

ACOL 202

29th April

$$f(n) = o(g(n))$$

$$f(n) = \Omega(g(n))$$

$$f(n) = \Theta(g(n))$$

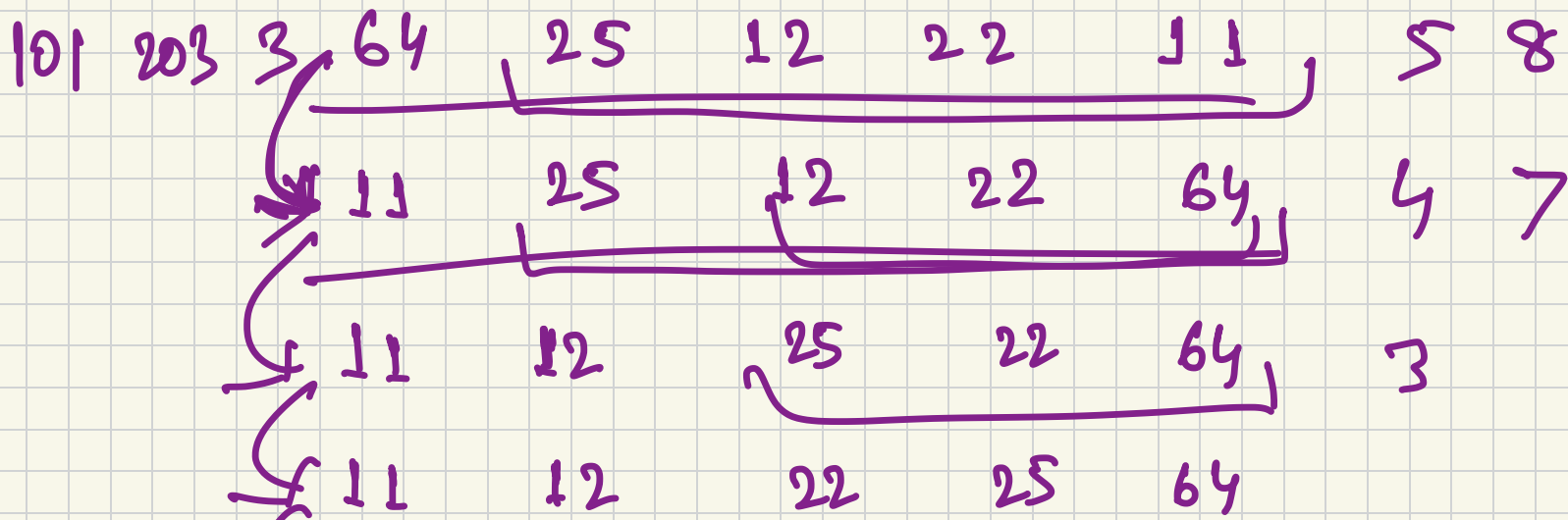
Selection Sort

64 25 12 22 11

find smallest, swap with the first element

find smallest among the remaining,
swap with the second element

and so on . . .



