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1 Combinatorial optimization

1.1 Max bipartite matchine

```
}
}
return false;
}
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
    mr = VI(w.size(), -1);
    mc = VI(w[0].size(), -1);
    int ct = 0;
    for (int i = 0; i < w.size(); i++) {
        VI seen(w[0].size());
        if (FindMatch(i, w, mr, mc, seen)) ct++;
    }
return ct;
}
</pre>
```

2 Geometry

2.1 Convex hull

```
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
#define ld long double
struct Point {
        ld x, y;
        bool operator < (const Point &p) const {</pre>
                 return x < p.x | | (x == p.x && y < p.y);
};
ld cross(const Point &O, const Point &A, const Point &B)
        return (A.x - 0.x) * (B.y - 0.y) - (A.y - 0.y) * (B.x - 0.x);
vector<Point> convex_hull(vector<Point> P)
        int n = P.size();
        int k = 0;
        vector<Point> H(2*n);
        sort(P.begin(), P.end());
        for (int i=0; i<n; i++)</pre>
                 while (k \ge 2 \&\& cross(H[k-2], H[k-1], P[i]) \le 0) k--;
                 H[k++] = P[i];
        for (int i=n-2, t = k+1; i>=0; i--)
                 while (k \ge t \&\& cross(H[k-2], H[k-1], P[i]) \le 0) k--;
                 H[k++] = P[i];
        H.resize(k-1);
        return H;
int main()
        ios_base::sync_with_stdio(false);
        return 0;
```

2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT (const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p)
                                 const
                                       { return PT(x+p.x, y+p.y);
  PT operator - (const PT &p)
PT operator * (double c)
                                       { return PT(x-p.x, y-p.y);
                                 const
                                         return PT(x*c,
                                 const
                                                            V*C
  PT operator / (double c)
                                 const { return PT(x/c,
                                                            y/c
                             { return p.x*q.x+p.y*q.y; }
double dot (PT p, PT q)
double dist2(PT p, PT q)
double cross(PT p, PT q)
                             { return dot(p-q,p-q); }
                           { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
  os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x);
                       { return PT(-p.y,p.x); }
PT RotateCW90 (PT p)
                        { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b // assuming a !=\ b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
```

```
&& fabs(cross(c-d, c-a)) < EPS;
}
  ^\prime determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
    if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-b, d-b) > 0)
      return false;
    return true;
     (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
  ^{\prime} segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
  b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact // tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
      p[j].y \le q.y \&\& q.y < p[i].y) \&\&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)
      return true;
    return false;
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
  b = b-a;
a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
```

```
if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++)
    for (int k = i+1; k < p.size(); k++) {
      int j = (i+1) % p.size();
      int \bar{1} = (k+1) % p.size();
      if (i == 1 \mid | j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
```

```
cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
      << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
      << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;</pre>
// expected: 6.78903
cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
      << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "</pre>
      << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 0 0 1
// expected: 1 1 1 0
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v;
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push\_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
      << PointInPolygon(v, PT(2,0)) << " "
      << PointInPolygon(v, PT(0,2)) << " "
      << PointInPolygon(v, PT(5,2)) << " "
      << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
      << PointOnPolygon(v, PT(2,0)) << " "
      << PointOnPolygon(v, PT(0,2)) << " "
      << PointOnPolygon(v, PT(5,2)) << " "
      << PointOnPolygon(v, PT(2,5)) << endl;
   expected: (1,6)
                (5, 4)
                        (4,5)
                blank line
                (4,5) (5,4)
                blank line
                (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr <math><< endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; <math>cerr << endl;
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;</pre>
cerr << "Centroid: " << c << endl;</pre>
```

