Linked Data Structures I: Singly-Linked Lists









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 - Similarly for the Java collections framework

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ramework

Any type whose abstract mathematical model involves a *string* or *set* or *multiset* (or *tree* or *binary tree*?) is a "collection" type.

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ramework

Examples are Queue, Stack, Sequence, Set, Map, SortingMachine, ...

- A Java array is not ideally suited as a data representation of the various collection types
 - Similarly for the Java collections framework

This is part of the package java.util, and includes many types similar to the OSU CSE components.

Collection Terminology

- Fixed size means the size/length of a collection is "inflexible", i.e., it is determined at initialization of the collection and cannot be incrementally adjusted
 - A classical synonym is static; this term unfortunately means other things in Java
- Dynamic means the size/length of a collection is "flexible", i.e., it can be incrementally adjusted by "adding" and "removing" entries, even from the middle

Collection Terminology

- Direct access means the entries of a collection (typically with a string model) may be accessed by providing an int position/index of any entry in the collection
 - A classical but unfortunate synonym is random access; nothing random about it!
- Sequential access means the entries of a collection (with a string model) may be accessed in increasing order of position by accessing the "next" entry in the collection

Collection Te

 Direct access means collection (typically wit may be accessed by position/index of accessed We might say any collection with an iterator allows sequential access, but this is about the *other* methods for access.

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Key Pros and Cons of Arrays

• Pros:

 Direct access is fast, i.e., it takes constant time independent of the length of the array

Cons:

- Its fixed size limits array utility where dynamic size is important: it can run out of room
- Adding and removing entries in the middle requires moving array entries, which is slow
- Initialization may be expensive, especially if many entries are not subsequently used

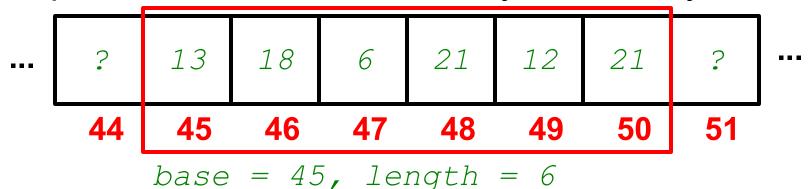
Fixed Size Can Support Fast Direct Access

A Java array is represented in a
 contiguous block of memory locations
 with consecutive memory addresses (IDs),
 so the memory address of the entry at
 index i can be directly calculated from
 the memory address of the first entry, by
 using simple arithmetic

Example

Client's view of an array

Implementer's view of an array in memory:



Examp

Client's view of an array

$$(entries = <13, 18,$$

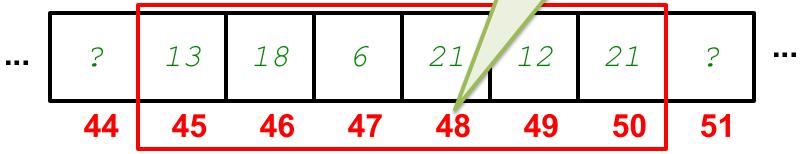


length = 6

If client wants to access the entry at position 3 of the array, how does implementer compute its memory address/ID?

Implementer's view of an array

memory:



$$base = 45$$
, $length = 6$

Examp

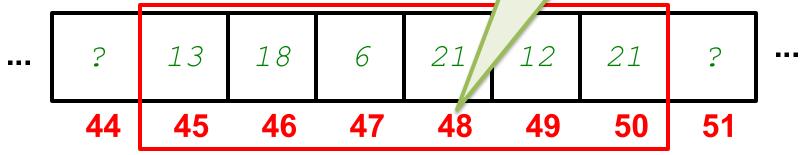
Client's view of an array

Every modern computer the JVM runs on provides constant-time access to any memory location given the memory address/ID.

Implementer's view of an array

length = 6

memory:



$$base = 45$$
, $length = 6$

Notice the Array Mismatches

Collection	Fixed Size?	Dynamic?	Direct Access?	Seq Access?
array	~		~	
Queue		✓		✓
Stack		✓		✓
Sequence		✓	✓	
Set		✓		
Мар		✓		
Sorting- Machine		✓		

What Can Be Done?

- To represent collections that are dynamic with sequential access, a different approach is needed: not arrays
 - Note: It is an open problem to represent a Sequence, which is dynamic and offers direct access, in a way that is efficient in both execution time of access operations (i.e., constant time) and memory "footprint" (i.e., constant factor overhead)

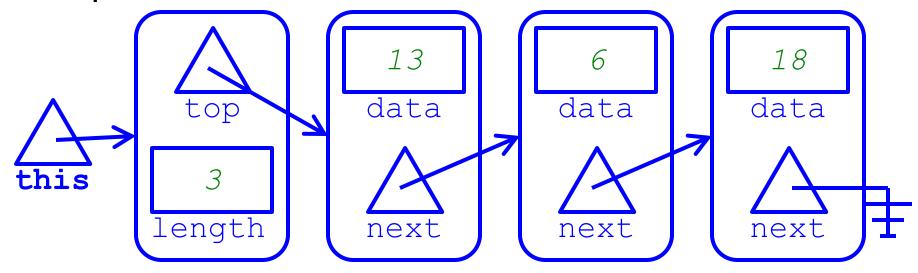
Dynamic Can Support Fast Sequential Access

- If we want a *dynamic* collection, then we should give up on storing all entries of the collection in contiguous memory locations
- If we want fast sequential access, then we should give up on fast direct access
 - Instead, for every entry in the collection, wherever it is in memory, simply keep a reference to (i.e., memory location of) the "next" entry

Example: Stack2

Client's view of a Stack:

Implementer's view of a Stack2:

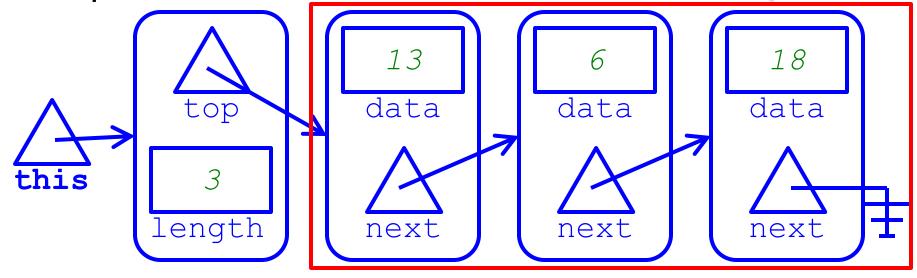


Example:

This is called a **singly- linked-list** data structure.

Client's view of a Stack:

Implementer's view of a Stack2:

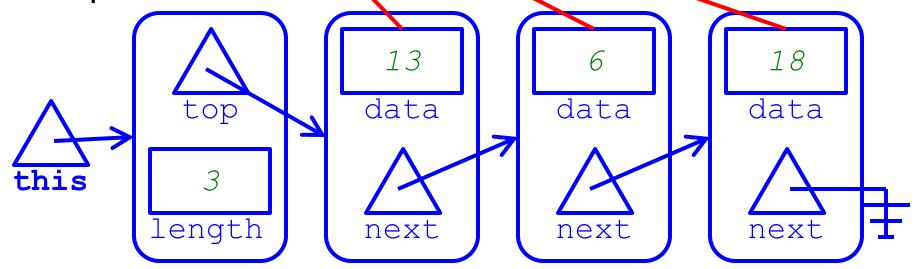


Example:

The abstraction function (correspondence) ...

