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# 1. Array

## Find duplicate in the array

### Problem

Given an array of integers, 1 ≤ a[i] ≤ *n* (*n* = size of array), some elements appear **twice** and others appear **once**.

Find all the elements that appear **twice** in this array.

Input:

[4,3,2,7,8,2,3,1]

Output:

[2,3]

### Reference

LEETCODE

### Approach

Since content is with-in array index range. We will use itself array as hashTable and when we found item first time we change sign to negative and if second time we get same negative it is duplicate. If data was not in-range we will use hashmap or set.

### Solution

public List<Integer> findDuplicates(int[] nums) {

List<Integer> list = new ArrayList<>();

for (int i = 0; i < nums.length; i++) {

int n=Math.abs(nums[i])-1;

if (nums[n] < 0) {

list.add(n+1);

} else {

nums[n] = -nums[n];

}

}

return list;

}

### Time and space complexity

Time - O(n)

Space – O(1).In space as list needed only for this question

## Find Numbers with Even Number of Digits

### Problem

Given an array nums of integers, return how many of them contain an even number of digits.

Example 1:

Input: nums = [12,345,2,6,7896]

Output: 2

Explanation:

12 contains 2 digits (even number of digits).

345 contains 3 digits (odd number of digits).

2 contains 1 digit (odd number of digits).

6 contains 1 digit (odd number of digits).

7896 contains 4 digits (even number of digits).

Therefore only 12 and 7896 contain an even number of digits.

### Reference

LEETCODE, MATH

### Approach

Approach 1 - can be to iterate over loop and convert each number to String and then check length is even or odd.

Approach 2 - can be to iterate over loop and use Math.log10 method and then check result%2==0. if it is true it is ODD else EVEN.

### Solution

public int findNumbers(int[] nums) {

int c=0;

for(int i=0;i<nums.length;i++){

int result = (int)Math.log10(nums[i]);

if(result %2!=0){

c++;

}

}

return c;

}

### Time and space complexity

Time - O(n)

Space – O(1)

## Largest Number

### Problem

Given a list of non-negative integers, arrange them such that they form the largest number.

Example 1:

Input: [10,2]

Output: "210"

Example 2:

Input: [3,30,34,5,9]

Output: "9534330"

### Reference

Leetcode, sort, array

### Approach

We need to sort data smartly i.e. write comparator smartly. So, for two string like 3 and 34 to check which one should come first just contact both combo like – 334 and 343. Now we know number larger can be made if 34 comes first and 3 after that.

So, we use above logic and sort the array.

### Solution

public String largestNumber(int[] nums) {

String[] str = new String[nums.length];

for(int i=0;i<nums.length;i++){

str[i]=String.valueOf(nums[i]);

}

Arrays.sort(str,(o1,o2)->{

String s1=o1+o2;

String s2=o2+o1;

return s2.compareTo(s1);

});

if("0".equals(str[0])){

return str[0];

}

StringBuilder sb=new StringBuilder(nums.length);

for(String s:str){

sb.append(s);

}

return sb.toString();

}

### Time and space complexity

Time - O(nlogn)

## Search Insert Position (Binary Search)

### Problem

Given a sorted array and a target value, return the index if the target is found. If not, return the index where it would be if it were inserted in order.

You may assume no duplicates in the array.

Example 1:

Input: [1,3,5,6], 5

Output: 2

Example 2:

Input: [1,3,5,6], 2

Output: 1

### Reference

LEETCODE, BINARY-SEARCH, ARRAY

### Approach

Binary search is best algorithm to search in a sorted array. It takes o(logn) time.

1. set start=0 and end=length-1

2. Iterate till start<=end

3. get mid of (start+end)/2 and check if target is in left or right or in the middle.

