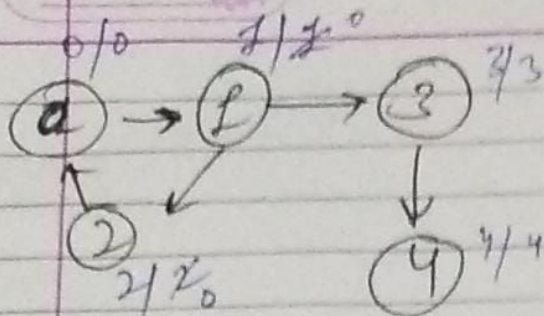


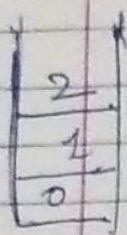
① Strongly connected components

Date _____
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(Tarjan's Algo)

(I)



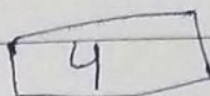
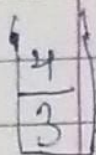
when (0/0) Point comes.
Pop from stack until
0 found.

[2, 1, 0] → SCC 1

(II)

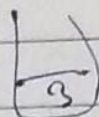
when (4/4) comes

Pop until 4 comes



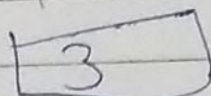
do
merge
Type

(IV)



when (3/3) comes

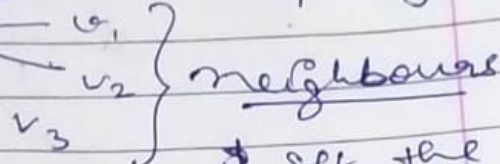
Pop until 3 comes



2) Largest Path in DAG

Shortest Path (Topological Sort)
(Problem Solving)

- 1) Get Topological order
- 2) Loop through the topological order (u)



if ($dis[u] \neq INF$ &
 $dis[u] + wt < dis[v]$)

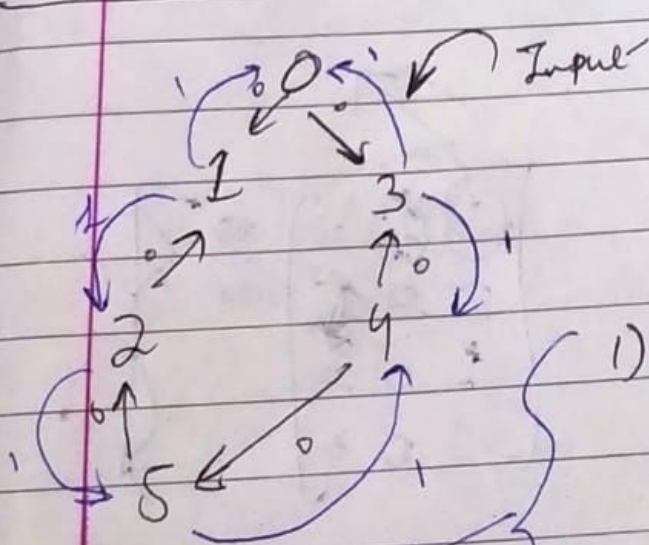
& see the
minimum
distance

2) Making wired connections Problem:

3) Minimum edges to Reverse to make Path from source to destination

1) Graph colouring (m-colouring)

Minimum edges to Reverse to Make a Path from source to Dest:



Input

output

min cost = 1

- 1) Provide a '0' weight to original edges
- 2) create Reverse edges with weight '1'

After this
Calc. shortest path

[0] BFS

③ Number of operation
to make network connected

Date _____

Page _____

$n=4$, $[0-1], [0-2], [1-2]$

1) Disjoint union set concept
is used.

→ $P(u) = \text{new DSU}(n)$;

So as you connect the component
Reduce them & don't waste
the cable.

0 — 1 \swarrow $P(0) \neq P(1)$

1

2

Comp --

$\Rightarrow 4-1=3$

$\Rightarrow 3-1=2$

\downarrow
 $P(0) \neq P(2)$

For connecting 3
Comp Require 2 edges
(Cables)

ans \Rightarrow Comp - 1

Graph Coloring : Chromatic no.
min color needed
to color a graph.

m-coloring Problem :

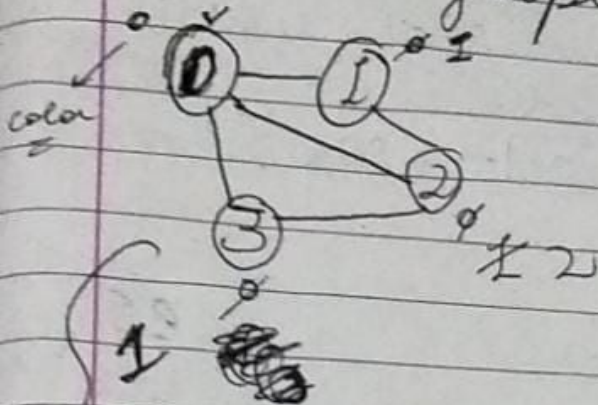
can we color a graph using M colors.

(4)

Approach \rightarrow BFS (not to Optimal)

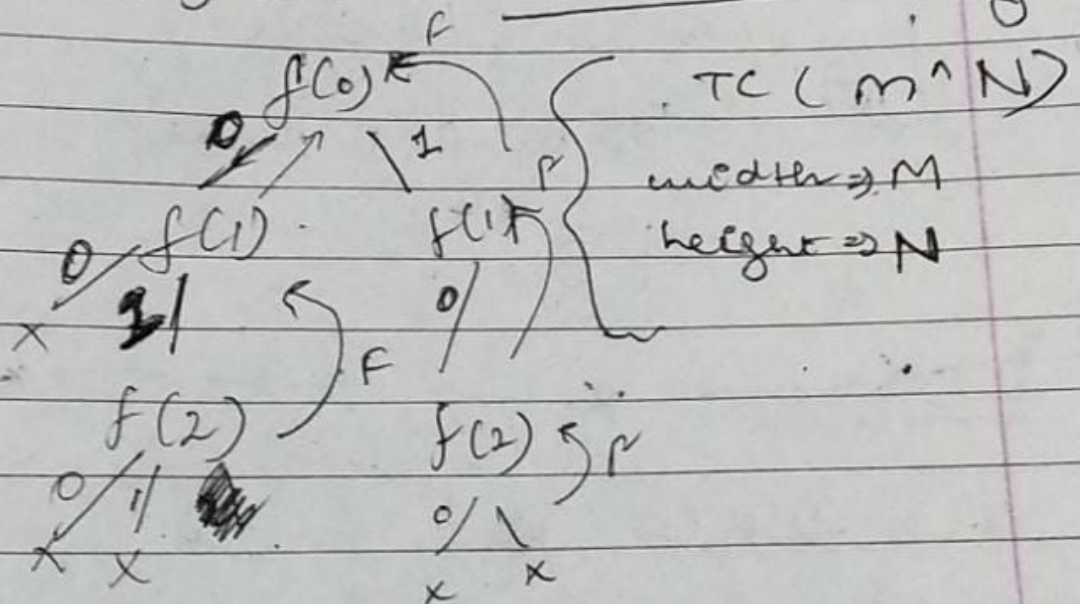
Q1: color a graph \rightarrow BFS
Q2: M coloring \rightarrow Backtracking

(I) color a graph (BFS)



- 1) Take node from queue & check all adjacent color if same increase their color
- 2) If adj nodes not visit add them to queue.

(II) M -coloring Prob. (Backtracking)



$M=2$
0, 1
only
2
poss