

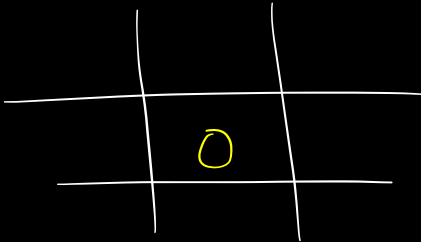
Today's Agenda :-

- 1) How to implement undo feature
- 2) Class Diagram of Tic Tac Toe
- 3) checkWinner() algo.

* Pointers: -

- 1) Design discussions are subjective

Global undo: → undo | remove the last move that was played on the board.



Move → which player made a turn at which location on the board.

$$\begin{array}{c} \hookrightarrow f(\text{player}, \text{location}) \\ \Downarrow \\ (x, c) \end{array}$$

Undo
↓

$$P_1 \quad P_2 \quad P_3 \quad P_1 \quad P_2 \quad P_3 \quad \dots$$

move will be dependant on (player, x,y)

we are storing moves so if players wants more than 1 undo

Ist way

①

List < Move >

Stack

$\left[\left\{ \begin{matrix} P_1 \\ 0 \\ 0 \end{matrix} \right\}, \left\{ \begin{matrix} P_2 \\ 1 \\ 2 \end{matrix} \right\}, \left\{ \begin{matrix} P_1 \\ 1 \\ 1 \end{matrix} \right\}, \dots \right]$

x

O(1)

→ Store moves in a list

→ Remove the last move from the list

→ Make that cell empty in the board as well.

P ₁		
	P ₁	P ₂



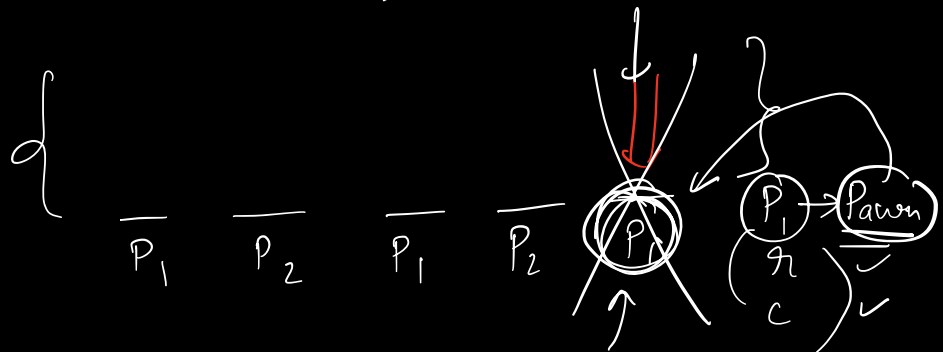
nextPlayer will update to that player.

