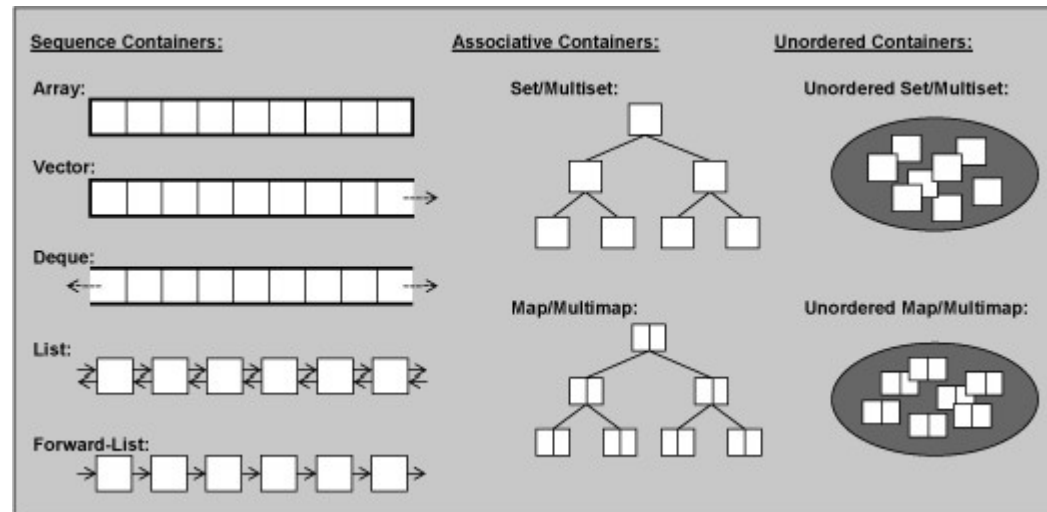


# CPSC 131 – Data Structures

## Array & Vector Abstract Data Types



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# Array Abstract Data Type

## Definitions:

- |          |                                                                         |
|----------|-------------------------------------------------------------------------|
| Capacity | - max number of elements that can be stored                             |
| Size     | - another name for Capacity – an array's size does not (can not) change |

## Fixed Capacity Array

- Capacity is constant
  - Set at container definition at design (compile) time

## Two flavors

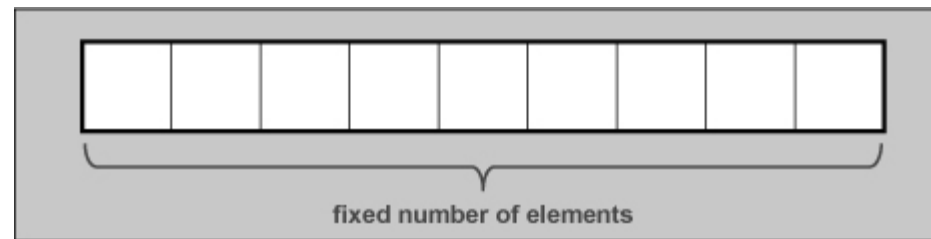
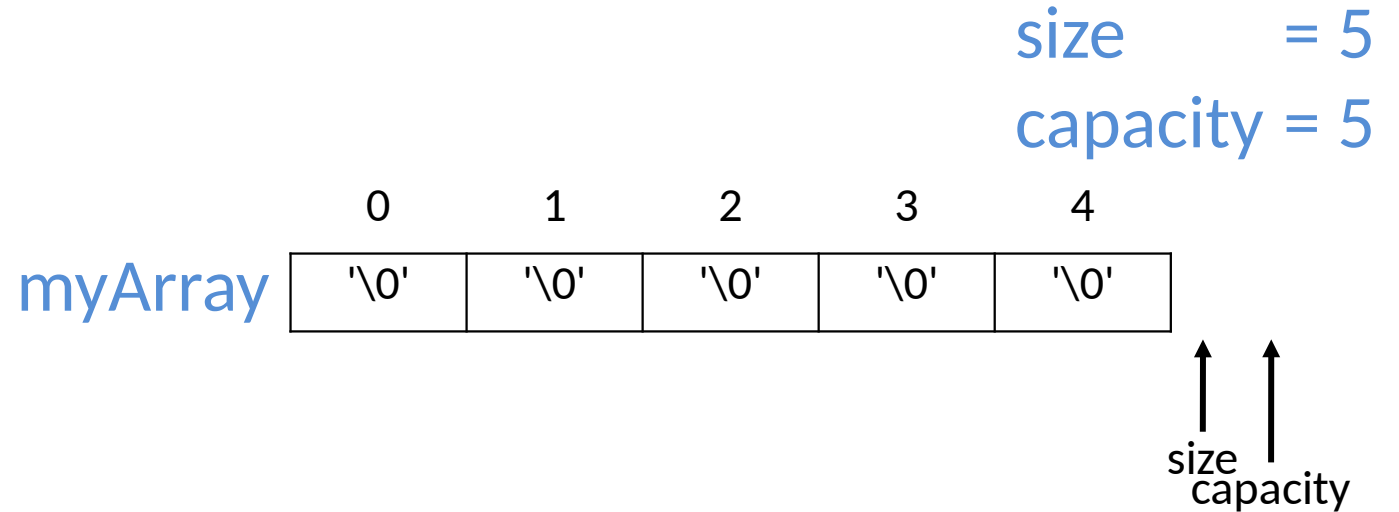
- Standard arrays
  - Smart wrapper around native array
  - `std::array` from the `<array>` library
  - Ex: `std::array<Student, 10> myArray;`
- Native arrays
  - aka C-Style or raw array
  - AVOID using these
  - Ex: `Student myArray[10];`

