

Lecture 20

The solution of Non-linear equations

One of the most frequently occurring problems in scientific work is to find the roots of equations of the form

$$f(x) = 0$$

In general we can only hope for "approximate solutions" i.e. a pt x^* for which $|f(x^*)|$ is "small".

Some iterative methods of not finding

1) Bisection method

Recall intermediate value theorem
If ^{f is cts} $f(a)f(b) < 0$ then f has
a zero in (a, b) .

Suppose $f(a)f(b) < 0$
set $a_0 = a$ $b_0 = b$

$$m = \frac{a_0 + b_0}{2}$$

* If $f(a_0)f(m) < 0$ then set
 $b_1 = m$ $a_1 = a_0$

otherwise set $a_1 = m$ $b_1 = b$
so not lies in $[a_1, b_1)$

Example $f(x) = x^2 - 2$

$$a_0 = 1 \quad f(a_0) = -1$$

$$b_0 = 2 \quad f(b_0) = 2$$

$$m = \frac{1+2}{2} = 1.5$$

$$f(1.5) = 2.25 - 2 = 0.25$$

So $a_0 = 1$ $b_1 = 1.5$
root lies in $[1, 1.5]$

Algorithm for Bisection method

Given a function continuous in $[a_0, b_0]$
and such that $f(a_0)f(b_0) < 0$

For $n = 0, 1, 2, \dots$ until satisfied do

$$\text{Set } m = \frac{a_n + b_n}{2}$$

If $f(a_n)f(m) < 0$ Then set $a_{n+1} = a_n$
 $b_{n+1} = m$

Otherwise set $a_{n+1} = m$, $b_{n+1} = b_n$

Then $f(x)$ has a root in the interval
 $[a_{n+1}, b_{n+1}]$

Example

$$f(x) = x^2 - 2$$

(4 sig digits)

n	a_n	b_n	$f(a_n)$	$f(b_n)$
0	1	2	-1	2
1	1	1.5	-1	2.5 E-1
2	1.25	1.5	-4.375 E-1	2.5 E-1
3	1.375	1.5	-1.094 E-1	2.5 E-1
4	1.375	1.438	-1.094 E-1	6.641 E-1
5	1.407	1.438	-2.035 E-2	6.641 E-1
6	1.407	1.423	-2.035 E-2	2.493 E-2
7	1.407	1.415	-2.035 E-2	2.225 E-3
8	1.411	1.415	-9.079 E-3	2.225 E-3
	1.413	1.415	-3.431 E-3	2.225 E-3
10	1.414	1.415	-6.040 E-4	2.225 E-3

$$\frac{1.414 + 1.415}{2} = 1.415 \quad \text{in 4 sig digits}$$

So algorithm ends.

root lies in $[1.414, 1.415]$

- Bisection method always converges to the root
 - Convergence is slow
-

One can hope to get to root faster by using fully the information about $f(x)$ available at each step.

In our example

$$f(x) = x^2 - 2$$

$$f(1) = -1 \quad \text{and} \quad f(2) = 2$$

Since $|f(1)|$ is closer to zero than $|f(2)|$ the root ξ is likely to be closer to 1 than 2.

Hence rather than check the midpt or avg value of 1 and 2, we check $f(x)$ at the weighted average

$$w = \frac{|f(2)| \cdot 1 + |f(1)| \cdot 2}{|f(2)| + |f(1)|}$$

Since $f(1), f(2)$ have opposite sign

$$w = \frac{f(2) \cdot 1 - f(1) \cdot 2}{f(2) - f(1)}$$

In our example

$$w = \frac{2 + 2}{3} = 1.333$$

$$f(w) < 0$$

So the root lies in $[1.333, 2]$

Repeating the process we get

$$w = 1.400 \quad \text{and so on}$$

This algorithm is known as regula-falsi or false-position method.

