

## ARTIFICIAL INTELLIGENCE

REVISION simple reflex agents reflex agents with state goal-based agents ytility-based agents All of these 4 types can be turned into learning agents

#### PROBLEM SOLVING

Problem Solving Agents -> Atomic Representation

Planning Agents -> Factored or Structured Representation

Assumption: Solution of a problem is a fixed sequence of actions and this sequence DOES NOT depend upon future percepts

(Chapter 4 relaxes above mentioned assumption)

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## INFORMED VS UNINFORMED SEARCH

- Problem Solv

**/uninformed** search algorithms—algorithms that are given no information about the problem other than its definition.

Informed search algorithms, on the other hand, can do quite well given some

guidance on where to look for solutions

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- PROBLEM SOLVING AGENT
- **Goal Formulation** 1)
- **Problem Formulation** 2)
- Searching Solution 3)
- **Executing Solution**

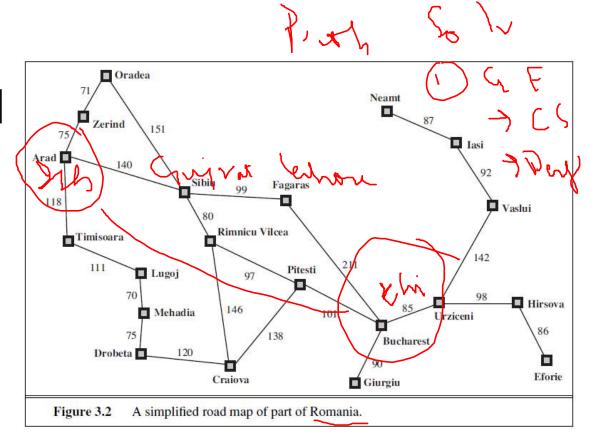
## 1- GOAL FORMULATION

Goal formulation, based on the current situation/state and the agent's performance measure, is the first step in problem solving.

Performance Measure: Visit as many cities as possible, spend as low as possible on fuel.

Current State: In Arad

Possible Goal: In Bucharest



A set of states in which goal is satisfied

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## 2- PROBLEM FORMULATION

**Problem formulation** is the process of deciding what actions and states to consider, given a goal.

The agent will consider actions at the level of driving from one major town to another. Each state therefore corresponds to being in a particular town.

Go Left, Go Right, Go forward, Go reverse. If these 4 actions are considered agent will never go out of the parking lot let alone reaching Bucharest.

Possible Action: Go Cityname

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# TATE SPACE GRAPH

On the right is the state space graph of our problem

State Space \- the set of all states reachable from the initial state by any sequence of actions.

A **path** in the state space is a sequence of states connected by a sequence of actions.

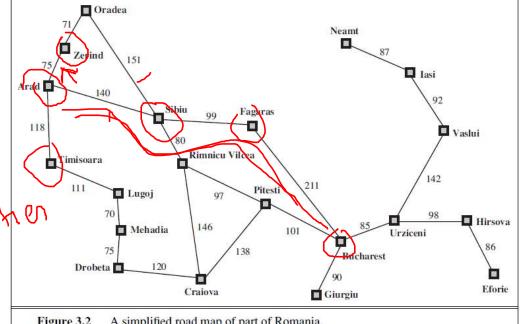
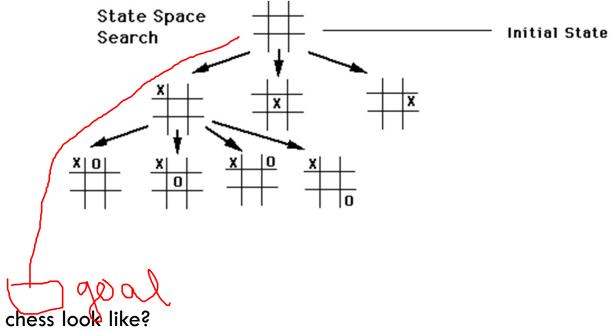


Figure 3.2 A simplified road map of part of Romania.

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## STATE SPACE GRAPH



How would a state space of chess look like?

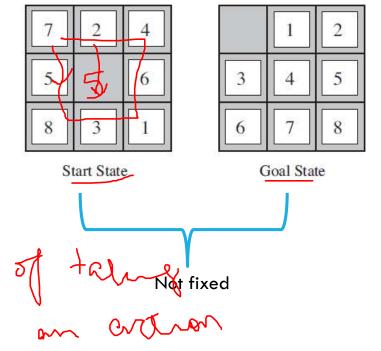
## TOY PROBLEM: 8 PUZZLE

Any state can be used as starting and Goal states.