```
return 0;
}
```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that // involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
         return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
     while (b) { int t = a%b; a = b; b = t; }
     return a;
}
// computes lcm(a,b)
int lcm(int a, int b)
         return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
         int ret = 1;
         while (b)
                  if (b & 1) ret = mod(ret*a, m);
                  a = mod(a*a, m);
                  b >>= 1;
         return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
         int xx = y = 0;
         int yy = x = 1;
         while (b) {
                  int q = a / b;
                  int t = b; b = a%b; a = t;
                  t = xx; xx = x - q*xx; x = t;
                  t = yy; yy = y - q*yy; y = t;
         return a;
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
         int x, y;
         VI ret;
         int g = extended_euclid(a, n, x, y);
         if (!(b%g)) {
                  x = mod(x*(b / g), n);
for (int i = 0; i < g; i++)
                           ret.push_back(mod(x + i*(n / g), n));
         }
```

```
return ret;
}
// computes b such that ab = 1 (mod n), returns -1 on failure
int mod_inverse(int a, int n) {
         int x, y;
int g = extended_euclid(a, n, x, y);
         if (g > 1) return -1;
         return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
         int s, t;
         int g = extended_euclid(m1, m2, s, t);
         if (r1%g != r2%g) return make_pair(0, -1);
         return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / q, m1*m2 / q);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is // unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
         PII ret = make_pair(r[0], m[0]);
         for (int i = 1; i < m.size(); i++) {</pre>
                  ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i
                     ]);
                  if (ret.second == -1) break;
         return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
         if (!a && !b)
                  if (c) return false;
                  x = 0; y = 0;
return true;
         if (!a)
                  if (c % b) return false;
                  x = 0; y = c / b; return true;
         if (!b)
                  if (c % a) return false;
                  x = c / a; y = 0;
                  return true;
         int g = gcd(a, b);
         if (c % g) return false;
x = c / g * mod_inverse(a / g, b / g);
         y = (c - a*x) / b;
         return true;
int main() {
    // expected: 2
         cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
         int x, y;
int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
         // expected: 95 451
         VI sols = modular_linear_equation_solver(14, 30, 100);
```

```
for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";
    cout << endl;

// expected: 8
    cout << mod_inverse(8, 9) << endl;

// expected: 23 105
    // 11 12

PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
    cout << ret.first << " " << ret.second << endl;
    ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
    cout << ret.first << " " << ret.second << endl;

// expected: 5 -15
    if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
    cout << x << " " << y << endl;
    return 0;
}</pre>
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
  Uses:
     (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
   Running time: O(n^3)
   INPUT:
              a[][] = an nxn matrix
              b[][] = an nxm matrix
   OUTPUT:
                     = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1;
  for (int i = 0; i < n; i++) {</pre>
    int pj = -1, pk = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])
        if (pj == -1 \mid | fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0);</pre>
    ipiv(pk)++;
    swap(a[pj], a[pk]);
swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
```

```
for (int p = 0; p < n; p++) a[pk][p] *= c;</pre>
    for (int p = 0; p < m; p++) b[pk][p] *= c;
     for (int p = 0; p < n; p++) if (p != pk) {
       c = a[p][pk];
       a[p][pk] = 0;
       for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
       for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  }
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
  return det;
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \}; double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
    a[i] = VT(A[i], A[i] + n);
    b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333 0.0666667
                  0.166667 0.166667 0.333333 -0.333333
  0.233333 0.833333 -0.133333 -0.0666667
                  0.05 - 0.75 - 0.1 0.2
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < n; j++)
  cout << a[i][j] << ' ';</pre>
     cout << end1;</pre>
  }
     expected: 1.63333 1.3
                  -0.166667 0.5
                  2.36667 1.7
  //
                  -1.85 - 1.35
  cout << "Solution: " << endl;</pre>
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < m; j++)
    cout << b[i][j] << ' ';</pre>
    cout << endl;</pre>
  }
```

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
//
// Running time: O(n^3)
//
// INPUT: a[][] = an nxm matrix
//
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
// returns rank of a[][]

