# Introduction to High-Performance Computing

Using clusters to speed up your research

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- ✓ slides and data files at http://bit.ly/introhpc
  - ► the link will download a file introHPC.zip (~3 MB)
  - ▶ unpack it to find codes/ and slides.pdf

#### Workshop outline

- To work on the remote cluster today, you will need
  - 1. Wi-Fi access (Eduroam? Campus guest account?)
  - 2. an SSH client
    - pre-installed on Linux/Mac
    - http://mobaxterm.mobatek.net Home Edition on Windows
  - 3. only on the production clusters:
    - a Compute Canada account http://bit.ly/2K0ol4S (approved by the PI) or a guest account
    - be added to the Slurm reservation
- Cluster hardware overview
- Basic tools for cluster computing
  - ► logging in, transferring files
  - software environment, modules
  - ► Linux command line, editing remote files
- Programming languages and tools
  - overview of languages from HPC standpoint
  - parallel programming environments
  - compilers
  - ▶ quick look at OpenMP, MPI, Chapel, make

- Working with Slurm scheduler
- Debugging and very briefly on profiling
- Best practices (common mistakes)

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# Hardware overview

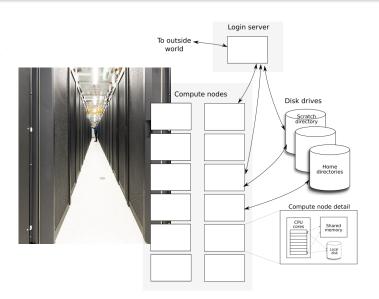
# Added to CC's national infrastructure in the past several years

- Arbutus @UVic is an extension to West Cloud (in production since September 2016)
- General-purpose clusters: Cedar @SFU and Graham @UofWaterloo (both in production since June 2017), Béluga @McGill (in production since April 2019)
- Large parallel cluster: Niagara @UofToronto (in production since April 2018)





- Mostly off-the-shelf components for individual nodes, everything rack-mounted
- Typically hundreds of nodes, wired by fast interconnect
- Shared vs. distributed memory
- Login vs. compute nodes
- Compute nodes: CPU-only, GPU nodes (accelerators)
- Job scheduler
- Development/visualization nodes



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	cedar.computecanada.ca	graham.computecanada.ca	beluga.computecanada.ca				
purpose	general-purpose cluster for a variety of workloads						
specs	https://docs.computecanada.ca/ wiki/Cedar	https://docs.computecanada.ca/ wiki/Graham	https://docs.computecanada.ca/ wiki/Beluga				
processor count	58,416 CPUs and 584 GPUs	35,520 CPUs and 320 GPUs	34,880 CPUs and 688 GPUs				
interconnect	100Gbps Intel OmniPath, non-blocking to 1024 cores	100Gbps Mellanox InfiniBand, non-blocking to 1024 cores	InfiniBand 100Gbps or better				
base nodes	<b>576 nodes</b> : 128GB, 32 cores <b>640</b> <i>SL</i> <b>nodes</b> : 192GB, 48 cores	<b>864 nodes</b> : 128GB, 32 cores	<b>172</b> <i>SL</i> <b>nodes</b> : 96GB, 40 cores <b>516</b> <i>SL</i> <b>nodes</b> : 192GB, 40 cores				
large-memory nodes	128 nodes: 256GB, 32 cores 24 nodes: 512GB, 32 cores 24 nodes: 1.5TB, 32 cores 4 nodes: 3TB, 32 cores	<b>56 nodes</b> : 256GB, 32 cores <b>24 nodes</b> : 512GB, 32 cores <b>3 nodes</b> : 3TB, 64 cores	<b>12</b> <i>SL</i> <b>nodes</b> : 768GB, 40 cores				
GPU base	114 nodes: 128GB, 24 cores, 4 NVIDIA P100 Pascal 12GB GPUs 32 nodes: 256GB, 24 cores, 4 NVIDIA P100 Pascal 16GB GPUs	160 nodes: 128GB, 32 cores, 2 NVIDIA P100 Pascal 12GB GPUs	172 nodes: 192GB, 40 cores, 4 NVIDIA Volta V100 16GB GPUs				

	niagara.computecanada.ca
purpose	for large parallel jobs, ideally $\geq$ 1,000 cores with an allocation
specs	https://docs.computecanada.ca/wiki/Niagara and https://docs.scinet.utoronto.ca
processor count	60,000 CPUs and no GPUs
interconnect	EDR Infiniband (Dragonfly+, completely connected topology, dynamic routing), 1:1 to 432 nodes, effectively 2:1 beyond that
base nodes	<b>1,500</b> <i>SL</i> <b>nodes</b> : 192GB, 40 cores

- ▶ No local disk, nodes booting off the network, small RAM filesystem
- All cores are Intel Skylake (2.4 GHz, AVX512)  $\Rightarrow$  SL = Skylake
- Authentication via CC accounts; for now, need to request a SciNet account; long-term regular access for all CC account holders
- **Scheduling is by node** (in multiples of 40 cores)

- Users with an allocation: job sizes up to 1000 nodes and 24h max runtime
- Users without an allocation: job sizes up to 20 nodes and 12h max runtime
- Maximum number of jobs per user: running 50, queued 150

# Accessing resources: RAS vs. RAC

- ~15% of compute cycles available via the Rapid Access Service (RAS)
  - ► available to all CC users via default queues
  - you can start using it as soon as you have a CC account
  - ► shared pool with resources allocated via "fair share" mechanism
  - will be sufficient to meet computing needs of many research groups
- ~85% of compute cycles allocated via annual Resource Allocation Competitions (RAC)
  - ► apply if you need >50 CPU-years or >10 GPU-years
  - only PIs can apply, allocation per research group
  - ▶ announcement in the fall of each year via email to all users
  - ▶ 2019 RAC: 507 applications, success rate (awarded vs. requested): 41% of CPUs, 20% of GPUs, 86% of storage, 95% of virtual CPUs

#### File systems

Overview

Details at https://docs.computecanada.ca/wiki/Storage\_and\_file\_management

filesystem	quotas	backed up?	purged?	performance	mounted on compute nodes?
\$HOME	50GB, 5e5 files per user, 100GB Niagara	nightly, latest snapshot	no	medium	yes
\$SCRATCH	20TB Cedar, 100TB Graham, 25TB Niagara, 1e6 files per user except when full	no	yes	high for large files	yes
/project (long-term disk storage)	1TB, 5e5 files per user could be increased via RAC, Niagara only via RAC	nightly	no	medium	yes
/nearline (tape archive)	5TB per group via RAC	no	no	medium to low	no
/localscratch	none	no	maybe	very high	local

- Wide range of options from high-speed temporary storage to different kinds of long-term storage
- On Cedar no longer allowed to submit jobs from \$HOME
- To check disk usage: quota command (aliased to diskusage\_report)
- To request more storage: support@computecanada.ca for small increases, RAC for large requests

# Basic tools for working with a cluster

#### Logging into the systems

• On Mac or Linux in terminal:

```
$ ssh userXXX@206.12.90.125  # guest account on our training cluster (where XX=01..60)
$ ssh yourUsername@cedar.computecanada.ca  # use your CC account on Cedar
$ ssh yourUsername@graham.computecanada.ca  # use your CC account on Graham
```

