

product
how many no of ways
both can occur

UNIT-2 PERMUTATION AND COMBINATION.

$$nC_r = \frac{n!}{(n-r)!r!} \quad nPr = \frac{n!}{(n-r)!}$$

sum rule
no of ways it occur
at same time

1. Suppose three persons enter into the car where there are 5 seats. In how many ways can they take their seats.

$$5P_3 = \frac{5!}{(5-3)!} = \frac{5 \times 4 \times 3 \times 2 \times 1}{2} = 5 \times 3 \times 4 \times 1 =$$

5 ways 4 ways 3 ways 2 ways 1 way

2. In how many ways can one select diff two books from diff subjects among 6 distinct computer, 3 Maths, 2 chemistry.

3. Suppose you have 6 new year cards sharing 4 friends

$$6P_4 = \frac{6!}{(6-4)!} = \frac{6!}{2!} = 6 \times 5 \times 4 \times 3 \times 1$$

4. In how many no of ways can 4 girls & 5 boys
are arranged in a row such that 4 girls are sitting together

$$6! \times 4!$$

girls $\underline{B_1 B_2 B_3 B_4 B_5}$

or we do a circle in group of 4, so $6! \times 4!$

and partitioning with pluriplies of 5 is in approaches

5. A student has to answer ten qns choosing atleast
4 from each from part A & B. If there are 6 in part A, 7 in
Part B in how many no of way a student choose

A	B
6	7
5	6
4	5

6. There are 18 cans of soup on a shelf. 3 of them
can contain tomato soup. Other contains something.

$$\frac{8!}{3! 1! 1! 1! 1! 1!} \quad 3 \rightarrow \text{are identical}$$

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Circulation

circular permutations in which objects are arranged in a circle.

To calc the no of ways in which n objects can be arranged in a circle be arbitrarily fix position of one object so the remaining $n-1$ objects can be arranged as if they were on a straight line (ie) the no arrangements in a circle is $(n-1)!$

NOTE : The no of circular permutation of n different objects taken all at a time is $\frac{n!}{2}$ [clockwise & anticlockwise order are treated as different]

No of circular permutation of n different objects taken all at a time $\frac{n!}{n} = \frac{n!}{n} = (n-1)!$

NOTE : The following situation leads to identical nature of clockwise & anti clock wise

1. formation of garland with flowers
2. formation of necklace with beads.
3. A person sitting around table with distinct neighbours.

1. Find the no of ways 6 people can be seated in a round table

$$(n-1)! = (6-1)! = 5! = 120.$$

2. Find the no of ways in which 5 people A, B, C, D, E seated in a round table such that A, B should sit together.

$$(4-1)! \times 2! = 3! 2! = 12.$$

3. In how many ways can 3 men & 3 women be seated in a round table such that no two men should sit together.

$$\text{men position } (3-1)! = 2!$$

$$\text{women position } 3! \text{ tot} = 3! \times 2! = 12$$

4. In how many ways can 5 people can be seated around circular table. two people insist to sit together

2. If two people refuse to sit each other.

$$1. \quad 3! 2! = 12$$

2. No of ways 2 people sit together = 12 ways

$$\text{No of ways 2 people refuse to sit} = (4-1)! - 12 \\ = 24 - 12$$

$$= 12.$$

4. Write all possible circular permutation using
1, 2, 3, 4

$$n(n-1)(n-2) \dots (n-(n-1)) = n!(n-1)$$

$$\frac{4P_3}{3} = 8$$

5. Find the no. of ways preparing a chain with
6 diff colored beads.

$$\frac{(n-1)!}{(n-2)} = \frac{(5-1)!}{(5-2)} =$$

6. Find the no. of ways of arranging chief minister
& 10 cabinet minister in a circular table so that
the chief minister should sit in one position.

7. n+1 pigeon occupies n pigeon holes at least.
One hole has

PIGEON HOLE PRINCIPLE

Generalization of pigeon hole principle.

If m pigeons occupies n holes then at least 1 hole
has more than $\lceil \frac{m-1}{n} \rceil + 1$ pigeons

$\lfloor x \rfloor \rightarrow$ floor function 8.5 \rightarrow 8.
only integer

1. Show that among 100 people atleast 9 of them were born in the same month.

no of people (pigeon) $m = 100$.

no of holes (no of month) $= 12$.

$$\left[\frac{100-1}{12} \right] + 1 = \left[\frac{99}{12} \right] + 1 = 8 + 1 = 9.$$

2. Show that if 7 colors are used to paint 50 bicycles atleast 8 bicycle have same color.

no of bicycles (pigeon) $m = 50$.

no of colors (holes) $n = 7$.

$$\left[\frac{50-1}{7} \right] + 1 = \left[\frac{49}{7} \right] + 1 = 7 + 1 = 8.$$

3. In a group of six or more person atleast 3 are mutual friends or atleast 3 are mutual strangers.

Consider the group has exactly 6 person A, B, C, D, E, F .

fix A and create a pigeon hole has consist of friends to A and consist of strangers to A.

m , no of pigeons (total people ext A) $= 5$

n , no of holes $= 2$.

$$\left[\frac{5-1}{2} \right] + 1 = \left[\frac{4}{2} \right] + 1 = 2 + 1 = 3$$

Case 1: If P (pigeon hole) departs out persons, then while any two of these 3 persons are finds then the hole 1 contains mutually 3 finds.

