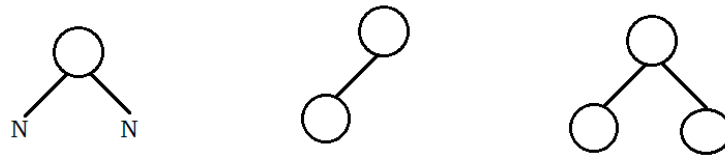
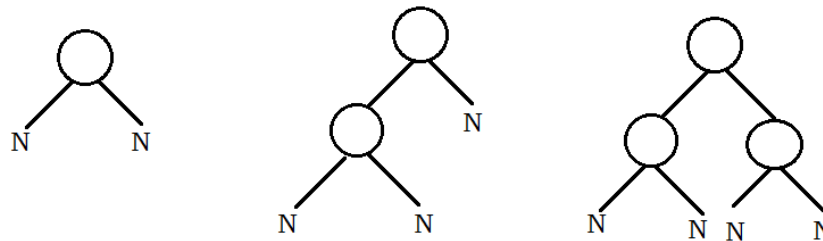


# Binary Search Tree

- BST is a non-linear data structure.
- One data element is connected to multiple elements in this structure.
- We represent the data in nodes.
- Every node has 3 fields
  1. Data field
  2. Left child
  3. Right child
- In BST, every node has at most 2 children.



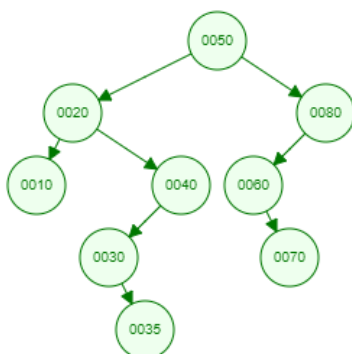
**The tree with N nodes having N+1 null nodes.**



- We store information into BST by comparing with Parent node.
- Least value is connected to left side of Parent node.
- Highest value is connected to right side of Parent node.
- BST not allow duplicates.
- Keys(elements) must be unique to store the data.

**Inserting elements into BST:**

50, 20, 80, 10, 40, 30, 35, 60, 70



**Node structure:** We represent the node using user data type called structure.

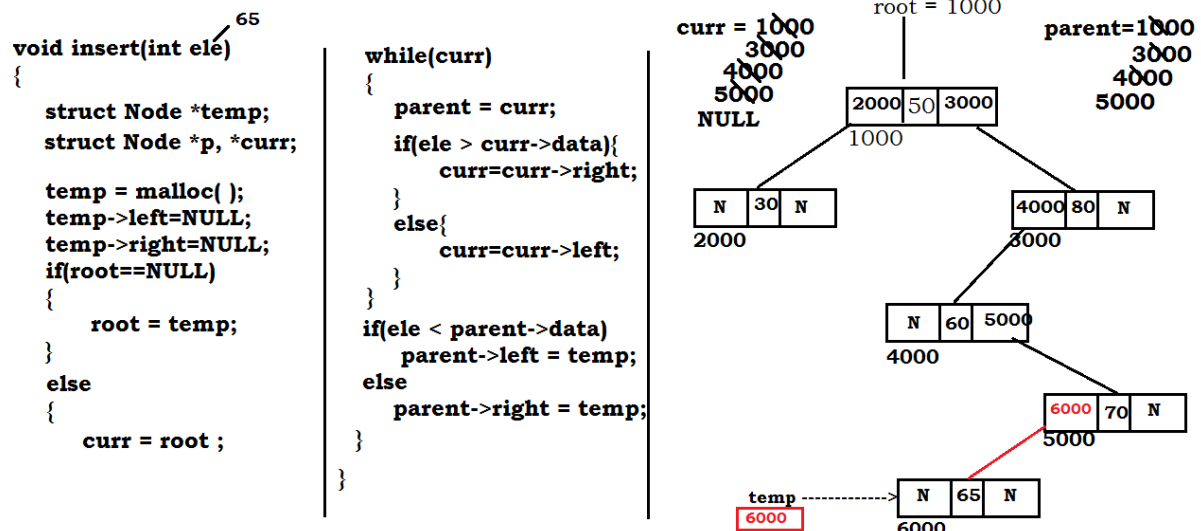
```
struct Node
{
    int data;
    struct Node *left;
    struct Node *right;
};
struct Node *root=NULL;
```

### Operations:

1. Insert
2. Delete
3. Traverse

### Insertion:

- Construct the Node.
- Place the Node data.
- Find the Parent node in the tree.
- Connect to Left or Right depends on the value of new node.