4. If target<arr[mid] it means target is present in left. So update end=mid-1.

5. So, by this approach we are dividing the items to be searched to half every time.

In this particular problem if element does not exist. In such case start will tell the position of element where it should supposed to be. In classic binary search if item does not found we return -1

### Solution

public int searchInsert(int[] nums, int target) {

int start = 0;

int end = nums.length - 1;

int mid = 0;

while (start <= end) {

mid = (start + end) / 2;

if (nums[mid] == target) {

return mid;

} else if (nums[mid] < target) {

start = mid + 1;

} else {

end = mid - 1;

}

}

return start;

}

### Time and space complexity

O(logn)

O(1)

## Count Negative Numbers in a Sorted Matrix

### Problem

Given m \* n matrix grid which is sorted in non-increasing order both row-wise and column-wise.

Return the number of negative numbers in grid.

Example 1:

Input: grid = [[4,3,2,-1],[3,2,1,-1],[1,1,-1,-2],[-1,-1,-2,-3]]

Output: 8

Explanation: There are 8 negatives number in the matrix.

### Reference

LEETCODE, ARRAY-2D, BINARY-SEARCH

### Approach

We use binary search algorithm row by row and find center if it’s negative update end=center-1 else start=center+1

When loop terminates start will be the index of first negative. So, total negative in that row is row.length – start.

Since we also know that column is also decreasing so, for second row we update end to start-1.so, that we will apply binary search to only 0 to last positive number in previous row.

And we keep it doing till last row.

### Solution

public int countNegatives(int[][] grid)

{

int c = 0;

for (int i = 0, end = grid[i].length - 1; i < grid.length; i++) {

int start = 0;

while (start <= end) {

int mid = (start + end) / 2;

if (grid[i][mid] < 0) {

end = mid - 1;

} else {

start = mid + 1;

}

}

c = c + grid[i].length - start;

end = start - 1;

}

return c;

}

### Time and space complexity

Time - O(n+m)

Space - O(1)

## Move Zeroes to end of array

### Problem

Given an array nums, write a function to move all 0's to the end of it while maintaining the relative order of the non-zero elements.

Example:

Input: [0,1,0,3,12]

Output: [1,3,12,0,0]

### Reference

LEETCODE, ARRAY

### Approach

Here to make code generic we move val to the end of array.

We will keep count of val in c. and if c> 0 means we have at least one val. We move current element to i-c location. And update arr[i] to val.

It works because we make sure we are shifting non zero element to next available index on left.whic will be i-c.

If we does not have any zero we will not shift.

e.g. – 010004 -> in this case 1 will be shift to zero index .

i.e. 100004. Now c=1 and i=1. So, c keep on incrementing to 4. For i=5,arr[5-4]=arr[5] .

so,op will be 140000.

### Solution

public void searchAndShift(int[] arr, int val) {

int c = 0;

for (int i = 0; i < arr.length; i++) {

if (arr[i] == val) {

c++;

} else if (c > 0) {

arr[i - c] = arr[i];

arr[i] = val;

}

}

}

### Time and space complexity

Time - O(n)

Space – O(1)

## Two Sum

### Problem

Given an array of integers, return indices of the two numbers such that they add up to a specific target.

You may assume that each input would have exactly one solution, and you may not use the same element twice.

Example:

Given nums = [2, 7, 11, 15], target = 9,

Because nums[0] + nums[1] = 2 + 7 = 9,

return [0, 1].

### Reference

LEETCODE, ARRAY, HASHMAP

### Approach

Take hashmap and check if current item is in map if yes return else put (target-current item) in a loop.

### Solution

public int[] twoSum(int[] numbers, int target) {

Map<Integer, Integer> map = new HashMap<>();

for(int i = 0; i < numbers.length; i++) {

if(map.containsKey(numbers[i])) {

return new int[] {map.get(numbers[i]), i};

}

map.put(target - numbers[i], i);

}

return null;

}

### Time and space complexity

Time - O(n)

Space – O(n) (hashmap)

# 2. String

## Check if a string is substring of source. (Rabin Karp Algorithm)

### Problem

Check whether a given string is substring of source.

Example 1-

Input –

helloji,loj

Output –

True

### Reference

STRING, RABIN-KARP, ABDUL BARI

### Approach

Naive approach is to check character by character starting from i=0 and if not matched go back and now check for i=1.

Better approach -

It uses hashcode of a string and instead of matching character one by one.

\* We just match hashcode and once hashcode matched we check the content.

