# Lecture 11: Krylov subspace, Generalized minimal residual method



School of Mathematical Sciences, Xiamen University

## 1. Krylov subspace

• Given  $\mathbf{A} \in \mathbb{C}^{m \times m}$  and nonzero  $\mathbf{r} \in \mathbb{C}^m$ , the jth Krylov subspace generated by  $\mathbf{A}$  and  $\mathbf{r}$  is defined by

$$\mathcal{K}_j(\mathbf{A}, \mathbf{r}) := \operatorname{span}\{\mathbf{r}, \mathbf{A}\mathbf{r}, \mathbf{A}^2\mathbf{r}, \cdots, \mathbf{A}^{j-1}\mathbf{r}\}.$$

Obviously,  $\mathcal{K}_j(\mathbf{A}, \mathbf{r}) \subseteq \mathcal{K}_{j+1}(\mathbf{A}, \mathbf{r})$  and  $\dim \mathcal{K}_j(\mathbf{A}, \mathbf{r}) \leq j$ .

# Proposition 1

Let  $\mathbb{P}_j$  denote the set of polynomials of degree  $\leq j$ . Then

$$\mathcal{K}_j(\mathbf{A}, \mathbf{r}) = \{ p(\mathbf{A})\mathbf{r} \mid p \in \mathbb{P}_{j-1} \}.$$

# Proposition 2

If the minimal polynomial of the matrix **A** has degree n, then for any j > n and any nonzero  $\mathbf{r} \in \mathbb{C}^m$ ,

$$\dim \mathcal{K}_i(\mathbf{A}, \mathbf{r}) \leq n.$$

### 1.1. Arnoldi process

## Algorithm: Arnoldi process generating orthonormal basis

Given 
$$\mathbf{A} \in \mathbb{C}^{m \times m}$$
 and nonzero  $\mathbf{r} \in \mathbb{C}^m$ 
 $\mathbf{q}_1 = \mathbf{r}/\|\mathbf{r}\|_2$ 
for  $j = 1, 2, 3, \ldots$ ,
 $\mathbf{v} = \mathbf{A}\mathbf{q}_j$ 
for  $i = 1$  to  $j$ 
 $h_{ij} = \langle \mathbf{v}, \mathbf{q}_i \rangle = \mathbf{q}_i^* \mathbf{v}$ 
 $\mathbf{v} = \mathbf{v} - h_{ij} \mathbf{q}_i$ 
end
 $h_{j+1,j} = \|\mathbf{v}\|_2$ 
 $\mathbf{q}_{j+1} = \mathbf{v}/h_{j+1,j}$ 
end

• At the end of step j, we obtain

$$\mathbf{v} = (\mathbf{I} - \mathbf{q}_j \mathbf{q}_j^*) \cdots (\mathbf{I} - \mathbf{q}_2 \mathbf{q}_2^*) (\mathbf{I} - \mathbf{q}_1 \mathbf{q}_1^*) \mathbf{A} \mathbf{q}_j.$$

• We call the Arnoldi process breaks down at step k if  $h_{k+1,k} = 0$ .

#### Remark 3

The Arnoldi process is the modified Gram-Schimidt orthogonalization applied to  $\{\mathbf{r}, \mathbf{Aq}_1, \mathbf{Aq}_2, \cdots\}$ . We have the Arnoldi relation

$$\mathbf{A} \begin{bmatrix} \mathbf{q}_1 & \cdots & \mathbf{q}_j \end{bmatrix} = \begin{bmatrix} \mathbf{q}_1 & \cdots & \mathbf{q}_{j+1} \end{bmatrix} \begin{bmatrix} h_{11} & \cdots & h_{1j} \\ h_{21} & \ddots & \vdots \\ & \ddots & h_{jj} \\ & & h_{j+1,j} \end{bmatrix},$$

that is  $\mathbf{AQ}_j = \mathbf{Q}_{j+1}\widetilde{\mathbf{H}}_j$ . ( $\widetilde{\mathbf{H}}_j$  is an upper Hessenberg matrix.) Let

$$\mathbf{H}_j := \mathbf{Q}_j^* \mathbf{A} \mathbf{Q}_j = \mathbf{Q}_j^* \mathbf{Q}_{j+1} \widetilde{\mathbf{H}}_j = \begin{bmatrix} h_{11} & \cdots & h_{1j} \\ h_{21} & \ddots & \vdots \\ & \ddots & h_{jj} \end{bmatrix}.$$

If the Arnoldi process breaks down at step k, then  $\mathbf{AQ}_k = \mathbf{Q}_k \mathbf{H}_k$ .