### Node deletion:

- We need to find whether the element is present or not.
- If the element is present, we need to delete the node and re-arrange other nodes.
- If not present, display "Element not Found".

```

delete(int ele)
{
    struct Node *curr,
    *parent ;
    int found=0 ;
    curr = root ;
    while(curr)
    {
        if(curr->data == ele)
        {
            found = 1 ;
            break ;
        }
        else
        {
            parent = curr ;
            if(ele>curr->data){
                curr=curr->right ;
            }
        }
    }
}

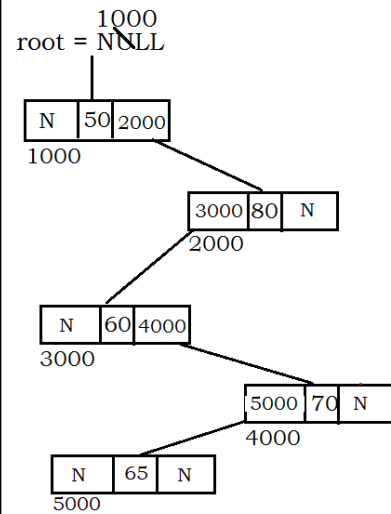
```

```

else{
    curr = curr->left;
}
}
if(!found)
{
    printf("No such element
    to delete \n");
    return ;
}

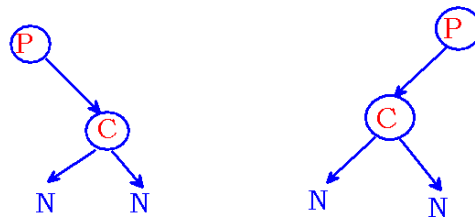
```

**3 cases to remove the element which is present in the BST**



**Case 1:** Deleting the element has No child

The Node has no child



```

if(curr->left==NULL && curr->right==NULL)
{
    if(curr==parent->right)
    {
        parent->right=NULL;
    }
    else
    {
        parent->left=NULL;
    }
    free(curr);
}

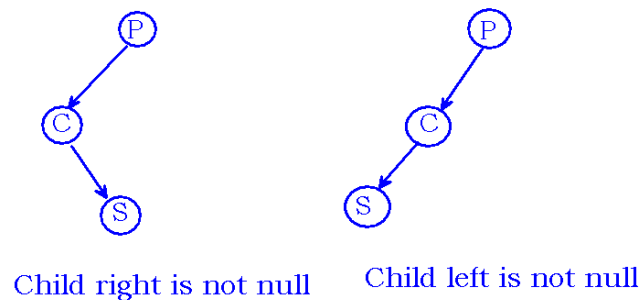
```

