

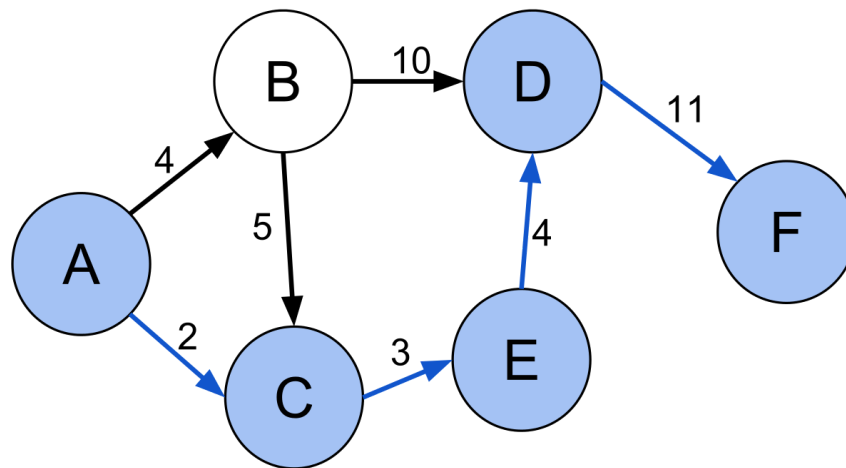
Lecture 19: Finding Shortest Path

Date: 10/24/2023

Shortest Path

A shortest path from **vertex s** to **vertex t** in an edge-weighted graph is a directed path from s to t with the property that no other such path has a lower weight.

⇒ One of the most intuitive use cases of graphs is to find the shortest path from one node to another.



reference: https://en.wikipedia.org/wiki/Shortest_path_problem

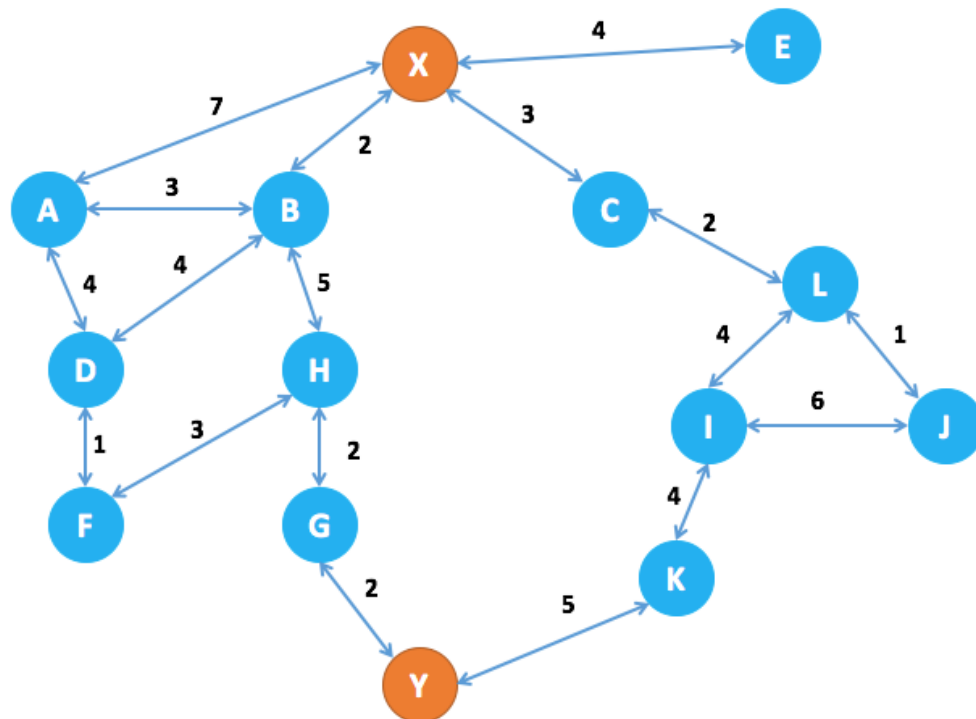
Applications

application	vertex	edge
<i>map</i>	intersection	road
<i>network</i>	router	connection
<i>schedule</i>	job	precedence constraint
<i>arbitrage</i>	currency	exchange rate