- On Windows many options:
  - ► MobaXTerm https://docs.computecanada.ca/wiki/Connecting\_with\_MobaXTerm
  - ► PuTTY https://docs.computecanada.ca/wiki/Connecting\_with\_PuTTY
  - ► Secure Shell Extension http://bit.ly/2FeLpOR in Chrome browser
  - ▶ bash from the Windows Subsystem for Linux (WSL) in Win10 need to enable developer mode and then WSL
- SSH key pairs are very handy, save you from typing passwords
  - ► implies secure handling of private keys, non-empty passphrases
  - ► https://docs.computecanada.ca/wiki/SSH\_Keys
  - ► https://docs.computecanada.ca/wiki/Using\_SSH\_keys\_in\_Linux
  - https://docs.computecanada.ca/wiki/Generating\_SSH\_keys\_in\_Windows
- GUI connection: X11 forwarding (through ssh), VNC, x2go
- Client-server workflow in selected applications, both on login and compute nodes

#### Linux command line

- All our systems run Linux (CentOS 7)  $\Rightarrow$  you need to know basic command line
  - ► separate 3-hour "Bash Scripting and Tools" session
  - $\blacksquare \ \ \, attend\ a\ Software\ Carpentry\ \textit{bash}\ session\ \texttt{https://software-carpentry.org/workshops} \\$
  - ▶ lots of tutorials online, e.g. tutorials 1 4 at http://bit.ly/2vH3j8v
- ullet Much typing can be avoided by using bash aliases, functions,  $\sim$ /.bashrc, hitting TAB

# FILE COMMANDS Is directory listing Is-aIF pass command arguments cd dir change directory to dir cd change to home pwd show current directory mkdir dir create a directory rm file delete file rm-r dir delete directory rm file force remove file rm-rf dir force remove directory cp file target copy file to target mv file target rename or move file to target In-s file link create symbolic link to file touch file create or update file

 $\frac{\text{PATHS}}{\text{relative vs. absolute paths}}$  meaning of  $\sim$  . ...

#### FILE COMMANDS

command > file redirect command output to file command ≫ file append command output to file more file page through contents of file cat file print all contents of file head -n file output the first n lines of file tail -n file output the contents of file tail -f file output the contents of file as it grows

#### PROCESS MANAGEMENT

top display your currently active processes
ps display all running processes
kill pid kill process ID pid

#### FILE PERMISSIONS

chmod -R u+rw,g-rwx,o-rwx file set permissions

#### ENVIRONMENT VARIABLES AND ALIASES

export VAR='value' set a variable
echo \$VAR print a variable
alias Is='Is -aFh' set an alias command

#### SEARCHING

grep pattern files search for pattern in files command | grep pattern example of a pipe find . -name '\*.txt' | wc -l another pipe

#### OTHER TOOLS

man command show the manual for command command --help get quick help on command df -kh. show disk usage du -kh. show directory space usage

#### COMPRESSION AND ARCHIVING tar cyf file tar files create a tar file

tar xvf file.tar extract from file.tar gzip file compress file gunzip file.gz uncompress file.gz

#### LOOPS

for i in \*tex; do wc -1 \$i; done loop example

# Bash walk-through on the cluster

Connect to the cluster

Basics 0000000000

- 2. Can you tell the difference between local and remote shells?
- Run whoami and hostname
- Bring up a manual page on some command, also try --help flag
- 5. Check files in your home directory: are there any hidden files? lists files by time of last change
- 6. Check file permissions
- 7. Check out different filesystems \$HOME, \$SCRATCH, \$PROJECT - what are their paths?
- 8. Play with paths, try to use both absolute and relative paths, use special characters  $\cdot \cdot \cdot \sim$
- 9. Create a directory, put a file with some contents into it (with a text editor), look at this file from the command line with more and cat
- 10. Copy this file into another file, try moving files, delete a file, delete a directory
- 11. Create several files, put them into a gzipped tar archive

- 12. Move this archive into another directory, unpack it
- 13. Create a new directory, download http://hpc-carpentry.github.io/ hpc-intro/files/bash-lesson.tar.gz into it and unpack it there
- 14. How much space do the unpacked files take?
- 15. Count the number of lines in all \*.fastq files
- 16. Try to redirect output from the last command into a file
- 17. Try pipes: construct a one-line command to display the name of the longest (by the number of lines) file
- 18. Search inside files with grep
- 19. Find files with find
- 20. Write and run a bash script (need to start with #!/bin/bash - called *shebang*) making it executable
- 21. Print out some shell variables
- Write a quick loop to display each \* . fastq file's name, its number of lines, and then its first two lines, separating individual files with an empty line
- 23. Write a quick loop to remove SRR from each \* . fastq file's name (use \$ {name: 3:14} syntax)

#### Editing remote files from the command line

- For novice users: nano (easiest option)
- For power users: emacs-nw (super configurable, lots of shortcuts)
- For die-hard fans: vi and vim (basic, difficult to use)
- Remote graphical emacs not recommended
  - ▶ you would connect via ssh with an X11 forwarding flag (-X, -Y)

Best practices

#### diting remote mes (cont.)

- My favourite option: local emacs on my laptop editing remote files via ssh with emacs's built-in package tramp
  - ightharpoonup need to add to your  $\sim$ /.emacs

```
(require 'tramp)
(setq tramp-default-method "ssh")
```

only makes sense with a working ssh-key pair

▶ your private key is your key to the cluster, so don't share it!

# Cluster software environment at a glance

- Programming languages: C/C++, Fortran 90, Python, R, Java, Matlab, Chapel several different versions and flavours for most of these
- **CPU** parallel development support: MPI, OpenMP, Chapel
- GPU parallel development support: CUDA, OpenCL, OpenACC
- Job scheduler: Slurm open-source scheduler and resource manager
- Popular software: installed by staff, listed at https://docs.computecanada.ca/wiki/Available\_software
  - ▶ lower-level, not performance sensitive packages installed via Nix package manager
  - general packages installed via EasyBuild framework
  - everything located under /cvmfs, loaded via modules (next slide)
- Other software
  - ► email *support@computecanada.ca* with your request, or
  - ► can compile in your own space (feel free to ask staff for help)

#### Software modules

 Use appropriate modules to load centrally-installed software (might have to select the right version)

```
$ module avail <name>  # search for a module (if listed)
$ module spider <name>  # will give a little bit more info
$ module list  # show currently loaded modules
$ module load moduleName
$ module unload moduleName
```

# show commands in the module

- All associated prerequisite modules will be automatically loaded as well
- Modules must be loaded before a job using them is submitted

\$ module show moduleName

▶ alternatively, can load a module from the job submission script

#### File transfer

In Mac/Linux terminal or in Windows MobaXterm or bash/WSL you have two good options:

- use scp to copy individual files and directories
  - \$ scp filename yourUsername@cedar.computecanada.ca:/path/to
    \$ scp yourUsername@cedar.computecanada.ca:/path/to/filename localPath
- use interactive sftp
- use rsync to sync files or directories

```
$ flags='-av --progress --delete'
```

- \$ rsync \$flags localPath/\*pattern\* yourUsername@cedar.computecanada.ca:/path/to
- \$ rsync \$flags yourUsername@cedar.computecanada.ca:/path/to/\*pattern\* localPath

Windows MobaXTerm allows GUI drag-and-drop Windows PuTTY uses pscp command for secure file transfer

