

Lecture 3: Projector, Classical/Modified Gram–Schmidt orthogonalization, QR factorization



School of Mathematical Sciences, Xiamen University

1. Projector

- A square matrix $\mathbf{P} \in \mathbb{C}^{m \times m}$ is called a *projector* if $\mathbf{P}^2 = \mathbf{P}$. Any projector is diagonalizable. (Eigenvalues?) **Example:** $\mathbf{P} = \begin{bmatrix} 0 & 0 \\ \alpha & 1 \end{bmatrix}$

Theorem 1

Let \mathbf{P} be a projector. Then,

- (1) for all $\mathbf{v} \in \text{range}(\mathbf{P})$, we have $\mathbf{P}\mathbf{v} = \mathbf{v}$;
- (2) $\text{range}(\mathbf{P})$ and $\text{null}(\mathbf{P})$ satisfy

$$\text{range}(\mathbf{P}) \cap \text{null}(\mathbf{P}) = \{\mathbf{0}\}, \quad \text{range}(\mathbf{P}) + \text{null}(\mathbf{P}) = \mathbb{C}^m;$$

- (3) $\mathbf{I} - \mathbf{P}$ is a projector, and

$$\text{range}(\mathbf{I} - \mathbf{P}) = \text{null}(\mathbf{P}), \quad \text{null}(\mathbf{I} - \mathbf{P}) = \text{range}(\mathbf{P}).$$

- (4) if $\mathbf{P} \neq \mathbf{0}, \mathbf{I}$, we have $\|\mathbf{I} - \mathbf{P}\|_2 = \|\mathbf{P}\|_2$. (See [Ref. 1](#) and [Ref. 2](#))

- Two subspaces $\mathcal{S}_1, \mathcal{S}_2 \subseteq \mathbb{C}^m$ are called *complementary subspaces* if they satisfy

$$\mathcal{S}_1 \cap \mathcal{S}_2 = \{\mathbf{0}\}, \quad \mathcal{S}_1 + \mathcal{S}_2 = \mathbb{C}^m.$$

Theorem 2

Let \mathcal{S}_1 and \mathcal{S}_2 be complementary subspaces. Then there exists a unique projector \mathbf{P} with $\text{range}(\mathbf{P}) = \mathcal{S}_1$ and $\text{null}(\mathbf{P}) = \mathcal{S}_2$.

Proof.

The existence is left as an exercise. Now we prove the uniqueness. Let \mathbf{e}_j denote the j th column of the identity matrix \mathbf{I} . Since \mathcal{S}_1 and \mathcal{S}_2 are complementary, we can assume $\mathbf{e}_j = \mathbf{s}_j^1 + \mathbf{s}_j^2$, where $\mathbf{s}_j^1 \in \mathcal{S}_1$, and $\mathbf{s}_j^2 \in \mathcal{S}_2$. Assume both \mathbf{P}_1 and \mathbf{P}_2 are desired projectors. Then we have

$$\begin{aligned} \forall 1 \leq j \leq m, \quad (\mathbf{P}_1 - \mathbf{P}_2)\mathbf{e}_j &= (\mathbf{P}_1 - \mathbf{P}_2)\mathbf{s}_j^1 + (\mathbf{P}_1 - \mathbf{P}_2)\mathbf{s}_j^2 \\ &= \mathbf{P}_1\mathbf{s}_j^1 - \mathbf{P}_2\mathbf{s}_j^1 = \mathbf{s}_j^1 - \mathbf{s}_j^1 = \mathbf{0}. \end{aligned}$$

Therefore, $\mathbf{P}_1 = \mathbf{P}_2$, i.e., uniqueness. □

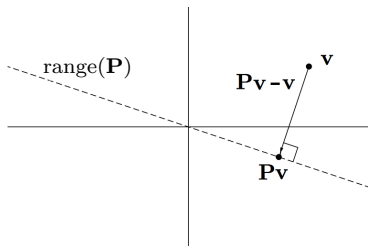
- Let \mathcal{S}_1 and \mathcal{S}_2 be complementary subspaces. The unique projector \mathbf{P} with $\text{range}(\mathbf{P}) = \mathcal{S}_1$ and $\text{null}(\mathbf{P}) = \mathcal{S}_2$ is called the *projector onto \mathcal{S}_1 along \mathcal{S}_2* .

