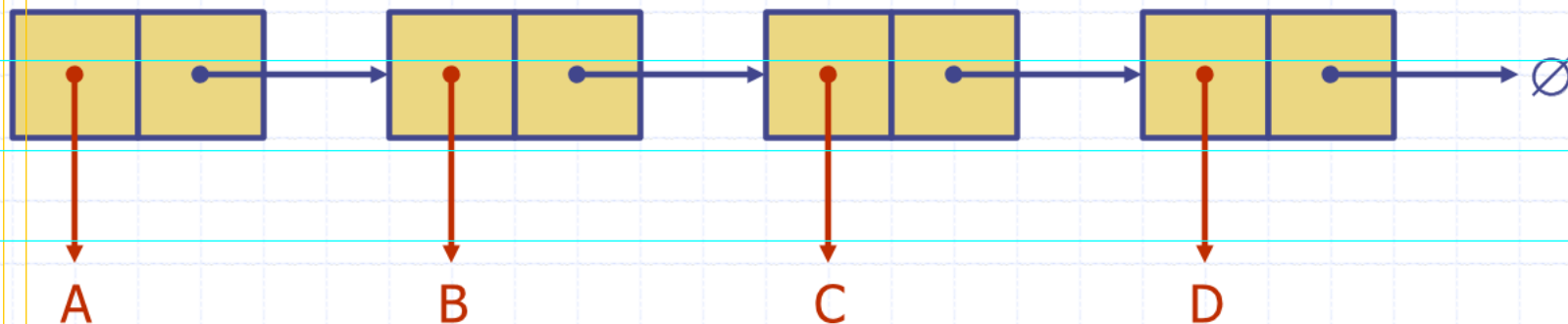
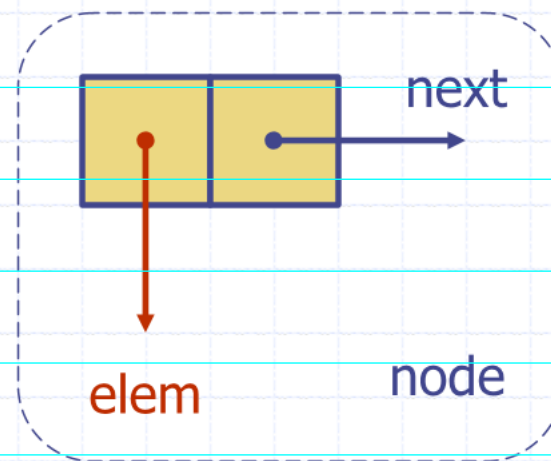


# Singly Linked List (§ 4.4.1)

◆ A singly linked list is a concrete data structure consisting of a sequence of nodes

