## Review the rest

# Selection Based on Scaling and Ranking Mechanisms

- Static Scaling Mechanisms
  - linear scaling

$$S(x) = ax + b$$

$$tf(x) = S(f(x))$$

Evaluation function: eval(x) = tf(x) = s(f(x)) = af(x) + b

Power law scaling

$$eval(x) = tf(x) = s(f(x)) = (f(x))^{U}$$

Logarithmic scaling

$$tf(x) = s(f(x)) = e^{\frac{f(x)}{T}}$$

T is controlling the selective pressure  $\rightarrow$ >decrease during search process

Selection pressure  $\rightarrow$  degree to which highly fit individuals are allowed to produce offspring (copies) in the next generation

Low values of selection pressure  $\rightarrow$  provide each individual in the population with a reasonable selection prob.

High values of selection pressure → strongly favor the best individuals in the population → population diversity decrease quickly

#### Dynamic Scaling mechanism

• Sigma truncation

$$T(x) = \begin{cases} x & \text{if } x > 0 \\ o & \text{else} \end{cases}$$

$$S(x) = T\left(x - \left(m_t - c\sigma_t\right)\right)$$

$$eval(x) = tf(x) = S(f(x)) = T\left(f(x) - \left(m_t - c\sigma_t\right)\right)$$

$$eval(x) = \begin{cases} f(x) - \left(m_t - c\sigma_t\right) & \text{if } f(x) > \left(m_t - c\sigma_t\right) \\ 0 & \text{if } f(x) \le \left(m_t - c\sigma_t\right) \end{cases}$$

$$OR$$

$$eval(x) = \begin{cases} 1 + \frac{f(x) - m_t}{\sigma_t} & \text{if } \sigma_t \neq 0 \\ 0 & \text{if } \sigma_t = 0 \end{cases}$$

Window scaling

$$eval(x) = f(x) - \min_{y \in P(t)} f(y)$$

$$OR$$

$$eval(x) = f(x) - \min_{y \in W} f(y)$$

where W is the set of all individuals contained in the last g generation,  $g \ge 1$ 

Fitness remapping for minimization problem

tf(x) = M - f(x) where M is the upperbound of the objective function

 $eval(x(t)) = tf(x(t)) = M_t - f(x(t))$  where  $x(t) \in P(t)$  and  $M_t$  is max value of function f found so far or max value of f found within generation P(t)

$$\underline{OR} \quad eval(x(t)) = \frac{1}{K + f(x(t)) - f_{\min}(t)} \text{ where } K \text{ is positive constant and } f_{\min}(t)$$

is the minimum observed value of function f up to time t

$$\underline{\mathsf{OR}} \qquad eval(x) = \frac{1}{K + f(x)}$$

Rank-based Selection

let number of individuals in each generation  $\rightarrow N$  (|P(t)|=N)

• Linear Ranking selection selection probability of  $x^i$ 

$$p_{i} = \frac{1}{N} \left[ Min + (Max - Min) \frac{r_{i} - 1}{N - 1} \right]$$

where  $r_i \rightarrow$  rank of individual  $x^i$  (ranked in increasing order) expected value of offspring of  $x^i$ 

$$n_{j} = Np_{j}$$
And 
$$\sum_{i=1}^{N} n_{i} = \sum_{i=1}^{N} Np_{i} = N \text{ Then } \sum_{i=1}^{N} Np_{i} = N$$

$$\sum_{i=1}^{N} N \frac{1}{N} \left[ Min + (Max - Min) \frac{r_{i} - 1}{N - 1} \right] = \sum_{i=1}^{N} Min + \sum_{i=1}^{N} (Max - Min) \frac{r_{i} - 1}{N - 1} = N$$

$$NMin + \frac{1}{N - 1} (Max - Min) \sum_{i=1}^{N} (r_{i} - 1) = N$$

But, since 
$$\sum_{i=1}^{N} (r_i - 1) = \sum_{i=1}^{N} (i - 1) = \sum_{i=0}^{N-1} i = \frac{N(N-1)}{2}$$

Hence, 
$$Min + \frac{1}{2}(Max - Min) = 1$$

$$Min = 2 - Max$$
 OR  $Max + Min = 2$ 

Since  $Min \ge 0$  and  $Min \le Max \longrightarrow Max \le 2$  and  $Max \ge 1 \longrightarrow 1 \le Max \le 2$ 

Hence 
$$p_N = \frac{Max}{N}$$
 and  $p_1 = \frac{2 - Max}{N} = \frac{Min}{N}$ 

Alternatively, 
$$p_i = \frac{1}{N} \left[ Max + 2(Max - 1) \left( 1 - \frac{r_i - 1}{N - 1} \right) \right]$$

## nonlinear ranking selection Selection probability will be

$$p_i = \frac{1}{c} \left( 1 - e^{1 - r_i} \right)$$

## • SUS Algorithm

- 1.  $ptr = Rand(); /*random within [0,1] \rightarrow (uniform distribution)*/$
- 2. For (sum=i=0; i<N;i++)
  - 1. For  $(sum +=n_i; sum>ptr; ptr++)$
  - Select(i);
  - 3. End For
- 3. End For

guarantee that  $x^i$  will reproduce at least  $\lfloor n_i \rfloor$  but not more than  $\lceil n_i \rceil$ 

