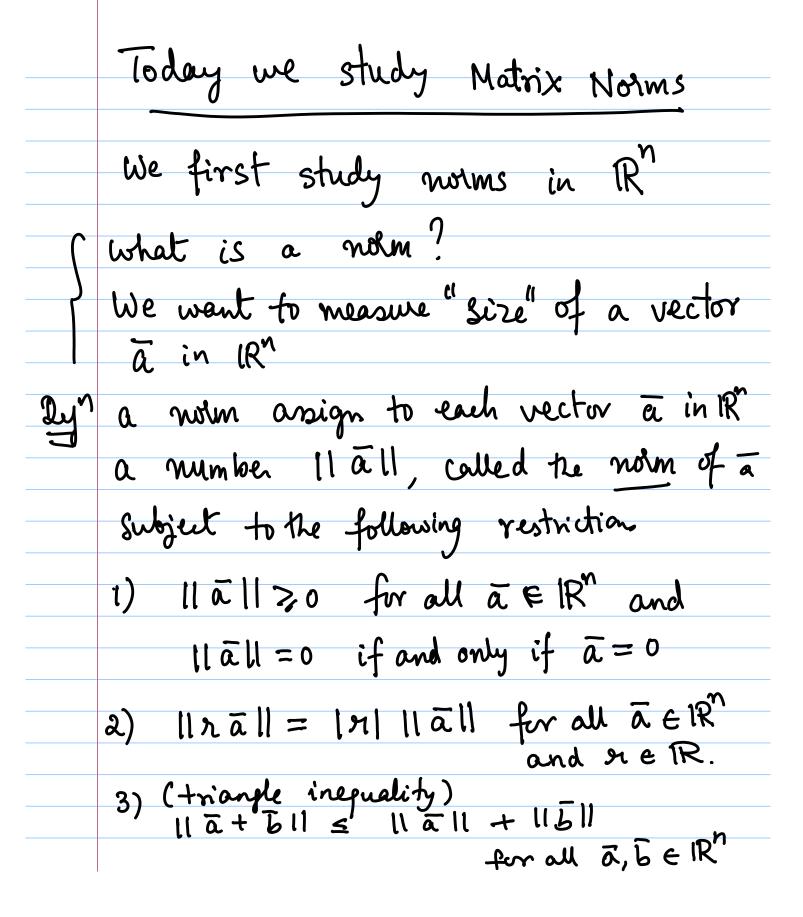
Lecture 17 Last time we did · Choleskey's factorization of a possible definite matrix A A = Llt where L is lower triangular. Their requires only 50% of the calculation needed for LU decomposition. . If GE has row changes Then we can factor A into PF = LU where P is a permutation matrix.

solve Ax=b PAX = Pb = b 1.08=6 set y = Ux solve Ly = 6 and then solve Ux = y Matrix factorization is useful if we have to solve Ax = b for many different values of bWe then turned one attention to round-off errors that can occur while doing Gt. we learn scaled partial pivoting for GE.

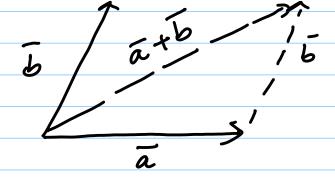


Note the first restriction forces all n-vectors but the zero vector to have possitive "length".

The second restriction states for exam

that \bar{a} and its negative $-\bar{a}$ have the same length and that the $3\bar{a}$ has length 3 times that of \bar{a}

The third restriction is called triangle inequality



Examples of norms 1) Euclidean norm a = (a, a2, -., an) $\|\bar{a}\|_{2} = \sqrt{\alpha_{1}^{2} + \alpha_{2}^{2} + \cdots + \alpha_{n}^{2}}$ Euclidean norm is also called the Easy to verify that lialla > 0 and = 0 iff a= 0 ותן וותו To prove triangle inequality we need the "Cauchy-Schwarz inepality