**Case 2:** Deleting element that has single child

Node is connected right to Parent



Node is connected to left of Parent

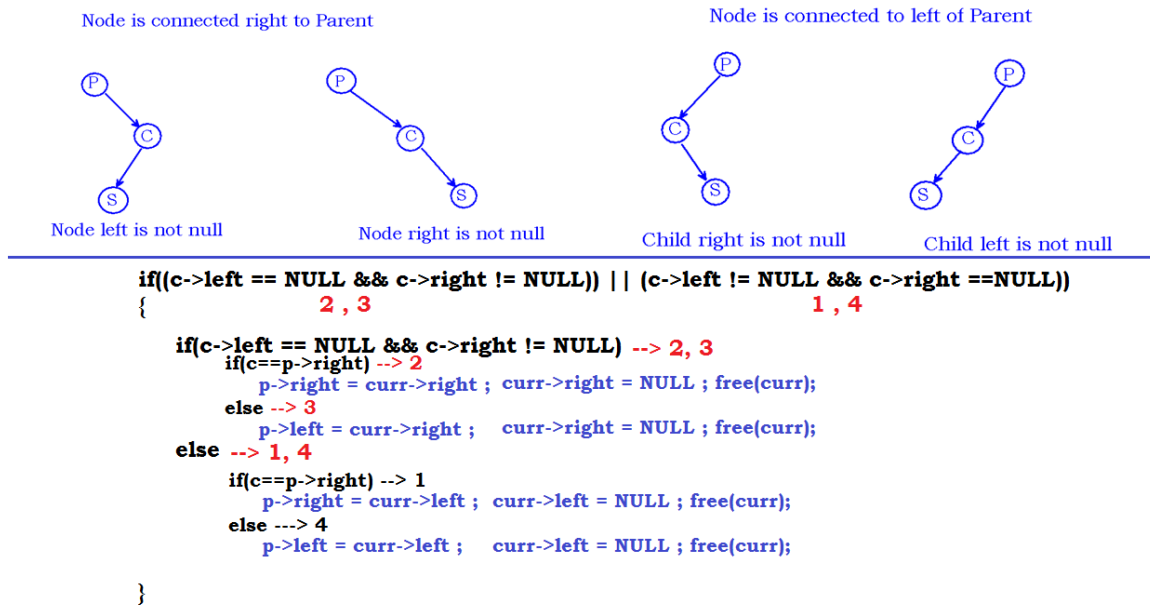


```
if((curr->left==NULL && curr->right!=NULL) ||
    (curr->left!=NULL && curr->right==NULL))
{
    if(curr->left==NULL && curr->right!=NULL) // diag - 2,3
    {
        if(curr==parent->right) // diag - 2
        {
            parent->right = curr->right ;
        }
        else // diag - 3
        {
            parent->left = curr->right ;
        }
        curr->right=NULL;
        free(curr);
    }
    else // diag - 1,4
    {
        if(curr==parent->right) // diag - 1
        {
            parent->right = curr->left;
        }
        else // diag - 4
```

```

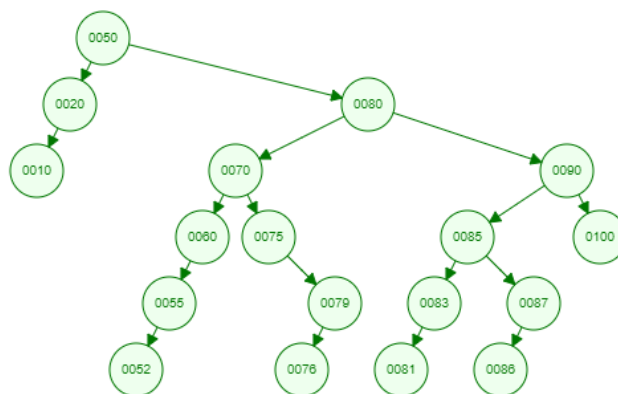
    {
        parent->left = curr->left;
    }
    curr->left = NULL;
    free(curr);
}

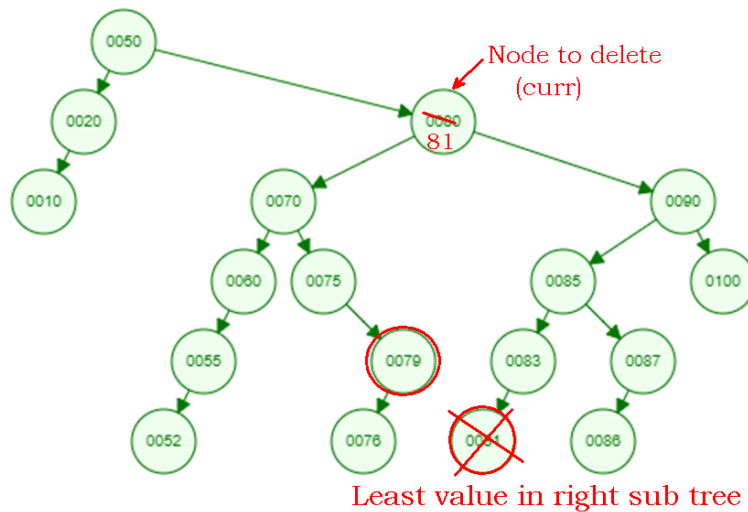
```



**Case 3:** Deleting the element that has 2 children.

- Replace the current node data with
  - Least element in the right sub tree node data or
  - Highest element in the left sub tree node data.
- Removes the data swapped node.