## • Example

Chomosome	f(x)=x <sup>2</sup>	r <sub>1</sub>	p <sub>i</sub>	n <sub>i</sub>
x <sup>1</sup> =10	100	1	$\frac{2-1.2}{3} = 0.27$	3(0.27)=0.81
x <sup>2</sup> =15	225	2	$\frac{1}{3} \left[ 0.8 + \left( 1.2 - 0.8 \right) \frac{2 - 1}{3 - 1} \right] = 0.33$	3(0.33)=0.99
x <sup>3</sup> =20	400	3	0.4	3(0.4)=1.2

Suppose ptr is  $0.8 \rightarrow$  the selection will be  $x^1, x^3$  and  $x^3$ , respectively

#### Tournament selection

Similar to rank selection in terms of selection pressure  $\rightarrow$  but more efficient

- Binary tournament (deterministic)
  - 2 individuals directly compete for selection
  - With or without reinsertion of the competing individual into the original population
  - based on direct pairwise comparison of individuals
    - 2 chromosomes are chosen at random from the population
    - The fitness of the chosen chromosomes is calculated
    - The most successful chromosome (the winner) is selected → to intermediate population

#### Probability tournament

• Step 3  $\rightarrow$  random  $R \in [0,1]$  if R < p then the fitter of the two is selected; otherwise the less fit is selected

where p is preselected or p varied with time  $\rightarrow$ 

$$p(t) = p_0 e^{-ct}$$
 or  $p(t) = p_0 \frac{c}{1 + \ln t}$  where  $p_0 \in (0,1)$  and  $c$  is positive constant

Boltzmann tournament

 $x \rightarrow$  current solution,  $y \rightarrow$  alternative solution  $\rightarrow$  binary tournament between x and y with

$$p(x) = \frac{1}{\underbrace{f(x) - f(y)}_{T}}$$

$$1 + e^{-T}$$

Where T is the temperature  $\rightarrow$  reduced in a predefined manner in successive iteration

The winner of the trial will be the current solution in the next trial

Elitism 

 selection of a set of individual from the current generation to survive to the next generation

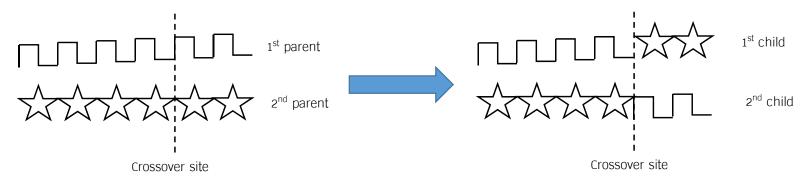
The fraction of new individual at each generation has been called generation gap

Recombination operator

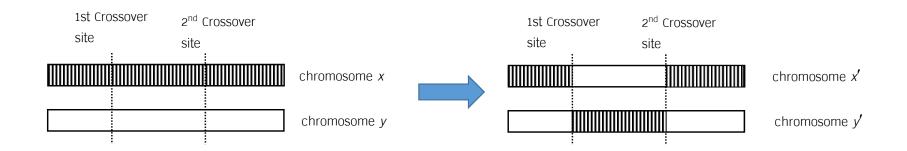
cross over:  $C: X \times X \to X \times X \to C(x,y) = (x',y')$ 

• One-point crossover

$$x = x_1 x_2 x_3 ... x_L$$
 and  $y = y_1 y_2 y_3 ... y_L$  after cross over  
 $x' = x_1 x_2 x_3 ... x_{k-1} x_k y_{k+1} ... y_L$  and  $y' = y_1 y_2 y_3 ... y_{k-1} y_k x_{k+1} ... x_L$ 



#### • Two-point cross over



#### • *N*-point cross over



- N-point cross over algorithm
- P1: for each chromosome from intermediate population  $(P^1)$ 
  - P1.1: random  $q \in [0,1] \rightarrow$  uniform distribution
- P1.2: if  $q < p_c$  (where  $p_c$  is mating probability) then respective chromosome is considered mating  $\rightarrow$ go to mating pool else not mating
- P2: let  $m^* \rightarrow$  number of chromosome selected at P1 if  $m^*$  is even then select pair chromosome (random) if  $m^*$  is odd then discard or select chromosome from  $P^1$
- P3: Cross over is applied on pairs in P2
  - P3.1: for each pair, generate k ( $\leq k \leq L$ ) and random crossing site  $k_1, k_2, ..., k_N$  for ( $N \geq 2$ )
- P3.2: descendants obtained at P3.1 become potential members of  $P(t+1) \rightarrow$  candidate for mutation
  - P3.3: parents of newly generated chromosome are discarded from  $P^1$
  - P3.4: chromosomes left in  $P^1$  are added to P(t+1)

