

Proposition 52:- If  $G$  is a planar graph, then  $\chi(G) \leq 6$ .

Lemma (this is actually a problem in the textbook):

If  $G$  is planar, then  $\delta(G) \leq 5$ .

Proof: We use Proposition 50.

Now, if  $\bar{d}$  = average degree in  $G$ , then  $\delta \leq \bar{d}$ .

$$\begin{aligned} \text{Hence, } \delta \leq \bar{d} &= \frac{\sum_{v \in V(G)} d(v)}{|V(G)|} = \frac{2|E(G)|}{|V(G)|} \leq \frac{2(3|V(G)| - 6)}{|V(G)|} \\ &= 6 - \frac{12}{|V(G)|} < 6, \text{ as reqd.} \end{aligned}$$

Proof of Main Proposition 52:-

By induction on  $n = |V(G)|$

Base Case: The result is obviously true for  $n \leq 6$ .

I.H: Suppose result holds for  $n$ , and suppose  $|V(G)| = n+1$ .

Now,  $v \in V(G)$  with  $d(v) \leq 5$ .

Consider  $G' = G - v$ . Then  $G'$  is planar, with  $n$  vertices.

Hence, by I.H, it is 6-colorable.

Since  $v$  has at most 5 neighbours, we can always give it the 6th color to get a proper 6-coloring of  $G$ .

Proof of the 5-color Theorem (Thm. 59):-

We will follow essentially the same logic as above; however, the difficulty comes if  $v$  has 5 neighbours which all have different colors! Anyway, we proceed by induction on  $n$ .

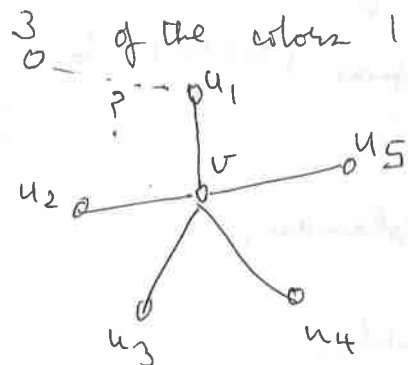
Base Case: Obviously, result holds if  $n \leq 5$ .

I.H: Suppose result holds for all graphs with  $n$  vertices, and suppose  $G$  has  $n+1$  vertices.

let  $v \in G$  with  $d(v) \leq 5$ ,  
 let  $G' = G - v$ , so by  
 induction  $G'$  is 5-colorable,  
 so we properly color  $G'$  with 5  
 colors. Now, if  $N(v)$  has  
 $\leq 4$  different colors, we  
 are done.

So we are reduced to the  
 case where  $d(v) = 5$  and  
 the 5 neighbours of  $v$  have  
 5 different colors. Since we  
 can't extend this coloring to  
 $G$ , we have to do some  
 recoloring.

let us have a crossing-free  
 embedding of  $G$  - every vertex  
 of  $G$  (except  $v$ ) is colored by one  
 of the colors 1 to 5.



let  $u_1$  to  $u_5$   
 be the 5 neighbours  
 of  $v$  in  
 anti-clockwise  
 order, and  
 wlog, let them  
 have colors 1 to 5

respectively.

We want to change the color of  
 one of  $v$ 's neighbours, say  
 of  $u_1$  from 1 to 3 - now  
 we can color  $v$  with  $u_1$   
 and we are done. However,  $u_1$   
 might have a neighbour with  
 color 3, that would lead to  
 a problem.

let  $H_{1,3}$  be the subgraph of  
 $G$  induced by vertices of color  
 1 and 3. Notice, that if  
 in one component of  $H_{1,3}$  we  
 exchange colors 1 and 3, the

coloring <sup>of  $G'$</sup>  would remain  
 proper (since  $v$  is not  
 yet colored). (2)

We therefore exchange the  
 colors 1 and 3 in the  
 component which contains  
 $u_1$ . So it would be possible  
 to color  $v$  with color 1.

The problem however is that  
 $u_3$  could have been in the  
 same component of  $H_{1,3}$  as  
 $u_1$ , so now it has color 3.  
 If they are in separate components,  
 we are done (since then  $u_3$   
 retains color 3, and color 1 is  
 available for  $v$ ).

So we are left with the case  
 in which  $u_1$  and  $u_3$  are in  
 the same component of  $H_{1,3}$ .

So now we try to recolor  
 vertex 2 with color 4.  
 Proceeding in the same way,  
 we are forced to the situation  
 that  $u_2$  and  $u_4$  are in  
 the same component of  $H_{2,4}$ .  
 We claim this cannot happen.

Now, since  $u_1$  and  $u_3$  are  
 in the same component of  $H_{1,3}$ ,  
 there exists a path  $P$  from  $u_1$  to  $u_3$   
 in  $H_{1,3}$ .

Similarly, there is a path  $Q$   
 in  $H_{2,4}$  from  $u_2$  to  $u_4$ .

Clearly,  $P$  and  $Q$  have no  
 vertices in common.

Furthermore, path  $P$  along with  $v$   
 forms a cycle, which becomes a  
 simple closed curve in the plane.  
 However vertices  $u_2$  and  $u_4$  are  
 on different sides of this closed curve,  
 i.e. the path  $Q$  must cross this  
 curve  $\Rightarrow u_2, u_4$  are  
 in different components, and we are  
 done.