Case 2: any two pigeon hole 2 has three person's namely $\rightarrow D, E, F$ two of these 3 are strangers; hole - 2 contains \therefore hole - 2 contains 3 mutual strangers.
 \rightarrow none of them are mutually strangers

If $A \& B$ are finite subset of finite set U then

$$|A \cup B| = |A| + |B| - |A \cap B|$$

where modus of $A \rightarrow$ cardinality of set A

This principle can be extended by finite no of finite sets

$$|A_1 \cup A_2 \cup \dots \cup A_n| = \sum |A_i| + \sum_{i < j} |A_i \cap A_j| + \sum_{i < j < k} |A_i \cap A_j \cap A_k| + \dots + (-1)^{n+1} |A_1 \cap A_2 \cap \dots \cap A_n|.$$

1. Find the no. of integers between 1 to 250 both inclusive that are not divisible by 2, 3, 5, 7

Let A, B, C, D be the set of integers that lie btwn 1 & 250 and are divisible by 2, 3, 5 & 7 respectively

$$|A| = \left\lfloor \frac{250}{2} \right\rfloor = 125 \quad |A \cap B \cap C \cap D| = \left\lfloor \frac{250}{210} \right\rfloor = 1$$

$$|B| = \left\lfloor \frac{250}{3} \right\rfloor = 83 \quad |B \cap C \cap D| = \left\lfloor \frac{250}{105} \right\rfloor = 2$$

$$|C| = \left\lfloor \frac{250}{5} \right\rfloor = 50$$

$$|B \cap C| = \left\lfloor \frac{250}{15} \right\rfloor = 16$$

$$|D| = \left\lfloor \frac{250}{7} \right\rfloor = 35 \quad |C \cap D| = \left\lfloor \frac{250}{35} \right\rfloor = 7$$

$$|A \cap B| = \left\lfloor \frac{250}{6} \right\rfloor = 41 \quad |A \cap B \cap C| = \left\lfloor \frac{250}{30} \right\rfloor = 8$$

$$|A \cap C| = \left\lfloor \frac{250}{10} \right\rfloor = 25 \quad |A \cap B \cap D| = \left\lfloor \frac{250}{42} \right\rfloor = 5$$

$$|A \cap D| = \left\lfloor \frac{250}{14} \right\rfloor = 17 \quad |A \cap C \cap D| = \left\lfloor \frac{250}{70} \right\rfloor = 3$$

$$|B \cap D| = \left\lfloor \frac{250}{21} \right\rfloor = 11$$

By the principle of Inclusion exclusion

The no of integers btwn 1 & 250 at least one

of the integers.

$$= 125 + 83 + 50 + 35 - 41 - 25 - 17 - 16 - 7 - 11 + 8 + 5 + 3$$

$$+ 2 - 1$$

$$= 250 - 193$$

$$= 57$$

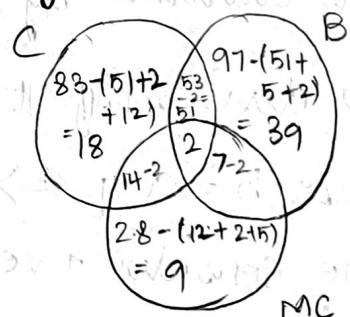
2. In a Survey of 150 clg students reveals that 83 own car, 97 own bikes, 28 own motor cycle, 53 own a car and a bike, 14 own car & motor cycle, 7 bike & motor 2 own all three. How many own a bike & nothing else.

How many do not own any of the three

$$n|C| = 83$$

$$n|B| = 97$$

$$n|M| = 150 - |C \cup B \cup MC|$$



$a = \left\lfloor \frac{a}{d} \right\rfloor \cdot d$ $a = \left\lfloor \frac{a}{d} \right\rfloor \cdot d$

Number theory is the part of mathematics to investigate more. Specifically the properties of positive integers.

Basic properties of divisibility theory and congruences

Congruences

DIVISIBILITY THEORY

Let a & b are two integers with $a \neq 0$.
a divides b or b is divisible by a or a is a divisor or factor of b or b is a multiple of a . ($a|b$)

$a|b$ if $b = aq$ for some $q \in \mathbb{Z}$.