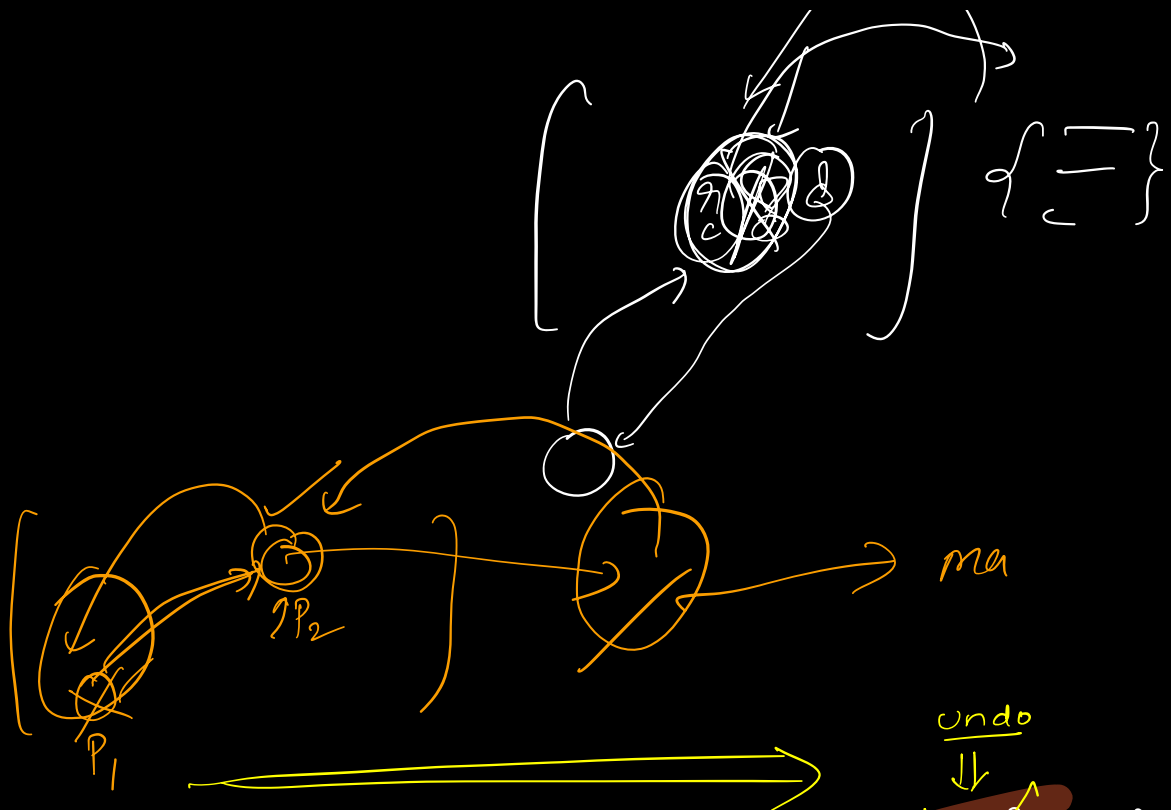
II

Chess

↳ In games like Chess, it's not easy

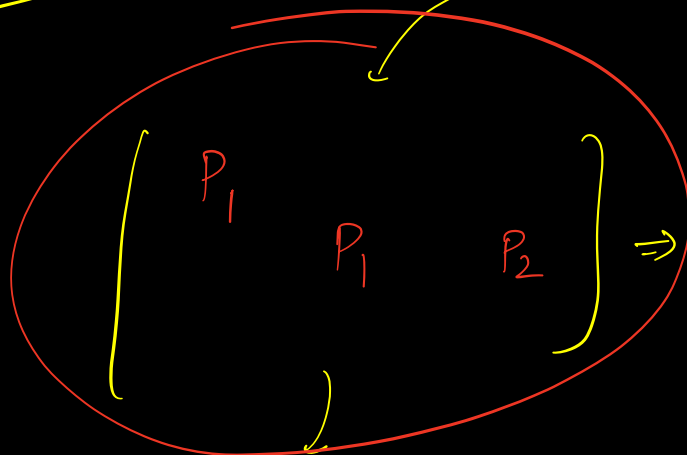
Chess





II hist < Move > : $\left[\begin{Bmatrix} P_1 \\ 0 \\ 0 \end{Bmatrix} \begin{Bmatrix} P_2 \\ 1 \\ 2 \end{Bmatrix} \begin{Bmatrix} P_3 \\ 1 \\ 1 \end{Bmatrix} \begin{Bmatrix} P_2 \\ 2 \\ 2 \end{Bmatrix} \dots \right]$

TC - higher



This move was never played ??
initial

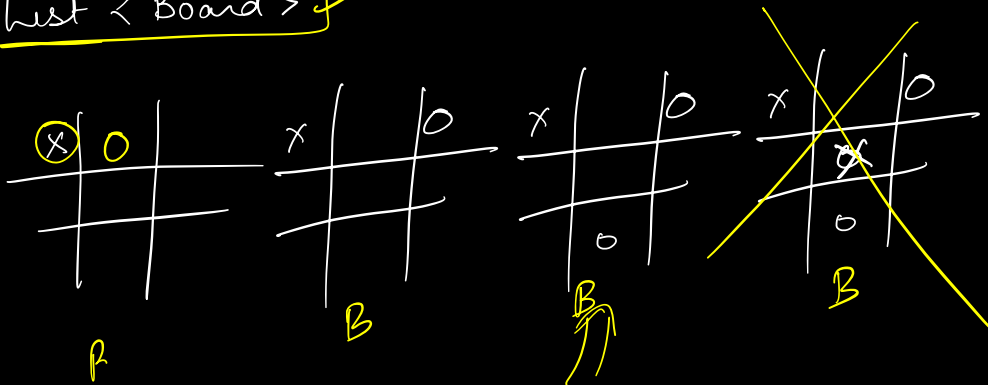
fresh board

exactly at our required slot

- ① Remove last move from list
- ② Clear the board
- ③ Execute all $(N-1)$ moves on empty board.

III Store the Snapshot of the game after every move.

Snapshot
List < Board > ✓



new state, we are done.

Easy to implement

Good TC

Bad SC \Rightarrow Space is v.v.v. cheap.

