Find Majority element > 1/3

Problem Statement

Suggest Edit

Given an array 'A' of 'N' integers, find the majority elements of the array.

A majority element in an array 'A' of size 'N' is an element that appears more than floor(N / 3) times.

Note: The floor function returns the number rounded down to the nearest integer.

Note: Return the array of majority elements in sorted order.

Example:

Input: 'N' = 9 'A' = [2, 2, 1, 3, 1, 1, 3, 1, 1]

Output: 1

Explanation: The frequency of '1' is 5, which is greater than floor(N / 3), hence '1' is the majority element.

Explanation of Questions element find and shore attility arril = { 1, 1, 1, 2, 2, 2} n . 6 $\frac{3}{3}$, $\frac{6}{3}$, $\frac{2}{2}$ arr[1"] (1)(1), 2,2,2)

Sorted areorder

Observahm.

$$\eta = 8$$

$$\rightarrow \left[\begin{array}{c} 8 \\ \hline 3 \end{array}\right] = 2$$

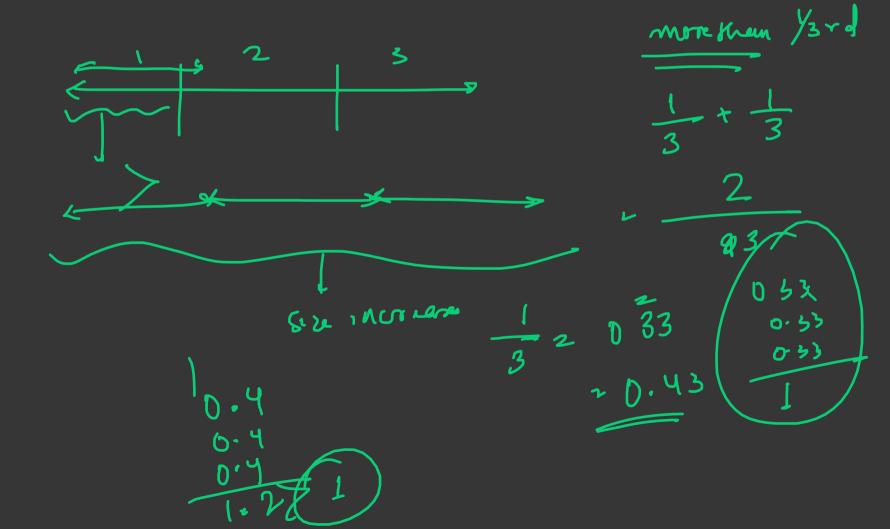
At max how many element can be there in an arroy?

- · At minimial a clement would have 3 Daws ence 20, if 3+3+3 = 9 > circ array
- · We to arrow et 3 element of 3 occurrence the size of the arroy is the but

 It is greaker than original size of arrow
- . So, At wex an well have only 2 element.

eg:
$$99 = \left[\frac{99}{3}\right] = \frac{33}{3}$$

We have said that the occurerence should be more than 1/3rd of array size n. 2 1/3



Boute force:

m = 6

J. (. = 0 (1)

Better Approach:

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if (it. seems > n/s)

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representations.

Ophmal Approach:

Modified verim - 2 majority element thongs.

So, we know that we can have only 2 elements which can be more than n/3 times.

So, Idea is to find 2 majority element to the array and that if the 2 majority elements to the array and that if the 2 majority elements are greater than n/s time.

Pseudo wale:

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- To avoid overlopping of majority clement.
ele L, eliz, court 1 = 0, count 2 10.
 or (120, 1<n, 1++)
           if ( count 1 20 % d' err[i) & 1 ele 2)
               elct a arr[.].
                cout = 1 = 1
           Similar fore ele 2.
            Moore Voly Also
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Cont

$$\frac{T \cdot C}{2N} := \frac{1}{2} \left(\frac{1}{2} \left(\frac{1}{2} \right) \right) + O(N) = O(2N) = O(N)$$

$$\mathbb{S}_{C}:\longrightarrow\widetilde{\mathbb{Q}(1)}$$