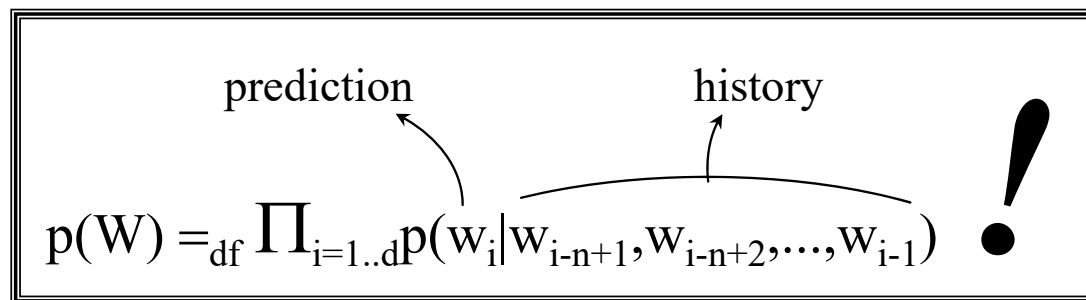
Markov Chain

- Unlimited memory (cf. previous foil):
 - for w_i , we know all its predecessors $w_1, w_2, w_3, \dots, w_{i-1}$
- Limited memory:
 - we disregard “too old” predecessors
 - remember only k previous words: $w_{i-k}, w_{i-k+1}, \dots, w_{i-1}$
 - called “ k^{th} order Markov approximation”
- + stationary character (no change over time):

$$p(W) \cong \prod_{i=1..d} p(w_i | w_{i-k}, w_{i-k+1}, \dots, w_{i-1}), d = |W|$$

n-gram Language Models

- $(n-1)^{\text{th}}$ order Markov approximation \rightarrow n-gram LM:



The diagram shows the equation $p(W) =_{\text{df}} \prod_{i=1..d} p(w_i | w_{i-n+1}, w_{i-n+2}, \dots, w_{i-1})$ enclosed in a double-bordered box. An arrow labeled "prediction" points from the word w_i to the left. An arrow labeled "history" points from the sequence $w_{i-n+1}, w_{i-n+2}, \dots, w_{i-1}$ to the left. A large exclamation mark is placed to the right of the equation.

$$p(W) =_{\text{df}} \prod_{i=1..d} p(w_i | w_{i-n+1}, w_{i-n+2}, \dots, w_{i-1})$$

- In particular (assume vocabulary $|V| = 60\text{k}$):
 - **0-gram LM: uniform model,** $p(w) = 1/|V|$, **1 parameter**
 - **1-gram LM: unigram model,** $p(w)$, **6×10^4 parameters**
 - **2-gram LM: bigram model,** $p(w_i | w_{i-1})$ **3.6×10^9 parameters**
 - **3-gram LM: trigram model,** $p(w_i | w_{i-2}, w_{i-1})$ **2.16×10^{14} parameters**

LM: Observations

- How large n ?
 - nothing is enough (theoretically)
 - but anyway: as much as possible (\rightarrow close to “perfect” model)
 - empirically: 3
 - **parameter estimation? (reliability, data availability, storage space, ...)**
 - **4 is too much: $|V|=60k \rightarrow 1.296 \times 10^{19}$ parameters**
 - **but: 6-7 would be (almost) ideal (having enough data): in fact, one can recover the original text ssequence from 7-grams!**
- Reliability $\sim (1 / \text{Detail})$ (\rightarrow need compromise)
- For now, keep word forms (no “linguistic” processing)

The Length Issue

- $\forall n; \sum_{w \in \Omega^n} p(w) = 1 \Rightarrow \sum_{n=1.. \infty} \sum_{w \in \Omega^n} p(w) >> 1 (\rightarrow \infty)$
- We want to model all sequences of words
 - for “fixed” length tasks: no problem - n fixed, sum is 1
 - **tagging, OCR/handwriting (if words identified ahead of time)**
 - for “variable” length tasks: have to account for
 - **discount shorter sentences**
- General model: for each sequence of words of length n,
define $p'(w) = \lambda_n p(w)$ such that $\sum_{n=1.. \infty} \lambda_n = 1 \Rightarrow$
$$\sum_{n=1.. \infty} \sum_{w \in \Omega^n} p'(w) = 1$$

e.g., estimate λ_n from data; or use normal or other distribution

Parameter Estimation

- Parameter: numerical value needed to compute $p(w|h)$
- From data (how else?)
- Data preparation:
 - **get rid of formatting etc. (“text cleaning”)**
 - **define words (separate but include punctuation, call it “word”)**
 - **define sentence boundaries (insert “words” $\langle s \rangle$ and $\langle /s \rangle$)**
 - **letter case: keep, discard, or be smart:**
 - name recognition
 - number type identification[these are huge problems per se!]
 - **numbers: keep, replace by $\langle \text{num} \rangle$, or be smart (form \sim pronunciation)**

Maximum Likelihood Estimate

- MLE: Relative Frequency...
 - ...best predicts the data at hand (the “training data”)
- Trigrams from Training Data T:
 - count sequences of three words in T: $c_3(w_{i-2}, w_{i-1}, w_i)$
[NB: notation: just saying that the three words follow each other]
 - count sequences of two words in T: $c_2(w_{i-1}, w_i)$:
 - **either use $c_2(y, z) = \sum_w c_3(y, z, w)$**
 - **or count differently at the beginning (& end) of data!**

$$p(w_i | w_{i-2}, w_{i-1}) =_{\text{est.}} c_3(w_{i-2}, w_{i-1}, w_i) / c_2(w_{i-2}, w_{i-1}) \quad !$$

Character Language Model

- Use individual characters instead of words:

$$p(W) =_{\text{df}} \prod_{i=1..d} p(c_i | c_{i-n+1}, c_{i-n+2}, \dots, c_{i-1})$$

- Same formulas etc.
- Might consider 4-grams, 5-grams or even more
- Good only for language comparison
- Transform cross-entropy between letter- and word-based models:

$$H_S(p_c) = H_S(p_w) / \text{avg. \# of characters/word in } S$$

LM: an Example

- Training data:

$\langle s \rangle \langle s \rangle$ He can buy the can of soda.

- Unigram: $p_1(\text{He}) = p_1(\text{buy}) = p_1(\text{the}) = p_1(\text{of}) = p_1(\text{soda}) = p_1(.) = .125$

$p_1(\text{can}) = .25$

- Bigram: $p_2(\text{He}|\langle s \rangle) = 1$, $p_2(\text{can}|\text{He}) = 1$, $p_2(\text{buy}|\text{can}) = .5$,

$p_2(\text{of}|\text{can}) = .5$, $p_2(\text{the}|\text{buy}) = 1, \dots$

- Trigram: $p_3(\text{He}|\langle s \rangle, \langle s \rangle) = 1$, $p_3(\text{can}|\langle s \rangle, \text{He}) = 1$,

$p_3(\text{buy}|\text{He}, \text{can}) = 1$, $p_3(\text{of}|\text{the}, \text{can}) = 1$, ..., $p_3(.|\text{of}, \text{soda}) = 1$.

- Entropy: $H(p_1) = 2.75$, $H(p_2) = .25$, $H(p_3) = 0$ ← Great?!

LM: an Example (The Problem)

- Cross-entropy:
- $S = \langle s \rangle \langle s \rangle$ It was the greatest buy of all.
- Even $H_S(p_1)$ fails ($= H_S(p_2) = H_S(p_3) = \infty$), because:
 - all unigrams but $p_1(\text{the})$, $p_1(\text{buy})$, $p_1(\text{of})$ and $p_1(.)$ are 0.
 - all bigram probabilities are 0.
 - all trigram probabilities are 0.
- We want: to make all (theoretically possible*) probabilities non-zero.

* in fact, all: remember our graph from day 1?

LM Smoothing (And the EM Algorithm)

The Zero Problem

- “Raw” n-gram language model estimate:
 - necessarily, some zeros
 - **!many: trigram model $\rightarrow 2.16 \times 10^{14}$ parameters, data $\sim 10^9$ words**
 - which are true 0?
 - **optimal situation: even the least frequent trigram would be seen several times, in order to distinguish it's probability vs. other trigrams**
 - **optimal situation cannot happen, unfortunately (open question: how many data would we need?)**
 - \rightarrow we don't know
 - we must eliminate the zeros
- Two kinds of zeros: $p(w|h) = 0$, or even $p(h) = 0$!

Why do we need Nonzero Probs?

- To avoid infinite Cross Entropy:
 - happens when an event is found in test data which has not been seen in training data
 - $H(p) = \infty$: prevents comparing data with > 0 “errors”
- To make the system more robust
 - low count estimates:
 - **they typically happen for “detailed” but relatively rare appearances**
 - high count estimates: reliable but less “detailed”

Eliminating the Zero Probabilities: Smoothing

- Get new $p'(w)$ (same Ω): almost $p(w)$ but no zeros
- Discount w for (some) $p(w) > 0$: new $p'(w) < p(w)$

$$\sum_{w \in \text{discounted}} (p(w) - p'(w)) = D$$

- Distribute D to all w ; $p(w) = 0$: new $p'(w) > p(w)$
 - possibly also to other w with low $p(w)$
- For some w (possibly): $p'(w) = p(w)$
- Make sure $\sum_{w \in \Omega} p'(w) = 1$
- There are many ways of *smoothing*

Smoothing by Adding 1

- Simplest but not really usable:
 - Predicting words w from a vocabulary V , training data T :
$$p'(w|h) = (c(h,w) + 1) / (c(h) + |V|)$$
 - **for non-conditional distributions: $p'(w) = (c(w) + 1) / (|T| + |V|)$**
 - Problem if $|V| > c(h)$ (as is often the case; even $\gg c(h)$!)
- Example: Training data: $\langle s \rangle$ what is it what is small ? $|T| = 8$
 - $V = \{ \text{what, is, it, small, ?, } \langle s \rangle, \text{ flying, birds, are, a, bird, .} \}, |V| = 12$
 - $p(\text{it})=.125, p(\text{what})=.25, p(.)=0$ $p(\text{what is it?}) = .25^2 \times .125^2 \cong .001$
$$p(\text{it is flying.}) = .125 \times .25 \times 0^2 = 0$$
 - $p'(\text{it})=.1, p'(\text{what})=.15, p'(.)=.05$ $p'(\text{what is it?}) = .15^2 \times .1^2 \cong .0002$
$$p'(\text{it is flying.}) = .1 \times .15 \times .05^2 \cong .00004$$

Adding *less than* 1

- Equally simple:
 - Predicting words w from a vocabulary V , training data T :
$$p'(w|h) = (c(h,w) + \lambda) / (c(h) + \lambda|V|), \lambda < 1$$
 - **for non-conditional distributions:** $p'(w) = (c(w) + \lambda) / (|T| + \lambda|V|)$
- **Example:** Training data: $\langle s \rangle$ what is it what is small ? $|T| = 8$
 - $V = \{ \text{what, is, it, small, ?, } \langle s \rangle, \text{flying, birds, are, a, bird, .} \}, |V| = 12$
 - $p(\text{it})=.125, p(\text{what})=.25, p(.)=0$ $p(\text{what is it?}) = .25^2 \times .125^2 \cong .001$
$$p(\text{it is flying.}) = .125 \times .25 \times 0^2 = 0$$
 - Use $\lambda = .1$:
 - $p'(\text{it}) \cong .12, p'(\text{what}) \cong .23, p'(\text{.}) \cong .01$ $p'(\text{what is it?}) = .23^2 \times .12^2 \cong .0007$
$$p'(\text{it is flying.}) = .12 \times .23 \times .01^2 \cong .000003$$

Good - Turing

- Suitable for estimation from large data
 - similar idea: discount/boost the relative frequency estimate:
$$p_r(w) = (c(w) + 1) \times N(c(w) + 1) / (|T| \times N(c(w))) ,$$

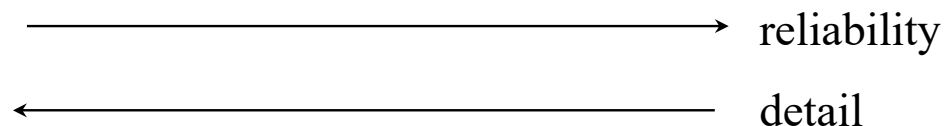
where $N(c)$ is the count of words with count c (count-of-counts)
specifically, for $c(w) = 0$ (unseen words), $p_r(w) = N(1) / (|T| \times N(0))$
 - good for small counts ($< 5-10$, where $N(c)$ is high)
 - variants (see MS)
 - normalization! (so that we have $\sum_w p'(w) = 1$)

Good-Turing: An Example

- **Example:** remember: $p_r(w) = (c(w) + 1) \times N(c(w) + 1) / (|T| \times N(c(w)))$
 Training data: $\langle s \rangle$ what is it what is small ? $|T| = 8$
 - $V = \{ \text{what, is, it, small, ?, } \langle s \rangle, \text{flying, birds, are, a, bird, .} \}, |V| = 12$
 $p(\text{it}) = .125, p(\text{what}) = .25, p(.) = 0$ $p(\text{what is it?}) = .25^2 \times .125^2 \cong .001$
 $p(\text{it is flying.}) = .125 \times .25 \times 0^2 = 0$
 - **Raw reestimation** ($N(0) = 6, N(1) = 4, N(2) = 2, N(i) = 0$ for $i > 2$):
 $p_r(\text{it}) = (1+1) \times N(1+1) / (8 \times N(1)) = 2 \times 2 / (8 \times 4) = .125$
 $p_r(\text{what}) = (2+1) \times N(2+1) / (8 \times N(2)) = 3 \times 0 / (8 \times 2) = 0$: keep orig. $p(\text{what})$
 $p_r(.) = (0+1) \times N(0+1) / (8 \times N(0)) = 1 \times 4 / (8 \times 6) \cong .083$
 - **Normalize** (divide by $1.5 = \sum_{w \in |V|} p_r(w)$) and compute:
 $p'(\text{it}) \cong .08, p'(\text{what}) \cong .17, p'(.) \cong .06$ $p'(\text{what is it?}) = .17^2 \times .08^2 \cong .0002$
 $p'(\text{it is flying.}) = .08 \times .17 \times .06^2 \cong .00004$

Smoothing by Combination: Linear Interpolation

- Combine what?
 - **distributions of various level of detail vs. reliability**
- n-gram models:
 - **use (n-1)gram, (n-2)gram, ..., uniform**



- Simplest possible combination:
 - sum of probabilities, normalize:
 - $p(0|0) = .8, p(1|0) = .2, p(0|1) = 1, p(1|1) = 0, p(0) = .4, p(1) = .6$:
 - $p'(0|0) = .6, p'(1|0) = .4, p'(0|1) = .7, p'(1|1) = .3$

Typical n-gram LM Smoothing

- Weight in less detailed distributions using $\lambda=(\lambda_0,\lambda_1,\lambda_2,\lambda_3)$:

$$p'_{\lambda}(w_i | w_{i-2}, w_{i-1}) = \lambda_3 p_3(w_i | w_{i-2}, w_{i-1}) + \\ \lambda_2 p_2(w_i | w_{i-1}) + \lambda_1 p_1(w_i) + \lambda_0 / |V|$$

- Normalize:

$$\lambda_i > 0, \sum_{i=0..n} \lambda_i = 1 \text{ is sufficient } (\lambda_0 = 1 - \sum_{i=1..n} \lambda_i) \text{ (n=3)}$$

- Estimation using MLE:

- fix the p_3, p_2, p_1 and $|V|$ parameters as estimated from the training data
- then find such $\{\lambda_i\}$ which minimizes the cross entropy
(maximizes probability of data): $-(1/|D|) \sum_{i=1..|D|} \log_2(p'_{\lambda}(w_i | h_i))$

Held-out Data

- What data to use?
 - try the training data T : but we will always get $\lambda_3 = 1$
 - **why?** (let p_{iT} be an **i-gram distribution estimated using r.f. from T**)
 - **minimizing $H_T(p'_\lambda)$ over a vector λ , $p'_\lambda = \lambda_3 p_{3T} + \lambda_2 p_{2T} + \lambda_1 p_{1T} + \lambda_0 / |V|$**
 - remember: $H_T(p'_\lambda) = H(p_{3T}) + D(p_{3T} \| p'_\lambda)$; p_{3T} fixed $\rightarrow H(p_{3T})$ fixed, and it is the best
 - which p'_λ minimizes $H_T(p'_\lambda)$? ... a p'_λ for which $D(p_{3T} \| p'_\lambda) = 0$
 - ...and that's p_{3T} (because $D(p \| p) = 0$, as we know).
 - ...and certainly $p'_\lambda = p_{3T}$ if $\lambda_3 = 1$ (maybe in some other cases, too).
 - $(p'_\lambda = 1 \times p_{3T} + 0 \times p_{2T} + 0 \times p_{1T} + 0 / |V|)$
 - thus: do not use the training data for estimation of λ !
 - **hold out part of the training data (*heldout \underline{H}*); remaining data: the true/raw training data, \underline{T}**
 - **the *test* data \underline{S} (e.g., for comparison purposes): still must be some different data!**

The Formulas

- Repeat: minimizing $-(1/|H|)\sum_{i=1..|H|}\log_2(p'_\lambda(w_i|h_i))$ over λ

$$p'_\lambda(w_i|h_i) = p'_\lambda(w_i|w_{i-2},w_{i-1}) = \lambda_3 p_3(w_i|w_{i-2},w_{i-1}) + \lambda_2 p_2(w_i|w_{i-1}) + \lambda_1 p_1(w_i) + \lambda_0/|V| \quad !$$

- “Expected Counts (of lambdas)”: $j = 0..3$

$$c(\lambda_j) = \sum_{i=1..|H|} (\lambda_j p_j(w_i|h_i) / p'_\lambda(w_i|h_i)) \quad !$$

- “Next λ ”: $j = 0..3$

$$\lambda_{j,next} = c(\lambda_j) / \sum_{k=0..3} (c(\lambda_k)) \quad !$$

The (Smoothing) EM Algorithm

1. Start with some λ , such that $\lambda_j > 0$ for all $j \in 0..3$.
 2. Compute “Expected Counts” for each λ_j .
 3. Compute new set of λ_j , using the “Next λ ” formula.
 4. Start over at step 2, unless a termination condition is met.
- Termination condition: convergence of λ .
 - Simply set an ε , and finish if $|\lambda_j - \lambda_{j,\text{next}}| < \varepsilon$ for each j (step 3).
 - Guaranteed to converge:
 - follows from Jensen’s inequality, plus a technical proof.