prefer
Time \gg Space

Maintainable code V-Opt TC
Bad SC VS V-Op SC
Good enough TC Complex

Class Diagram :-

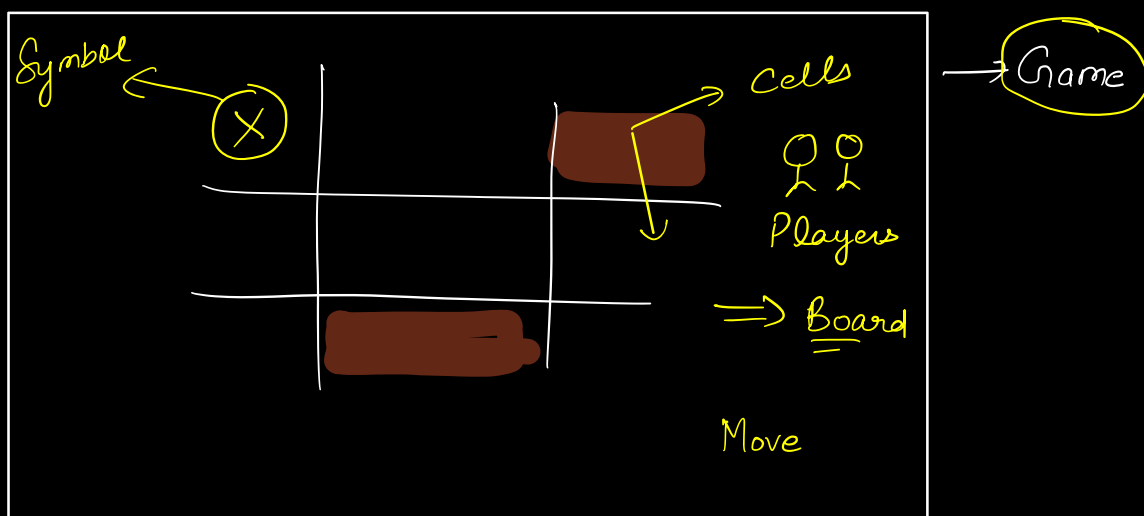
How to identify entities

identify Nouns

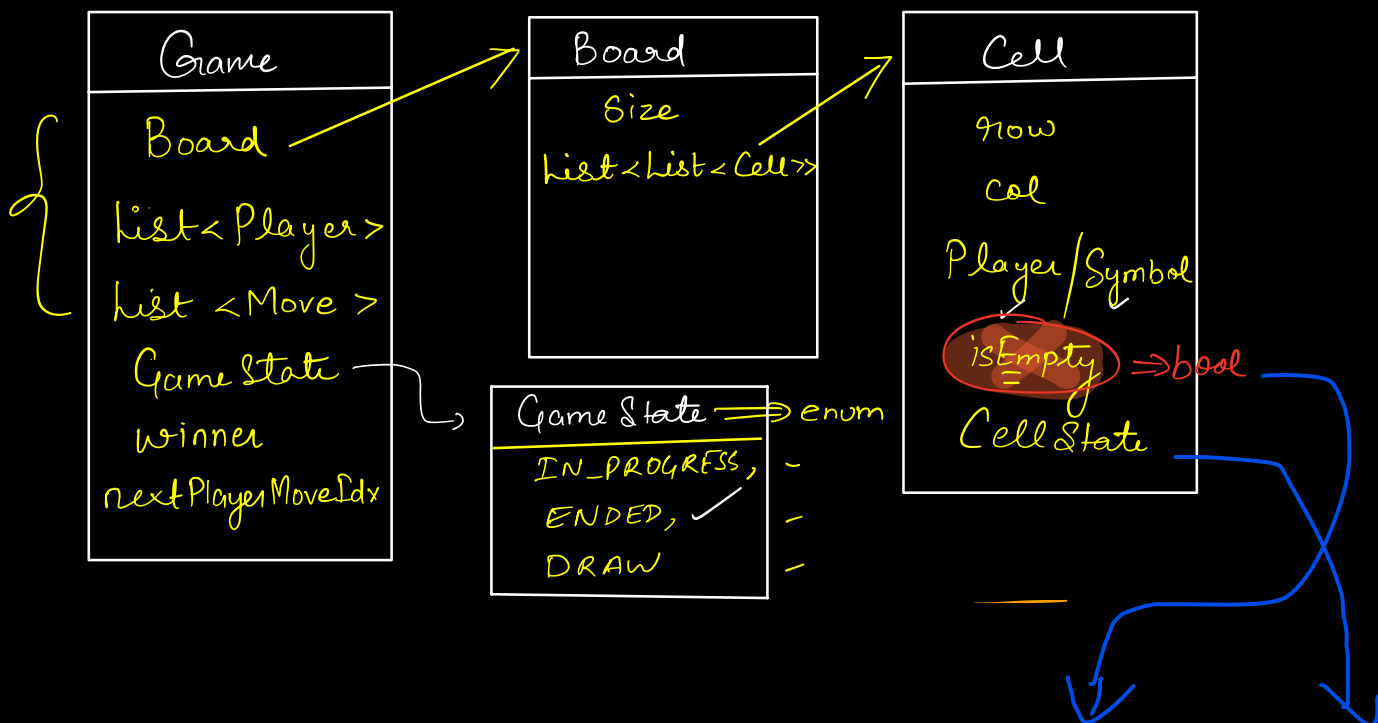
from Req. Gath.