Algorithm (Regula-falsi)

Given a function $f(x)$ continuous on the interval $[a_0, b_0]$ and such that $f(a_0)f(b_0) < 0$

For $n = 0, 1, 2, \dots$, until satisfied do

Calculate
$$w = \frac{f(b_n)a_n - f(a_n)b_n}{f(b_n) - f(a_n)}$$

If $f(a_n)f(w) \leq 0$, set $a_{n+1} = a_n$, $b_{n+1} = w$

Otherwise set $a_{n+1} = w$, $b_{n+1} = b_n$

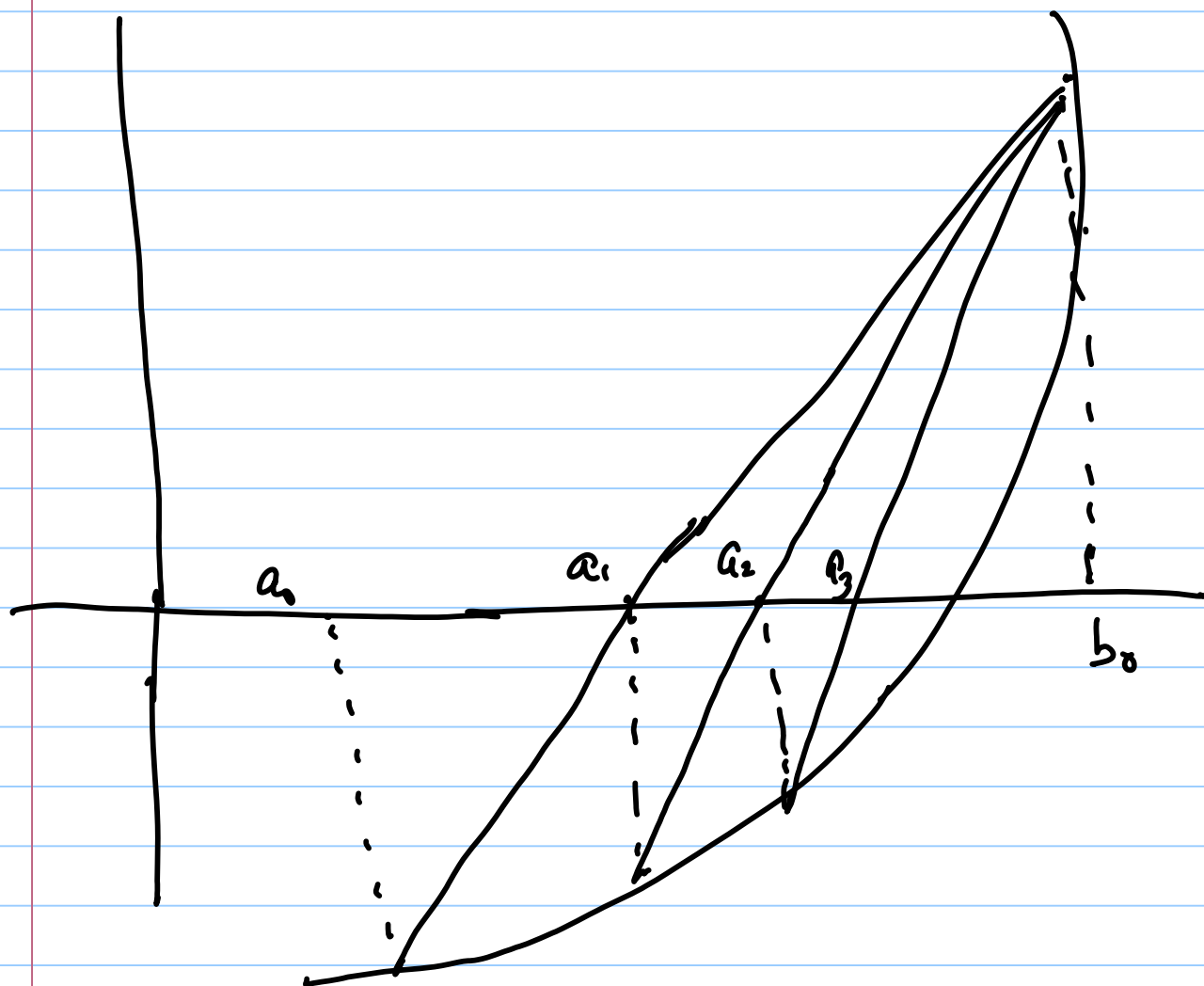
Example

$$f(x) = x^2 - 2$$

4 sig digits

n	a_n	b_n	$f(a_n)$	$f(b_n)$	w_n
0	1	2	-1	2	1.3333
1	1.3333	2	$-2.222E-1$	2	1.4000
2	1.4000	2	$-4.086E-2$	2	1.412
3	1.412	2	$-6.256E-2$	2	1.414
4	1.414	2	- -	- -	

