

## Boolean Evaluation to true (similar to MCM)

Ex expression:  $T \mid T \& F$

OR  $\rightarrow \mid$   
AND  $\rightarrow \&$   
XOR  $\rightarrow \wedge$

### Problem Statement

[Suggest Edit](#)

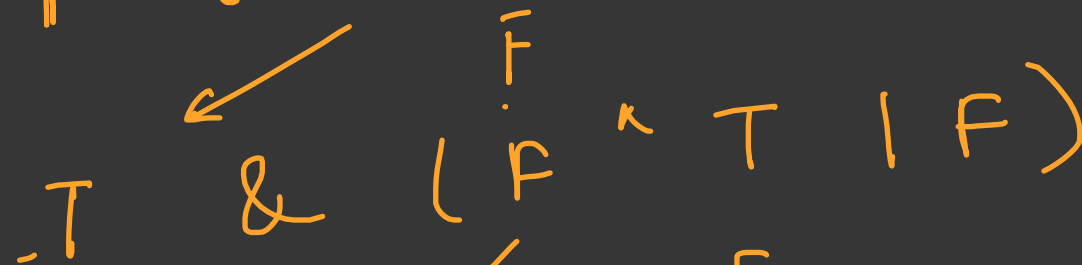
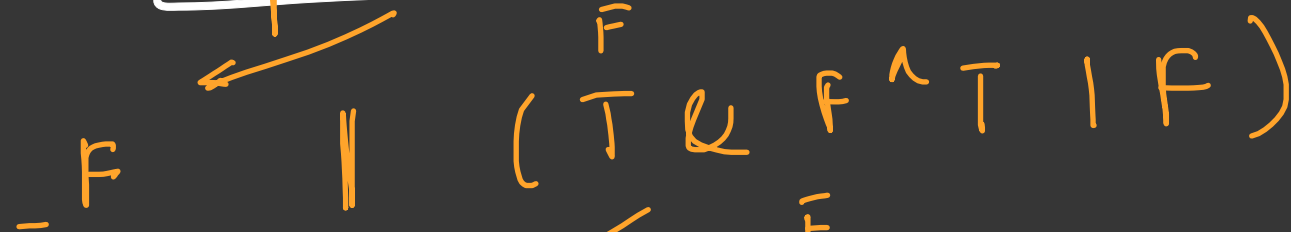
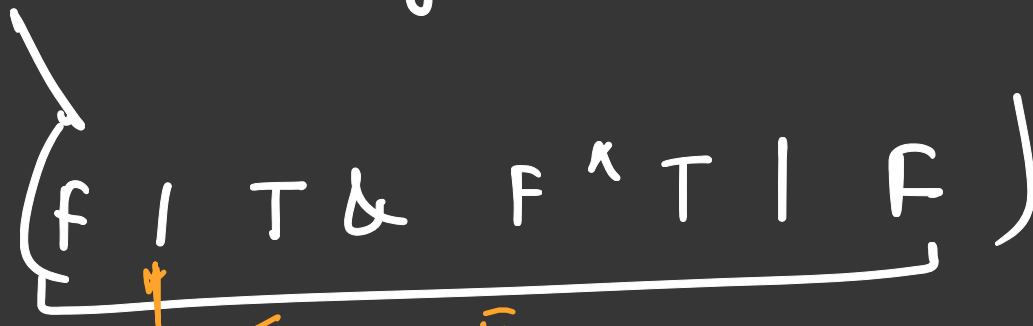
You are given an expression 'exp' in the form of a string where operands will be : (TRUE or FALSE), and operators will be : (AND, OR or XOR).

Now you have to find the number of ways we can parenthesize the expression such that it will evaluate to TRUE.

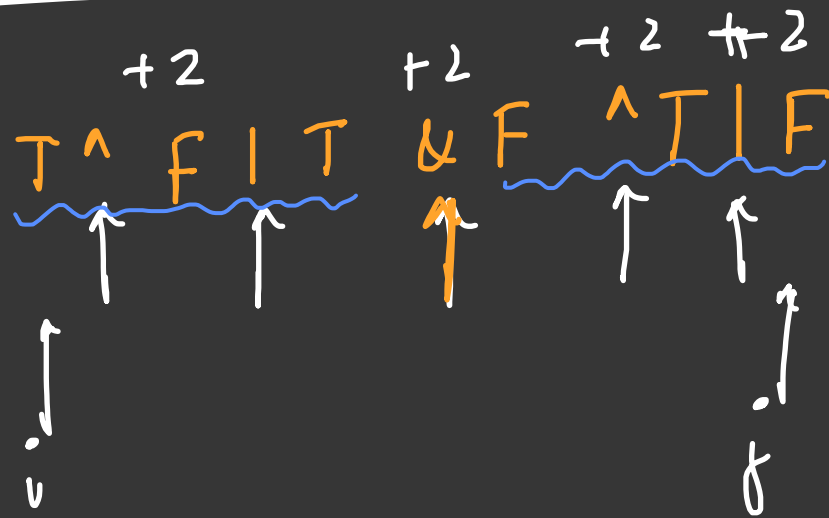
As the answer can be very large, return the output modulo 1000000007.