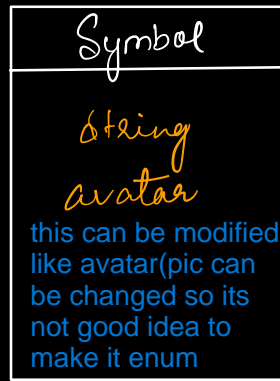
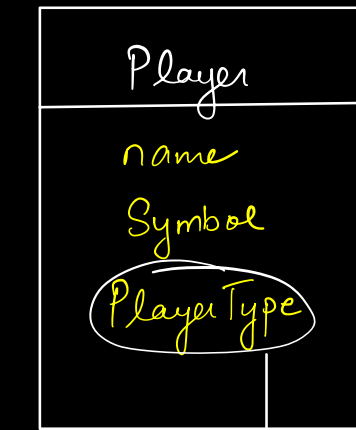
Visualization

1) Game 2) Board 3) Cells



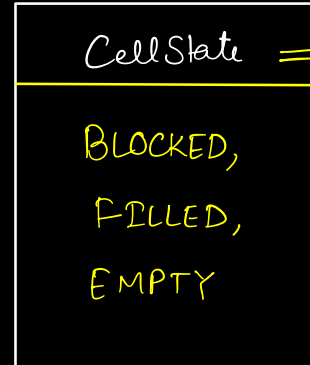
whenever you feel if any data type is not primitive or non-primitive and new class is required then immediately create another class





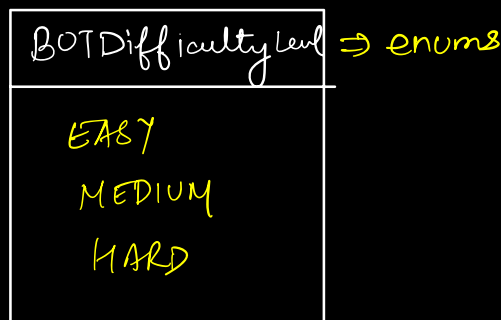
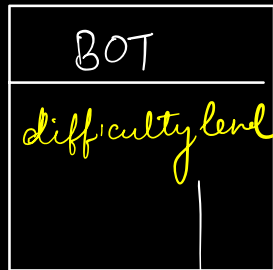
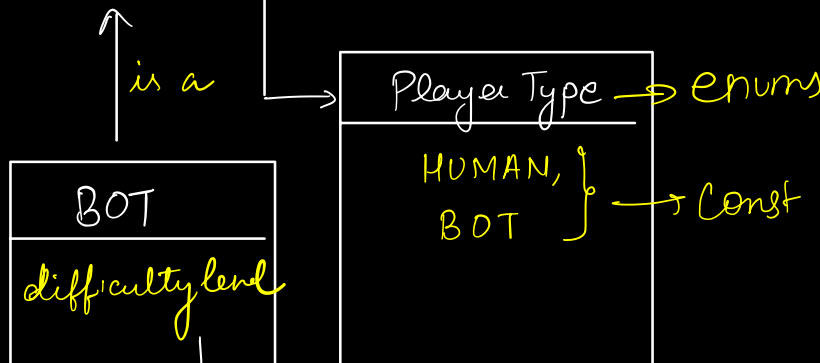
Boolean attr are not extensible

Instead of booleans, you can use enums.



enum

if another requirement comes like few cells are blocked and we can not play there.. so one boolean datatype is not enough.. so make enums



this can be optional if you're sure you have player put in this cell.

