

Laboratory Exercise 1

Giới thiệu về công cụ MARS

Goals

Sau bài thực hành này, bạn sẽ cài đặt được công cụ MARS, viết thử chương trình đơn giản để thử nghiệm công cụ MARS như lập trình hợp ngữ, chạy giả lập, gỡ rối, và các phương tiện nhằm hiểu rõ hơn về bản chất và các hoạt động thực sự đã xảy ra trong bộ xử lý MIPS.

Literature

- Tài liệu Kiến trúc MIPS.pptx
- Tài liệu MARS features, file .doc
- Tài liệu MARS Tutorial, file .doc

About MIPSIT

- MIPS (**M**icroprocessor **w**ithout **I**nterlocked **P**ipeline **S**tages) hình thành trên cơ sở RISC.
- Năm 1981: John L. Hennessy đứng đầu một nhóm bắt đầu một công trình nghiên cứu về *bộ xử lý MIPS* đầu tiên tại **Stanford University**
- Một thiết kế chủ chốt trong MIPS là yêu cầu các câu lệnh phải hoàn thành trong 1 chu kỳ máy.
- Một số ứng dụng
 - Pioneer DVR-57-H
 - Kenwood HDV-810 Car Navigation System
 - HP Color Laser Jet 2500 Printer

Kick-off

Tải về và chạy

1. Tải về Java Runtime Environment, JRE, để chạy công cụ MARS
<http://www.oracle.com/technetwork/java/javase/downloads/index.html>
2. Cài đặt JRE
3. Tải về công cụ MARS, bao gồm
 - a. Phiên bản mới nhất của MARS, và nên lấy thêm 2 tài liệu:
 - b. MARS features, và
 - c. MARS Tutorialở URL sau
<http://courses.missouristate.edu/KenVollmar/MARS/download.htm>

Xong. Công cụ MARS có thể thực hiện ngay mà không cần cài đặt. Click đúp vào file Mars.jar để chạy.

Chú ý: để tiện lợi, các sinh viên có thể tải về các tài nguyên liên quan tại URL duy nhất sau: <ftp://dce.hust.edu.vn/tiennd/ict4/Mars/>

Cơ bản về giao diện lập trình IDE

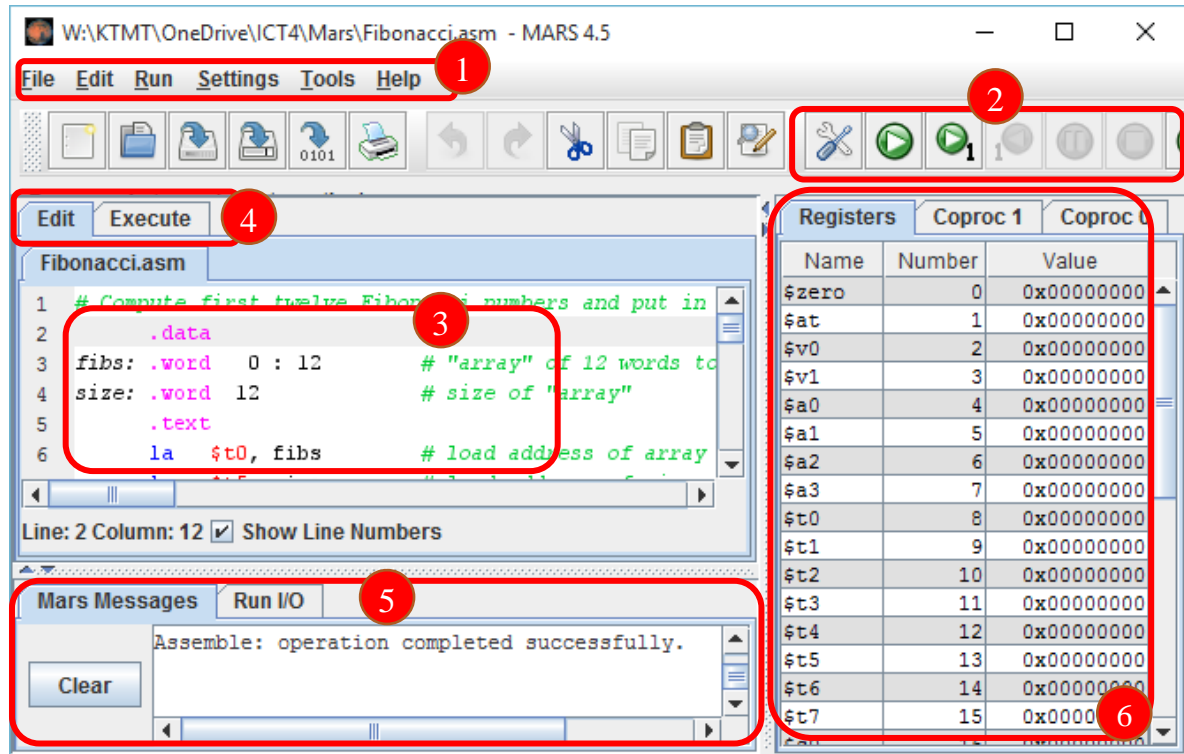


Figure 1. IDE of MARS Tool

1. **Menus:** Hầu hết các mục trong menu đều có các icon tương ứng
 - Di chuyển chuột lên trên của icon → tooltip giải thích về chức năng tương ứng sẽ hiển thị.

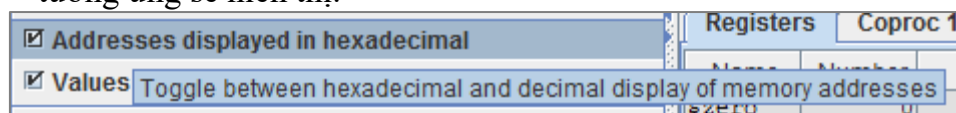


Figure 2. Tooltip giải thích chức năng trong các Menu

- Các mục trong menu cũng có phím tắt tương ứng.
2. **Toolbar:**
 - Chứa một vài tính năng soạn thảo cơ bản như: copy, paste, open..
 - Các tính năng gỡ rối (trong hình chữ nhật màu đỏ)
 - Run: chạy toàn bộ chương trình
 - Run one step at a time: **chạy từng lệnh và dừng (rất hữu ích)**
 - Undo the last step: **khôi phục lại trạng thái ở lệnh trước đó (rất hữu ích)**
 - Pause: tạm dừng quá trình chạy toàn bộ (Run)
 - Stop: kết thúc quá trình gỡ rối
 - Reset MIPS memory and register: Khởi động lại
 3. **Edit tab:** MARS có bộ soạn thảo văn bản tích hợp sẵn với tính **tô màu theo cú pháp**, giúp sinh viên dễ theo dõi mã nguồn. Đồng thời, khi sinh

viên gõ lệnh mà chưa hoàn tất, một popup sẽ hiện \$rad để trợ giúp. Vào menu Settings / Editor... để thay đổi màu sắc, font...