How many primitive steps

$$\boxed{n + (n-1) + (n-2) + \dots + 1}$$

$$\frac{n(n+1)}{2} \quad O(n^2)$$

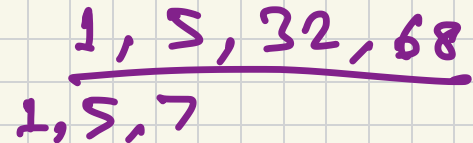
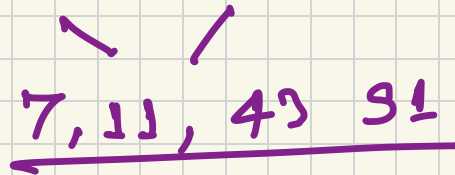
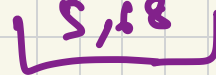
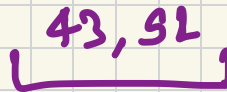
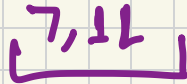
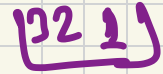
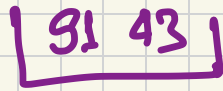
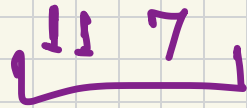
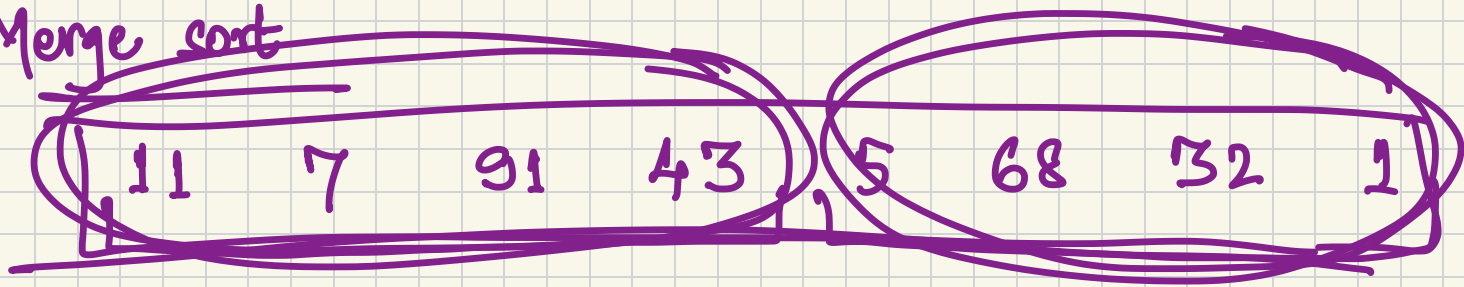
Find the smallest in the array / list
is $O(n)$ operation.