#include <iostream>
#include <vector>
#include <cmath>
```

```
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
  int n = a.size();
  int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m \&\& r < n; c++) {
     int j = r;
for (int i = r + 1; i < n; i++)</pre>
       if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
     if (fabs(a[j][c]) < EPSILON) continue;</pre>
     swap(a[j], a[r]);
     T s = 1.0 / a[r][c];
     for (int j = 0; j < m; j++) a[r][j] *= s;
for (int i = 0; i < n; i++) if (i != r) {</pre>
       T t = a[i][c];
       for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
     <u>r</u>++;
  return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
    },
                      8},
       5, 11, 10,
       9, 7, 6, 12},
4, 14, 15, 1},
     {13, 21,
                21,
                      13}};
  VVT a(n);
  for (int i = 0; i < n; i++)</pre>
     a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
  // expected: 3
cout << "Rank: " << rank << endl;</pre>
   // expected: 1 0 0 1
                    0
                      1 0 3
                    0 \ 0 \ 1 \ -3
                    0 0 0 3.10862e-15
                    0
                      0 0 2.22045e-15
  cout << "rref: " << endl;</pre>
  for (int i = 0; i < 5; i++) {
  for (int j = 0; j < 4; j++)
    cout << a[i][j] << ' ';</pre>
     cout << endl;</pre>
  }
```

3.4 Fast Fourier transform

```
#include <cassert>
#include <cstdio>
#include <cmath>

struct cpx
{
    cpx(){}
    cpx(double aa):a(aa),b(0){}
    cpx(double aa, double bb):a(aa),b(bb){}
    double a;
    double b;
    double modsq(void) const
```

```
return a * a + b * b;
  cpx bar (void) const
    return cpx(a, -b);
};
cpx operator + (cpx a, cpx b)
  return cpx(a.a + b.a, a.b + b.b);
cpx operator *(cpx a, cpx b)
  return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator / (cpx a, cpx b)
  cpx r = a * b.bar();
  return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP (double theta)
  return cpx(cos(theta), sin(theta));
const double two_pi = 4 * acos(0);
// in:
            input array
            output array
  out:
// step:
            {SET TO 1} (used internally)
            length of the input/output {MUST BE A POWER OF 2}
  dir:
// dir: either plus or minus one (direction of the FFT) // RESULT: out[k] = \sum_{j=0}^{s} \sin_j x = 1 in[j] * exp(dir * 2pi * i * j * k /
   size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
  if(size < 1) return;</pre>
  if(size == 1)
    out[0] = in[0];
    return;
  FFT(in, out, step * 2, size / 2, dir);
FFT(in + step, out + size / 2, step * 2, size / 2, dir);
  for(int i = 0; i < size / 2; i++)
    cpx even = out[i];
    cpx odd = out[i + size / 2];
    out[i] = even + EXP(dir * two_pi * i / size) * odd;
    out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
// Usage:
  f[0...N-1] and g[0...N-1] are numbers
// Want to compute the convolution h, defined by
// h[n] = sum \ of \ f[k]g[n-k] \ (k = 0, ..., N-1).
// Here, the index is cyclic; f[-1] = f[N-1], \ f[-2] = f[N-2], \ etc.
// Let F[0...N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[] in O(N log ar{	exttt{N}}) time, do the following:
     1. Compute F and G (pass dir = 1 as the argument).
     2. Get H by element-wise multiplying F and G.
     3. Get h by taking the inverse FFT (use dir = -1 as the argument)
         and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.
int main(void)
  printf("If rows come in identical pairs, then everything works.\n");
```

```
 \begin{array}{l} {\rm cpx} \ {\rm a[8]} \ = \ \{0,\ 1,\ {\rm cpx}(1,3),\ {\rm cpx}(0,5),\ 1,\ 0,\ 2,\ 0\}; \\ {\rm cpx} \ {\rm b[8]} \ = \ \{1,\ {\rm cpx}(0,-2),\ {\rm cpx}(0,1),\ 3,\ -1,\ -3,\ 1,\ -2\}; \\ \end{array} 
cpx A[8];
cpx B[8];
FFT(a, A, 1, 8, 1);
FFT(b, B, 1, 8, 1);
for (int i = 0; i < 8; i++)
  printf("%7.21f%7.21f", A[i].a, A[i].b);
printf("\n");
for (int i = 0; i < 8; i++)
   cpx Ai(0,0);
  for (int j = 0; j < 8; j++)
     Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
  printf("%7.21f%7.21f", Ai.a, Ai.b);
printf("\n");
cpx AB[8];
for(int i = 0; i < 8; i++)
AB[i] = A[i] * B[i];
cpx aconvb[8];
FFT (AB, aconvb, 1, 8, -1);
for (int i = 0; i < 8; i++)
  aconvb[i] = aconvb[i] / 8;
for(int i = 0; i < 8; i++)
  printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
printf("\n");
for (int i = 0; i < 8; i++)
  cpx aconvbi(0,0);
  for (int j = 0; j < 8; j++)
     aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
  printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
printf("\n");
return 0;
```

4 Graph algorithms

4.1 Bellman-Ford shortest paths with negative edge weights (C++)