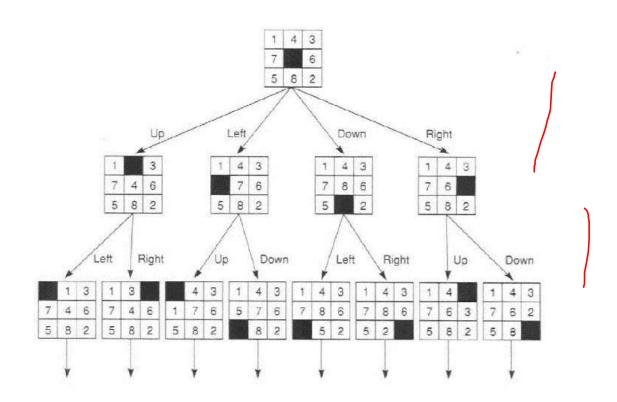
Actions: Left, Right, Up, or Down.

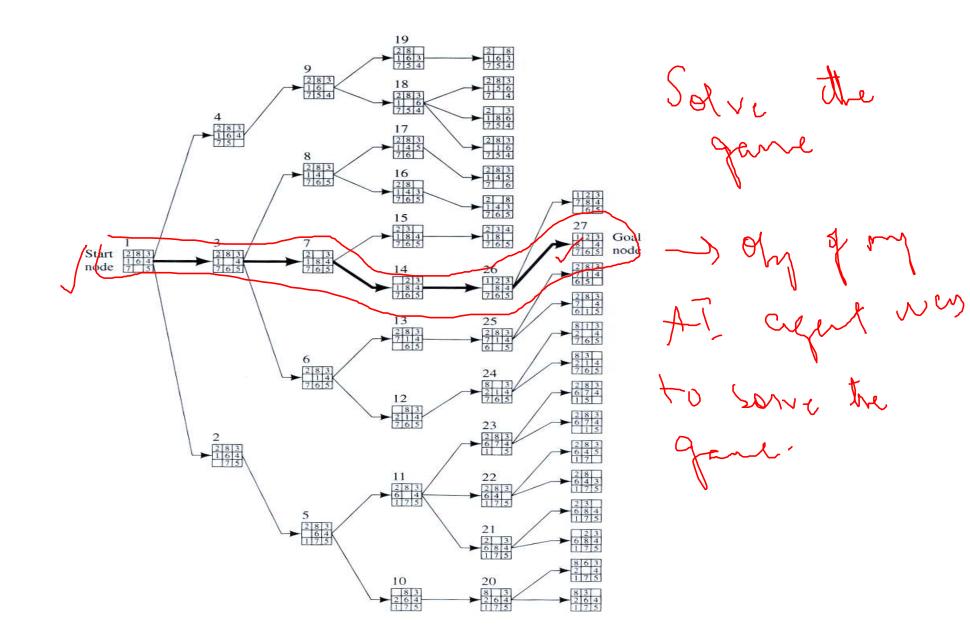
**Transition model:** Given a state and action, this returns the resulting state; for example, if we apply *Left* to the start state in adjacent figure, the resulting state has the 5 and the blank switched.

Path Cost: Each step costs 1, so the path cost is the number of steps in the path.



## TOY PROBLEM: 8 PUZZLE

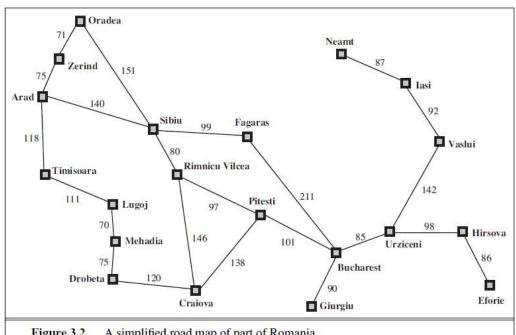




## PROBLEM DEFINITION

A problem consists of five components,

- Initial State
- Possible Actions
- Transition Model
- Goal Test
- Path Cost



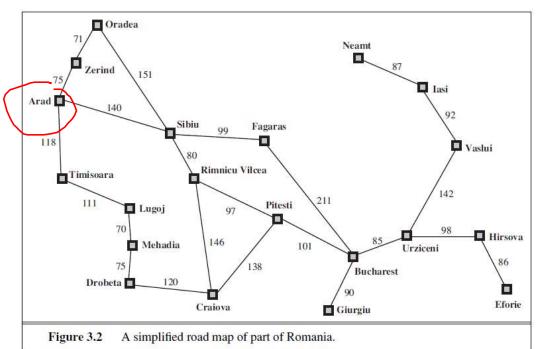
A simplified road map of part of Romania. Figure 3.2