#### Globus file transfer

Details at https://docs.computecanada.ca/wiki/Globus

- The CC Globus Portal https://globus.computecanada.ca is a fast, reliable, and secure service for big data transfer (log in with your CC account)
- Easy-to-use web interface to automate file transfers between any two endpoints
  - ▶ an *endpoint* could be a CC system, another supercomputing facility, a campus cluster, a lab server, a personal laptop (requires Globus Connect Personal app)
  - runs in the background: initialize transfer and close the browser, it'll email status
- Uses GridFTP transfer protocol: much better performance than scp, rsync
  - ▶ achieves better use of bandwidth with multiple simultaneous TCP streams
  - ▶ some ISPs and Eduroam might not always play well with the protocol (throttling)
- Automatically restarts interrupted transfers, retries failures, checks file integrity, handles recovery from faults
- Command-line interface available as well

#### ligh-level overview of programming models

- Installed serial compilers and interpreters: C, C++, Fortran, Python, R, Java
- For more details see https://docs.computecanada.ca/wiki/Programming\_Guide

#### In HPC speed matters. Not all languages are built equal in terms of performance!

- Native loops and arithmetic in Python are 80-200X slower than optimized compiled C/C++/Fortran
  - ▶ Python compilers and accelerators (Cython, Numba, Nuitka, etc.) try to improve things to some extent
  - use precompiled numerical libraries such as numpy and scipy
  - ► call C/C++/Fortran functions from Python
- Native R is even slower ... designed for desktop-scale statistical computation and graphics
  - very popular in engineering, mathematics, statistics, bioinformatics
  - ▶ there are some ways to accelerate R to an "acceptable balance" of coding time investment vs. performance
- Matlab, Java are also slowish (3-10X compared to optimized compiled C/C++/Fortran)
- Later (when submitting serial jobs) we'll do an exercise to time pi.c vs. pi.py
  - ► you can find both codes in codes / directory
  - ightharpoonup make sure to use the same value of n

# Parallel programming environment

#### CPU parallel development support: OpenMP (since 1997), MPI (since 1994)

- ▶ OpenMP is a language extension for C/C++/Fortran provided by compilers, implements shared-memory parallel programming
- ► MPI is a library with implementations for C/C++/Fortran/Python/R/etc, designed for distributed-memory parallel environments, also works for CPU cores with access to common shared memory
- ► industry standards for the past 20+ years

#### Chapel is a open-source parallel programming language

- ► ease-of-use of Python + performance of a traditional compiled language
- combines shared- and distributed-memory models; data and task parallelism for both
- ► multi-resolution: high-level parallel abstractions + low-level controls
- ▶ in my opinion, by far the best language to learn parallel programming ⇒ we teach it as part of HPC Carpentry, in summer schools and full-day Chapel workshops
- experimental support for GPUs
- ▶ relative newcomer to HPC, unfortunately still rarely used outside its small/passionate community

#### • GPU parallel development support: CUDA, OpenCL, OpenACC

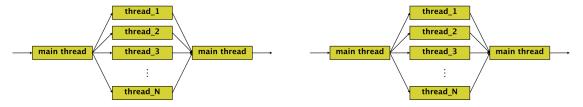
#### Installed compilers

	I	ntel	GN	ľU	PGI		
	intel/2016.4 ar loaded	nd openmpi/2.1.1 by default	module load g	gcc/5.4.0 <sup>(*)</sup>	module load pgi/17.3 (*)		
С	icc	mpicc	gcc -O2	mpicc	pgcc	mpicc	
Fortran 90	ifort	mpifort	gfortran -O2	mpifort	pgfortran	mpifort	
C++	icpc	mpiCC	g++ -O2	mpiCC	pgc++	mpiCC	
OpenMP flag	-qopenmp		-fopenmp		-mp		

- $^{(*)}$  in both cases intel/2016.4 will be unloaded and openmpi/2.1.1 reloaded automatically
- mpiXX scripts invoke the right compiler and link your code to the correct MPI library
- use mpiXX --show to view the commands they use to compile and link

# OpenMP quick look

- OpenMP is a language extension (C, C++, Fortran) for parallel programming in a SMP environment ⇒ pure OpenMP is always limited to a single node
- Programmer uses compiler directives to define "parallel regions" in code which are executed in separate threads
  - uns the master thread until the first parallel region is encountered
  - 2 creates a team of parallel threads
  - 3 when the team threads complete all commands in the parallel region, they synchronize and terminate, leaving only the master thread



# MPI quick look

- MPI library available for all popular programming languages (C, C++, Fortran, Python, R, Java, ...)
- Each processor runs exactly the same copy of the code
- Pseudo-code to exchange variables between two processors (point-to-point operation), starting with A on proc0 and B on proc1:

```
rank <- MPI function to find the current task
if rank == 0
send A to 1
receive B from 1
else if rank == 1
receive A from 0
send B to 0
endif</pre>
```

- Point-to-point or collective communications
- Each processor runs exactly the same copy of the code
- Pseudo-code to calculate a sum using collective reduce operation:

```
sum = 0, partialSum = 0
np <- MPI function to find the total number of tasks
rank <- MPI function to find the current task
decide if I am MASTER (rank=0) or WORKER (rank=1, ..., np-1)
compute partialSum: 1/np-th of the total work based on rank
if I am MASTER
receive from WORKERS their partialSum
compute sum from all partialSum's
print sum
else if I am WORKER
send to MASTER partialSum
endif</pre>
```

#### MPI quick look (cont.)

- More complex MPI operations
  - collective communication routines
  - ② derived data types
  - communicators and virtual topologies

command		send buf	fer		receive buffer			
Communa	$P_1$	$P_2$	$P_3$	$P_4$	$P_1$	$P_2$	$P_3$	$P_4$
MPI_Send()+ MPI_Recv()		A					A	
MPI_Sendrecv		A	В			В	A	
MPI_Bcast		A			A	A	A	A
MPI_Gather	A	В	C	D		A,B,C,D		
MPI_Scatter		A,B,C,D			A	В	С	D
MPI_Allgather	A	В	C	D	A,B,C,D	A,B,C,D	A,B,C,D	A,B,C,D
MPI_Reduce	A	В	C	D		r(A,B,C,D)		
MPI_Allreduce	A	В	С	D	r(A,B,C,D)	r(A,B,C,D)	r(A,B,C,D)	r(A,B,C,D)

# Chapel quick look

High-level abstractions for task and data parallelism

	single locale shared memory parallelism	multiple locales distributed memory parallelism likely shared memory parallelism		
task parallel	config var numtasks = 2; coforall taskid in 1numtasks do writeln("this is task ", taskid);	forall loc in Locales do on loc do writeIn("this locale is named ", here.name);		
data parallel	var A, B, C: [11000] real; forall (a,b,c) in zip(A,B,C) do c = a + b;	use BlockDist; const mesh = {1100,1100} dmapped Block(boundingBox={1100,1100}); var T: [mesh] real; forall (i,j) in T.domain do $T[i,j] = i + j;$		

Locality and parallelism are orthogonal concepts in Chapel: can even have serial execution on mutiple locales

#### Build tools: make

- Tool for automating builds, typically in a workflow with multiple dependencies
- ullet Most frequent usage: source files changed  $\Rightarrow$  recompile parts of the code
  - ▶ consider a large software project with hundreds of source code files, e.g., *Enzo*: 426 C++ files, 6 C files, 121 fortran77 files, 10 fortran90 files, 48 header files
  - ▶ typically work on a small section of the program, e.g., debugging a single function, with much of the rest of the program unchanged  $\Rightarrow$  would be a waste of time to recompile everything (with *Enzo* typically 30-40 mins with heavy optimization) every time you want to compile/run the code
- Another example: updated data files  $\Rightarrow$  redraw the figure  $\Rightarrow$  rebuild the paper
- Hard or impossible to keep track of:
  - ▶ what depends on what
  - ▶ what's up-to-date and what isn't (don't want to redo everything from scratch every time)