or even

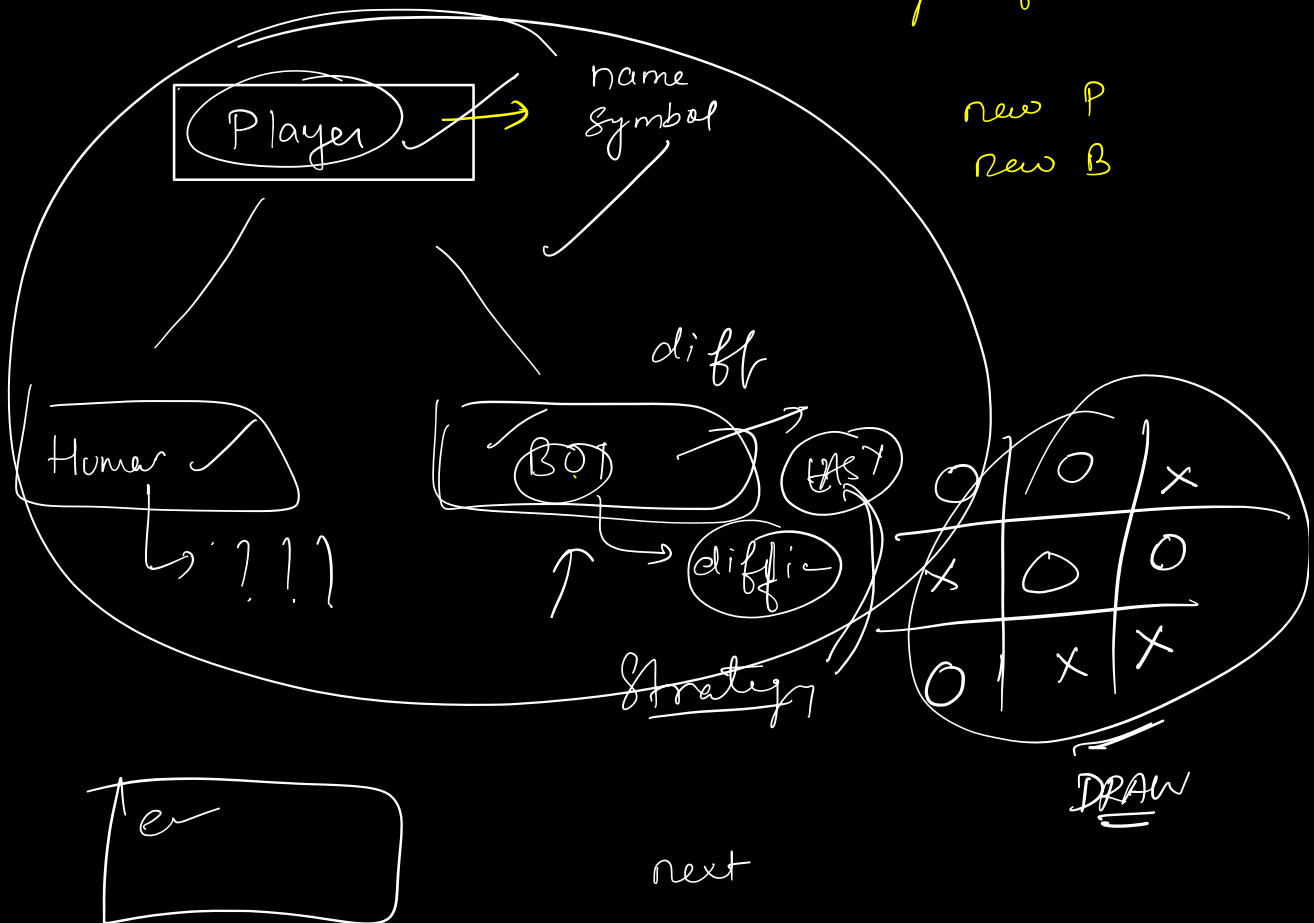
state = en
 C = 0
 n = 0

Cell c1 = new Cell()

Board
 hist < hist < Cell > >



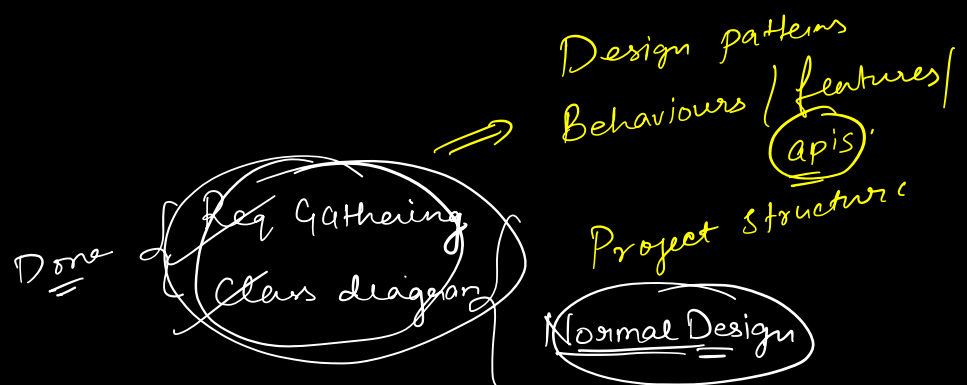
Objects of cell class



create Game — makeMove()
└
initialize the game

Game Builder ⇒ perform all validations in the game.

Strategy.



Pimp

Check if someone won the game ? \Rightarrow algo

20 min

Assume somebody has the won the game.

Game should ideally stop as soon as somebody wins.

After any move, the game can end. Should ideally check winner?

(P₁) \rightarrow cell

(P₂)

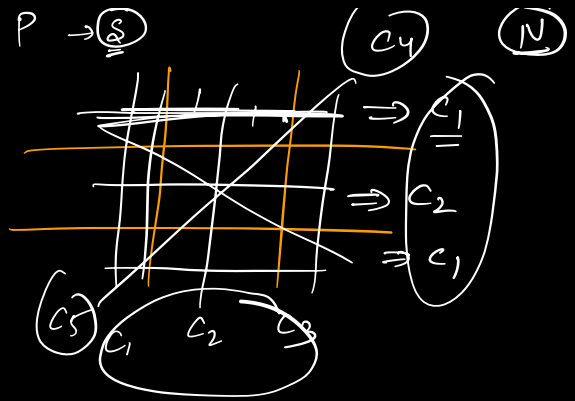
P₃

After every move, check if that player has won the game. If yes, stop the game, else continue

\rightarrow symbol

$P \rightarrow (x, y)$
 only checking
 for current player
 whether he won
 the game