```
tor<int> dist(V, INF); dist[s] = 0; // INF = 2B not MAX_INT to avoid overflow
REP (i, 0, V - 1) // relax all E edges V-1 times, O(V)
REP (u, 0, V - 1) // these two loops = O(E)
TRvii (AdjList[u], v) // has edge and can be relaxed
    dist[v->first] = min(dist[v->first], dist[u] + v->second);
REP (i, 0, V - 1)
    printf("SSSP(%d, %d) = %d\n", s, i, dist[i]);

bool negative_cycle_exist = false;
REP (u, 0, V - 1) // one more pass to check
TRvii (AdjList[u], v)
if (dist[v->first] > dist[u] + v->second) // should be false, but if
        possible
    negative_cycle_exist = true; // then negative cycle exists!
    printf("Negative Cycle Exist? %s\n", negative_cycle_exist? "Yes" : "No")
```

4.2 Floyd-Warshall

```
EP (k, 0, V - 1) // recall that #define REP(i, a, b) for (int i = int(a); i <=
   int(b); i++)
        REP (i, 0, V - 1)
        REP (j, 0, V - 1)
        AdjMatrix[i][j] = min(AdjMatrix[i][j], AdjMatrix[i][k] + AdjMatrix[k][j])</pre>
```

4.3 Fast Dijkstra's algorithm

4.4 Strongly connected components

```
tack<int> dfs_scc; // additional information for SCC
set<int> in_stack; // for dfs_low update check
void tarjanSCC(int u) {
         dfs_low[u] = dfs_num[u] = dfsNumberCounter++; // dfs_low[u] <= dfs_num[u]</pre>
         dfs_scc.push(u); in_stack.insert(u); // stores u based on order of
            visitation
         TRvii (AdjList[u], v) {
                  if (dfs_num[v->first] == DFS_WHITE) // a tree edge
                           tarjanSCC(v->first);
                  if (in_stack.find(v->first) != in_stack.end()) // condition for
                     update
                           dfs_low[u] = min(dfs_low[u], dfs_low[v->first]); //
                              update dfs_low[u]
         if (dfs_low[u] == dfs_num[u]) { // if this is a root of SCC
                  printf("SCC: ");
                  while (!dfs_scc.empty() && dfs_scc.top() != u) {
                           printf("%d", dfs_scc.top()); in_stack.erase(dfs_scc.top
                               ()); dfs_scc.pop();
                  printf("%d\n", dfs_scc.top()); in_stack.erase(dfs_scc.top());
                     dfs_scc.pop();
         } }
```

4.5 Eulerian path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
{
    int next_vertex;
    iter reverse_edge;
```

```
Edge(int next_vertex)
                :next_vertex(next_vertex)
};
const int max_vertices = ;
int num_vertices;
list<Edge> adj[max_vertices];
                                        // adjacency list
vector<int> path;
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
                find_path(vn);
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

4.6 Kruskal's alternative

```
#include<bits/stdc++.h>
using namespace std;
typedef pair<int, int> iPair;
struct Graph{
    int V, E;
    vector< pair<int, iPair> > edges;
    Graph(int V, int E) {
        this->V = V;
        this->E = E;
    void addEdge(int u, int v, int w) {
        edges.push_back({w, {u, v}});
    int kruskalMST();
};
struct DisjointSets{
    int *parent, *rnk;
    int n;
    DisjointSets(int n) {
        this->n = n;
        parent = new int[n+1];
        rnk = new int[n+1];
        for (int i = 0; i <= n; i++) {
            rnk[i] = 0;
            parent[i] = i;
    int find(int u) {
        if (u != parent[u])
            parent[u] = find(parent[u]);
        return parent[u];
    void merge(int x, int y) {
        x = find(x), y = find(y);
```