Client's view of a Stack:

Implementer's view of a Stack2:



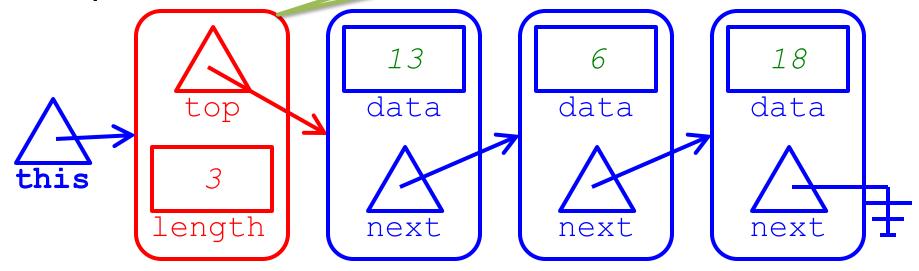
Example

Client's view of a Stacl

this = <13, 6,

The *instance variables*(*fields*) of the data
representation for Stack2
are shown here.

Implementer's view : a Stack2:

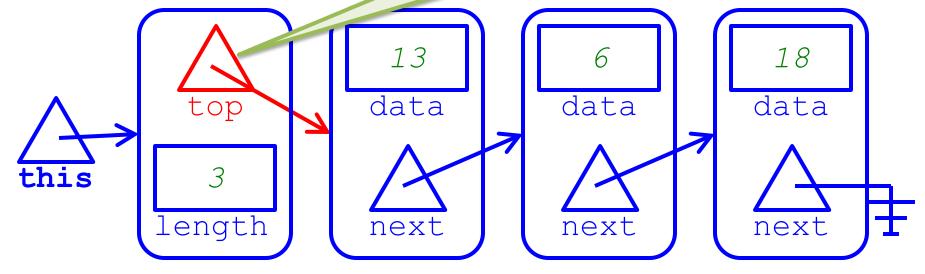


Example

Client's view of a Stacl this = <13, 6,

The Stack methods only require access to the first entry, i.e., the top, so we keep a reference to its node.

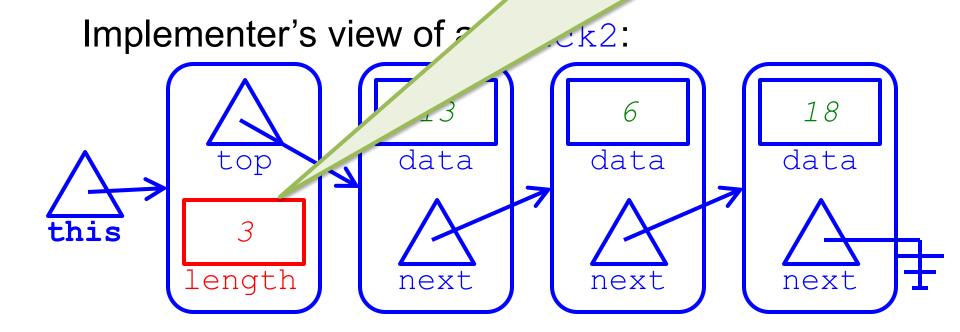
Implementer's view of __cack2:



Examp

Client's view of a State this = <13, 6,

The Stack methods include length, so we keep this direct count of the number of nodes in the linked list.

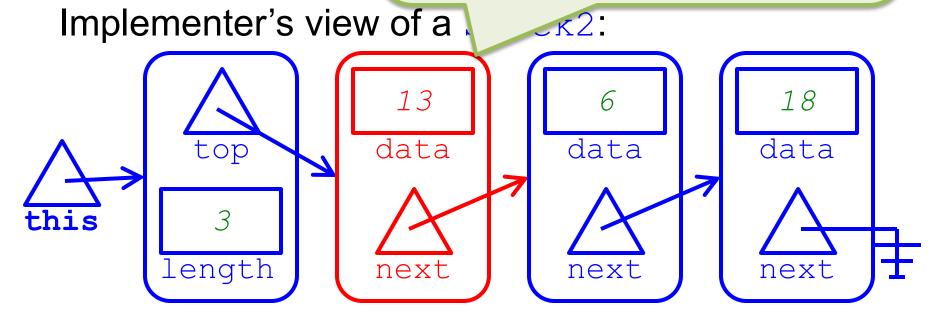


Exam

Client's view of a S

this = <13,

Each of these objects (a pair of variables) is called a *node* in this *singly-linked-list data structure*: a variable of type T, and a reference to the "next" node.



Declaration of Node Class

 A Node class is declared as a nested class inside the kernel implementation that uses it in a data representation (e.g., Stack2, Queue2, and similar classes) private final class Node { private T data; private Node next;

Declaratio

A Node class is
 class inside the
 that uses it in a data

```
This declaration is recursive, and may seem circular!
One of the instance variables of a Node is of type Node.
How can this work?
```

entation (e.g., هناستن

```
private final class Node
  private T data;
  private Node next;
}
```

Stack2, Queue2, an

It works because the instance variable next is a reference variable, hence is a reference to a Node object rather than a "nested" object.