#### Theorem 4

Suppose the Arnoldi process breaks down at step k. Then,

$$span{\mathbf{q}_1, \mathbf{q}_2, \cdots, \mathbf{q}_j} = \mathcal{K}_j(\mathbf{A}, \mathbf{r}), \quad j = 1, 2, \dots, k,$$

and the set  $\{\mathbf{q}_1, \mathbf{q}_2, \cdots, \mathbf{q}_k\}$  is orthonormal.

# Corollary 5

The matrices  $\{\mathbf{Q}_j\}_{j=1}^k$  generated by the Arnoldi process are Q-factors of reduced QR factorizations of the Krylov matrices,

$$\mathbf{K}_j := \begin{bmatrix} \mathbf{r} & \mathbf{A}\mathbf{r} & \cdots & \mathbf{A}^{j-1}\mathbf{r} \end{bmatrix} = \mathbf{Q}_j\mathbf{R}_j, \quad j = 1, 2, \dots, k.$$

#### Remark 6

Both  $\mathbf{K}_j$  and  $\mathbf{R}_j$  are not formed explicitly in the Arnoldi process.

# 2. Generalized minimal residual method (GMRES)

• Idea of GMRES: Consider a nonsingular linear system

$$\mathbf{A}\mathbf{x} = \mathbf{b}, \quad \mathbf{A} \in \mathbb{C}^{m \times m}, \quad \mathbf{b} \in \mathbb{C}^m.$$

For any initial guess  $\mathbf{x}_0$ , at step j, GMRES finds the jth approximate solution

$$\mathbf{x}_j = \underset{\mathbf{x} \in \mathbf{x}_0 + \mathcal{K}_j(\mathbf{A}, \mathbf{r}_0)}{\operatorname{argmin}} \|\mathbf{b} - \mathbf{A}\mathbf{x}\|_2,$$

where  $\mathbf{r}_0 := \mathbf{b} - \mathbf{A}\mathbf{x}_0$  and

$$\mathcal{K}_j(\mathbf{A}, \mathbf{r}_0) = \operatorname{span}\{\mathbf{r}_0, \mathbf{A}\mathbf{r}_0, \dots, \mathbf{A}^{j-1}\mathbf{r}_0\}.$$

For the residual  $\mathbf{r}_i = \mathbf{b} - \mathbf{A}\mathbf{x}_i$ , we have

$$\|\mathbf{r}_j\|_2 = \min_{\mathbf{x} \in \mathbf{x}_0 + \mathcal{K}_j(\mathbf{A}, \mathbf{r}_0)} \|\mathbf{b} - \mathbf{A}\mathbf{x}\|_2 \quad \text{and} \quad \mathbf{r}_j \perp \mathbf{A}\mathcal{K}_j.$$

Assume that the Arnoldi process for the orthonormal basis of  $\mathcal{K}_j(\mathbf{A}, \mathbf{r}_0)$  breaks down at step k. For  $1 \leq j < k$ , we have

$$\begin{aligned} \|\mathbf{r}_{j}\|_{2} &= \min_{\mathbf{x} \in \mathbf{x}_{0} + \mathcal{K}_{j}(\mathbf{A}, \mathbf{r}_{0})} \|\mathbf{b} - \mathbf{A}\mathbf{x}\|_{2} = \min_{\mathbf{y} \in \mathbb{C}^{j}} \|\mathbf{r}_{0} - \mathbf{A}\mathbf{Q}_{j}\mathbf{y}\|_{2} \\ &= \min_{\mathbf{y} \in \mathbb{C}^{j}} \|\mathbf{r}_{0} - \mathbf{Q}_{j+1}\widetilde{\mathbf{H}}_{j}\mathbf{y}\|_{2} \quad \text{(by Arnoldi relation)} \\ &= \min_{\mathbf{y} \in \mathbb{C}^{j}} \|\|\mathbf{r}_{0}\|_{2}\mathbf{e}_{1} - \widetilde{\mathbf{H}}_{j}\mathbf{y}\|_{2}. \end{aligned}$$

For j = k, we have

$$\|\mathbf{r}_{k}\|_{2} = \min_{\mathbf{x} \in \mathbf{x}_{0} + \mathcal{K}_{k}(\mathbf{A}, \mathbf{r}_{0})} \|\mathbf{b} - \mathbf{A}\mathbf{x}\|_{2} = \min_{\mathbf{y} \in \mathbb{C}^{k}} \|\mathbf{r}_{0} - \mathbf{A}\mathbf{Q}_{k}\mathbf{y}\|_{2}$$
$$= \min_{\mathbf{y} \in \mathbb{C}^{k}} \|\mathbf{r}_{0} - \mathbf{Q}_{k}\mathbf{H}_{k}\mathbf{y}\|_{2} \quad \text{(by Arnoldi relation)}$$
$$= \min_{\mathbf{y} \in \mathbb{C}^{k}} \|\|\mathbf{r}_{0}\|_{2}\mathbf{e}_{1} - \mathbf{H}_{k}\mathbf{y}\|_{2}.$$

Once  $\mathbf{y}_i$  is found, set  $\mathbf{x}_i = \mathbf{x}_0 + \mathbf{Q}_i \mathbf{y}_i$ .