B-F



checkWinner(Board, Player) {

 for (i = 0; i < N; i++)

 int cnt = 0

 for (int j = 0; j < N; j++) {

 if (board[i][j] == Player.symbol) {

 cnt++

 }

 if (cnt == N) {

 return true;

 }

Rows
 N^2

 for (int j = 0; j < N; j++)

 int cnt = 0

 for (int i = 0; i < N; i++) {

 if (board[i][j] == Player.symbol) {

 cnt++

 }

Colⁿ
 N^2

```

    }
    if (cnt == N) {
        return true;
    }
}

```

```

int cnt = 0;
for (int i = 0; i < N; i++) {
    if (board[i][i] == Player.symbol)
        cnt++;
}
if (cnt == N) return true;

```

L-R
Diag
N

```

int cnt = 0;
for (int i = 0; i < N; i++) {
    if (board[i][N-i] == Player.symbol)
        cnt++;
}
if (cnt == N) return true;

```

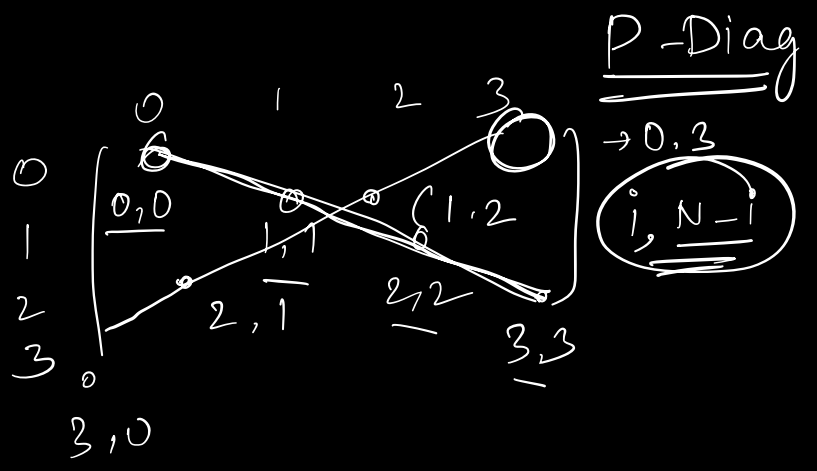
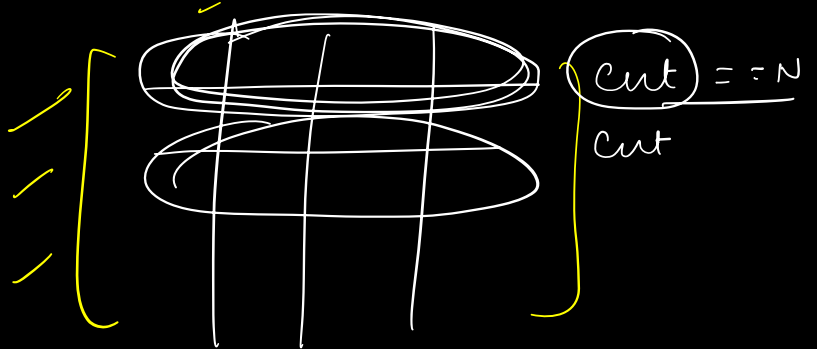
R-L
Diag
N

```

return false;
}

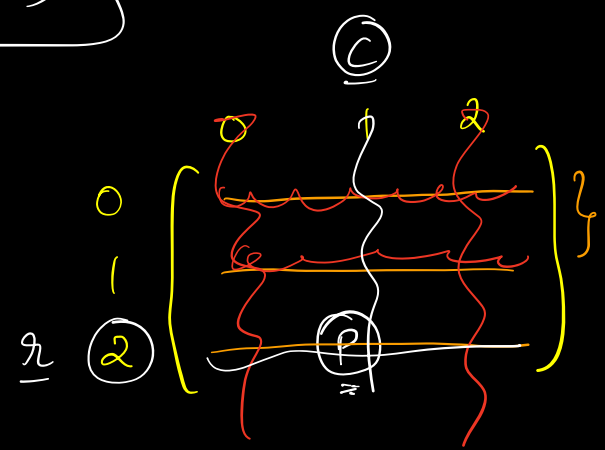
```