\* If not matched we subtract hashcode of first character and add hashcode of new character

\* It saves time of un-necessary comparison all the time.

\* But in worst case it might be possible that we might get hashcode same on every check.

\* To calculate hashcode again we just subtract hashcode of first character and add hashcode of next character in previous value.

\* For better performance make hash code function better to avoid un-necessary collision.

### Solution

int hSource = 0;

int hStr = 0;

//calculate hashcode of both source and string for first comparison

for (int i = 0; i < str.length(); i++) {

hSource = hSource + hashFunction(source.charAt(i));

hStr = hStr + hashFunction(str.charAt(i));

}

// we compare hash first and if matched return true.

// calculate hash again except for last value of i as we are generating hash in

// advanced.

for (int i = 0; i <= source.length() - str.length(); i++) {

if (hStr == hSource) {

int j = 0;

for (j = 0; j < str.length(); j++) {

if (source.charAt(j + i) != str.charAt(j)) {

break;

}

}

if (j == str.length()) {

return true;

}

}

//to avoid calculation after last index

if (i < source.length() - str.length()) {

hSource = hSource - hashFunction(source.charAt(i)) + hashFunction(source.charAt(i + str.length()));

}

}

return false;

### Time and space complexity

\* worst case - o(n\*m)

\* Best case - o(n+m)

## Anagram

### Problem

An anagram is a word formed by rearranging the letters of a different word. typically using all the original letters exactly once.

Given two strings s and t, write a function to determine if t is an anagram of s.

You may assume the string contains only lowercase alphabets.

Example 1:

Input: s = "anagram", t = "nagaram"

Output: true

Example 2:

Input: s = "rat", t = "car"

Output: false

### Reference

LEETCODE, ARRAY

### Approach

Take array with 26 size and from first string increment counter and for second decrement counter.

After loop finished iterate over table array and check if any non-zero value exists it’s not anagram.

### Solution

public boolean isAnagram(String s, String t) {

int[] table=new int[26];

if(s.length()!=t.length()){

return false;

}

for(int i=0;i<s.length();i++){

table[s.charAt(i)-'a']+=1;

table[t.charAt(i)-'a']-=1;

}

for(int i=0;i<table.length;i++){

if(table[i]!=0){

return false;

}

}

return true;

}

### Time and space complexity

Time - O(n+26)

Space – O(26) means constant

# 3. Math

## 1. Palindrome Number

### Problem

Determine whether an integer is a palindrome. An integer is a palindrome when it reads the same backward as forward.

Example 1:

Input: 121

Output: true

Example 2:

Input: -121

Output: false

Explanation: From left to right, it reads -121. From right to left, it becomes 121-

### Reference

LEETCODE, MATH, MOD

### Approach

Reverse the original number by adding remainder to the original number – res\*10+(num%10);

### Solution

public boolean isPalindrome(int num) {

int res = 0;

int num1 = num;

while (num > 0) {

res = res \* 10 + (num % 10);

num = num / 10;

}

return num1 == res;

}

### Time and space complexity

Time - o(n)

Space – o(1)

# 4. Dynamic Programming

## 1. Maximum sum in Contiguous Sub-Array

### Problem

Given an integer array nums, find the contiguous subarray (containing at least one number) which has the largest sum and return its sum.

Example:

Input: [-2,1,-3,4,-1,2,1,-5,4],

Output: 6

Explanation: [4,-1,2,1] has the largest sum = 6.

### Reference

LEETCODE, UDEMY, DP

### Approach

\*Take global\_max which hold the max overall

\* And curr\_max will hold the max till curr iteration.

\* We will update curr\_max by this - store max of (current element, curr\_max+current element)

\* By this we make sure that either current is taken or previous one is included in contiguous space.

### Solution

public int maxSubArray(int[] nums) {

int curr\_max = nums[0];

int global\_max = nums[0];

for (int i = 1; i < nums.length; i++) {

curr\_max = MathUtil.max(nums[i], nums[i] + curr\_max);

if (curr\_max > global\_max) {

global\_max = curr\_max;

}

}

return curr\_max;

}

### Time and space complexity

O(n)

O(n)