Regula falsi method produces a point at which $|f(x)|$ is "small" somewhere faster than the bisection method, it fails completely to give a "small interval" where the root is known to lie.



$$w = \frac{f(b_n)a_n - f(a_n)b_n}{f(b_n) - f(a_n)}$$

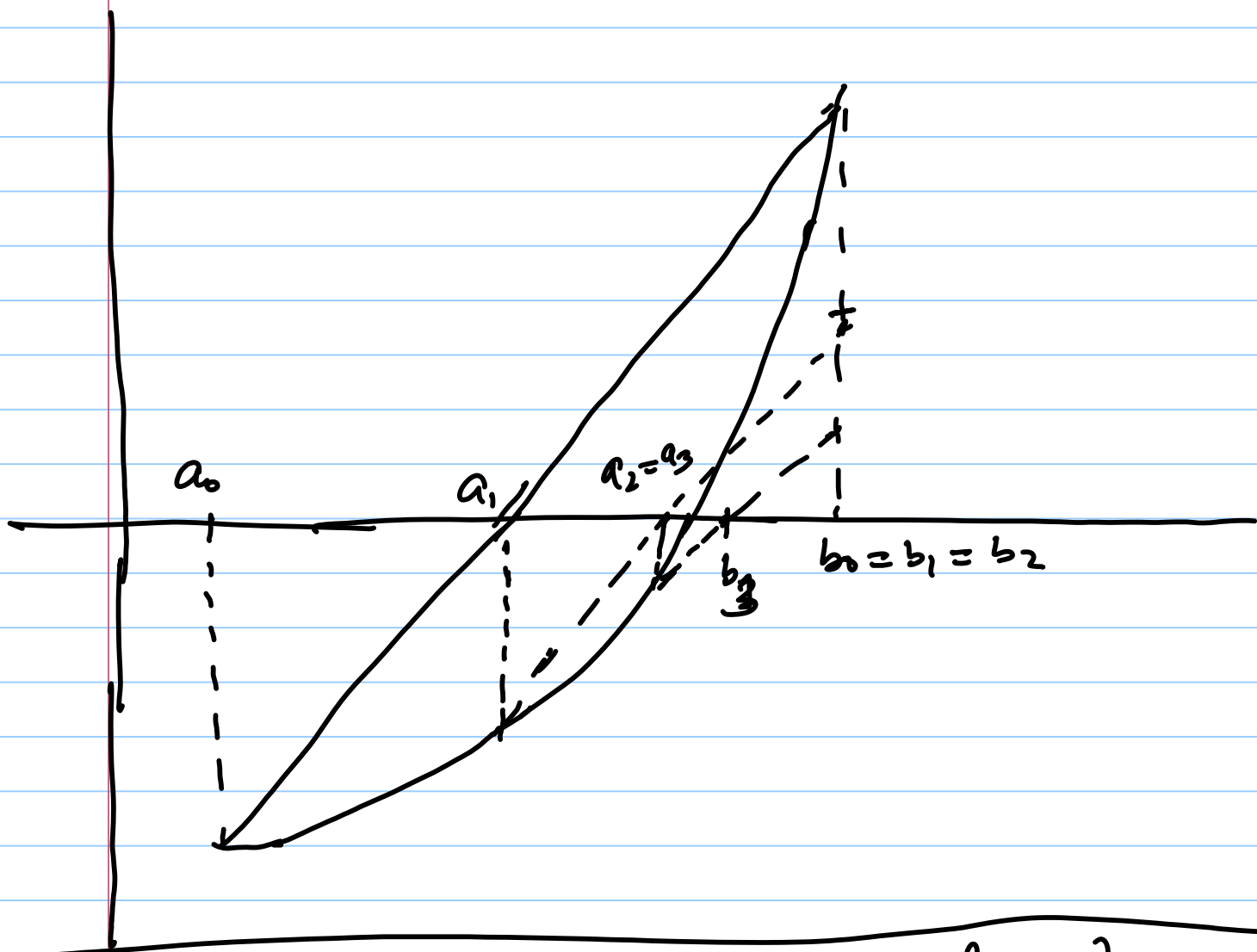
is the pt at which the st-line passing through $(a_n, f(a_n))$, $(b_n, f(b_n))$ intersects the x-axis

Such a st-line is secant to $f(x)$ and in our example $f(x)$ is concave upward and increasing (in the interval $[1, 2]$ of interest), hence the secant is always above the graph of $f(x)$. Consequently w always lies to the left of the zero (in our example).