```
if (rnk[x] > rnk[y])
                parent[y] = x;
                 parent[x] = y;
           if (rnk[x] == rnk[y])
                 rnk[y]++;
} ;
int Graph::kruskalMST() {
     int mst_wt = 0; // Initialize result
     sort(edges.begin(), edges.end());
     DisjointSets ds(V);
     vector< pair<int, iPair> >::iterator it;
     for (it=edges.begin(); it!=edges.end(); it++){
           int u = it->second.first;
           int v = it->second.second;
           int set_u = ds.find(u);
           int set_v = ds.find(v);
           if (set_u != set_v) {
    cout << u << " - " << v << endl;</pre>
                 mst wt += it->first;
                 ds.merge(set_u, set_v);
     return mst_wt;
int main(){
     int V = 9, E = 14;
     Graph g(V, E);
     g.addEdge(0, 1, 4); g.addEdge(1, 2, 8); g.addEdge(1, 7, 11); g.addEdge(2, 3, 7); g.addEdge(2, 8, 2); g.addEdge(2, 5, 4); g.addEdge(3, 4, 9); g.addEdge(3, 5, 14); g.addEdge(4, 5, 10); g.addEdge(5, 6, 2); g.addEdge(6, 7, 1); g.addEdge(6, 8, 6); g.addEdge(7, 8, 7); g.addEdge(0, 7, 8); cout << "Edges of MST are \n";
     int mst_wt = g.kruskalMST();
     cout << "\nWeight of MST is " << mst_wt;</pre>
     return 0;
```

4.7 Prim alternative

```
#include<bits/stdc++.h>
using namespace std;
# define INF 0x3f3f3f3f
typedef pair<int, int> iPair;
class Graph{
    int V;
               // No. of vertices
    list< pair<int, int> > *adj;
public:
    Graph(int V); // Constructor
    void addEdge(int u, int v, int w);
    void primMST();
Graph::Graph(int V) {
    this->V = V;
    adj = new list<iPair> [V];
void Graph::addEdge(int u, int v, int w) {
    adj[u].push_back(make_pair(v, w));
    adj[v].push_back(make_pair(u, w));
void Graph::primMST() {
    priority_queue< iPair, vector <iPair> , greater<iPair> > pq;
int src = 0; // Taking vertex 0 as source
    vector<int> key(V, INF);
    vector<int> parent(V, -1);
```

```
vector<bool> inMST(V, false);
     pq.push(make_pair(0, src));
     key[src] = 0;
     while (!pq.empty())
          int u = pq.top().second;
          pq.pop();
          inMST[u] = true; // Include vertex in MST
list< pair<int, int> >::iterator i;
          for (i = adj[u].begin(); i != adj[u].end(); ++i)
               int v = (*i).first;
               int weight = (*i).second;
               if (inMST[v] == false && key[v] > weight)
                     key[v] = weight;
                     pq.push(make_pair(key[v], v));
                     parent[v] = u;
                }
          }
     for (int i = 1; i < V; ++i)
          printf("%d - %d\n", parent[i], i);
int main(){
     int V = 9;
     Graph g(V); g.addEdge(0, 1, 4); g.addEdge(0, 7, 8); g.addEdge(1, 2, 8); g.addEdge(1, 7, 11); g.addEdge(2, 3, 7);
     g.addEdge(2, 8, 2); g.addEdge(2, 5, 4); g.addEdge(3, 4, 9); g.addEdge(3, 5, 14); g.addEdge(4, 5, 10); g.addEdge(5, 6, 2); g.addEdge(6, 7, 1); g.addEdge(6, 8, 6); g.addEdge(7, 8, 7);
     g.primMST();
     return 0;
```

4.8 Bridge Finding

