

Question1:

According to the problem, we have a set of N deliveries and their due date (set T). And we can sort the due dates of these deliveries from earliest to latest such that

$$T(n) = \{ T_1 < T_2 < T_3 < T_4 < T_5 \dots T_{(n-1)} < T_n \}$$

Base case:

$N=1$ and since $T \leq N$, we can easily see that there exists two options: 1) $T=0$, this is an undefined behavior, deliveries will never due on day 0. 2) $T=1$, the teleportation machine will deliver the only item on that particular day. Therefore the S is feasible.

Induction Hypothesis:

When $N = K$, S should be feasible for T such that $T \leq K$, the number of deliveries due within T days in set S should be less or equal to T

Inductive Step:

Need to prove that when $N = K+1$, S should be feasible for T such that $T \leq K+1$, the number of deliveries due within T days in set S should be less or equal to T . Since we have sorted the $T(n)$ array in an increasing order, and the first K items would be feasibly delivered according to the inductive hypothesis, we would have the $(K+1)$ th element left to deliver. And because the first K th items are feasibly delivered within $T \leq K$ days, so for the $(K+1)$ th element, it would still have at least $K+1 - K = 1$ days to deliver. Therefore, we would guarantee the $(K+1)$ th element to be delivered within its due date. Therefore, set S is feasible if and only if for all days $T \leq N$.

Question2:

Description of the algorithm:

- 1) We need to arrange the delivery date array T in ascending order. (can use Merge sort or Heap sort to achieve optimal complexity) $\rightarrow O(n \log n)$
- 2) Sort payment array P according to array T .
- 3) Find $T[k]$, $T[k+1]$, $T[k+2]$ such that they are the deliveries that due on the same day, and sort the corresponding subarray P_{same} in payment array P in a decreasing order that $P_{same\ a} > P_{same\ b} > P_{same\ c} \dots$
- 4) At this point we have a array T that is in an ascending order, and an array P with any potential "same day subarray P_{same} " sorted to a descending order.
- 5) Now we can start iterating through the delivery due date array T starting from index 0.
- 6) When we encountered a subarray T_{same} such that it contain the same due date X :

- a) If the left neighbor of such T_{same} has a due date = $x-1$, if so, we only include $P_{same}[0]$ of that subarray so we can take the highest payment. Add $P_{same}[0]$ to P_{total} .
- b) If the left neighbor of such T_{same} has a due date = $x-n$, if so, we include $P_{same}[0...n]$. For example: $T = [1,2,5,5,5,5,6,7,8]$, in this case, we can include the payment for the first 3 items that are due on day 5. Add $P_{same}[0...n]$ to P_{total} .
- c) Print out P_{total} .

Running Time Analysis:

Using merge/heap sort to sort the original T into the ascending form, and sorting the payment array P into the corresponding order would take $O(n \log n)$. Sorting the sub-array P_{same} of payment array P would take, in the worst case, $O(n \log n)$, the other operations would generally take some constant time. The overall running time would be $O(n \log n)$

Proof of correctness:

Base case: $N = 1$, we only have one item to deliver in one day (refer to question 1), and the max payment is fixed \rightarrow true!

Inductive Hypothesis: $N = K$, the algo would produce the correct max payment amount.