## PROBLEM DEFINITION

A problem consists of five components,

- ✓Initial State : in (Arad)
  - Possible Actions (given a state): e.g., from Arad possible actions are {Go(Sibiu), Go(Timisoara), Go(Zerind)}.
  - Transition Model: specifies the relationship between a state, a possible action and the resulting successor state e.g.,

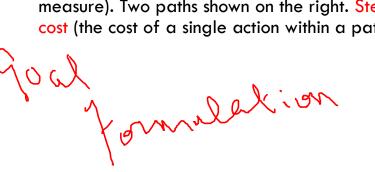
RESULT(In(Arad),Go(Zerind)) = In(Zerind)

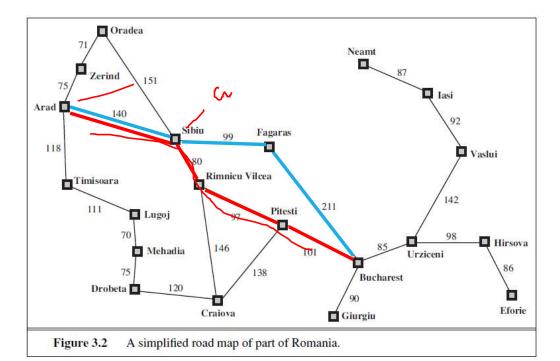


## PROBLEM DEFINITION

A problem consists of five components,

- Initial State
- Possible Actions
- Transition Model
- Goal Test (whether a given state is a goal state) e.g. {In(Bucharest )}.
- Path Cost (based upon agent's performance measure). Two paths shown on the right. Step cost (the cost of a single action within a path)

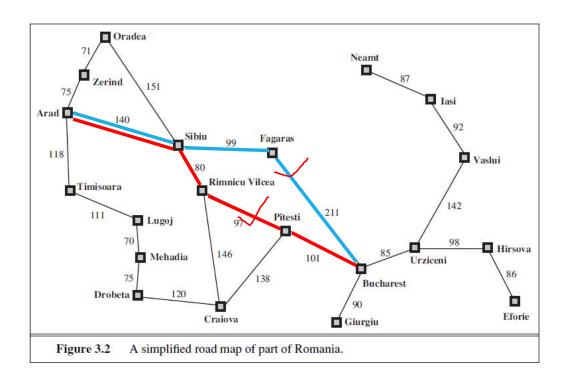




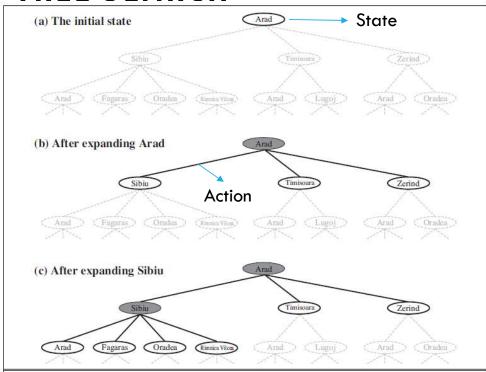
## **SOLUTION DEFINITION**

A **solution** to a problem is an action sequence that leads from the initial state to a goal state.

Solution quality is measured by the path cost function, and an eptimal solution has the lowest path cost among all solutions.



## TREE SEARCH



**Figure 3.6** Partial search trees for finding a route from Arad to Bucharest. Nodes that have been expanded are shaded; nodes that have been generated but not yet expanded are outlined in bold; nodes that have not yet been generated are shown in faint dashed lines.