◆ Each node stores

- element
- link to the next node

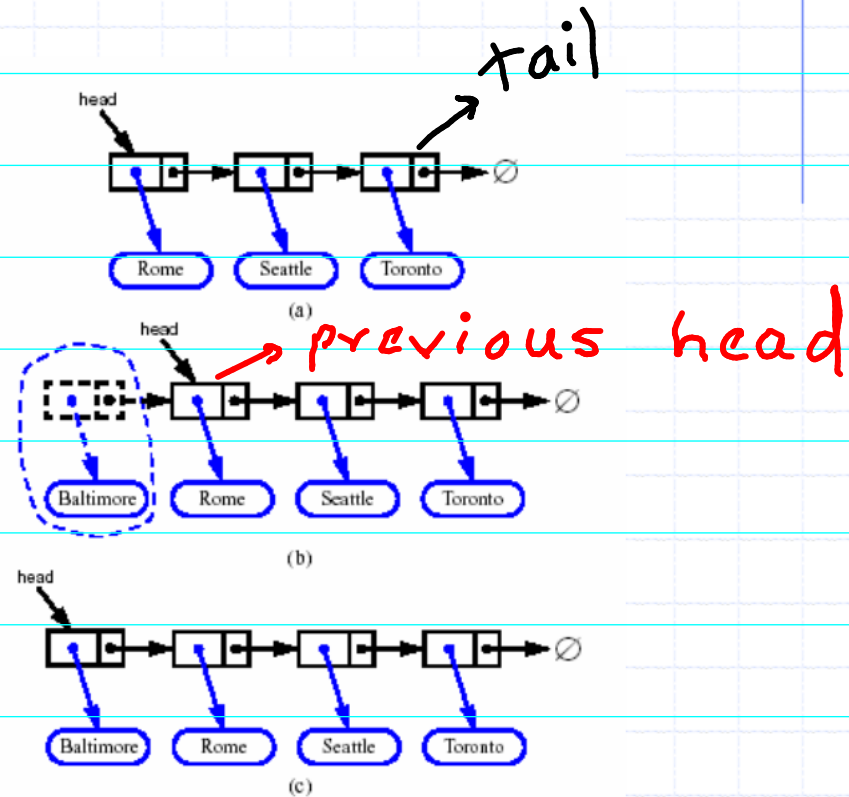


# The Node Class for List Nodes

```
public class Node {
    // Instance variables:
    private Object element;
    private Node next;
    /** Creates a node with null references to its element and next node. */
    public Node() {
        this(null, null);
    }
    /** Creates a node with the given element and next node. */
    public Node(Object e, Node n) {
        element = e;
        next = n;
    }
    // Accessor methods:
    public Object getElement() {
        return element;
    }
    public Node getNext() {
        return next;
    }
    // Modifier methods:
    public void setElement(Object newElem) {
        element = newElem;
    }
    public void setNext(Node newNext) {
        next = newNext;
    }
}
```

# Inserting at the Head

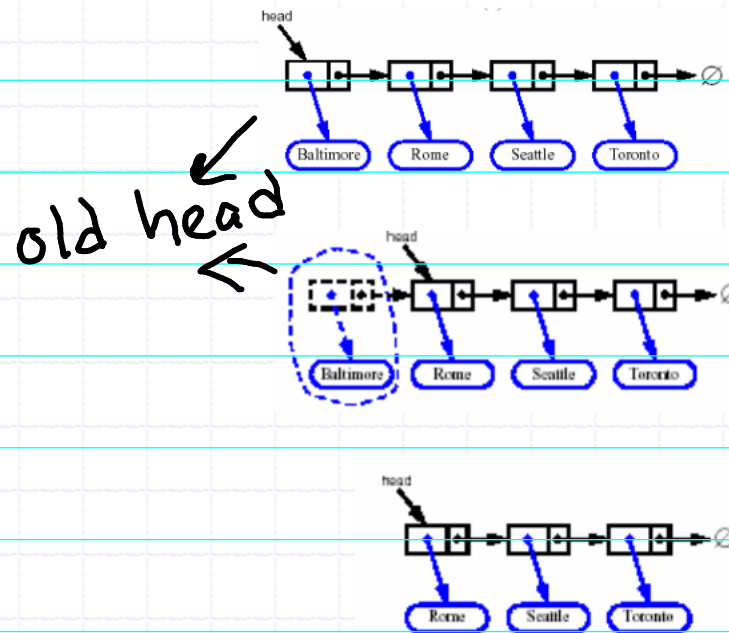
1. Allocate a new node
2. Insert new element
3. Have new node point to old head
4. Update head to point to new node



head : first element  
 $O(1)$

# Removing at the Head

1. Update head to point to next node in the list
2. Allow garbage collector to reclaim the former first node



# Delete method for linked list

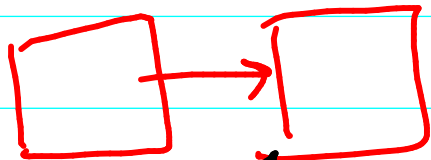
```
public void deletehead ( ) {
```

```
    head = head.getNext ( ) ;
```



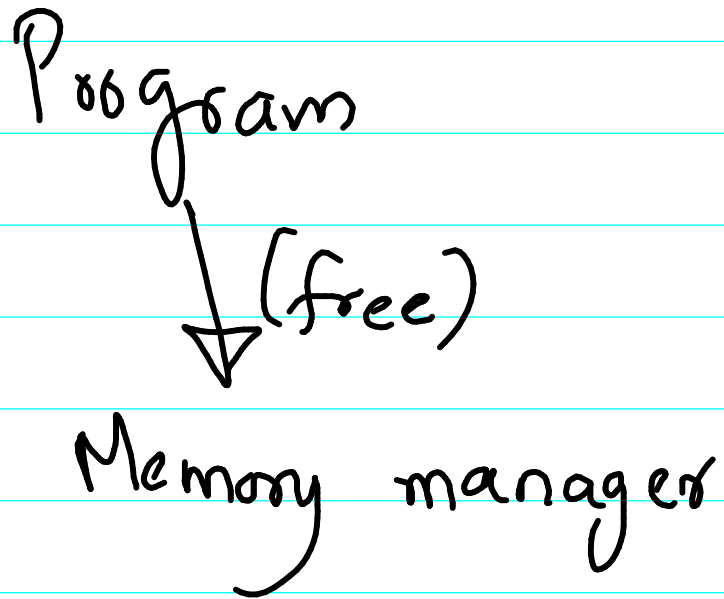
In C++ you would next "free" head

In Java, garbage collector, run periodically  
by the Virtual M/c (JVM) does this  
freeing

