

## DP on Stocks

1 Best Time to Buy & Sell Stock

arr[] = { 7, 1, 5, 3, 6, 4 }

'If you are selling on  $i^{\text{th}}$  day,  
you buy on the minim price from  $1st \rightarrow (i-1)$

# Buy and Sell Stock-II

## Problem Statement

[Suggest Edit](#)

You have been given stock values/prices for N number of days. Every i-th day signifies the price of a stock on that day. Your task is to find the maximum profit which you can make by buying and selling the stocks.

### Note :

You may make as many transactions as you want but can not have more than one transaction at a time i.e, if you have the stock, you need to sell it first, and then only you can buy it again.

Sample Input:

val = { 1 2 3 4 5 6 7 }

you can buy as many times I want

and sell as many times

## Understanding and Approach:

prices = { 7, 1, 5, 3, 6, 4 }

↑  
b s      b s  
4      7, 3 = 7

Recursion:

$f(i, \text{buy})$

0	1	2	3	4	5
7	1	5	3	6	4

↑  
 $f(0, 1)$

↓  
starting on 0<sup>th</sup> day  
with buy, what  
maximum profit?

# Recursion Solution:

$f(i, \text{buy})$

base/line  $\rightarrow$  if  $(i = n)$  return 0,

if (buy)  $\max$  profit =  $\left\{ \begin{array}{l} \text{price}[ind] + f(i+1, 0) \end{array} \right\}$  take  
 $\left\{ \begin{array}{l} 0 + f(i+1, 1) \end{array} \right\}$  not take

else  
profit =  $\max \left\{ \begin{array}{l} \text{price}[ind] + f(i+1, 1) \\ 0 + f(i+1, 0) \end{array} \right\}$

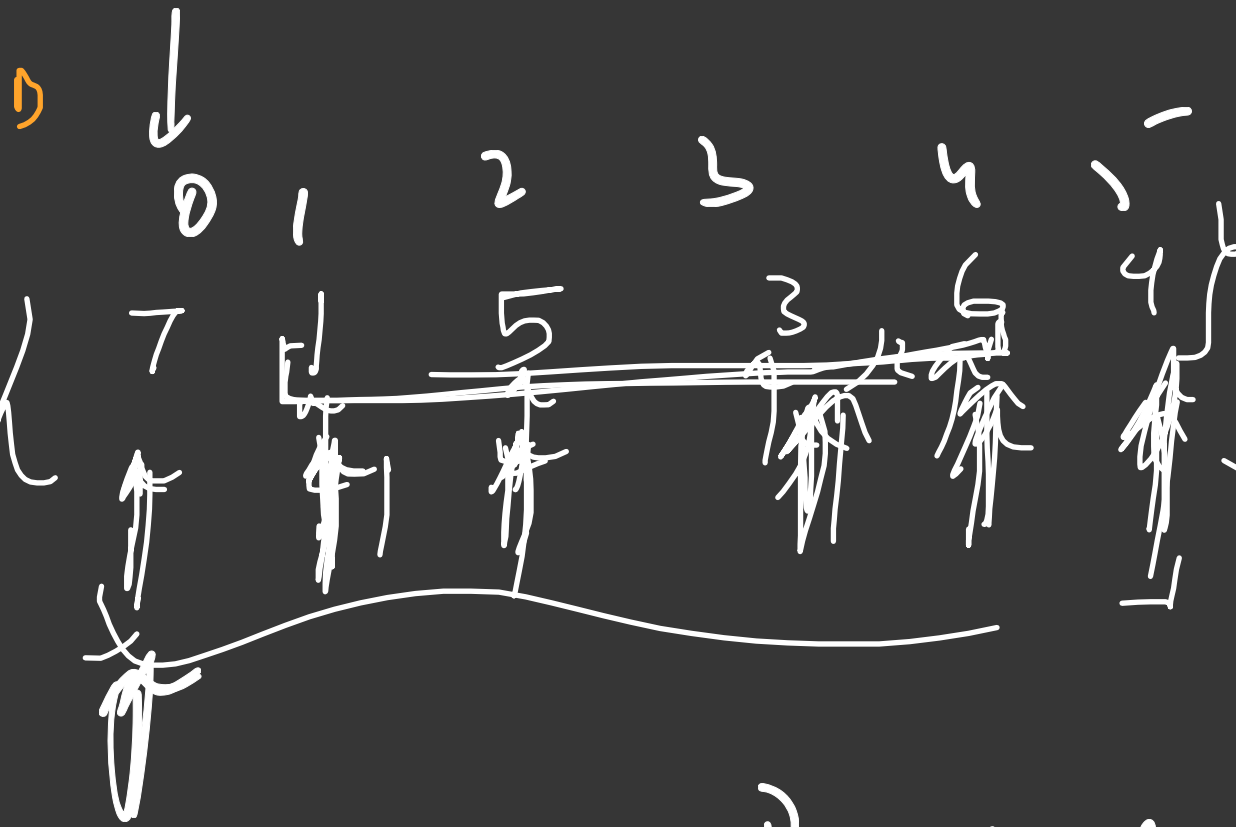
return profit.

Approach:

{ 7 1 5 3 6 4 }

↑ ↑  
3 4

at any and as many times  
and sell it any day  
as many time



5 Perms  
Try all ways

Recursion

Profit-max Subseq. in . prices or not take.