#### Punctuated crossover

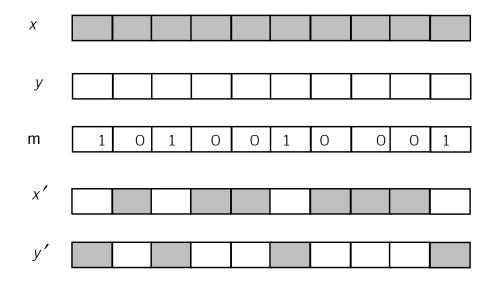
- Record the crossover points in the chromosome
- If a certain crossover point has generated a very poor offspring this crossover point will not be considered further
- If the fitness of the descendant is good, the crossover point is maintained as active
- Cross-over point self adjust according to the previous dynamics of the search process  $C = x_1 x_1 \dots x_l k_1 k_1 \dots k_l$  when  $k_i = 1 \rightarrow$  active crossover point,  $k_i = 0 \rightarrow$  not a crossover point

During initialization probability of generating  $k_i = 1$  is small

Set of crossover point used in the recombination of 2 chromosomes is the union of the parents

#### • Uniform crossover

- 1. Set  $m_i = 0$  for i = 1, 2, ..., L
- 2. For each i
  - 1.  $random \xi \sim U(0,1)$
  - 2. If  $\xi \leq p_{at\_i}$  then  $m_i = 1$  where  $p_{at\_i}$  is the crossover probability at  $i^{th}$  position



#### Mutation

Average number of position that will undergo a mutation  $\rightarrow B = NLp_m$  where  $p_m$   $\rightarrow$  mutation probability ( $p_m \in [0.001, 0.01]$ ),  $N \rightarrow$  number of chromosomes,  $L \rightarrow$  length of chromosome

#### Algo

- 1. For each chromosome une For each position of the chromosome
  - 1. Generate a random number q from U(0,1)
  - 2. If  $q < p_m$  invert that position If  $q \ge p_m$  do not mutate

Weak mutation step 2 changed to if  $q < p_m$  then 1 of 0 or 1 is chosen at random

Non-uniform mutation

$$p_m(t) = p_m e^{-\beta t} \quad \text{where } \beta \ge 1$$

<u>OR</u>

$$p_m(t) = \left(\frac{\alpha}{\delta}\right)^{\frac{1}{2}} \left(\frac{\frac{-\beta t}{e^{\frac{2}{2}}}}{\frac{1}{L^2 N}}\right) \quad \text{where } \alpha, \beta, \delta \rightarrow \text{Real-valued positive}$$

parameter

<u>OR</u>

$$p_m(t) = \frac{1}{2 + \frac{L - 2}{\tau}t}$$
 where  $0 \le t \le T \to p_m(0) = 0.5$  and  $p_m(\tau) = \frac{1}{L}$ 

<u>OR</u>

$$p_m(x) = \frac{1}{2(f(x)+1)-L}$$

## Inversion operator

Acts on a single chromosome. It inverts the values between 2 position of a chromosome. The 2 positions are chosen at random

### Simple genetic algorithm

1. Set t = 0

Initialize (P(0))  $\rightarrow$  random

a selection mechanism is chosen and if necessary a scaling mechanism

- 2. The chromosome of P(t) are evaluated using the fitness function. The most successful individual from P(t) is kept.
- 3. Selection operator is applied N times

Selected chromosome make up an intermediate population  $P^1$  (having N members as well) representing candidate from the mating pool (MP). Some chromosomes from P(t) may have more copies in  $P^1$  while some others may have none

4. Crossover operator is applied to the chromosomes from MP

Newly generated chromosomes forn P<sup>2</sup>

Parents of the chromosomes obtained through crossover are discarded from P<sup>1</sup>

Chromosomes that remained in  $P^1$  become member of  $P^2$ 

- 5. Mutation operator is applied to  $P^2$  outcome is the new generation(P(t+1))
- 6. set t = t+1

if t < T (T is the maximum number of generation) then go to step 2 otherwise stop

- Solution 
   best member of the last population or best individual in all generation
- Can involve inversion or other types of operators

## • Example

$$\max_{x_1, x_2} f(x_1, x_2)$$
 where  $f(x_1, x_2) = \frac{1}{1 + x_1^2 + x_2^2}$ 

chromos ome	<b>x</b> <sub>1</sub>	X <sub>2</sub>	Fitness value
$C_0$	-1	2	0.167
C <sub>1</sub>	-2	3	0.007
C <sub>2</sub>	1.5	0.0	0.31
C <sub>3</sub>	0.5	-1.0	0.44

