

0-1 knapsack

✓ subset sum

✓ equal sum partition

count of subset sum

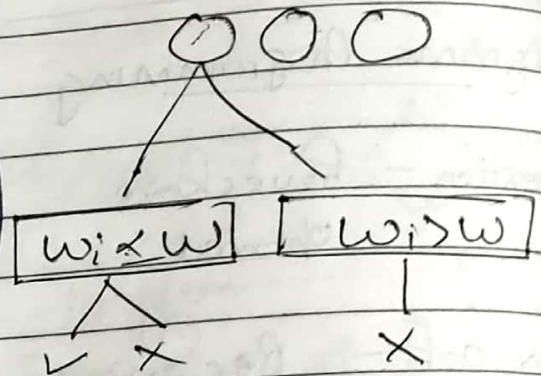
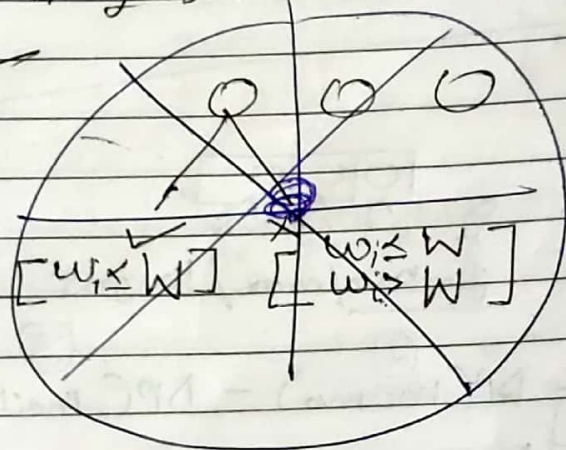
✓ min subset sum diff

✓ target sum

of subset with given diff

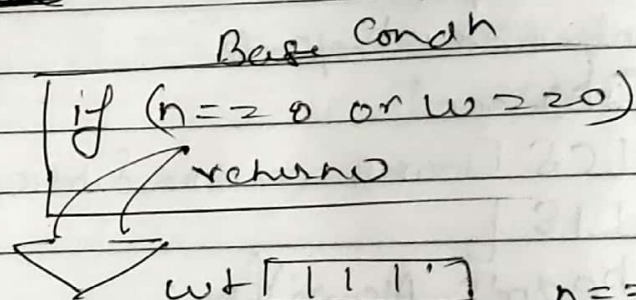


~~Recursive~~



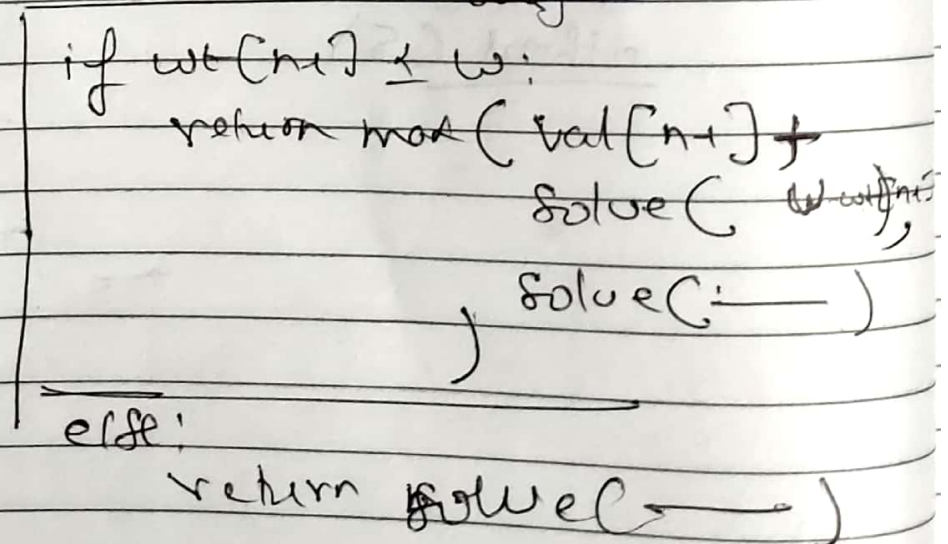
~~Base Cond~~

~~Base Cond~~



wt [] n = 0
val []
w w = 0

choice diagram



memoization

Base Case

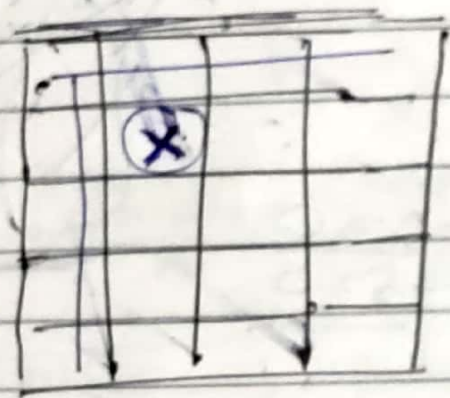
if memo[i][j] != None
return memo[i][j]

choice tree

memo[i][j] = ans
return memo[i][j]

disadvantage: Stack overflow / little slower

tabulation



if cut[i][j] -

① initialization → base case

② solve → memo

i =
j =

01 knapsack bottom-up



take them according to method

(1) Solve (n, w)

memo = [None for (nw)]
for (w, h)

(2) # base condition
if $(n=0 \text{ or } w=0)$
return 0

base condn
for $(h=0 \rightarrow n+1)$
for $(w=0 \rightarrow w+1)$
if $n=0 \text{ or } w=0$
memo[i][j] = 0

(3) # choice diagram

$w = w + wt[n]$
if $w \geq 0$:
return $\max(\text{val}[n] + \text{solve}(w - wt[n], n-1))$

choice diagram
for n in range(1, n+1)
for w in range(1, w+1)

if $w - wt[n] \geq 0$
if $(w - wt[n]) \geq 0$
memo[n][w] = $\max(\text{val}[n] + \text{memo}[n-1][w - wt[n]])$

else:
return solve(n-1, w)

else:
memo[n][w] = memo[n-1][w]