1.1. Orthogonal and oblique projectors

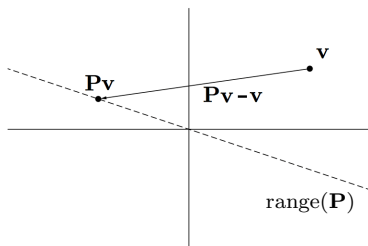
- For a projector \mathbf{P} , if $\text{range}(\mathbf{P})$ and $\text{null}(\mathbf{P})$ are orthogonal, then it is called an *orthogonal* projector. Otherwise, *oblique*.

Warning: orthogonal projectors are not orthogonal matrices!!!

- Geometric interpretation



The orthogonal projection



An oblique projection

Theorem 3

A projector \mathbf{P} is an orthogonal projector if and only if $\mathbf{P} = \mathbf{P}^*$.

- $\mathbf{P} = \begin{bmatrix} 0 & 0 \\ \alpha & 1 \end{bmatrix}$: oblique (if $\alpha \neq 0$) or orthogonal (if $\alpha = 0$) projector.

Theorem 4

Let the columns of \mathbf{Q}_r be an orthonormal basis of an r -dimensional subspace \mathcal{S} . Then the orthogonal projector onto \mathcal{S} is given by $\mathbf{Q}_r \mathbf{Q}_r^*$, and the orthogonal projector onto \mathcal{S}^\perp is given by $\mathbf{I} - \mathbf{Q}_r \mathbf{Q}_r^*$.

- $\mathbf{a} \neq \mathbf{0}$, $\mathbf{P}_{\mathbf{a}} = \frac{\mathbf{a}\mathbf{a}^*}{\mathbf{a}^*\mathbf{a}}$, $\mathbf{P}_{\mathbf{a}^\perp} = \mathbf{I} - \frac{\mathbf{a}\mathbf{a}^*}{\mathbf{a}^*\mathbf{a}}$
- Let $\mathbf{A} \in \mathbb{C}^{m \times n}$. The orthogonal projector onto $\text{range}(\mathbf{A})$ is given by $\mathbf{U}_r \mathbf{U}_r^*$, where \mathbf{U}_r is the matrix in SVD of \mathbf{A} .
- Others: $\mathbf{A}\mathbf{A}^\dagger$ onto $\text{range}(\mathbf{A})$, $\mathbf{A}^\dagger\mathbf{A}$ onto $\text{range}(\mathbf{A}^*)$

1.2. General definitions

- Suppose $\langle \cdot, \cdot \rangle$ denotes an inner product on \mathbb{C}^m . A linear mapping $\mathbf{T} : \mathbb{C}^m \mapsto \mathbb{C}^m$ is called
 - a *projector* if it is idempotent, i.e., if $\mathbf{T}^2 = \mathbf{T}$;
 - an *orthogonal projector* (with respect to $\langle \cdot, \cdot \rangle$) if for all $\mathbf{x}, \mathbf{y} \in \mathbb{C}^m$,

$$\langle \mathbf{x} - \mathbf{T}\mathbf{x}, \mathbf{T}\mathbf{y} \rangle = 0;$$

- *self-adjoint* (with respect to $\langle \cdot, \cdot \rangle$) if for all $\mathbf{x}, \mathbf{y} \in \mathbb{C}^m$,

$$\langle \mathbf{T}\mathbf{x}, \mathbf{y} \rangle = \langle \mathbf{x}, \mathbf{T}\mathbf{y} \rangle.$$

- It is easy to see that a projector is orthogonal (with respect to $\langle \cdot, \cdot \rangle$) if and only if it is self-adjoint (with respect to $\langle \cdot, \cdot \rangle$).
- If \mathbf{T} is self-adjoint (with respect to $\langle \cdot, \cdot \rangle$) so is $\mathbf{I} - \mathbf{T}$ and vice versa.
- **Exercise:** Prove if $\forall \mathbf{x}, \mathbf{y} \in \mathbb{C}^m$, $\langle \mathbf{x} - \mathbf{T}\mathbf{x}, \mathbf{T}\mathbf{y} \rangle = 0$, then $\mathbf{T}^2 = \mathbf{T}$.