The statement a divides b is written and its negation is denoted by $a \nmid b$.

ex, 66 is divisible by 11 $\Rightarrow 11|66$

$$66 = 11 \times 6$$

ex, -7 divides 42 $\Rightarrow 7|-42$

$$42 = (-7) \times 6$$

0 is divisible by every integers

If a divides b then $-a$ is also divides b

because $b = aq \Rightarrow b = -a \times -q$. it is enough if we consider +ve divisor of integers

Theorem 1:

Let $a, b, c \in \mathbb{Z}$ a set of integers

- i) if $a|b$ & $a|c$ then $a|b+c$.
- ii) if $a|b$ & $b|c$ then $a|c$.
- iii) if $a|b$ & $a|c$ then $a|m(b+c)$.
- iv) if $a|b$ then $a|m b$ for $m \in \mathbb{Z}$.
- v) Since $a|b, a|c$ it follows from the definition

$b = ma$, $c = na$ where m, n are integers

$$b+c = (m+n)a.$$

$$a|b+c.$$

ii) $a|b$ & $b|c$

$$b = ma \text{ & } c = mb \quad m, n \in \mathbb{Z}.$$

$$c = nb$$

$$= n(ma)$$

$$= (nm)a$$

$$\therefore a|c.$$

iii) Since $a|b \Rightarrow b = ma$

$$mb = m(na)$$

$$mb = (mn)a$$

$$a|m b.$$

$$a|b \Rightarrow b = ma$$

$$a|c \Rightarrow c = na$$

$$a|m b \text{ & } a|n c \Rightarrow a|m b + n c$$

$$v) a|b \Rightarrow a|m b$$

$$a|c \Rightarrow a|n c$$

$$\therefore a|m b + n c.$$

A positive integer is $b > 1$ is called prime if the only divisor of b is $1 \& b$.

A positive integer is $b > 1$ is called composite if the divisor of b is more than $1 \& b$.

NOTE :-

i) A number that is not a prime is divisible by prime.

ii) A number that is prime

FUNDAMENTAL THEOREM OF ARITHMETIC.

Every positive integer $n > 1$ can be written uniquely as that product of prime numbers where the prime factors are written in order of increasing size.

$n=2$ since 2 is a prime number.

if n is a prime

if n is composite

Let us assume that the theorem holds good for +ve integer $< n$

$$n = ab \quad \because a, b < n$$

each of $a \& b$ can be expressed as product of primes

1. If $n > 1$ the composite integer and p is a prime factor of n then $p \leq \sqrt{n}$ or prove that a positive integer n is a prime number if no prime $p \leq \sqrt{n}$ divides n

Prove n is composite

PROOF :-

$n = ab$ for some integer $a \neq b$ ($a, b < n$)

without loss of generality $a \leq b$

let $a \geq \sqrt{n}$.

then $b \geq a \geq \sqrt{n}$. implies $ab \geq \sqrt{n} \cdot \sqrt{n} \geq n$.

which is a contradiction

so one of the factors, a should not be greater than or equal to \sqrt{n} .

Method of repetition $a \leq \sqrt{n}$

This divisor is either prime or by the fundamental theorem of arithmetic

n has a divisor $\leq \sqrt{n}$.

ze

NOTE: To test the given integer n is prime it is enough to see that it is not divisible by any prime $\leq \sqrt{n}$.

1. Find the prime factorization of 7007.

$$\frac{7007}{7} = \frac{1001}{7} = \frac{143}{13} = \frac{13}{13}$$

Prime factors of 7007 = $7^2 \times 11 \times 13$.

DIVISION ALGORITHM

When $a \in b$ are any 2 integers $b > 0$, there exist a unique integer q such that $a = bq + r$ when $0 \leq r < b$.

Proof:

i. what are the quotient and remainder when 77 is divided by 3 .

$$77 = 7(4)(3) + 2$$

Let $a \in b$ be any two integers we say that b is a divisor of a written as $b | a$ if $a = qb$ for integer q .

Let $a \in b$ be any two integers then an integer d is called a common divisor of $a \in b$ by $d | a \& d | b$.

GREATEST COMMON DIVISOR (GCD)

Let a, b are non zero integers then an integer d is called Greatest Common divisor if the following condition satisfies

i) $d | a \& d | b$.

ii) For any integer n , if $n | a \& n | b$ for $n | d$.

If $a \in b$ are set to be $\{5, 15\}$

If GCD of $a_i, b_j = 1$ then $1 \leq i < j \leq n$ are pairwise

EUCLID'S ALGORITHM

Euclid's algorithm

If γ_1 is the remainder when a is divided by b , γ_2 is the remainder when b is divided by γ_1 , γ_3 is remainder when γ_2 is divided by γ_1 .
 $\gamma_{k+1} = 0$ then the last non-zero remainder γ_k is the GCD of a, b .

