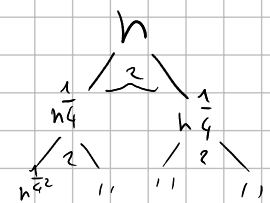


$$T(n) = \begin{cases} 1, & \text{se } n \leq 2; \\ 2 \cdot T(\sqrt[n]{n}) + \log(2n), & \text{altrimenti.} \end{cases}$$

LIV	INPUT		CONT. MOD6	RAMI	TOT
0	$n$		$\log(2n)$	1	$\log(2n)$
1	$n^{\frac{1}{4}}$		$\log(2n^{\frac{1}{4}})$	$2^1$	$2 \log(2n^{\frac{1}{4}})$
2	$n^{\frac{1}{4^2}}$		$\log(2n^{\frac{1}{4^2}})$	$2^2$	$2^2 (\log(2n^{\frac{1}{4^2}}))$
$\vdots$					
i	$n^{\frac{1}{4^i}}$		$\log(2n^{\frac{1}{4^i}})$	$2^i$	$2^i (\log(2n^{\frac{1}{4^i}}))$
					complessita

CASO BASE

$$K \cdot \frac{1}{n^{\frac{1}{4^K}}} \leq 2$$

$$\frac{1}{4^K} \log_2 n \leq 1 = \log_2 n \leq 4^K = \log_4 \log_2 n = K$$

COMPLESSITA  
CASO K

$$2^i \log(2n^{\frac{1}{4^i}}) = 2^i \log 2 + 2^i \log n^{\frac{1}{4^i}} = \underbrace{(2^i \log 2)}_{\text{certific}} + \frac{2^i}{4^i} \log n = \log n \sum_{i=0}^K \frac{1}{2} = \log n \frac{1}{\frac{1}{2}} = \log n^2 = \Theta \log n$$

#LIVELLI