Linear search

[2 3 5 13 23 13 11] , 4

$O(n)$

Merge sort



$$\underline{T(n)} = \underline{2 T\left(\frac{n}{2}\right)} + c \cdot n$$

$$\boxed{T(n) = T(n-1) + c \cdot n}$$

$$T(n) = T\left(\frac{n}{2}\right) + c$$

$$T(n) = T(\cancel{n-1}) + c \cdot n$$

$$T(\cancel{n-1}) = T(\cancel{n-2}) + c \cdot (n-1)$$

$$T(\cancel{n-2}) = T(\cancel{n-3}) + c \cdot (n-2)$$

$$T(\cancel{n-3}) = T(\cancel{n-4}) + c \cdot (n-3)$$

\vdots

\vdots

\vdots

$T(\cancel{2})$

$=$

$T(1)$

$+$

$c \cdot 2$

$T(n)$

$=$

$T(1)$

$+$

$c \cdot 2$

$+$

$c \cdot 3$

$+$

\dots

$+$

$c \cdot n$

$$T(n) = O(n^2)$$

$=$

c

$+$

$c \cdot 2$

$+$

$c \cdot 3$

$+$

\dots

$+$

$c \cdot n$

$$= c \frac{n(n+1)}{2}$$

$$T(n) = T(n/2) + c$$

$$= T(n/4) + c + c$$

$$= T(n/8) + c + c + c$$

$$= T(n/16) + c + c + c + c$$

$$= T(n/32) + 5c$$

$$= T(n/2^6) + \underline{6c}$$

$$| T(1) + c \cdot \log(n)$$

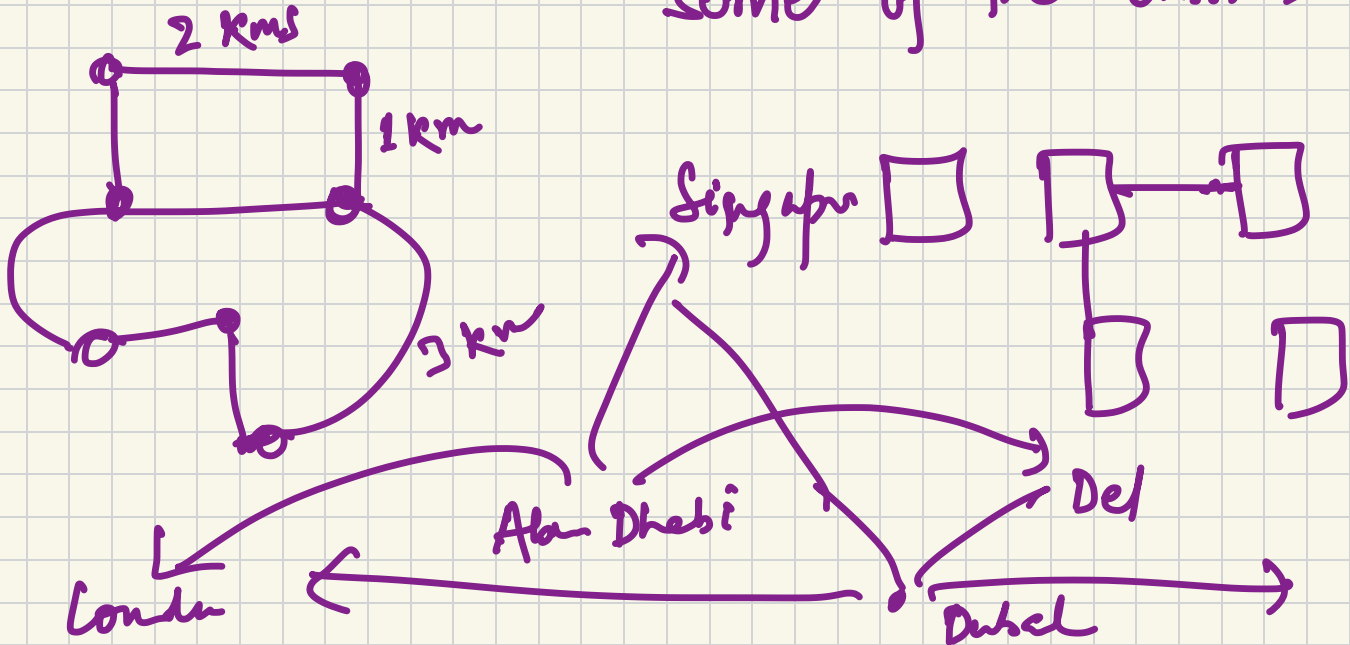
$$T(n) = 2T(n/2) + c \cdot n$$

Exercise Solve this recurrence.

Graphs are networks

Graph Theory

- a collection of entities
- pairwise relationship between some of the entities.



An undirected graph $G = (V, E)$
is a pair (V, E)

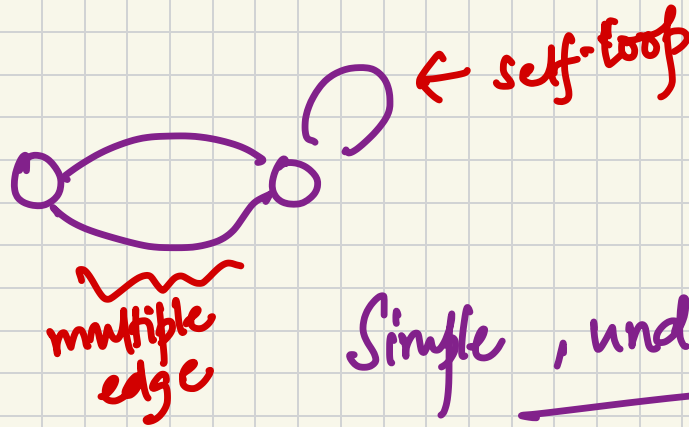
where V is the set of vertices
 E is the set of edges.

(u, v)

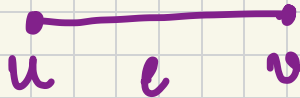


This is an ordered
pair if G
is a directed
graph.