Finding GCD using prime factorization

If the prime factorization of $a \& b$ are $a = P_1^{a_1}, P_2^{a_2}, \dots, P_n^{a_n}$
 $b = P_1^{b_1}, P_2^{b_2}, \dots, P_n^{b_n}$ Each exponent is a non-negative integer where all the prime occurring in prime factorization either a or b included with both factorization with 0 exponent if necessary.

$$\text{gcd}(a, b) = P_1^{\min(a_1, b_1)}, P_2^{\min(a_2, b_2)}, \dots, P_n^{\min(a_n, b_n)}$$

1. $\text{gcd}(120, 500)$ using Euclid's algorithm, prime factorization.

$$120 = 2^3 \times 3^1 \times 5^1$$

$$500 = 2^2 \times 5^3 \times 3^0$$

$$\begin{aligned}\text{gcd}(120, 500) &= 2^{\min(2, 2)} \times 3^{\min(0, 1)} \times 5^{\min(1, 3)} \\ &= 2^2 \times 3^0 \times 5^1\end{aligned}$$

Euclid algorithm,

$$500 \div 120 = 20 \quad \text{remainder}$$

$$120 \div 20 = 0$$

20 is the gcd of $(a, b) \Rightarrow (500, 120)$

2. Find $(1575, 231)$ by using Euclid algorithm.

$$a = 1575 \div 231 =$$

$$1575 = (231)q + 189$$

$$231 = (189)q + 42$$

$$189 = (42)q + 21$$

$$42 = 21(q) + 0$$

GCD of $(1575, 231)$ is 21

3. GCD of (a, b) can be expressed as integral of linear combination of a, b such that $\text{GCD}(a, b) = ma + nb$

3. Expressing gcd of $(252, 198) = 18$

$$252 = 198(1) + 54$$

$$198 = 3 \times 54 + 36$$

$$54 = 36 \times 1 + 18$$

$$36 = 18 \times 2 + 0$$

18 is the gcd of $(252, 198)$.

$$ma + nb = \text{GCD}(a, b)$$

$$(3 \times 1, 0 \times 1) \leftarrow (1, 0)$$

(whilst remainders) and

If a, b, c are positive integers such that

$\gcd(a, b) = 1$ and $a \mid bc$ then $a \mid c$.

Proof:

Let $a \& b$ are relatively prime integers $\gcd(a, b) = 1$,
there exist an integer such that $ma + nb = 1$ —①

$$mac + nbc = c$$

$$mac + nba = c$$

$$(mc + nb)a = c \quad bc = ka \quad k \in \mathbb{Z}$$

$$a \mid c.$$

3. $\gcd(ka, kb) = k\gcd(a, b)$

A. If $\gcd(a, b) = d$, then $\gcd\left(\frac{a}{d}, \frac{b}{d}\right) = 1$

5. $\gcd(a, b) = 1$ then for any integer c ,

$$\gcd(ac, b) = \gcd(c, b)$$

6. If each of a_1, a_2, \dots, a_n is a coprime of b

is $\gcd(a, b) = 1$ then the product $a_1, a_2, a_3, \dots, a_n$
is also coprime to b is $\gcd(a_1, a_2, \dots, a_n, b) = 1$.

7. If p is a prime & $p \mid a_1, a_2, \dots, a_n$ where each a_i
is an integer then $p \mid a_i$ for some i .

LCM (Least Common Multiple)

If a & b are +ve integers then the smallest +ve integer that is divisible by both a & b is called least common multiple of a & b . & its denoted by $\text{LCM}(a, b)$

NOTE! even if either both a & b are -ve $\text{LCM}(a, b)$ is always positive.

Alternate definition of $\text{LCM}(a, b)$

If the prime factorization of a & b are

$$a = P_1^{a_1} P_2^{a_2} P_3^{a_3} \dots P_n^{a_n}$$

$$b = P_1^{b_1} P_2^{b_2} P_3^{b_3} \dots P_n^{b_n}$$

Condition stated in the alternate definition:

$$\text{LCM}(a, b) = P_1^{\max(a_1, b_1)} P_2^{\max(a_2, b_2)} P_3^{\max(a_3, b_3)} \dots P_n^{\max(a_n, b_n)}$$

for example $\text{LCM}(24, 30)$

$$24 = 2^3 \times 3^1 \times 5^0$$

$$30 = 2^1 \times 3^1 \times 5^1$$

If a & b are two +ve integers then

$$\therefore \text{GCD}(a, b) \cdot \text{LCM}(a, b) = ab$$

Proof

Use prime factorization find LCM & GCD of (231, 1575) and also verify $\text{GCD}(m, n) \times \text{LCM}(m, n) = mn$.

$$231 = 3 \times 11 \times 7 \times 5^0$$

$$1575 = 3^2 \times 5^2 \times 7 \times 11^0$$

$$\text{GCD}(231, 1575) = 3 \times 11^0 \times 7 \times 5^0 = 21$$

$$\text{LCM}(231, 1575) = 3^2 \times 5^2 \times 11 \times 7 = 17325$$

$$\text{GCD}(231, 1575) \times \text{LCM}(231, 1575) = 363,825 \quad \text{--- (1)}$$

$$231 \times 1575 = 363,825 \quad \text{--- (2)}$$

3	1575	31
5	535	231
3	105	77
5	35	01
	7	

2. Find the integers m & n such that $512m + 320n = 64$

$$m + n = \text{GCD}(a, b)$$

$$\text{GCD}(512, 320) = 64$$

$$512 = 320 + 192$$

$$320 = 192 + 128$$

$$192 = 128 + 64$$

$$128 = 64 + 0$$

$$\text{GCD}(512, 320) = 64$$

$$64 = 192 - 128$$

$$= 192 - (320 - 1 \times 192)$$

$$= 2 \times 192 - 320$$

$$= 2(512 - 320) - 320$$

$$= 2 \times 512 - 3 \times 320$$

$$m = 2 \quad n = -3$$

UNIT 3 - MATHEMATICAL LOGIC.