Typical shortest-paths applications

Types

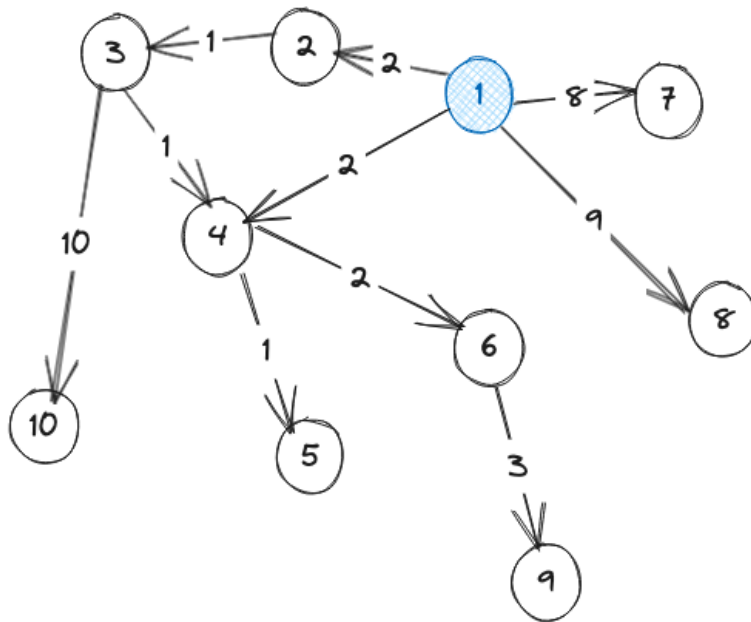
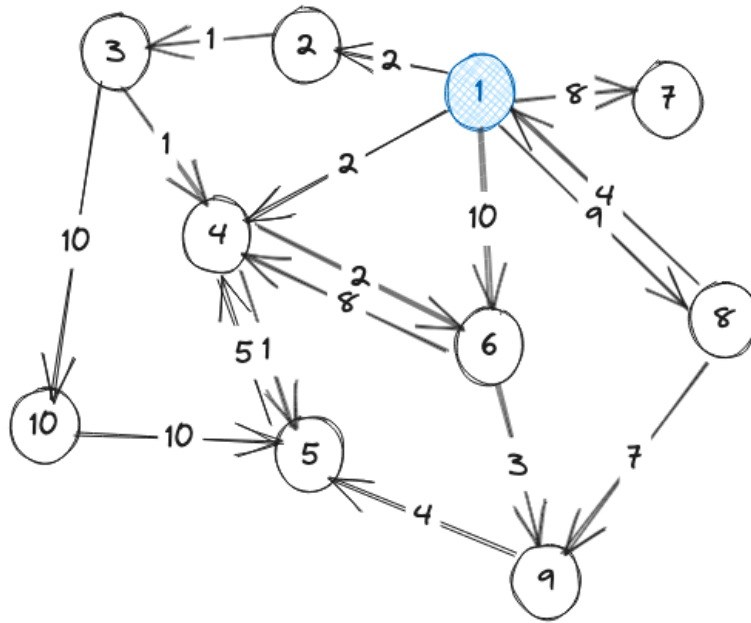
⇒ Single Source Single Sink: find path between source and sink



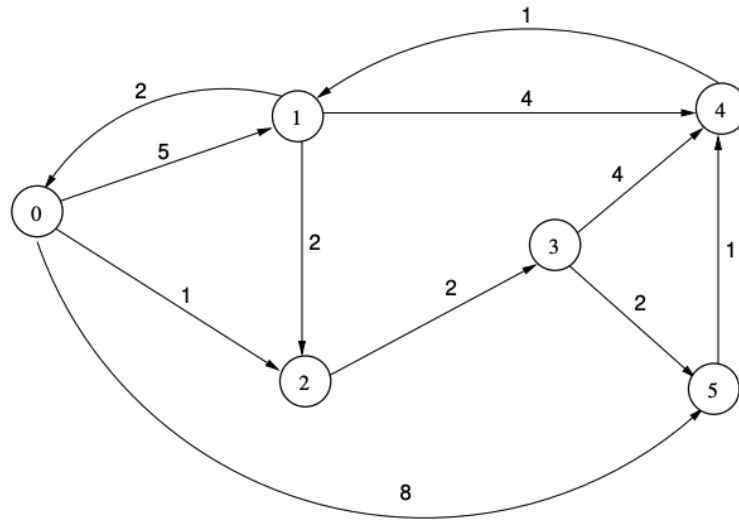
reference:

<https://benalexkeen.com/implementing-dijkstras-shortest-path-algorithm-with-python/>