**The Node has 2 children condition:**

```
if(curr->left!=NULL && curr->right != NULL)
{
    Logic...
}
```

**If curr-> right has no left child and right child:**

```
t1 = curr->right ;
if(t1->left == NULL && t1->right == NULL)
{
    curr->data = t1->data;
    curr->right = NULL;
    free(t1);
}
```

**If curr-> right has no left child but right child is present:**

```
if(t1->right!=NULL && t1->left==NULL)
{
    curr->data=t1->data;
    curr->right=t1->right;
    t1->right=NULL;
    free(t1);
}
```

**If curr->right has left child:**

```
t1 = curr->right;
if(t1->left != NULL)
{
    t2 = t1->left;
}
while(t2->left)
{
```

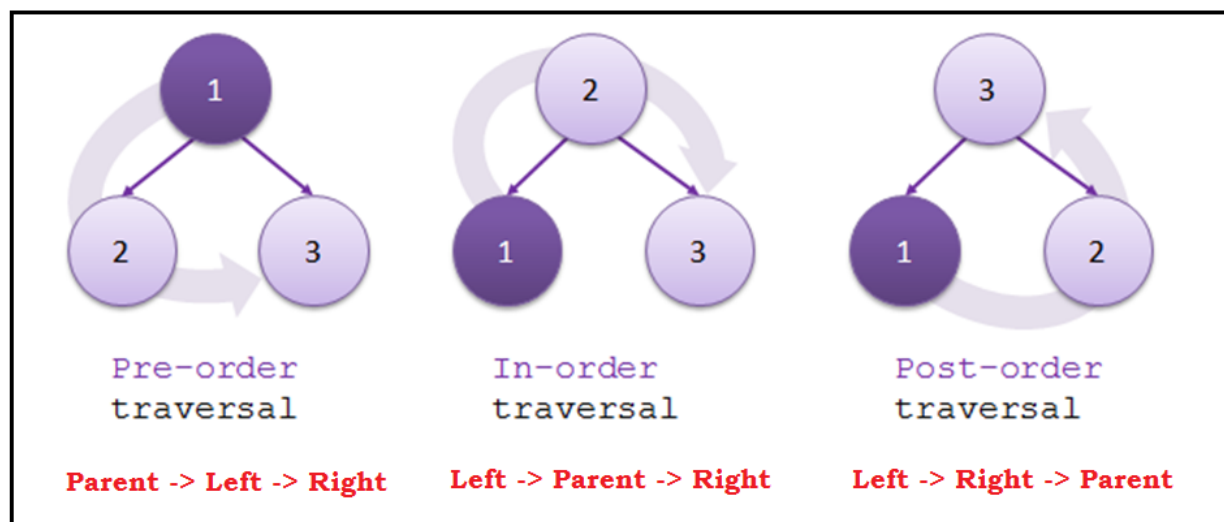
```

    t1 = t1->left;
    t2 = t2->left;
}
curr->data = t2->data;
t1->left = t2->right;
t2->right = NULL;
free(t2);

```

### Traversal:

- **We can traverse the tree in 3 ways**
  - **In order traversal**
  - **Pre order traversal**
  - **Post order traversal**



### Recursion:

- Function calling itself.
- Calling the function from the definition of same function.