① Express all in terms of index

$f(i, buy)$

sell  
price[i] (if

② Try all ways

if (Buy) { sell [i] = buy [i] }  
profit = max

③ Max of all ways

else {

④ Base Case-

}

# Recursion Tree

- T#

buy



$$f(1, 0)$$

$$f(i, buy)$$

not buy 0 +

$$f(1, 1)$$

$$f(\underline{0}, \underline{1})$$

max profit

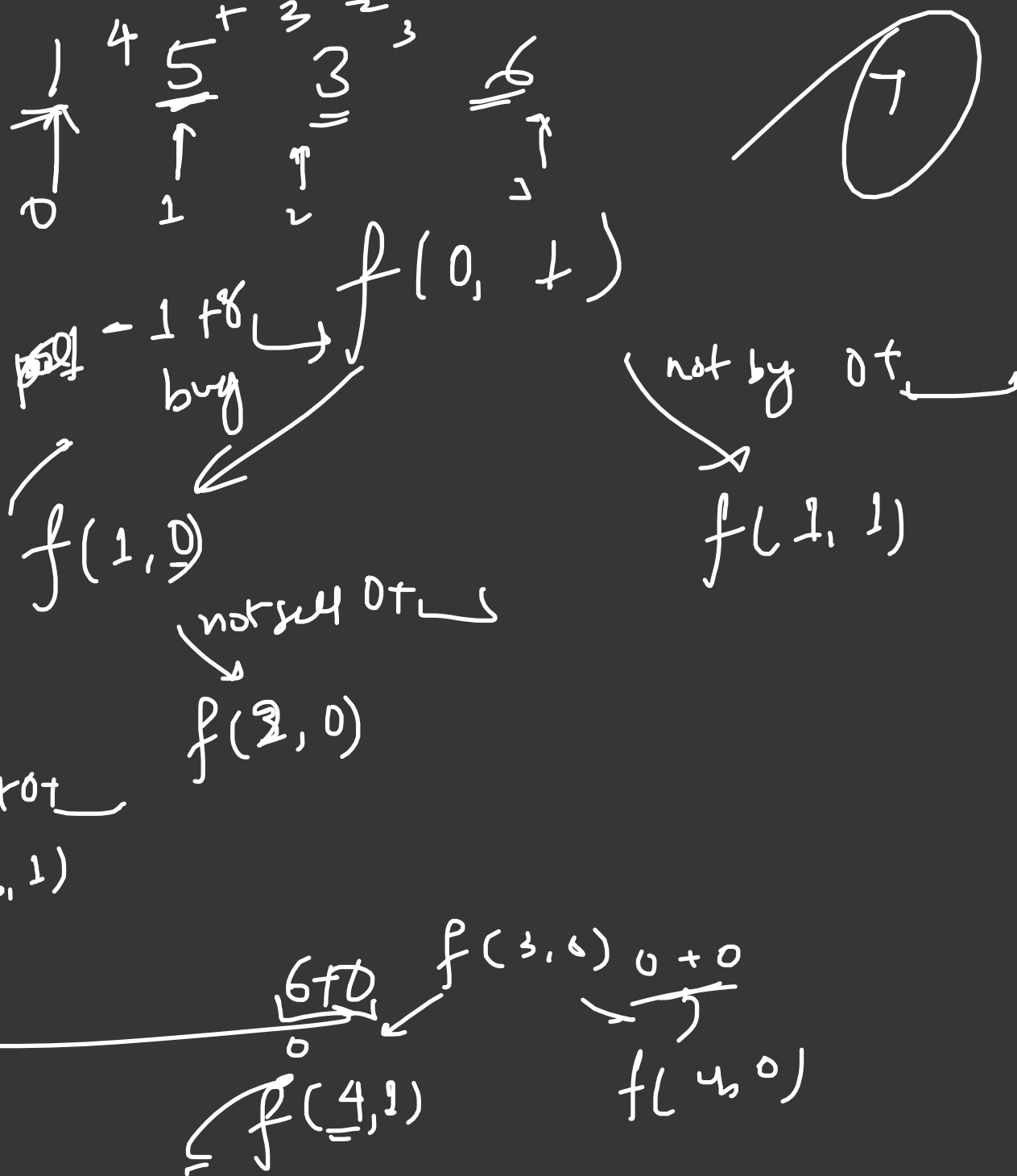
buy true // not buy  
profit = max  $\left( \underbrace{-\text{price}[i] + f(i+1, 0)}, \underbrace{0 + f(i+1, 1)} \right);$

buy false // sell // not sell  
profit = max  $\left( \text{price}[i] + f(i+1, 0), 0 + f(i+1, 1) \right);$

return profit



Test Case:



Memorization

$f(i, buy)$

$dp[i][buy]$

Tabulation:

Base Case  $\rightarrow$

$i$

$buy$

Copy the recursion

$dp[n+1]$

Base Case:

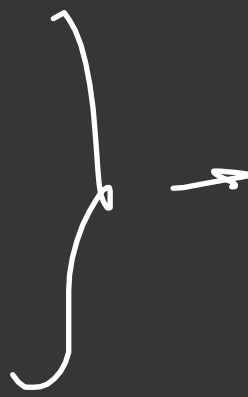
$dp[n] = 0$

for ( $i = n-1; i \geq 0; i--$ )  
for ( $int buy = 0; buy < 1; buy++$ )

# Space Optimisation:

36 → space Optimisation with  
variable and note

37  
28  
39  
40



7:30

8 → 10 → Code Chef Contest

11:00 → 12:00  
pp.

12.00 → 1:00 → Android

1 → 3:00 → DSA sheet