```
ulationPointAndBridge(int u) {
       dfs_low[u] = dfs_num[u] = dfsNumberCounter++; // dfs_low[u] <= dfs_num[u]</pre>
       TRvii (AdjList[u], v)
               if (u == dfsRoot) // special case
                               rootChildren++; // count children of root
                       articulationPointAndBridge(v->first);
                       if (dfs_low[v->first] >= dfs_num[u]) // for articulation
                          point
                              articulation_vertex[u] = true; // store this
                                 information first
                       if (dfs_low[v->first] > dfs_num[u]) // for bridge
                              printf(" Edge (%d, %d) is a bridge\n", u, v->
                                 first);
                       dfs_low[u] = min(dfs_low[u], dfs_low[v->first]); //
                          update dfs_low[u]
               else if (v->first != dfs_parent[u]) // a back edge and not direct
                   cycle
                       dfs_low[u] = min(dfs_low[u], dfs_num[v->first]); //
                          update dfs_low[u]
// inside int main()
dfsNumberCounter = 0;
memset(dfs_num, DFS_WHITE, sizeof dfs_num);
printf("Bridges:\n");
REP (i, 0, V - 1)
       if (dfs_num[i] == DFS_WHITE) {
               dfsRoot = i; rootChildren = 0;
```

5 Data structures

5.1 Suffix array

```
#define MAXXN 65536
#define MAXLG 17
char A[MAXN];
struct entry
                                           int nr[2];
                                                          int p;
} L[MAXN];
int P[MAXLG][MAXN];
int N,i;
int stp, cnt;
int cmp(struct entry a, struct entry b)
                                           return a.nr[0] == b.nr[0] ?(a.nr[1] < b.nr[1] ?1: 0): (a.nr[0] < b.nr[0]
                                                      ?1: 0);
}
int main()
                             gets(A);
                            for (N=strlen(A), i = 0; i < N; i++)
        P[0][i] = A[i] - 'a';
for(stp=1, cnt = 1; cnt < N; stp++, cnt *= 2)</pre>
                                                          for (i=0; i < N; i++)
                                                                                       L[i].nr[0]=P[stp- 1][i];
                                                                                       L[i].nr[1]=i +cnt <N? P[stp -1][i+ cnt]:-1;
                             sort(L, L+N, cmp);
                             for (i=0; i < N; i++)
                                                          P[stp][L[i].p] = i > 0 && L[i].nr[0] == L[i-1].nr[0] && L[i].nr[1] == L[i-1].nr[1] && L[i].nr[1] == L[i-1].nr[1] && L[i].nr[1] == L[i-1].nr[1] && L[i].nr[1] &
                                                                        L[i-1].nr[1] ? P[stp][L[i-1].p] : i;
                            return 0;
          LCP Pseudocode :-
          FindLCP (x, y) 
answer = 0
                                                      for k = ceil(log N) to 0
                                                                                    if SortIndex[k][x] = SortIndex[k][y]
                                                                                               // sort-index is same if the first k characters are
                                                                                                          same
                                                                                                                           answer += 2k
                                                                                                          // now we wish to find the characters that are same
                                                                                                                    in the remaining strings
                                                                                                                        x += 2k
                                                                                           v += 2k
*/
```

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x \le N)
    tree[x] += v;
    x += (x \& -x);
  }
}
// get cumulative sum up to and including x
int get(int x) {
  int res = 0;
  while(x) {
    res += tree[x];
    x = (x \& -x);
  return res;
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x)
  int idx = 0, mask = N;
  while (mask && idx < N)</pre>
    int t = idx + mask;
    if(x \ge tree[t]) {
      idx = t;
      x -= tree[t];
    mask >>= 1;
  return idx;
```

5.3 Union-find set

```
#include <iostream>
#include <vector>
using namespace std;
int find(vector<int> &C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]);
}
void merge(vector<int> &C, int x, int y) { C[find(C, x)] = find(C, y); }
int main()
{
    int n = 5;
    vector<int> C(n);
    for (int i = 0; i < n; i++) C[i] = i;
    merge(C, 0, 2);
    merge(C, 1, 0);
    merge(C, 3, 4);
    for (int i = 0; i < n; i++) cout << i << " " << find(C, i) << endl;
    return 0;
}</pre>
```