Theorem For 2, 7 E 1R" 7. 7 < 11211, 114112 y=0 or n=0 then nothing to prom So suppre 72 to and 7 to For $x \in \mathbb{R}$ $0 \le \| x - y \|_{2}^{2} = \sum_{i=1}^{n} (x_{i} - y_{i})^{2}$ $= \sum_{i=1}^{\infty} \chi_i^2 - 2\chi \sum_{i=1}^{\infty} \chi_i \chi_i + \chi^2 \sum_{i=1}^{\infty} \chi_i^2$ 2> \[\frac{1}{2} \times \frac{1}{2} + \frac{1}{2} \ldots \ Since 11201/2>0 and 11811/2>0 take $\lambda = \frac{||\lambda||_2}{||\lambda||_2}$

$$\left(2\frac{\|(\chi\|_{2})}{\|\eta\|_{2}}\right)\left(\sum_{i=1}^{n}\chi_{i}\chi_{i}\right) \leq \|\chi\|_{1}^{2} + \|\chi\|_{1}^{2} \|\eta\|_{1}^{2}$$

$$= 2\|\chi_{2}\|^{2}$$

$$\vdots = \|\chi\|_{1}\|\eta\|_{2} \|\eta\|_{2}$$
We now prove triangle inequality
$$\|\chi + \chi\|_{2}^{2} = \sum_{i=1}^{n} (\chi_{i} + \chi_{i})^{2}$$

$$= \sum_{i=1}^{n} \chi_{i}^{2} + 2 \sum_{i=1}^{n} \chi_{i} \chi_{i} + \sum_{i=1}^{n} \chi_{i}^{2}$$

$$\leq \|\chi\|_{2}^{2} + 2 \|\chi\|_{2} \|\eta\|_{2} + \|\eta\|_{2}^{2}$$

$$= (\|\chi\|_{2} + \|\chi\|_{2})^{2}$$
Thus $\|\chi + \chi\|_{2} \leq \|\chi\|_{2} + \|\eta\|_{2}$

1- norm

 $\overline{a} = (a_1, a_2, -, a_n)$

11 all, = la,1+la21+---+lan)

It is easy to see that

llāll, >0 and =0 iff a=0

also Urālly = Inlial,

11a+b11, = la,+b,1+ laz+b2l+ + - - + la,+b,1

< 1911+161+ + 1021+1621+ --- +191+169

= 11a41 + 1161/2

Thus latbly 5 lall, + libly

00-nom

a= (a,,, a,)

llāll_o = max [ai]

One can easily verify that ILallo is also a norm on IR".

If II II is a norm in IR then we can define distance between two vectors \bar{x}, \bar{y} as $dist(\bar{x}, \bar{y}) = 11\bar{x} - \bar{y}11$

Matrix norm

A matrix norm on the set of all nxn matrices is a real valued function, II II defined on this set satisfying for all nxn matrices A, B and all real numbers of

(1811+11A11 = 1181+A11 (iii)

(iv) NABII < NANUBII

Theorem If IIII is a norm in IR then

II A II = max II A x II ハメリニュ is a matrix norm. Front clearly NAIL > 0 11All=0 => 11Ax11=0 for all 11211=1 $\Rightarrow A \times = 0$ for all \times with $11 \times 11 = 1$ Claim A = 0 otherwise some column say i is non-zero P. be i'th co-ordinate vech Ae; #0 $u = \frac{ei}{110.11}$ has not 1 ાા ૯.ના

$$(AB)x = A(Bx)$$

for
$$\overline{z} \neq 0$$
 $\overline{x} = \frac{\overline{z}}{||\overline{z}||}$ is a with vector

$$\frac{||A|| = \max ||A \times || = \max ||A \left(\frac{Z}{||Z||}\right)||}{||X|| = ||Z \neq 0||}$$

$$\begin{aligned} \|A\|_2 &= \max \|Ax\|_2 & \text{Euclidean norm} \\ \|A\|_1 &= \max \|Ax\|_1 & \text{I-norm} \\ \|x\|_1 &= 1 \end{aligned}$$

$$\|A\|_1 &= \max \|Ax\|_2 & \text{I-norm} \\ \|x\|_2 &= 1 \end{aligned}$$

Eudidean norm is difficult to calculate However one can calculate both the I-norm and the os-norm || All = max \ \frac{n}{5=1} $||A|| = \max_{1 \le j \le n} \sum_{i=1}^{N} |a_{ij}|$ [(A((Sums yours 11 Ally sums columns Example 11 All = max {8, 6, 6} = 8

