

Running time of programs

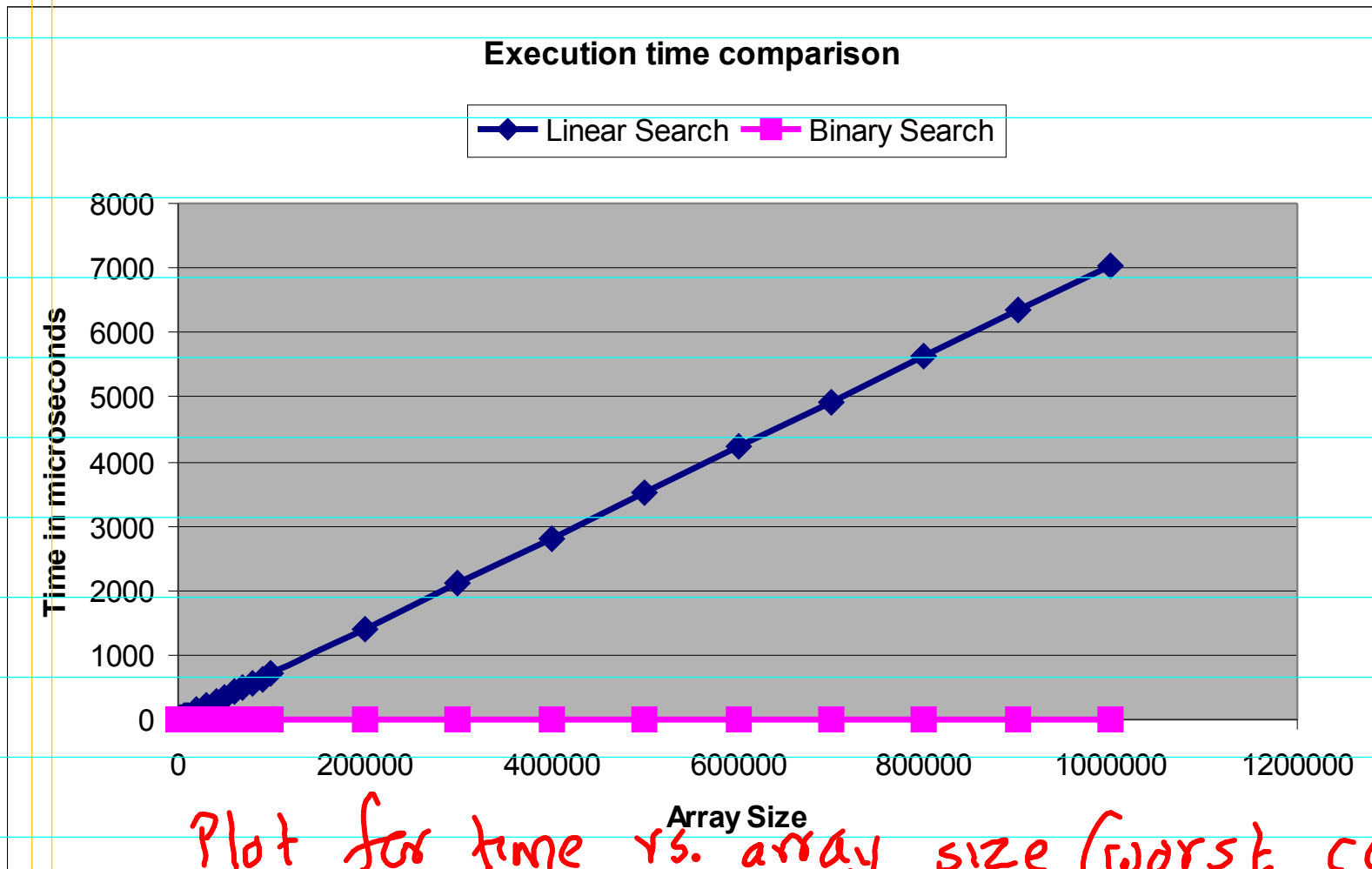
- Running time is a function of?
 - Array size
 - Number position
- Rigorous, scientific way to do find running time?
 - Run the program and time it (must properly design running time experiments)
 - Determine "input" (Array A, number to search for)

"Proper" for fixed i/p's for fair comparisons

- ① No background processes
- ② Comparable data types (structures)
- ③ Single threaded / no wasted steps

④ To normalise wrt OS/kernel related background processes ⑤ average across multiple runs for each program for fixed i/p's or ⑥ use virtual machines with guaranteed resources