5.4 Lazy segment tree

```
#include <bits/stdc++.h>
using namespace std;
const int N = 1e5+10;
int A[N];
int tree[4*N];
```

```
void build(int node, int start, int end)
        if(start == end) {
                 tree[node] = A[start];
        else{
                 int mid = (start + end) / 2;
                 build(2*node, start, mid);
                 build(2*node+1, mid+1, end);
                 tree[node] = tree[2*node] + tree[2*node+1];
        }
}
void update(int node, int start, int end, int idx, int val){
        if(start == end) {
                 A[idx] += val;
                 tree[node] += val;
        else{
                 int mid = (start + end) / 2;
                 if(start <= idx and idx <= mid) {</pre>
                          update(2*node, start, mid, idx, val);
                 else{
                          update(2*node+1, mid+1, end, idx, val);
                 tree[node] = tree[2*node] + tree[2*node+1];
int query(int node, int start, int end, int 1, int r){
        if(r < start or end < 1) {
                 return 0;
        if(1 <= start and end <= r){
                 return tree[node];
        int mid = (start + end) / 2;
        int p1 = query(2*node, start, mid, 1, r);
        int p2 = query(2*node+1, mid+1, end, 1, r);
        return (p1 + p2);
void updateRange(int node, int start, int end, int 1, int r, int val){
        if(lazy[node] != 0) {
                 tree[node] += (end - start + 1) * lazy[node];
                                                                      // Update it
                 if(start != end) {
                          lazy[node*2] += lazy[node];
                                                                           // Mark
                             child as lazy
                          lazy[node*2+1] += lazy[node];
                                                                           // Mark
                             child as lazy
                 lazy[node] = 0;
                                                                      // Reset it
        if(start > end or start > r or end < 1)</pre>
                                                                 // Current segment
            is not within range [1, r]
                 return;
        if(start >= 1 and end <= r){
                 tree[node] += (end - start + 1) * val;
                 if(start != end) {
                          lazy[node*2] += val;
                          lazy[node*2+1] += val;
                 return;
        int mid = (start + end) / 2;
updateRange(node*2, start, mid, l, r, val);
updateRange(node*2 + 1, mid + 1, end, l, r, val);
                                                               // Updating left child
                                                                // Updating right
            child
        tree[node] = tree[node*2] + tree[node*2+1];
                                                               // Updating root with
            max value
}
```

```
int queryRange(int node, int start, int end, int 1, int r){
        if(start > end or start > r or end < 1)</pre>
                return 0;
                                   // Out of range
        if(lazy[node] != 0){
                tree[node] += (end - start + 1) * lazy[node];
                                                                           //
                   Update it
                if(start != end)
                         lazy[node*2] += lazy[node];
                                                              // Mark child as lazy
                                                         // Mark child as lazy
                        lazy[node*2+1] += lazy[node];
                lazy[node] = 0;
                                                 // Reset it
        if(start >= 1 and end <= r)</pre>
                                                 // Current segment is totally
           within range [1, r]
                return tree[node];
        int mid = (start + end) /
                                   2;
        int p1 = queryRange(node*2, start, mid, 1, r);
                                                                // Query left
        int p2 = queryRange(node*2 + 1, mid + 1, end, 1, r); // Query right child
        return (p1 + p2);
int main(){
        return 0;
```

5.5 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                          // children[i] contains the children of
   node i
                                         // A[i][j] is the 2^j-th ancestor of node
int A[max_nodes][log_max_nodes+1];
    i, or -1 if that ancestor does not exist
int L[max_nodes];
                                         // L[i] is the distance between node i
   and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
    if(n==0)
        return -1;
    int p = 0;
    if (n >= 1 << 16) \{ n >>= 16; p += 16; \}
                                       8; }
    if (n >= 1<< 8) { n >>=
                             8; p +=
    if (n >= 1<< 4) { n >>=
                             4; p +=
                                        4; }
    if (n >= 1<< 2) { n >>=
                              2; p +=
                                        2;
                                          }
    if (n >= 1 << 1) {
    return p;
void DFS(int i, int 1)
    L[i] = 1;
    for(int j = 0; j < children[i].size(); j++)</pre>
        DFS(children[i][j], l+1);
}
int LCA(int p, int q)
    // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node {
m p} situated on the same level as {
m q}
    for(int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1 << i) >= L[q])
            p = A[p][i];
    if(p == q)
        return p;
    // "binary search" for the LCA
```