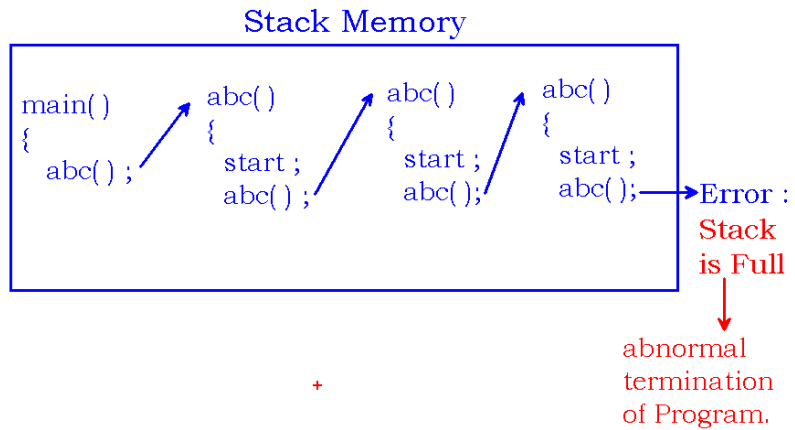
### Notes:

- When we invoke a function, memory will be allocated inside the Stack.
- While the application is executing, if the memory is full, the program terminates abnormally with Runtime error: Stack is Full.

```

1 #include<stdio.h>
2 void abc();
3 int main()
4 {
5     abc();
6 }
7 void abc()
8 {
9     printf("Start \n");
10    abc();
11    printf("End \n");
12 }

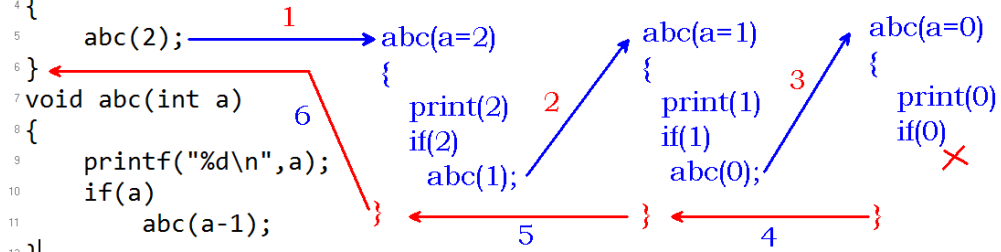
```



```

1 #include<stdio.h>
2 void abc(int);
3 int main()
4 {
5     abc(2);
6 }
7 void abc(int a)
8 {
9     printf("%d\n",a);
10    if(a)
11        abc(a-1);
12 }

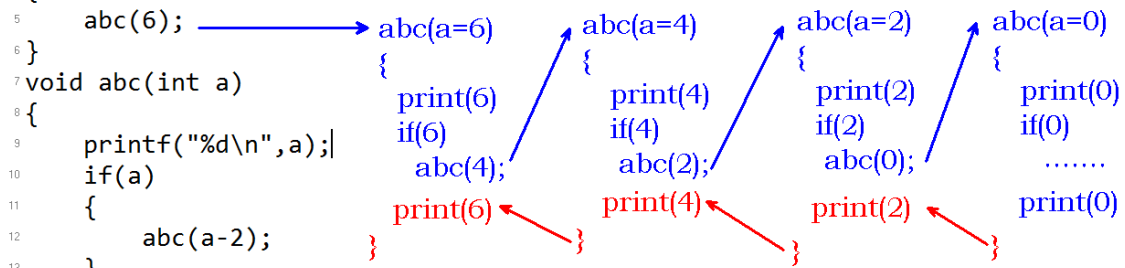
```



```

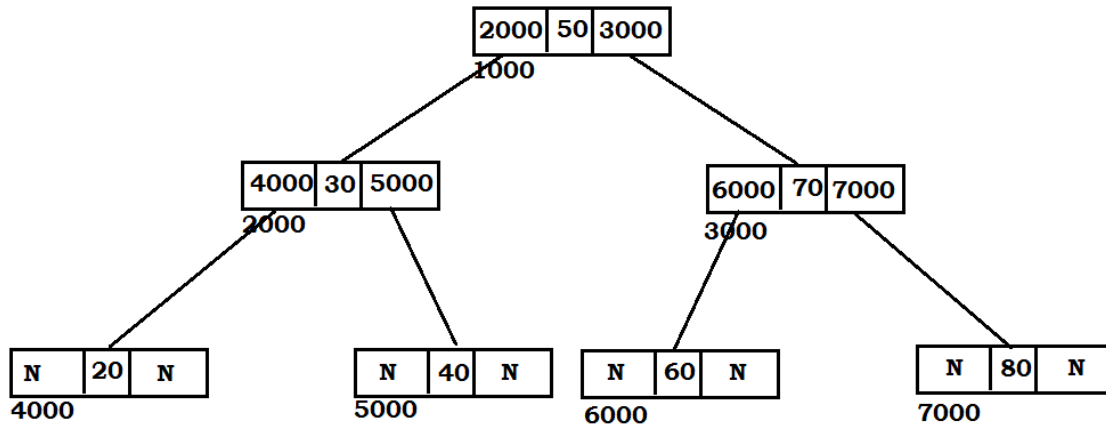
1 #include<stdio.h>
2 void abc(int);
3 int main()
4 {
5     abc(6);
6 }
7 void abc(int a)
8 {
9     printf("%d\n",a);
10    if(a)
11    {
12        abc(a-2);
13    }
14    printf("%d\n",a);
15 }