ATOMIC PREPOSITION. (primary / primitive / simple) ^{declarative Sentence}

Preposition that does not contain any of the logical operators or connectives are called atomic preposition.

MOLECULAR OR COMPOUND PREPOSITION, OR COMPOSITE STATEMENT

Statement constructed by combining one or more atomic statements using connectives are called compound preposition

Truth value of preposition.

If a preposition is true its truth value is true denoted by T or 1.

If a preposition is false its truth value is false denoted by F or 0.

LOGICAL OPERATORS OR CONNECTIVES

Used to combine two or more proposition
five logical operators or connectives

	Name	Symbol
NOT	Negation	\sim or \neg

AND	Conjunction	\wedge
-----	-------------	----------

OR	Disjunction	\vee
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If then	Conditional	\rightarrow
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If and only if	Biconditional	\leftrightarrow
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NEGATION

Let P be a preposition that simply says that

If P is true then $\neg P$ is false

$\neg P \Rightarrow$ "negation of P ".

If P is false then $\neg P$ is true.

"It is not the case that P ".

P	$\neg P$
T	F
F	T

p	q	$p \wedge q$	p	q	$p \wedge q$
T	T	T	T	T	T
T	F	F	T	F	F
F	T	F	F	T	F
F	F	F	F	F	F

CONJUNCTION

Let $p \wedge q$ be two prepositions is the preposition that is true when both $p \wedge q$ are true and false otherwise.

ex; Today is Sunday $\rightarrow p$

p	q	\wedge
F	F	F
F	T	F
T	F	F
T	T	T

$p \wedge q$ is Today is Sunday and it is raining today

DISJUNCTION

Let p, q be any preposition The preposition $p \vee q$ is false only when both are false

Otherwise true.

p	q	\vee
F	F	F
F	T	T
T	F	T
T	T	T

CONDITIONAL STATEMENT

NOTATION

If P & q are any two statement then the Statement $P \rightarrow q$ has "If P , then q " is called Conditional statement. It is false only when P is true and q is false otherwise true.

P	q	$P \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

P	q	$P \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

NOTATION

BI CONDITIONAL

If P & q are any two statement then bi conditional statement is $P \leftrightarrow q$ "if P is true whenever both P & q has same truth value".

P	q	$P \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

P	q	$P \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

NOTATION

Conditional
Alternate way to express $P \rightarrow q$, $\rightarrow P$ implies q

$\rightarrow P$ is sufficient for q .

$\rightarrow P$ only if q

$\rightarrow q$ necessary of P

$\rightarrow q$ if P

$\rightarrow q$ whenever P

We shall avoid the word implies ~~implies~~

Variation in Conditional statement.

→ If P conditional q is Conditional preposition

the q conditional p is called converse of $P \rightarrow q$.

2. $\neg P \rightarrow \neg q$ is called inverse of $P \rightarrow q$,

3. $\neg q \rightarrow \neg P$ is called contrapositive of $P \rightarrow q$

for example,

If it rains today then I will stay at home.

P: It rains today q: I will stay at home.

Converse : If I will stay at home then it rains today

Inverse : If I does not rains today then it will not stay at home.

Contrapositive : If It will not stay at home then it does not rain today.

Different ways to represent biconditional

1. P is necessary & sufficient for q & vice versa.

2. If P then q and conversely

3. P if and only if q

Order of Precedence for logical operators.

Name	Operator	Precedence.	Conditional or Implication
\neg	Negation	1	
\wedge	Conjunction	2	
\vee	disjunction	3	
\rightarrow	Conditional	4	
\leftrightarrow	Biconditional	5	

$P \vee \neg P$
 $T \vee F \quad T$

Tautology

An expression involving logical variables that is

always true for all possible truth values of
variables in it.

example.

P		$\neg P$		$P \vee \neg P$	
T	F	F	T	T	F

Contradiction

An expression involving logical variables that is

always false for all possible truth values of
variables in it.

P		$\neg P$		$P \wedge \neg P$	
T	F	F	T	F	F

Contingency

A Compound proposition that is neither tautology
and nor contradiction is called contingency.