Remark on Linear Interpolation Smoothing

- “Bucketed” smoothing:
 - use several vectors of λ instead of one, based on (the frequency of) history: $\lambda(h)$
 - e.g. for $h = (\text{micrograms, per})$ we will have
$$\lambda(h) = (.999, .0009, .00009, .00001)$$
(because “cubic” is the only word to follow...)
 - actually: not a separate set for each history, but rather a set for “similar” histories (“bucket”):
$$\lambda(b(h)), \text{ where } b: V^2 \rightarrow N \text{ (in the case of trigrams)}$$
$$\underline{b} \text{ classifies histories according to their reliability } (\sim \text{frequency})$$

Bucketed Smoothing: The Algorithm

- First, determine the bucketing function \underline{b} (use heldout!):
 - decide in advance you want e.g. 1000 buckets
 - compute the total frequency of histories in 1 bucket ($f_{\max}(b)$)
 - gradually fill your buckets from the most frequent bigrams so that the sum of frequencies does not exceed $f_{\max}(b)$ (you might end up with slightly more than 1000 buckets)
- Divide your heldout data according to buckets
- Apply the previous algorithm to each bucket and its data

Simple Example

- Raw distribution (unigram only; smooth with uniform):

$$p(a) = .25, p(b) = .5, p(\alpha) = 1/64 \text{ for } \alpha \in \{c..r\}, = 0 \text{ for the rest: s,t,u,v,w,x,y,z}$$

- Heldout data: baby; use one set of λ (λ_1 : unigram, λ_0 : uniform)

- Start with $\lambda_1 = .5$; $p'_\lambda(b) = .5 \times .5 + .5 / 26 = .27$

$$p'_\lambda(a) = .5 \times .25 + .5 / 26 = .14$$

$$p'_\lambda(y) = .5 \times 0 + .5 / 26 = .02$$

$$c(\lambda_1) = .5 \times .5 / .27 + .5 \times .25 / .14 + .5 \times .5 / .27 + .5 \times 0 / .02 = 2.72$$

$$c(\lambda_0) = .5 \times .04 / .27 + .5 \times .04 / .14 + .5 \times .04 / .27 + .5 \times .04 / .02 = 1.28$$

$$\text{Normalize: } \lambda_{1,\text{next}} = .68, \lambda_{0,\text{next}} = .32.$$

Repeat from step 2 (recompute p'_λ first for efficient computation, then $c(\lambda_i)$, ...)

Finish when new lambdas almost equal to the old ones (say, < 0.01 difference).

Some More Technical Hints

- Set $V = \{\text{all words from training data}\}$.
 - **You may also consider $V = T \cup H$, but it does not make the coding in any way simpler (in fact, harder).**
 - **But: you must *never* use the test data for you vocabulary!**
- Prepend two “words” in front of all data:
 - **avoids beginning-of-data problems**
 - **call these index -1 and 0: then the formulas hold exactly**
- When $c_n(w, h) = 0$:
 - **Assign 0 probability to $p_n(w|h)$ where $c_{n-1}(h) > 0$, but a uniform probability ($1/|V|$) to those $p_n(w|h)$ where $c_{n-1}(h) = 0$ [this must be done both when working on the heldout data during EM, as well as when computing cross-entropy on the test data!]**

Words and the Company They Keep

Motivation

- Environment:
 - mostly “not a full analysis (sentence/text parsing)”
- Tasks where “words & company” are important:
 - word sense disambiguation (MT, IR, TD, IE)
 - lexical entries: subdivision & definitions (lexicography)
 - language modeling (generalization, [kind of] smoothing)
 - word/phrase/term translation (MT, Multilingual IR)
 - NL generation (“natural” phrases) (Generation, MT)
 - parsing (lexically-based selectional preferences)

Collocations

- Collocation
 - Firth: “word is characterized by the company it keeps”; collocations of a given word are statements of the habitual or customary places of that word.
 - non-compositionality of meaning
 - cannot be derived directly from its parts (heavy rain)
 - non-substitutability in context
 - for parts (red light)
 - non-modifiability (& non-transformability)
 - kick the yellow bucket; take ~~exceptions~~ to

X

Association and Co-occurrence; Terms

- Does not fall under “collocation”, but:
- Interesting just because it does often [rarely] appear together or in the same (or similar) context:
 - (doctors, nurses)
 - (hardware, software)
 - (gas, fuel)
 - (hammer, nail)
 - (communism, free speech)
- Terms:
 - need not be > 1 word (notebook, washer)

Collocations of Special Interest

- Idioms: really fixed phrases
 - **kick the bucket, birds-of-a-feather, run for office**
- Proper names: difficult to recognize even with lists
 - **Tuesday (person's name), May, Winston Churchill, IBM, Inc.**
- Numerical expressions
 - containing “ordinary” words
 - **Monday Oct 04 1999, two thousand seven hundred fifty**
- Phrasal verbs
 - Separable parts:
 - **look up, take off**

Further Notions

- Synonymy: different form/word, same meaning:
 - **notebook / laptop**
- Antonymy: opposite meaning:
 - **new/old, black/white, start/stop**
- Homonymy: same form/word, different meaning:
 - **“true” (random, unrelated): can (aux. verb / can of Coke)**
 - **related: polysemy; notebook, shift, grade, ...**
- Other:
 - **Hyperonymy/Hyponymy: general vs. special: vehicle/car**
 - **Meronymy/Holonymy: whole vs. part: body/leg**

How to Find Collocations?

- Frequency
 - plain
 - filtered
- Hypothesis testing
 - t test
 - χ^2 test
- Pointwise (“poor man’s”) Mutual Information
- (Average) Mutual Information

Frequency

- Simple
 - Count n-grams; high frequency n-grams are candidates:
 - **mostly function words**
 - **frequent names**
- Filtered
 - Stop list: words/forms which (we think) cannot be a part of a collocation
 - **a, the, and, or, but, not, ...**
 - Part of Speech (possible collocation patterns)
 - **A+N, N+N, N+of+N, ...**

Hypothesis Testing

- Hypothesis
 - something we test (against)
- Most often:
 - compare possibly interesting thing vs. “random” chance
 - “Null hypothesis”:
 - **something occurs by chance (that’s what we suppose).**
 - **Assuming this, prove that the probability of the “real world” is then too low (typically < 0.05 , also 0.005, 0.001)... therefore reject the null hypothesis (thus confirming “interesting” things are happening!)**
 - **Otherwise, it’s possible there is nothing interesting.**

t test (Student's t test)

- Significance of difference
 - compute “magic” number against normal distribution (mean μ)
 - using real-world data: (\bar{x} real data mean, s^2 variance, N size):
 - $t = (\bar{x} - \mu) / \sqrt{s^2 / N}$
 - find in tables (see MS, p. 609):
 - **d.f. = degrees of freedom (parameters which are not determined by other parameters)**
 - **percentile level $p = 0.05$ (or better)**
 - the bigger t :
 - **the better chances that there is the interesting feature we hope for (i.e. we can reject the null hypothesis)**
 - **t : at least the value from the table(s)**

t test on words

- null hypothesis: independence
 - mean μ : $p(w_1) p(w_2)$
- data estimates:
 - \mathbf{x}' = MLE of joint probability from data
 - s^2 is $p(1-p)$, i.e. almost p for small p ; N is the data size
- Example: (d.f. \sim sample size)
 - ‘general term’ (homework corpus): $c(\text{general}) = 108$, $c(\text{term}) = 40$
 - $c(\text{general}, \text{term}) = 2$; expected $p(\text{general})p(\text{term}) = 8.8\text{E-}8$
 - $t = (9.0\text{E-}6 - 8.8\text{E-}8) / (9.0\text{E-}6 / 221097)^{1/2} = 1.40$ (not > 2.576) thus ‘general term’ is not a collocation with confidence 0.005
 - ‘true species’: $(84/1779/9)$: $t = 2.774 > 2.576$!!

Pearson's Chi-square test

- χ^2 test (general formula): $\sum_{i,j} (O_{ij} - E_{ij})^2 / E_{ij}$
 - where O_{ij}/E_{ij} is the observed/expected count of events i, j
- for two-outcomes-only events:

$w_{\text{right}} \setminus w_{\text{left}}$	= true	\neq true
= species	9	1,770
\neq species	75	219,243

$$\chi^2 = \frac{221097(219243 \times 9 - 75 \times 1770)^2}{1779 \times 84 \times 221013 \times 219318} = 103.39 > 7.88$$

(at .005 thus we can reject the independence assumption)

Pointwise Mutual Information

- This is **NOT** the MI as defined in Information Theory
 - (IT: average of the following; not of values)

- ...but might be useful:

$$I'(a,b) = \log_2 (p(a,b) / p(a)p(b)) = \log_2 (p(a|b) / p(a))$$

- Example (same):

$$I'(\text{true}, \text{species}) = \log_2 (4.1\text{e-}5 / 3.8\text{e-}4 \times 8.0\text{e-}3) = 3.74$$

$$I'(\text{general}, \text{term}) = \log_2 (9.0\text{e-}6 / 1.8\text{e-}4 \times 4.9\text{e-}4) = 6.68$$

- measured in bits but it is difficult to give it an interpretation
- used for ranking (\neq the null hypothesis tests)

Mutual Information and Word Classes

The Problem

- Not enough data
 - **Language Modeling: we do not see “correct” n-grams**
 - solution so far: smoothing
 - **suppose we see:**
 - short homework, short assignment, simple homework
 - **but not:**
 - simple assignment
 - **What happens to our (bigram) LM?**
 - $p(\text{homework} \mid \text{simple}) = \text{high probability}$
 - $p(\text{assignment} \mid \text{simple}) = \text{low probability (smoothed with } p(\text{assignment}))$
- They should be much closer!

Word Classes

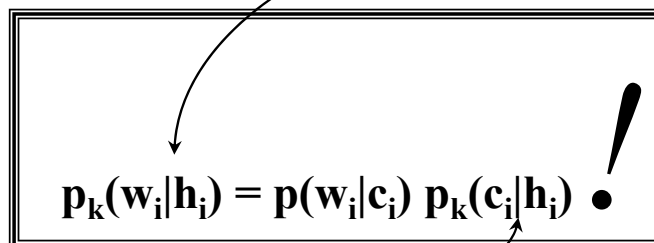
- Observation: similar words behave in a similar way
 - trigram LM:
 - trigram LM, conditioning:
 - a ... homework (any attribute of homework: short, simple, late, difficult),
 - ... the woods (any verb that has the woods as an object: walk, cut, save)
 - trigram LM: both:
 - a (short,long,difficult,...) (homework,assignment,task,job,...)

Solution

- Use the Word Classes as the “reliability” measure
- Example: we see
 - **short homework, short assignment, simple homework**
 - but not:
 - **simple assignment**
 - Cluster into classes:
 - **(short, simple) (homework, assignment)**
 - covers “simple assignment”, too
- Gaining: realistic estimates for unseen n-grams
- Loosing: accuracy (level of detail) within classes

The New Model

- Rewrite the n-gram LM using classes:
 - Was: $[k = 1..n]$
 - $p_k(w_i|h_i) = c(h_i, w_i) / c(h_i)$ [history: (k-1) words]
 - Introduce classes:


$$p_k(w_i|h_i) = p(w_i|c_i) p_k(c_i|h_i) \quad \bullet$$

- history: classes, too: [for trigram: $h_i = c_{i-2}, c_{i-1}$, bigram: $h_i = c_{i-1}$]
- Smoothing as usual
 - over $p_k(w_i|h_i)$, where each is defined as above (except uniform which stays at $1/|V|$)

Training Data

- Suppose we already have a mapping:
 - $r: V \rightarrow C$ assigning each word its class ($c_i = r(w_i)$)
- Expand the training data:
 - $T = (w_1, w_2, \dots, w_{|T|})$ into
 - $T_C = (<w_1, r(w_1)>, <w_2, r(w_2)>, \dots, <w_{|T|}, r(w_{|T|})>)$
- Effectively, we have two streams of data:
 - word stream: $w_1, w_2, \dots, w_{|T|}$
 - class stream: $c_1, c_2, \dots, c_{|T|}$ (def. as $c_i = r(w_i)$)
- Expand Heldout, Test data too

Training the New Model

- As expected, using ML estimates:
 - $p(w_i|c_i) = p(w_i|r(w_i)) = c(w_i) / c(r(w_i)) = c(w_i) / c(c_i)$
 - **!!! $c(w_i, c_i) = c(w_i)$ [since c_i determined by w_i]**
 - $p_k(c_i|h_i)$:
 - $p_3(c_i|h_i) = p_3(c_i|c_{i-2}, c_{i-1}) = c(c_{i-2}, c_{i-1}, c_i) / c(c_{i-2}, c_{i-1})$
 - $p_2(c_i|h_i) = p_2(c_i|c_{i-1}) = c(c_{i-1}, c_i) / c(c_{i-1})$
 - $p_1(c_i|h_i) = p_1(c_i) = c(c_i) / |T|$
- Then smooth as usual
 - not the $p(w_i|c_i)$ nor $p_k(c_i|h_i)$ individually, but the **$p_k(w_i|h_i)$**

Classes: How To Get Them

- We supposed the classes are given
- Maybe there are in [human] dictionaries, but...
 - dictionaries are incomplete
 - dictionaries are unreliable
 - do not define classes as equivalence relation (overlap)
 - do not define classes suitable for LM
 - **small, short... maybe; small and difficult?**
- → we have to construct them from data (again...)