• The least squares problem about  $\mathbf{y}$  can be solved inexpensively with Givens rotations, exploiting the upper Hessenberg structure of  $\widetilde{\mathbf{H}}_j$ , costing just  $\mathcal{O}(j^2)$  or  $\mathcal{O}(j)$  instead of  $\mathcal{O}(j^3)$ .

## 2.1. Convergence of GMRES

#### Theorem 7

Assume that the Arnoldi process for the orthonormal basis of  $\mathcal{K}_j(\mathbf{A}, \mathbf{r}_0)$  breaks down at step k.

(1) For  $1 \le j < k$ , the residual  $\mathbf{r}_j$  satisfies  $(\mathbf{AQ}_j)^* \mathbf{r}_j = \mathbf{0}$ , i.e.,

$$\mathbf{r}_j \perp \mathbf{A} \mathcal{K}_j$$
.

(2) For  $0 \le j \le k$ , the residual  $\mathbf{r}_j$  satisfies

$$\|\mathbf{r}_0\|_2 \ge \|\mathbf{r}_1\|_2 \ge \dots \ge \|\mathbf{r}_{k-1}\|_2 > \|\mathbf{r}_k\|_2 = 0.$$

That is to say GMRES converges monotonically and finds the exact solution at step k.

#### Theorem 8

Suppose **A** is diagonalizable, i.e.,  $\mathbf{A} = \mathbf{V} \mathbf{\Lambda} \mathbf{V}^{-1}$  for some nonsingular matrix **V** and diagonal matrix **\Lambda**. At step j of the GMRES iteration, the residual  $\mathbf{r}_j$  satisfies

$$\frac{\|\mathbf{r}_j\|_2}{\|\mathbf{r}_0\|_2} \le \min_{p \in \mathbb{P}_j, p(0) = 1} \|p(\mathbf{A})\|_2 \le \kappa(\mathbf{V}) \min_{p \in \mathbb{P}_j, p(0) = 1} \max_{\lambda \in \Lambda(\mathbf{A})} |p(\lambda)|,$$

where  $\Lambda(\mathbf{A})$  is the set of eigenvalues of  $\mathbf{A}$ , and  $\kappa(\mathbf{V}) = \|\mathbf{V}\|_2 \|\mathbf{V}^{-1}\|_2$ .

• Y. Saad and M.H. Schultz

GMRES: A generalized minimal residual algorithm for solving nonsymmetric linear systems

SIAM J. Sci. Stat. Comput., 7: 856–869, 1986.

• Y. Saad

A Flexible Inner-Outer Preconditioned GMRES Algorithm SIAM J. Sci. Comput., 14: 461–469, 1993.

• Exercise: Assume the Arnoldi process for  $\{A, r_0\}$  breaks down at step k > 1. For  $1 \le j < k$ , we have the Arnoldi relation

$$\mathbf{A}\mathbf{Q}_j = \mathbf{Q}_{j+1}\widetilde{\mathbf{H}}_j.$$

For  $1 \le j < k$ , prove the following:

(a) The jth residual  $\mathbf{r}_i$  of GMRES can be uniquely expressed as

$$\mathbf{r}_j = p_j(\mathbf{A})\mathbf{r}_0, \qquad \deg(p_j) \le j, \qquad p_j(0) = 1.$$

(b) Let  $\mathbf{H}_j = \mathbf{Q}_j^* \mathbf{A} \mathbf{Q}_j$ . The unique polynomial  $p_j$  in (a) is given by

$$p_j(z) = \prod_{i=1}^{j} \left(1 - \theta_i^{(j)} z\right),$$

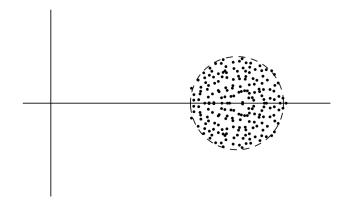
where  $\theta_i^{(j)}$ ,  $i=1,2,\ldots,j$ , are the eigenvalues of  $(\widetilde{\mathbf{H}}_j^*\widetilde{\mathbf{H}}_j)^{-1}\mathbf{H}_j^*$ .

