Correctness

- C1. For every v in Q' , $d[v]$ is the length of the shortest path between s and v .
- C2. For every v in Q , $d[v]$ is the length of the shortest path between s and v that only uses vertices in Q' (except the last vertex).

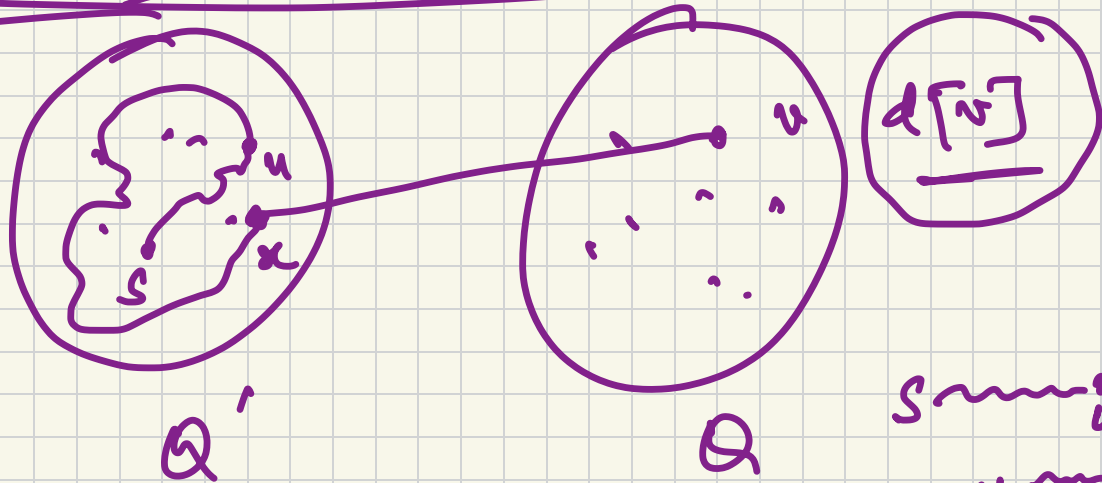
Both these claims are true in the beginning.

Suppose u has the smallest d -value in Q .

$$Q' \leftarrow Q' \cup \{u\}$$

$$Q \leftarrow Q \setminus \{u\}$$

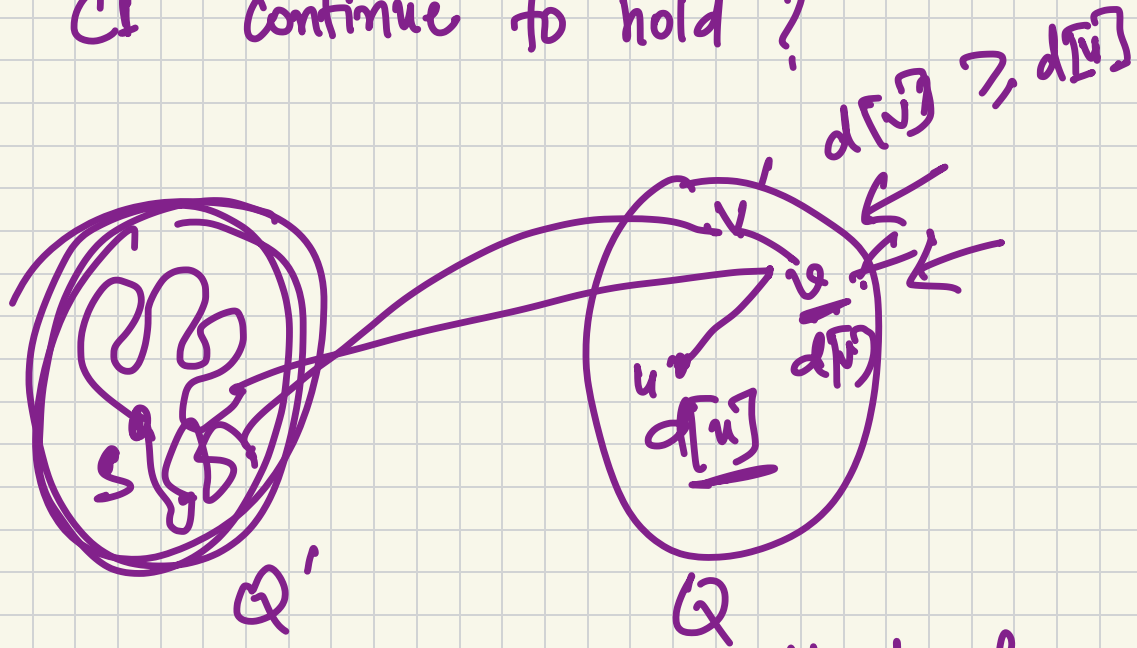
Edge-relaxation //



$$\begin{array}{l} s \rightsquigarrow x \rightsquigarrow v \\ s \rightsquigarrow u \rightsquigarrow \underline{x} \rightsquigarrow \underline{v} \end{array}$$

