

For Thursday 20/5/2020

(1)

(this is a simplified and tightened version of what was done in the class).

Proof of Theorem 3 (Konig - Egervary):-

Let G be a fixed X - Y bipartite graph, and let M be a maximum matching in G . We will construct a vertex cover Q with $|Q| = |M|$. In view of earlier remark, the result follows.

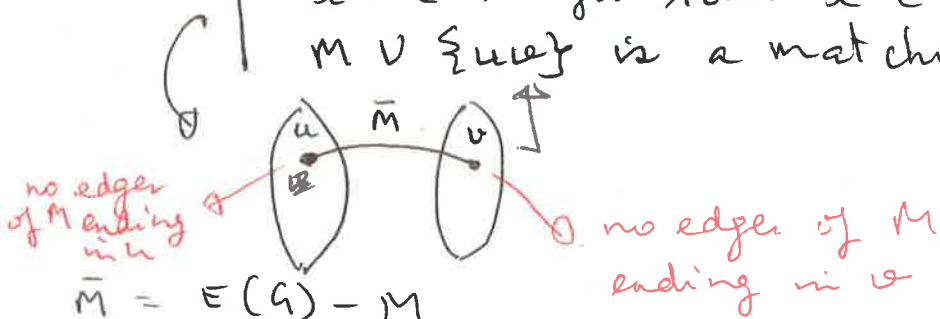
For each edge $xy \in M$, select one of its end-points :- select y if y is the end-point of an alternating path starting in X s.t. its first edge (from X to Y) is not in M ; i.e. if $P = x_1 y_1 \dots x_k y$ and $x_1 y_1 \notin M$. Else, select x .

Put $Q =$ set of selected vertices, so $|Q| = |M|$. We need to show Q is a vertex cover.

So let $uv \in E(G)$, $u \in X$, $v \in Y$.

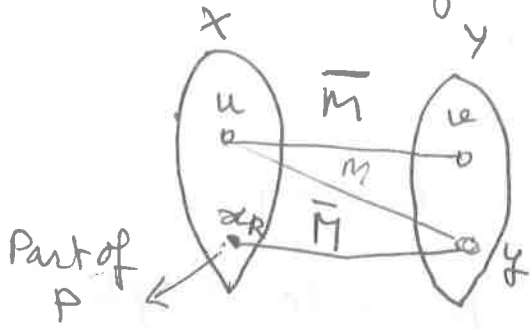
If $uv \in M$, then either u or $v \in Q$ and we are done. So we may assume $uv \notin M$.

But, either $uy \in M$ for some $y \in Y$ or $xv \in M$ for some $x \in X$; if not, then $M \cup \{uv\}$ is a matching larger than $M \Rightarrow \Leftarrow$



(P.T.O.)

Case 1. $u y \in M$ for some $y \in Y$.



Then either $u \in Q$ or $y \in Q$.

If $u \in Q$, we are done,

so ~~assume~~ $y \in Q$.

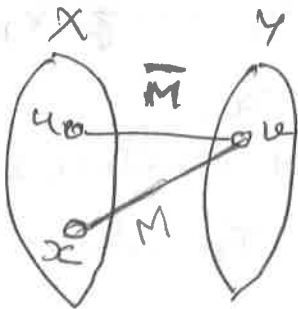
$\therefore y$ is the end-point of an alternating path $P: x_1 y_1 \dots x_k y$.

If $x_k \neq u$, then $u y x_k$ is an M -augmenting path $\Rightarrow \Leftarrow$

~~$x_k = u$~~

But if $x_k = u$, then since P is an alternating path with $x_1 y_1 \notin M$, it follows that ~~$x_k y \notin M$~~
 $x_k y = u y \notin M \Rightarrow \Leftarrow$

Case 2. $x v \in M$ for some ~~$x \in X$~~ $x \in X$.



Then v being an end-point of the alternating path $P: u v$ has to be the vertex for Q selected on behalf of the edge $x v \in M$,

i.e. $v \in Q$.