50.020 Security Lecture 5: Password and Rainbow Tables

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#### **Password-based Authentication**

# **Terminology**

- Access control
  - Allow/deny users access to resources
  - Sometimes, delegation is possible
- Authentication verifies correctness of data and source
  - In this context: verifying the *identity* of login request
  - identification itself does not include verification

#### Identification schemes

50.020 Security Lecture 5: Password and Rainbow Tables Identifaction only needs to provide an identity

- No direct security requirements
- Identifier should not be secret
  - E.g.: Social security numbers (bad)

Common Identification schemes:

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Common Identification schemes:

- Email addresses
- Usernames/Real names
- Phone numbers
- Credit card numbers
- Bank cards (IDs on cards)
- National ID number, FIN, etc.
- Face recognition
  - Most other biometrics not good enough

### Authentication schemes

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■ Which authentication schemes do you know?

#### Authentication schemes

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Which authentication schemes do you know?

- Knowledge-based authentication
  - Passwords
  - Patterns (Android lock screen)
  - Banking PINs
- Token-based authentication
  - Physical keys
  - Cryptographic tokens (e.g., from bank)
  - Certificate-based (e.g., TPM)
- Biometrics-based authentication
  - Fingerprint, voice, retina
- Multi-factor authentication

## Password-based Authentication

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Passwords are the most common way to authenticate users

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- Advantages:
  - Can be changed
  - User is free to choose
- Disadvantages
  - Can be forgotten
  - User can create bad passwords
  - Can be re-used

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■ But: how to remember 30+ passwords?

# Why is guessing passwords so easy (compared to keys)?

- Passwords use string.printable
- Passwords are somewhat short
- Some passwords are used more frequently
  - Or have frequently used components
- This enables semi-intelligent brute-forcing
  - dictionaries
  - hybrid attacks

# Dictionary attacks

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_	11		-:	
	Users	preter	simple	passwords

- Dictionary attacks produce lists of popular passwords
- Ordered by popularity
- This maximises likelyhood of success with minimal tries
- Often based on sets of passwords that became public

Rank	Password	
1	123456	
2	password	
3	12345678	
4	qwerty	
5	abc123	
6	123456789	
7	111111	
8	1234567	
9	iloveyou	
10	adob123	

Popular Passwords 2013 (according to SplashData)

# Hybrid attacks

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- Users have heard about dictionary attacks
- "p4sSw0Rd" might not be in dictionary
  - But it is still pretty similar
- Hybrid attacks also try combinations and popular substitutions
  - E.g. replacements such as "a-> 4" and "o-> 0", case
- Interesting estimation of effort for attacks:

https://www.bennish.net/password-strength-checker/

# Finding Passwords in practise

- Both dictionary attacks and hybrid attacks can be used to build long lists of likely passwords
- If there is an API to submit unlimited password attempts, this could be called to break into a system
  - In practise, accounts are quickly locked to prevent this
- In most cases, dictionary and hybrid attacks are used to attempt to find preimages of hashes
  - Password hashes were stolen in some attack
  - Attacker has unlimited attempts to find preimage
- Is your account compromised? Check at https://haveibeenpwned.com/

### HavelBeenPwned Example

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Tables

#### Breaches you were pwned in

A "breach" is an incident where a site's data has been illegally accessed by hackers and then released publicly. Review the types of data that were compromised (email addresses, passwords, credit cards etc.) and take appropriate action, such as changing passwords.

#### lost.fm

Last fin: In March 2012, the music website last fin was hacked and 43 million user accounts were exposed. Whilst last fin knew of an incident back in 2012, the scale of the hack was not known until the data was released publicly in September 2016. The breach included 37 million unique email addresses, usernames and passwords stored as unsalted MDS hashes.

Compromised data: Email addresses, Passwords, Usernames, Website activity



UnkedIn. In May 2016, UnkedIn had 164 million email addresses and passwords exposed. Originally hacked in 2012, the data remained out of sight until being offered for sale on a dark market site 4 years later. The passwords in the breach were stored as SHA1 hashes without salt, the vast majority of which were quickly cracked in the days following the release of the data.

Compromised data: Email addresses, Passwords



MySpace. In approximately 2008, MySpace suffered a data breach that exposed almost 360 million accounts. In May 2016 the data was offered up for sale on the "Seal Deal" dark market website and included email addresses, usernames and SHA1 hashes of the first 10 characters of the password converted to lowercase and stored without a salt. The exact breach date is unknown, but analysis of the data suggests it was 8 years before being made public.

Compromised data: Email addresses, Passwords, Usernames



Patron: In October 2015, the crowfunding site Patron was hacked and over 16GB of data was released publicly. The dump included almost 14GB of database records with more than 2.3M unique email addresses and millions of personal messages.

