

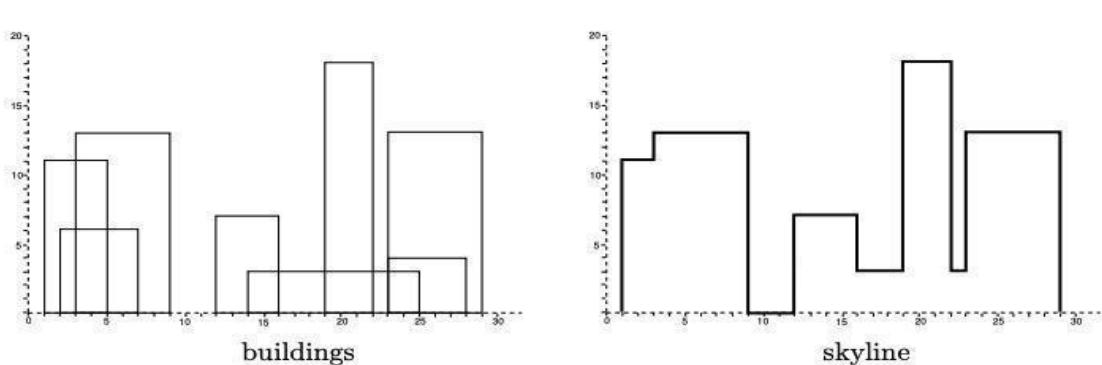
# Homework 5

## All ungraded

- For each part below, solve the following recurrences by giving tight  $\Theta$ -notation bounds in terms of  $n$  for sufficiently large  $n$ , and briefly describe the steps. Assume that  $T(n)$  is a positive and non-decreasing function of  $n$  and represents the running time of an algorithm. In some cases, we shall need to invoke the Master Theorem with a generalization of case 2:

If the recurrence  $T(n) = aT(n/b) + f(n)$  is satisfied with  $f(n) = \Theta(n^{\log_b a} \log^k n)$  for some  $k \geq 0$ , then  $T(n) = \Theta(n^{\log_b a} \log^{k+1} n)$ .

- $a.$   $T(n) = 49T(n/7) + n^2 \log n$
  - $b.$   $T(n) = 4.001 T(n/2) + n^2 \log^4 n$
  - $c.$   $T(n) = 100T(n/2) + n^{50}$
  - $d.$   $T(n) = 10T(n/2) + 2^n$
  - $e.$   $T(n) = 2T(\sqrt{n}) + \log n$
- Assume that you have a blackbox that can multiply two integers. Describe an algorithm that when given an  $n$ -bit positive integer  $a$  and an integer  $x$ , computes  $x^a$  with at most  $O(n)$  calls to the blackbox. (a ‘blackbox’ simply means a ‘function’ whose implementation specifics are *unknown*)
  - Solve Kleinberg and Tardos, Chapter 5, Exercise 3.
  - A city's skyline is the outer contour of the silhouette formed by all the buildings in that city when viewed from a distance. A building  $B_i$  is represented as a triplet  $(L_i, H_i, R_i)$  where  $L_i$  and  $R_i$  denote the left and right  $x$  coordinates of the building, and  $H_i$  denotes the height of the building. Describe an  $O(n \log n)$  algorithm for finding the skyline of  $n$  buildings.  
For example, the skyline of the buildings  $\{(3, 13, 9), (1, 11, 5), (12, 7, 16), (14, 3, 25), (19, 18, 22), (2, 6, 7), (23, 13, 29), (23, 4, 28)\}$  is  $\{(1, 11), (3, 13), (9, 0), (12, 7), (16, 3), (19, 18), (22, 3), (23, 13), (29, 0)\}$ . (Note that the  $x$  coordinates in a skyline are sorted)



- Solve Kleinberg and Tardos, Chapter 5, Exercise 5.