#### Build tools: make (cont.)

- Idea: use a build manager to automate the process
- Need to describe the following in a build file (often called a *makefile*)
  - dependencies for each target, e.g. an executable depends on source code files
  - commands used to update targets
- The manager program will aid you in your large workflow
  - checks whether sources are older than targets
  - ightharpoonup if not  $\Rightarrow$  rebuild
- Most widely used build manager is make
  - ▶ invented in 1975, evolved into a programming language of its own
  - ► https://docs.computecanada.ca/wiki/Make
  - ► for large projects there are even pre-processor/build tools for make (CMake, etc.)

```
main.f90
program main
implicit none
real*8 :: a, b, add, sub
a = 4.
b = 1.
print*, add(a,b)
print*, sub(a,b)
end program main
```

```
add.f90
function add(a,b)
  implicit none
  real*8, intent(in) :: a, b
  real*8 :: add
  add = a + b
  return
end function add
```

```
sub.f90

function sub(a,b)

implicit none

real*8, intent(in) :: a, b

real*8 :: sub

sub = a - b

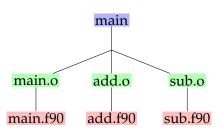
return

end function sub
```

Compiler stage: each .f90 file is converted into an object code (.o) which the computer understands directly

```
gfortran -c add.f90
gfortran -c sub.f90
gfortran -c main.f90
```

2 Linker stage: linking all object codes to optional libraries to produce an executable program main gfortran main.o add.o sub.o -o main



- Start with three files main.f90, add.f90, sub.f90
- Create a file called Makefile, put the following in it:

```
main: main.o add.o sub.o
<TAB> gfortran main.o add.o sub.o -o main
main.o: main.f90
<TAB> gfortran -c main.f90
add.o: add.f90
<TAB> gfortran -c add.f90
sub.o: sub.f90
<TAB> gfortran -c sub.f90
clean:
<TAB> /bin/rm -rf *.o main
```

#### • The format is:

```
target: prerequisites
<TAB> rule to make the target
```

• Compile by typing make or make main

#### Make: first take (shorter version)

Let's remove redundancy using wildcards (%) and predefined makefile variables

```
<TAB> gfortran $^ -o $@
%.o: %.f90
<TAB> gfortran -c $^
clean:
<TAB> @/bin/rm -rf *.o main
```

main: main.o add.o sub.o

- Predefined makefile variables:
  - ► \$@ is "the target of this rule"
  - ► \$^ is "all prerequisites of this rule"
  - ► \$< is "the first prerequisite of this rule"
  - ▶ \$? is "all out-of-date prerequisites of this rule"
- "@" means silent run (without echoing the command)

#### Make: exercise

• Let's create a makefile for compiling  $\pi$  to replace the following lines:

initial command	replacement		
gcc -02 pi.c -o serial	make serial		
gcc -02 -fopenmp sharedPi.c -o openmp	make openmp		
mpicc distributedPi.c -o mpi	make mpi		

On Cedar no need to load any modules before compiling (defaults to Intel compilers)

#### Other essential tools

- Version control (git or mercurial) normally taught as a 3-hour Software Carpentry course
- Terminal multiplexer (*screen* or *tmux*)
  - ► share a physical terminal between several interactive shells
  - access the same interactive shells from many different terminals
  - very useful for persistent sessions, e.g., for compiling large codes
- *VNC* and *x2go* clients for remote interactive GUI work
  - ▶ on Cedar in \$HOME/.vnc/xstartup can switch from twm to mwm/etc. as your default window manager
  - ► can run VNC server on compute nodes
- Understand client-server workflows in other, non-VNC applications: Jupyter notebook, ParaView, VisIt, etc.
  - ▶ we'll take a look at running a Python notebook shortly

#### Python

Details at https://docs.computecanada.ca/wiki/Python

Initial setup:

```
module avail python  # several versions available
module load python/3.7.0
virtualenv astro  # install Python tools in your $HOME/astro
source ~/astro/bin/activate
pip install numpy jupyter pandas  # all these will go into your $HOME/astro
...
```

• Usual workflow:

```
source ~/astro/bin/activate # load the environment
python
...
deactivate
```

- On the same machine create as many virtual environments as you want, switch between them easily
- Ideally you would work on the command line (graphics would be saved to a file) and *eventually* script all your workflows, and then run them as *offscreen batch jobs* on compute nodes
  - few interactive nodes; interactive workflows might not be fully utilizing resources
  - ► for those (hopefully) rare occasions when you need interactivity and on-screen graphics, you can start a remote Jupyter notebook (next slide) you can run it on a *login* (poor security + login node policies and limitations) or *compute* (better) node

#### Jupyter notebook inside a scheduled job

Details at https://docs.computecanada.ca/wiki/Jupyter

- Before you do this... stop and think whether you should be running a Jupyter notebook on a cluster
  - One-time setup:

• Run the Jupyter notebook server on a compute node:

```
source ~/astro/bin/activate # load the environment
salloc --time=1:0:0 --ntasks=1 --cpus-per-task=1 --mem-per-cpu=1024M
~/astro/bin/notebook.sh
```

• Write down the <u>node's name</u>, the Jupyter notebook server <u>listening port</u>, and the <u>token</u>. Next, use them to connect your laptop's browser to this notebook:

```
ssh userXXX@206.12.90.125 -L 8888:nodel:8888 open http://localhost:8888/?token=308cbf318cda9880829445f9baae225ea3d85169d98e43b0
```

- Connecting to a remote Jupyter notebook from Windows
   https://docs.computecanada.ca/wiki/Jupyter#From\_Windows
- A Jupyter notebook can start an RStudio session via RStudio Launcher
  - ▶ adds an RStudio Session option to the Jupyter notebook's New dropdown list
  - uses Jupyter notebook's token authentication system
  - ► details at https://docs.computecanada.ca/wiki/Jupyter#RStudio\_Launcher

#### R on cluster's command line

Details at https://docs.computecanada.ca/wiki/R

Languages/tools

```
$ module spider r  # several versions available
$ module load r/3.4.3
$ R
> install.packages("sp")  # install packages from cran.r-project.org; it'll suggest
# installing into your personal library $HOME/R/
$ R CMD INSTALL -1 $HOME/myRLibPath package.tgz  # install non-CRAN packages
```

- Running R scripts: Rscript script.R
- Installing and running Rmpi: see our documentation
- pbdR (*Programming with Big Data in R*): high-performance, high-level interfaces to MPI, ZeroMQ,
   ScaLAPACK, NetCDF4, PAPI, etc. http://r-pbd.org
- Launching multiple serial R calculations via array jobs (details in Scheduling)
  - ▶ inside the *job submission script* use something like

```
Rscript script${SLURM_ARRAY_TASK_ID}.R
    or
export params=${SLURM_ARRAY_TASK_ID}
Rscript script.R
    and then inside script.R:
s <- Sys.getenv('params')
filename <- paste('/path/to/input', s, '.csv', sep='')</pre>
```

Best practices

# Scheduling and job management

figures and some material in this section borrowed from Kamil Marcinkowski

#### Frequently asked cluster questions

- Why does my job take such a long time to start?
- Is there anything that can be done to make my job start more quickly?
- Why does my job's start time estimate keep moving into the future?