2. Gram–Schmidt orthogonalization (GS)

- For n linearly independent vectors $\{\mathbf{a}_i\}_{i=1}^n$: at the j th step, Gram–Schmidt orthogonalization finds a unit vector \mathbf{q}_j that is orthogonal to $\mathbf{q}_1, \dots, \mathbf{q}_{j-1}$, lies in $\text{span}\{\mathbf{a}_1, \dots, \mathbf{a}_j\}$ as follows:

$$\tilde{\mathbf{q}}_j = \mathbf{a}_j - \sum_{i=1}^{j-1} \mathbf{q}_i^* \mathbf{a}_j \mathbf{q}_i, \quad \mathbf{q}_j = \frac{\tilde{\mathbf{q}}_j}{\|\tilde{\mathbf{q}}_j\|_2}.$$

- Equivalently, Gram–Schmidt orthogonalization can be represented via projectors:

$$\tilde{\mathbf{q}}_j = \mathbf{P}_j \mathbf{a}_j, \quad \mathbf{q}_j = \frac{\tilde{\mathbf{q}}_j}{\|\tilde{\mathbf{q}}_j\|_2}.$$

where $\mathbf{P}_j = \mathbf{I} - \mathbf{Q}_{j-1} \mathbf{Q}_{j-1}^*$ and $\mathbf{Q}_{j-1} = [\mathbf{q}_1 \quad \mathbf{q}_2 \quad \dots \quad \mathbf{q}_{j-1}]$.

2.1. Classical Gram–Schmidt orthogonalization (ClGS)

- ClGS is based on the use of

$$\tilde{\mathbf{q}}_j = \mathbf{a}_j - \mathbf{q}_1^* \mathbf{a}_j \mathbf{q}_1 - \mathbf{q}_2^* \mathbf{a}_j \mathbf{q}_2 \cdots - \mathbf{q}_{j-1}^* \mathbf{a}_j \mathbf{q}_{j-1}$$

and calculates \mathbf{q}_j by evaluating the following formulas in order:

$$\begin{aligned} \mathbf{q}_j^{(0)} &= \mathbf{a}_j, \\ \mathbf{q}_j^{(1)} &= \mathbf{q}_j^{(0)} - \mathbf{q}_1^* \mathbf{a}_j \mathbf{q}_1, \\ \mathbf{q}_j^{(2)} &= \mathbf{q}_j^{(1)} - \mathbf{q}_2^* \mathbf{a}_j \mathbf{q}_2, \\ &\vdots \\ \mathbf{q}_j^{(j-1)} &= \mathbf{q}_j^{(j-2)} - \mathbf{q}_{j-1}^* \mathbf{a}_j \mathbf{q}_{j-1}, \\ \mathbf{q}_j &= \mathbf{q}_j^{(j-1)} / \|\mathbf{q}_j^{(j-1)}\|_2. \end{aligned}$$

2.2. Modified Gram–Schmidt orthogonalization (MGS)

- MGS is based on the use of

$$\tilde{\mathbf{q}}_j = \mathbf{P}_j \mathbf{a}_j = (\mathbf{I} - \mathbf{q}_{j-1} \mathbf{q}_{j-1}^*) \cdots (\mathbf{I} - \mathbf{q}_2 \mathbf{q}_2^*) (\mathbf{I} - \mathbf{q}_1 \mathbf{q}_1^*) \mathbf{a}_j$$

and calculates \mathbf{q}_j by evaluating the following formulas in order:

$$\begin{aligned} \mathbf{q}_j^{(0)} &= \mathbf{a}_j, \\ \mathbf{q}_j^{(1)} &= \mathbf{q}_j^{(0)} - \mathbf{q}_1^* \mathbf{q}_j^{(0)} \mathbf{q}_1, \\ \mathbf{q}_j^{(2)} &= \mathbf{q}_j^{(1)} - \mathbf{q}_2^* \mathbf{q}_j^{(1)} \mathbf{q}_2, \\ &\vdots \\ \mathbf{q}_j^{(j-1)} &= \mathbf{q}_j^{(j-2)} - \mathbf{q}_{j-1}^* \mathbf{q}_j^{(j-2)} \mathbf{q}_{j-1}, \\ \mathbf{q}_j &= \mathbf{q}_j^{(j-1)} / \|\mathbf{q}_j^{(j-1)}\|_2. \end{aligned}$$

