The Apriori Algorithm for Finding Association Rules

function apriori $(I, T, s_{\min}, c_{\min}, k_{\max})$ (∗ apriori algorithm for association rules ∗)
begin
    k := 1;   (∗ — find frequent item sets ∗)
    $C_k := \bigcup_{i \in I} \{i\}$;   (∗ start with single element sets ∗)
    $F_k := \mathrm{prune}(C_k, T, s_{\min})$;   (∗ and determine the frequent ones ∗)
    while $F_k \neq \emptyset$ **and** $k \leq k_{\max}$ **do begin**   (∗ while there are frequent item sets ∗)
       $C_{k+1} := \mathrm{candidates}(F_k)$;   (∗ create item sets with one item more ∗)
       $F_{k+1} := \mathrm{prune}(C_{k+1}, T, s_{\min})$;   (∗ and determine the frequent ones ∗)
       k := $k + 1$;   (∗ increment the item counter ∗)
    end;
    $R := \emptyset$;   (∗ — generate association rules ∗)
    forall $f \in \bigcup_{j=2}^{k} F_j$ **do begin**   (∗ traverse the frequent item sets ∗)
       m := 1;   (∗ start with rule heads (consequents) ∗)
       $H_m := \bigcup_{i \in f} \{i\}$;   (∗ that contain only one item ∗)
       repeat   (∗ traverse rule heads of increasing size ∗)
          forall $h \in H_m$ **do**   (∗ traverse the possible rule heads ∗)
             if $\frac{s(f)}{s(f-h)} \geq c_{\min}$ (∗ if the confidence of the rule ∗)
             then R := $R \cup \{[(f-h) \to h]\}$;   (∗ is high enough, add it to the result, ∗)
             else $H_m := H_m - \{h\}$;   (∗ otherwise discard the rule head ∗)
          $H_{m+1} := \mathrm{candidates}(H_m)$;   (∗ create rule heads with one item more ∗)
          m := $m + 1$;   (∗ increment the head item counter ∗)
       until $H_m = \emptyset$ **or** $m \geq |f|$;   (∗ until there are no more rule heads ∗)
    end;   (∗ or the antecedent would become empty ∗)
    return R;   (∗ return the rules found ∗)
end (∗ apriori ∗)

function candidates (F_k) (∗ generate candidates with $k+1$ items ∗)
begin
    $C := \emptyset$;   (∗ initialize the set of candidates ∗)
    forall $f_1, f_2 \in F_k$ (∗ traverse all pairs of frequent item sets ∗)
    with $f_1 = \{i_1, \ldots, i_{k-1}, i_k\}$ (∗ that differ only in one item and ∗)
    and $f_2 = \{i_1, \ldots, i_{k-1}, i_k'\}$ (∗ are in a lexicographic order ∗)
    and $i_k < i_k'$ **do begin**   (∗ (the order is arbitrary, but fixed) ∗)
       $f := f_1 \cup f_2 = \{i_1, \ldots, i_{k-1}, i_k, i_k'\}$;   (∗ the union of these sets has $k+1$ items ∗)
       if $\forall i \in f : \ f - \{i\} \in F_k$ (∗ only if all k element subsets are frequent, ∗)
       then $C := C \cup \{f\}$;   (∗ add the new item set to the candidates ∗)
    end;   (∗ (otherwise it cannot be frequent) ∗)
    return C;   (∗ return the generated candidates ∗)
end (∗ candidates ∗)

function prune $(C, T, s_{\min})$ (∗ prune infrequent candidates ∗)
begin
    forall $c \in C$ **do**   (∗ initialize the support counters ∗)
       $s(c) := 0$;   (∗ of all candidates to be checked ∗)
    forall $t \in T$ **do**   (∗ traverse the transactions ∗)
       forall $c \in C$ **do**   (∗ traverse the candidates ∗)
          if $c \in t$ (∗ if the transaction contains the candidate, ∗)
          then $s(c) := s(c) + 1$;   (∗ increment the support counter ∗)
end (∗ prune ∗)