#### Why job scheduler?

- Tens of thousands of CPUs, many thousands of simultaneous jobs ⇒ need an automated solution to manage a queue of pending jobs, allocate resources to users, start/stop/monitor jobs ⇒ we use Slurm open-source scheduler/resource manager
  - ▶ efficiency and utilization: we would like all resources (CPUs, GPUs, memory, disk, bandwidth) to be all used as much as possible, and minimize gaps in scheduling between jobs
  - minimize turnaround for your jobs
- Submit jobs to the scheduler when you have a calculation to run; can specify:
  - ► walltime: maximum length of time your job will take to run
  - ▶ number of CPU cores, perhaps distribution across nodes
  - memory (per core or total)
  - ► if applicable, number of GPUs
  - ► Slurm partition, reservation, software licenses, ...
- Your job is automatically started by the scheduler when enough resources are available
  - ► standard output and error go to file(s)

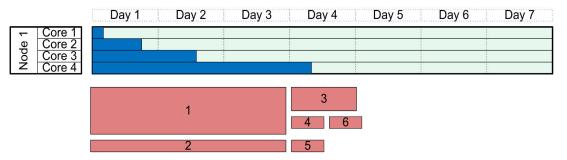
#### Fairshare mechanism

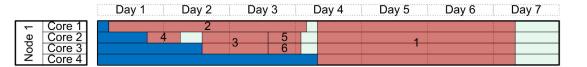
Allocation based on your previous usage and your "share" of the cluster

- Priority: one per research group (not per user!), ranges from 6 (= high, default) to 0 (= low)
- Each group has a share target
  - for regular queues: *share*  $\propto$  the number of group members
  - ightharpoonup in RAC: *share*  $\propto$  the awarded allocation (important projects get a larger allocation)
- If a research group has used more than its *share* during a specific interval (typically 7-10 days) ⇒ its priority will go down, and vice versa
  - ▶ the exact formula for computing priority is quite complex and includes some adjustable weights and optionally other factors, e.g., how long a job has been sitting in the queue
  - ightharpoonup no usage during during the current fairshare interval  $\Rightarrow$  recover back to level 6
- Higher priority level can be used to create short-term bursts
- Reservations (specific nodes) typically only for special events

#### Job packing: simplified view

- Consider a cluster with 4 running jobs and 6 newly submitted jobs
- Scheduled jobs are arranged in order of their users' priority, starting from the top of the priority list (jobs from users with the highest priority)
- Consider 2D view: cores and time
- In reality a multi-dimensional rectangle to fit on a cluster partition: add memory (3rd dimension), perhaps GPUs (4th dimension), and so on, but let's ignore these for simplicity

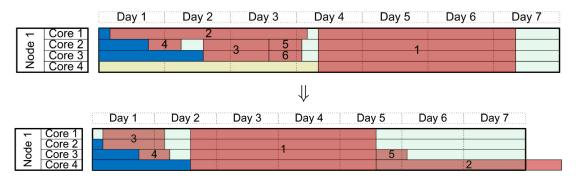




 Backfill: small lower-priority jobs can run on processors reserved for larger higher-priority jobs (that are still accumulating resources), if they can complete before the higher-priority job begins

#### Why does my job's start time estimate keep moving into the future?

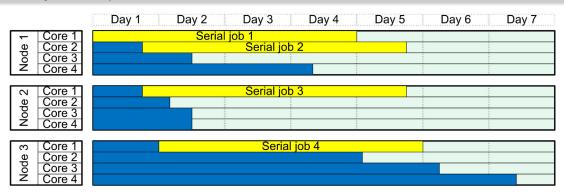
- If a running job finishes early, or a waiting job is canceled, or a new higher priority job is added to the queue ⇒ all waiting jobs are rescheduled from scratch in the next cycle, again in the order of their priority
- This will change the estimated start time, and not always in the obvious direction (notice what happens to job #2 in the graph below!)



#### Job billing: goes into determining your priority

- Recall: base nodes have 128GB and 32 cores per node ⇒ effectively 4GB per core
- Job billing is by core and memory (via core-equivalents: 4GB = 1 core), whichever is larger
  - ► this is fair: large-memory jobs use more resources
  - ▶ a 8GB serial job running for one full day will be billed 48 core-hours

#### Many serial jobs on a cluster



- Each job uses a single CPU: easiest and most efficient to schedule, excellent scaling linear speedup
- Submitting many serial jobs is called "serial farming" (perfect for filling in the parameter space, running Monte Carlo ensembles, etc.)
- In your job script you can ask for a serial job with #SBATCH --ntasks=1 (this is the default if not specified)

./serial

```
$ icc pi.c -o serial
$ sbatch [other flags] job_serial.sh
$ squeue -u username [-t RUNNING] [-t PENDING] # list all current jobs
$ sacct -j jobID [--format=jobid,maxrss,elapsed] # list resources used by completed job
#!/bin/bash
#SBATCH --time=00:05:00 # walltime in d-hh:mm or hh:mm:ss format
#SBATCH --job-name="quick test"
#SBATCH --mem=100 # 100M
```

- --account=... needed only on production systems if you have more than one allocation (RAS / RAC / reservations), used for "billing" purposes (not the same as your cluster account!)
- --reservation=... used only for special events
- On Cedar/Graham in this school we will use --account=ucss19-wa\_cpu and --reservation=ucss19-wr\_cpu
- On the training cluster no need to specify --account, and there are no reservations
- It is good practice to put all flags into a job script (and not the command line)
- Could specify number of other flags (more on these later)

#### Exercises: simple serial jobs

- 1. Submit a serial job that:
  - ► runs the hostname command
- 2. Monitor your job with squeue -u \$USER, check your email, print output file
- 3. How much memory did the job use?
- 4. Now try timing pi.c vs. pi.py on a compute node, using either:
  - ▶ time <command>, or
  - ► Slurm's walltime reporting

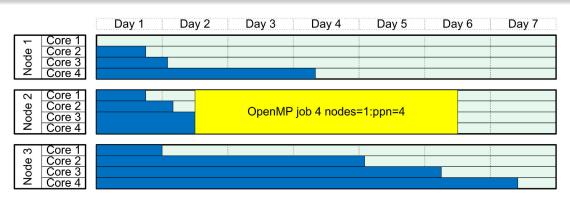
\$ sbatch job\_array.sh [other flags]

Overview

- Job arrays are a handy tool for submitting many serial jobs that have the same executable and might differ only by the input they are receiving through a file
- Job arrays are preferred as they don't require as much computation by the scheduling system to schedule, since they are evaluated as a group instead of individually
- In the example below we want to run 30 times the executable "myprogram" that requires an input file; these files are called input1.dat, input2.dat, ..., input30.dat, respectively

```
#!/bin/bash
#SBATCH --array=1-30
                    # 30 iobs
#SBATCH -- job-name=myprog # single job name for the array
#SBATCH --time=02:00:00 # maximum walltime per job
#SBATCH --mem=100
                           # maximum 100M per job
#SBATCH --output=myprog%A%a.out # standard output
#SBATCH --error=myprog%A%a.err # standard error
# in the previous two lines %A" is replaced by jobID and "%a" with the array index
./myprogram input$SLURM_ARRAY_TASK_ID.dat
```