Creating the Word-to-Class Map

- We will talk about bigrams from now
- Bigram estimate:
 - $p_2(c_i|h_i) = p_2(c_i|c_{i-1}) = c(c_{i-1}, c_i) / c(c_{i-1}) = c(r(w_{i-1}), r(w_i)) / c(r(w_{i-1}))$
- Form of the model:
 - just raw bigram for now:
 - $P(T) = \prod_{i=1..|T|} p(w_i|r(w_i)) p_2(r(w_i)|r(w_{i-1})) \quad (p_2(c_1|c_0) =_{df} p(c_1))$
- Maximize over r (given $r \rightarrow$ fixed p, p_2):
 - define objective $L(r) = 1/|T| \sum_{i=1..|T|} \log(p(w_i|r(w_i)) p_2(r(w_i)|r(w_{i-1})))$
 - $r_{best} = \operatorname{argmax}_r L(r)$ ($L(r)$ = norm. logprob of training data... as usual)

Simplifying the Objective Function

- Start from $L(r) = 1/|T| \sum_{i=1..|T|} \log(p(w_i|r(w_i)) p_2(r(w_i)|r(w_{i-1})))$:

$$1/|T| \sum_{i=1..|T|} \log(p(w_i|r(w_i)) \frac{p(r(w_i))}{p(r(w_i))} p_2(r(w_i)|r(w_{i-1})) / \frac{p(r(w_i))}{p(r(w_i))}) =$$

$$1/|T| \sum_{i=1..|T|} \log(\frac{p(w_i)r(w_i)}{p(r(w_i))} p_2(r(w_i)|r(w_{i-1})) / p(r(w_i))) =$$

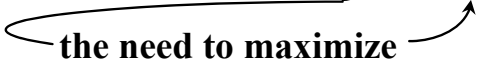
$$1/|T| \sum_{i=1..|T|} \log(p(w_i)) + 1/|T| \sum_{i=1..|T|} \log(p_2(r(w_i)|r(w_{i-1})) / p(r(w_i))) =$$

$$-H(W) + 1/|T| \sum_{i=1..|T|} \log(p_2(r(w_i)|r(w_{i-1})) \frac{p(r(w_{i-1}))}{(p(r(w_{i-1})) p(r(w_i)))}) =$$

$$-H(W) + 1/|T| \sum_{i=1..|T|} \log(\frac{p(r(w_i),r(w_{i-1}))}{(p(r(w_{i-1})) p(r(w_i)))}) =$$

$$-H(W) + \sum_{d,e \in C} p(d,e) \log(p(d,e) / (p(d) p(e))) =$$

$$-H(W) + I(D,E)$$

(event E picks class adjacent (to the right) to the one picked by D)
- Since W does not depend on r, we ended up with $I(D,E)$.


the need to maximize

Maximizing Mutual Information

(dependent on the mapping r)

- Result from previous foil:
 - Maximizing the probability of data amounts to maximizing $I(D,E)$, the mutual information of the adjacent classes.
- Good:
 - We know what a MI is, and we know how to maximize.
- Bad:
 - There is no way how to maximize over so many possible partitionings: $|V|^{|V|}$
 - no way to test them all.

Training or Heldout?

- Training:
 - best $I(D,E)$: all words in a class of its own
 - **will not give us anything new.**
- Heldout: ok, but:
 - must smooth to test any possible partitioning (unfeasible):
 - **using raw model: 0 probability of heldout (almost) guaranteed**
 - will not be able to compare anything
 - some smoothing estimates? (to be explored...)
- Solution:
 - use training anyway, but only keep $I(D,E)$ as large as possible

The Greedy Algorithm

- Define merging operation on the mapping $r: V \rightarrow C$:
 - $\text{merge}: R \times C \times C \rightarrow R' \times C^{-1}: (r, k, l) \rightarrow r', C'$ such that
 - $C^{-1} = \{C - \{k, l\} \cup \{m\}\}$ (throw out k and l , add new $m \notin C$)
 - $r'(w) = \dots m$ for $w \in r_{\text{INV}}(\{k, l\})$,
..... $r(w)$ otherwise.
- 1. Start with each word in its own class ($C = V$), $r = \text{id}$.
- 2. Merge two classes k, l into one, m , such that
$$(k, l) = \text{argmax}_{k, l} I_{\text{merge}(r, k, l)}(D, E).$$
- 3. Set new $(r, C) = \text{merge}(r, k, l)$.
- 4. Repeat 2 and 3 until $|C|$ reaches predetermined size.

Word Classes in Applications

- Word Sense Disambiguation: context not seen [enough(-times)]
- Parsing: verb-subject, verb-object relations
- Speech recognition (acoustic model): need more instances of [rare(r)] sequences of phonemes
- Machine Translation: translation equivalent selection [for rare(r) words]

Word Classes: Programming Tips & Tricks

The Algorithm (review)

- Define $\text{merge}(r, k, l) = (r', C')$ such that
 - $C' = C - \{k, l\} \cup \{m \text{ (a new class)}\}$
 - $r'(w) = r(w)$ except for k, l member words for which it is m .
- 1. Start with each word in its own class ($C = V$), $r = \text{id}$.
- 2. Merge two classes k, l into one, m , such that
$$(k, l) = \operatorname{argmax}_{k, l} I_{\text{merge}(r, k, l)}(D, E).$$
- 3. Set new $(r, C) = \text{merge}(r, k, l)$.
- 4. Repeat 2 and 3 until $|C|$ reaches a predetermined size.

Complexity Issues

- Still too complex:
 - $|V|$ iterations of the steps 2 and 3.
 - $|V|^2$ steps to maximize $\text{argmax}_{k,l}$ (selecting k,l freely from $|C|$, which is in the order of $|V|^2$)
 - $|V|^2$ steps to compute $I(D,E)$ (sum within sum, all classes, also: includes log)
 - \Rightarrow total: $|V|^5$
 - i.e., for $|V| = 100$, about 10^{10} steps; \sim several hours!
 - but $|V| \sim 50,000$ or more

Trick #1: Recomputing The MI the Smart Way: Subtracting...

- Bigram count table:

$l \setminus r$	c_1	c_2	c_3	c_4
c_1	10	2	0	1
c_2	0	0	5	2
c_3	0	2	0	3
c_4	2	3	0	0

- Test-merging c_2 and c_4 : recompute only rows/cols 2 & 4:
 - subtract column/row (2 & 4) from the MI sum (intersect.!)
 - add sums of merged counts (row & column)

...and Adding

- Add the merged counts:

$l \setminus r$	c_1	c_2'	c_3
c_1	10	3	0
c_2'	2	5	5
c_3	0	5	0



- Be careful at intersections:
 - (don't forget to add this:)

	c_2	c_3	c_4
c_2	0	5	2
c_3	2	0	3
c_4	3	0	0



Trick #2: Precompute the Counts-to-be-Subtracted

- Summing loop goes through i, j
- ...but the single row/column sums do not depend on the (resulting sums after the) merge
- \Rightarrow can be precomputed
 - **only $2k$ logs to compute at each algorithm iteration, instead of k^2**
- Then for each “merge-to-be” compute only add-on sums, plus “intersection adjustment”

Formulas for Tricks #1 and #2

- Let's have \underline{k} classes at a certain iteration. Define:

$$q_k(l,r) = p_k(l,r) \log(p_k(l,r) / (p_{kl}(l) p_{kr}(r)))$$

now the same, but using counts:

$$q_k(l,r) = c_k(l,r)/N \log(N c_k(l,r)/(c_{kl}(l) c_{kr}(r)))$$

- Define further (row+column \underline{i} sum): intersection adjustment
 $s_k(a) = \sum_{l=1..k} q_k(l,a) + \sum_{r=1..k} q_k(a,r) - q_k(a,a)$
←
- Then, the subtraction part of Trick #1 amounts to
 $sub_k(a,b) = s_k(a) + s_k(b) - q_k(a,b) - q_k(b,a)$
← remaining intersect. adj.

precomputed

Formulas - cont.

- After-merge add-on:

$$\text{add}_k(a,b) = \sum_{l=1..k, l \neq a,b} q_k(l, a+b) + \sum_{r=1..k, r \neq a,b} q_k(a+b, r) + q_k(a+b, a+b)$$

- What is it a+b? Answer: the new (merged) class.
- Hint: use the definition of q_k as a “macro”, and then
$$p_k(a+b, r) = p_k(a, r) + p_k(b, r) \text{ (same for other sums, equivalent)}$$
- The above sums cannot be precomputed
- After-merge Mutual Information (I_k is the “old” MI, kept from previous iteration of the algorithm):

$$I_k(a,b) \text{ (MI after merge of cl. } a,b) = I_k - \text{sub}_k(a,b) + \text{add}_k(a,b)$$

Trick #3: Ignore Zero Counts

- Many bigrams are 0
 - (see the paper: Canadian Hansards, $< .1$ % of bigrams are non-zero)
- Create linked lists of non-zero counts in columns and rows (similar effect: use perl's hashes)
- Update links after merge (after step 3)

Trick #4: Use Updated Loss of MI

- We are now down to $|V|^4$: $|V|$ merges, each merge takes $|V|^2$ “test-merges”, each test-merge involves order-of- $|V|$ operations ($\text{add}_k(i,j)$ term, foil #8)
- Observation: many numbers (s_k, q_k) needed to compute the mutual information loss due to a merge of $i+j$ ***do not change***: namely, those which are not in the vicinity of neither i nor j .
- Idea: keep the MI loss matrix for all pairs of classes, and (after a merge) update only those cells which have been influenced by the merge.

Formulas for Trick #4 (s_{k-1}, L_{k-1})

- Keep a matrix of “losses” $L_k(d,e)$.¹
- Init: $L_k(d,e) = \text{sub}_k(d,e) - \text{add}_k(d,e)$ [then $I_k(d,e) = I_k - L_k(d,e)$]
- Suppose a,b are now the two classes merged into a :
- Update ($k-1$: index used for the next iteration; $i,j \neq a,b$):
 - $s_{k-1}(i) = s_k(i) - q_k(i,a) - q_k(a,i) - q_k(i,b) - q_k(b,i) + q_{k-1}(a,i) + q_{k-1}(i,a)$
 - ${}^2L_{k-1}(i,j) = L_k(i,j) - s_k(i) + s_{k-1}(i) - s_k(j) + s_{k-1}(j) +$
 $+ q_k(i+j,a) + q_k(a,i+j) + q_k(i+j,b) + q_k(b,i+j) -$
 $- q_{k-1}(i+j,a) - q_{k-1}(a,i+j)$ [NB: may substitute even for s_k, s_{k-1}]

NB ¹ L_k is symmetrical $L_k(d,e) = L_k(e,d)$ (q_k is something different!)

²The update formula $L_{k-1}(l,m)$ is wrong in the Brown et. al paper

Completing Trick #4

- $s_{k-1}(a)$ must be computed using the “Init” sum.
- $L_{k-1}(a,i) = L_{k-1}(i,a)$ must be computed in a similar way, for all $i \neq a,b$.
- $s_{k-1}(b)$, $L_{k-1}(b,i)$, $L_{k-1}(i,b)$ are not needed anymore (keep track of such data, i.e. mark every class already merged into some other class and do not use it anymore).
- Keep track of the minimal loss during the $L_k(i,j)$ update process (so that the next merge to be taken is obvious immediately after finishing the update step).

Efficient Implementation

- Data Structures: (N - # of bigrams in data [fixed])
 - $\text{Hist}(k)$ history of merges
 - **$\text{Hist}(k) = (a,b)$ merged when the remaining number of classes was k**
 - $c_k(i,j)$ bigram class counts [updated]
 - $c_{kl}(i), c_{kr}(i)$ unigram (marginal) counts [updated]
 - $L_k(a,b)$ table of losses; upper-right triangle [updated]
 - $s_k(a)$ “subtraction” subterms [optionally updated]
 - $q_k(i,j)$ subterms involving a log [opt. updated]
 - **The optionally updated data structures will give linear improvement only in the subsequent steps, but at least $s_k(i)$ is necessary in the initialization phase (1st iteration)**

Implementation: the Initialization Phase

- 1 Read data in, init counts $c_k(l,r)$; then $\forall l,r,a,b; a < b$:
- 2 Init unigram counts:

$$c_{kl}(l) = \sum_{r=1..k} c_k(l,r), \quad c_{kr}(r) = \sum_{l=1..k} c_k(l,r)$$

– complicated? remember, must take care of start & end of data!

- 3 Init $q_k(l,r)$: use the 2nd formula (count-based) on foil 7,

$$q_k(l,r) = c_k(l,r) / N \log(N c_k(l,r) / (c_{kl}(l) c_{kr}(r)))$$

- 4 Init $s_k(a) = \sum_{l=1..k} q_k(l,a) + \sum_{r=1..k} q_k(a,r) - q_k(a,a)$
- 5 Init $L_k(a,b) = s_k(a) + s_k(b) - q_k(a,b) - q_k(b,a) - q_k(a+b,a+b) +$
 $- \sum_{l=1..k, l \neq a,b} q_k(l,a+b) - \sum_{r=1..k, r \neq a,b} q_k(a+b,r)$

Implementation: Select & Update

- 6 Select the best pair $(\underline{a}, \underline{b})$ to merge into \underline{a} (watch the candidates when computing $L_k(a, b)$); save to $\text{Hist}(k)$
- 7 Optionally, update $q_k(i, j)$ for all $i, j \neq b$, get $q_{k-1}(i, j)$
 - remember those $q_k(i, j)$ values needed for the updates below
- 8 Optionally, update $s_k(i)$ for all $i \neq b$, to get $s_{k-1}(i)$
 - again, remember the $s_k(i)$ values for the “loss table” update
- 9 Update the loss table, $L_k(i, j)$, to $L_{k-1}(i, j)$, using the tabulated q_k , q_{k-1} , s_k and s_{k-1} values, or compute the needed $q_k(i, j)$ and $q_{k-1}(i, j)$ values dynamically from the counts: $c_k(i+j, b) = c_k(i, b) + c_k(j, b)$; $c_{k-1}(a, i) = c_k(a+b, i)$

Towards the Next Iteration

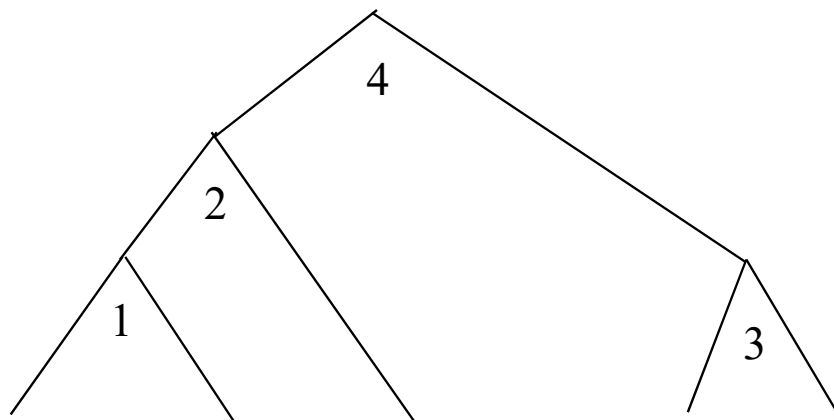
- 10 During the $L_k(i,j)$ update, keep track of the minimal loss of MI, and the two classes which caused it.
- 11 Remember such best merge in $\text{Hist}(k)$.
- 12 Get rid of all s_k , q_k , L_k values.
- 13 Set $k = k - 1$; stop if $k == 1$.
- 14 Start the next iteration
 - either by the optional updates (steps 7 and 8), or
 - directly updating $L_k(i,j)$ again (step 9).