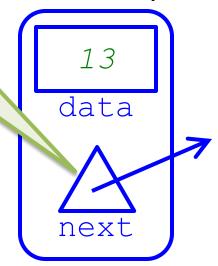
that uses it in a

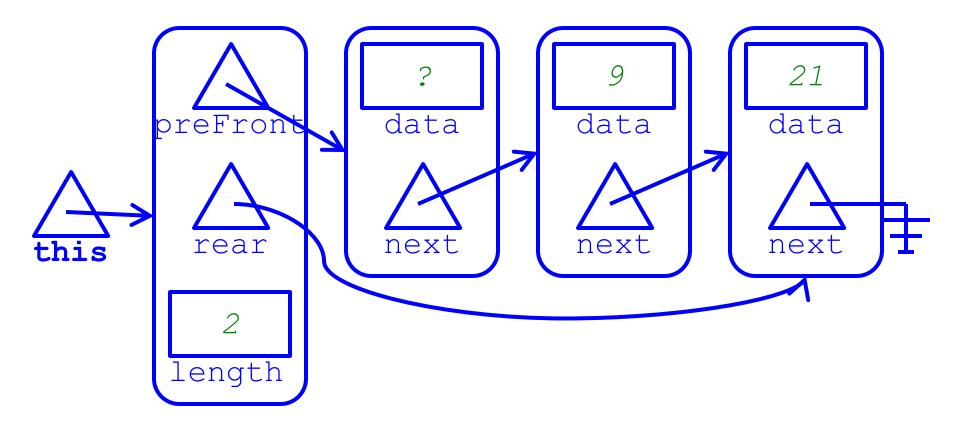
Stack2, Queue2, &

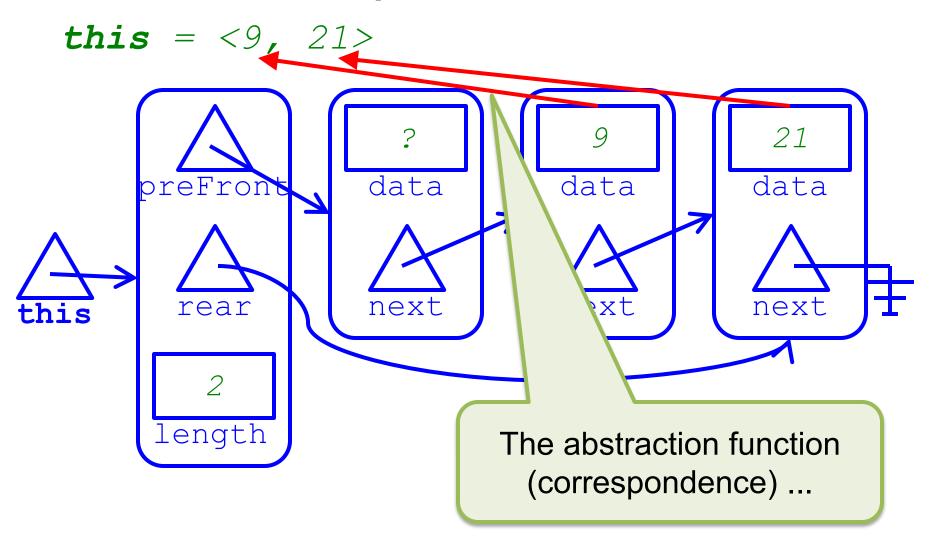
```
private final class No
  private T data;
  private Node next;
```

bde Class

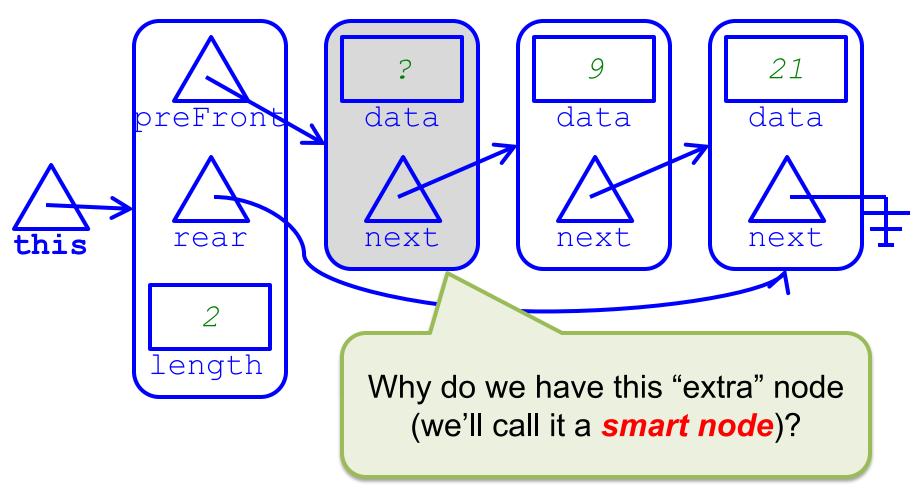
as a **nested**hplementation
presentation (e.g.,
imilar classes)