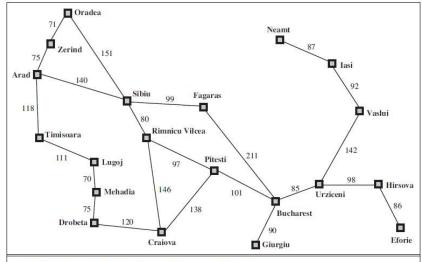
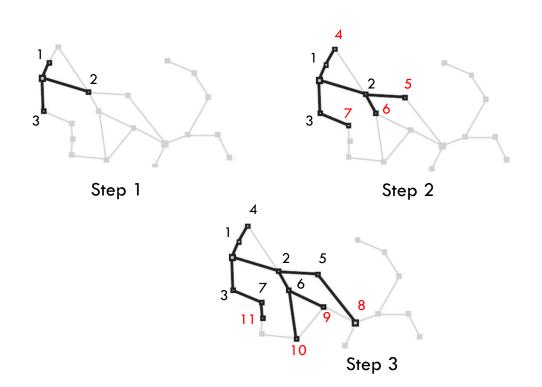
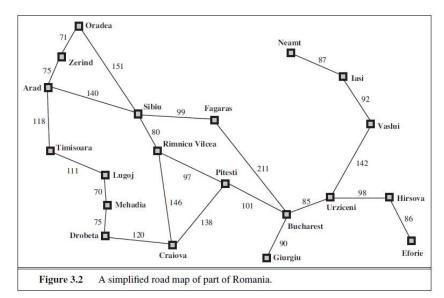


Figure 3.2 A simplified road map of part of Romania.

## **GRAPH SEARCH**

Use an explore set / closed list / frontier to remember the states already visited





## INFRASTRUCTURE FOR SEARCH ALGORITHMS

- n.STATE: the state in the state space to which the node corresponds;
- n.PARENT: the node in the search tree that generated this node;
- n.Action: the action that was applied to the parent to generate the node;
- n.PATH-COST: the cost, traditionally denoted by g(n), of the path from the initial state to the node, as indicated by the parent pointers.

Frontier / expanded nodes / unexpanded nodes are stored in the form of a queue (can be FIFO, LIFO (stack) or priority queue)

## MEASURING PROBLEM-SOLVING PERFORMANCE

**Completeness:** Is the algorithm guaranteed to find a solution when there is one?

**Optimality:** Does the strategy find the optimal solution.

Time complexity: How long does it take to find a solution?

**Space complexity:** How much memory is needed to perform the search?

## MEASURING PROBLEM-SOLVING PERFORMANCE

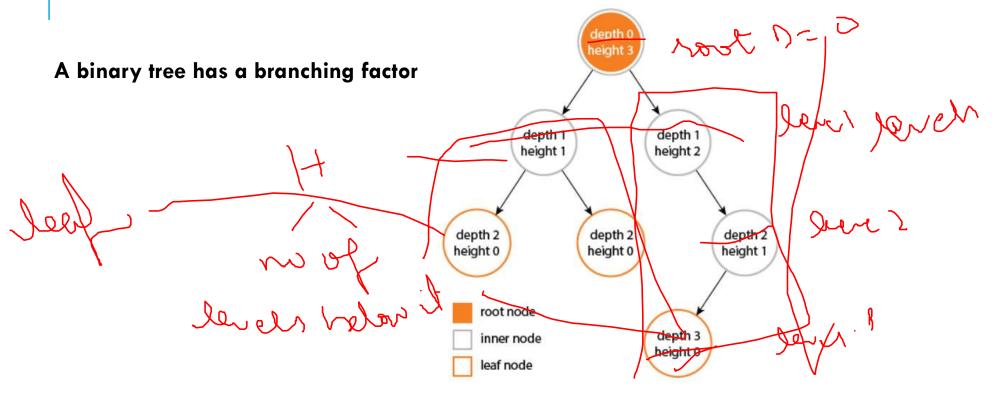
#### Complexity in a tree or a graph is measured by:

- 1)b, the branching factor or maximum number of children of any node
- 2)d, the depth of the goal node (i.e., the number of steps along the path from the root)
- 3)m, the maximum possible depth of entire tree.

Time is often measured in terms of the number of nodes generated during the search.

Space in terms of the maximum number of nodes stored in memory.