eg:  $F \mid T \wedge F$

$(F \mid T) \wedge F \rightarrow T \wedge F \rightarrow T$  (1)  $\rightarrow$  way  
 $F \mid (T \wedge F) \rightarrow F \mid T \rightarrow T$  (2)  $\rightarrow$  way  
} 2 ways



We use partition algorithm

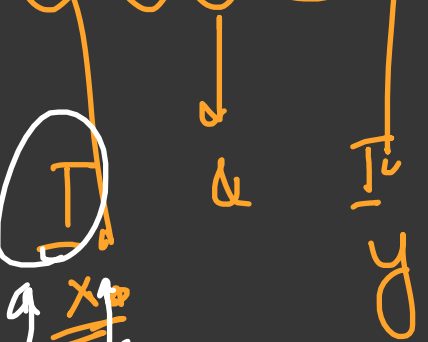


$part(x)$   
left & right

for ( $k = i+1; k < j; k += 2$ )

& (T)

left & right

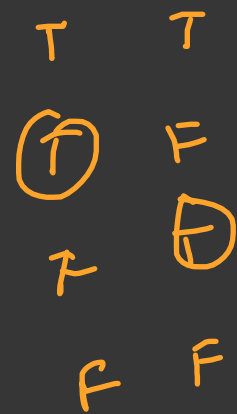


~~x \* y~~



|

left | right

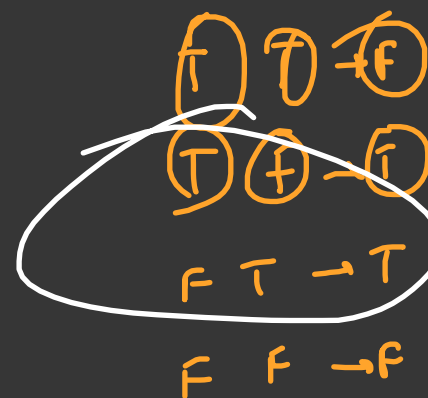
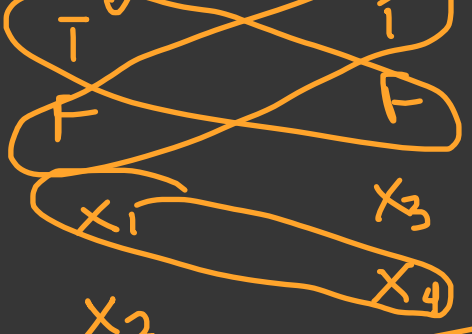


(OR)



^

left ^ right



## Recursive Relation

$f(0, n-1, 1)$

$f(i, j, \text{isTrue})$

if  $(i > j)$  return 0,

if  $(i == j)$  {

if  $(\text{isTrue} == 0)$  return  $\text{arr}[i] == \text{'T'}$ ,

If we looking for True

if we are looking for false.

else return  $\text{arr}[i] == \text{'F'}$ ;

for  $(k = i+1; k < j; k++)$

leftT =  $f(i, k-1, 1)$

leftF =  $f(i, k-1, 0)$

right T =  $f(k+1, i, l)$

right F =  $f(k+1, j, 0)$

if ( $arr[k] == 'Q'$ )

if (isTrue) ways +=  $LT \times RT$ ,

else ways +=  $(T \times RF) + (LF \times RT) + (LF \times RF)$ ,

else if ( $arr[k] == 'I'$ )

if (isTrue) ways +=  ~~$LF \times$~~   $(LT \times RT) + (LF \times RF) + (LT + RF)$

else ways +=  $(F \times RF)$ ;

else { if (isTrue) ways +=  $(RF \times LT) + (LF \times RT)$

else ways +=  $(RF \times LF) + (RT \times LT)$ ,

return ways;