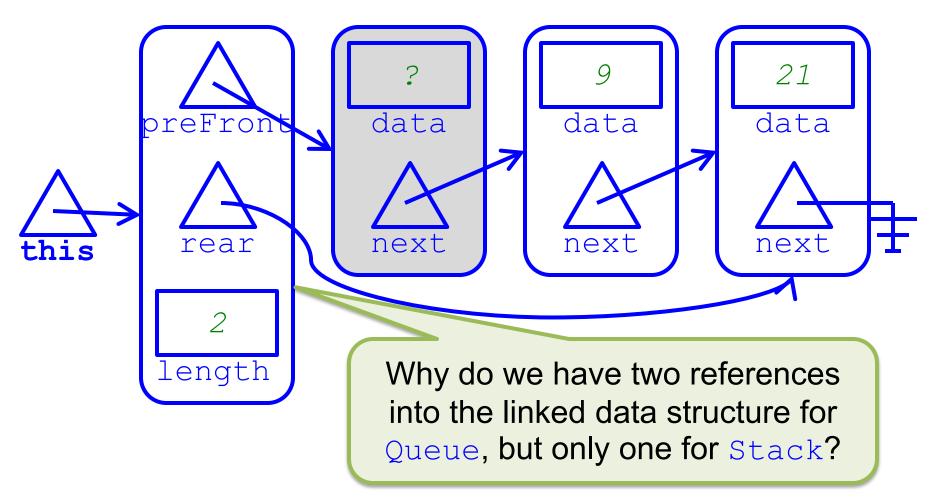






$$this = <9, 21>$$

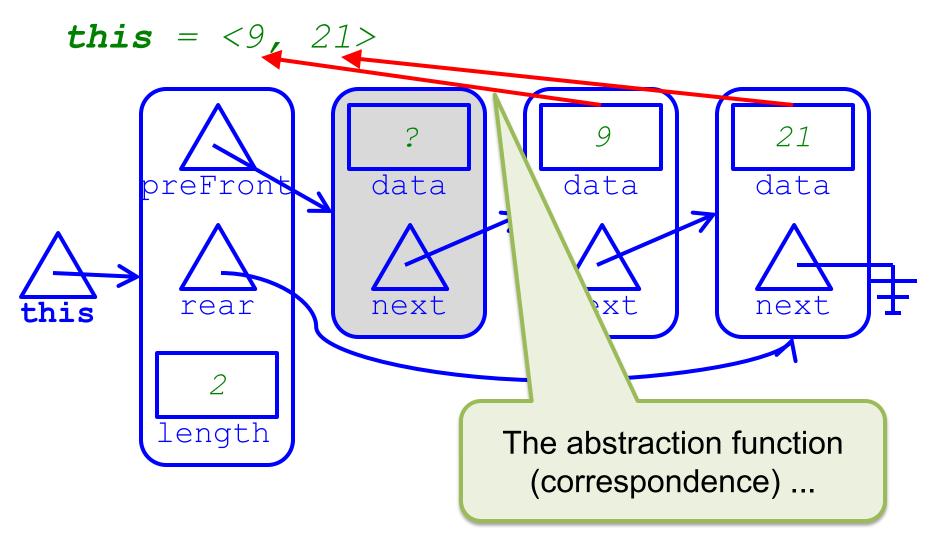


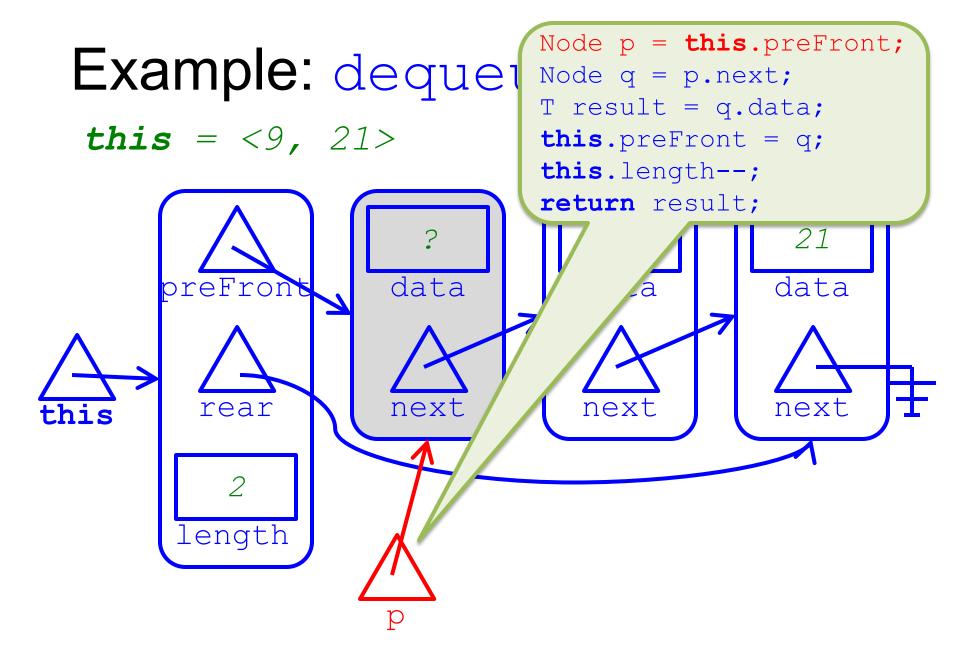


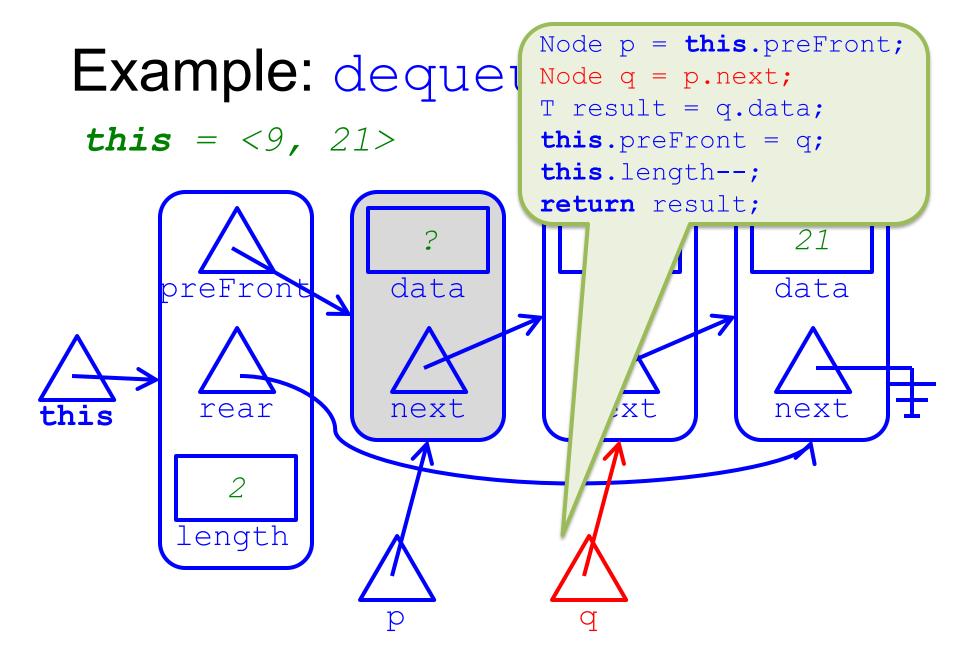
Example: dequeue for Queue2

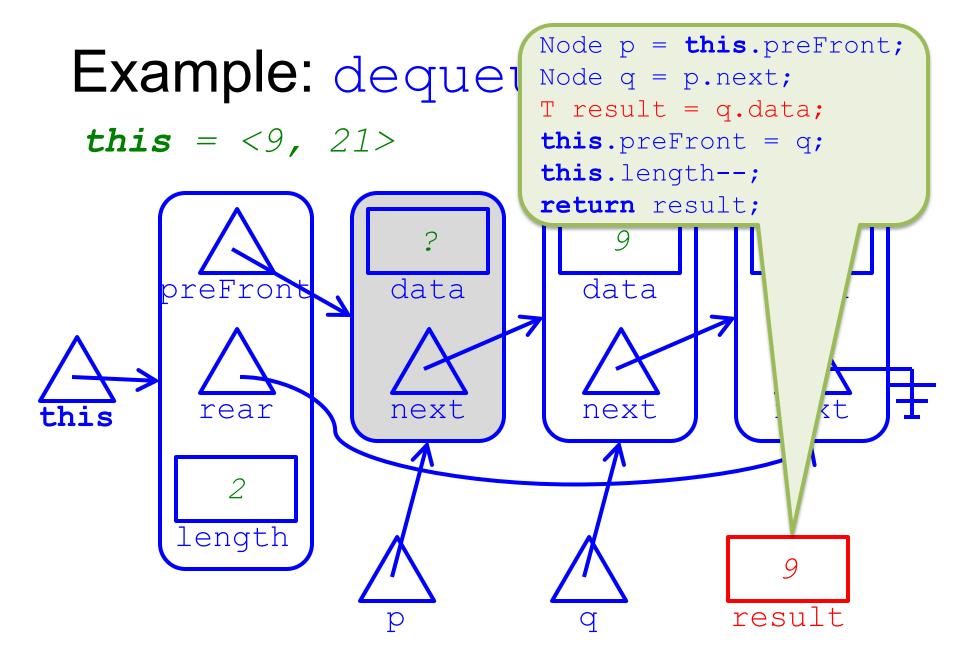
```
public final T dequeue() {
  Node p = this.preFront;
  Node q = p.next;
  T result = q.data;
  this.preFront = q;
  this.length--;
  return result;
```