#### Single-node, muti-core job



- All threads are part of the same process, share single memory address space
- OpenMP is one of the easiest methods of parallel programming
- Always limited to a single node
- Does not have to occupy an entire node

./openmp

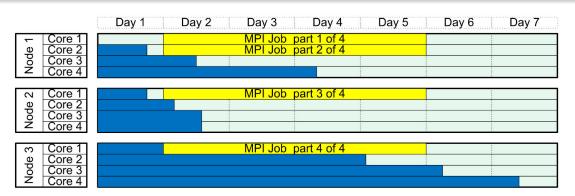
#### Scheduler: submitting OpenMP or threaded jobs

```
$ icc -qopenmp sharedPi.c -o openmp
$ sbatch job_openmp.sh [other flags]
$ squeue -u username [-t RUNNING] [-t PENDING] # list all current jobs
$ sacct -j jobID [--format=jobid,maxrss,elapsed] # list resources used by
# your completed job

#!/bin/bash
#SBATCH --cpus-per-task=2 # number of cores
#SBATCH --time=0-00:05 # walltime in d-hh:mm or hh:mm:ss format
#SBATCH --mem=100 # 100M for the whole job (all threads)
export OMP_NUM_THREADS=$SLURM_CPUS_PER_TASK # passed to the program
echo running on $SLURM_CPUS_PER_TASK cores
```

• Did you get any speedup running this calculation on four cores?

#### MPI job



- Distributed memory: each process uses a different memory address space
- Communication via messages
- More difficult to write MPI-parallel code than OpenMP
- Can scale to much larger number of processors (clusters are larger than SMP machines)

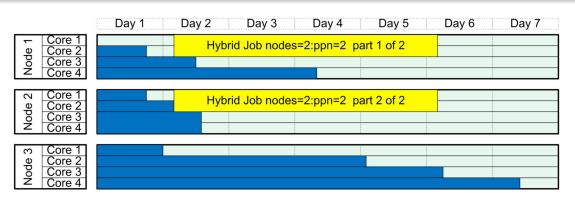
#### Scheduler: submitting MPI jobs

\$ mpicc distributedPi.c -o mpi
\$ sbatch job\_mpi.sh [other flags]

```
$ squeue -u username [-t RUNNING] [-t PENDING] # list all current jobs
$ sacct -j jobID [--format=jobid,maxrss,elapsed] # list resources used by completed job
#!/bin/bash
#SBATCH --ntasks=4 # number of MPI processes
#SBATCH --time=0-00:05 # walltime in d-hh:mm or hh:mm:ss format
#SBATCH --mem-per-cpu=100 # in MB
mpirun -np $SLURM_NTASKS ./mpi # or 'srun ./mpi'
```

- Did you get any speedup running this calculation on four processors?
- What is the code's *parallel efficiency*? Why is it not 100%?

#### Hybrid job



- Combining OpenMP and MPI for running on clusters of SMP machines
  - inside each node: shared-memory communication without MPI programming/overhead
  - scaling to larger resources (CPUs, memory) on the cluster
- In practice might be easier to code with MPI only
- New advanced languages such as Chapel combine shared- and distributed-memory models

```
#!/bin/bash
\#SBATCH --ntasks=4 # number of MPI tasks
#SBATCH --cpus-per-task=12 # number of cores per task
#SBATCH --time=12:00:00 # maximum walltime
if [ -n "$SLURM_CPUS_PER_TASK" ]; then
 omp threads=$SLURM CPUS PER TASK
else
 omp_threads=1
fi
export OMP_NUM_THREADS=$omp_threads
srun ./mpi_openmp_program
```

- More fine-grained control with #SBATCH --ntasks-per-node=...
- Do not run this script: only an example

#### Scheduler: submitting GPU jobs

```
#!/bin/bash
\#SBATCH --nodes=3
                       # number of nodes
#SBATCH --gres=gpu:1
                         # GPUs per node
\#SBATCH --mem=4000M
                         # memory per node
#SBATCH --time=0-05:00
                         # walltime in d-hh:mm or hh:mm:ss format
#SBATCH --output=%N-%j.out
                             # %N for node name, %j for jobID
srun ./gpu_program
```

Do not run this script: only an example

#### Scheduler: interactive jobs

```
$ salloc --time=1:0:0 --ntasks=2  # submit a 2-core interactive job for 1h
$ echo $SLURM_...  # can check out Slurm environment variables
$ ./serial  # this would be a waste: we have allocated 2 cores
$ srun ./mpi  # run an MPI code, could also use mpirun/mpiexec
$ exit  # terminate the job
```

Should automatically go to one of the Slurm interactive partitions

```
$ sinfo -a | grep interac
```

- Useful for debugging or for any interactive work, e.g., GUI visualization
  - ▶ interactive CPU-based ParaView client-server visualization on Cedar and Graham https://docs.computecanada.ca/wiki/Visualization
  - ▶ we use salloc in our hands-on Chapel course
- Make sure to only run the job on the processors assigned to your job this will happen automatically if you use srun, but not if you just ssh from the headnode

#### Slurm environment variables

- Available inside running jobs
- You can start an interactive job, type echo \$SLURM and then hit Tab to see all defined variables inside your job
- Useful to pass parameters to your program at runtime or to print out job information

Environment variable	Description
SLURM_JOB_ID SLURM_NNODES SLURM_NTASKS SLURM_NPROCS SLURM_MEM_PER_CPU SLURM_JOB_NODELIST SLURM_JOB_CPUS_PER_NODE	unique slurm jobID number of nodes allocated to the job number of tasks allocated to the job number of tasks allocated to the job memory allocated per CPU list of nodes on which resources are allocated number of CPUs allocated per node
SLURM_JOB_ACCOUNT	account under which this job is run

#### Exercise: simple interactive job

- Try to print some Slurm environment variables from an interactive job
- Run your "pi.c vs. pi.py" timing exercise as an interactive job

#### Exercise: Jupyter notebook server (more advanced interactive)

- Start a Jupyter notebook server on training VM's compute node
  - ▶ use Python virtual environment from /home/centos/astro
  - ▶ hopefully, not too many people run it: the cluster has only 6 nodes and 2 cores / 3GB per node
- Or you can do the same on Cedar or Graham
  - ▶ there you will have to roll out your own Python virtual environment

#### Slurm script flags

- Use them inside a job script (with #SBATCH) or in the command line (after sbatch)
- We already saw many examples in previous slides
- For full list of flags, type man sbatch

Slurm script command	Description
#SBATCHntasks=X #SBATCHnodes=X #SBATCHnodes=X-Y #SBATCHcpus-per-task=X #SBATCHtasks-per-node=X #SBATCHmail-type=ALL #SBATCHmail-user=	request X tasks; with cpus-per-task=1 (the default) this requests X cores request a minimum of X nodes request a minimum of X nodes and a maximum of Y nodes request a minimum of X CPUs per task request a minimum of X tasks be allocated per node notify via email about ALL, NONE, BEGIN, END, FAIL, REQUEUE set email address

### Slurm script flags (cont.)

Slurm script command	Description
#SBATCHoutput=name%j.out #SBATCHerror=name.err #SBATCHmem=2000 #SBATCHmem-per-cpu=2000 #SBATCHgres=gpu:1 #SBATCHexclusive #SBATCHdependency=after:jobID1 #SBATCHdependency=afterok:jobID1 #SBATCHarray=0-4 #SBATCHarray=0-4	standard output and error log standard error log request 2000 MB of memory in total request 2000 MB of memory per CPU request 1 GPU per node request node(s) with no other running job(s) request that the job starts after jobID1 has started request that the job starts after jobID1 has finished successfully request job array of 5 jobs with indices 0-4 array of 5 jobs with a maximum of 2 jobs running at the same time
#SBATCHarray=1,3,5,9,51	request job array of 5 jobs with indexes 1, 3, 5, 9, 51

#### It is very important to specify memory correctly!

- If you don't ask for enough, and your job uses more, your job will be killed
- If you ask for too much, it will take a much longer time to schedule a job, and you will be wasting resources
- If you ask for more memory than is available on the cluster your job will never run; the scheduling system will not stop you from submitting such a job or even warn you
  - always ask for slightly less than total memory on a node as some memory is used for the OS, and your job will not start until enough memory is available
- Can use either #SBATCH --mem=4000 or #SBATCH --mem-per-cpu=2000
- What's the best way to find your code's memory usage?