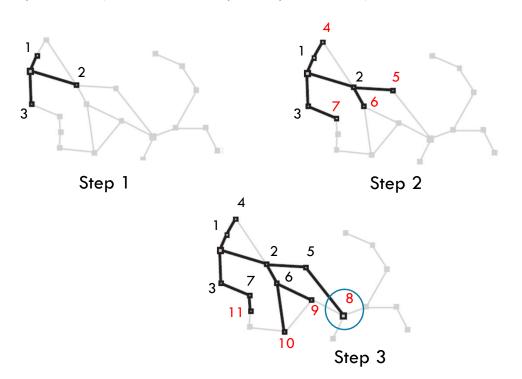
## MEASURING PROBLEM-SOLVING PERFORMANCE

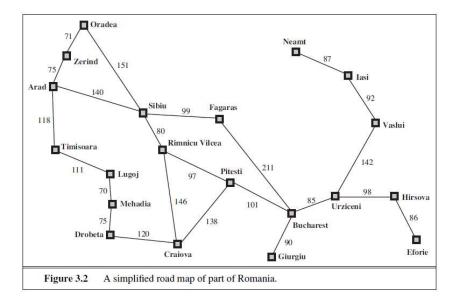


## **COMPLEXITY**

Branching Factor, b=3 (Max num of children for any node is 3)

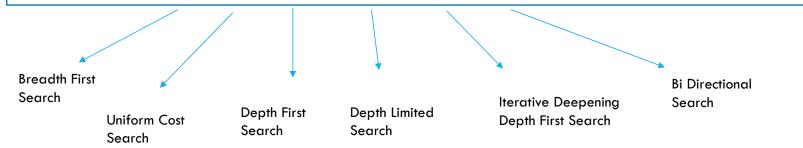
Depth, d = 3 (Bucharest is 3 steps away from Arad)





## INFORMED VS UNINFORMED SEARCH

**uninformed** search algorithms—algorithms that are given no information about the problem other than its definition.



**Informed** search algorithms, on the other hand, can do quite well given some guidance on where to look for solutions

Summary.

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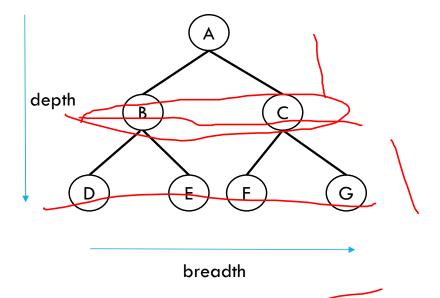
**Expand Shallowest Node First** 

Frontier (or fringe): nodes in queue to be explored

Reached: Nodes that are already explored

Frontier is a first-in-first-out (FIFO) queue, i.e., new successors go at end of the queue.