Figure 3. Popup trợ giúp hoàn tất lệnh và giải thích lệnh

4. Edit/Execute:

Mỗi file mã nguồn ở giao diện soạn thảo có 2 cửa sổ - 2 tab: Edit và Execute

- **Edit tab:** viết chương trình hợp ngữ với tính năng **tô màu theo cú pháp**.
- **Execute tab:** biên dịch chương trình hợp ngữ đã viết ở Edit tab thành mã máy, **chạy và gỡ rối**.

5. Message Areas: Có 2 cửa sổ message ở cạnh dưới của giao diện IDE

- **The Run I/O tab** chỉ có tác dụng khi đang chạy run-time
 - **Hiển thị các kết quả xuất \$raconsole, và**
 - **Nhập dữ liệu vào cho chương trình qua console.**
- MARS có tùy chọn để bạn có thể mọi thông tin nhập liệu vào qua console sẽ được hiển thị lại \$ramessage area.
- **MARS Messages tab** được dùng để hiển thị cho các thông báo còn lại như là các báo lỗi trong quá trình biên dịch hay trong quá trình thực hiện run-time. Bạn có thể **click vào thông báo lỗi để chương trình tự động nhảy tới dòng lệnh** gây \$ralỗi.

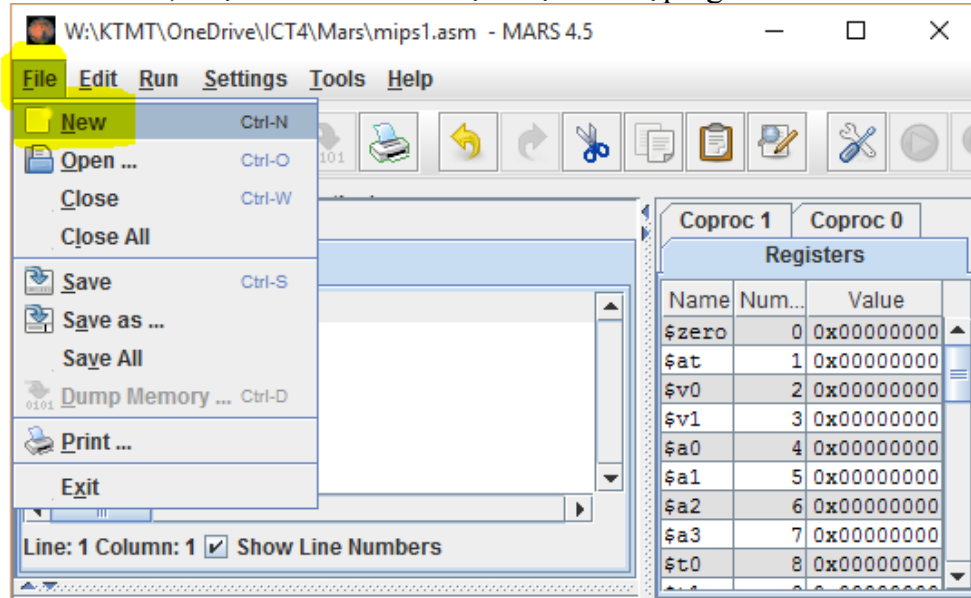
6. MIPS Registers: Bảng hiển thị giá trị của các thanh ghi của bộ xử lý MIPS, luôn luôn được hiển thị, bất kể chương trình hợp ngữ có được chạy hay không. Khi viết chương trình, bảng này sẽ giúp bạn nhớ được tên của các thanh ghi và địa chỉ của chúng. Có 3 tab trong bảng này:

- **the Register File:** các thanh ghi số nguyên với địa chỉ từ \$0 tới \$31, và cả 3 thanh ghi đặc LO, HI và thanh ghi Program Counter
- **the Coprocessor 0 registers:** các thanh ghi của bộ đồng xử lý C0, phục vụ cho xử lý ngắt
- **the Coprocessor 1 registers:** các thanh ghi số dấu phẩy động

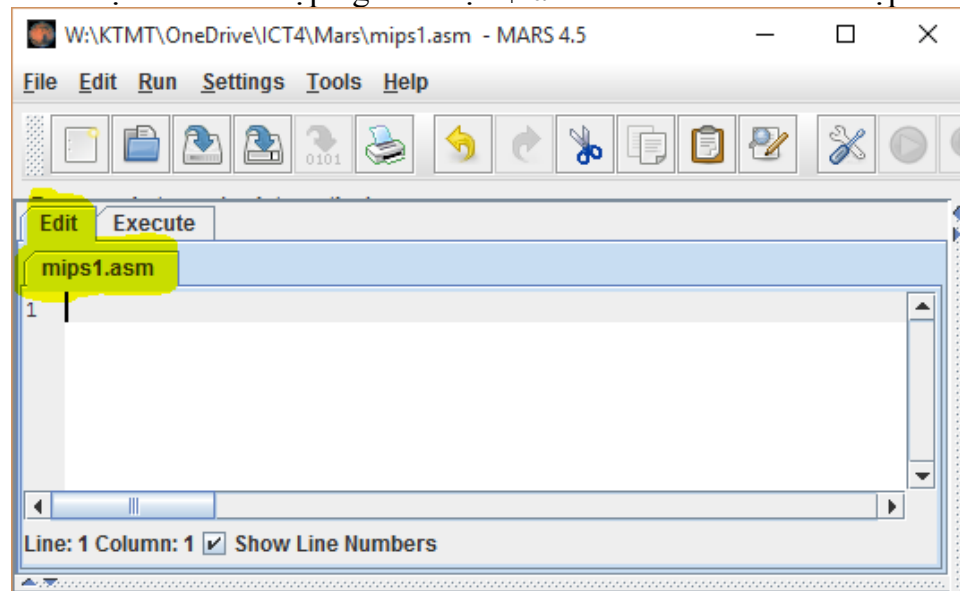
Bắt đầu lập trình và hiểu các công cụ với chương trình Helloworld

1. Click vào file mars.jar để bắt đầu chương trình

2. Ở thanh menu, chọn File / New để tạo một file hợp ngữ mới



3. Cửa sổ soạn thảo file hợp ngữ sẽ hiện \$ranhư hình bên. Bắt đầu lập trình



4. Hãy gõ đoạn lệnh sau vào cửa sổ soạn thảo

```
.data                                # Vùng du lieu, chua cac khai bao bien
x:      .word      0x01020304      # bien x, khoi tao gia tri
message: .asciiz    "Bo mon Ky thuat May tinh"
.text                                # Vùng lệnh, chua cac lệnh hop ngu
la $a0, message                      #Dua dia chi bien message vao thanh ghi a0
li $v0, 4                            #Gan thanh ghi $v0 = 4
syscall                              #Goi ham so v0, ham so 4, la ham print

addi $t1,$zero,2                     #Thanh ghi $t1 = 2
addi $t2,$zero,3                     #Thanh ghi $t2 = 3
add $t0, $t1, $t2                    #Thanh ghi t- = $t1 + $t2
```