# Array Abstract Data Type

## The Abstraction - What can I do to an Array

[`std::array`](#)

- Construct, destruct, assign
- Copy, compare
- Iterate
- Access elements
  - `at`, `operator[]`, `front`, `back`
- Query
  - `empty`, `size`
- Operations
  - ----



# Vector Abstract Data Type

## Definitions:

- |          |                                             |
|----------|---------------------------------------------|
| Capacity | - max number of elements that can be stored |
| Size     | - number of elements that are stored        |

## Fixed Capacity Vector

- Capacity is constant
  - Set at container construction during runtime, or
  - Set at container definition at design (compile) time

## Extendable Capacity Vector

- Capacity is dynamic and changes during runtime
  - Initialized at container construction during runtime
  - Grows and shrinks during runtime



# Vector Abstract Data Type

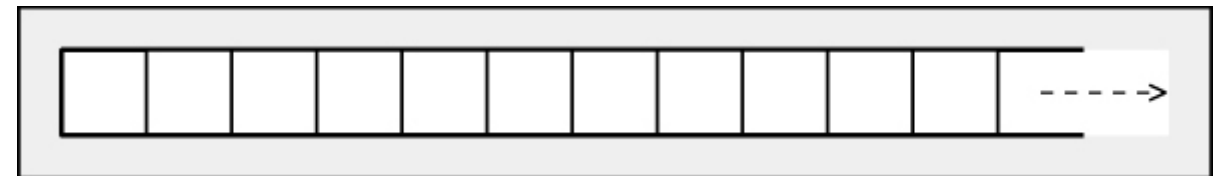
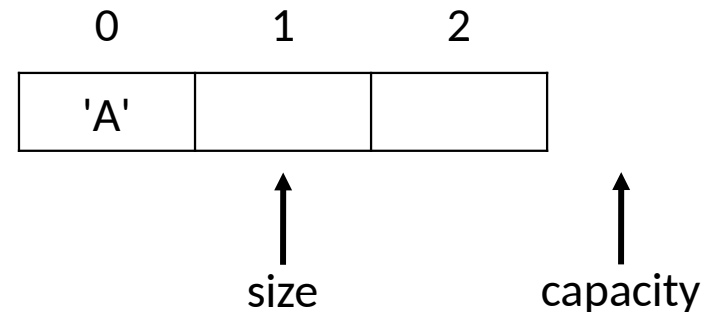
## The Abstraction - What can I do to a Vector

