Hw1  $\alpha$  A[0...4]= $\langle 10,7,0,3,8 \rangle$  $I(A) = \{(0,1), (0,2), (0,3), (0,4), (1,2), (1,3)\}$ for i=0 to n-1: a,b,c, j:=i while j+1>0: : [i] A ([i] A fi add (j,i) to the set Cs j := j-1Let tibe the number of times line 5 runs. Thus, ti Li-1 Lines  $a_1 + nb_1 + (n-1)(C_1 + C_2)$ 1,2 iC3 (i-1) (Cu+Cb) Runtime  $T(A) = a_1 + nb_1 + (n-1)(c_1+c_2) + (\sum_{i=0}^{n-1} i) c_3 + (\sum_{i=0}^{n-1} (i-1))(c_4+c_6) + \sum_{i=0}^{n-1} t_i$  $= (a,+b,)+(n-1)(b,+C,+C_2)+\frac{n(n-1)}{2}(c_3+c_4+c_6)-n(c_4+c_6)+\sum_{i=1}^{n-1}i-1$ = (a+b,)+(n-1)(b+C,+C2-(Cu+C6+1))+n(n-1)(C3+Cu+C6)-(Cu+C6+1) = A+(n-1)(B)+n(n-1) C Qualitie

 $O(|IAN|)=n^2$ 

|I(A)| is minimized when there are no elements in A greater than any subsequent elements. This makes ti=0

|I(A)| is maximized when all elements are greater than all subsequent elements. This makes  $t_i = i-1$ 

Selection sort pseudo-code for i=0 in n-1 a, b, c, j:= i+1 a:=A[i] 3 Cy 4 min := a while j < n-1: 5 if A[i] < min: Cb 6 min:=A[i] CI j := j+1CB A[i] := min 9 C10 min:=a Lines 1...4,9,10: a,+nb,+(n-1)(C,+C2+C3+C4+Ca+C10) 5: (n-i+1) Cs 6,8: (n-i)(C,+C,) 7: ti C7 I(h) = a,+nb,+(n-1)(c,+c2+c3...) + (n(n+1) + n)C + (n(n+1) c6c3) (n(n+1) C7) A=a,+b,+C3 B=b,+C3+C,+C2+C3+C4+C4+C10 C=C5+C6+C7+C8  $|I(A)| = A + (n-1)B + \frac{n(n+1)}{2}C$ Quadratic time; 0=n2