Kết quả như sau.

```

Edit  Execute
-----
mips1.asm*
1  .data                # Vung du lieu, chua cac khai bao bien
2  x: .word             0x01020304  # bien x, khoi tao gia tri
3  message: .asciiz     "Bo mon Ky thuat May tinh"
4
5  .text                # Vung lenh, chua cac lenh hop ngu
6  la $a0, message      #Dua dia chi bien message vao thanh ghi a0
7  li $v0, 4            #Gan thanh ghi v0 = 4
8  syscall              #Goi ham so v0, ham so 4, la ham print
9
10 addi $t1,$zero,2      #Thanh ghi t1 = 2
11 addi $t2,$zero,3      #Thanh ghi t2 = 3
12 add $t0, $t1, $t2     #Thanh ghi t- = t1 + t2

```

5. Để biên dịch chương trình hợp ngữ trên thành mã máy, thực hiện một trong các cách sau:

- vào menu Run / Assemble, hoặc
- trên thanh menu, bấm vào biểu tượng , hoặc
- bấm phím tắt F3.

6. Nếu đoạn hợp ngữ đúng, MARS sẽ chuyển từ Edit tab sang Execute tab.

Execute

Text Segment

Bkpt	Address	Code	Basic	Source
	0x00400000	0x3c011001	lui \$1,0x00001001	6: la \$a0, message #Dua d...
	0x00400004	0x34240004	ori \$4,\$1,0x00000004	
	0x00400008	0x24020004	addiu \$2,\$0,0x00000004	7: li \$v0, 4 #Gan t...
	0x0040000c	0x0000000c	syscall	8: syscall #Goi h...
	0x00400010	0x20090002	addi \$9,\$0,0x00000002	10: addi \$t1,\$zero,2 #Thanh...
	0x00400014	0x200a0003	addi \$10,\$0,0x00000003	11: addi \$t2,\$zero,3 #Thanh...
	0x00400018	0x012a4020	add \$8,\$9,\$10	12: add \$t0, \$t1, \$t2 #Thanh...

Labels

Label	Address
mips1.asm	
x	0x10010000
message	0x10010004

Data Segment

Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	m o B	K n o	h t y	t a u	y a M	h n i t	\0 \0 \0 \0
0x10010020	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010040	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010060	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010080	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x100100a0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x100100c0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x100100e0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010100	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0

0x10010000 (.data) Hexadecimal Addresses Hexadecimal Values ASCII

Chú ý: nếu đoạn hợp ngữ có lỗi, cửa sổ Mars Messages sẽ hiển thị chi tiết lỗi. Bấm vào dòng thông báo lỗi để trình soạn thảo tự động nhảy tới dòng code bị lỗi, rồi tiến hành sửa lại cho đúng.

Mars Messages Run I/O

```
Assemble: assembling W:\KIMT\OneDrive\ICT4\Mars\mips1.asm  
Error in W:\KIMT\OneDrive\ICT4\Mars\mips1.asm line 7 column 4: "la": Too few or incorrectly  
Assemble: operation completed with errors.  
Assemble: assembling W:\KIMT\OneDrive\ICT4\Mars\mips1.asm  
Error in W:\KIMT\OneDrive\ICT4\Mars\mips1.asm line 8 column 8: Symbol "printf" not found in s  
Assemble: operation completed with errors.
```

7. Ở Execute tab, có 2 cửa sổ chính là **Text Segment**, và **Data Segment**

The screenshot shows the Mars Execute tab with the 'Execute' sub-tab selected. It displays two main windows: 'Text Segment' and 'Data Segment'. The 'Text Segment' window shows assembly code with annotations in Vietnamese. The 'Data Segment' window shows memory addresses and values. Arrows indicate the mapping between the assembly code and the data segment.

Text Segment

Bkpt	Address	Code	Basic	Source
	0x00400000	0x3c011001	lui \$1,0x00001001	6: la \$a0, message
	0x00400004	0x34240004	ori \$4,\$1,0x00000004	
	0x00400008	0x24020004	addiu \$2,\$0,0x00000004	7: li \$v0, 4
	0x0040000c	0x0000000c	syscall	8: syscall
	0x00400010	0x20090002	addi \$9,\$0,0x00000002	10: addi \$t1,\$zero,2
	0x00400014	0x200a0003	addi \$10,\$0,0x00000003	11: addi \$t2,\$zero,3
	0x00400018	0x012a4020	add \$8,\$9,\$10	

Data Segment

Address	Value (-0)	Value (+4)	Value (+8)	Value (+12)
0x10010000	m o B K		
0x10010020	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010040	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010060	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010080	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x100100a0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x100100c0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x100100e0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0
0x10010100	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0	\0 \0 \0 \0

Annotations:

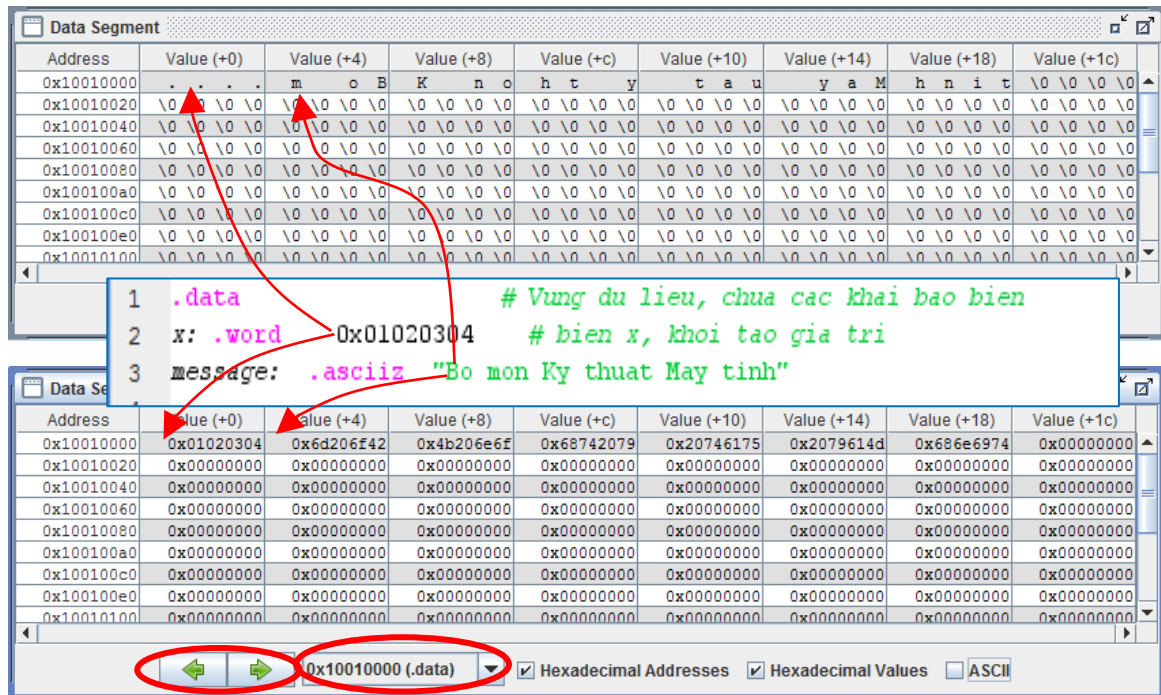
- .data** # Vùng dữ liệu, chứa các khai báo biến
x: .word 0x01020304 # biến x, khởi tạo giá trị
message: .asciiz "Bo mon Ky thuat May tinh"
- .text** # Vùng lệnh, chứa các lệnh hợp ngữ
la \$a0, message #Đưa địa chỉ biến message vào thanh ghi a0
li \$v0, 4 #Gán thanh ghi v0 = 4
syscall #Gọi hàm sys_v0, hàm sys_4, là hàm print
addi \$t1,\$zero,2 #Thanh ghi t1 = 2
addi \$t2,\$zero,3 #Thanh ghi t2 = 3
add \$t0,\$t1,\$t2 #Thanh ghi t = t1 + t2