Moving Words Around

- Improving Mutual Information
 - take a word from one class, move it to another (i.e., two classes change: the moved-from and the moved-to), compute $I_{\text{new}}(D,E)$; keep change permanent if
$$I_{\text{new}}(D,E) > I(D,E)$$
 - keep moving words until no move improves $I(D,E)$
- Do it at every iteration, or at every m iterations
- Use similar “smart” methods as for merging

Using the Hierarchy

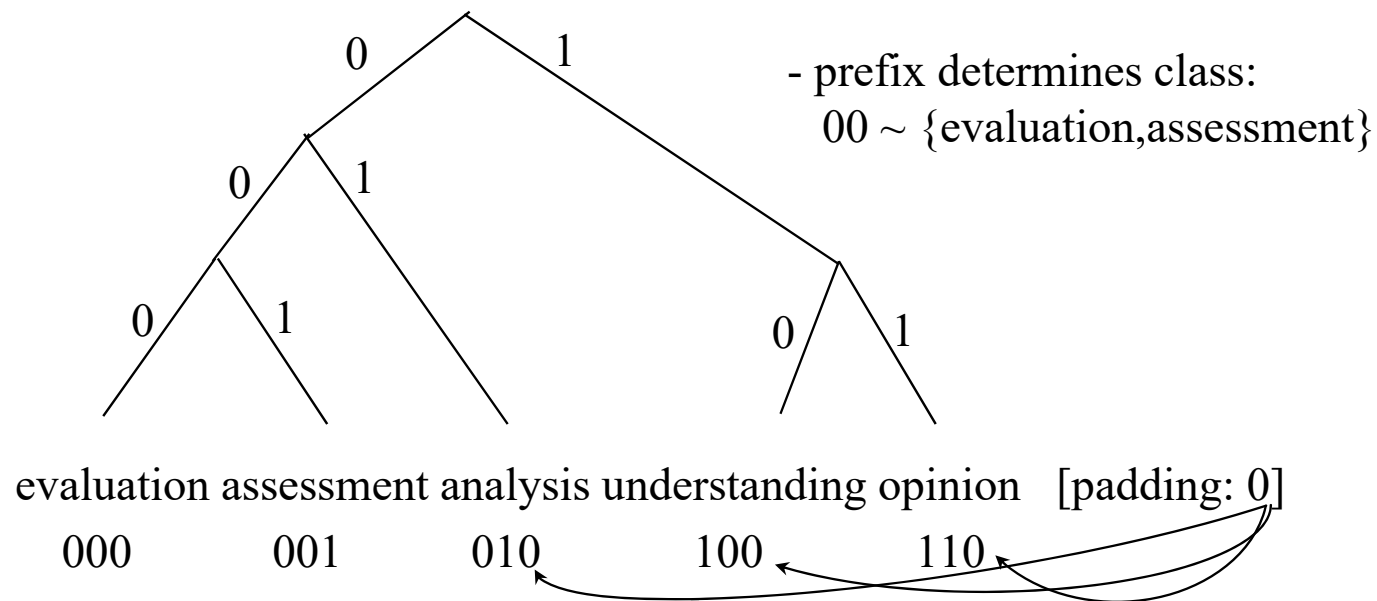
- Natural Form of Classes
 - follows from the sequence of merges:



evaluation assessment analysis understanding opinion

Numbering the Classes (within the Hierarchy)

- Binary branching
- Assign 0/1 to the left/right branch at every node:



Markov Models

Review: Markov Process

- Bayes formula (chain rule):

$$P(W) = P(w_1, w_2, \dots, w_T) = \prod_{i=1..T} p(w_i | w_1, w_2, \dots, w_{i-n+1}, \dots, w_{i-1})$$

- n-gram language models:

- Markov process (chain) of the order n-1:

$$P(W) = P(w_1, w_2, \dots, w_T) = \prod_{i=1..T} p(w_i | w_{i-n+1}, w_{i-n+2}, \dots, w_{i-1})$$

Using just one distribution (Ex.: trigram model: $p(w_i | w_{i-2}, w_{i-1})$):

Positions: 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16

Words: My car **broke down** . and within hours Bob 's car **broke down** . too .

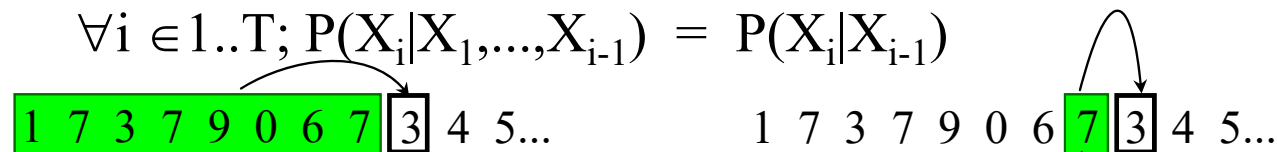
$$p(., \mathbf{broke\ down}) = p(w_5 | w_3, w_4) = p(w_{14} | w_{12}, w_{13})$$

Markov Properties

- Generalize to any process (not just words/LM):
 - Sequence of random variables: $X = (X_1, X_2, \dots, X_T)$
 - Sample space S (*states*), size N : $S = \{s_0, s_1, s_2, \dots, s_N\}$

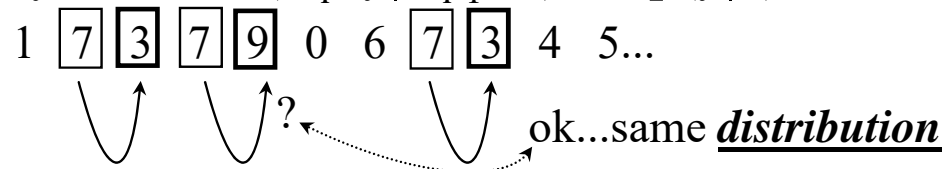
1. Limited History (Context, Horizon):

$$\forall i \in 1..T; P(X_i | X_1, \dots, X_{i-1}) = P(X_i | X_{i-1})$$



2. Time invariance (M.C. is stationary, homogeneous)

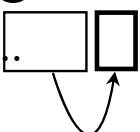
$$\forall i \in 1..T, \forall y, x \in S; P(X_i=y | X_{i-1}=x) = p(y|x)$$



Long History Possible

- What if we want trigrams:

1 7 3 7 9 0 6 7 3 4 5.



- Formally, use transformation:

Define new variables Q_i , such that $X_i = \{Q_{i-1}, Q_i\}$:

Then

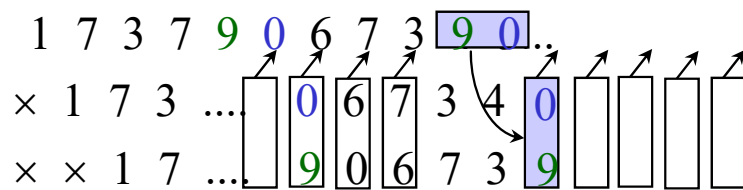
$$P(X_i | X_{i-1}) = P(Q_{i-1}, Q_i | Q_{i-2}, Q_{i-1}) = P(Q_i | Q_{i-2}, Q_{i-1})$$

Predicting (X_i):

1 7 3 7 9 0 6 7 3 9 0 . . .

× 1 7 3 ... 0 6 7 3 4 0

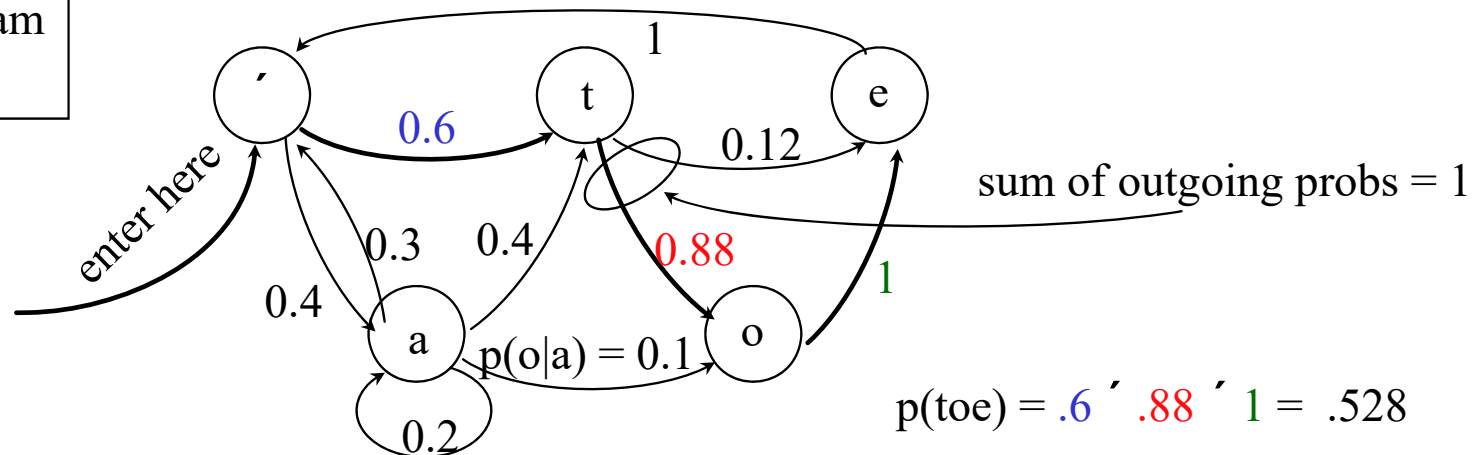
History ($X_{i-1} = \{Q_{i-2}, Q_{i-1}\}$): × × 1 7 ... 9 0 6 7 3 9



Graph Representation: State Diagram

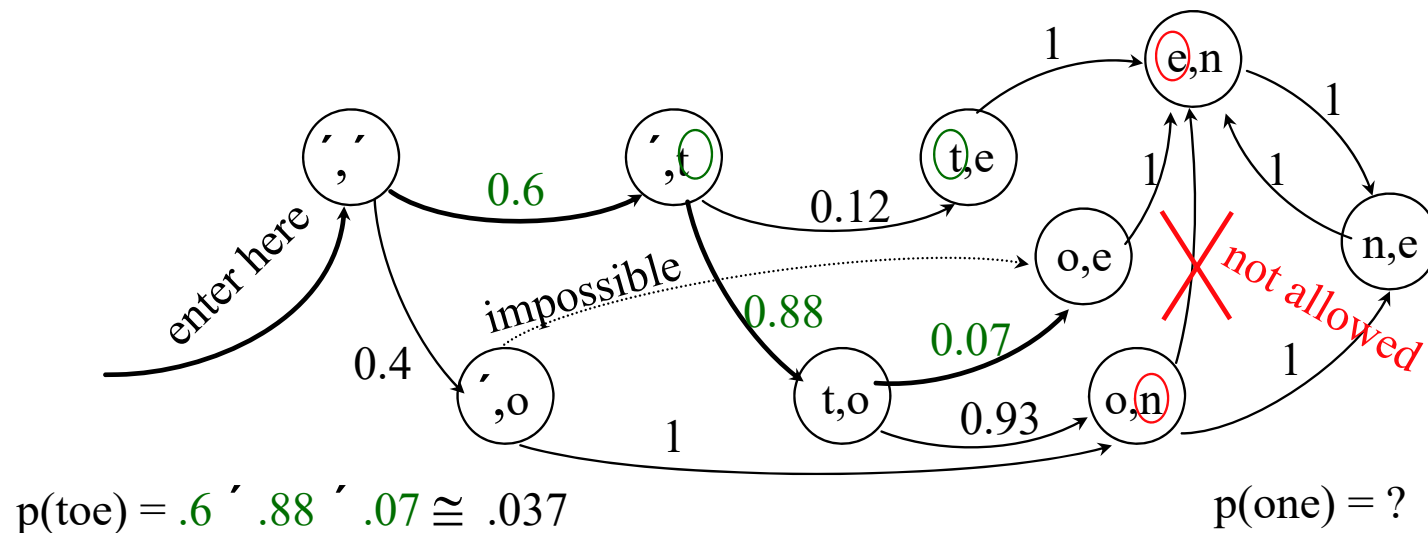
- $S = \{s_0, s_1, s_2, \dots, s_N\}$: states
- Distribution $P(X_i|X_{i-1})$:
 - transitions (as arcs) with probabilities attached to them:

Bigram
case:

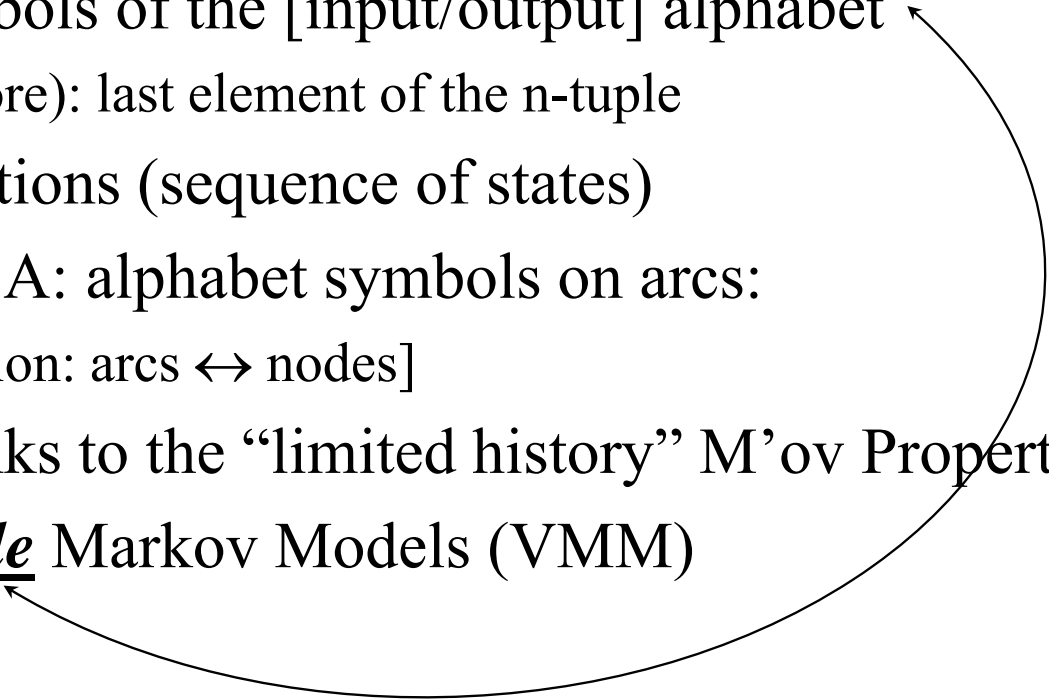


The Trigram Case

- $S = \{s_0, s_1, s_2, \dots, s_N\}$: states: pairs $s_i = (x, y)$
- Distribution $P(X_i | X_{i-1})$: (r.v. X : generates pairs s_i)

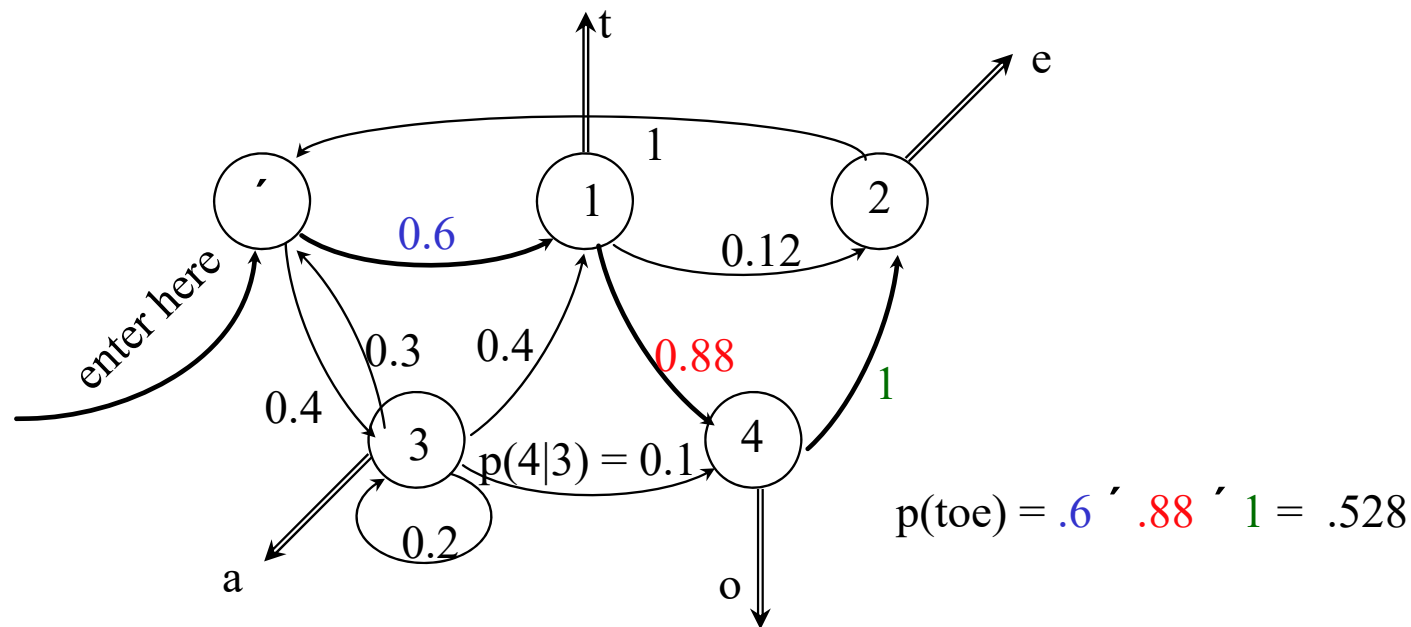


Finite State Automaton

- States \sim symbols of the [input/output] alphabet
 - pairs (or more): last element of the n-tuple
 - Arcs \sim transitions (sequence of states)
 - [Classical FSA: alphabet symbols on arcs:
 - transformation: arcs \leftrightarrow nodes]
 - Possible thanks to the “limited history” Markov Property
 - So far: Visible Markov Models (VMM)
- 