### Vector

- Construct, destruct, assign
- Copy, compare
- Iterate
- Access elements
  - at, operator[], front, back
- Query
  - empty, size, capacity
- Operations
  - insert, erase, clear, push\_back, pop\_back

size = 1  
capacity = 3

myVector



# Array <sup>vs</sup> Vector – What's the difference?

Arrays	Vectors
Capacity is constant	Two flavors, Fixed and Extendable Capacity
Size is constant	Size changes
Capacity and Size are always the same	Capacity and Size usually differ
Every cell always contains an element	Some cells do not contain an element
No insert and erase operations	Elements can be inserted and erased
Two template parameters	One template parameter



# Array <sup>vs</sup> Vector – What's the same?

## Arrays and Vectors

Consecutive locations in memory

Indexed the same way

operator[] may over index  
avoidance is client's responsibility

Comparison, assignment, initialization



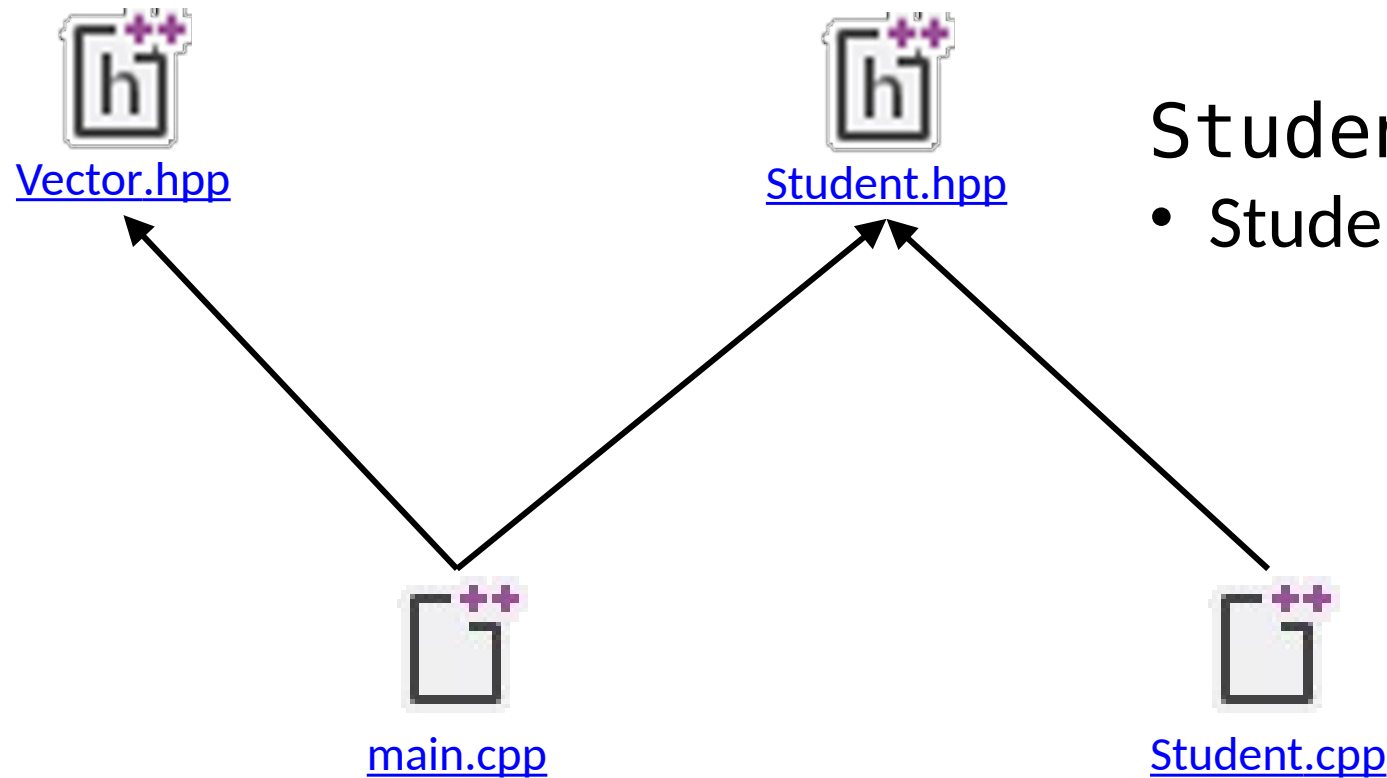
# Vector Implementation Example

main.cpp includes both