- **Text Segment:** là vùng không gian bộ nhớ chứa các mã lệnh hợp ngữ. Tương ứng với mã nguồn hợp ngữ, các dòng nào viết sau chỉ thị **.TEXT** tức là lệnh và sẽ thuộc Text Segment.
- **Data Segment:** là vùng không gian bộ nhớ chứa các biến. Tương ứng với mã nguồn hợp ngữ, các dòng nào viết sau chỉ thị **.DATA** tức là lệnh và sẽ thuộc Text Segment.

Chú ý: vì lý do nào đó, nếu ta khai báo biến sau chỉ thị .TEXT hoặc ngược lại thì trình biên dịch sẽ báo lỗi hoặc bỏ qua khai báo sai đó.


8. Ở Execute tab, sử dụng checkbox bên dưới để thay đổi cách hiển thị dữ liệu cho dễ quan sát

- ☐ **Hexadecimal Addresses** : hiển thị địa chỉ ở dạng số nguyên hệ 16
- ☐ **Hexadecimal Values** : hiển thị giá trị thanh ghi ở dạng số nguyên hệ 16
- ☐ **ASCII** : hiển thị giá trị trong bộ nhớ ở dạng ký tự ASCII



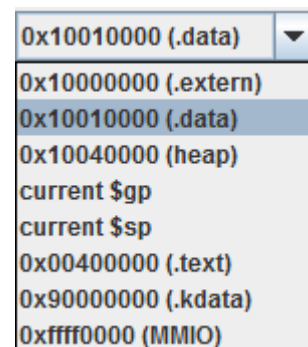
- **Address:** địa chỉ của dữ liệu, biến ở dạng số nguyên. Giá trị mỗi dòng tăng 32 đơn vị (ở hệ 10, hoặc $20_{(16)}$) bởi vì mỗi dòng sẽ trình bày 32 byte ở các địa chỉ liên tiếp nhau
- **Các cột Value:** mỗi cột Value chứa 4 byte, và có 8 cột Value, tương ứng với 32 byte liên tiếp nhau.

Trong hình ảnh trên, có thể thấy rõ giá trị của biến $x = 0x01020304$ được hiển thị chính xác trong Data Segment khi hiển thị dữ liệu ở dạng số ☐ ASCII, và giá trị của chuỗi “Bo mon Ky thuat May tinh” khi hiển thị ở dạng kí tự ☒ ASCII. *Lưu ý rằng việc lưu trữ chuỗi trong bộ nhớ ở dạng little-endian là do cách lập trình hàm phần mềm syscall, chứ không phải do bộ xử lý MIPS qui định. Có thể thấy, ở công cụ giả lập MIPS IT, hàm print lại qui định chuỗi theo kiểu big-endian.*

Bấm vào cặp nút  để dịch chuyển tới vùng địa chỉ lân cận

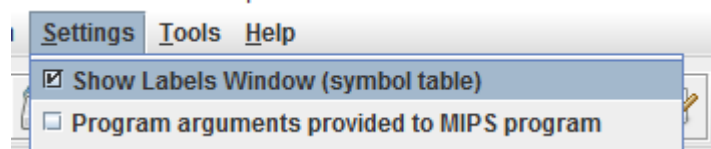
Bấm vào combobox để dịch tới vùng bộ nhớ chứa loại dữ liệu được chỉ định. Trong đó lưu ý

- .data: vùng dữ liệu
- .text: vùng lệnh
- \$sp: vùng ngăn xếp



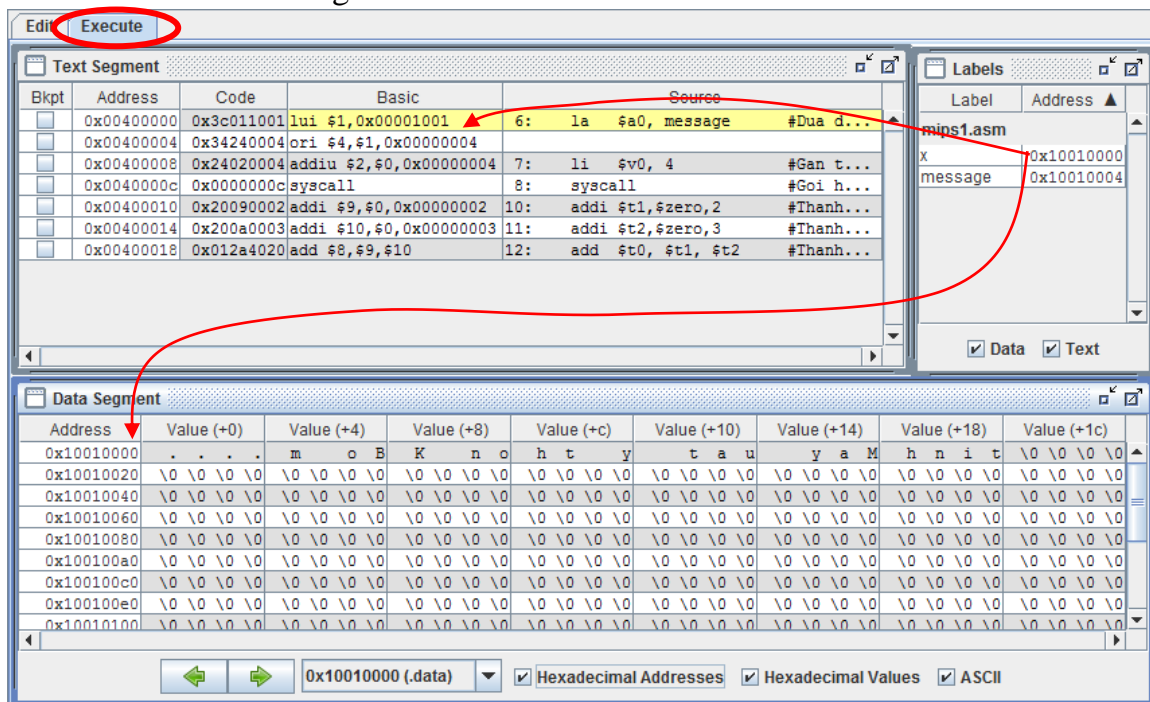
11. Cửa sổ Label: hiển thị tên nhãn và hằng số địa chỉ tương ứng với nhãn khi được biên dịch \$ramã máy.