Goal-Test when inserted



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Goal-Test when inserted

```
function BREADTH-FIRST-SEARCH(problem) returns a solution node or failure

node ← Node(problem.Initial)

if problem.Is-Goal(node.State) then return node

frontier ← a FIFO queue, with node as an element

reached ← {problem.Initial}

while not Is-Empty(frontier) do

node ← Pop(frontier)

for each child in Expand(problem, node) do

s ← child.State

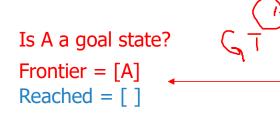
if problem.Is-Goal(s) then return child

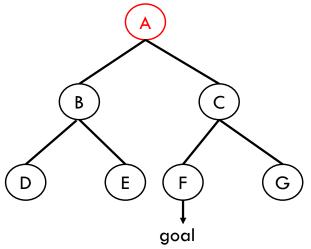
if s is not in reached then

add s to reached

add child to frontier

return failure
```





```
function BREADTH-FIRST-SEARCH(problem) returns a solution node or failure

node ← NODE(problem INITIAL)

if problem.IS-GOAL(node.STATE) then return node

frontier ← a FIFO queue, with node as an element

reached ← {problem.INITIAL}

while not IS-EMPTY(frontier) do

node ← POP(frontier)

for each child in EXPAND(problem, node) do

s ← child.STATE

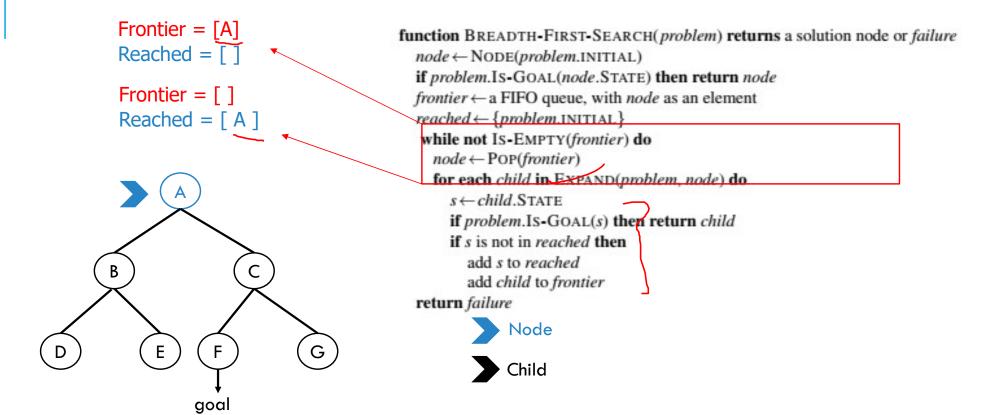
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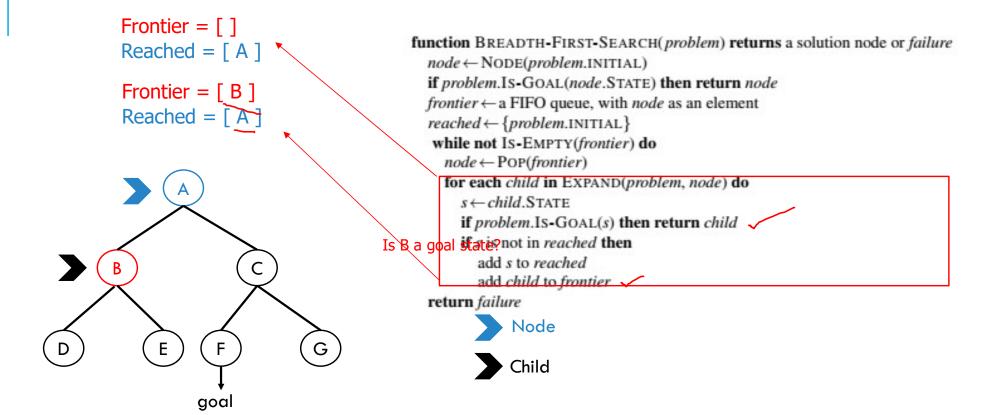
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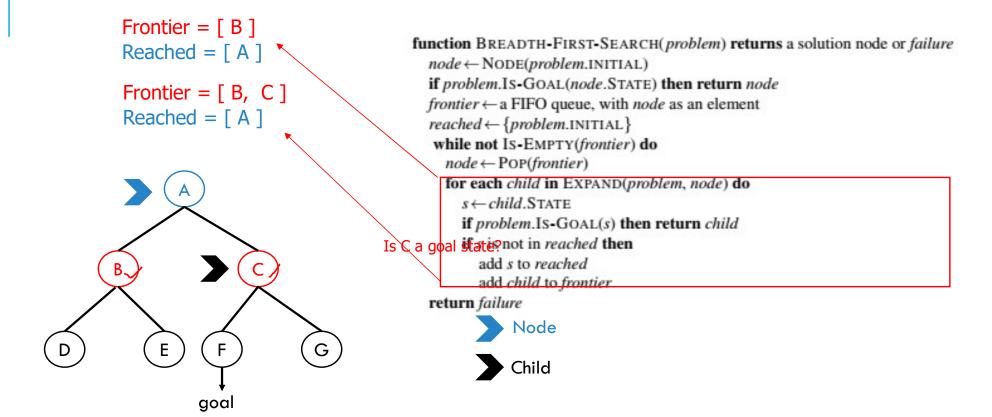
add s to reached

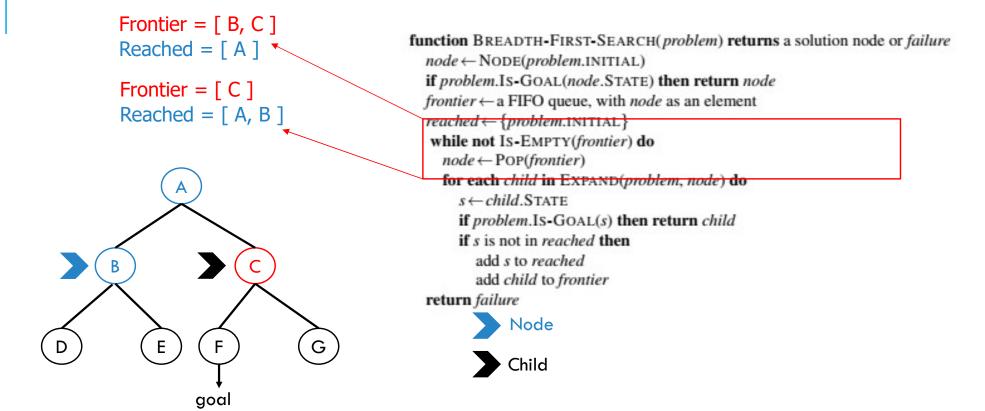
add child to frontier

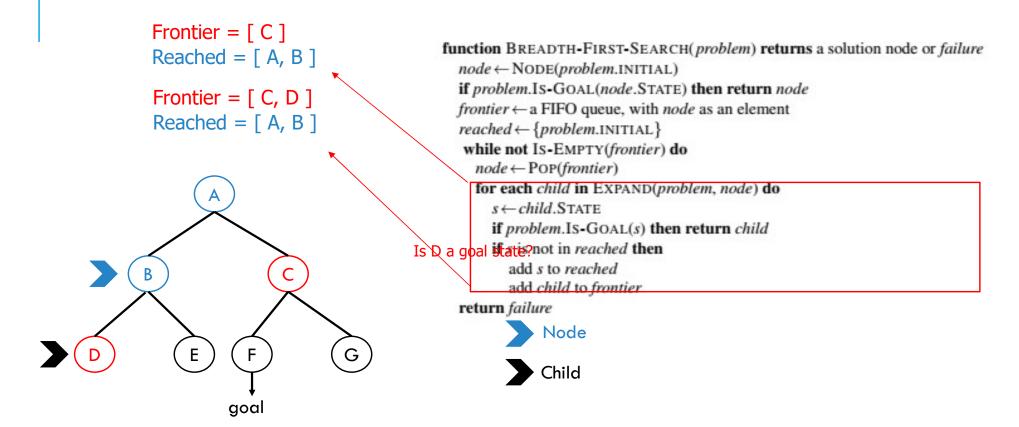
return failure
```