A graph is simple if it has no multiple edges or self-loops.



Simple, undirected graphs



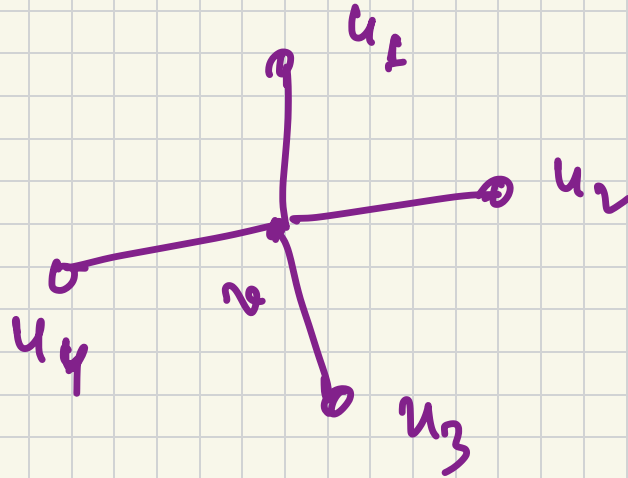
u and v are adjacent

u is a neighbor of v

u and v are endpoints of e

Neighborhood

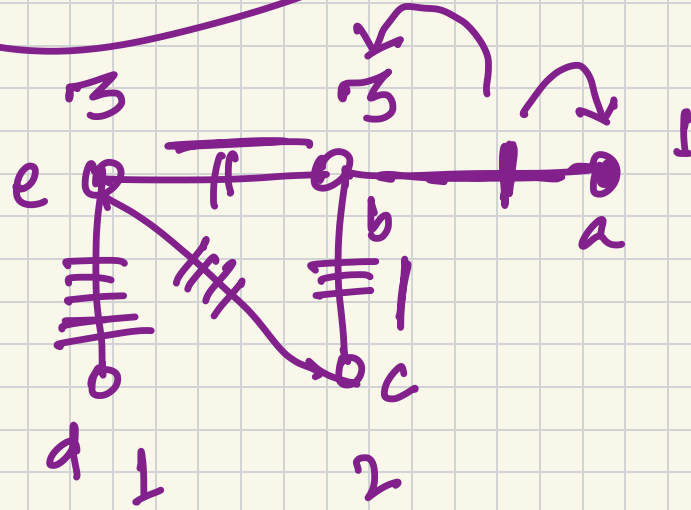
$$N(v) = \{u \mid (u,v) \in E\}$$



$\text{degree}(v)$
 $=$ size of the
neighborhood of
 v .

Claim

$$\sum_{v \in V} \text{degree}(v) = 2 |E|$$



$$\begin{aligned} \sum \text{degrees} &= 3 + 3 + 1 + 2 + 2 \\ &= 10 \end{aligned}$$

$$= 2 * 5$$

If n_{odd} is the number of nodes with an odd degree, then n_{odd} must be even.