Cửa sổ Label không tự động hiển thị. Phải vào menu Settings / chọn Show Labels Windows



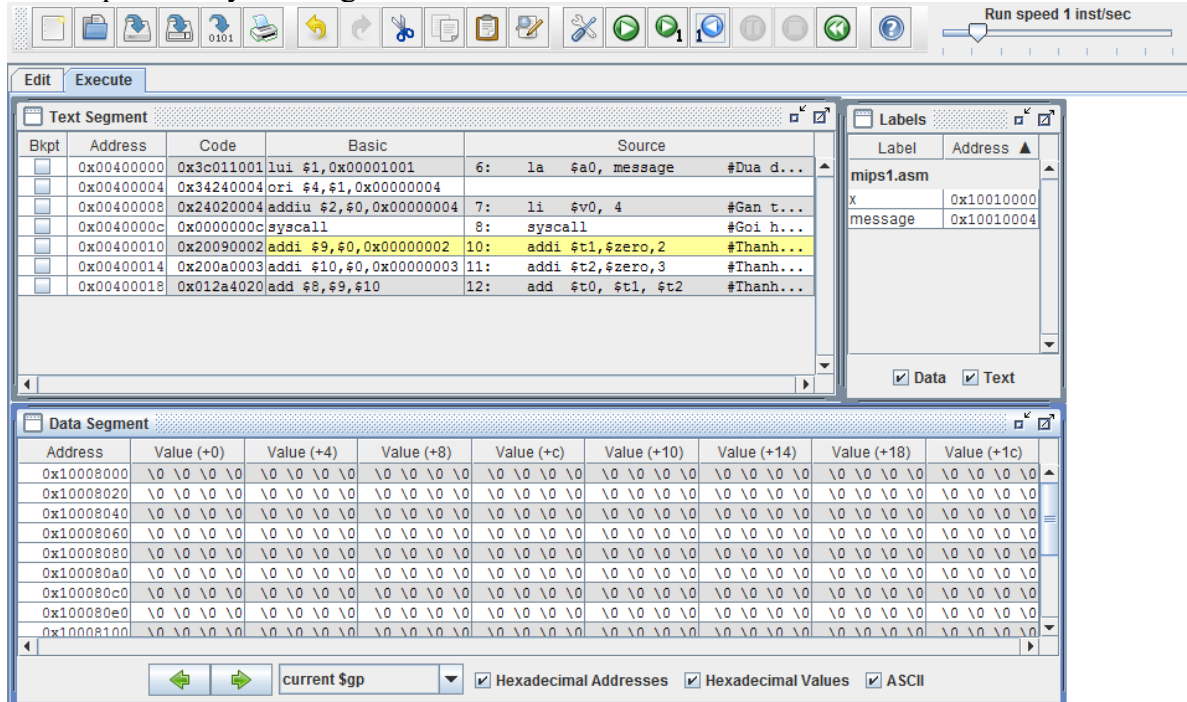
Trong ảnh sau, ta thấy các thông tin quan trọng :

- Trong cửa sổ Labels cho biết :
 - o X chỉ là tên gọi nhớ, x sẽ được qui đổi thành hằng số 0x10010000.
 - o Message cũng chỉ là tên gọi nhớ, sẽ được qui đổi thành hằng số 0x10010004
 - o Click đúp vào tên biến, sẽ tự động chuyển sang vị trí tương ứng trong cửa sổ Data Segment.
- Trong cửa sổ Text Segment cho biết
 - o Quả thực, ở lệnh gán *la \$a0, message* tên gọi nhớ message đã được chuyển thành hằng số 0x10010004 thông qua cặp lệnh *lui, ori*
- Trong cửa sổ Data Segment cho biết
 - o Quả thực, để giám sát giá trị của biến X, ta mở Data Segment ở hằng số 0x10010000 sẽ nhìn thấy giá trị của X.
 - o Quả thực, để giám sát giá trị của biến message, ta mở Data Segment ở hằng số 0x10010004 sẽ nhìn thấy giá trị của message.

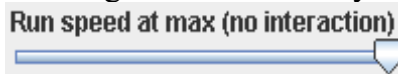


Chạy giả lập

1. Tiếp tục chạy chương trình Hello World ở trên.



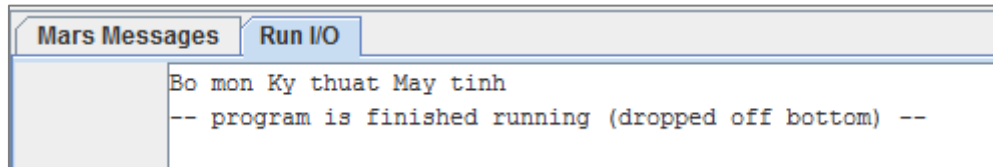
2. Sử dụng slider bar để thay đổi tốc độ thực thi lệnh hợp ngữ.



Mặc định, tốc độ thực thi là tối đa, và ở mức này, ta không thể can thiệp được nhiều vào quá trình hoạt động của các lệnh và kiểm soát chúng. Có thể dịch chuyển slider bar về khoảng 2 lệnh/giây để dễ quan sát.

3. Ở Execute tab, chọn cách để thực thi chương trình

- Bấm vào icon , Run, để thực hiện toàn bộ chương trình. Khi sử dụng Run, ta quan sát dòng lệnh được tô màu vàng cho biết chương trình hợp ngữ đang được xử lý tới chỗ nào. Đồng thời, quan sát sự biến đổi dữ liệu trong cửa sổ Data Segment và cửa sổ Register.
- Bấm vào icon , Reset, để khởi động lại trình giả lập về trạng thái ban đầu. Tất cả các ngăn nhớ và các thanh ghi đều được gán lại về 0.
- Bấm vào icon , Run one step, để thực thi chỉ duy nhất 1 lệnh rồi chờ bấm tiếp vào icon đó, để thực hiện lệnh kế tiếp.
- Bấm vào icon , Run one step backwards, để khôi phục lại trạng thái và quay trở lại lệnh đã thực thi trước đó.
- Sau khi chạy xong tất cả các lệnh của phần mềm Hello Word, sẽ thấy cửa sổ Run I/O hiển thị chuỗi.