⇒ Shortest-paths Tree: given source, find shortest path to all other nodes



⇒ All Pair: find shortest path from all nodes to all other nodes

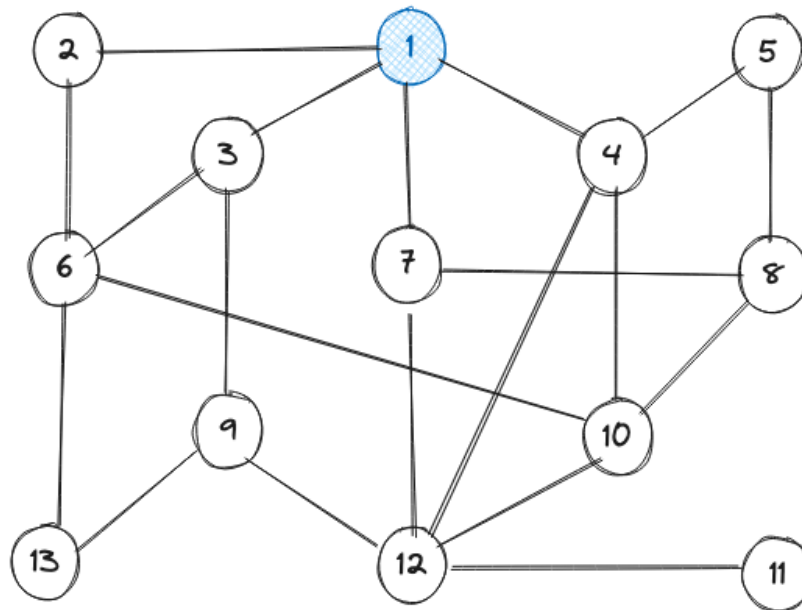


	0	1	2	3	4	5
0	0	5	1	3	6	5
1	2	0	2	4	4	6
2	8	6	0	2	5	4
3	6	4	6	0	3	2
4	3	1	3	5	0	7
5	4	2	4	6	1	0

Table 2: The shortest-paths matrix for the graph in Figure 1

reference: <https://codeahoy.com/learn/graphalgorithms/ch8/>

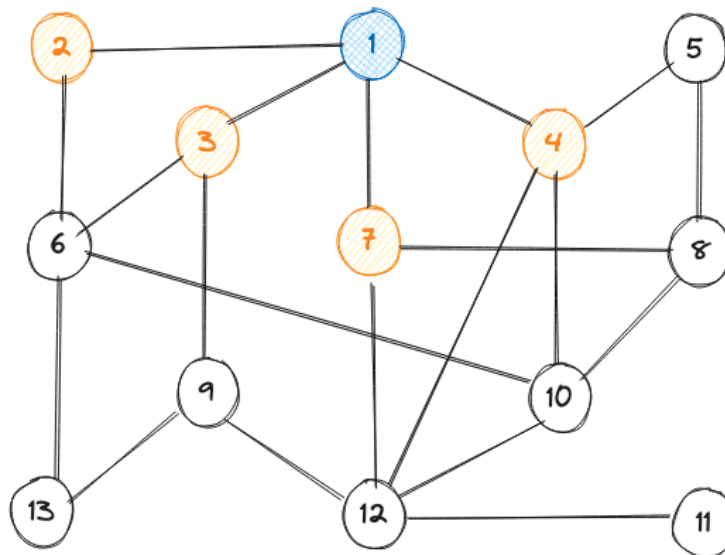
Unweighted Graph Shortest Path



If the graph is unweighted then we can assume each edge weight as one and find the shortest path using that new graph.

Breadth First Traversal

⇒ This traversal starts from the source node and in each iteration finds its first connected node.



⇒ After the first iteration, we found the shortest path from source node 1 to nodes 2,3,7,4.

⇒ The same way, in each iteration we find shortest paths to their immediately connected nodes.

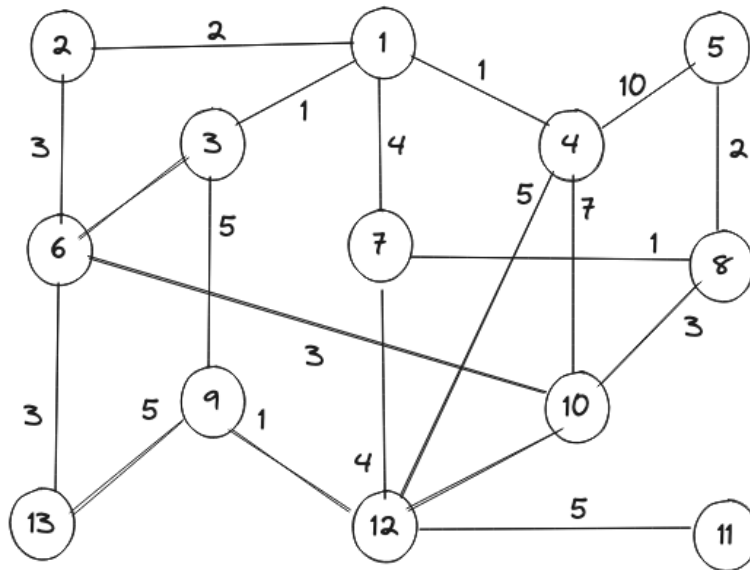
⇒ At the end of the algorithm, we will have shortest distances for all the nodes from the source node.