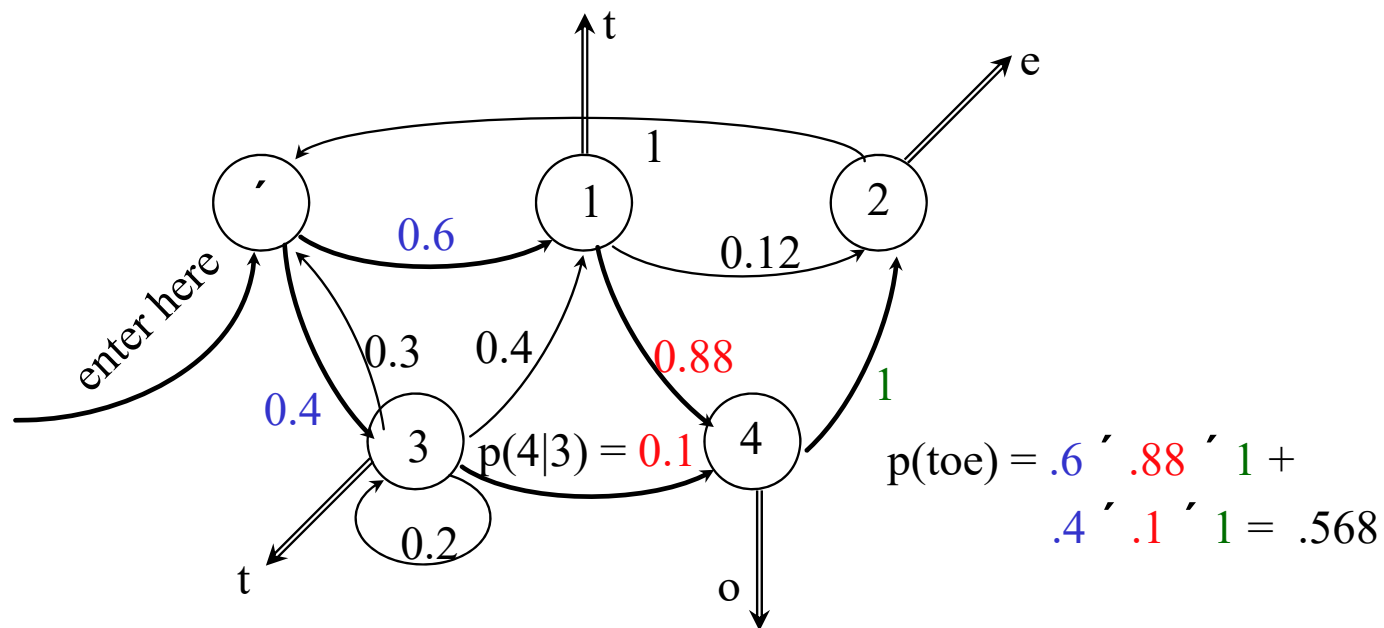
Hidden Markov Models

- The simplest HMM: states generate [observable] output (using the “data” alphabet) but remain “invisible”:



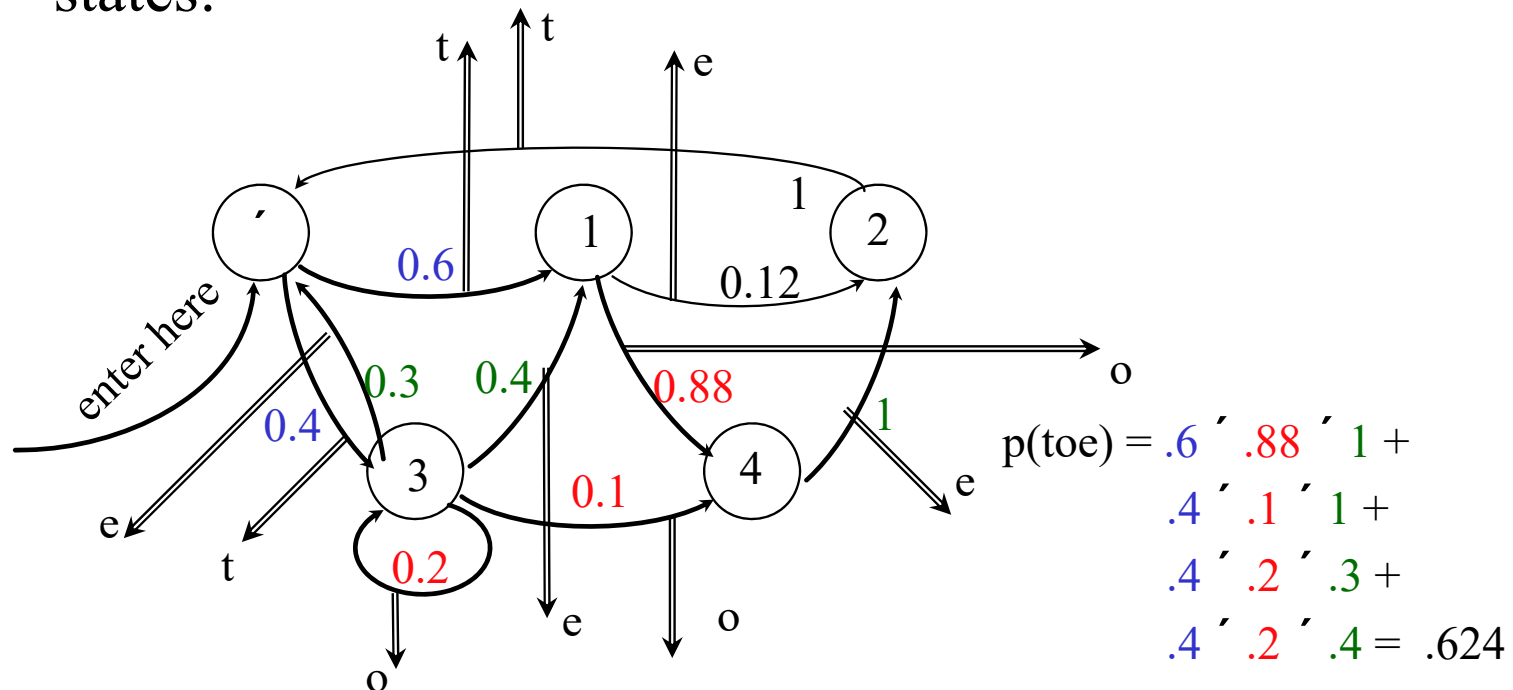
Added Flexibility

- So far, no change; but different states may generate the same output (why not?):



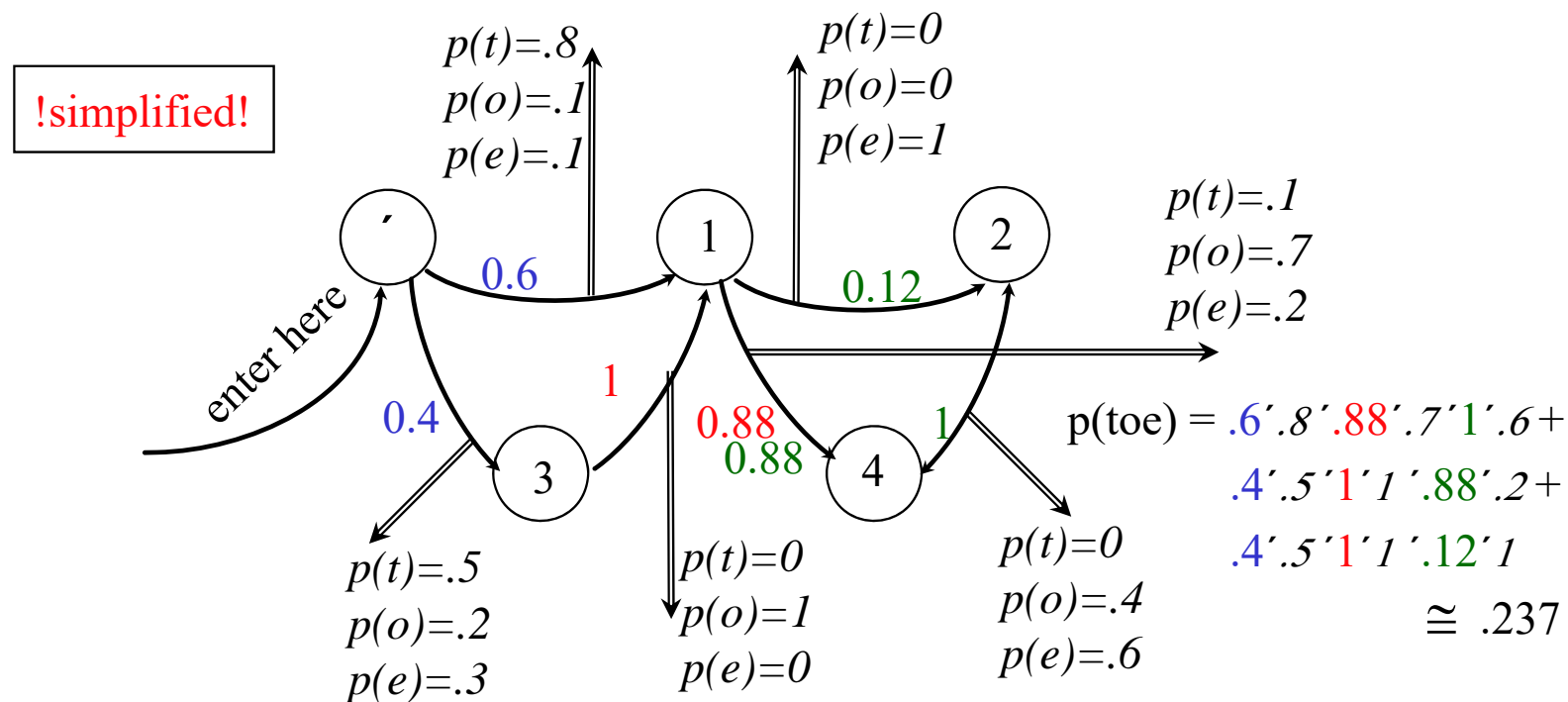
Output from Arcs...

- Added flexibility: Generate output from arcs, not states:



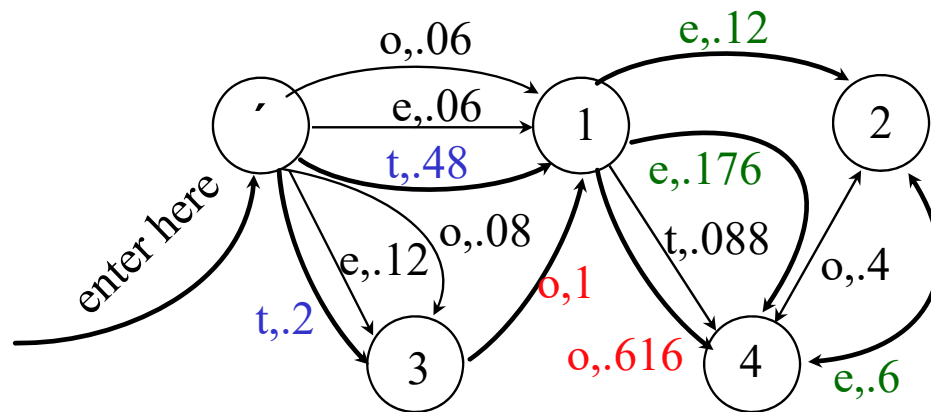
... and Finally, Add Output Probabilities

- Maximum flexibility: [Unigram] distribution (sample space: output alphabet) at each output arc:



Slightly Different View

- Allow for multiple arcs from $s_i \rightarrow s_j$, mark them by output symbols, get rid of output distributions:



$$p(\text{toe}) = .48' .616' .6 + \\ .2' 1' .176 + \\ .2' 1' .12 \cong .237$$

In the future, we will use the view more convenient for the problem at hand.

Formalization

- HMM (the most general case):
 - five-tuple (S, s_0, Y, P_S, P_Y) , where:
 - $S = \{s_0, s_1, s_2, \dots, s_T\}$ is the set of states, s_0 is the initial state,
 - $Y = \{y_1, y_2, \dots, y_V\}$ is the output alphabet,
 - $P_S(s_j|s_i)$ is the set of prob. distributions of transitions,
 - size of P_S : $|S|^2$.
 - $P_Y(y_k|s_i, s_j)$ is the set of output (emission) probability distributions.
 - size of P_Y : $|S|^2 \times |Y|$
- Example:
 - $S = \{x, 1, 2, 3, 4\}$, $s_0 = x$
 - $Y = \{t, o, e\}$

Formalization - Example

- Example (for graph, see foils 11,12):

– $S = \{x, 1, 2, 3, 4\}$, $s_0 = x$

– $Y = \{e, o, t\}$

– P_S :

$P_{Y|S}$:

	x	1	2	3	4
x	0	.6	0	.4	0
1	0	0	.12	0	.88
2	0	0	0	0	1
3	0	1	0	0	0
4	0	0	1	0	0

$\Rightarrow \Sigma = 1$

	e	x	1	2	3	4
o		x	1	2	3	4
t		x	1	2	3	4
x			.8		.5	
1				0		.1
2					0	
3			0			
4				0		

$\Sigma = 1$

Using the HMM

- The generation algorithm (of limited value :-)):
 1. **Start in $s = s_0$.**
 2. **Move from s to s' with probability $P_s(s'|s)$.**
 3. **Output (emit) symbol y_k with probability $P_s(y_k|s,s')$.**
 4. **Repeat from step 2 (until somebody says enough).**
- More interesting usage:
 - Given an output sequence $Y = \{y_1, y_2, \dots, y_k\}$, compute its probability.
 - Given an output sequence $Y = \{y_1, y_2, \dots, y_k\}$, compute the most likely sequence of states which has generated it.
 - ...plus variations: e.g., n best state sequences

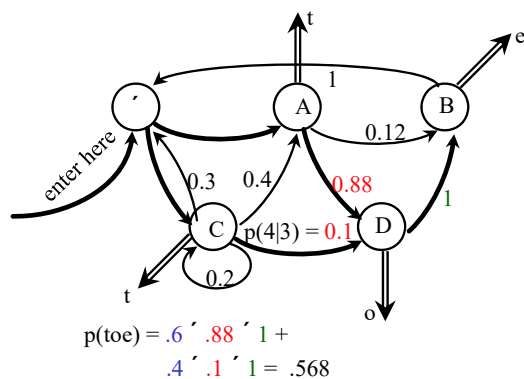
HMM Algorithms: Trellis and Viterbi

HMM: The Two Tasks

- HMM (the general case):
 - five-tuple (S, S_0, Y, P_S, P_Y) , where:
 - $S = \{s_1, s_2, \dots, s_T\}$ is the set of states, S_0 is the initial state,
 - $Y = \{y_1, y_2, \dots, y_V\}$ is the output alphabet,
 - $P_S(s_j | s_i)$ is the set of prob. distributions of transitions,
 - $P_Y(y_k | s_i, s_j)$ is the set of output (emission) probability distributions.
- Given an HMM & an output sequence $Y = \{y_1, y_2, \dots, y_k\}$:
 - (Task 1) compute the probability of Y ;
 - (Task 2) compute the most likely sequence of states which has generated Y .