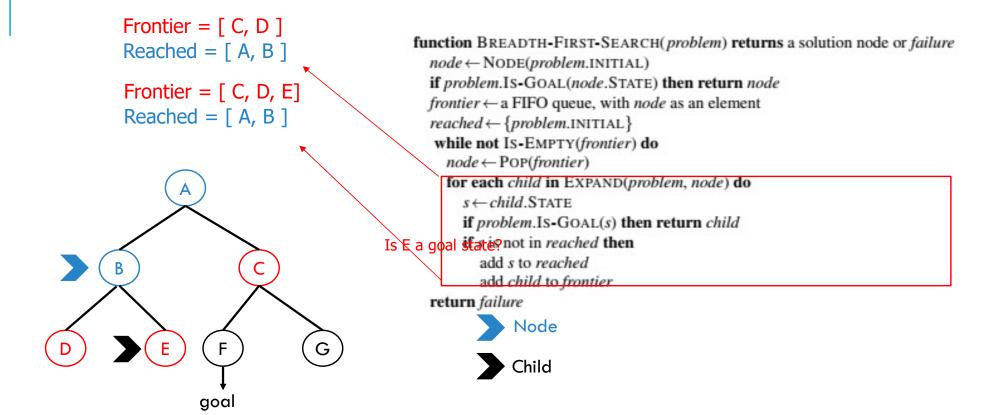


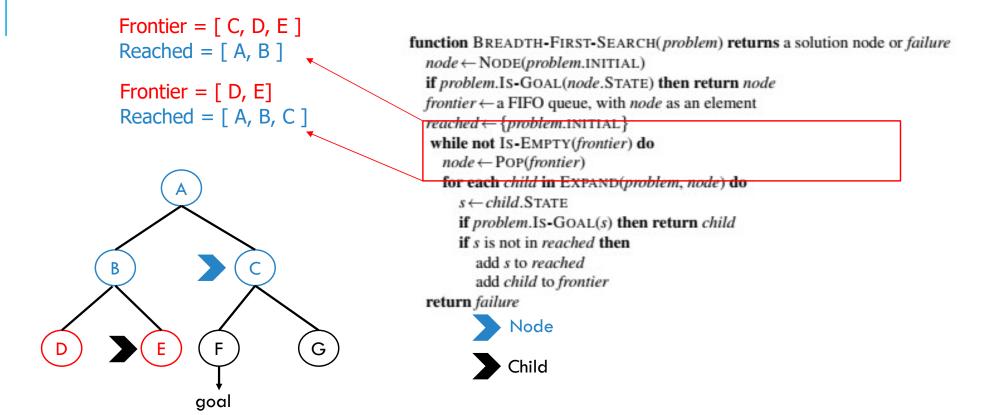


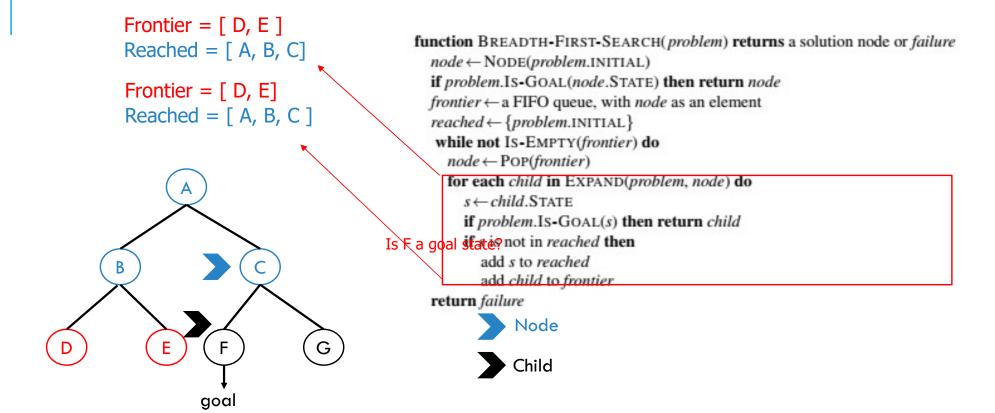




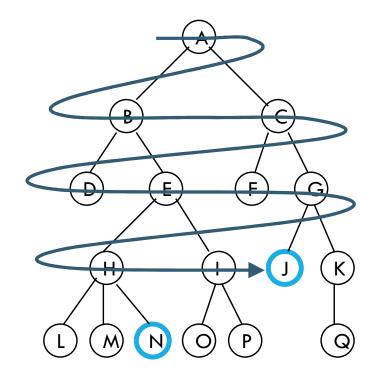








## **BFS: SUMMARY**

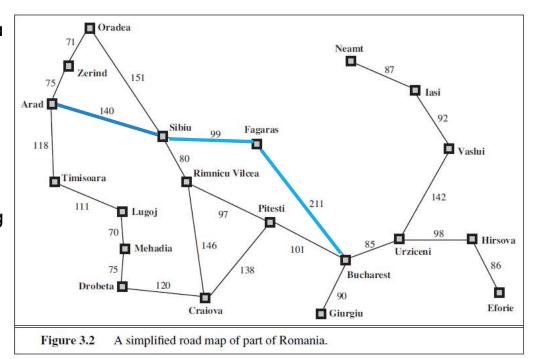


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## **BFS: COMPLETENESS**

Completeness: Yes (guaranteed to find a solution if there exists one)

if the shallowest goal node is at some finite depth d, breadth-first search will eventually find it after generating all shallower nodes (provided the branching factor b is finite).

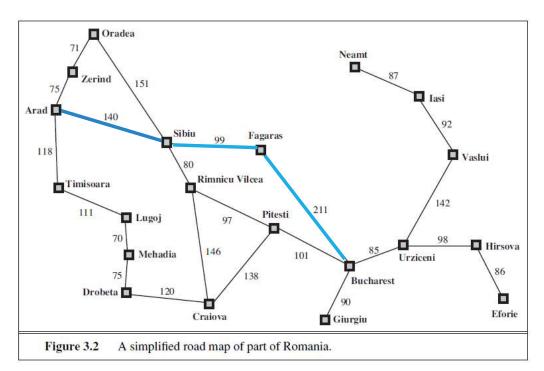


BFS goal path

## BFS: OPTIMAL?

Not necessarily optimal

Only optimal if every action has same cost.



BFS goal path

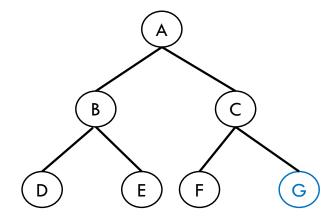
## BFS: TIME COMPLEXITY?

Worst complexity is when G is a goal state.

In this case, total number of nodes generated is

$$b+b^2+b^3+\ldots+b^d=O(\overline{b^d})$$
 (The d<sup>th</sup> layer contains nodes much larger than all the nodes in previous layers combined!)

So the time complexity is O(bd)



A binary tree, b=2, d=2

## **BFS: SPACE COMPLEXITY?**

There will be  $O(b^{d-1})$  nodes in the explored set and  $O(b^d)$  nodes in the frontier.

So the space complexity is  $O(b^d)$ , i.e., it is dominated by the size of the frontier.

Exponential time complexity can be accepted but exponential space complexity is BAD!