 $(C_3, C_2)$  and  $(C_3, C_0)$   $(C_3, C_2)$  and  $(C_3, C_0)$ 

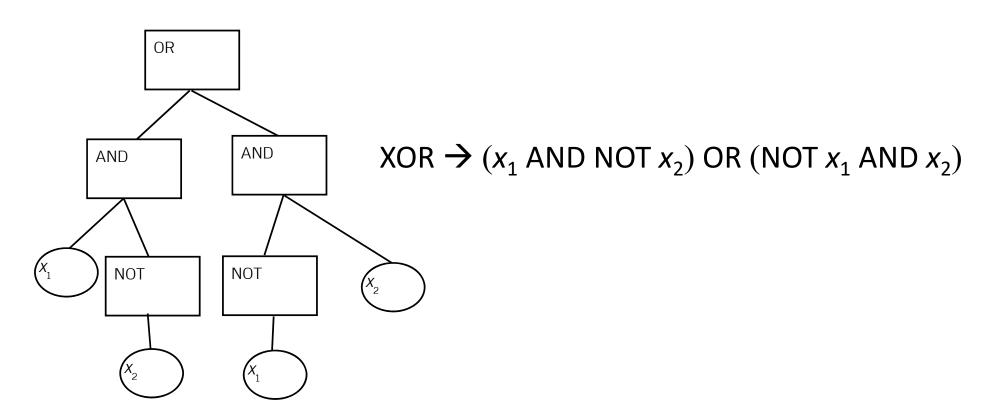
chromos ome	<b>X</b> <sub>1</sub>	X <sub>2</sub>	Fitness value
C <sub>0</sub>	0.5	0	0.8
<b>C</b> <sub>1</sub>	1.5	-1.0	0.24
C <sub>2</sub>	0.5	2.0	0.19
C <sup>3</sup>	-1.0	-1.0	0.33

## Genetic Programming

- Each individual/chromosome → 1 computer program represented using a tree structure
- Terminal set → specifies all variables and constant → elements of the terminal set form the leaf modes
- Function set → all functions that can be applied to the elements of the terminal set may be mathematical arithmetic and/or Boolean function or if-then-else → form the non-leaf nodes.

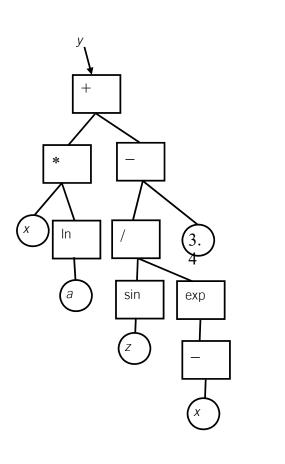
## Example

function {AND, OR, NOT} and terminal set  $\{x_1, x_2\}$  where  $x_1, x_2 \in \{0,1\}$ 



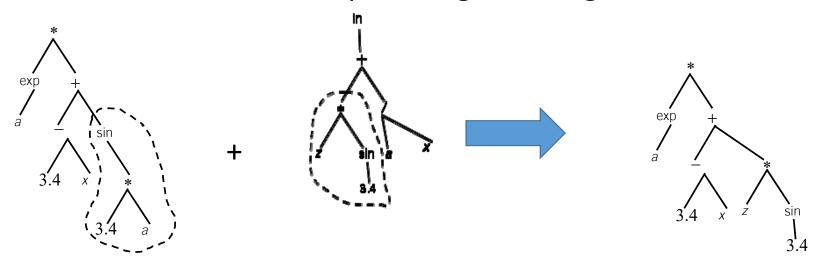
## • Example

Function  $\{-,+,*,/\}$  and terminal set  $\{a,x,z,3.4\}$  where  $a,x,z \in \Re$ 

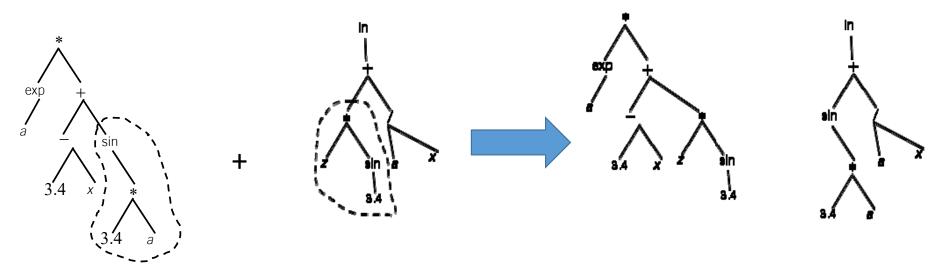


$$y = x * \ln(a) + \sin(z) / \exp(-x) - 3.4$$

## Crossover → 2 parents generating 1 child



## Crossover → 2 parents generating 2 children



- Mutation 

  choose a random point in a tree and replacing the subtree beneath that point by a randomly generated subtree
- Calculate the fitness of each program → by running it on a set of "fitness case"

## Evolutionary programming

Algo

```
1. Let current generation be t=0
2. Initialize all parameters
    Initialize population C(0) with n_s individual
    For each individual \mathbf{x}(t) \in C(t)
         calculate fitness value f(x,(t))
      End
     While no convergence
      • For each individual \mathbf{x}(t) \in C(t)
               mutate \mathbf{x}_{i}(t) to produce offspring \mathbf{x}_{i}'(t)
              calculate fitness value (f(x_i(t)))
              add \mathbf{x}_{i}'(t) in offspring set C'(t)
           End
          Select new population C(t+1) from C(t) \cup C'(t) using selection operator
          t = t + 1
      End
```