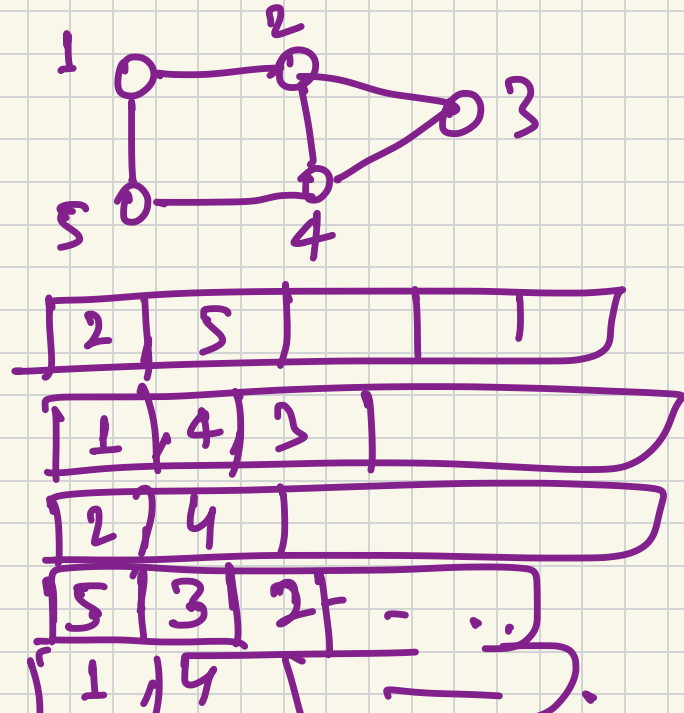
Representing Graphs

Adjacency Matrix

	1	2	3	4	5
1	0	1	0	0	1
2	1	0	1	1	0
3	0	1	0	1	0
4	0	1	1	0	1
5	1	0	0	1	0

Adjacency list

1
2
3
4
5



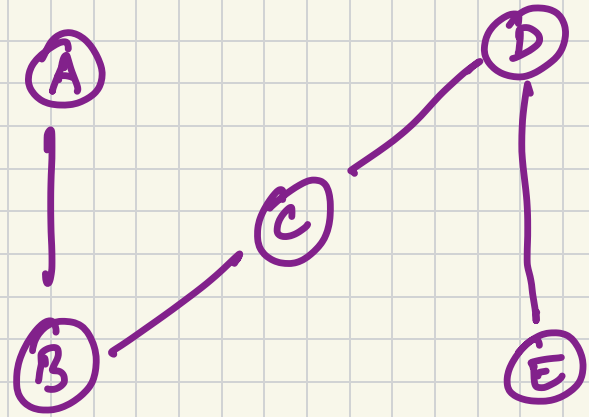
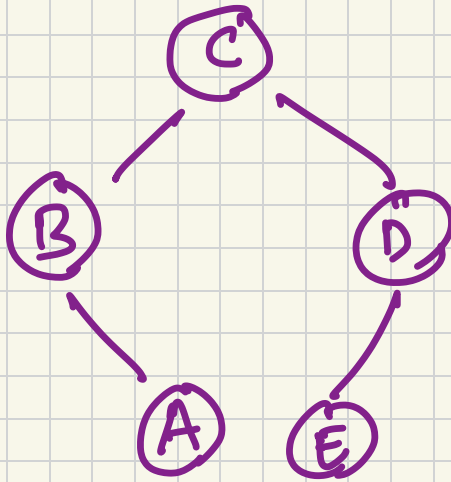
Graph Isomorphism

$G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$
are isomorphic if there exists a
bijection $f: V_1 \rightarrow V_2$

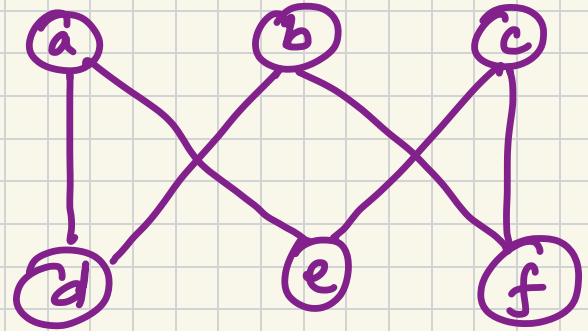
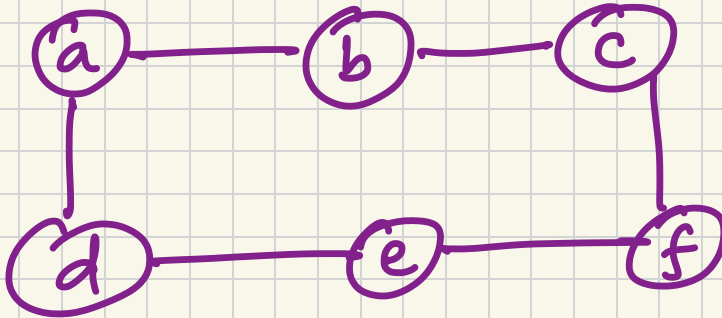
such that