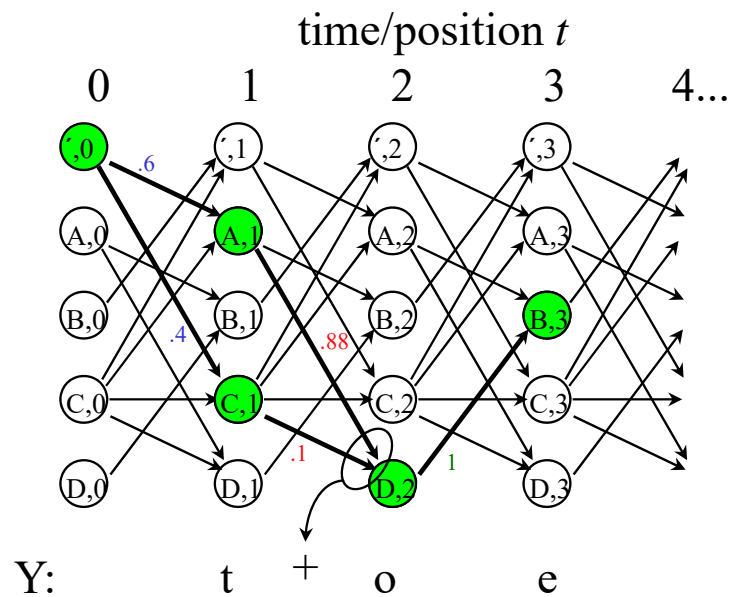
Trellis - Deterministic Output

HMM:



Trellis:

“rollout”

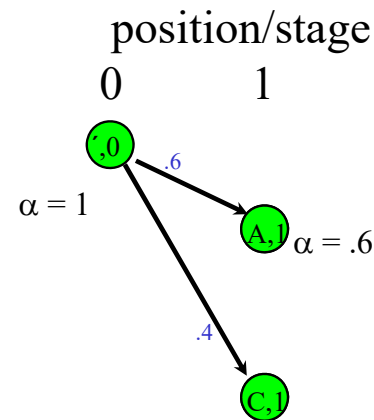
- trellis state: (HMM state, position)
- each state: holds one number (prob): α
- probability or Y: $\sum \alpha$ in the last state


$$\alpha(,0) = 1 \quad \alpha(A,1) = .6 \quad \alpha(D,2) = .568 \quad \alpha(B,3) = .568$$

$$\alpha(C,1) = .4$$

Creating the Trellis: The Start

- Start in the start state (\times),
 - set its $\alpha(\times, 0)$ to 1.
- Create the first stage:
 - get the first “output” symbol y_1
 - create the first stage (column)
 - but only those trellis states
 - which generate y_1
 - set their $\alpha(state, 1)$ to the $P_S(state|\times) \alpha(\times, 0)$
- ...and forget about the 0-th stage



$y_1:$ t


Trellis: The Next Step

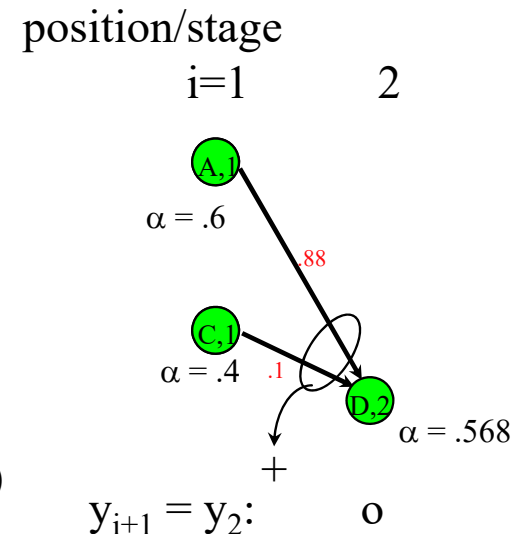
- Suppose we are in stage i
- Creating the next stage:

- create all trellis states in the next stage which generate y_{i+1} , but only those reachable from any of the stage- i states
- set their $\alpha(state, i+1)$ to:

$$P_S(state|prev.state) \times \alpha(prev.state, i)$$

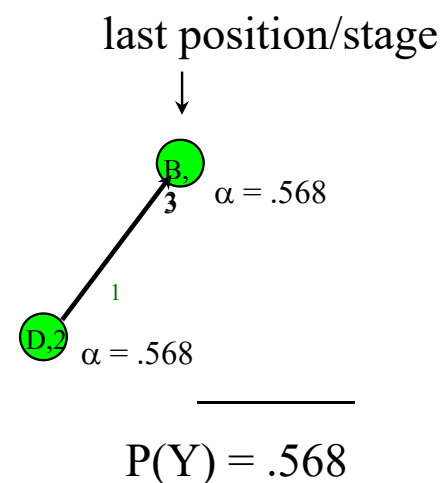
(add up all such numbers on arcs going to a common trellis state)

- ...and forget about stage i



Trellis: The Last Step

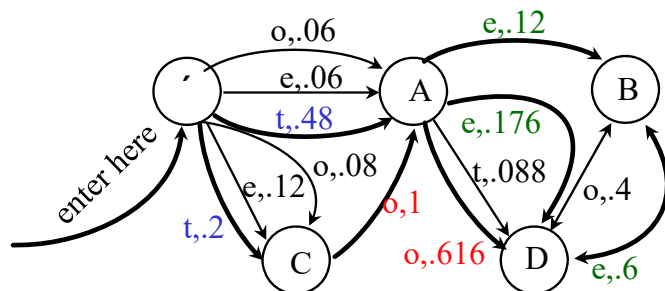
- Continue until “output” exhausted
 - $|Y| = 3$: until stage 3
- Add together all the $\alpha(state, /Y/)$
- That’s the $P(Y)$.
- Observation (pleasant):
 - memory usage max: $2|S|$
 - multiplications max: $|S|^2|Y|$



Trellis: The General Case (still, bigrams)

- Start as usual:
 - start state ($'$), set its $\alpha(' , 0)$ to 1.

$\alpha = 1$

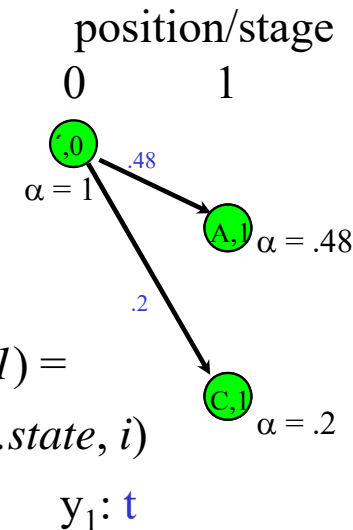
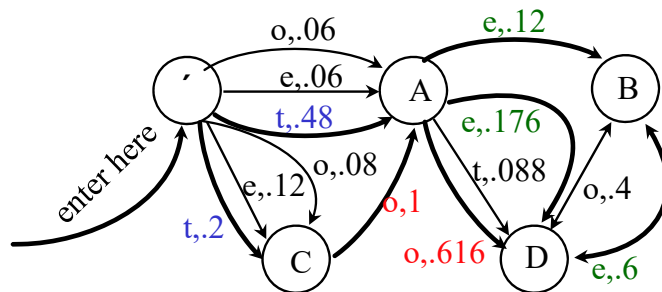


$$\begin{aligned}
 p(\text{toe}) &= .48' .616' .6 + \\
 &\quad .2' 1' .176 + \\
 &\quad .2' 1' .12 \cong .237
 \end{aligned}$$

General Trellis: The Next Step

- We are in stage i :
 - Generate the next stage $i+1$ as before (except now arcs generate output, thus use only those arcs marked by the output symbol y_{i+1})

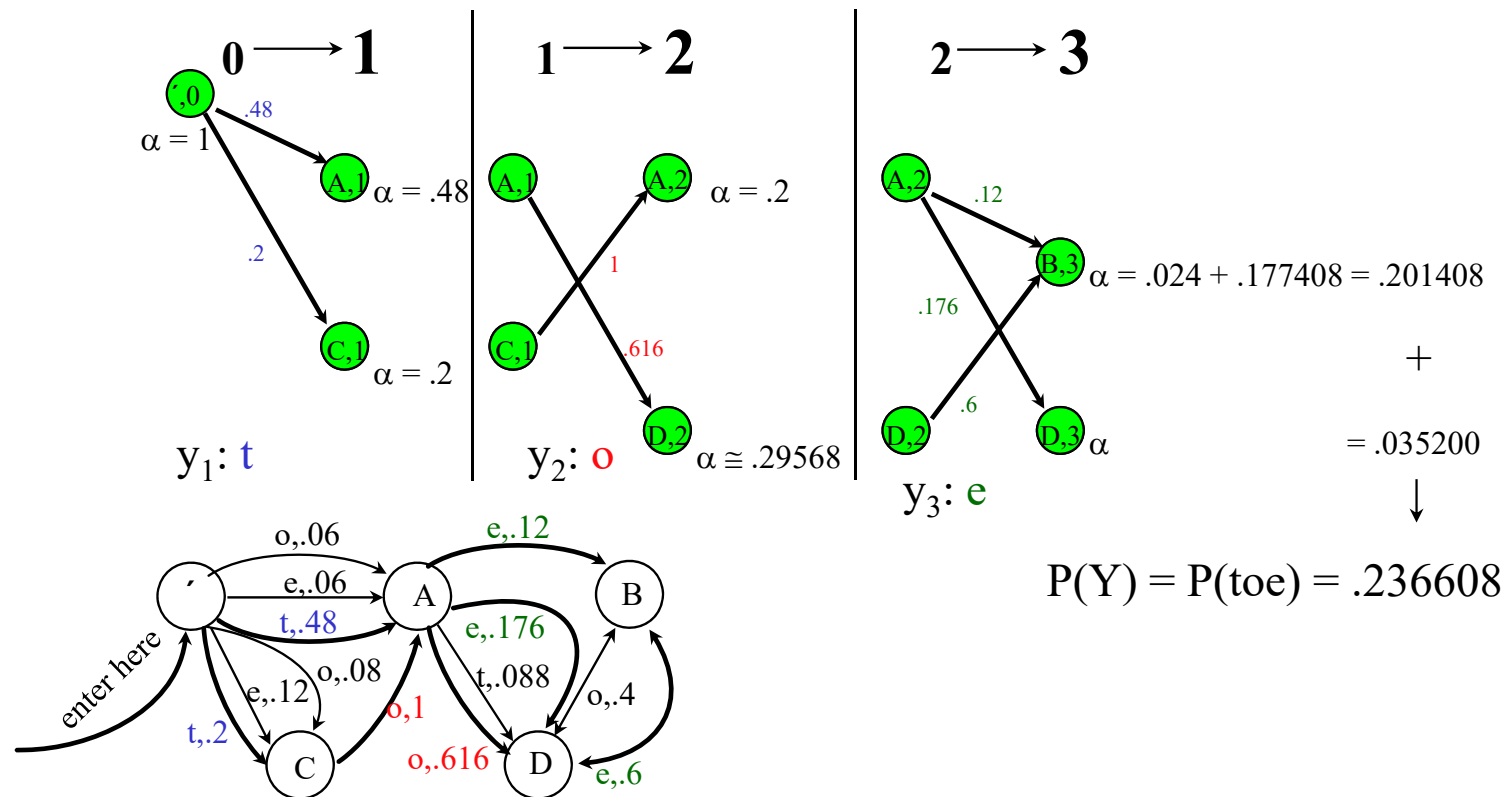
- For each generated *state*, compute $\alpha(state, i+1) = \sum_{\text{incoming arcs}} P_Y(y_{i+1} | state, prev.state) \times \alpha(prev.state, i)$



...and forget about stage i as usual.

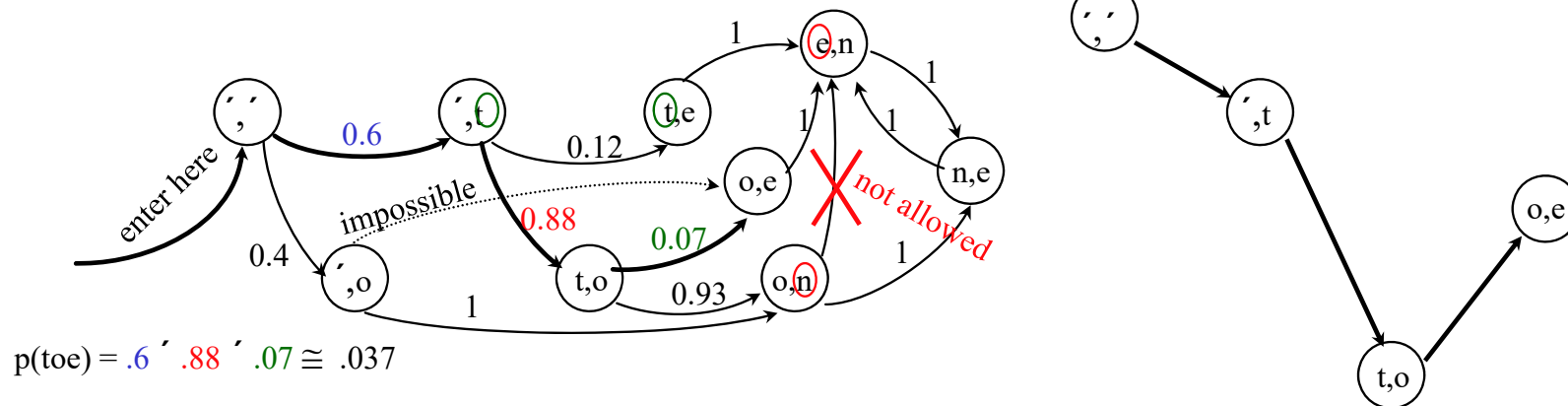
Trellis: The Complete Example

Stage:



The Case of Trigrams

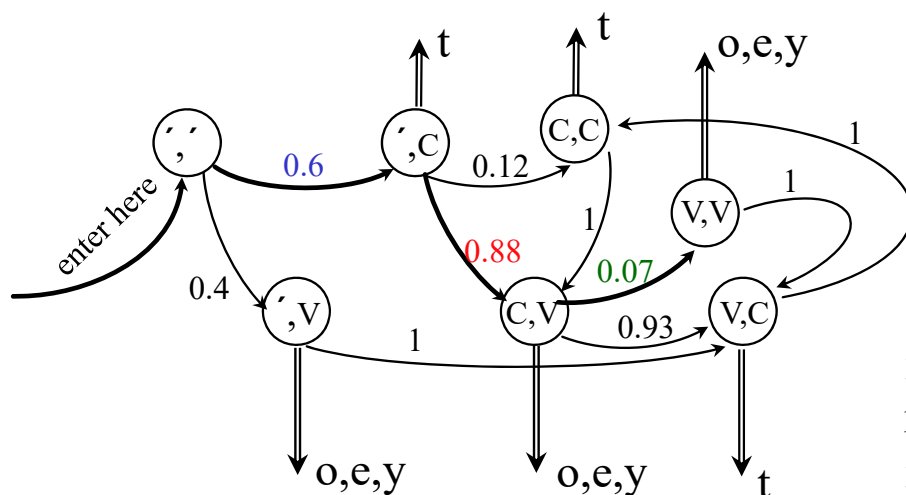
- Like before, but:
 - states correspond to bigrams,
 - output function always emits the second output symbol of the pair (state) to which the arc goes:



Multiple paths not possible → trellis not really needed

Trigrams with Classes

- More interesting:
 - n-gram class LM: $p(w_i|w_{i-2}, w_{i-1}) = p(w_i|c_i) p(c_i|c_{i-2}, c_{i-1})$
 - states are pairs of classes (c_{i-1}, c_i) , and emit “words”:
(letters in our example)



$p(t|C) = 1$ usual,
 $p(o|V) = .3$ non-
 $p(e|V) = .6$ overlapping
 $p(y|V) = .1$ classes

$$p(\text{toe}) = .6 \cdot 1 \cdot .88 \cdot .3 \cdot .07 \cdot .6 \cong .00665$$