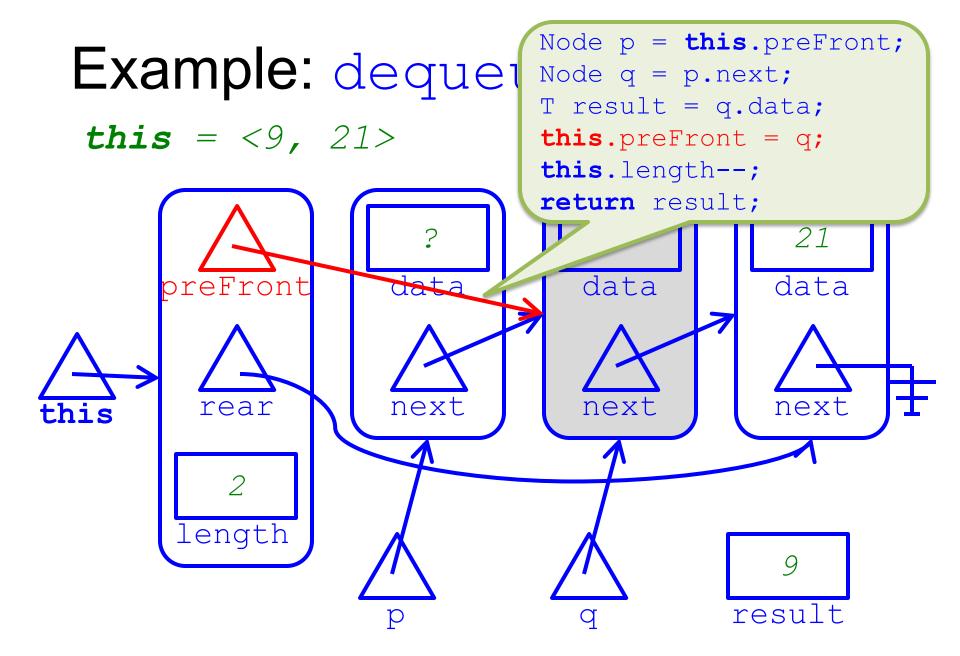
Example: dequeue for Queue2



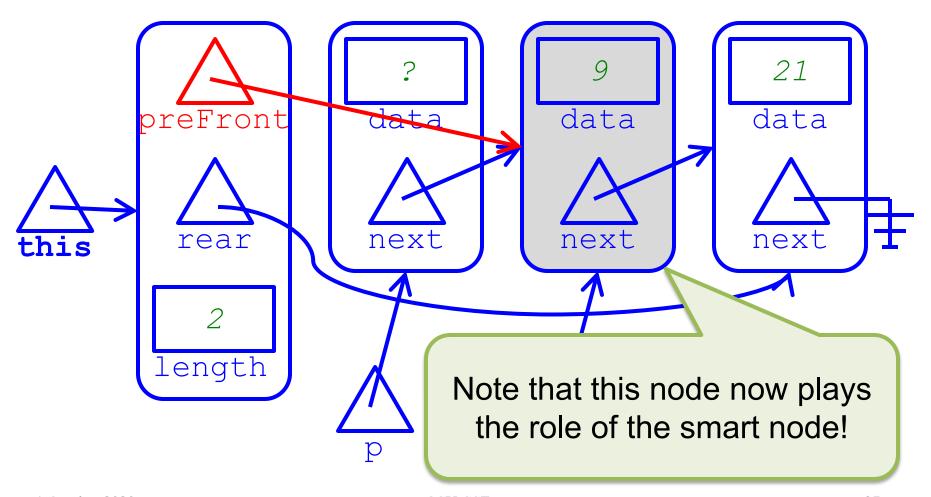


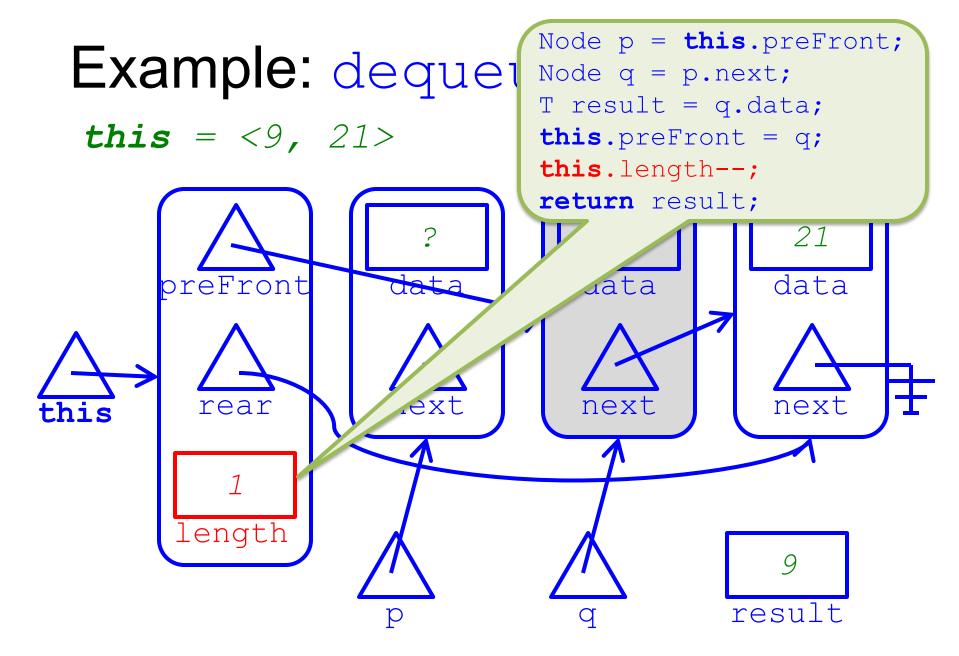


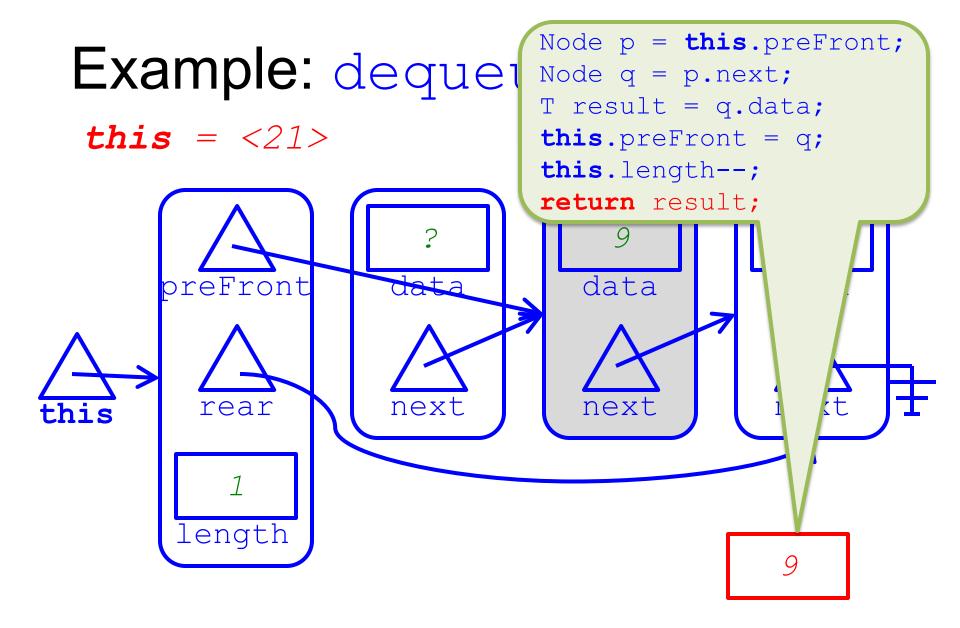


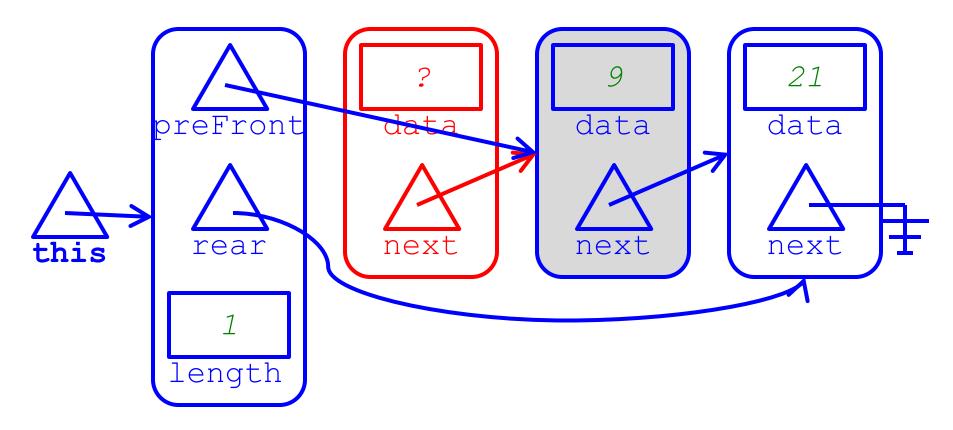