Pseudocode

```
function find_shortest_path(graph, source){
    shortest_paths = {}
    level = {source}
    while level is not empty:
        // iterate through all nodes in level
        for node in level:
            // iterate through all connected nodes
            for connected in node->connected:
                // update the shortest distance
                shortest_path[connected] = shortest_path[node] + 1
                add connected node to new_level

        update level to new level
    }
```

Weighted Graph Shortest Path



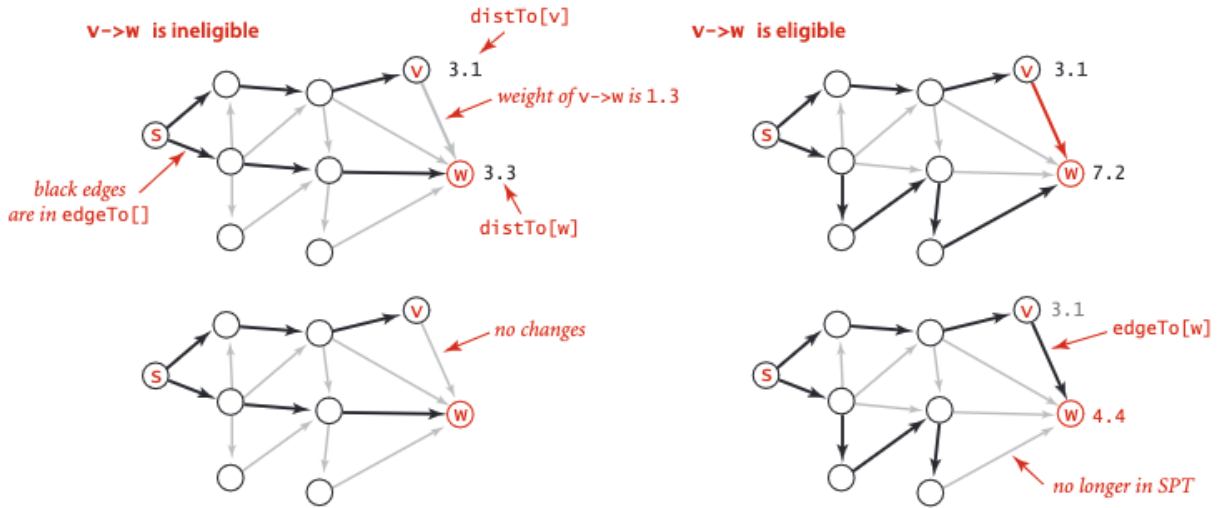
⇒ Using Breadth First Traversal won't work here as it does not take edge weight into consideration.

i.e. From node 1 to node 5, BFS will give 1->4->5 as the shortest path. Which is **not** the shortest. Correct shortest path will be, 1->7->8->5

Relaxation

Edge Relaxation

⇒ To relax an edge e , from v to w , means to figure out if the best way to reach from source to w is from using edge $v \rightarrow w$.



Edge relaxation (two cases)

```
private void relax(DirectedEdge e)
{
    int v = e.from(), w = e.to();
    if (distTo[w] > distTo[v] + e.weight())
    {
        distTo[w] = distTo[v] + e.weight();
        edgeTo[w] = e;
    }
}
```

Vertex Relaxation

⇒ Relax every edge going out from a specific vertex.

```
private void relax(EdgeWeightedDigraph G, int v)
{
    for (DirectedEdge e : G.adj(v))
    {
        int w = e.to();
        if (distTo[w] > distTo[v] + e.weight())
        {
            distTo[w] = distTo[v] + e.weight();
            edgeTo[w] = e;
        }
    }
}
```


Dijkstra's

⇒ Until we don't have every node in the shortest path tree, find the minimum distance node from source and relax it.

⇒ Similar to Prim's algorithm, greedy.

Pseudocode

```
function find_shortest_path(graph, source){  
    priority_queue = {}  
    initialize all distances to Math.inf except source  
  
    priority_queue.insert(source)  
    while priority_queue not empty  
        relax_vertex(priority_queue.getMin())  
}
```

Problems