$N \times N$



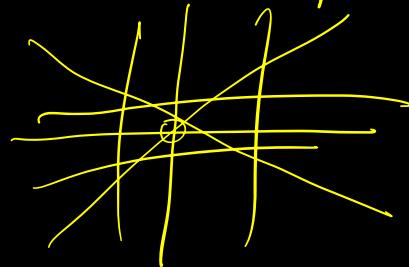
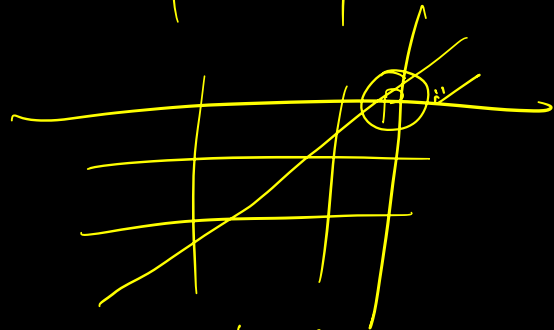
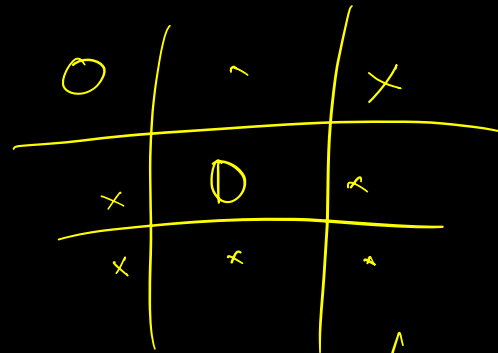
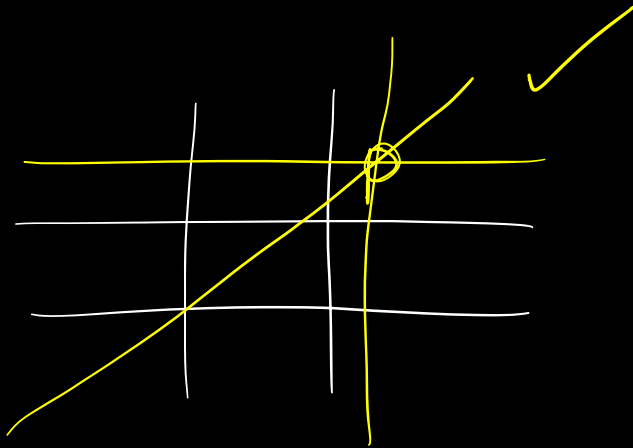
$$N^2 + N^2 + N + N$$

$O(N^2)$

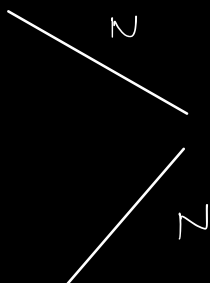


Obs

If player is making a move on (r, c) index then winning can only happen either on r th row or c th col or diag if that (r, c) is part of that diagonal.



N N
 (r, c)



```

int cnt = 0
for (int j = 0; j < N; j++) {
    if (board[r][j] == Player.symbol) {
        cnt++
    }
}
if (cnt == N) {
    return true;
}

```

N
 Check
 r^{th}
 row
 (counting
 no of player's
 symbol in
 r^{th} row)

```

int cnt = 0
for (int i = 0; i < N; i++) {
    if (board[i][c] == Player.symbol) {
        cnt++
    }
}
if (cnt == N) {
    return true;
}

```

N
 Check
 c^{th}
 col

```

int cnt = 0
for (int i = 0; i < N; i++) {
    if (board[i][i] == Player.symbol) {
        cnt++
    }
}

```

N

if (cnt == N) return true

int cnt = 0

for (int i = 0; i < N; i++) {

if (board[i][N-i] == Player

cnt++)