2.3. ClGS and MGS algorithms

Algorithm: GS for n linearly independent vectors $\{\mathbf{a}_i\}_{i=1}^n$.

```
for  $j = 1$  to  $n$ 
     $\mathbf{q}_j = \mathbf{a}_j$ 
    for  $i = 1$  to  $j - 1$ 
         $\begin{cases} r_{ij} = \mathbf{q}_i^* \mathbf{a}_j & \text{ClGS} \\ r_{ij} = \mathbf{q}_i^* \mathbf{q}_j & \text{MGS} \end{cases}$ 
         $\mathbf{q}_j = \mathbf{q}_j - r_{ij} \mathbf{q}_i$ 
    end
     $r_{jj} = \|\mathbf{q}_j\|_2$ 
     $\mathbf{q}_j = \mathbf{q}_j / r_{jj}$ 
end
```

- The computational cost: $\sim 2mn^2$ (leading term) for $\mathbf{a}_i \in \mathbb{C}^m$

Remark 5

ClGS and MGS are mathematically equivalent. In finite precision arithmetic, MGS introduces smaller errors than ClGS.

- Experiment: ClGS vs. MGS

```
m = 512;
```

```
n = 256;
```

```
A = rand(m,n);
```

```
[QC,RC] = clgs(A);
```

```
ClGS_error = norm(QC'*QC-eye(n))
```

```
[QM,RM] = mgs(A);
```

```
MGS_error = norm(QM'*QM-eye(n))
```

3. QR factorization

- **Definition:** Let m and n be arbitrary positive integers ($m \geq n$ or $m < n$). Given $\mathbf{A} \in \mathbb{C}^{m \times n}$, not necessarily of full rank, a *full QR factorization* of \mathbf{A} is a factorization

$$\mathbf{A} = \mathbf{Q}\mathbf{R}$$

where $\mathbf{Q} \in \mathbb{C}^{m \times m}$ is unitary, and $\mathbf{R} \in \mathbb{C}^{m \times n}$ is upper triangular. For $m \geq n$, a *reduced QR factorization* of \mathbf{A} is a factorization

$$\mathbf{A} = \mathbf{Q}_n \mathbf{R}_n$$

where $\mathbf{Q}_n \in \mathbb{C}^{m \times n}$ has orthonormal columns, and

$$\mathbf{R}_n = \begin{bmatrix} r_{11} & r_{12} & \cdots & r_{1n} \\ & r_{22} & \ddots & \vdots \\ & & \ddots & \vdots \\ & & & r_{nn} \end{bmatrix} \cdot \begin{array}{c} \text{[gray rectangle]} \\ \mathbf{A} \end{array} = \begin{array}{c} \text{[gray rectangle]} \\ \mathbf{Q} \end{array} \cdot \begin{array}{c} \text{[upper triangular rectangle]} \\ \mathbf{R} \end{array}$$

Theorem 6 (Existence of QR)

Every matrix $\mathbf{A} \in \mathbb{C}^{m \times n}$ ($m \geq n$) has a reduced QR factorization and a full QR factorization.

Proof.

- Existence of reduced QR factorization.

For the full column rank case, Gram–Schmidt orthogonalization produces a sequence of reduced QR factorizations for $\mathbf{A} \in \mathbb{C}^{m \times n}$:

$$\mathbf{A}_j := [\mathbf{a}_1 \quad \mathbf{a}_2 \quad \dots \quad \mathbf{a}_j] = \mathbf{Q}_j \mathbf{R}_j, \quad j = 1:n.$$

For the rank-deficient case, $\tilde{\mathbf{q}}_j = \mathbf{0}$ at one or more steps j , GS fails to produce \mathbf{q}_j . At this moment, we pick \mathbf{q}_j arbitrarily to be any unit vector orthogonal to $\text{span}\{\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_{j-1}\}$, set $r_{jj} = 0$, and then continue the Gram–Schmidt orthogonalization until we obtain a reduced QR factorization.

- Existence of full QR factorization.