- Vector.hpp and Student.hpp

Student.cpp includes only

- Student.hpp





# In Class Sketching Activity

Using the Implementation Examples, for each of the major operations of Arrays and Vectors, step through the code and sketch the resulting structure

# Analysis of the Vector Abstract Data Type

## Complexity Analysis (1)

Function	Analysis – <code>std::array&lt;T, S&gt;</code>	Analysis – <code>std::vector&lt;T&gt;</code> (Extendable Vector)
<code>at()</code>	$O(1)$ Elements directly indexable	same
<code>size()</code>	$O(1)$ Always returns $S$ , as in <code>std::array&lt;T, S&gt;</code>	$O(1)$ Returns the number of elements held
<code>empty()</code>	$O(1)$	same
<code>clear()</code>	Not available <code>std::array&lt;T, S&gt;</code> will always have $S$ elements	$O(n)$ All elements are destroyed and size set to zero <i><math>O(1)</math> if only size set to zero, as in zyBook</i>

# Analysis of the Vector Abstract Data Type

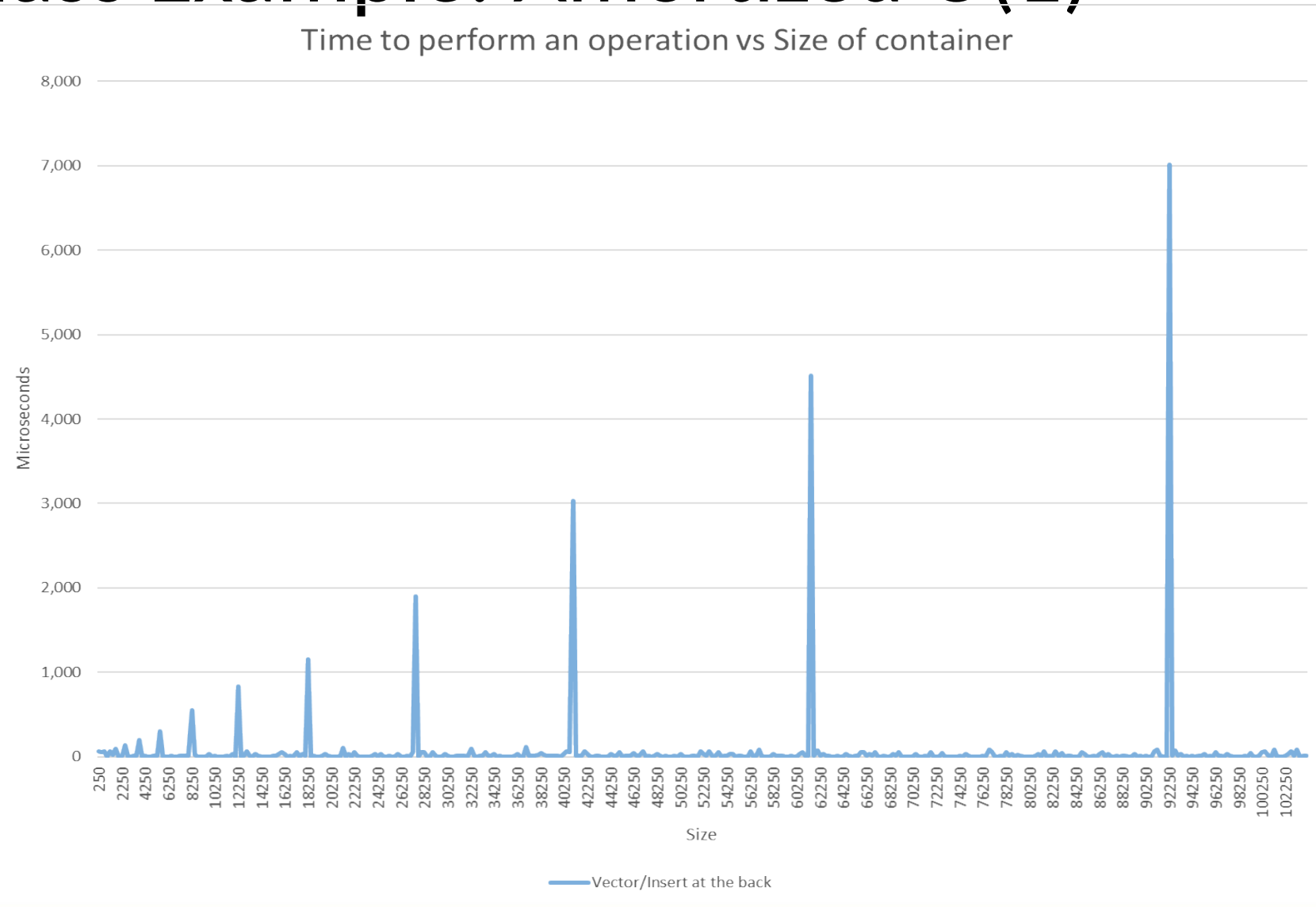
## Complexity Analysis (2)