#### Slurm jobs and memory (cont.)

 Second best way: use Slurm command to estimate your completed code's memory usage

```
sacct -j jobID [--format=jobid,maxrss,elapsed]
                 # list resources used by a completed job
```

- Use the measured value with a bit of a cushion, maybe 15-20%
- Be aware of the discrete polling nature of Slurm's measurements
  - sampling at equal time intervals might not always catch spikes in memory usage
  - ▶ sometimes you'll see that your running process is killed by the Linux kernel (via kernel's cgroups https://en.wikipedia.org/wiki/Cgroups) since it has exceeded its memory limit but Slurm did not poll the process at the right time to see the spike in usage that caused the kernel to kill the process, and reports lower memory usage
  - sometimes sacct output in the memory field will be empty, as Slurm has not had time to poll the job (job ran too fast)

Overview

```
$ squeue -u username [-t RUNNING] [-t PENDING] # list all current jobs
$ squeue -p partitionName
                                                     # list all jobs in a partition
$ squeue -P --sort=-p,i --states=PD --format="%.10A %.18a %.12P %.8C %.12m %.15l %.25S"
                       # show all queued (=PD) jobs sorted by their current priority
$ scancel [-t PENDING] [-u username] [iobID] # kill/cancel iobs
$ sinfo
                                    # view information about Slurm partitions
$ partition-stats
                                    # similar in a tabular format
$ scontrol show partition
                                    # similar with more details
$ sinfo --states=idle
                                    # show idle node(s) on cluster
$ sinfo -n gra10 -o "%n.%c.%m"
                                    # list node's name, core count and memory
$ sacct -i iobID --format=iobid.maxrss.elapsed # resources used by a completed iob
$ sacct -u username --format=jobid, jobname, avecpu, maxrss, maxvmsize, elapsed
                                        # show details of all your jobs
$ sacct -aX -S 2018-04-25 --format=account %20, partition %20, jobid %12, submit %22, \
 start%22,end%15,timelimit%15,regmem,ncpus%8,nnodes%8,state | more
              # show all jobs (all users) on the system since a specific date
$ scontrol show job jobID # produce a very detailed report for the job
$ sprio [-j jobID1, jobID2] [-u username] # list job priority information
```

#### Slurm partitions

- Idea: restrict jobs of specific shapes to node sets
  - obvious for large-memory / GPU / interactive jobs
- By-node vs. by-core
  - ► *full-node* (or *by-node*) jobs can run on most nodes
  - ▶ partial-node (or by-core) jobs have access to fewer nodes
    - by-core jobs can fit into little, sparse places ⇒ too many of them will easily fill the cluster, leaving no resources for "denser" jobs ⇒ hence separate cluster partitions
- Short vs. long jobs
  - ▶ jobs 3 hours and shorter can run on most nodes
  - ► longer jobs have access to fewer nodes
- Backfill jobs can run on most nodes

## In no particular order:

- Other jobs have greater priority
- ② There is a problem with the job / resources are not available
  - resources do not exist on the cluster?
  - ▶ did not allow for OS memory (whole-node runs)?
- The job is blocked
  - ► disk quota / other policy violation?
  - dependency not met?
- There is a problem with the scheduling system or cluster
  - ▶ most frequent: someone just submitted 10,000 jobs via a script, then cancelled them, then resubmitted them, rendering Slurm unresponsive for a while, even if resources are available for your job right now

## Debugging and profiling

#### Identifying bugs and errors in your code

- Methodical process of finding and fixing flaws in software
- Typical signs that your program is buggy include:
  - ▶ it fails to complete (crashes), usually with an error message in the output file ("Segmentation fault", "Floating point exception", etc) or with a numeric *job exit code*
  - ► it produces incorrect output (NaNs)
  - $\blacktriangleright$  it fails to progress (hangs), often showing  $\sim 100\%$  CPU usage

signal name	OS signal	OS signal	description
	#	name	
floating point	8	SIGFPE	the program attempted an arithmetic operation with val-
exception			ues that do not make sense (e.g., divide by zero)
segmentation	11	SIGSEGV	the program accessed memory incorrectly (e.g., access-
fault			ing an array beyond its declared bounds, using incorrect
			pointers
aborted	6	SIGABRT	generated by the runtime library of the program or a li-
			brary it uses, after having detected a failure condition

#### Common bugs

- Arithmetic: infinities (division by zero), out of range
- Memory access: index out of range, uninitialized pointers
- Logic: infinite loop, corner cases (sloppy condition evaluation)
  - example: try evaluating 0.1 + 0.2 == 0.3 in Python or R
- Misuse: wrong initial conditions, variable initialization (forgot to set to zero?), implicit variable declarations (⇒ wrong type)
- Syntax: wrong operator, function arguments (variable number/types must match)
- Resource starvation: memory leak
- Parallel: race conditions, deadlock, nonmatching send/receive

# Debugging

- Debugger is a program to manipulate and inspect your program as it is running
  - ▶ not a magic bullet you are the real debugger!
- Write better code, use existing libraries instead of your own code
- Test individual parts of your code
- Once you are convinced there is a problem
  - $\buildrel {}^{\blacktriangleright}$  try to reproduce the problem in the simplest situation possible
  - try to reverse your steps to a working state (version control!) and then make one change at a time
  - ► try smaller problem size
  - ▶ use compiler flags to turn off floating point exception masking (the code will stop)
  - ► turn on compiler warnings (GNU: -Wall)
  - (mostly Fortran) use compiler flags to enable runtime checking for various conditions (array indices within bounds, uninitialized variables, proper pointer usage)
  - ensure that variables are defined with sufficient precision (overflow/underflow)
  - use a debugger
  - ▶ use print statements? ... not a good strategy

# Debuggers for compiled languages

- Command-line serial debugger gdb standard on Linux systems
- Commercial parallel GUI debuggers: DDT, TotalView
  - ▶ we have DDT on Graham (installed under \cvmfs on Cedar but no license ...)
- Prepare your code for debugging
  - ► compile your program with "-g" flag to include a symbol table, if you are going to run it in a debugger
  - disable all processor optimizations these might produce misleading debugger behaviour
  - turn off floating point masking behaviour for the program to stop when a NaN or an Inf is computed
    - Intel C/C++ compiler: need explicit signal handling in the code

```
bugs1.c
#include <stdio.h>
void divide(float e, float d) {
  printf("%f\n", e/d);
void array(float f[], int index) {
  printf("%f\n", f[index]);
int main(int argc, char **argv) {
  float a = 0., b = 1., c[10]:
  int i;
  for (i = 0; i < 10; i++)
    c[i] = (float)i;
  divide(b, a);
  array(c, 1200);
  return(0);
```