Let $\mathbf{A} = \mathbf{Q}_n \mathbf{R}_n$ be a reduced QR factorization of \mathbf{A} . A full QR factorization can be constructed via

$$\mathbf{A} = \mathbf{Q}\mathbf{R} := [\mathbf{Q}_n \quad \mathbf{Q}_c] \begin{bmatrix} \mathbf{R}_n \\ \mathbf{0} \end{bmatrix},$$

where $\mathbf{Q}_c \in \mathbb{C}^{m \times (m-n)}$ has orthonormal columns orthogonal to $\text{span}\{\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_n\}$. □

Theorem 7

Every matrix $\mathbf{A} \in \mathbb{C}^{m \times n}$ ($m \geq n$) of full column rank has a unique reduced QR factorization $\mathbf{A} = \mathbf{Q}_n \mathbf{R}_n$ with $r_{jj} > 0$.

Proof.

$r_{11}\mathbf{q}_1 = \mathbf{a}_1$ and $r_{11} > 0 \Rightarrow r_{11}$ and \mathbf{q}_1 unique $\Rightarrow r_{12}$ and $r_{22}\mathbf{q}_2$ unique, by $r_{22} > 0 \Rightarrow r_{22}$ and \mathbf{q}_2 unique, and so on. □

3.1. When vectors become continuous functions

- Replace \mathbb{C}^m by $C[-1, 1]$, a linear space of real-valued continuous functions on $[-1, 1]$ with the L^2 inner product

$$\forall f(x), g(x) \in C[-1, 1], \quad \langle f(x), g(x) \rangle_{L^2} = \int_{-1}^1 f(x)g(x)dx,$$

and the norm

$$\|f(x)\|_{L^2} = \sqrt{\langle f(x), f(x) \rangle_{L^2}}.$$

Gram–Schmidt orthogonalization (GS) with respect to the L^2 inner product $\langle f(x), g(x) \rangle_{L^2}$ is: At step j ,

$$\begin{aligned}\tilde{q}_j(x) &= a_j(x) - \sum_{i=1}^{j-1} \langle a_j(x), q_i(x) \rangle_{L^2} q_i(x), \\ q_j(x) &= \tilde{q}_j(x) / \|\tilde{q}_j(x)\|_{L^2}.\end{aligned}$$

The functions $q_j(x)$ satisfy

$$\langle q_i(x), q_j(x) \rangle_{L^2} = \int_{-1}^1 q_i(x) q_j(x) dx = \delta_{ij} = \begin{cases} 1 & \text{if } i = j, \\ 0 & \text{if } i \neq j. \end{cases}$$

Then we have “continuous QR factorization”

$$A = QR = \left[\begin{array}{c|c|c|c} & & & \\ q_1(x) & q_2(x) & \cdots & q_n(x) \end{array} \right] \left[\begin{array}{cccc} r_{11} & r_{12} & \cdots & r_{1n} \\ & r_{22} & & \vdots \\ & & \ddots & \\ & & & r_{nn} \end{array} \right]$$

where

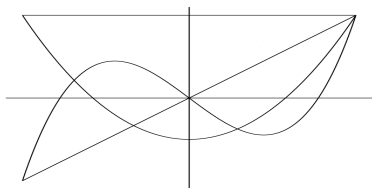
$$A = [a_1(x) \quad a_2(x) \quad \cdots \quad a_n(x)]$$

and

$$r_{jj} = \|\tilde{q}_j(x)\|_{L^2}, \quad r_{ij} = \langle a_j(x), q_i(x) \rangle_{L^2}.$$

- **Example:** $a_j(x) = x^{j-1}$, $j = 1, 2, \dots, n$

$$\mathbf{A} = \begin{bmatrix} 1 & x & x^2 & \dots & x^{n-1} \end{bmatrix}$$



Legendre polynomials $P_j(x) = q_j(x)/q_j(1)$:

$$P_1(x) = x, \quad P_2(x) = \frac{3}{2}x^2 - \frac{1}{2}, \quad P_3(x) = \frac{5}{2}x^3 - \frac{3}{2}x.$$

Experiment: Discrete Legendre polynomials

```
x = (-128:128)'/128;
A = [x.^0 x.^1 x.^2 x.^3]
[Q,R] = qr(A,0);
scale = Q(257,:);
Q = Q*diag(1./scale);
plot(x,Q)
```

