# Search in Matrix

| | | | | |
|---|---|---|---|---|
| 5 | 3 | 4 | 8 | 7 |
| 1 | 9 | 2 | 6 | 10 |
| 12 | 11 | 14 | 15 | 13 |
| 20 | 19 | 17 | 18 | 16 |
| 22 | 23 | 21 | 25 | 24 |

target = 14

Linear search

1 D → array

Binary Search

O(log N)

1 2 3 4 5

start end

start end

target = 4

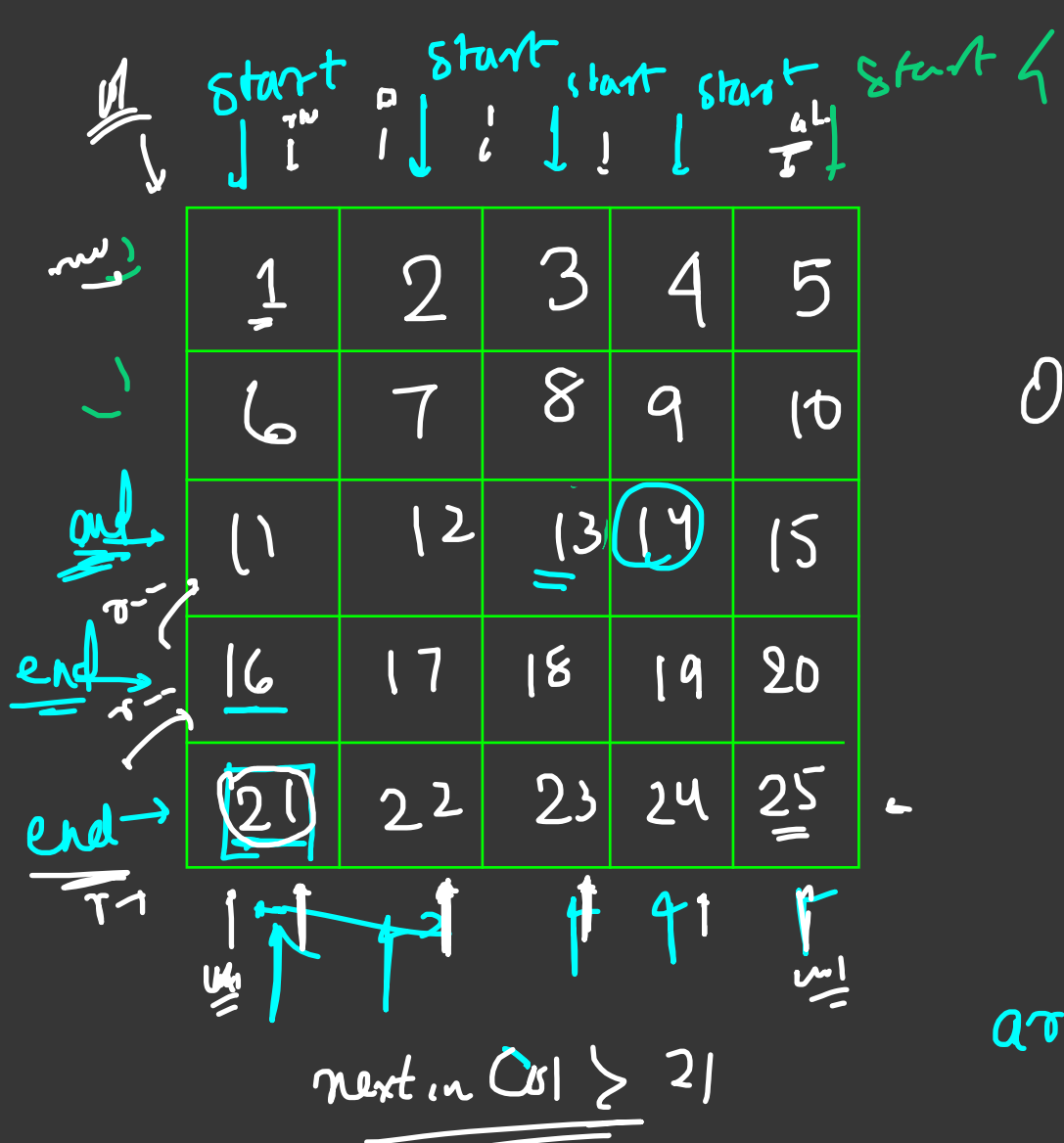
arr[mid] > target , target

return true

| | | | | |
|---|---|---|---|---|
| 5 | 3 | 4 | 8 | 7 |
| 1 | 9 | 2 | 6 | 10 |
| 12 | 11 | 14 | 15 | 13 |
| 20 | 19 | 17 | 18 | 16 |
| 22 | 23 | 21 | 25 | 24 |

→

| | | | | |
|---|---|---|---|---|
| 1 | 2 | 3 | 4 | 5 |
| 6 | 7 | 8 | 9 | 10 |
| 11 | 12 | 13 | 14 | 15 |
| 16 | 17 | 18 | 19 | 20 |
| 21 | 22 | 23 | 24 | 25 |




row → sorted

col → sorted

$O(\log(m \times n)) + O(n)$

$O(n)$

start = 0 (col)

end = n-1 (row-1)

target = 14

arr[start][end] > K {

end--;

arr[start][end] < K

Pseudo code:

Sort the matrix (O(N log N))

```
int r = n-1, c = 0,
while( start < n && end > 0)
    if ( arr[ r ][ c ] == k)
        return true,
    else if ( arr[r][c] > k)
        r--,
    else
        c--,
```

# Optimised Binary Search :

| | 0 | 1 | 2 | 3 | 4 |
|---|---|---|---|---|---|
| 0 | 1 | 2 | 3 | 4 | 5 |
| 1 | 6 | 7 | 8 | 9 | 10 |
| 2 | 11 | 12 | 13 | 14 | 15 |
| 3 | 16 | 17 | 18 | 19 | 20 |
| 4 | 21 | 22 | 23 | 24 | 25 |

(2,1)