}

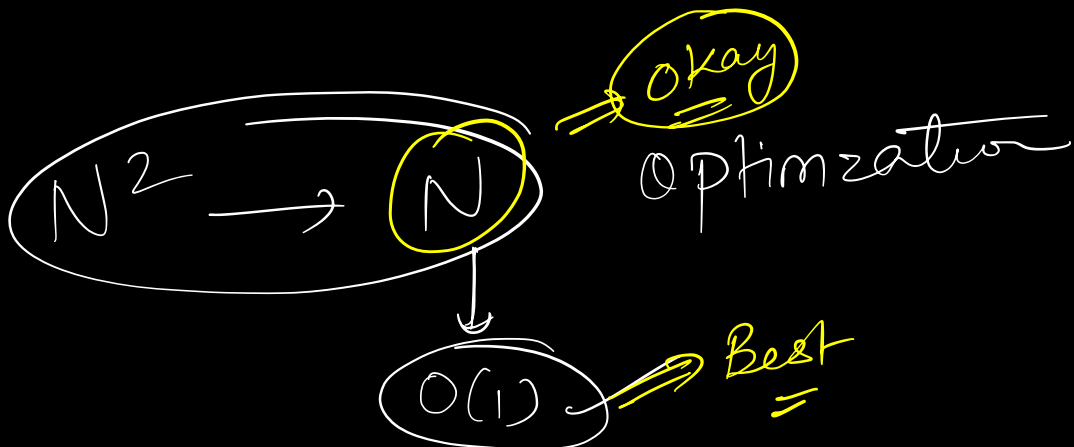
if (cnt == N) return true

return false

} N

Total TC - $N \times 4$

$O(N)$



$P(r, c)$

At the end of the day,

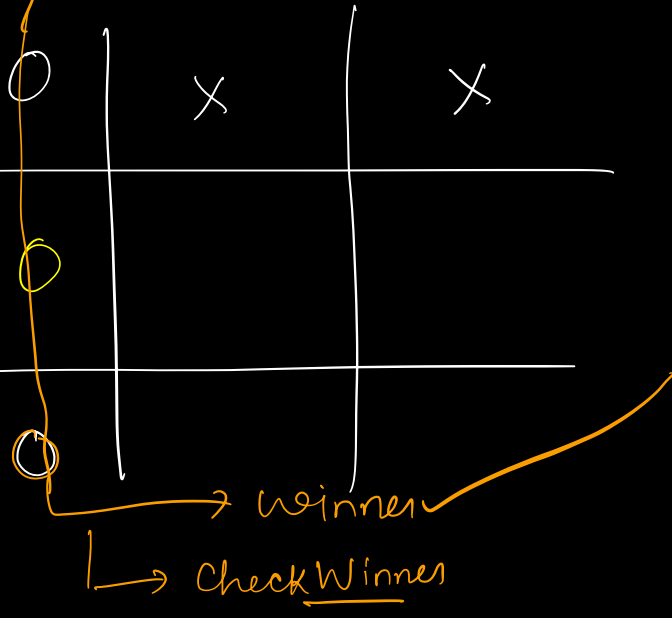
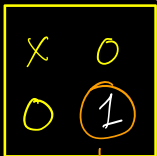
aren't we counting no. of Player's
symbol in r th row & c th col
& P & S diag

\Downarrow

What if we maintain
this info on the fly??

Hashmap \Leftarrow $\left\{ \begin{array}{l} P, \textcircled{r}, \Rightarrow \text{cnt} \\ P, \textcircled{c}, \Rightarrow \text{cnt} \\ P, \textcircled{L-R}, \Rightarrow \text{cnt} \\ P, \textcircled{R-L}, \Rightarrow \text{cnt} \end{array} \right\}$ $\textcircled{\underline{O(1)}}$

X	1
0	1



extra space
to
optimize
check
win

For each row, column, diagonals separate hashmap those have key as players(X,O) and value will be count of the cells in which that player symbol is present.

$SC \rightarrow \underline{N^2}$

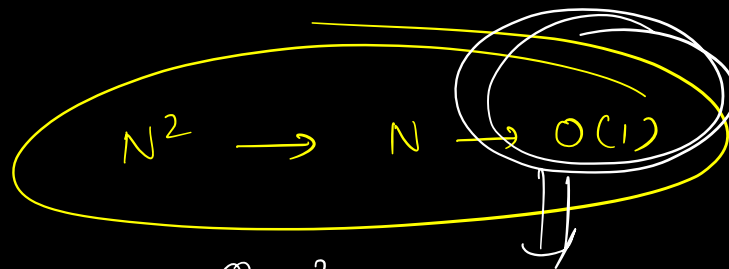
$(N) \quad (P, C) \Rightarrow \begin{matrix} \cancel{f} \text{ row,} \\ \cancel{f} \text{ col} \\ \cancel{f} \text{ 2} \end{matrix}$
 $N * (N + N + 2)$

$\times 1$

$\begin{matrix} 0 & 1 \\ \times & 2 \end{matrix}$	$\begin{matrix} \times & 1 \\ 0 & 1 \end{matrix}$	$\begin{matrix} \times & \\ 0 & 2 \end{matrix}$	$\begin{matrix} \times & 2 \\ 0 & \end{matrix}$	$\begin{matrix} \times & 1 \\ 0 & 2 \end{matrix}$ $P_{91, C}$
$\begin{matrix} \times & 2 \\ 0 & 1 \end{matrix}$	X	O	X	
$\begin{matrix} \times & 1 \\ 0 & 1 \end{matrix}$		O	X	N
$\begin{matrix} \times \\ 0 & 2 \end{matrix}$	O			

$N + N + 2$
 $2N + 2$

$O(1)$



Implement this
in our code

0	X	0

$\begin{matrix} 0 & 2 \\ \times & \\ \text{circle} & 1 \end{matrix}$

Dmplt will start from next class

Analysis

{ Last sol has
 N^2 as SC
are we ok?

$N = 1000$ \Rightarrow 1000×1000
 10^6 $\times 10B$

10^5 $O(1)$
 $4B$ $4B$ \Rightarrow $10B$

10^5 player
 \downarrow
 $10^5 \times 10^5$
 $10^{10} \times 10$
 10^{11} \Rightarrow $100GB$ \Rightarrow $10MB$ \Rightarrow $O(1)$
 \downarrow
 $100GB$ \times

$100GB$ of RAM \Rightarrow $1TB$ SSD
 $5TB$ HDD 1000

Cost

TC \rightarrow 10^5 \Rightarrow open
SC \rightarrow $O(1)$

10⁵

Board
10⁵ x 10⁵

Good idea