$$\forall u, v \in V_1$$

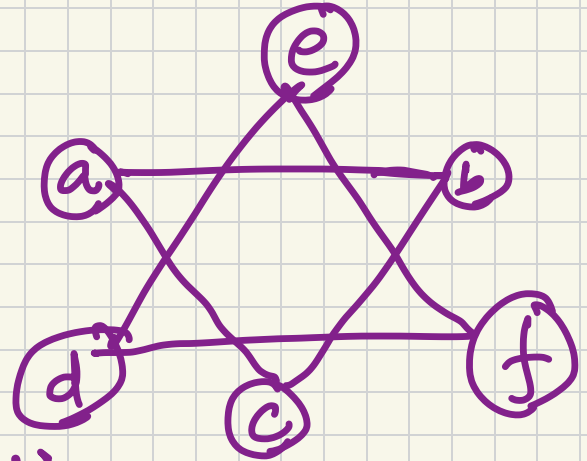
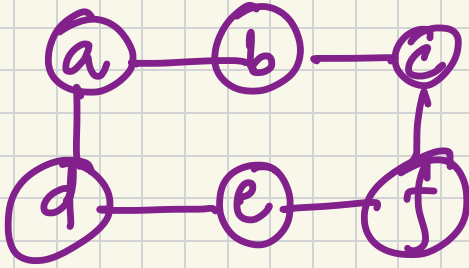
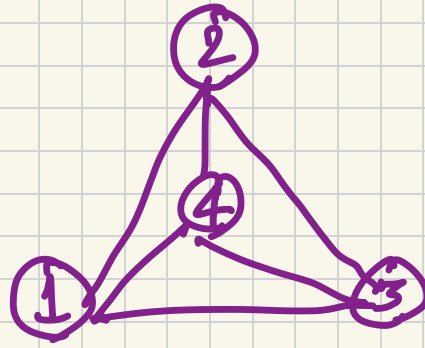
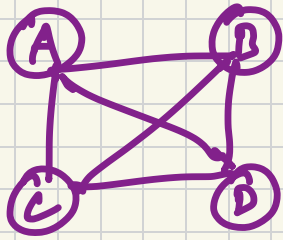
edge $(u, v) \in E_1$ iff
edge $(f(u), f(v)) \in E_2$.



isomorphic



isomorphic



Not isomorphic

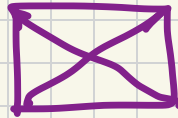
Subgraph of a graph $G = (V, E)$
is another graph $G' = (V', E')$
where $V' \subseteq V$
and $E' \subseteq E$
and for every $e \in E'$,
the endpoints of e
are in V' .

Special types of graphs

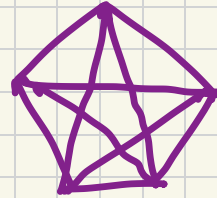
Complete graph (clique)



K_3



K_4



K_5

How many edges does a complete graph of n vertices have: $\binom{n}{2}$.