P		q		$P \vee q$	
T	F	T	F	T	F

P		q		$P \wedge q$	
T	F	T	F	T	F

1. Write a symbolic form, you can access internet from college only if you are a computer science major or a fresher.

P: You can access internet from college

q: You are Computer Science major or a fresher.

P only if q = $P \rightarrow q$.

$$P \rightarrow (q \vee \neg \bar{q})$$

2. It is necessary to have a valid password to login to the server.

P: to login to the server

q: valid password.

$$P \rightarrow q$$

3. Home team wins whenever it is raining.

P: Raining

q: Home team wins whenever P.

$$P \rightarrow q$$

4. for you to get an 'O' grade in this course, it is necessary & sufficient that you have to learn discrete maths.

P: for you to get an 'O' grade

$$P \leftrightarrow q$$

Construct the truth table for the following.

1. $(P \vee q) \vee \neg P$.

P	q	$\neg P$	$P \vee q$	$(P \vee q) \vee \neg P$
T	T	F	T	T
T	F	F	T	T
F	T	T	T	T
F	F	T	F	T

2. truth table for $p \vee q \rightarrow p \wedge q$.

P	q	$P \vee q$	$P \wedge q$	$P \vee q \rightarrow P \wedge q$
T	T	T	T	T
T	F	T	F	F
F	T	T	F	F
F	F	F	F	T

3. Show that the following are tautology.

$$\neg q \wedge (P \rightarrow q) \rightarrow \neg P$$

P	q	$\neg P$	$\neg q$	$(P \rightarrow q)$	$\neg q \wedge (P \rightarrow q)$	$\neg q \wedge (P \rightarrow q) \rightarrow \neg P$
T	T	F	F	T	F	T
T	F	F	T	F	F	T
F	T	T	F	F	F	T
F	F	T	T	T	T	T

$$2. \neg(\neg(p \rightarrow q) \wedge p \wedge (\neg p \rightarrow q)) \in F$$

P	q	r	$q \rightarrow r$	$\neg(q \rightarrow r)$	$\neg(q \rightarrow r) \wedge r$	$p \rightarrow q$	$\neg(p \rightarrow q)$	$\neg(\neg(p \rightarrow q) \wedge p \wedge (\neg p \rightarrow q))$	F
T	T	T	T	F	F	T	F	T	T
T	T	F	F	T	F	T	F	T	T
T	F	T	F	T	F	F	T	F	T
F	T	F	F	T	F	F	T	F	T
F	F	T	F	T	F	F	T	F	T
F	F	F	F	T	F	F	T	F	T

$$q = T \vee P$$

$$q = T \vee q$$

$$\text{not } p \vee q$$

$$T = T \wedge q$$

$$T = T \wedge q$$

$$\text{not } q \wedge q$$

Duality - law

The dual of compound preposition that contains only logical operators \vee, \wedge, \neg is the preposition obtained by replacing each joint by \wedge , each \wedge by \vee , each \neg by \neg , each T by F , each F by T where $T \& F$ are special variable representing compound preposition that are Tautology & contradiction resp. Dual of preposition A is denoted by A^* .

$$\text{ex, } A = (P \vee q) \wedge T \quad A = (P \wedge q) \vee T \quad \text{not } p \wedge q \wedge T$$

$$A^* = (P \wedge q) \vee F$$

Logical Equivalence.

The compound preposition $p \& q$ are said to be logically equivalent if $p \leftrightarrow q$ is a tautology
logical equivalence is denoted by \equiv or \leftrightarrow

$$P \wedge T \equiv P. \quad T \wedge T \equiv T$$

Tautological implication

Let $R \& S$ be two compound preposition Then

R is said to be logically imply S if

$\Rightarrow R \rightarrow S$ is a tautology is denoted by $R \Rightarrow S$

NOTE.

$R \rightarrow S$ is tautology $R \Rightarrow S$

$R \Leftarrow S$ is tautology $R \Rightarrow S$ or $R \Leftarrow S$.

Name of The law	Primal form	Dual form
1. Identity law	$P \vee F \equiv P$	$P \wedge T \equiv P$
2. Domination law (Null law)	$P \vee T \equiv T$	$P \wedge F \equiv F$
3. Idempotent (Self repeat)	$P \vee P \equiv P$	$P \wedge P \equiv P$
4. Double Negation (Involution law)	$\neg(\neg P) \equiv P$	
5. Commutative law	$P \vee q \equiv q \vee P$	$P \wedge q \equiv q \wedge P$
6. Associative law	$(P \vee q) \vee r \equiv p \vee (q \vee r)$	$(P \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
7. Distributive law	$P \vee (q \wedge r) \equiv (P \vee q) \wedge (P \vee r)$	$P \wedge (q \vee r) \equiv (P \wedge q) \vee (P \wedge r)$
8. De morgan's law	$\neg(P \vee q) \equiv \neg P \wedge \neg q$	$\neg(P \wedge q) \equiv \neg P \vee \neg q$
9. Absorption law	$P \vee (P \wedge q) \equiv P$	$P \wedge (P \vee q) \equiv P$
10. Negation law (Complement)	$P \vee \neg P \equiv T$	$P \wedge \neg P \equiv F$

