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TUTORIAL-7

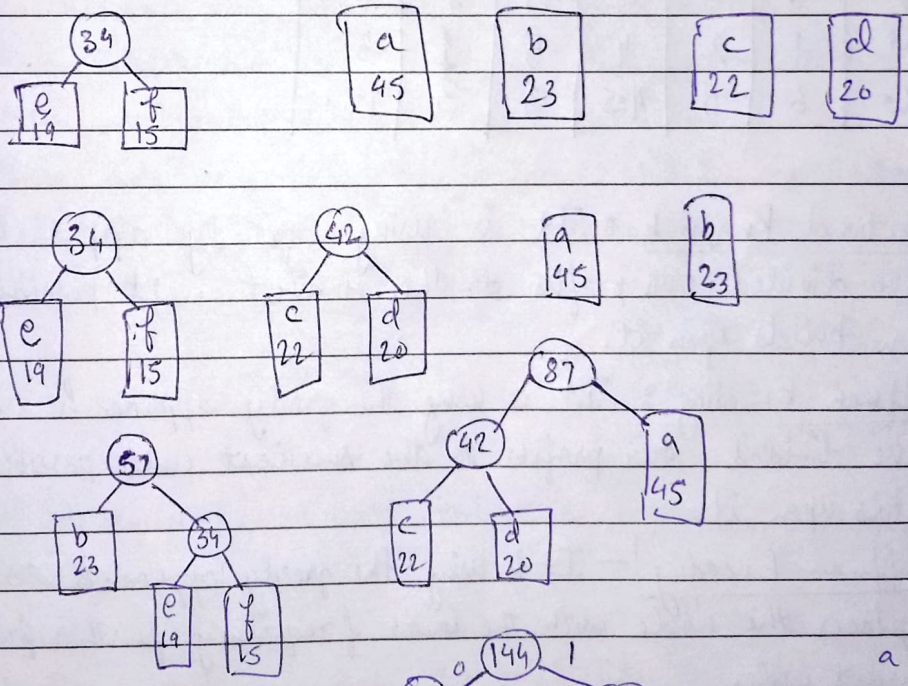
Sol1. It is an algorithmic paradigm that builds up a solution by adjoining smaller pieces together, always choosing the next piece that offers the most obvious & immediate benefit.

We should use greedy approach whenever a locally optimal solution is also globally optimal.

Sol2.

Name	TC	SC
Activity Selection	$O(n\log n) - O(n)$	$O(n)$
Job Sequencing	$O(n^2) - O(n\log n)$	$O(n)$
Fractional Knapsack	$O(n\log n) - O(n)$	$O(n)$
Huffman Encoding	$O(n\log n) - O(\log n)$	$O(n)$

Sol3. a=45 , d=20 , b=23 , e=19 , c=22, f=15



Total bits = (45×2)+(23×2) +(22×3)+(20×3)+ (19×3)+ (15×3) = 364 bits

a = 11
b = 00
c = 100
~~d = 101~~
e = 010
f = 011

Sol 4. A 2-tree is used to implement Huffman encoding algorithm. It is a binary tree where every node has either 2-child or no child.

- Applications of Huffman Encoding –
  - Data compression in long files without any loss.
  - To implement traffic routes with traffic magnitude.

Sol 5.

v	10	5	15	7	6	18	3
w	2	3	5	7	1	4	1
v/w	5	1.67	3	1	6	4.5	3

$$k = 15 - 1 - 2 - 4 - 5 - 1 - 2 = 0$$

$$\text{Profit} = 30 + 10 + 18 + 15 + 3 + 3.34$$
$$= 79.34$$

v	6	10	18	15	3	5
w	1	2	4	5	1	3
v/w	6	5	4.5	3	3	1.67

Sol 6. Fractional Knapsack : It is using a greedy approach as we have divided our profits to the smallest unit possible & then builds upon it.

Huffman Encoding : It is using the greedy approach as we have divided our profits to the smallest unit possible & then builds upon it.

Huffman Encoding – It is using the greedy approach as it always places the node with the lower frequency further from the parent node.

Sol7.

Start	1	2	0	6	9	10
End	3	5	7	8	11	12
Index	0	1	2	3	4	5

Jobs To Do ⇒ [0], [3], [4] or [5]

i.e. ⇒ Max = 4

Sol8.

	Profit	Deadline
a	20	2
b	15	2
c	10	1
d	5	3
e	1	3

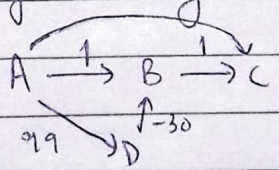
0	1	2
b	a	d

0 1 2 3

Profit = 20 + 15 + 5
= 40

Sol9. Times when not to use greedy algorithm :-
- When approach involves a lot of assumptions, such as "pick always the ..."
- When we need complex implementation.
- When we are making performance - critical application.

eg- Dijkstra's algo is very unoptimised for graphs with negative edges



Here, we cannot find distance of pair [A, C] → It gives 0, though it is −200.

Sol10. TC of job sequencing = $O(n^2)$, but we can improve it using priority queue by using (Max Heap).

Algorithm :-
1. Sort based on deadlines
2. Iterate the end & calculate the available slots b/w two consecutive deadlines. Put everything in Max Heap.
3. If slots available & there are jobs in MaxHeap, include Job ID with max profit & deadlines in the result.
4. Sort the array based on deadlines.

TC = $O(n \log n)$; SC = $O(n)$