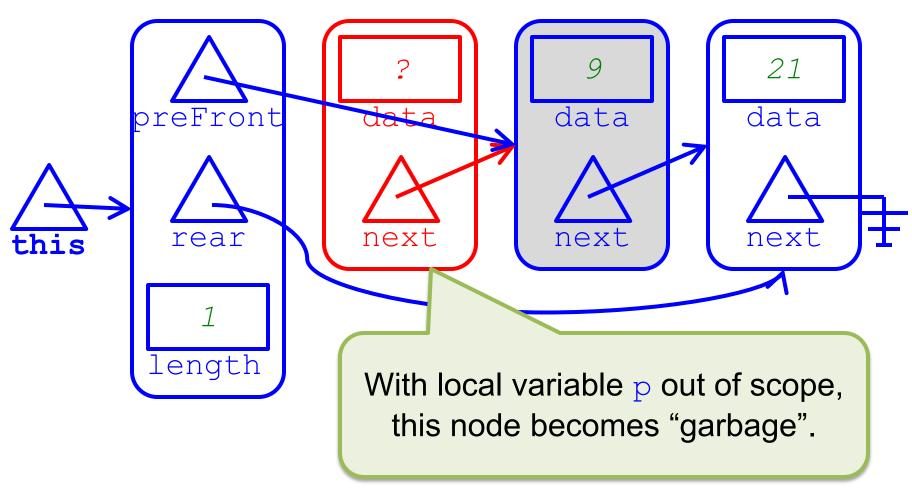
$$this = <9, 21>$$

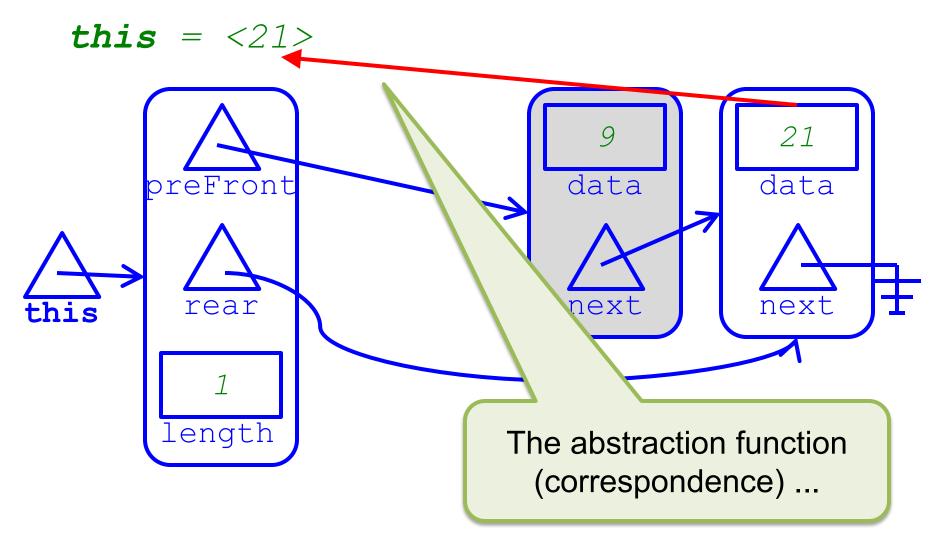




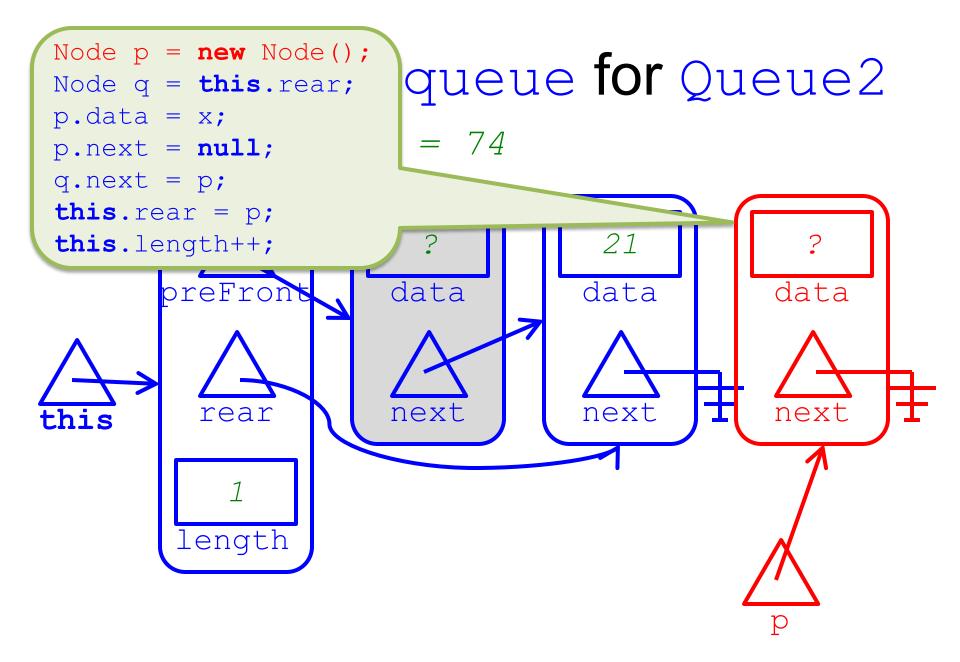


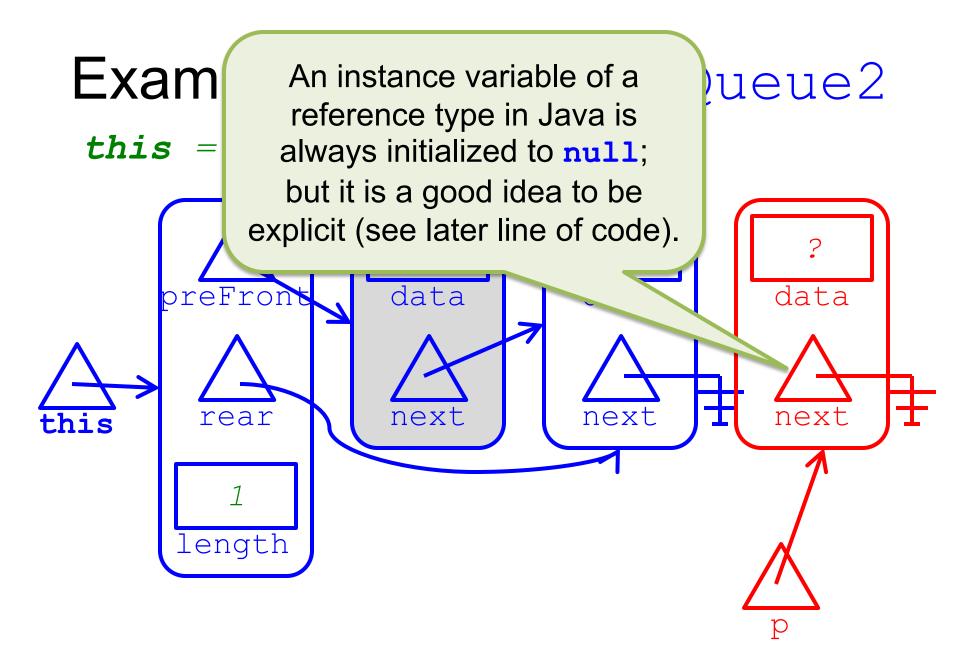


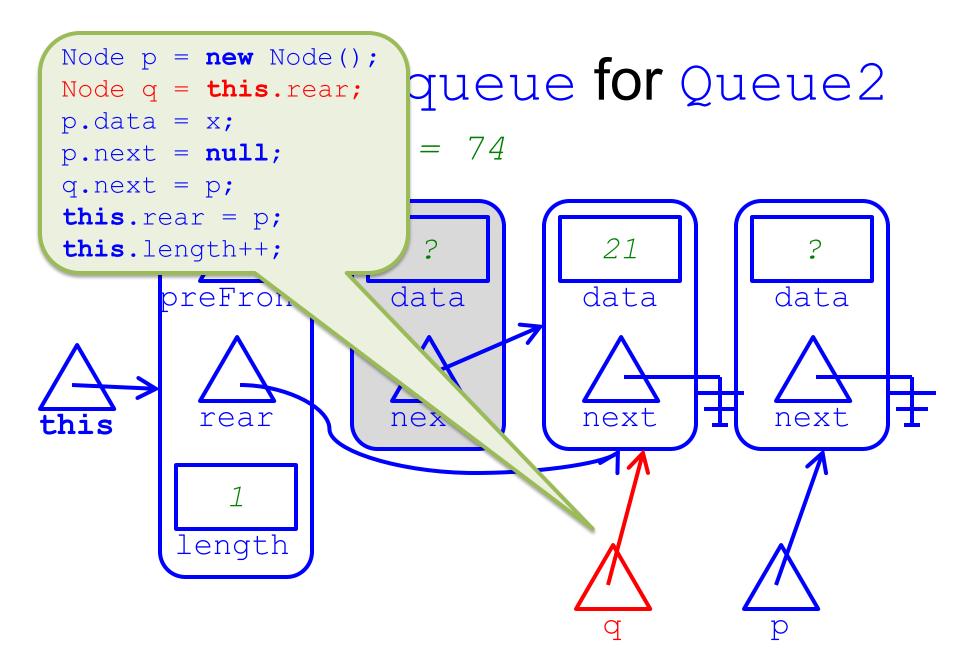