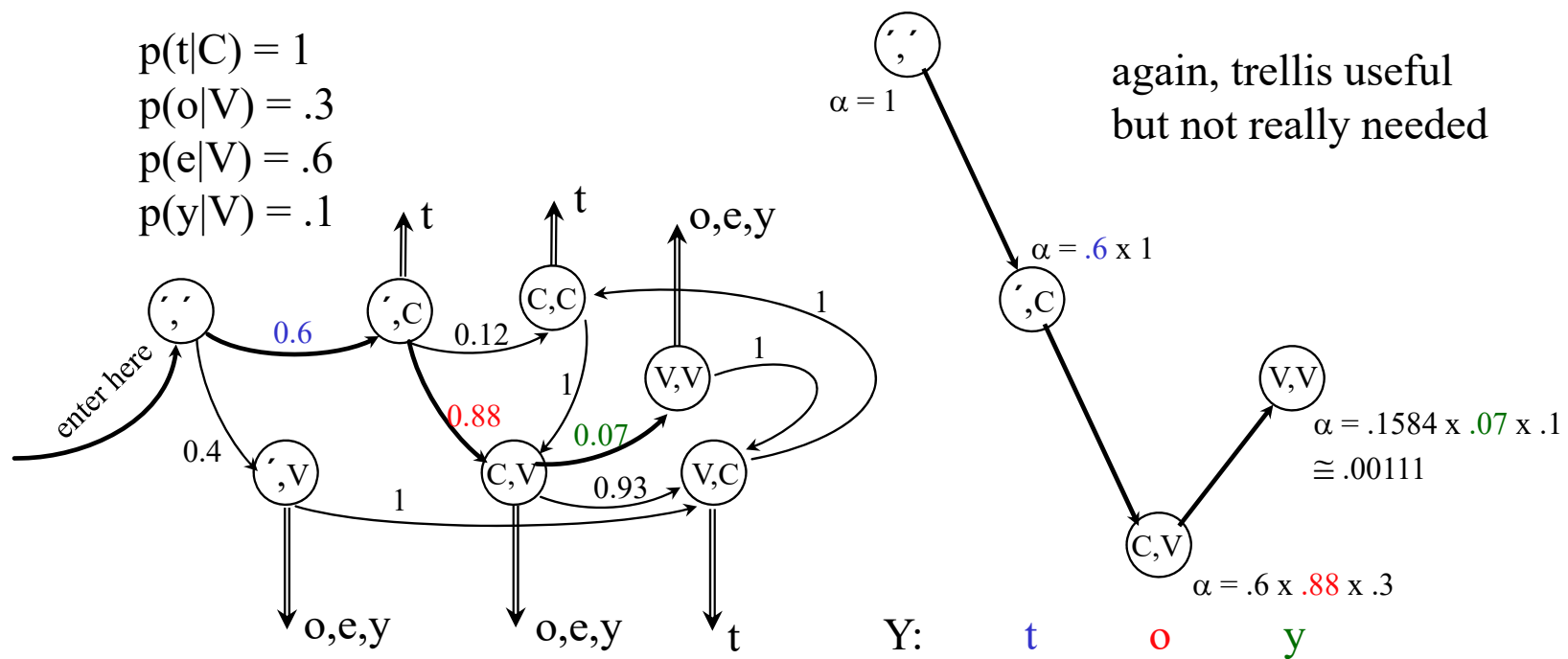
$$p(\text{teo}) = .6 \cdot 1 \cdot .88 \cdot .6 \cdot .07 \cdot .3 \cong .00665$$

$$p(\text{toy}) = .6 \cdot 1 \cdot .88 \cdot .3 \cdot .07 \cdot .1 \cong .00111$$

$$p(\text{tty}) = .6 \cdot 1 \cdot .12 \cdot 1 \cdot 1 \cdot .1 \cong .0072$$

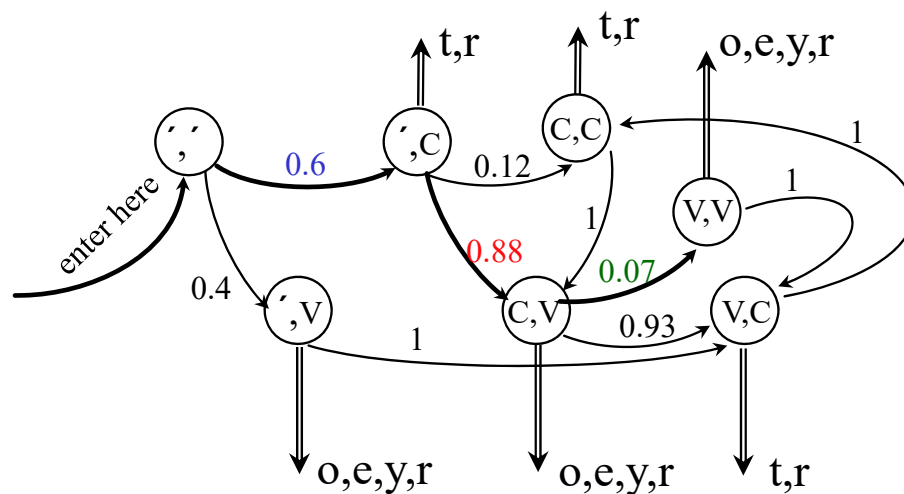
Class Trigrams: the Trellis

- Trellis generation ($Y = \text{"toy"}$):



Overlapping Classes

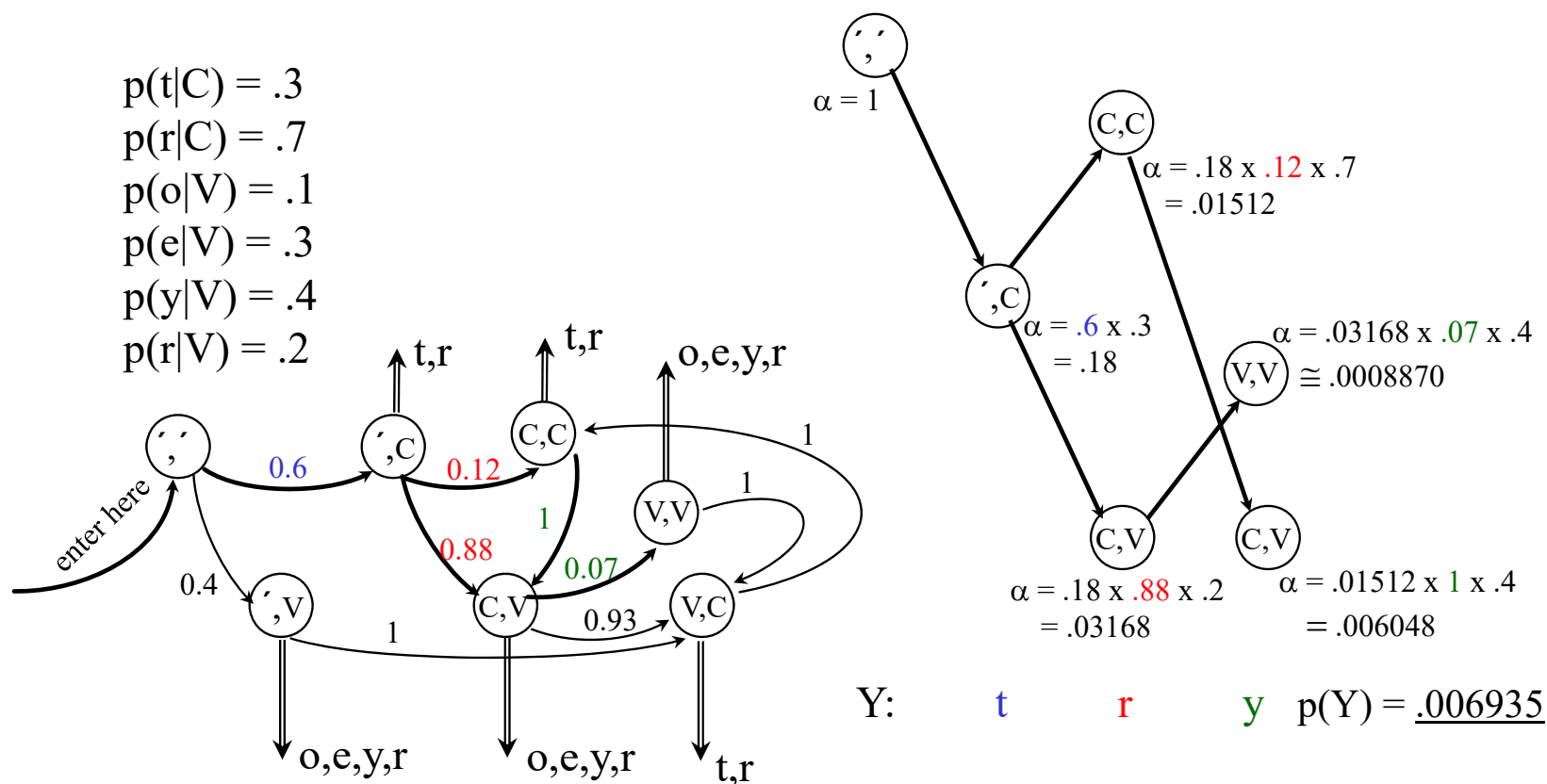
- Imagine that classes may overlap
 - e.g. 'r' is sometimes vowel sometimes consonant, belongs to V as well as C:



$$\begin{aligned}
 p(t|C) &= .3 \\
 p(r|C) &= .7 \\
 p(o|V) &= .1 \\
 p(e|V) &= .3 \\
 p(y|V) &= .4 \\
 p(r|V) &= .2
 \end{aligned}$$

$$p(\text{try}) = ?$$

Overlapping Classes: Trellis Example



Trellis: Remarks

- So far, we went left to right (computing α)
- Same result: going right to left (computing β)
 - supposed we know where to start (finite data)
- In fact, we might start in the middle going left and right
- Important for parameter estimation
(Forward-Backward Algorithm alias Baum-Welch)
- Implementation issues:
 - scaling/normalizing probabilities, to avoid too small numbers
& addition problems with many transitions

The Viterbi Algorithm

- Solving the task of finding the most likely sequence of states which generated the observed data
- i.e., finding

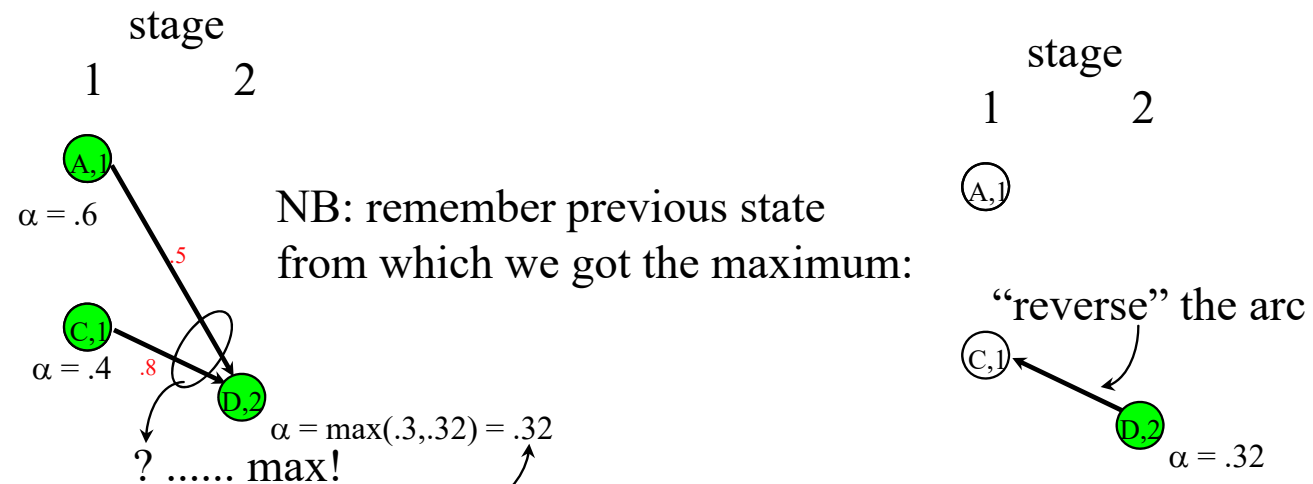
$$S_{\text{best}} = \operatorname{argmax}_S P(S|Y)$$

which is equal to (Y is constant and thus $P(Y)$ is fixed):

$$\begin{aligned} S_{\text{best}} &= \operatorname{argmax}_S P(S, Y) = \\ &= \operatorname{argmax}_S P(s_0, s_1, s_2, \dots, s_k, y_1, y_2, \dots, y_k) = \\ &= \operatorname{argmax}_S \prod_{i=1..k} p(y_i | s_i, s_{i-1}) p(s_i | s_{i-1}) \end{aligned}$$

The Crucial Observation

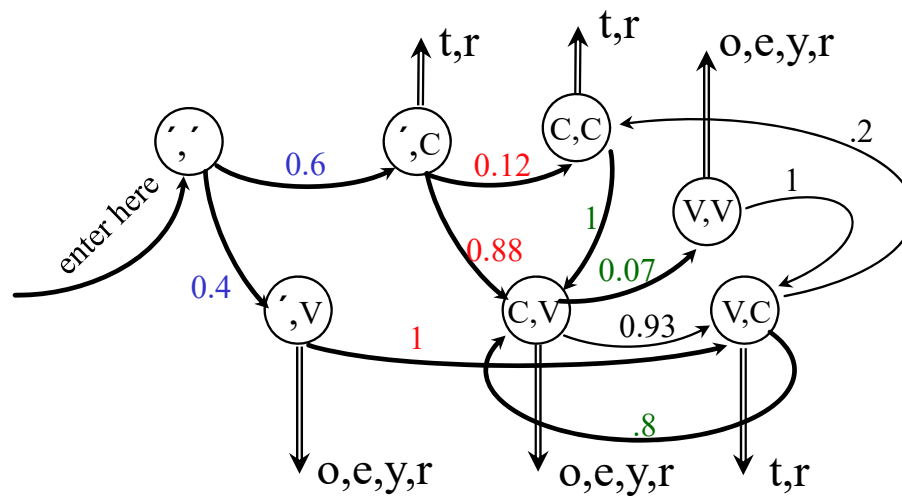
- Imagine the trellis build as before (but do not compute the α s yet; assume they are o.k.); stage i :



this is certainly the “backwards” maximum to $(D,2)$... but
it cannot change even whenever we go forward (M. Property: Limited History)

Viterbi Example

- ‘r’ classification (C or V?, sequence?):



$$\begin{aligned}
 p(t|C) &= .3 \\
 p(r|C) &= .7 \\
 p(o|V) &= .1 \\
 p(e|V) &= .3 \\
 p(y|V) &= .4 \\
 p(r|V) &= .2
 \end{aligned}$$

$$\operatorname{argmax}_{XYZ} p(rry|XYZ) = ?$$

Possible state seq.: (',v)(v,c)(c,v)[VCV], (',c)(c,c)(c,v)[CCV], (',c)(c,v)(v,v) [CVV]

α in trellis
state:
best prob
from start
to here

Diagram illustrating a Markov Decision Process (MDP) with 7 states and transitions. The states are represented by circles containing a pair of symbols (e.g., 'C,C', 'V,V'). The transitions are labeled with probabilities in different colors: blue for 0.6, 0.4, and 0.12; red for 0.88 and 1; green for 1, 0.07, and 0.8. The diagram includes a start state (smiley face) and a goal state (V,V).

States and Transitions:

- Start State (Smiley Face):**
 - Transition to 'C,C' with probability 0.6 (blue).
 - Transition to 'C,V' with probability 0.4 (blue).
- 'C,C' State:**
 - Transition to 'C,V' with probability 0.88 (red).
 - Transition to 'V,V' with probability 1 (green).
- 'C,V' State:**
 - Transition to 'V,C' with probability 0.93 (green).
 - Transition to 'V,V' with probability 0.07 (green).
- 'V,V' State (Goal State):**
 - Self-loop with probability 0.2 (black).
 - Transition to 'V,C' with probability 1 (black).
- 'V,C' State:**
 - Transition to 'V,V' with probability 0.8 (green).
- 'V,V' State (Goal State):**
 - Transition to 'C,V' with probability 1 (red).

Additional labels in the diagram include 'enter here' pointing to the start state, and various action labels (e.g., 't,r', 'o,e,y,r') associated with transitions.

Diagram illustrating the calculation of α values for various node pairs in a sequence model. The nodes represent pairs of characters, and the edges represent transitions between them. The values are calculated as products of transition probabilities and the α value of the target state.