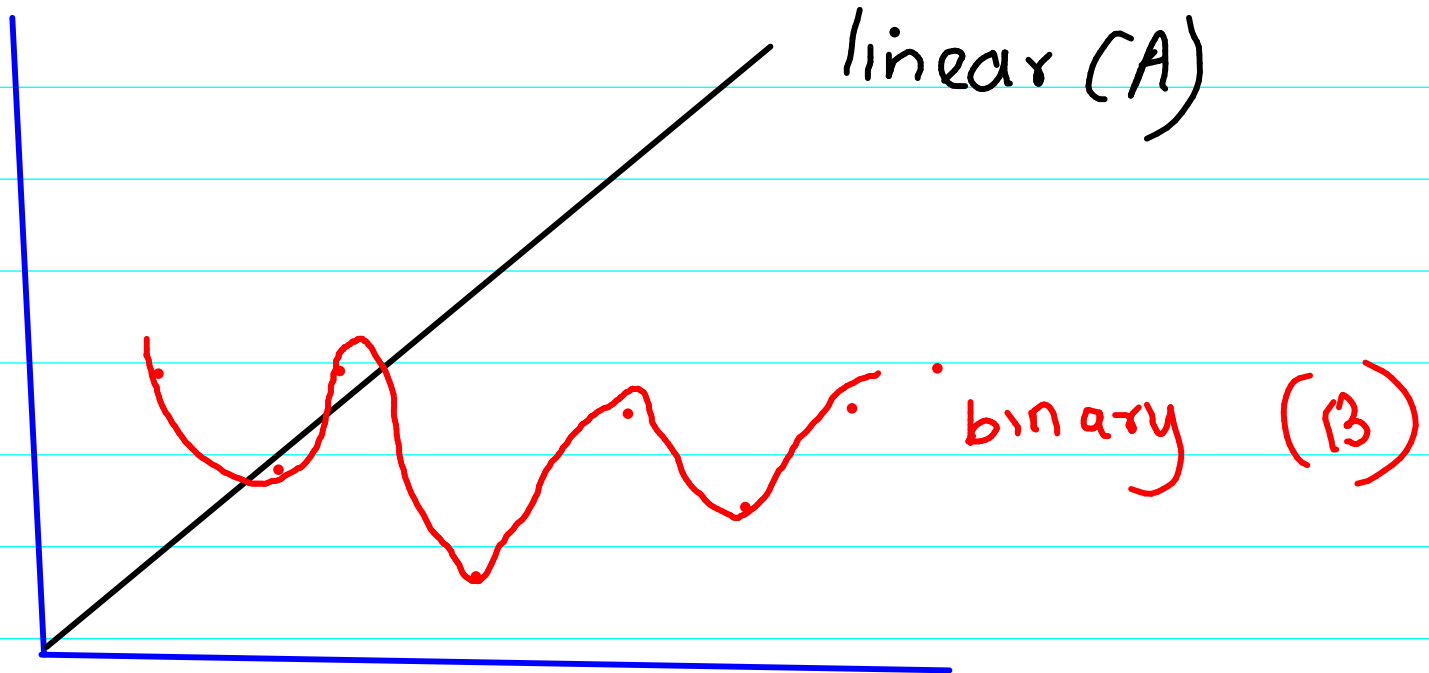
Results of experiment



*Worst case

Plot for time vs. array size (worst case over position of num being searched)