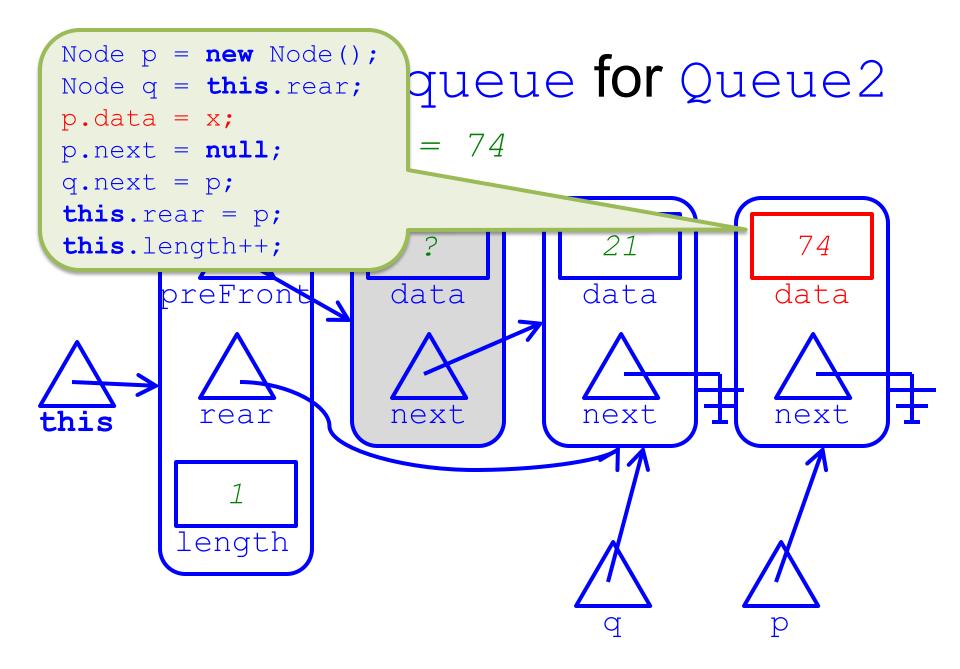


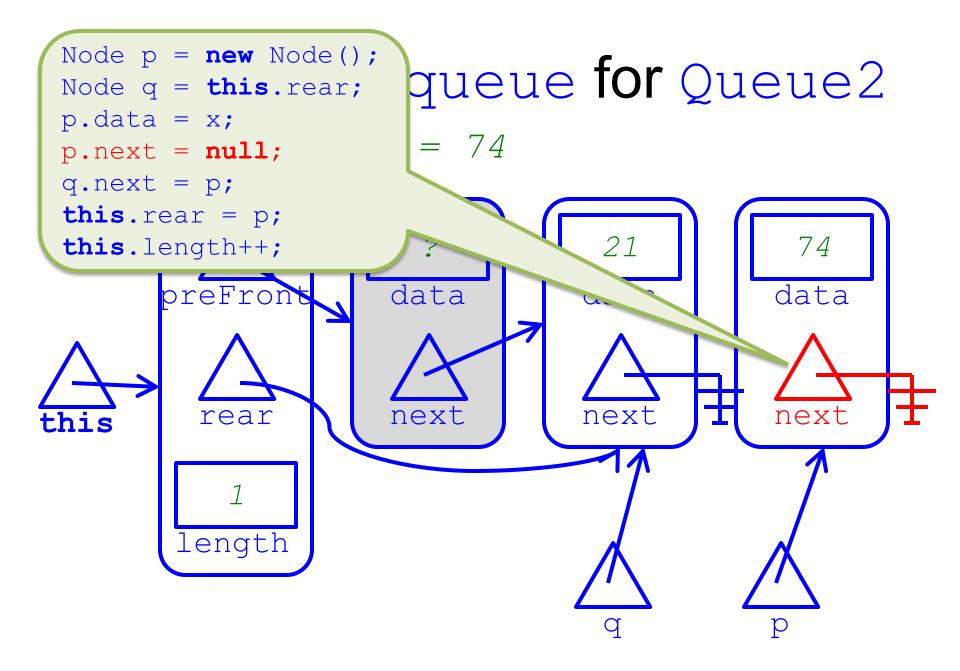
```
public final void enqueue(T x) {
  Node p = new Node();
  Node q = this.rear;
  p.data = x;
  p.next = null;
  q.next = p;
  this.rear = p;
  this.length++;
```