#### 2.2. Numerical examples

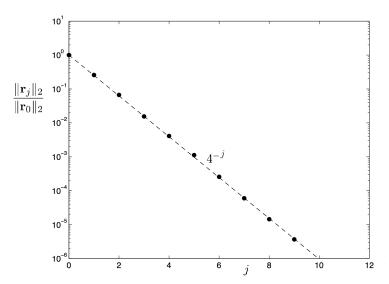
• Example 1:  $\mathbf{A}$ ,  $200 \times 200$  entries from real normal distribution of mean 2 and standard deviation  $0.5/\sqrt{200}$ 

$$m = 200; A = 2*eye(m)+0.5*randn(m)/sqrt(m);$$

$$\mathbf{A}\mathbf{x} = \mathbf{b}, \quad \mathbf{x}_0 = \mathbf{0}, \quad \mathbf{b} = \begin{bmatrix} 1 & 1 & \cdots & 1 \end{bmatrix}^{\mathsf{T}}$$



## Convergence history of Example 1

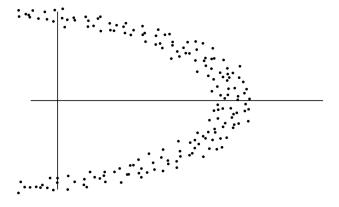


• Example 2:

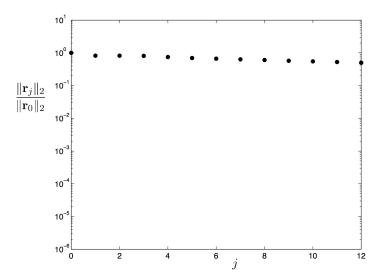
$$m = 200; B = 2*eye(m)+0.5*randn(m)/sqrt(m);$$

A = B + D, D is the diagonal matrix with complex entries

$$d_i = (-2 + 2\sin\theta_i) + i\cos\theta_i, \quad \theta_i = \frac{(i-1)\pi}{m-1}, \quad 1 \le i \le m.$$



# Convergence history of Example 2



## 2.3. Preconditioning (see Lecture 40 of NLA)

- To improve the convergence of Krylov subspace methods, it is important to have a preconditioner (suitable approximation for the original coefficient matrix **A**), denoted by **M**.
- Left preconditioning, i.e.,

$$\mathbf{M}^{-1}\mathbf{A}\mathbf{x} = \mathbf{M}^{-1}\mathbf{b}.$$

• Right preconditioning is often used, i.e.,

$$\mathbf{A}\mathbf{M}^{-1}\mathbf{z} = \mathbf{b}, \quad \mathbf{x} = \mathbf{M}^{-1}\mathbf{z},$$

because it produces the same residual as that of the original system in exact precision arithmetic.

• Note that we never explicitly form  $\mathbf{M}^{-1}$ . Only the action of applying the preconditioner solve operation  $\mathbf{M}^{-1}$  to a given vector is computed in iterative methods. So  $\mathbf{M}^{-1}\mathbf{z}$  must be cheap.

• How to find a good preconditioner? It's problem dependent.

Example. Let 
$$\mathcal{A} = \begin{bmatrix} \mathbf{A} & \mathbf{B}^* \\ \mathbf{C} & \mathbf{0} \end{bmatrix}$$
 and  $\mathcal{M} = \begin{bmatrix} \mathbf{A} & \mathbf{0} \\ \mathbf{0} & \mathbf{C}\mathbf{A}^{-1}\mathbf{B}^* \end{bmatrix}$ , where  $\mathbf{A} \in \mathbb{C}^{m \times m}$  is invertible, and  $\mathbf{B}, \mathbf{C} \in \mathbb{C}^{n \times m}$  with  $m \geq n$ . Assume that  $-\mathbf{C}\mathbf{A}^{-1}\mathbf{B}^*$  is invertible.

The preconditioned matrix  $\mathcal{M}^{-1}\mathcal{A}$  is diagonalizable and has at most three distinct eigenvalues 1,  $(1+\sqrt{5})/2$ , and  $(1-\sqrt{5})/2$ .

#### 2.4. Restarted GMRES

- For larger values of j, the cost of GMRES in operations and storage may be prohibitive. In such circumstances a method called l-step restarted GMRES or GMRES(l) is often employed.
- GMRES(l): After l steps, the GMRES iteration is started anew with the current vector  $\mathbf{x}_l$  as an initial guess.
- Note that GMRES(l) can be expected fail to converge, whereas GMRES always succeeds for exact arithmetic. (Embree's paper)
- GMRES-IR, GMRES-DR, FGMRES-DR, etc.