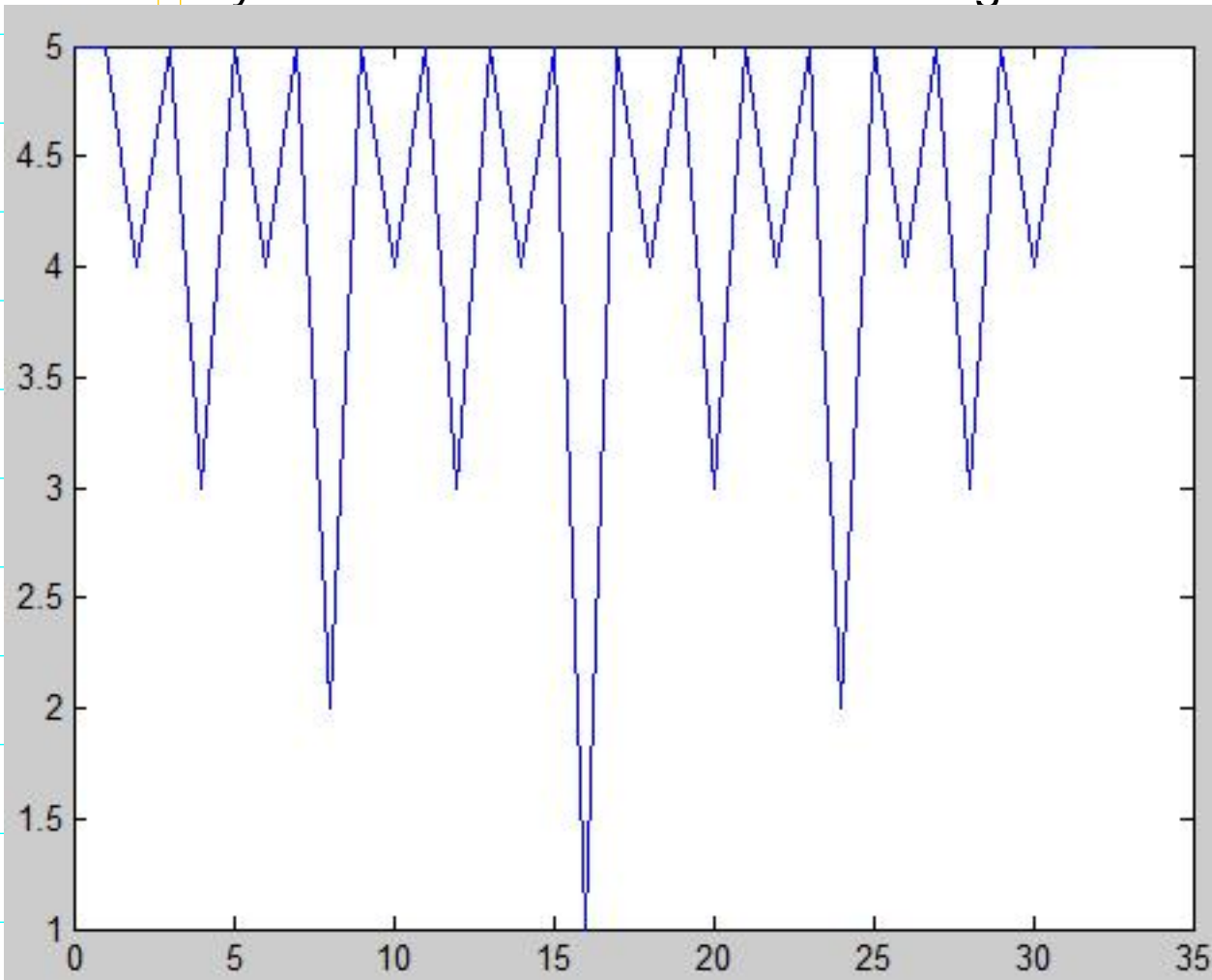
Q: How will plot(s) of time vs position of num being searched for look like (for fixed array size)



Look at answer posted by Sasank Chilamkurthy

at <https://piazza.com/class/hq4vn7sblku32q?cid=11>

Assume length to be of 2^n form. Write index of our element in binary and count n = no. of trailing zero. n is our answer.



Disadvantages

- Can take too much time
- Affected by too many factors (the hardware, the compiler, etc)
- Too much detailed – we just need to know the “essential behaviour”
- And most importantly
 - Never do an experiment before trying to reason about the outcome!!

i.e first set your expectation right

Algorithm Analysis

- For any system, if pen-paper mathematical analysis is possible *one must always do this. Why?*
 - Usually much faster
 - Zooms in on the relevant details, allows to ignore unnecessary details
 - Gives a *fundamental* “environment free” understanding of a system
- But such analysis *always* requires a “model” of the system under study

Model for Algorithm Analysis

- “Normal” computer – sequential instructions
- All “basic” instructions take one unit of time – addition, multiplication, comparison, assignment
- Computer has infinite memory
- Clearly, some aspects are ignored – disk read times, “paging”, context switching, etc
 - But this is intentional
 - We want to understand the *algorithm*, not implementation issues

Imagine running program in “debug” mode,
of “return” presses = # of lines executed

Analysis of search algorithm A

Let size be N . Assume element is at index $n = 0 \dots N-1$

```
for (i=0; i < size; i++) {  
    if (A[i] == num) {  
        found = true;  
        break;  
    }  
}
```

Assignment: 1

Array access: 1
Comparison: 1

$(0 \dots, n) \times (n+1)$

Assignment: 1

+1

for break

Comparison: 1

$\times (n+1)$

Increment: 1

$\times n$

$\text{int } i$: primitive
 $\text{int}[] A$: non-primitive

Some of these are done multiple times – which? How many times?

Search algorithm A: time for successful search

- Total time ($T_s(n)$), when element is at index $n=0 \dots N-1$ (successful search)

- $1 + 2 \times (n+1) + 1 + 1 \times n + 1 \times (n+1)$

- $T_s(n) = 4n + 5 + \underbrace{1}_{\text{for break}}$

Analysis of search algorithm A

Now assume search element is not present

```
for (i=0; i < size; i++) {  
    if (A[i] == num) {  
        found = true;  
        break;  
    }  
}
```

Assignment: 1

Array access: 1
Comparison: 1

$\times N$

Comparison: 1

$\times (N+1)$

Increment: 1

$\times N$

Note: The time reqd for this case is the same as the time reqd when $n = N-1$

Search Algorithm A: time for unsuccessful search

- Total time for unsuccessful search:

- $1 + 2 \times N + 1 \times N + 1 \times (N+1)$

- $T_u(N) = 4N + 2$

Next Question: [H/W]

How would you compute the average number of instructions executed by program A, where you average across all possible values of 'num' while holding the length of the list (that is, the value of end-begin) as a constant