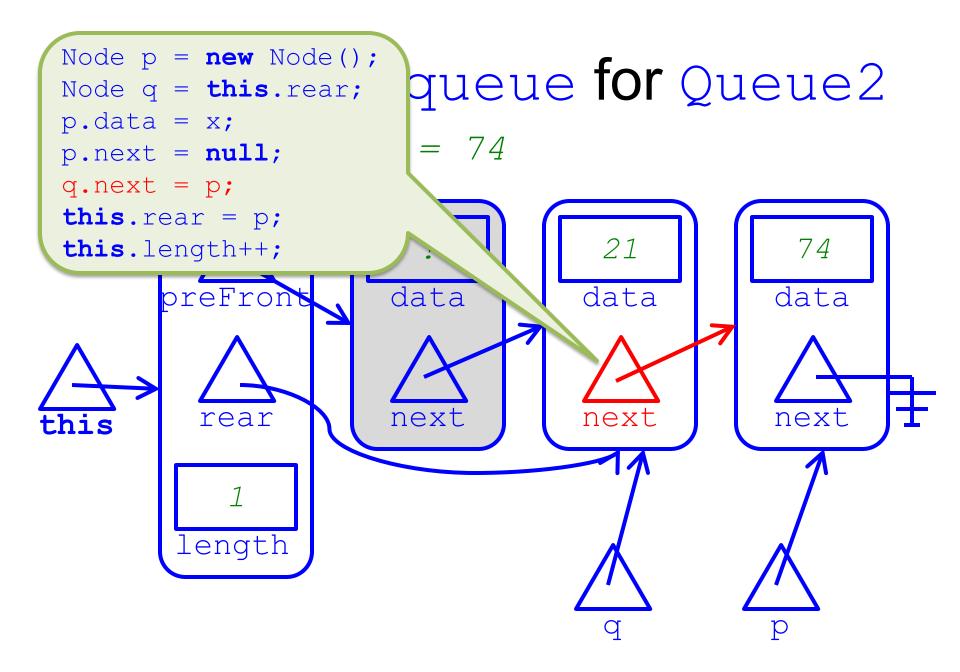


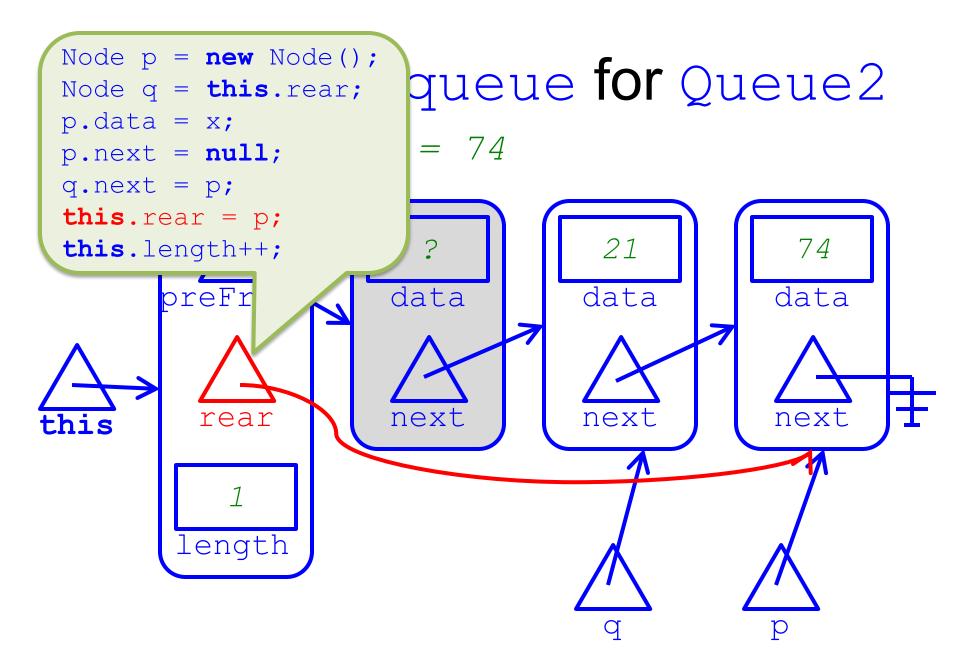


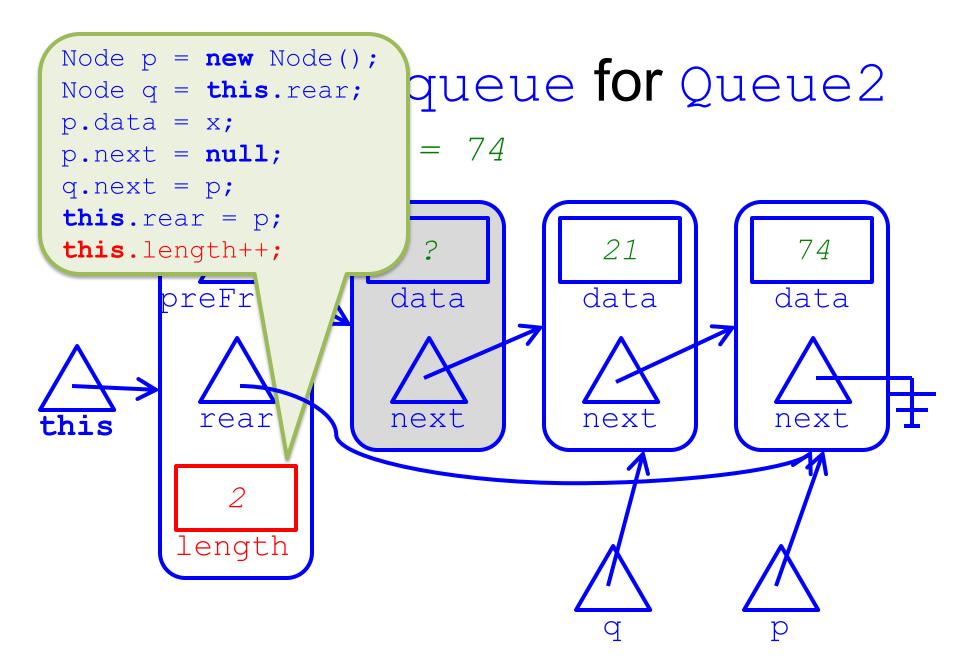




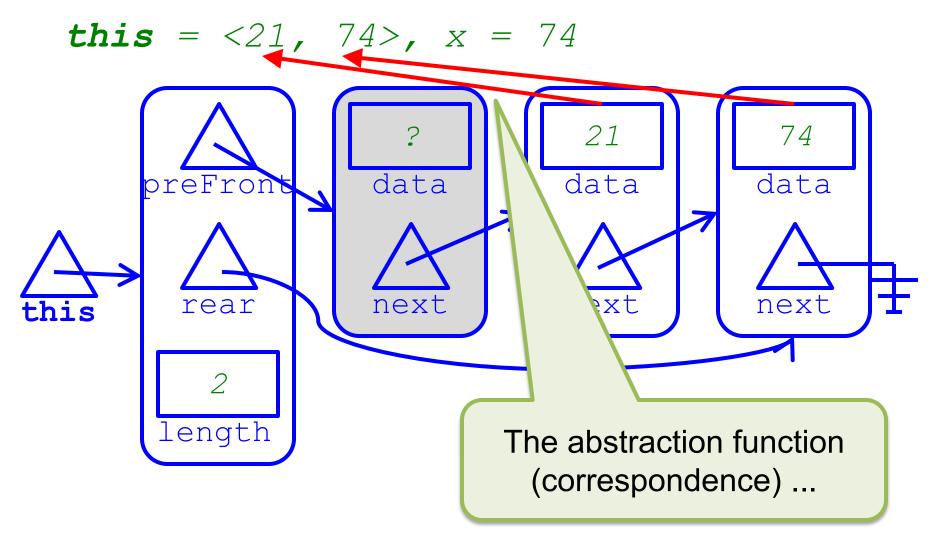








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The "Smart" Node

- To see why we want the extra node at the beginning of the linked list, write the code for enqueue without it (but be careful)
 - You should be able to see why it's a smart node rather than a dummy node
- Why is the smart node helpful in the representation of a Queue, but not of a Stack?

Resources

- Wikipedia: Linked Data Structure
 - http://en.wikipedia.org/wiki/Linked_data_structure
- Big Java (4th ed), Sections 15.1, 15.2 (but not the part about iterators)
 - https://library.ohio-state.edu/record=b8540788~S7