Compromised data: Email addresses, Payment histories, Private messages, Website activity

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### Strengthening Passwords

# Multi-factor authentication (MFA)

- Simple, most commong form: Two-factor authentication (TFA)
- Combines username/password with second way, e.g. text messages
  - Example: DBS login into account



First factor login

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- MFA is an application of "defense in depth"

# Summary: "Best practises" for passwords

- Do not re-use the same password between two services
- Change passwords periodically (but not too often)
- Do not base your passwords on dictionary words
  - Dialects or made-up words are better
- How to find good passwords?
  - Find a good scheme to adapt your password to each site
  - Example: "p4ssw0rd(amazon)", "p4ssw0rd(sutd)" etc
- Use random passwords, and management software
- Use hashes/ your scheme here

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# **Finding Hash Preimages**

# **Brute Forcing Hashes**

- Common cryptographic hashes have length 128 (e.g., MD5), 160 (SHA1), 224-512(SHA3)
- Brute-forcing SHA1 takes about  $O(2^{160})$  computations
- How could we speed this up?
  - Precompute some/all values!
- If we precompute 2<sup>160</sup> hashes, we can directly look up preimage
  - Unfortunately, this takes 160 \* 2<sup>120</sup> Terabyte of storage

# Improving Brute Force

- So clearly, computing hashes on demand or full precomputation is infeasible
- Can we mix both?
  - Do precomputations, but only store a subset of the found hashes
  - Just store as many hash values as you have storage space
  - Ensure that you can recover preimage of the hash values
- This is the idea behind rainbow tables, which are based on hash chains

# Improving Brute Force (continue)

- In the following, we ignore *hash collisions* to simplify things
- We also assume that the input space is smaller than the output space, e.g. if only 10 character inputs are considered...
- Note: Rainbow tables are not computed "on the fly" to look up one hash
  - Direct brute force would be more efficient in that case
  - Rainbow tables allow you to re-use brute force effort for many hashes

#### Hash Chains

- Hash chains trade storage space vs. computational effort
- General operations of hash chain:
  - H() is hashing function from input domain to hash domain
  - R() is some reduction function, mapping from hash to input domain
    - If you only care about string printable of length < 10, then R() should map into that
- Lets initiate a hash chain with  $I_1$  as first input

$$I_1 \xrightarrow{H()} O_1 \xrightarrow{R()} I_2 \xrightarrow{H()} O_2 \xrightarrow{R()} \cdots \xrightarrow{H()} O_t$$

- After t operations, we get hash output  $O_t$
- If we store  $I_1$  and  $O_t$ , how can we find  $I_t$ ?

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Solution to be provided in class.

#### Rainbow tables

- Rainbow tables trade time vs. space in hash reversal
- $\blacksquare$  A rainbow table consists of m hash chains of length t
- For each chain, only the first input (i.e., plaintext)  $I_1$  and the last hash  $O_t$  are stored.
- Overall space requirement: (|I| + |O|) \* m
- Even more space-efficient:
  - Each of the *m* chains can use its index as starting input *I*
  - For each chain, only the last hash  $O_t$  is stored
  - Overall space requirement: |O| \* m
- The product m \* t must be  $\geq$  number of possible input values
  - E.g. if 10 characters [a-Z]: 26<sup>10</sup>
- Runtime of hash lookup: O(t/2 + t/2) if comparisons are free, and hashing is only expensive operation

# Rainbow table operation

- How to use hash chain for lookup: We want to find X:Y=H(X)
- Check if Y is in list of last chain elements
  - if yes  $(Y = O_z)$ : regenerate chain z using the input value  $I_z$ . Then find the  $I_z$  that was used to compute  $O_z$ , this is our X
  - if no: compute H(R(Y))=Y', see if this is in list of last chain elements
    - if yes  $(Y' = O_z)$ : regenerate chain z using the input value  $I_z$ . Then find the  $I_z$  that was used to compute  $O_z$ , this is our X
    - $lack {f i}$  if no: apply further iterations of reduction and hashing on  ${f Y}$

# Rainbow table operation (continue)

- Estimated effort: (only counting hashing, as most expensive op)
  - expected t/2 reductions to find a matching last value + t/2 average effort to regenerate chain  $\Rightarrow O(t/2 + t/2)$
- Brute force effort
  - Either O(n/2 = m \* t/2) (on-the-fly computation) or
  - O(1) computation and O(m \* t) space

#### RainbowCrack

- Rainbow tables can be generated using RainbowCrack tool for MD5, SHA1, NTLM, . . .
- If you have enough space and time, you can open a commercial rainbowtable:
  - https://www.cloudcracker.com/ [was still working in 2017, now offline]
  - \$17 per hash lookup

# Defending against Rainbow tables

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To make rainbow tables infeasible, a salt (random number) is added to each hash

- E.g.: x = H(m||s) with salt s
- The salt can be stored with hashed password (x,s)
- The salt should be different for each user
- What is the benefit?
- The attacker cannot just use the same rainbow table
- If attacker would want to pre-compute rainbow tables: n bit salt increases effort for attacker by  $2^n$ 
  - Each salt requires own rainbow table of same size as original one
- Some people say rainbow tables are dead. . .

#### Hashcat

- If salts are used, brute force might be the only solution to find preimages
  - In particular, if non-salt part of input is from small space
  - In particular, if hashing function is computationally cheap
- Effort for the attacker directly depends on cost of hash
  - If hashing can be done in 1% of time, attack is 100 times faster
  - Bitcoin caused a lot of specialized hashing hardware to appear
  - Modern GPUs can also be used for hashing (e.g. NVIDIA CUDA)
- Hashcat is an example tool to do such online attacks
  - http://hashcat.net/oclhashcat/
  - Hashcat leverages GPUs for hashing using OpenCL

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#### **Detecting compromise**

# Honeywords

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- How to detect if an attacker is trying to break into your system?
- Create password entries for dummy users
- Trigger alarms when these users try to log in
- Trigger big alarm if these users log in with correct password!
- More reading by Rivest:

http://people.csail.mit.edu/rivest/honeywords/

#### Conclusions

- Password-based authentication
- Strengthening password
- Finding Hash pre-images (rainbow table, etc.)