Clearly $C2$ continues to hold after this step (assuming that $C1$ was true earlier).

Why does $C1$ continue to hold?



Exercise You will complete the proof.

Graphs

Edges

Binary Relation

Given sets A and B , a binary relation $R: A \rightarrow B$ is a subset of $A \times B$.

$(a, b) \in R$ is denoted by $a R b$
 $a \sim_R b$

Every relation $R: A \rightarrow B$ can be represented as a bipartite graph.

Matrix representation

(seen example \rightarrow adjacency matrix)

Inverse relation

$$b R^{-1} a \text{ iff } a R b$$

Composition

$$R: B \rightarrow C \mid S: A \rightarrow B \mid R \circ S: A \rightarrow C$$

$$a (R \circ S) c \text{ iff } \exists b \in B \text{ such that}$$

$$a S b \text{ and } b R c.$$

Relations on one set

$$R \subseteq A \times A$$

set of cities

$a R a'$ if there is a direct flight from a to a' .

\mathbb{Z}

$$a R a' \text{ iff } a \equiv a' \pmod{7}$$

 \mathbb{N}

$$a R a' \text{ iff } a \leq a'$$

Reflexive

$$\forall a \in A$$

$$a R a$$

Symmetric

$$\forall x, y \in A$$

$$x R y \Rightarrow y R x$$

Transitive

$$\forall x, y, z \in A$$

$$x R y \text{ and } y R z$$

$$\Rightarrow x R z$$

 R, \leq

reflexive ✓
 Symmetric ✗
 transitive ✓

(antisymmetric)
 $x R y \text{ and } y R x \Rightarrow x = y$ ✓

$$\sim \text{ on } \mathbb{Z}$$

$x \sim y$ if x and y have the same parity (even or odd)

reflexive
symmetric
transitive

equivalence relation

$$a \sim b \quad \text{iff} \quad a \equiv b \pmod{5}$$

$\sim \sim \sim$ 0, 5, 10, 15, 20, . . . -
 $\sim \sim \sim$ 1, 6, 11, 16, 21,
 $\sim \sim \sim$ 2, 7, 12, 17, . . .
 $\sim \sim \sim$ 3, 8, 13, 18
 $\sim \sim \sim$ 4, 9, 14, 19 =

Partial Order

Irreflexive

$$\forall x \in X \quad (x, x) \notin R$$

xRx
is not true

Let A be a set.

[A relation on A is a partial order if it is reflexive, antisymmetric and transitive.

[strict partial order if it is irreflexive, antisymmetric and transitive.

\nsubseteq

$a \mid b$

$a < b$

partial order

strict partial order $a \leq b$ partial order

Comparability

Let \leq be a partial order on A .