Giải lập & gỡ rối: quan sát sự thay đổi của các biến

Trong quá trình chạy giả lập, hãy chạy từng lệnh với chức năng Run one step. Ở mỗi lệnh, quan sát sự thay đổi giá trị trong cửa sổ Data Segment và cửa sổ Register, và hiểu rõ ý nghĩa của sự thay đổi đó.

Giải lập & gỡ rối: thay đổi giá trị biến khi đang chạy run-time

Trong khi đang chạy giả lập, ta có thể thay đổi giá trị của một ngăn nhớ bất kỳ bằng cách

1. Trong Data Segment, click đúp vào một ngăn nhớ bất kỳ

Address	Value (+0)	Value (+4)
0x10010000	0x01020304	0x6d206f42
0x10010020	0x00000000	0x00000000
0x10010040	0x00000000	0x00000000

2. Nhập giá trị mới

Address	Value (+0)	Value (+4)
0x10010000	0x00000008	0x6d206f42
0x10010020	0x00000000	0x00000000
0x10010040	0x00000000	0x00000000

3. Tiếp tục chạy chương trình

Giải lập & gỡ rối: thay đổi giá trị thanh ghi khi đang chạy run-time

Cách làm tương tự như thay đổi giá trị của biến, áp dụng cho cửa sổ Registers

Name	Number	Value
\$zero	0	0x00000000
\$at	1	0x10010000
\$v0	2	0x00000004
\$v1	3	0x00000000
\$a0	4	0x10010004

Tra cứu Help



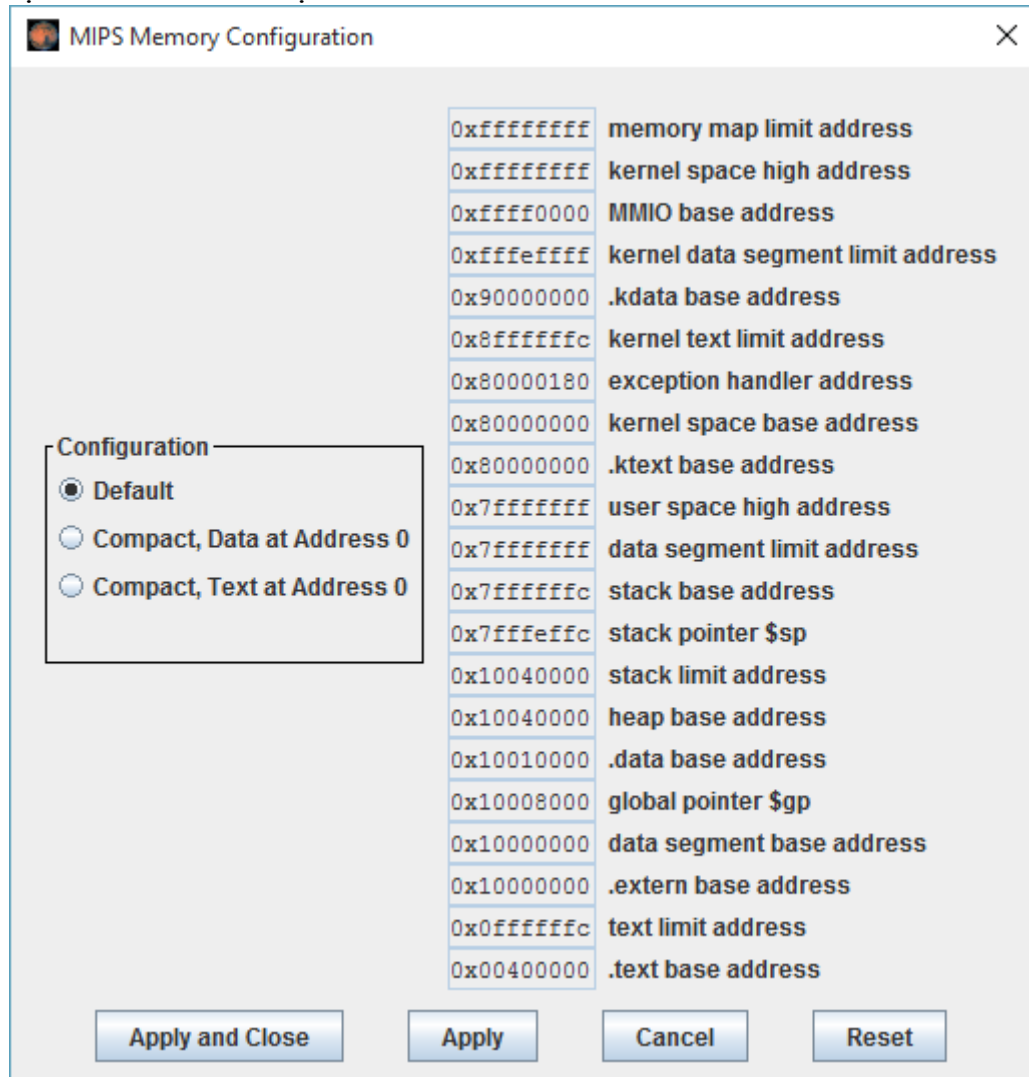
Bấm nút Help để xem giải thích các lệnh của MIPS, các giả lệnh, chỉ dẫn biên dịch và các hàm của syscalls.

Các hằng địa chỉ

Chọn menu Settings / Memory Configuration...

Cửa sổ MIPS Memory Configuration chứa bảng qui định các hằng địa chỉ mà công cụ MARS sử dụng.

Theo bảng này, có thể thấy các mã lệnh luôn bắt đầu từ địa chỉ 0x00400000, còn dữ liệu luôn bắt đầu từ địa chỉ 0x10000000.



The image shows a 'MIPS Memory Configuration' dialog box. On the left, there is a 'Configuration' section with three radio buttons: 'Default' (selected), 'Compact, Data at Address 0', and 'Compact, Text at Address 0'. On the right, there is a list of memory addresses and their corresponding labels. At the bottom, there are four buttons: 'Apply and Close', 'Apply', 'Cancel', and 'Reset'.

Address	Label
0xffffffff	memory map limit address
0xffffffff	kernel space high address
0xffff0000	MMIO base address
0xffffefff	kernel data segment limit address
0x90000000	.kdata base address
0x8ffffffc	kernel text limit address
0x80000180	exception handler address
0x80000000	kernel space base address
0x80000000	.ktext base address
0x7fffffff	user space high address
0x7fffffff	data segment limit address
0x7ffffffc	stack base address
0x7ffffeffc	stack pointer \$sp
0x10040000	stack limit address
0x10040000	heap base address
0x10010000	.data base address
0x10008000	global pointer \$gp
0x10000000	data segment base address
0x10000000	.extern base address
0x0fffffc	text limit address
0x00400000	.text base address

Laboratory Exercise 2

Instruction Set, Basic Instructions, Directives

Goals

Sau bài thực hành này, sinh viên sẽ nắm được nguyên lý cơ bản về tập lệnh của bộ xử lý MIPS; sử dụng được các lệnh hợp ngữ cơ bản và sử dụng công cụ gỡ rối để kiểm nghiệm lại các kiến thức về tập lệnh và hợp ngữ. Sinh viên cũng thành thạo với các chỉ thị biên dịch (Directives) để công cụ MARS có thể dịch hợp ngữ thành mã máy một cách đúng đắn.