If $f(x)$ were concave downward and increasing w would always lie to the right of the zero.

Two ways of improving regula-falsi method.

The first one, called "modified regula-falsi" replaces secants by st-lines of ever-smaller slope until w falls to the opp side of the root.



Algorithm (Modified regula falsi)

Given $f(x)$ continuous on $[a_0, b_0]$
and such that $f(a_0)f(b_0) < 0$

Set $F = f(a_0)$ $G = f(b_0)$ $w_0 = a_0$

For $n = 0, 1, 2, \dots$ until satisfied do

Calculate $w_{n+1} = \frac{G a_n - F b_n}{G - F}$

If $f(a_n) f(w_{n+1}) \leq 0$ set $a_{n+1} = a_n$
 $b_{n+1} = w_{n+1}$
 $G = f(w_{n+1})$

If also $f(w_n) f(w_{n+1}) > 0$ set $F = F/2$.

Otherwise set $a_{n+1} = w_{n+1}$, $F = f(w_{n+1})$
 $b_{n+1} = b_n$

If also $f(w_n) f(w_{n+1}) > 0$ set $G = G/2$

Then $f(x)$ has a zero in the interval
 $[a_{n+1}, b_{n+1}]$

Example

$$f(x) = x^2 - 2$$

$$a_0 = 1 \quad b_0 = 2 \quad w_0 = 1$$

n	a_n	b_n	F	G	w_n
0	1	2	-1	2	1
1	1.333	2	-2.222G	1	1.333
2	1.333	1.454	-2.222E-1	1.141E-1	1.454
3	1.413	1.454	-3.431E-3	1.141E-1	1.413
4	1.414	1.454	-6.040E-4	5.705E-2	1.414
5	1.414	1.454	-6.040E-4	2.853E-2	1.414
6	1.414	1.415	-6.040E-4	1.742E-2	1.415

not lies in $[1.414, 1.415]$

A second, very popular modification of the regula-falsi, called the secant method, retains the use of secants throughout, but may give up the bracketing of the root

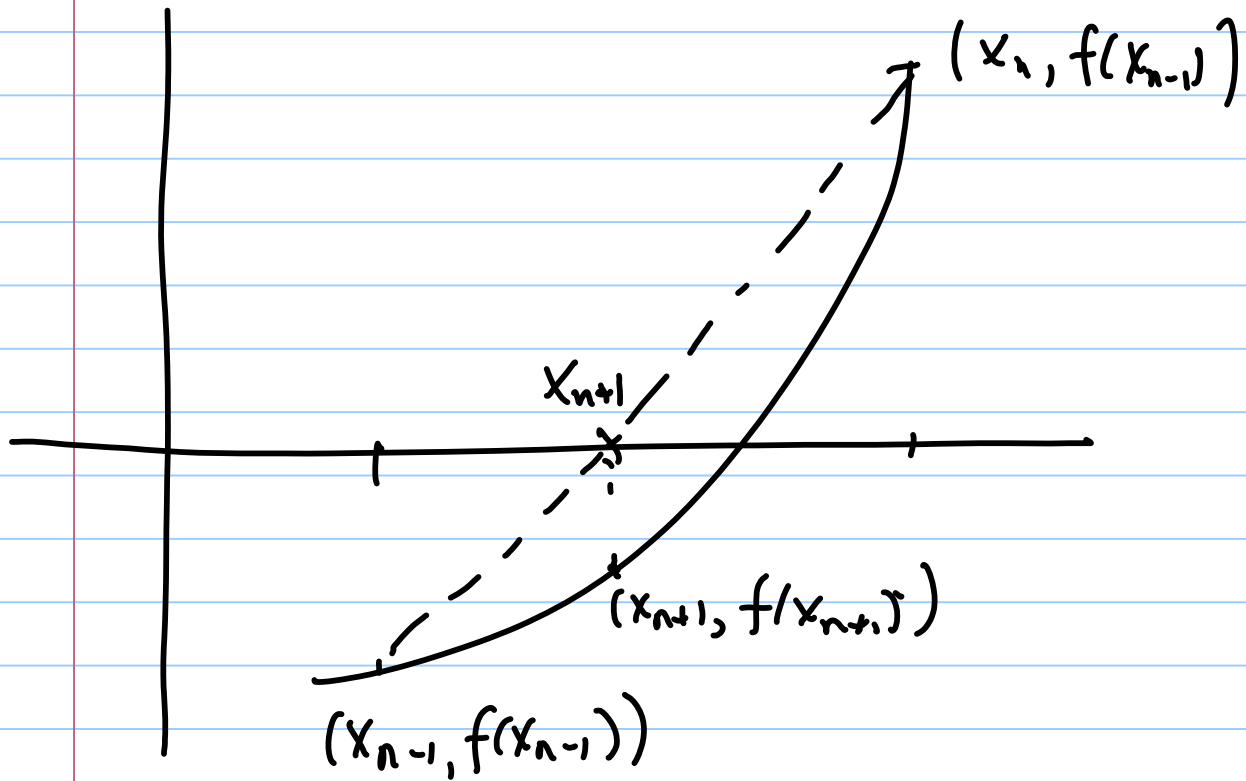
Algorithm (secant method)