```





## BST Traversal:



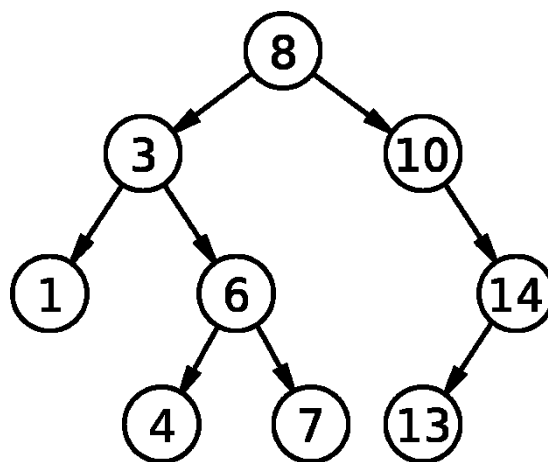
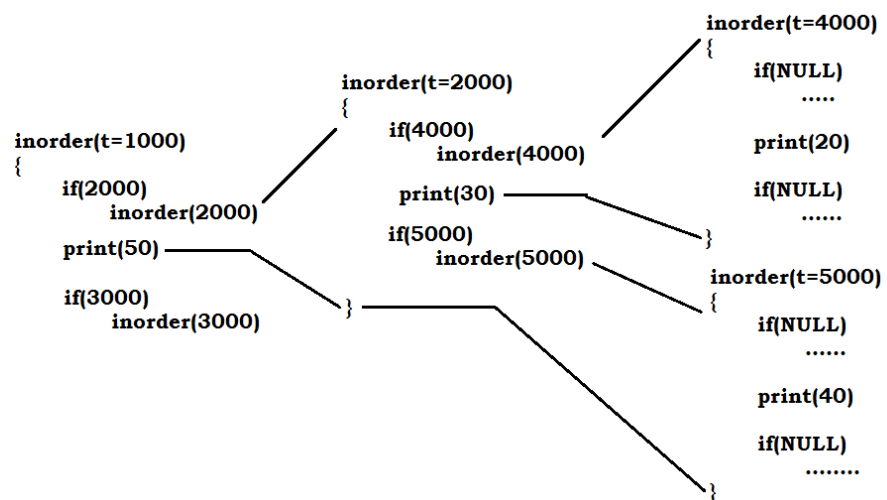
```

PrintInOrder()
{
    if(root==NULL)
        BST is empty;
    else
        inorder(root);
}

inorder(struct Node* t)
{
    if(t->left)
        inorder(t->left);

    print(t->data);

    if(t->right)
        inorder(t->right);
}
  
```



InOrder:  
 PreOrder:  
 PostOrder:

Elements : 6,9,2,8,4,0,7,1,6,3  
Construct BST and Traverse Post order.