Literature

- Tài liệu tóm tắt về Kiến trúc MIPS¹, file pptx
- Bảng tra cứu tập lệnh MIPS², file .doc

Assignments at Home and at Lab

Home Assignment 1

Đọc tài liệu về Kiến trúc MIPS và ghi nhớ các kiến thức cơ bản sau

- Tên và ý nghĩa của 32 thanh ghi
- Các thanh ghi đặc biệt PC, HI, LO
- Khuôn dạng của 3 loại lệnh I, J, R

MiniMIPS Instruction Formats

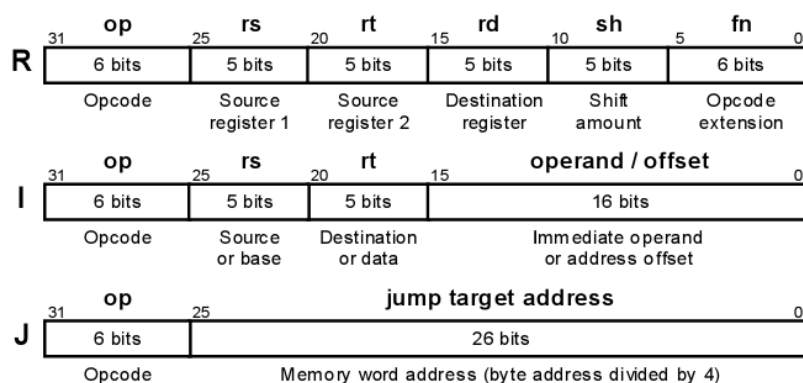


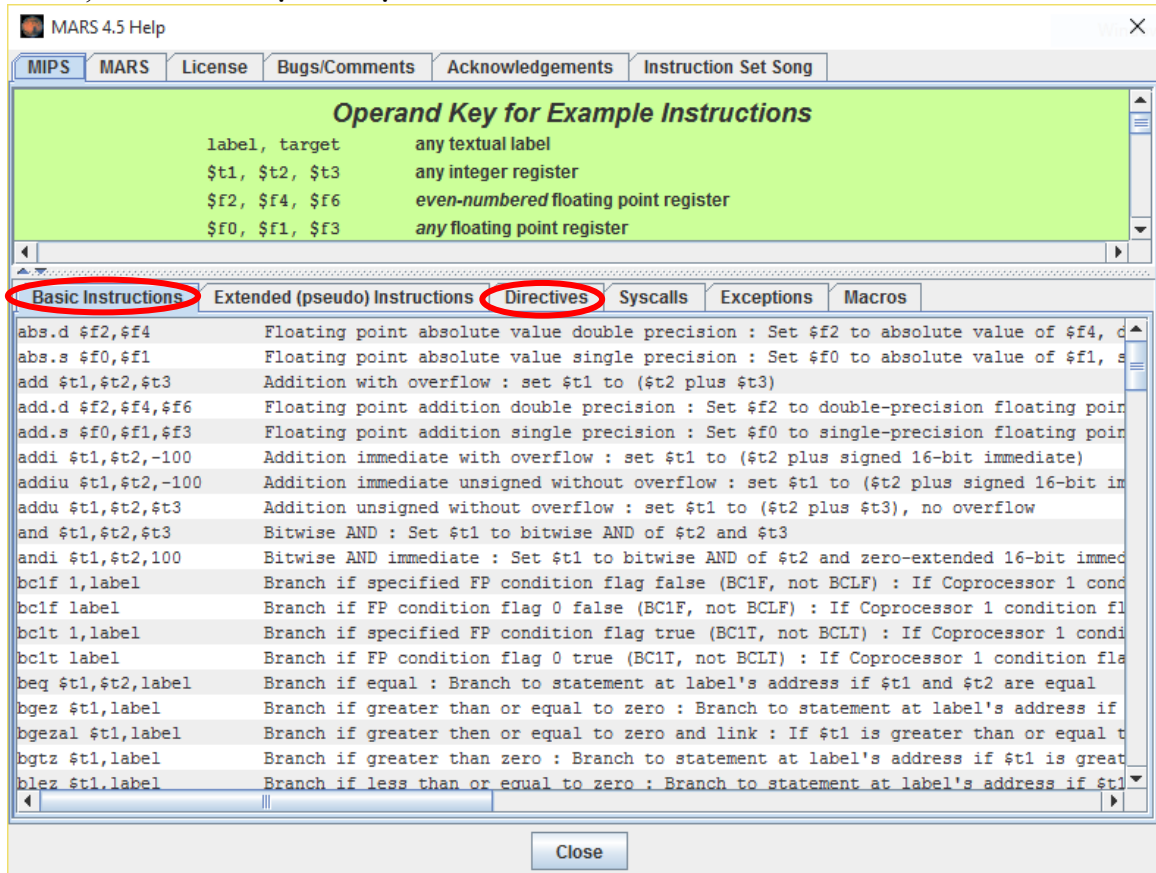
Figure 5.4 MiniMIPS instructions come in only three formats: register (R), immediate (I), and jump (J).

¹ Download tại: [ftp://dce.hust.edu.vn/tiennd/ict4/References/Kien truc MIPS.pptx](ftp://dce.hust.edu.vn/tiennd/ict4/References/Kien%20truc%20MIPS.pptx)

² Download tại: <ftp://dce.hust.edu.vn/tiennd/ict4/References/MIPS32-InstructionSet-QuickReference.pdf>

Home Assignment 2

Sử dụng công cụ MARS, tra cứu Help và tìm hiểu về các lệnh cơ bản trong MIPS, và các chỉ thị biên dịch




Assignment 1: lệnh gán số 16-bit

Gõ chương trình sau vào công cụ MARS.

```
#Laboratory Exercise 2, Assignment 1
.text
    addi    $s0, $zero, 0x3007 # $s0 = 0 + 0x3007 = 0x3007 ;I-type
    add     $s0, $zero, $0      # $s0 = 0 + 0 = 0 ;R-type
```

Sau đó:

- Sử dụng công cụ gỡ rối, Debug, chạy từng lệnh và dừng lại, 
 - Ở mỗi lệnh, quan sát cửa sổ Register và chú ý
 - o Sự thay đổi giá trị của thanh ghi \$s0
 - o Sự thay đổi giá trị của thanh ghi \$pc
 - Ở cửa sổ Text Segment, hãy so sánh mã máy của các lệnh trên với khuôn dạng lệnh để chứng tỏ các lệnh đó đúng như tập lệnh đã qui định
 - Sửa lại lệnh lui như bên dưới. Chuyện gì xảy ra sau đó. Hãy giải thích
- ```
addi $s0, $zero, 0x2110003d
```