1. Prove without truth table.

$$(P \wedge q) \rightarrow P \vee q \Leftrightarrow T$$

$$\begin{aligned}(P \wedge q) \rightarrow P \vee q &\equiv \neg(P \wedge q) \vee P \vee q && (\text{logical equivalence}) \\&\equiv (\neg P \vee \neg q) \vee P \vee q && (\neg P \equiv \neg P \vee q) \\&\equiv \neg P \vee (\neg q \vee (P \vee q)) && (\text{De Morgan's law}) \\&\equiv \neg P \vee ((\neg q \vee p) \vee p) && (\text{Associative law}) \\&\equiv \neg P \vee (T \vee p) && (\neg q \vee p \equiv T) \\&\equiv \neg P \vee T \\&\equiv T.\end{aligned}$$

2. Prove $(\neg P \vee q) \wedge (P \wedge (P \wedge q)) \equiv P \wedge q$ using laws.

$$\begin{aligned}\text{LHS} : (\neg P \vee q) \wedge (P \wedge (P \wedge q)) &\equiv (\neg P \vee q) \wedge ((P \wedge P) \wedge q) && (\text{Associativity}) \\&\equiv (\neg P \vee q) \wedge (P \wedge q) && (\text{Idempotency}) \\&\equiv (\neg P \vee q) \wedge (q \wedge P) && (\text{Commutativity})\end{aligned}$$

and elements P, q, r, s with p (Associative)

$$\equiv (\neg p \vee q \wedge r) \wedge p$$

order of operations is significant

$$\equiv (\neg p \vee q) \wedge p$$

$$\equiv (\neg p \vee q) \wedge p$$

$$\equiv \neg p \vee (q \wedge p)$$

$$= q \wedge p$$

if $\neg p \wedge q$, then if $\neg p$ is reduced, substitute

exist of $\neg p$ and reduce AND but $\neg p$ is not

3. Prove $P \rightarrow (q \rightarrow p) \equiv \neg p \rightarrow (p \rightarrow q)$

LHS: $P \rightarrow (q \rightarrow p) \equiv \neg p \vee (q \rightarrow p)$ (equivalence

involving Conditional)

$$\equiv \neg p \vee (\neg q \vee p)$$

$$\equiv \neg p \vee (p \vee \neg q)$$

$$\equiv (\neg p \vee p) \vee \neg q$$

$$\equiv T \vee \neg q$$

$$\equiv T \quad \text{--- } ① \quad (\text{Domination law})$$

RHS: $\neg p \rightarrow (p \rightarrow q) \equiv \neg(\neg p \vee p \rightarrow q)$

$$\equiv \neg(\neg p) \vee (\neg p \vee q) \quad (\text{Double Negation})$$

$$\equiv p \vee (\neg p \vee q)$$

$$\equiv (p \vee \neg p) \vee q$$

$$\equiv T \vee q$$

$$\equiv (T) \vee q \quad \text{--- } ②$$

from ① & ②, $T = (T) \vee q$

$$\textcircled{1} \equiv \textcircled{2}$$

$$\therefore P \rightarrow (q \rightarrow p) \equiv \neg p \rightarrow (p \rightarrow q)$$

All tautologies are equivalent to one another.

4. Determine which of the follow statements are tautologies or contradiction without truth table.

7. PRINCIPLE OF MATHEMATICAL INDUCTION

Let $P(n)$ be a statement involving positive integer prove that If $P(1)$ is true.

Induction: Consider $P(k)$ is true. Then $P(k+1)$ is also true. Then we can conclude the $P(k)$ is true for all positive integer.

1. Prove that $1 + 3 + 5 + \dots + (2n-1) = n^2$ by using mathematical induction.

$$\text{Let } P(n) = 1 + 3 + 5 + \dots + (2n-1) = n^2.$$

$$\text{Put } n=1$$

$$\text{And } P(1) = 2(1)-1 = 1$$

$$n^2 = 1^2 = 1$$

$\therefore P(1)$ is true.

Assume that the above statement is true for $P(k)$.

To prove that $P(k+1)$ is true.

$$\begin{aligned} P(k+1) &= 1 + 3 + 5 + \dots + (2k-1) + (2(k+1)-1) = (k+1)^2 \\ &= 2(k+1) - 1 + k^2 \\ &= 2k + 2 - 1 + k^2 \\ &= 2k + 1 + k^2 = (k+1)^2 \end{aligned}$$

$$2. 1+2 \cdot 3+3 \cdot 4+\dots+n(n+1) = \frac{n(n+1)(n+2)}{3}$$

$$P(n) = 1 \cdot 2 + 2 \cdot 3 + 3 \cdot 4 + \dots + n(n+1) = \frac{n(n+1)(n+2)}{3}.$$

$$P(1) = 1 \cdot 2 + 2 \cdot 3 + 3 \cdot 4 + \dots + 1(1+1) = \frac{1(1+1)(1+2)}{3}.$$

$$2 = \frac{2(3)}{3}$$

$$2 = 2.$$

$P(1)$ is true.

$$P(k) = 1 \cdot 2 + 2 \cdot 3 + 3 \cdot 4 + \dots + k(k+1) = \frac{k(k+1)(k+2)}{3}.$$

To prove $P(k)$ is true.