## Coevolution

- Predator-pray competitive coevolution →
  - Insects that eat plant
    - To survive plant needs to evolve mechanism to defend itself from the insects, and
    - The insects need the plants as food source to survive.
  - Inverse fitness interaction between 2 species → win for 1 species → failure for the other
- Symbiosis
  - Different species cooperate instead of compete → success in 1 species improve the survival strength of the other species → positive feedback among the species that take part in this cooperating process

- Coevolutonary algorithm (CoEA)
  - Competitive coevolution
    - Competition
    - Amensalism
  - Cooperative coevolution
    - Mutualism
    - Commensalism
    - Parasitism
- Competitive coevolution (CCE)
  - Evolve 2 population simultaneously
    - individual in 1 population → solution (evolve to solve as many test cases as possible) →
      fitness ∞ the number of test cases solved by the solution
    - individual in another population → test case (evolve to present an incrementally increasing level of difficult to the solution individuals) → fitness is inversely proportional to the number of strategies that solve it

- - Which individual from the competing population is used
  - Exactly how these competing individuals are used to compute the relative fitness
- Sampling scheme
  - All versub all sampling
  - Random sampling
  - Tournament sampling
  - All versus best sampling
  - Shared sampling
- Relative fitness function  $\rightarrow f(C_1.x_i)$ 
  - Simple function

2 population  $\rightarrow C_1$  and  $C_2 \rightarrow$  aim to calculate the relative fitness of each individual  $C_1.x_i$  of  $C_1$ 

Sample of individuals is taken from  $C_2 \rightarrow S$ 

sum  $C_1.x_i$  = count the number of individuals in S that  $C_1.x_i$  win  $f(C_1.x_i)$  = sum  $C_1.x_i$ 

Fitness sharing

sim  $C_1.x_i$  = similarities of  $C_1.x_i$  with all the other individuals in that population  $\rightarrow$  e.g., the number of individuals that also beats individual from  $C_2$  samples

$$f\left(C_{1}.x_{i}\right) = \frac{sumC_{1}.x_{i}}{simC_{1}.x_{i}}$$

Competitive fitness sharing

$$f(C_1.x_i) = \sum_{l=1}^{C_2.n_s} \frac{1}{C_1.n_l}$$

Where  $C_2.x_1$ ,  $C_2.x_2$ ,...,  $C_2.x_{C2.ns}$  form  $C_2$  and  $C_1.n_I$  = total number of individuals in  $C_1$  that defeat individual  $C_2.x_I$ 

- Tournament fitness
  - Binary tournament determine a relative fitness ranking
  - Tournament fitness results in a tournament tree with root element as the best individual
    - Each level → 2 opponent are randomly selected from that level and the best of the 2 advances to next level
  - If odd number → a single individual from the level moves to the next level
  - After tournament → use any selection operator

#### Hall of frame

- At each generation the best individual of a population is stored in that population's hall of frames
- May have limited size
- New individual → inserted in hall of frame will replace the worst or the oldest individual
- Individual from 1 population → complete against a sample of the current opponent population and its hall of frame
- Prevents overspecialization

- Competitive coevolution with 2 population Algo
  - Initialize  $C_1$  (solution),  $C_2$  (test) population
  - While stopping condition(s) not true do
    - For each  $C_1.x_i$  for  $i=1,..., C_1.n_s$  do
      - Select a sample of opponents from C<sub>2</sub>
      - Evaluate the relative fitness of  $C_1.x_i$  w.r.t. this sample
    - End
    - For each  $C_2.x_i$  for  $i=1,..., C_2.n_s$  do
      - Select a sample of opponents from C<sub>1</sub>
      - Evaluate the relative fitness of  $C_2.x_i$  w.r.t. this sample
    - End
    - Evolve population  $C_1$  for 1 generation
    - Evolve population C<sub>2</sub> for 1 generation
  - end
  - Select the best individual from  $C_1$  as solution

- Competitive coevolution with 1 population Algo
  - Initialize 1 population *C*
  - While stopping condition(s) not true do
    - For each  $C.x_i$  for  $i=1,..., C.n_s$  do
      - Select a sample of opponents from C (exclude C.x<sub>i</sub>)
      - Evaluate the relative fitness of *C.x<sub>i</sub>* w.r.t. the sample
    - End
    - Evolve population *C* for 1 generation
  - end
  - Select the best individual from C as solution

### Cooperative coevolution

- Similar to divide and conquer
- Only mutualism here
- Individual from different species or subpopulation have to cooperate in some way to solve a global task
- Fitness depends on that individual's ability to collaborate with individual from other species
- Ex. For an  $n_x$  dimensional problem  $\rightarrow n_x$  subpopulation (1 for each)
  - each subpopulation → responsible for optimizing one of the parameters of problem and no subpopulation can form a complete solution by itself
  - Merge representative from each subpopulation
  - Consider j subpopulation  $(C_j) \rightarrow$  each  $C_j \cdot x_i$  perform a single collaboration with the best individual from each other subpopulation  $\rightarrow$  complete solution
  - Credit assign to  $C_{i}$ . $x_{i}$   $\rightarrow$  fitness of the complete solution
  - Evolve independently from one another → separate population is used to evolve each subcomponent using some EA

