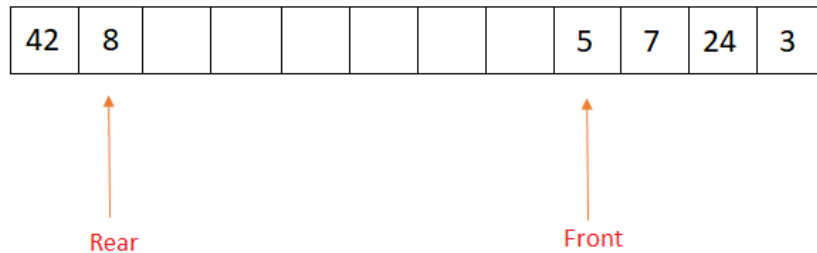


Introduction to Deques

A deque, also known as the **double-ended queue** is an ordered list in which elements can be inserted or deleted at either end. It is also known as a **head-tail linked list** because elements can be added to or removed from either the front (head) or the back (tail).



Properties of Deques

- Deque can be used both as **stack** and **queue** as it allows insertion and deletion of elements from both ends.
- Deque does not require LIFO and FIFO orderings enforced by data structures like stacks and queues.
- There are two variants of double-ended queues -
 - **Input restricted deque:** In this deque, insertions can only be done at one of the ends, while deletions can be done from both ends.
 - **Output restricted deque:** In this deque, deletions can only be done at one of the ends, while insertions can be done on both ends.

Operations on Deques

- **enqueue_front(data):** Insert an element at the front end.
- **enqueue_rear(data):** Insert an element at the rear end.
- **dequeue_front():** Delete an element from the front end.
- **dequeue_rear():** Delete an element from the rear end.
- **front():** Return the front element of the dequeue.
- **rear():** Return the rear element of the dequeue.
- **isEmpty():** Returns true if the deque is empty.
- **isFull():** Returns true if the deque is full.

Implementation of Deques

Dequeues can be implemented using data structures like **circular arrays** or **doubly-linked lists**.

Below is the circular array implementation for dequeues, the same approach can be used to implement dequeues using doubly-linked lists.

We maintain two variables: **front** and **rear**, front represents the **front end** of the deque and rear represents the **rear end** of the deque.

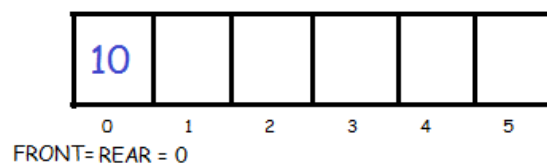
The circular array is represented as "**carr**", and size represented by **size**, having elements indexed from **0** to **size - 1**.

- **Inserting** an element at the **front** end involves **decrementing** the front pointer.
- **Deleting** an element from the **front** end involves **incrementing** the front pointer.
- **Inserting** an element at the **rear** end involves **incrementing** the rear pointer.
- **Deleting** an element from the **rear** end involves **decrementing** the rear pointer.
- **NOTE:** However, front and rear pointers need to be maintained, such that they remain within the bounds of indexing of **0** to **size - 1** of the circular array.
- Initially, the deque is empty, so front and rear pointers are initialized to **-1**, denoting that the deque contains no element.
- **Enqueue_front operation**

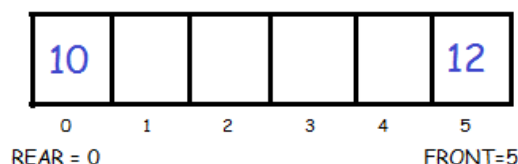
Steps:

- If the array is full, the data can't be inserted.
- If there are not any components within the Deque(or array) it means the front is equal to **-1**, increment front and rear, and set carr[front] as data.
- Else decrement front and set carr[front] as data.

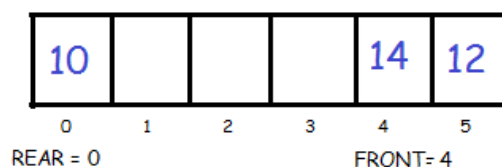
WHEN ONE ELEMENT IS ADDED
LET'S SAY 10,



INSERT 12 AT FRONT.



NOW INSERT 14 AT FRONT



```
function enqueue_front(data)

    // Check if deque is full
    if (front equals 0 and rear equals size-1) or (front equals rear+1)
        print ("Overflow")
        return

    /*
        Check if the deque is empty, insertion of an element from the
        front or rear will be equivalent.
        So update both front and rear to 0.
    */
    if front equals -1
        front = 0
        rear = 0
        carr[front] = data
        return

    // Otherwise check if the front is 0
    if front equals 0
        /*
            Updating front to size-1, so that front remains within the
            bounds of the circular array.
        */
        front = size-1
    else
        front = front-1

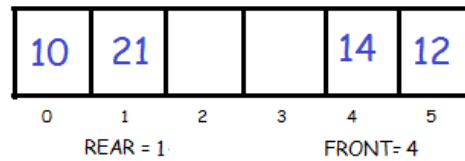
    carr[front] = data
    return
```

- **Enqueue_rear operation**

Steps:

- If the array is already full then it is not possible to insert more elements.
- If there are not any elements within the Deque i.e. rear is equal to -1, increase front and rear and set carr[rear] as data.
- Else increment rear and set carr[rear] as data.