Given a function $f(x)$ and two points x_{-1} and x_0

For $n = 0, 1, 2, \dots$ until satisfied do

Calculate

$$x_{n+1} = \frac{f(x_n) x_{n+1} - f(x_{n-1}) x_n}{f(x_n) - f(x_{n-1})}$$



Example

$$f(x) = x^2 - 2$$

$$x_{-1} = 1$$

$$x_0 = 2$$

n	x_n	$f(x_n)$
-1	1	-1
0	2	2
1	1.333	-2.222 E-1
2	1.400	-4.000 E-2
3	1.415	2.225 E-3
4	1.414	-6.040 E-4
5	1.414	-6.040 E-4

process ends

"note algorithm ends if $f(x_n) = f(x_{n-1})$

This makes the calculation of x_{n+1}
impossible

The expression

$$x_{n+1} = \frac{f(x_n) x_{n-1} - f(x_{n-1}) x_n}{f(x_n) - f(x_{n-1})}$$

is prone to round-off error since $f(x_n)$ and $f(x_{n-1})$ need not be of opposite signs.

It is better to calculate x_{n+1} from the equivalent expression

$$x_{n+1} = x_n - f(x_n) \cdot \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})}$$

in which x_{n+1} is obtained from x_n by adding the correction term

$$\frac{-f(x_n)}{f(x_n) - f(x_{n-1})} = \frac{-f(x_n)}{[f(x_n), f(x_{n-1})]}$$
$$\frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})}$$

If x_n close to x_{n-1} then we get-

$$[f(x_n), f(x_{n-1})] \approx f'(x_n)$$

Newton's method

Given $f(x)$ continuously diff and a
pt x_0

For $n = 0, 1, 2, \dots$ until satisfied do

$$\text{Calculate } x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

example

$$f(x) = x^2 - 2$$

$$x_0 = 1$$

$$x_1 = 1.5$$

$$x_2 = 1.416666667$$

$$x_3 = 1.414215686$$

$$x_4 = 1.414213562$$

$$x_n = x_4 \quad \text{for } n \geq 4.$$

Newton's method is special example
of "fixed-pt iteration"

$$g(x) = x - \frac{f(x)}{f'(x)}$$

$$x_{n+1} = g(x_n)$$

$g: I \rightarrow I$
 I closed interval contains root

If the seq $x_1, x_2, \dots \rightarrow \dots$ so general

converges to some pt ξ

and $g(x)$ is cts

$$\text{then } \xi = \lim_{n \rightarrow \infty} x_{n+1}$$

$$= \lim_{n \rightarrow \infty} g(x_n)$$

$$= g(\lim_{n \rightarrow \infty} x_n)$$

$$= g(\xi)$$

So ξ is a fixed pt of g .