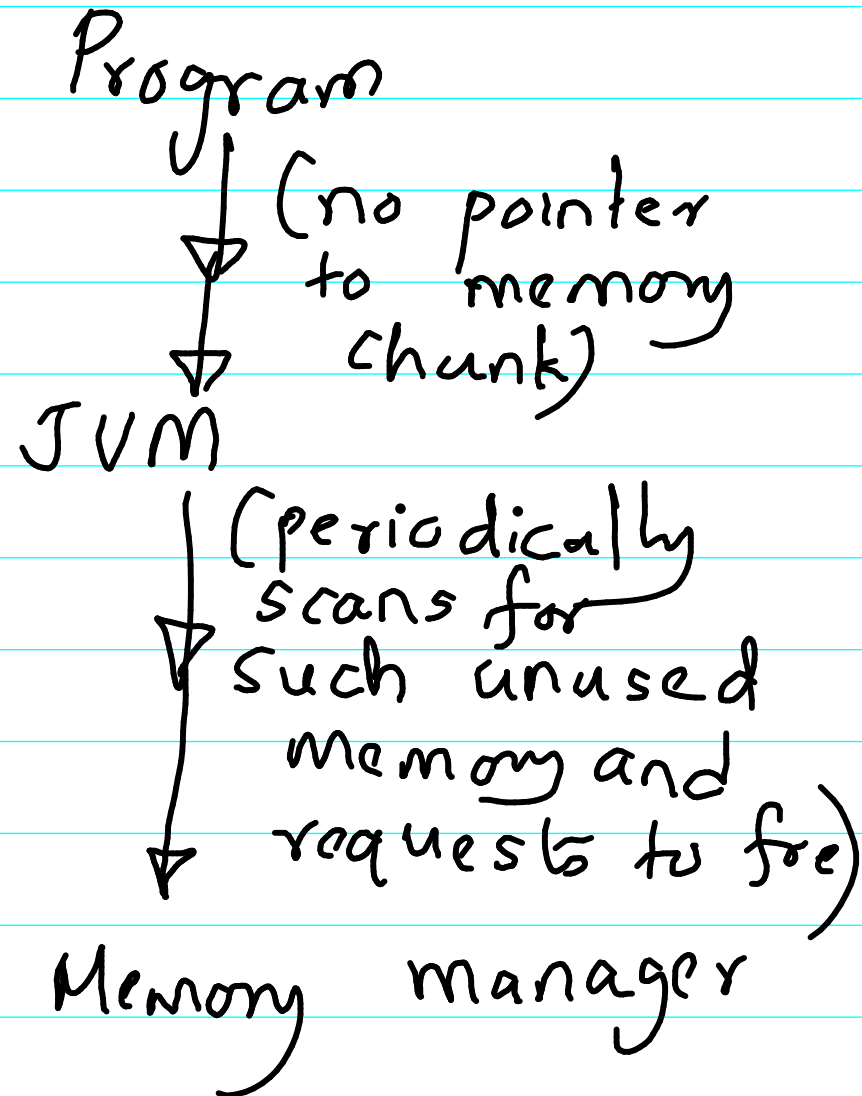


head → head.next

C++



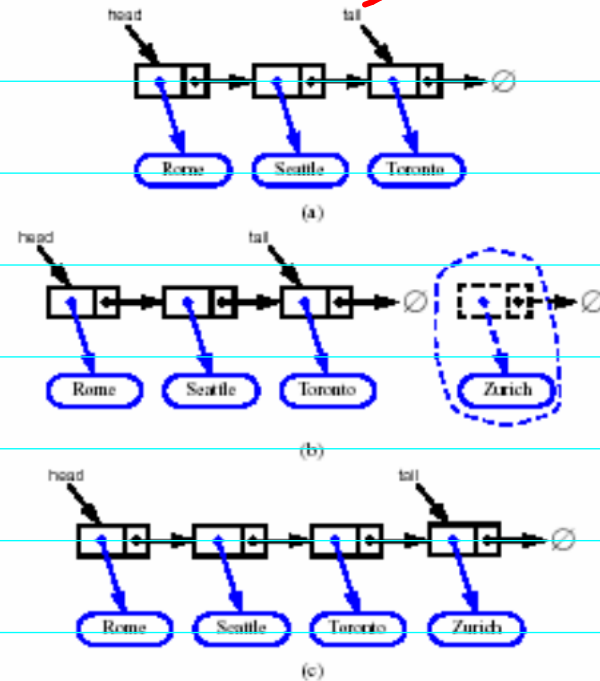
Java



# Inserting at the Tail

(Assuming you have a pointer to tail)