Nodes and their associated α values:

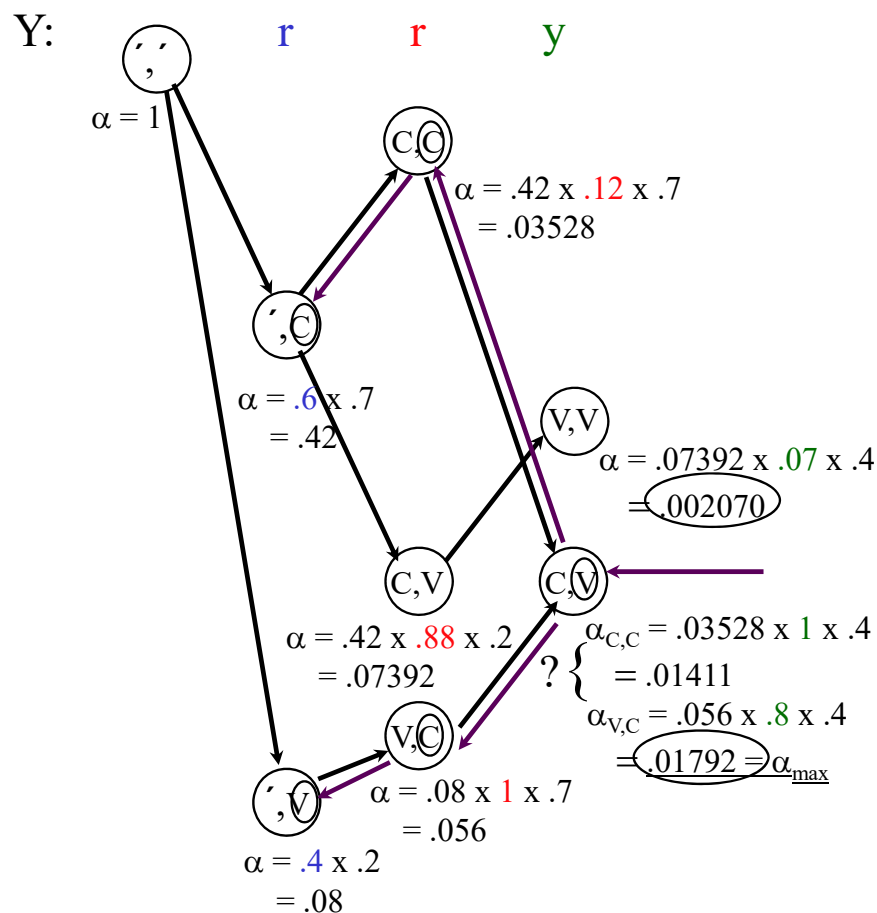
- Node $(', ')$: $\alpha = 1$
- Node $(', C)$: $\alpha = .6 \times .7 = .42$
- Node (C, C) : $\alpha = .42 \times .12 \times .7 = .03528$
- Node (C, V) : $\alpha = .42 \times .88 \times .2 = .07392$
- Node (V, V) : $\alpha = .07392 \times .07 \times .4 = .002070$
- Node (V, C) : $\alpha_{C,C} = .03528 \times 1 \times .4 = .01411$
 $\alpha_{V,C} = .056 \times .8 \times .4 = .01792 = \alpha_{\max}$
- Node $(', V)$: $\alpha = .08 \times 1 \times .7 = .056$
 $\alpha = .4 \times .2 = .08$

Transitions and associated values:

- From $(', ')$ to $(', C)$: $\alpha = .6 \times .7 = .42$
- From $(', ')$ to $(', V)$: $\alpha = .4 \times .2 = .08$
- From $(', C)$ to (C, C) : $\alpha = .42 \times .12 \times .7 = .03528$
- From $(', C)$ to (C, V) : $\alpha = .42 \times .88 \times .2 = .07392$
- From (C, C) to (V, V) : $\alpha = .07392 \times .07 \times .4 = .002070$
- From (C, V) to (V, V) : $\alpha = .07392 \times .07 \times .4 = .002070$
- From (C, V) to (V, C) : $\alpha_{C,C} = .03528 \times 1 \times .4 = .01411$
 $\alpha_{V,C} = .056 \times .8 \times .4 = .01792 = \alpha_{\max}$
- From (V, V) to (V, C) : $\alpha_{V,C} = .056 \times .8 \times .4 = .01792 = \alpha_{\max}$
- From (V, C) to $(', V)$: $\alpha = .08 \times 1 \times .7 = .056$
- From $(', V)$ to $(', V)$: $\alpha = .4 \times .2 = .08$

n-best State Sequences

- Keep track of n best “back pointers”:
- Ex.: $n=2$:
Two “winners”:
VCV (best)
CCV (2nd best)

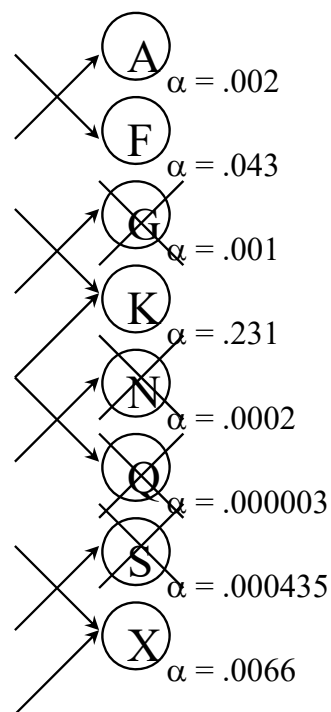


Tracking Back the n-best paths

- Backtracking-style algorithm:
 - **Start at the end, in the best of the n states (s_{best})**
 - **Put the other n-1 best nodes/back pointer pairs on stack, except those leading from s_{best} to the same best-back state.**
- Follow the back “beam” towards the start of the data, spitting out nodes on the way (backwards of course) using always only the best back pointer.
- At every beam split, push the diverging node/back pointer pairs onto the stack (node/beam width is sufficient!).
- When you reach the start of data, close the path, and pop the top-most node/back pointer(width) pair from the stack.
- Repeat until the stack is empty; expand the result tree if necessary.

Pruning

- Sometimes, too many trellis states in a stage:



- criteria:
- (a) $\alpha < \text{threshold}$
 - (b) $\Sigma\pi < \text{threshold}$
 - (c) # of states $> \text{threshold}$
(get rid of smallest α)

HMM Parameter Estimation: the Baum-Welch Algorithm

HMM: The Tasks

- HMM (the general case):
 - five-tuple (S, S_0, Y, P_S, P_Y) , where:
 - $S = \{s_1, s_2, \dots, s_T\}$ is the set of states, S_0 is the initial state,
 - $Y = \{y_1, y_2, \dots, y_V\}$ is the output alphabet,
 - $P_S(s_j | s_i)$ is the set of prob. distributions of transitions,
 - $P_Y(y_k | s_i, s_j)$ is the set of output (emission) probability distributions.
- Given an HMM & an output sequence $Y = \{y_1, y_2, \dots, y_k\}$:
 - ✓ (Task 1) compute the probability of Y ;
 - ✓ (Task 2) compute the most likely sequence of states which has generated Y .
 - (Task 3) Estimating the parameters (transition/output distributions)

A Variant of EM

- Idea (\sim EM, for another variant see LM smoothing):
 - Start with (possibly random) estimates of P_S and P_Y .
 - Compute (fractional) “counts” of state transitions/emissions taken, from P_S and P_Y , given data Y .
 - Adjust the estimates of P_S and P_Y from these “counts” (using the MLE, i.e. relative frequency as the estimate).
- Remarks:
 - many more parameters than the simple four-way smoothing
 - no proofs here; see Jelinek, Chapter 9

Setting

- HMM (without P_S, P_Y) (S, S_0, Y), and data $T = \{y^i \in Y\}_{i=1..|T|}$
 - **will use $T \sim |T|$**
 - HMM structure is given: (S, S_0)
 - P_S : Typically, one wants to allow “fully connected” graph
 - **(i.e. no transitions forbidden ~ no transitions set to hard 0)**
 - **why? → we better leave it on the learning phase, based on the data!**
 - **sometimes possible to remove some transitions ahead of time**
 - P_Y : should be restricted (if not, we will not get anywhere!)
 - **restricted ~ hard 0 probabilities of $p(y|s, s')$**
 - **“Dictionary”: states \leftrightarrow words, “m:n” mapping on $S \times Y$ (in general)**

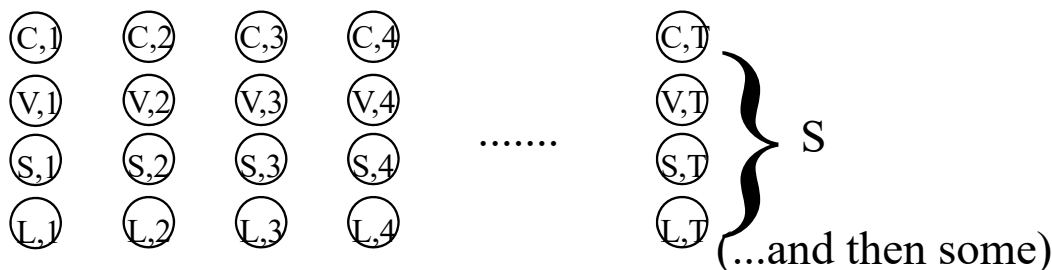
Initialization

- For computing the initial expected “counts”
- Important part
 - EM guaranteed to find a *local* maximum only (albeit a good one in most cases)
- P_Y initialization more important
 - fortunately, often easy to determine
 - **together with dictionary \leftrightarrow vocabulary mapping, get counts, then MLE**
- P_S initialization less important
 - e.g. uniform distribution for each $p(.|s)$

Data Structures

- Will need storage for:
 - The predetermined structure of the HMM
(unless fully connected → need not to keep it!)
 - The parameters to be estimated (P_S, P_Y)
 - The expected counts (same size as P_S, P_Y)
 - The training data $T = \{y^i \in Y\}_{i=1..T}$
 - The trellis (if f.c.): $\uparrow T$ Size: $T' \times S$ (Precisely, $|T'| \times |S|$)

Each trellis state:
two [float] numbers
(forward/backward)



The Algorithm Part I

1. Initialize P_S, P_Y

2. Compute “forward” probabilities:

- follow the procedure for trellis (summing), compute $\alpha(s,i)$
- use the current values of P_S, P_Y ($p(s'|s), p(y|s,s')$):

$$\alpha(s',i) = \sum_{s \rightarrow s'} \alpha(s,i-1) \times p(s'|s) \times p(y_i|s,s')$$

- NB: do not throw away the previous stage!

3. Compute “backward” probabilities

- start at all nodes of the last stage, proceed backwards, $\beta(s,i)$
- i.e., probability of the “tail” of data from stage i to the end of data

$$\beta(s',i) = \sum_{s \leftarrow s'} \beta(s,i+1) \times p(s|s') \times p(y_{i+1}|s',s)$$

- also, keep the $\beta(s,i)$ at all trellis states

The Algorithm Part II

4. Collect counts:

- for each output/transition pair compute

$$c(y, s, s') = \sum_{i=0..k-1, y=y_{i+1}} \alpha(s, i) \underbrace{p(s'|s) p(y_{i+1}|s, s')}_{\text{this transition prob ' output prob}} \beta(s', i+1)$$

one pass through data, only stop at (output) y prefix prob. tail prob

$$c(s, s') = \sum_{y \in Y} c(y, s, s') \text{ (assuming all observed } y_i \text{ in } Y)$$

$$c(s) = \sum_{s' \in S} c(s, s')$$

5. Reestimate: $p'(s'|s) = c(s, s')/c(s)$ $p'(y|s, s') = c(y, s, s')/c(s, s')$

6. Repeat 2-5 until desired convergence limit is reached.

Baum-Welch: Tips & Tricks

- Normalization badly needed
 - long training data \rightarrow extremely small probabilities
- Normalize α, β using the same norm. factor:

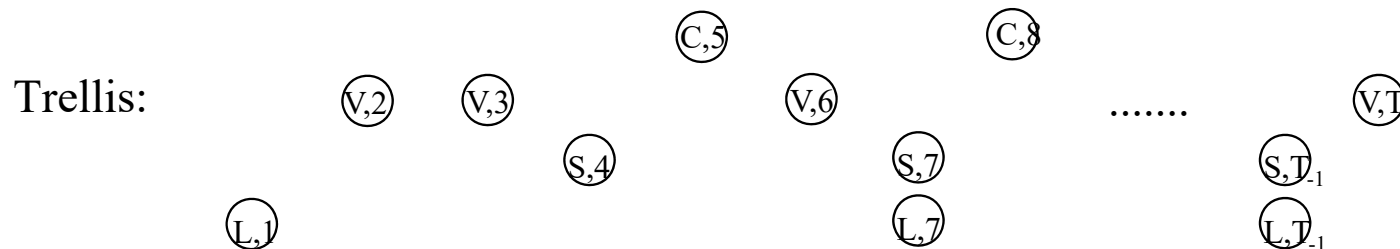
$$N(i) = \sum_{s \in S} \alpha(s, i)$$

as follows:

- **compute $\alpha(s, i)$ as usual (Step 2 of the algorithm), computing the sum $N(i)$ at the given stage i as you go.**
- **at the end of each stage, recompute all α s (for each state s):**
 - $\alpha^*(s, i) = \alpha(s, i) / N(i)$
- **use the same $N(i)$ for β s at the end of each backward (Step 3) stage:**
 - $\beta^*(s, i) = \beta(s, i) / N(i)$

Example

- Task: pronunciation of “the”
- Solution: build HMM, fully connected, 4 states:
 - S - short article, L - long article, C,V - starting w/consonant, vowel
 - thus, only “the” is ambiguous (a, an, the - not members of C,V)
- Output from states only ($p(w|s,s') = p(w|s')$)
- Data Y: an egg and a piece of the big the end



Example: Initialization

- Output probabilities:
 $p_{\text{init}}(w|c) = c(c,w) / c(c)$; where $c(S,\text{the}) = c(L,\text{the}) = c(\text{the})/2$
(other than that, everything is deterministic)
- Transition probabilities:
 - $p_{\text{init}}(c'|c) = 1/4$ (uniform)
- Don't forget:
 - about the space needed
 - initialize $\alpha(X,0) = 1$ (X : the never-occurring front buffer st.)
 - initialize $\beta(s,T) = 1$ for all s (except for $s = X$)

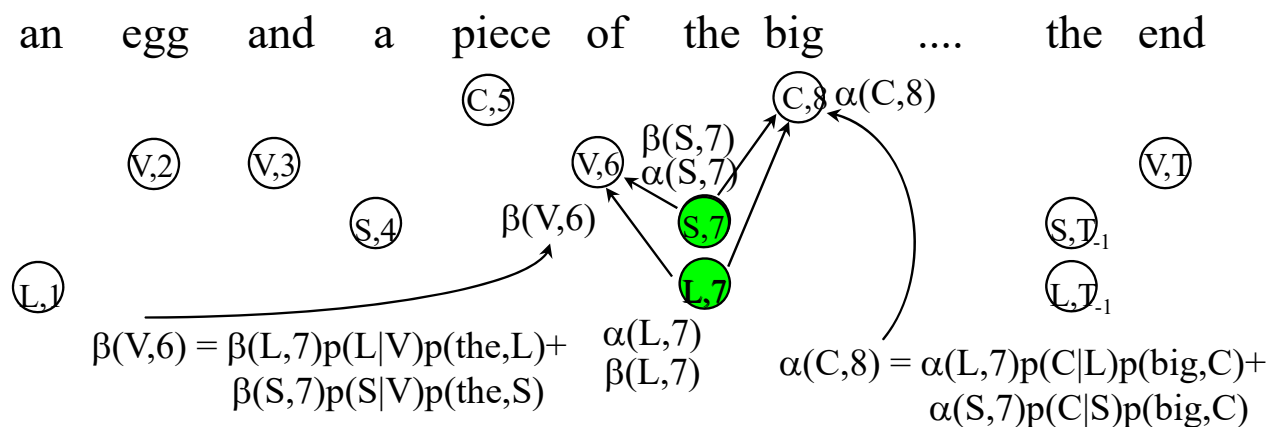
Fill in alpha, beta

- Left to right, alpha:

$$\alpha(s', i) = \sum_{s \rightarrow s'} \alpha(s, i-1) \times p(s'|s) \times p(w_i|s')$$

- Remember normalization (N(i)).
- Similarly, beta (on the way back from the end).

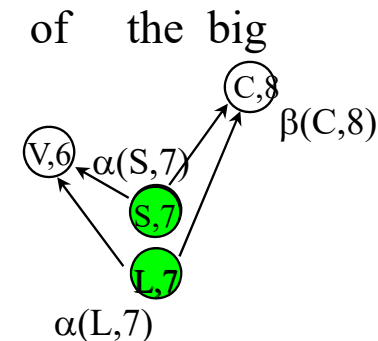
output from states



Counts & Reestimation

- One pass through data
 - At each position i , go through all pairs (s_i, s_{i+1})
 - Increment appropriate counters by frac. counts (Step 4):
 - $\text{inc}(y_{i+1}, s_i, s_{i+1}) = a(s_i, i) p(s_{i+1}|s_i) p(y_{i+1}|s_{i+1}) b(s_{i+1}, i+1)$
 - $c(y, s_i, s_{i+1}) += \text{inc}$ (for y at pos $i+1$)
 - $c(s_i, s_{i+1}) += \text{inc}$ (always)
 - $c(s_i) += \text{inc}$ (always)
- $$\text{inc}(\text{big}, L, C) = \alpha(L, 7) p(C|L) p(\text{big}, C) \beta(C, 8)$$

$$\text{inc}(\text{big}, S, C) = \alpha(S, 7) p(C|S) p(\text{big}, C) \beta(C, 8)$$
- Reestimate $p(s'|s)$, $p(y|s)$
 - and hope for increase in $p(C|S)$ and $p(V|L)$...!!



HMM: Final Remarks

- Parameter “tying”:
 - keep certain parameters same (\sim just one “counter” for all of them)
 - any combination in principle possible
 - ex.: smoothing (just one set of lambdas)
- Real Numbers Output
 - Y of infinite size (\mathbb{R} , \mathbb{R}^n):
 - **parametric (typically: few) distribution needed (e.g., “Gaussian”)**
- “Empty” transitions: do not generate output
 - **\sim vertical arcs in trellis; do not use in “counting”**