Proof of Theorem

We want to show

$$|A||_{\infty} = \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|$$

We first show

 $|A||_{\infty} \le \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|$
 $|A||_{\infty} \le \max_{1 \le i \le n} |x_{i}|$
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Thus
$$||A|| = \max_{1 \le i \le n} ||a_{ij}|| \le \sum_{j=1}^{n} |a_{ij}||$$

Thus $||A \times ||_{\mathcal{S}} = \max_{1 \le i \le n} |\sum_{j=1}^{n} |a_{ij}||$
 $\leq \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}||$

Thus $||A|| = \max_{1 \le i \le n} ||A \times ||_{\mathcal{S}} = ||a_{ij}||$
 $\leq \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}||$
 $\leq \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}||$

We will now show the apposite inequality

Let p be an integer with

 $\sum_{1 \le i \le n} ||a_{ij}|| = \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}||$
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Let
$$\overline{U}$$
 be the vector with components

 $U_j = \begin{cases} 1 & \text{if } ap_j \ge 0 \\ -1 & \text{if } ap_j < 0 \end{cases}$
 $U_j = \begin{cases} 1 & \text{if } ap_j < 0 \end{cases}$
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$$||A||_{\infty} = \max_{||A||_{\infty}} ||A||_{\infty}$$

$$||A||_{\infty} = ||A||_{\infty}$$

$$||A||_{\infty} = \max_{|S| \le i \le n} \sum_{j=1}^{\infty} |a_{ij}|$$

$$||A||_{\infty} = \max_{|S| \le i \le n} \sum_{j=1}^{\infty} |a_{ij}|^{2}$$

$$||A||_{E} = \left(\sum_{i=1}^{\infty} \sum_{j=1}^{\infty} |a_{ij}|^{2}\right)^{\frac{1}{2}}$$
One can from that the Frobenius norm is a mateix norm

Reasons for studying metrix nolms if x is computed solution of then its error is the difference ē = x - x This evol is unknown to since we do not know But we can always compute the " residual erun" $\vec{\lambda} = A \times - A \hat{\lambda} = b - A \hat{\lambda}$ If I = 0 then 2 is exact sol. One would expect it to be small if

n is close to n This is not so as the following example shows Example 1.012, + 0-99x2 = 2 0-99×1+1-01×2 = 2 unique sol n = n2=1 $\hat{\chi} = \begin{bmatrix} 2 \\ 6 \end{bmatrix}$ has even $\bar{e} = \begin{bmatrix} -1 \\ 1 \end{bmatrix}$ but residual error 2) By taking a diff right side we can achieve the opp effect

$$| \cdot 0 | \chi + 0 - 99 \chi_2 = 2$$

$$| \cdot 0 | \chi + | \cdot 0 | \chi_2 = -2$$

exact answer $x_1 = 100$, $x_2 = -100$.

$$\hat{\lambda} = \begin{bmatrix} 101 \\ -99 \end{bmatrix}$$
 has small error

but residual

$$\bar{x} = \begin{bmatrix} -2 \\ -2 \end{bmatrix}$$
 is large

Thus the residuel $\bar{x} = \bar{b} - A\bar{x}$ is notalways a reliable indicator of the size of the error $\bar{e} = \bar{x} - \hat{x}$.

It depends on the size of the matrix
$$A$$
 and A^{-1} .

 $\overline{\lambda} = A \times - A \hat{\chi} = A (x - \hat{\chi})$
 $= A \overline{e}$
 $\overline{e} = A^{-1} \overline{x}$

Yemark If matrix A is invertible

 $\overline{u} = A^{-1}(A u)$
 $\overline{u} = A^{-$

on the relative error Itell in 11211 terms of relative residuel 1/11/ 11611 11 All 11211 11611 4 11811 5 11211 11611 11/11 1161) 11611 5 11×11 5 1(A-1111161) (* Applied to u A11 are put this in above equation to get cond (A) = ||A|| ||A^-||) is called condition number

$$\frac{||A||}{||A||} \leq \frac{||A||}{||A||} \leq \frac{||A||}{|$$