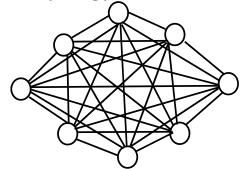
# Swarm Intelligence (SI)

- Structured collection of interacting organisms (or agents)
- Individual organism
  - Simple in structure but their collective behavior can become complex
- Tight coupling between individual behavior and the behabior of the entire swarm
- Social interaction > direct (through visual audio, chemical context), or indirect (occur when one changes the environment and the other individuals respond to the new environment)
- 2 concepts
  - Self-organization
    - Positive feedback

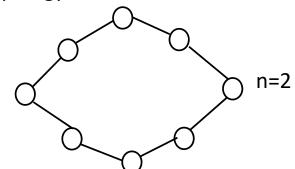
## Swarm Intelligence

- 2 concepts
  - Self-organization
    - Positive feedback
    - Negative feedback
    - Fluctuation
    - Multiple interaction
  - Division of labor
- SI→no one is in charge, no one is giving orders, which the swarm individuals should obey or follow
- Social networkstructureof swarm → form an integral part of the existence of that swarm → provide the communication channel, able to self-organize to form the optimal nest structures

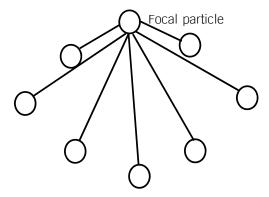
- Partical swarm optimization (PSO)
  - Simulation of social behavior of birds within flock
  - Particles are "flown" through hyperdimensional search space
  - Neighborhood topology
    - Star topology → first version PSO → gbest



Ring topology



## Wheel topology



- Individual best (pbest)
  - Initialize the swarm P(t) at t=0 of particles such that the position  $(\mathbf{x}_i(t))$  of  $i^{th}$  particle  $(P_i \in P(t))$  is random within the hyperspace
  - Evaluate the performance F of each particle using  $\mathbf{x}_i(t)$
  - Compare the performance of each individual to its best performance

• If 
$$F(\mathbf{x}_i(t)) < pbest_i$$
  

$$pbest_i = F(\mathbf{x}_i(t))$$

$$\mathbf{x}_{pbest_i}(t) = \mathbf{x}_i(t)$$

- End
- Change the velocity vector of each particle

$$v_i(t) = v_i(t-1) + \rho(x_{pbest_i} - x_i(t))$$
 where  $\rho$  is positive random number

Move each particle to a new position

$$\mathbf{x}_{i}(t) = \mathbf{x}_{i}(t-1) + \mathbf{v}_{i}(t)$$

- Set *t*=*t*+1
- Repeat until converge (go to step 2)

### Global best (gbest)

- Initialize the swarm P(t) at t=0 of particles such that the position  $(\mathbf{x}_i(t))$  of  $i^{th}$  particle  $(P_i \in P(t))$  is random within the hyperspace
- Evaluate the performance F of each particle using  $\mathbf{x}_i(t)$
- Compare the performance of each individual to its best performance
  - If  $F(\mathbf{x}_i(t)) < pbest_i$  $pbest_i = F(\mathbf{x}_i(t))$   $\mathbf{x}_{pbest_i}(t) = \mathbf{x}_i(t)$
  - End
- Compare the performance of each individual to its best performance
  - If  $F(\mathbf{x}_i(t)) < gbest$   $gbest = F(\mathbf{x}_i(t))$  $x_{ghest}(t) = x_i(t)$
- Change the velocity vector of each particle

$$\mathbf{v}_{i}(t) = \mathbf{v}_{i}(t-1) + \rho_{1}\left(\mathbf{x}_{pbest_{i}} - \mathbf{x}_{i}(t)\right) + \rho_{2}\left(\mathbf{x}_{gbest} - \mathbf{x}_{i}(t)\right)$$

• Move each particle to a new position

$$\mathbf{x}_{i}(t) = \mathbf{x}_{i}(t-1) + \mathbf{v}_{i}(t)$$

- Set *t*=*t*+1
- Repeat until converge (go to step 2)

Where  $\rho_1$  and  $\rho_2$  are random number  $\Rightarrow \rho_1 = r_1 C_1$  and  $\rho_2 = r_2 C_2$  with  $r_1$  and  $r_2 \sim U(0,1)$ ,  $c_1 + c_2 \leq 4$ 

#### Local best (lbest)

- Initialize the swarm P(t) at t=0 of particles such that the position  $(\mathbf{x}_i(t))$  of  $i^{th}$  particle  $(P_i \in P(t))$  is random within the hyperspace
- Evaluate the performance F of each particle using  $\mathbf{x}_i(t)$
- Compare the performance of each individual to its best performance
  - If  $F(\mathbf{x}_i(t)) < pbest_i$  $pbest_i = F(\mathbf{x}_i(t))$   $\mathbf{x}_{pbest_i}(t) = \mathbf{x}_i(t)$
  - End
- Compare the performance of each individual to its best performance
  - If  $F(\mathbf{x}_i(t)) < lbest$   $lbest = F(\mathbf{x}_i(t))$  $x_{lhest}(t) = x_i(t)$
- · Change the velocity vector of each particle

$$\mathbf{v}_{i}(t) = \mathbf{v}_{i}(t-1) + \rho_{1}\left(\mathbf{x}_{pbest_{i}} - \mathbf{x}_{i}(t)\right) + \rho_{2}\left(\mathbf{x}_{gbest} - \mathbf{x}_{i}(t)\right)$$