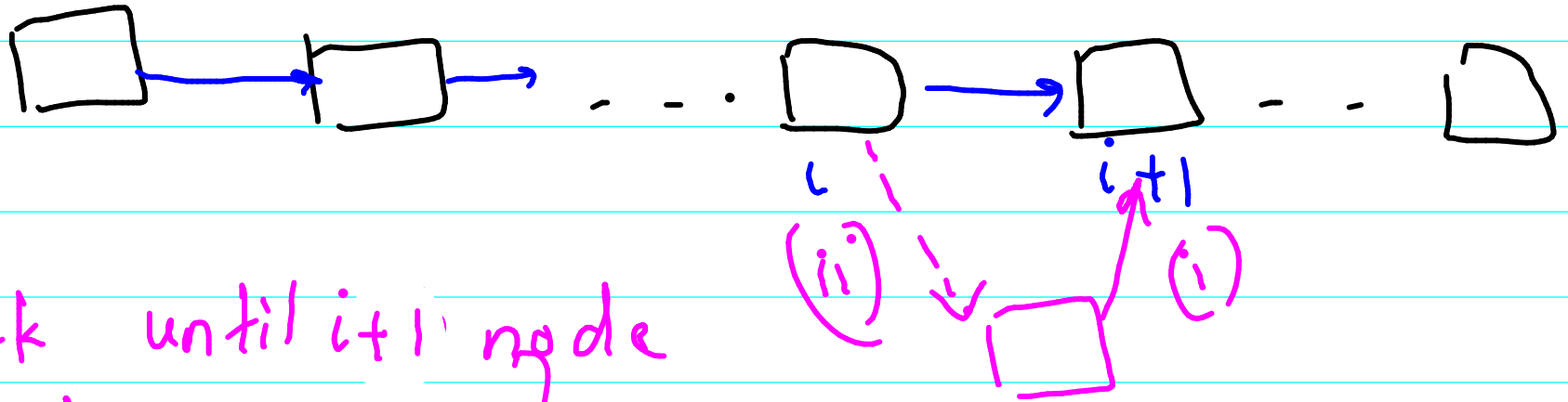
1. Allocate a new node
2. Insert new element
3. Have new node point to null
4. Have old last node point to new node
5. Update tail to point to new node



$O(1)$

$O(N)$  if you don't keep track of tail.

Element Insertion between  $i^{\text{th}}$  node &  $i+1^{\text{th}}$  node



(1) Seek until  $i+1$  node

(2) Create a new node for new element

(3) Set "next" pointers of  $i^{\text{th}}$  node & new node

Approx  $i+3$  operations

Complexity specified in terms of

input (size)  $N = O(N)$  since

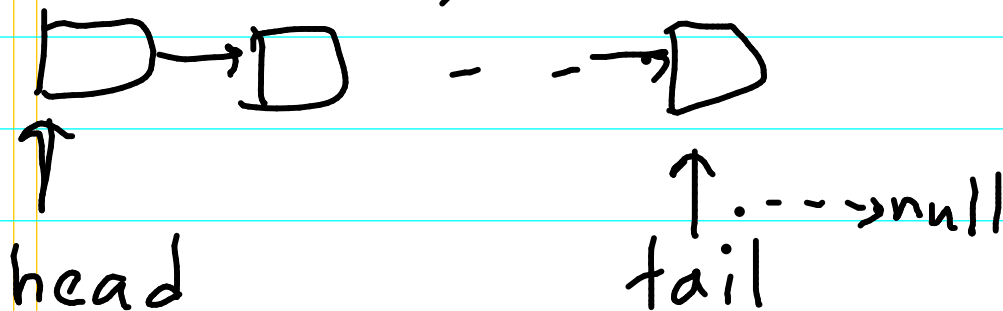
$i$  can be  $N-1$  in worst case



# Removing at tail

①  $tail = null$

(NO!)



② Traverse till

a)  $node \cdot getNext()$ .

$getNext()$

$= null$

b)  $node \cdot setNext(null)$

c)  $tail = node$

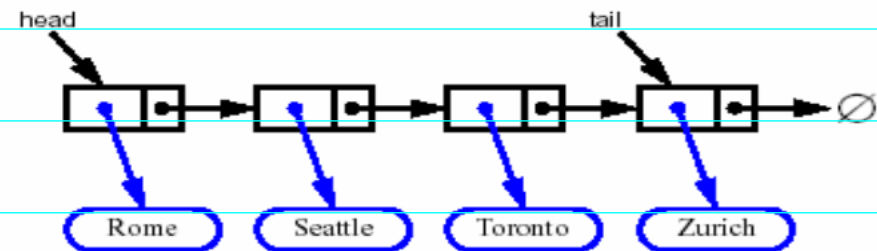
$O(N)$

(YES)

First check for corner cases: ie  $head = tail$   
(list containing single element)

# Removing at the Tail

- ◆ Removing at the tail of a singly linked list is not efficient!
- ◆ There is no constant-time way to update the tail to point to the previous node



Q: Suppose you want to implement a "list", not as Linked list but as an array of Nodes, i.e

Node[] nodeList

so that I could (a) insert at head

(b) insert at tail

(c) delete at head

(d) delete at tail.

Design a data structure that uses Node[] to implement above