

①

QUIZ – 15 minutes – open notes

1. Prove that every simple planar graph with fewer than 12 vertices has a vertex of degree at most four.
2. Give an example of the following or explain why no such example exists: a nonplanar graph with chromatic number 3.

Each question worth 5 marks.

Saturday, April 11, 2015

QUIZ. THURSDAY
2015 04 09

SOLUTION

1. Assume B.W.O.C. that there exists a planar graph G of order $n \leq 11$ s.t.

$$\delta(G) \geq 5.$$

$$\therefore 2|E(G)| = \sum_{v \in V(G)} d(v) \geq 5n \quad (1)$$

Putting $|E(G)| = m$, we also have

$$m \leq 3n - 6 \quad \text{since } G \text{ is planar}$$

$$\therefore 5n \leq 2m \leq 2(3n - 6) = 6n - 12$$

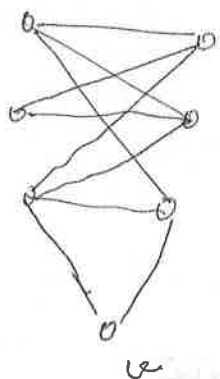
$$\Rightarrow 12 \leq n$$

$$\Rightarrow \text{contradiction}$$

Q2. There can be non-planar graphs with chromatic number 3.

(2)

Example 1: G consists of a $K_{3,3}$ and an additional vertex v which is joined to precisely one vertex of the partite sets of the $K_{3,3}$.



Since G contains a $K_{3,3}$, it is nonplanar by Kuratowski's Theorem. Since G is not bipartite, $\chi(G) > 2$.

OTOH, we can properly color G with 3 colors by giving v a color different from those used to color the $K_{3,3}$ subgraph.

$$\therefore 2 < \chi(G) \leq 3 \Rightarrow \chi(G) = 3$$

as required.

Example 2: The Petersen Graph, P

Since P is a 3-regular, $\chi(G) \leq 3$ by Brooks' Theorem.

OTOH, P contains a 5-cycle so it is not bipartite.

$$\therefore \chi(G) > 2$$

$$\text{Hence, again } \chi(G) = 3.$$

Since P contains a subdivision of K_5 , it is nonplanar by Kuratowski's Theorem.