```
for (int i = log_num_nodes; i >= 0; i--)
         if(A[p][i]] = -1 & A[p][i] != A[q][i]
             p = A[p][i];
             q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
    // read num_nodes, the total number of nodes
    log_num_nodes=lb(num_nodes);
    for (int i = 0; i < num_nodes; i++)</pre>
         int p;
         // read p, the parent of node i or -1 if node i is the root
         A[i][0] = p;
         if (p != -1)
             children[p].push_back(i);
         else
             root = i;
     // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)
    for(int i = 0; i < num_nodes; i++)</pre>
             if (A[i][j-1] !=-1)
                  A[i][j] = A[A[i][j-1]][j-1];
                  A[i][j] = -1;
    // precompute L
DFS(root, 0);
    return 0;
}
```

5.6 Sparse Table

```
#include <bits/stdc++.h>
using namespace std;
#define 11 long long
const int MAX_SIZE = 1e5+10;
const int K = 16;
const int ZERO = 0; //define zero according to F
11 arr[MAX_SIZE];
11 Table[MAX_SIZE][K+1];
ll F(ll a, ll b)
{
        return max(a, b);
void buildTable(int n)
         for (int i=0; i<n; i++)</pre>
                 Table[i][0] = arr[i];
         for(int j = 1; j<=K; j++)
                 for (int i=0; i<= n - (1<<j); i++)</pre>
                          Table[i][j] = F(Table[i][j-1], Table[i + (1 << (j-1))][j]
                 }
11 query(int 1, int r)
        11 \text{ answer} = ZERO;
         for (int j=K; j>=0; j--)
```

5.7 Segmented Sieve

```
#include <bits/stdc++.h>
using namespace std;
//all primes strictly smaller than limit -> stored in vectpr prime
void simpleSieve(int limit, vector<int> &prime)
        bool mark[limit + 1];
        memset(mark, true, sizeof(mark));
        for (int p=2; p*p < limit; <math>p++)
                 if(mark[p] == true)
                          prime.push_back(p);
                          for(int i=p*p; i<limit; i += p)</pre>
                                  mark[i] = false;
                 }
void segmentedSieve(int n)
        int limit = floor(sqrt(n)) + 1;
        vector<int> prime;
        simpleSieve(limit, prime);
        int low = limit;
        int high = 2*limit;
        while (low < n)
        {
                 bool mark[limit + 1];
                 memset(mark, true, sizeof(mark));
                 //mark non-primes in bucket
                 for(int i=0; i<prime.size(); i++)</pre>
                          //get smallest non-prime divisible by prime[i];
                          int loLim = floor(low/prime[i]) * prime[i];
                          if(loLim < low)</pre>
                                  loLim += prime[i];
                          //mark all multiples of prime[i]
                          //print primes
                 for (int i=low; i<high; i++)</pre>
                          if (mark[i - low] == true)
                                  cout << i << " ";
                 // update low and high
                 low += limit;
                 high += limit;
                 if(high >= n) high = n;
        }
```

```
int main()
{
          return 0;
}
```

6 Miscellaneous

6.1 Vimrc file

```
set autoindent
set smartindent
set cindent
set shiftwidth=4
set tabstop=4
set autoread
set cmdheight=1
set number
set splitright
set splitbelow
set makeprg=g++\ -std=c++14\ %\ -o\ out
map <F8>:!./out < input > output <ENTER><map <F7<:make \| cope 8 \| wincmd J <ENTER><map <C-Right> <C-W><Right>
map <C-Left> <C-W><Left>
map <C-Up> <C-W><Up>
map <C-Down> <C-W><Down>
only
40vsp input
20vsp output
wincmd w
```