

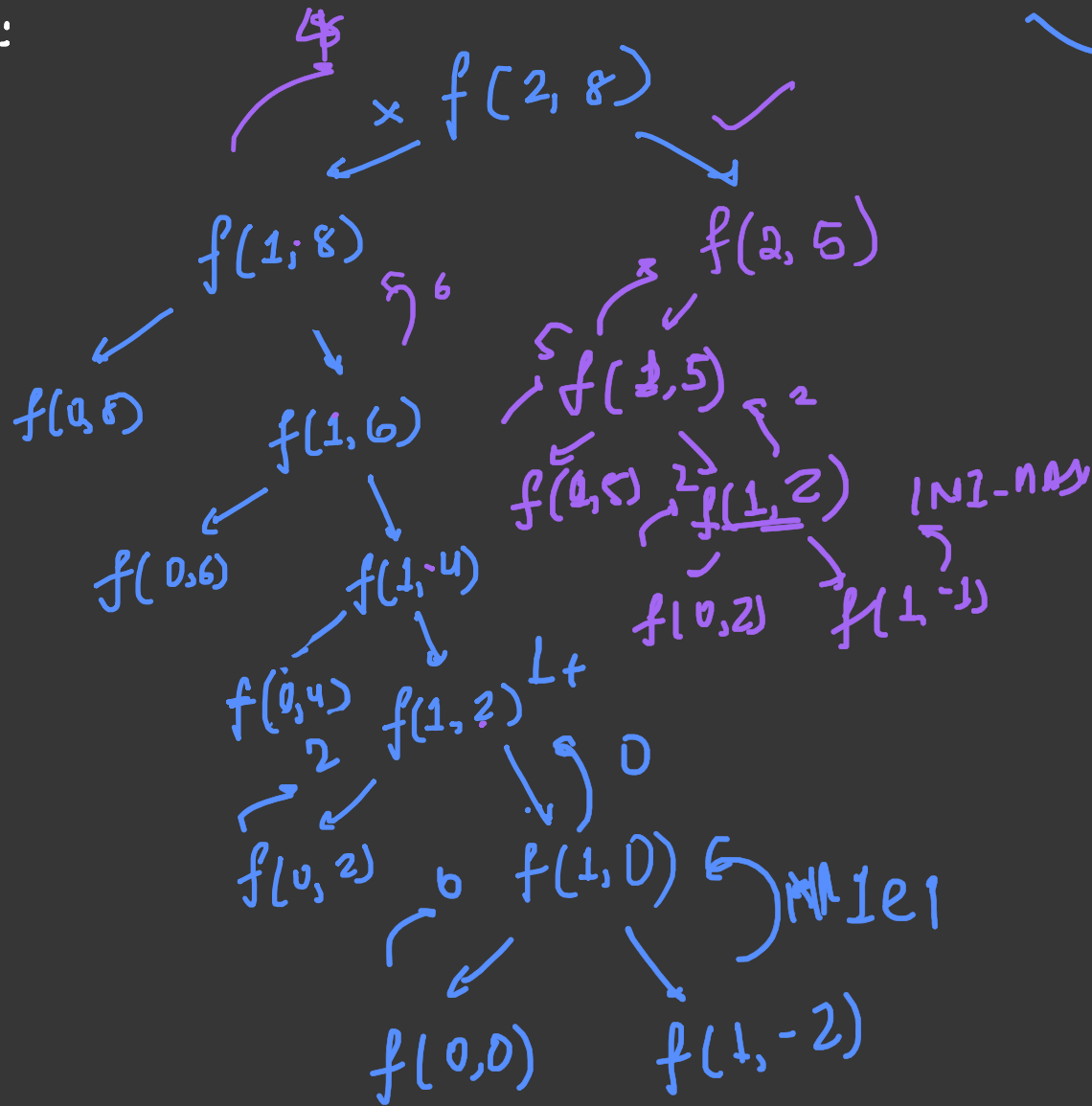
Minimum Coin

Recursion Tree:

arr = {1, 2, 3}

tar = 8

6



arr[i] >= target
pick

Recursion Relation:

inc []

back case
if (i == 0) {
if (arr[0] == target) return target / arr[0];
else return 1e9;

not pick = fun(i-1, target);

pick = 1e9;

if (arr[i] >= target) pick = 1 + fun(i, target - arr[i]);

return min(pick, not pick);

T.C: $\gg O(2^N)$ Exponential

S.C: $O(N)$ → Auxillary Space

arr[i] > target

This is different from normal subsequence as we have given infinite value.

Memorization:

we have 2 changing variable we use $\text{dp}[n][k]$
 \downarrow Target

Just add

if ($dp[i][K]_j = -1$) return $dp[i][K]_j$

- recursion relation

```
return dp[i][k] = min(pick, notpick);
```

$$T.C.: O(N \times T)$$

S.L. : $O(N^2) + \underbrace{O(N)}_{\text{auxiliary space}}$

Tabulation:

So what is the base case of recursion if (arr[0] < target) = target/arr[0];
↓
means for every target at arr[0]
if it divisible

Base Case:

```
for (T = 0; T <= target; T++)
```

```
if (T % arr[0] == 0)
```

```
dp[0][T] = T/arr[0];
```

```
else dp[0][T] = -1;
```

Nested loop

```
for (i = 1; i < n; i++)
```

```
for (T = 0; T <= target; T++)
```

```
notPick = dp[i-1][T];
```

```
pick = -1;
```

```
if (arr[i] >= T) pick = dp[i-1][T - arr[i]];
```

$dp[i][1] = \min(\text{pick}, \text{notpick})$

Space Optimise this,

prev,

curr