• Move each particle to a new position

$$\mathbf{x}_{i}(t) = \mathbf{x}_{i}(t-1) + \mathbf{v}_{i}(t)$$

- Set *t*=*t*+1
- Repeat until converge (go to step 2)

Where  $\rho_1$  and  $\rho_2$  are random number  $\Rightarrow \rho_1 = r_1 C_1$  and  $\rho_2 = r_2 C_2$  with  $r_1$  and  $r_2 \sim U(0,1)$ ,  $c_1 + c_2 \leq 4$ 

- Fitness calculation  $\rightarrow$  can be objective function  $\rightarrow$  same idea as in GA
- Convergence  $\rightarrow$  executed for a fixed number of iteration or velocity changes are close to 0 for all particles
- PSO system parameter
  - Dimension of problem
  - Number of individual
  - Upper limit of  $\rho$
  - Upper limit on the maximum velocity  $(v_{max})$ If  $v_{ij}(t) > v_{max}$  then  $v_{ij}(t) = v_{max}$  and If  $v_{ij}(t) < -v_{max}$  then  $v_{ij}(t) = -v_{max}$ where  $v_{ij}(t)$  is a velocity of ith particle at jth dimension at time tOR if do not want to use  $v_{max}$

$$\mathbf{v}_{i}(t) = \kappa \left( \mathbf{v}_{i}(t-1) + \rho_{1} \left( \mathbf{x}_{pbest_{i}} - \mathbf{x}_{i}(t) \right) + \rho_{2} \left( \mathbf{x}_{gbest} - \mathbf{x}_{i}(t) \right) \right)$$
Where

 $\kappa = 1 - \frac{1}{\rho} + \frac{\sqrt{\rho^2 - 4\rho}}{2}$ 

and 
$$\rho = \rho_1 + \rho_2 > 4.0$$

- Neighborhood size
  - gbest → simply lbest with the entire swarm as the neighborhood
  - gbest → more susceptible to local minima
- Inertia weight

$$\mathbf{v}_{i}(t) = \phi \mathbf{v}_{i}(t-1) + \rho_{1} \left( \mathbf{x}_{pbest_{i}} - \mathbf{x}_{i}(t) \right) + \rho_{2} \left( \mathbf{x}_{gbest} - \mathbf{x}_{i}(t) \right)$$

Where  $\phi$  is inertia weight ( $\phi \le 1$ ) and

$$\phi > \frac{1}{2}(c_1 + c_2) - 1$$

- Modified PSO
  - Using selection → perform selection before velocity update
     Algo
    - for each particle in the swarm, score the performance of that particle w.r.t a random selected group of k particles
    - Rank the particles according to these performance scores
    - Select the top half of the particle and copy their current position onto that of the bottom half of the swarm without changing the personal best values of the bottom half of the swarm

#### Breeding PSO

#### Algo

- Calculate the particle velocities and new positions
- For each particle assign a breeding probability  $p_b$
- Select 2 particles assume  $P_a$  and  $P_b$  and produce 2 offsprings using

$$\mathbf{x}_{a}(t+1) = r_{1}\mathbf{x}_{a}(t) + (1 - r_{1})\mathbf{x}_{b}(t)$$

$$\mathbf{x}_{b}(t+1) = r_{2}\mathbf{x}_{b}(t) + (1 - r_{2})\mathbf{x}_{a}(t)$$

$$\mathbf{v}_{a}(t+1) = \frac{\mathbf{v}_{a}(t) + \mathbf{v}_{b}(t)}{\|\mathbf{v}_{a}(t) + \mathbf{v}_{b}(t)\|} \|\mathbf{v}_{a}(t)\|$$

$$\mathbf{v}_{b}(t+1) = \frac{\mathbf{v}_{a}(t) + \mathbf{v}_{b}(t)}{\|\mathbf{v}_{a}(t) + \mathbf{v}_{b}(t)\|} \|\mathbf{v}_{b}(t)\|$$

Where  $r_1, r_2 \sim U(0,1)$ 

 Set personal best position of each particle involved in the breeding process to its current position

- Neighborhood topology
  - Use of indices
  - Use spatial neighbor
     P<sub>b</sub> is said to be in the neighborhood of P<sub>a</sub> if

$$\frac{\left\|\mathbf{x}_{a} - \mathbf{x}_{b}\right\|}{d_{\max}} < \xi$$

where  $d_{max}$  is the largest distance between any 2 particles and

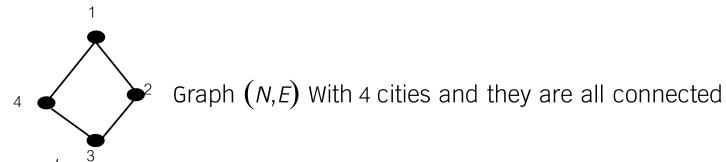
$$\xi = \frac{3t + 0.6t_{\text{max}}}{t_{\text{max}}}$$