$a, a' \in A$ are said to be comparable

if either $a \leq a'$ or $a' \leq a$ | otherwise
incomparable

Total order

Strict total order

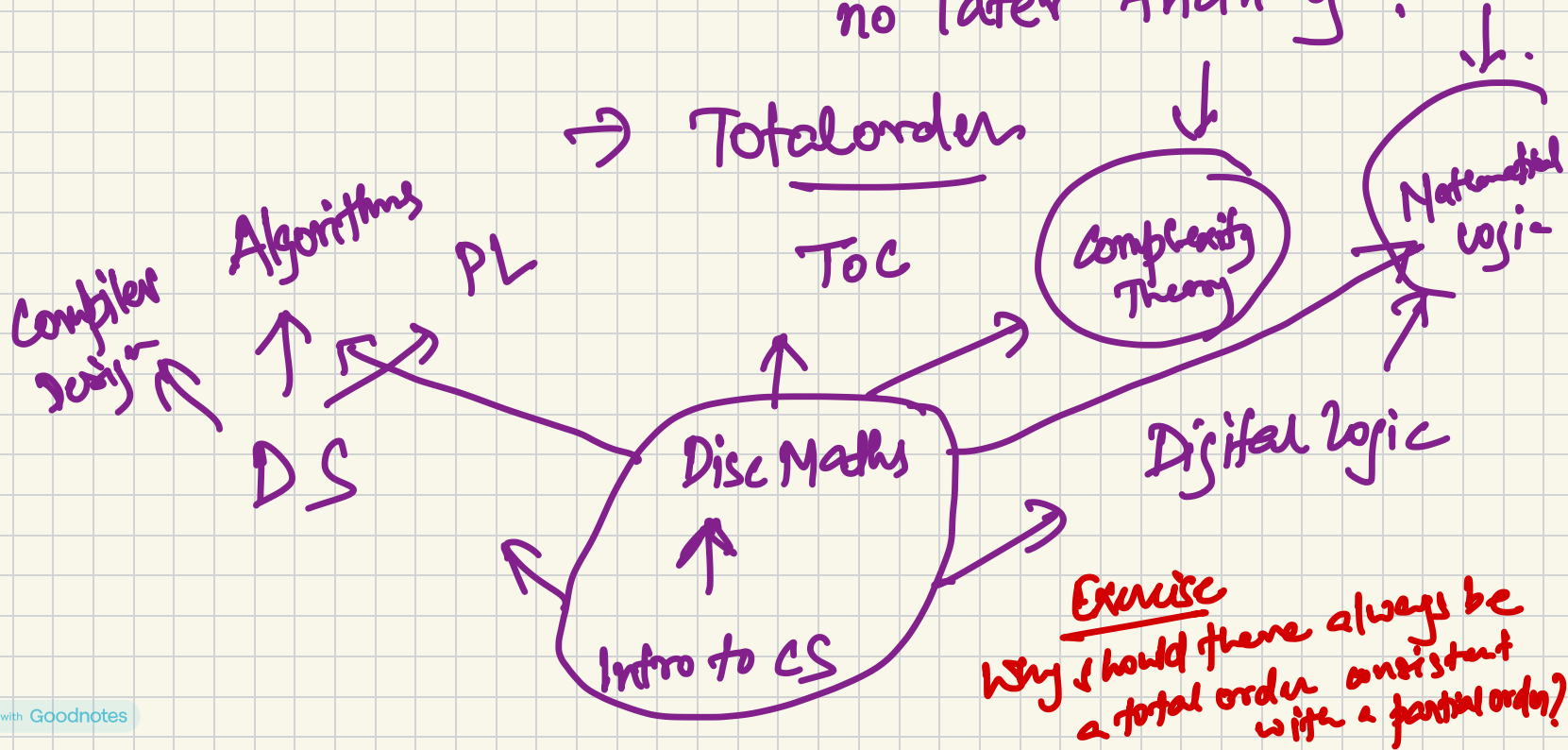
partial order

+ every pair of elements comparable

Strict partial order
+ every pair of distinct elements comparable.

$x R y$

if x comes alphabetically
no later than y .



The Well-ordering Principle

Every nonempty set of non-negative integers has a smallest element.

We have used it many times without explicitly talking about it.

For example: while proving that $\sqrt{2}$ is irrational, we assumed, for contradiction, that $\sqrt{2} = p/q$ in lowest terms.

How do we know that a fraction can be converted to one in lowest terms?

Suppose not. Let there be positive integers m and n such that m/n cannot be written in lowest terms.

Let C denote the set of numerators of all such fractions.

Clearly, $m \in C$. Thus, C is non-empty.

\therefore It must have a smallest element, say m_0 .

By defⁿ of C , there must be a positive integer n_0 such that $\frac{m_0}{n_0}$ cannot be written in lowest terms.

This means that m_0 and n_0 must have a common prime factor $p > 1$.

Consider $\frac{m_0/p}{n_0/p}$ which is equal to $\frac{m_0}{n_0}$.

If $\frac{m_0/p}{n_0/p}$ could be expressed in lowest terms then we could use the same expression for m_0/n_0 .

Therefore, it must be the case that $\frac{m_0/p}{n_0/p}$ also cannot be expressed in lowest terms. But then m_0/p (the numerator) must have been in C . But $m_0/p < m_0$. contradiction!