Function	Analysis – <code>std::array&lt;T, S&gt;</code>	Analysis – <code>std::vector&lt;T&gt;</code> (Extendable Vector)
<code>push_back()</code>	Not available <code>std::array&lt;T, S&gt;</code> will always have S elements	$O(n)$ amortized to $O(1)$ Special case of <code>insert()</code>
<code>erase()</code>	Not available <code>std::array&lt;T, S&gt;</code> will always have S elements	$O(n)$ Have to “close the gap” which means N copies (worst case, $N/2$ copies average case)
<code>splice</code>	Not available	Not available

# Efficiency Class Example: Amortized $O(1)$

vector.push\_back(data)

- The “norm” is constant time
- But every now and then consumes linear time
- The interval between spikes doubles each time
- The severity of the spike increases each time



# Analysis of the Vector Abstract Data Type

## Complexity Analysis (3)

Function	Analysis – <code>std::array&lt;T, S&gt;</code>	Analysis – <code>std::vector&lt;T&gt;</code> (Extendable Vector)
<code>insert()</code>	Not available <code>std::array&lt;T,S&gt;</code> will always have S elements	O(n) (worst case) If space is not available, <ul style="list-style-type: none"> <li>• get more space and copy N elements</li> <li>• Destroy N elements</li> </ul> “Open a gap” which means N copies
default construction	O(n) container is never empty	O(1) creates an empty container
Equality $C_1 == C_2$	O(n)	same

# Analysis of the Vector Abstract Data Type

## Complexity Analysis (4)

Function	Analysis – <code>std::array&lt;T, S&gt;</code>	Analysis – <code>std::vector&lt;T&gt;</code> (Extendable Vector)
<code>push_front</code>	Not available <code>std::array&lt;T, S&gt;</code> will always have $S$ elements	Not available
<code>resize</code>	Not available <code>std::array&lt;T, S&gt;</code> will always have $S$ elements	$O(n)$
<code>find</code>	$O(n)$ linear search from <code>begin()</code> to <code>end()</code> (i.e. $a[0]$ to $a[size()-1]$ )	same

# Analysis of the Vector Abstract Data Type

## Complexity Analysis (4)

Function	Analysis – <code>std::array&lt;T, S&gt;</code>	Analysis – <code>std::vector&lt;T&gt;</code> (Extendable Vector)
Visit every element e.g. <code>print()</code>	$O(n)$ Visiting every node from <code>begin()</code> to <code>end()</code>	same
Visit in reverse e.g. <code>print_reverse()</code>	$O(n)$ Visiting every node from <code>rbegin()</code> to <code>rend()</code> Direction doesn't matter	same