INSERT 21 AT REAR



```
function enqueue_rear(data)

    // Check if deque is full
    if (front equals 0 and rear equals size-1) or (front equals rear+1)
        print ("Overflow")
        return

    /*
        Check if the deque is empty, insertion of an element from the
        front or rear will be equivalent.
        So update both front and rear to 0.
    */

    if front equals -1
        front = 0
        rear = 0
        carr[front] = data
        return

    // Otherwise check if rear equals size-1

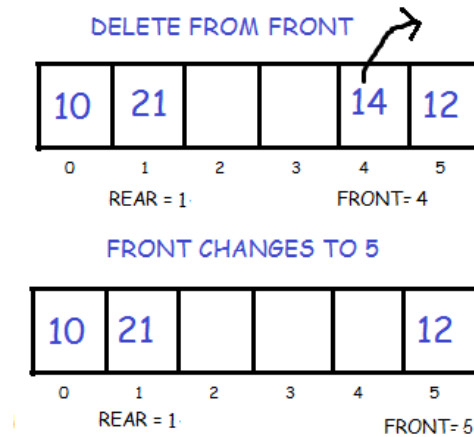
    if rear equals size-1
        /*
            Updating rear to 0, so that rear remains within the bounds
            of the circular array.
        */
        rear = 0
    else
        rear = rear+1

    carr[rear] = data
    return
```

- Dequeue_front operation

Steps:

- If the Deque is empty, return.
- If there is only one element in the Deque, that is, front equals rear, set front and rear as -1.
- Else increment front by 1.



```
function dequeue_front()
```

```
    // Check if deque is empty
```

```
    if front equals -1
```

```
        print ("Underflow")
```

```
        return
```

```
    /*
```

```
        Otherwise, check if the deque has a single element i.e front and
        rear are equal and will be non-negative as the deque is
        non-empty.
```

```
    */
```

```
    if front equals the rear
```

```
        /*
```

```
            Update front and rear back to -1, as the deque becomes
            empty.
```

```
        */
```

```
        front = -1
```

```
        rear = -1
```

```
        return
```

```
    // If the deque contains at least 2 elements.
```

```
    if front equals size-1
```

```
        // Bring front back to the start of the circular array.
```

```

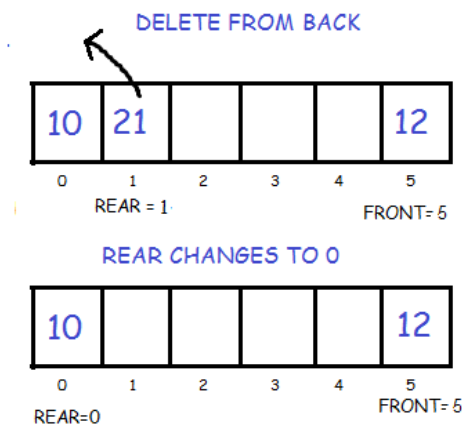
        front = 0
    else
        front = front+1
    return

```

- **Dequeue_rear operation**

Steps:

- If the Deque is empty, return.
- If there's just one component within the Deque, that is, rear equals front, set front and rear as -1.
- Else decrement rear by one.



```

function dequeue_rear()

```

```

    // Check if deque is empty
    if front equals -1
        print ("Underflow")
        return

```

```

    /*

```

```

        Otherwise, check if the deque has a single element i.e front and
        rear are equal and will be non-negative as the deque is
        non-empty.

```

```

    */

```

```

    if front equals the rear

```

```

        /*

```

```

            Update front and rear back to -1, as the deque becomes
            empty.

```

```

        */

```

```
front = -1
rear = -1
return

// If the deque contains at least 2 elements.

if rear equals 0
    // Bring rear back to the last index i.e size-1 of the circular array.
    rear = size-1
else
    rear = rear-1
return
```

- **Front operation**

Steps:

- If the Deque is empty, return.
- Else return carr[front].

```
function front()

// Check if deque is empty
if front equals -1
    print ("Deque is empty")
    return

// Otherwise return the element present at the front end
return carr[front]
```

- **Rear operation**

Steps:

- If the Deque is empty, return.
- Else return carr[rear].

```
function rear()

// Check if deque is empty
if front equals -1
    print ("Deque is empty")
    return
```

```
// Otherwise return the element present at the rear end  
return carr[rear]
```

- **Is Empty operation**

Steps:

- If front equals -1, the Deque is empty, else it's not.

```
function isEmpty()
```

```
// Check if front is -1 i.e no elements are present in deque.  
if front equals -1  
    return true  
else  
    return false
```

- **Is Full operation**

Steps:

- If front equals 0 and rear equals size - 1, or front equals rear + 1, then the Deque is full, else it's not. Here **"size"** is the size of the circular array.

```
function isFull()
```

```
/*  
    Check if the front is 0 and rear is size-1 or front == rear+1, in both  
    cases, we cannot move front and rear to perform any insertions  
*/  
  
if (front equals 0 and rear equals size-1) or (front equals rear+1)  
    return true  
else  
    return false
```

Time complexity of various operations

Let 'n' be the number of elements in the deque. The time complexities of deque operations in the worst case can be given as:

Operations	Time Complexity
Enqueue_front(data)	O(1)
Enqueue_rear(data)	O(1)
Dequeue_front()	O(1)
Dequeue_rear()	O(1)
Front()	O(1)
Rear()	O(1)
isEmpty()	O(1)
isFull()	O(1)

Applications of Deques

- Since deques can be used as stack and queue, they can be used to perform **undo-redo** operations in software applications.
- Deques are used in the **A-steal job scheduling algorithm** that implements task scheduling for multiple processors (multiprocessor scheduling).
- They are also helpful in finding max/min values of all **subarrays of size k** in the array in O(n) time, where n is the size of the array