Bipartite graph

$$G = (V, E)$$

where V can be partitioned into

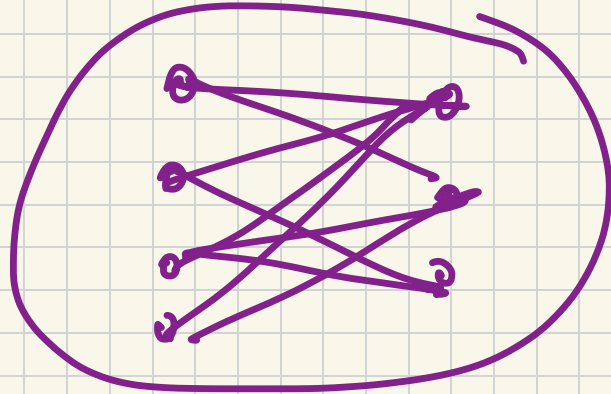
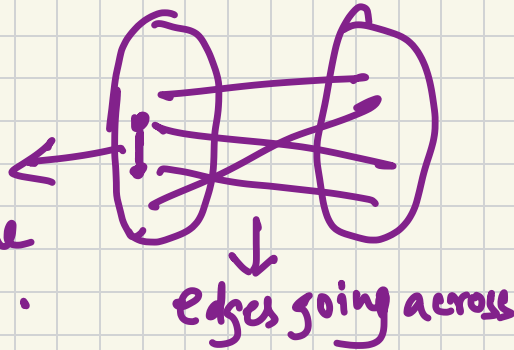
L and R such that for

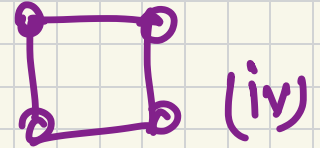
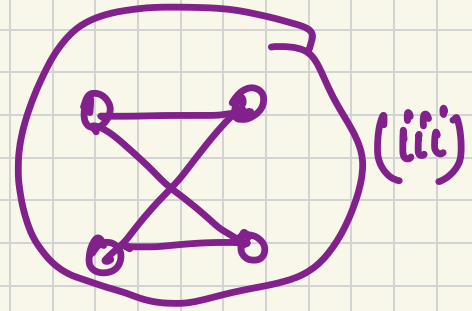
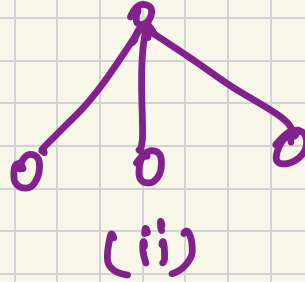
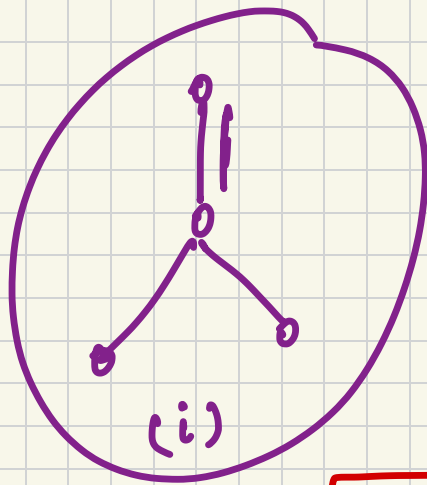
every edge $e \in E$

one endpoint of e is in L

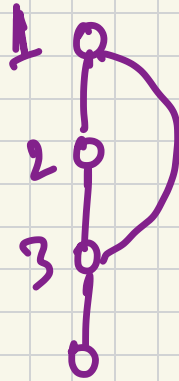
and the other endpoint is in R .

No edges
between two
vertices on the
same side.





i) and ii) are bipartite. They are isomorphic. So are (iii) and (iv).



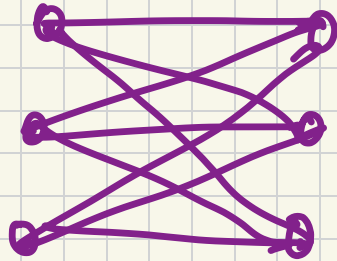
Not bipartite

Suppose, $1 \in L$.

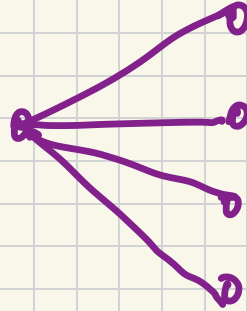
Then both 2 and 3 must be in R .

But that is not possible because $(2,3)$ is an edge!

Complete bipartite graphs



$K_{3,3}$



$K_{1,4}$