Then $P(k+1)$ should be true.

$$P(k+1) = 1 \cdot 2 + 2 \cdot 3 + 3 \cdot 4 + \dots + k(k+1) + (k+1)(k+1+1) \\ = \frac{(k+1)(k+2)(k+3)}{3}$$

$$k(k+1) + (k+1)(k+2)$$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

After addition of $k^2 + k + k^2 + 3k + 2$

3. Prove $n! \geq 2^{n-1}$ for all $n \geq 1$

$$P(n) = n! \geq 2^{n-1} \quad \forall n \geq 1$$

$$P(1) = 1! \geq 2^0$$

$$P(1) = 1! \geq 1$$

$\therefore P(1)$ is true.

Let $P(k)$ is true.

$$P(k) = k! \geq 2^{k-1}$$

The $P(k+1)$ must be true.

$$P(k+1) = (k+1)! \geq 2^{k+1-1}$$

$$\Rightarrow k! (k+1) \geq 2^k$$

$$\text{we know } k! \geq 2^{k-1}$$

$$2^{k-1} (k+1) \geq 2^k$$

$$2^{k-1} \times k + 2^{k-1} \geq 2^k$$

$$2 \cdot 2^{k-1} \geq 2^k$$

$$2^k \geq 2^k$$

4. Prove $8^n - 3^n$ is divisible by 5.

Let $P(n) = 8^n - 3^n$ is divisible by 5

$$P(1) = 8^1 - 3^1 = 5 \text{ is divisible by 5.}$$

$P(1)$ is true.

Assume $P(k) = 8^k - 3^k$ is divisible by 5 is true

$$P(k+1) = 8^{k+1} - 3^{k+1}$$

$$= 8^k \cdot 8 - 3^k \cdot 3$$

$$= 8^k (5+3) - 3^k \cdot 3$$

$$= 3(8^k - 3^k) + 5 \cdot 8^k = 3(5^k) + 5(8^k) = 5 [(5^{k-1}) 3 + 8^k]$$

is divisible by 5

ARGUMENT - Sequence of statement that ends in conclusion. Set of one or more premises and conclusion. The argument is true only if all premises true and a conclusion false.

THEORY OF INFERENCE

1. If I love cat then I love dog : $P \rightarrow q : H$

I love cat $\rightarrow P$

I love dog = q

If I love cat then I love dog : $P \rightarrow q$.

$$\frac{P \rightarrow q \wedge T}{\therefore q}$$

$P \rightarrow q \wedge T$
 $P \rightarrow q$ is T and

2. Determine whether $H_1 : P \rightarrow q \quad H_2 : P \cdot H_1 \wedge H_2$

P	q	$P \rightarrow q$	$H_1 \wedge H_2$: q
T	T	T	$T \wedge T$	T
T	F	F	$F \wedge P$	F
F	T	T	$F \wedge F$	T
F	F	T	$F \wedge F$	F

Rule P use the P

1- If it rains heavily then travelling will be difficult. If students arrive on time then travelling will not be difficult. They arrive on time \vdash

Rule T use the rules of inference

\neg It rains heavily : P.

Travelling is difficult : q

Rule CP \rightarrow Conditional Proof

Student arrive on time \vdash

H₁: $P \rightarrow q$

H₂: $\gamma \rightarrow \neg q$

H₃: γ

$C \rightarrow \neg P$

$P \rightarrow q$

$\gamma \rightarrow \neg q$

γ

$\neg P$

Step no	Statement	Reason
1	$P \rightarrow q$	Rule P
2	$\neg q \rightarrow \neg P$	Rule T Contrapositive
3	$\gamma \rightarrow \neg q$	Rule P
4	$\gamma \rightarrow \neg P$	Rule T using hypothetical Syllogism
5	γ	Rule P
6	$\neg P$	Rule T modus Ponens

\therefore The Conclusion is Valid.

2. show that $t \wedge s$ can be derived from the premises, $P \rightarrow q$, $q \rightarrow t \wedge s$, γ , $P \vee (\neg t \wedge s)$

Step no	Statement	Reason
1.	$P \rightarrow q$	Rule P.
2.	$q \rightarrow t \wedge s$	Rule P
3.	$P \rightarrow t \wedge s$	Rule T, hypothetical Syllogism.
4.	γ	Rule P
5.	propositio $\neg t \wedge s$	Rule T, modus ponen
6.	$P \vee (\neg t \wedge s)$	Rule P
7.	$\neg t \wedge s$	Rule T.