→ Flatten the array in 1D-array hypothetically

Size of 1D array = $n * m$

0 ——————————→

25 size

mid

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17

17 18 19 20 21 22 23 24

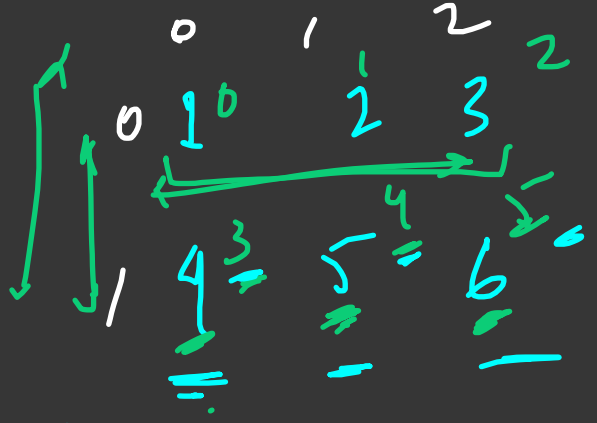
18 19 20 21 22 23 24 25

e

Start = 0

end = (n × m) − 1

mid = (25 + 0) / 2 = 12




$n = 2$

$m = 3$

$n \times m = 6$

4 → array

matrix

col

first index completely divisible by col size

0 3 6 9

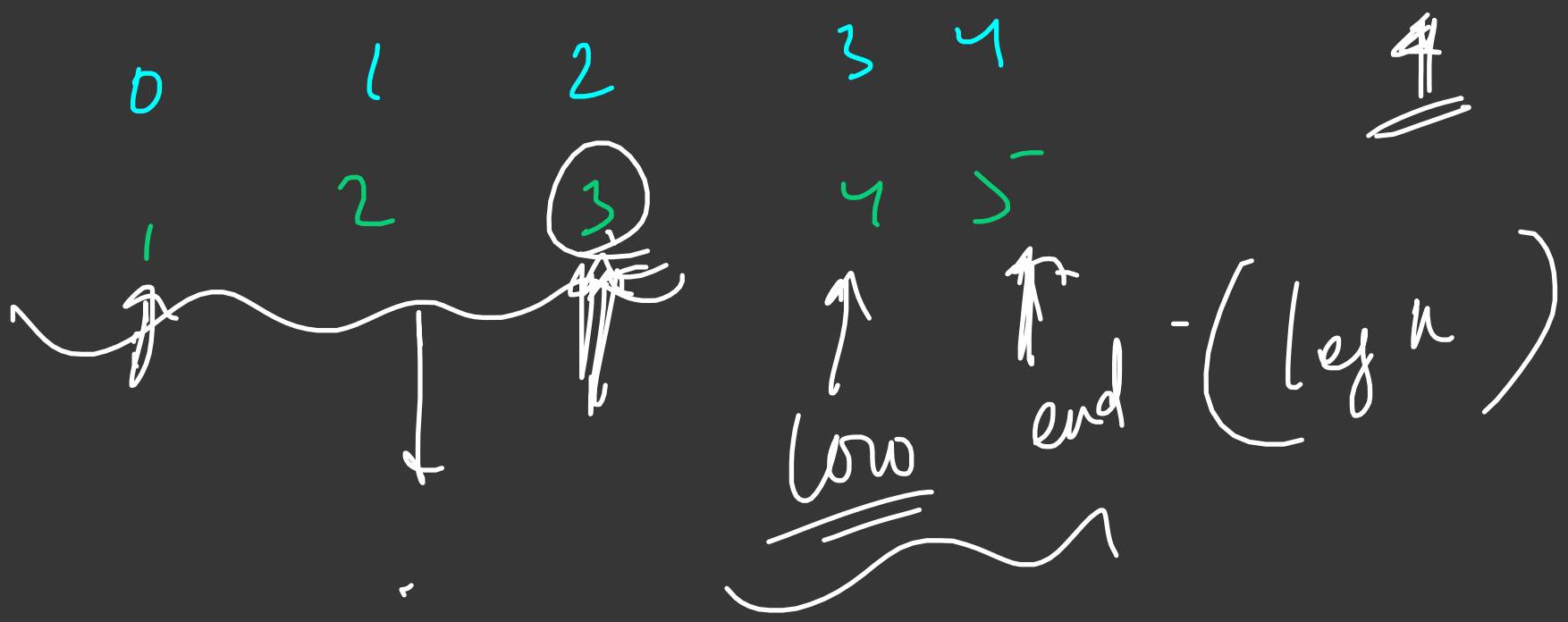
col = ind % m ;

row = ind / m ;

Binary

low mid high

$\frac{5+0}{2} = 2$

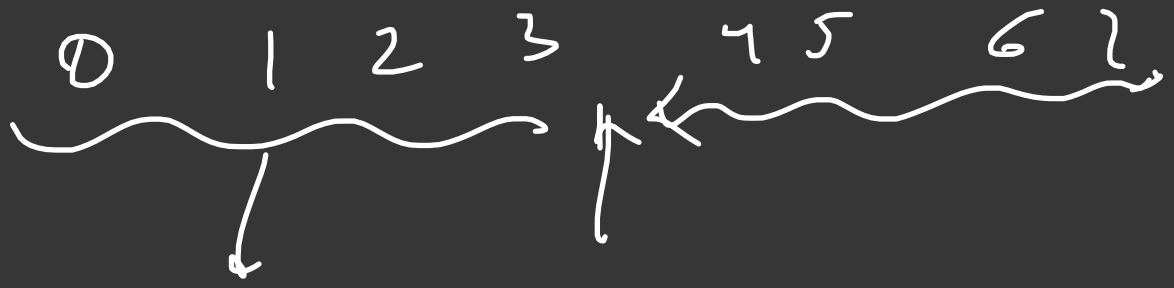


Matrix

$$T.C = \log_2 (n \times m)$$

$$S.C = O(1)$$

# Binary Search Time Complexity analysis




$$O(\log(n))$$