## Ant Colony Optimization

- tasks
  - Reproduction
  - Defense
  - Foor collection
  - Brood care
  - Nest brooming
  - Nest building
- Natural stigmergy
  - Lack of central coordination
  - Communication and coordination among individuals in a colony are based on local modifications of the environmeny
  - Positive feedback

- Shortest path problem → traveling salesman problem (TSP)
- Virtual pheromone
- Negative feedback  $\rightarrow$  avoid premature feedback
- Time scale → not too large and not too short
- Ant System (AS)
  - Find a closed tour of minimal length connected *n* given cities

$$d_{ij} = \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2}$$
 distance between city *i* and city *j*



• Tabu list and  $J_i^k$  set of all cities that ant k is has to visit when it is on city i

#### AS algo

$$p_{ij}^{k}(t) = \begin{cases} \frac{\left[\tau_{ij}(t)\right]^{\alpha} \left[\eta_{ij}\right]^{\beta}}{\sum\limits_{l \in J_{i}^{k}} \left[\tau_{il}(t)\right]^{\alpha} \left[\eta_{il}\right]^{\beta}} & \text{if } j \in J_{i}^{k} \\ 0 & \text{else} \end{cases}$$
where  $\alpha$  and  $\beta$  are adjustable and visibility  $\eta_{ij} = \frac{1}{d_{ij}}$ 

```
End For

For k= 1 to m

compute length L<sup>k</sup>(t) of tour T<sup>k</sup>(t) produced by ant k

End For

If an improved tour is found then

update T<sup>+</sup> and L<sup>+</sup>

End If
```

For every edge (i,j)

$$\tau_{ij}(t+1) = (1-\rho)\tau_{ij}(t) + \Delta\tau_{ij}(t) + e\Delta\tau_{ij}^{e}(t)$$

where

$$\Delta \tau_{ij}^{e}(t) = \begin{cases} \frac{Q}{L^{+}} & \text{if } (i,j) \in T^{+} \\ 0 & \text{if } (i,j) \notin T^{+} \end{cases}$$

$$\Delta \tau_{ij}(t) = \sum_{k=1}^{m} \Delta \tau_{ij}^{k}(t)$$

$$\Delta \tau_{ij}(t) = \sum_{k=1}^{m} \Delta \tau_{ij}^{k}(t)$$

$$\Delta \tau_{ij}^{k}(t) = \begin{cases} \frac{Q}{L^{k}(t)} & \text{if } (i,j) \in T^{k}(t) \\ 0 & \text{if } (i,j) \notin T^{k}(t) \end{cases}$$

End For End For

- Ant colony system (ACS)
  - Different transition rule
  - Different pheromone trail update rule
  - Use of local updates of pheromone trail to favor exploration
  - Use of candidate list→ cl closest cities ordered by increasing distance→ ant first restricts
    the choice of the next city to those in the list, and consider other cities if all in the list
    have been visited

For every edge (i,j)

$$\tau_{ii}(0) = \tau_{0}$$
 #initialize pheromone

End For

For k=1 to m #for each ant k

Place ant k on a randomly chosen city

End For

Set  $T^+$  and  $L^+$ 

For t=1 to  $t_{max}$ 

For k=1to m

build tour  $T^k(t)$  by applying n-1 times the following steps: if exits at least one city  $j \in C$  and idate list then choose the next city  $t \in C_j^k$  among cl citiesy in the candicate list as

$$j = \begin{cases} \arg\max_{u \in J_i^k} \left\{ \left[ \tau_{iu}(t) \right] \right] \left[ \eta_{iu} \right]^{\beta} \right\} & \text{if } q \le q_0 \\ J & \text{if } q > q_0 \end{cases}$$

where  $q \sim U(0,1)$  ,  $0 \le q_0 \le 1$ 

where  $J \in J_i^k$  is chosen according to  $p_{iJ}^k(t) = \frac{\left[\tau_{iJ}(t)\right] \left[\eta_{iJ}\right]^{\beta}}{\sum\limits_{l \in J_i^k} \left[\tau_{il}(t)\right] \left[\eta_{il}\right]^{\beta}}$ 

else choose the closest city  $j \in J_i^k$ 

end if

after each transition ant k applies the local update rule

$$\tau_{ij}(t) \leftarrow (1-\rho)\tau_{ij}(t) + \rho\tau_0$$
 where  $\tau_0 = (nL_{nn})^{-1}$ 

End For

For k=1 to mcompute length  $L^k(t)$  of tour  $T^k(t)$  produced by ant kEnd For
If an improved tour is found then
update  $T^+$  and  $L^+$ End If
For every edge (i,j)  $\tau_{ij}(t) \leftarrow (1-\rho)\tau_{ij}(t) + \rho \Delta \tau_{ij}(t)$ 

where

$$\Delta \tau_{ij}(t) = \frac{1}{L^{+}}$$

End For

End For