#### bugs2.c (same with floating point exception handling)

```
#include <fenv.h>
#include <signal.h>
#include <stdio.h>
void divide(float e, float d) {
 printf("%f\n", e/d);
void array(float f[], int index) {
 printf("%f\n", f[index]);
void fpehandler(int sig num) {
 signal (SIGFPE, fpehandler);
 printf("floating point exception, exiting\n");
 abort();
int main(int argc, char **argv) {
 int feenableexcept();
 feenableexcept (FE ALL EXCEPT);
 signal (SIGFPE, fpehandler);
  float a = 0., b = 1., c[10];
 int i;
 for (i = 0; i < 10; i++)
   c[i] = (float)i;
 divide(b, a);
 array(c, 1200);
 return(0);
```

### Typical gdb session

```
$ icc bugs1.c -o bugs
$ ./bugs
inf
70310426758547368093419569152.000000
$ icc bugs2.c -o bugs
$ ./bugs
floating point exception, exiting
Aborted
$ icc -q buqs2.c -o buqs
$ qdb bugs
(gdb) r
Starting program: bugs
Program received signal SIGFPE, Arithmetic exception.
0x00000000004006a5 in divide (e=1, d=0) at bugs2.c:5
          printf("%f\n",e/d);
5
(qdb) where
#0 0x000000000004006a5 in divide (e=1, d=0) at bugs2.c:5
   0x000000000004007d0 in main (argc=1,
          argv=0x7ffffffffc978) at bugs2.c:23
```

```
(qdb) 1 5
        #include <fenv.h>
        #include <signal.h>
        #include <stdio.h>
        void divide(float e, float d) {
          printf("%f\n",e/d);
        void array(float f[], int index) {
8
          printf("%f\n",f[index]);
9
        void fpehandler (int sig num) {
(qdb) p e
$1 = 1
(gdb) p d
$2 = 0
```

Debug/profile 000000000000

# Continue debugging

- Resolve the bug, and then repeat
- The next bug (index=1200) might or might not produce a segmentation fault (that memory address is used by something else inside the code at runtime?)
  - ► if it does, gdb will track it

- ▶ if it does not, need a more powerful memory debugger (valgrind)
- Other gdb commands:
  - insert breakpoints at certain lines or at the beginning of certain functions, then run until the next breakpoint
  - commands to show and delete breakpoints
  - run until the breakpoint
  - ▶ step one line at a time, step into a function or out of a function
  - run help and help className to get more info

# Using core files

- Often it takes a long time before the program reaches the error condition ⇒ cannot debug interactively
- In this case submit the job (debugging-instrumented "-g" executable) to the scheduler, tell the OS to produce a **core file** when it crashes
  - ▶ need to set up the Linux environment to produce core files (set core limit to be non-zero):
     ulimit -c unlimited ← do this on compute node as part of your running job!
     (put it into your job script)
- A core file contains the state of the program at the time it crashed ⇒ can load this
  file into the debugger to inspect the state and determine what caused the problem
- Once a core is produced, load it into gdb
  - \$ qdb ./programName core.jobID

# Debugging summary

- Prepare the code:
  - ► compile with "-g"
  - turn off optimization
  - turn off floating point masking behaviour for the program to stop when a NaN or an Inf is computed (Intel C/C++ compiler: need explicit signal handling in the code)
- gdb for *serial debugging*: will show the function and line in which the error occurred, the reason behind the error, can print variables, step outside the function
  - ► interactive debugging: gdb ./programName
  - ▶ post-processing core files: gdb ./programName core.jobID
- GUI and/or parallel debugging: use a commercial debugger (DDT, TotalView)
  - ► for MPI or threaded codes; can debug individual threads or processes
  - ► again, use "-g" to include a symbol table
  - ▶ parallel bugs: race conditions, deadlock, nonmatching send/receive

# Code profiling

- Profiling tools perform analysis of actual program execution providing very fine-grained information regarding program operation
  - number of times a function is called
  - ▶ amount of time spent in each function ⇒ identify functions that use most CPU time and try to optimize them
- Tools for code profiling
  - inline timing
    - ▷ C: time (seconds only), gettimeofday (seconds and microseconds separately)
    - ▶ Fortran 90: date\_and\_time (ms granularity)
    - ightharpoonup MPI\_Wtime (use MPI\_Wtick to check accuracy, usually  $\mu s$ )
  - 2 gprof (GNU command line serial code profiler)
  - 3 commercial profilers for parallel MPI codes, e.g., OPT
  - ④ IPM, open-source profiler for parallel MPI codes, and other tools from HPC centers
  - 5 Compilers often come with basic profilers too

# Code profiling (cont.)

Using a profiler typically involves three steps:

- (1) instrumenting the source code to collect data: extra compile flags and/or linking to special libraries
- (2) running the instrumented binary
- (3) performing analysis of the collected data after program execution using special GUI tools

# Best practices

# Best practices: computing

- Production runs: only on compute nodes via the scheduler
  - do not run anything intensive on login nodes or directly on compute nodes
- Only request resources (memory, running time) needed
  - ▶ with a bit of a cushion, maybe 115-120% of the measured values
  - use Slurm command to estimate your completed code's memory usage
- Test before scaling, especially parallel scaling
  - ► small tests Ok on login nodes
  - ► larger tests on "interactive" nodes (requested with salloc)
  - test parallel scaling (frequently big surprises!)
- For faster turnaround, if applicable, request whole nodes (--nodes=...) and short runtimes (--time=...)
  - by-node jobs can run on the "entire cluster" partitions, as opposed to smaller partitions for longer and "by-core" jobs
- Do not run unoptimized codes (use compilation flags -O2 or -O3 if needed)
- Be smart in your programming language choice, use precompiled libraries

#### Best practices: file systems

Filesystems in CC are a shared resource and should be used responsibly!

- Do not store millions of small files
  - organize your code's output
  - ▶ use tar or even better dar (https://docs.computecanada.ca/wiki/Dar, supports indexing, per-file compression, differential/incremental backups, encryption, more resilient to data corruption)
- Do not store large data as ASCII (anything bigger than a few MB): waste of disk space and bandwidth
  - use a binary format
  - ▶ use scientific data formats (NetCDF, HDF5, etc.): portability, headers, binary, compression, parallel
  - compress your files
- Use the right filesystem
- Learn and use parallel I/O
- If searching inside a file, might want to read it first
- Have a backup plan
- Regularly clean up your data in \$SCRATCH, \$PROJECT, possibly archive elsewhere

# Summary

# Documentation and getting help

- Official documentation https://docs.computecanada.ca/wiki
- WestGrid training materials https://westgrid.github.io/trainingMaterials
- Compute Canada Getting started videos http://bit.ly/2sxG033
- Compute Canada YouTube channel http://bit.ly/2ws0JDC
- Email support@computecanada.ca (goes into the ticketing system)
  - ▶ try to include your full name, CC username, institution, the cluster name, copy and paste as much detail as you can (error messages, jobID, job script, software version)
- Email me alex.razoumov@westgrid.ca for questions on any of today's topics
- Please get to know your local support
  - ► difficult problems are best dealt face-to-face
  - might be best for new users