### Assignment 2: lệnh gán số 32-bit

Gõ chương trình sau vào công cụ MARS.

```
#Laboratory Exercise 2, Assignment 2
```



```
.text
 lui $s0, 0x2110 #put upper half of pattern in $s0
 ori $s0, $s0, 0x003d #put lower half of pattern in $s0
```

Sau đó:



- Sử dụng công cụ gỡ rối, Debug, chạy từng lệnh và dừng lại,
- Ở mỗi lệnh, quan sát cửa sổ Register và chú ý
  - o Sự thay đổi giá trị của thanh ghi \$s0
  - o Sự thay đổi giá trị của thanh ghi \$pc
- Ở cửa sổ Data Segment, hãy click vào hộp combo để chuyển tới quan sát các byte trong vùng lệnh .text.
  - o Kiểm tra xem các byte đầu tiên ở vùng lệnh trùng với cột nào trong cửa sổ Text Segment.

### Assignment 3: lệnh gán (giả lệnh)

Gõ chương trình sau vào công cụ MARS.

```
#Laboratory Exercise 2, Assignment 3
.text
 li $s0, 0x2110003d #pseudo instruction=2 basic instructions
 li $s1, 0x2 #but if the immediate value is small, one ins
```

Sau đó:

- Biên dịch và quan sát các lệnh mã máy trong cửa sổ Text Segment. Giải thích điều bất thường?

### Assignment 4: tính biểu thức $2x + y = ?$

Gõ chương trình sau vào công cụ MARS.

```
#Laboratory Exercise 2, Assignment 4
.text
 # Assign X, Y
 addi $t1, $zero, 5 # X = $t1 = ?
 addi $t2, $zero, -1 # Y = $t2 = ?
 # Expression Z = 2X + Y
 add $s0, $t1, $t1 # $s0 = $t1 + $t1 = X + X = 2X
 add $s0, $s0, $t2 # $s0 = $s0 + $t2 = 2X + Y
```

Sau đó:



- Sử dụng công cụ gỡ rối, Debug, chạy từng lệnh và dừng lại,
- Ở mỗi lệnh, quan sát cửa sổ Register và chú ý
  - o Sự thay đổi giá trị của các thanh ghi
  - o Sau khi kết thúc chương trình, xem kết quả có đúng không?
- Ở cửa sổ Text Segment, xem các lệnh **addi** và cho biết điểm tương đồng với hợp ngữ và mã máy. Từ đó kiểm nghiệm với khuôn mẫu của kiểu lệnh I




```
Y : .word -1 # Variable Y, word type, init value =
Z : .word # Variable Z, word type, no init value

.text # DECLARE INSTRUCTIONS
Load X, Y to registers
la $t8, X # Get the address of X in Data Segment
la $t9, Y # Get the address of Y in Data Segment
lw $t1, 0($t8) # $t1 = X
lw $t2, 0($t9) # $t2 = Y

Calculate the expression Z = 2X + Y with registers only
add $s0, $t1, $t1 # $s0 = $t1 + $t1 = X + X = 2X
add $s0, $s0, $t2 # $s0 = $s0 + $t2 = 2X + Y

Store result from register to variable Z
la $t7, Z # Get the address of Z in Data Segment
sw $s0, 0($t7) # Z = $s0 = 2X + Y
```

Sau đó:

- Biên dịch và quan sát các lệnh mã máy trong cửa sổ Text Segment.
  - o Lệnh **la** được biên dịch như thế nào?
- Ở cửa sổ Label và quan sát địa chỉ của X, Y, Z.
  - o So sánh chúng với hằng số khi biên dịch lệnh **la** thành mã máy
  - o Click đúp vào các biến X, Y, Z để công cụ tự động nhảy tới vị trí của biến X, Y, Z trong bộ nhớ ở cửa sổ Data Segment. Hãy bảo đảm các giá trị đó đúng như các giá trị khởi tạo.
- Sử dụng công cụ gỡ rối, Debug, chạy từng lệnh và dừng lại, 
- Ở mỗi lệnh, quan sát cửa sổ Register và chú ý
  - o Sự thay đổi giá trị của các thanh ghi
  - o Xác định vai trò của lệnh **lw** và **sw**
- Ghi nhớ qui tắc xử lý
  - o Đưa tất cả các biến vào thanh ghi bằng cặp lệnh **la, lw**
  - o Xử lý dữ liệu trên thanh ghi
  - o Lưu kết quả từ thanh ghi trở lại biến bằng cặp lệnh **la, sw**
- Tìm hiểu thêm các lệnh **lb, sb**

# Laboratory 3

## Load/ Store , Jump & Branch instructions

---

### Goals

After this laboratory exercise, you should know how to use load/store, jump and branch instructions. You should also be able to implement high level programming language structures such as conditional statement (if-then-else), loop and selection statement (switch-case)

---

### Literature

Behrooz Parhami (CAMS): Section 5.4, 5.5

---

### Preparation

Before start this laboratory, you should review textbook and read the entire laboratory exercise in detail. You also need to review Laboratory Exercise 2 and try to experience MARS by yourself.

---

### Assignments at Home and at Lab

#### Home Assignment 1

This home assignment implements “if-then-else” statement using some fundamental instructions, such as slt, addi, jump and branch.

```
if (i<=j)
 x=x+1;
 z=1;
else
 y=y-1;
 z=2*z;
```

At first, you should draw the algorithm chart for this statement. After that, you read this example carefully, try to clarify the function of each instructions.

```
#Laboratory Exercise 3, Home Assignment 1
start:
 slt $t0,$s2,$s1 # j<i
 bne $t0,$zero,else # branch to else if j<i
 addi $t1,$t1,1 # then part: x=x+1
 addi $t3,$zero,1 # z=1
 j endif # skip "else" part
else: addi $t2,$t2,-1 # begin else part: y=y-1
 add $t3,$t3,$t3 # z=2*z
endif:
```

## Home Assignment 2

The following example demonstrates how to implement loop statement. This program computes the sum of elements of array A.

```
loop: i=i+step;
 Sum=sum+A[i];
 If(I !=n) goto loop;
```

Assuming that the index i, the starting address of A, the comparison constant n, step and sum are found in registers \$s1, \$s2, \$s3, \$s4 and \$s5, respectively. You should try to understand each line in this code.

```
#Laboratory 3, Home Assignment 2
.text
loop: add $s1,$s1,$s4 #i=i+step
 add $t1,$s1,$s1 #t1=2*s1
 add $t1,$t1,$t1 #t1=4*s1
 add $t1,$t1,$s2 #t1 store the address of A[i]
 lw $t0,0($t1) #load value of A[i] in $t0
 add $s5,$s5,$t0 #sum=sum+A[i]
 bne $s1,$s3,loop #if i != n, goto loop
```

## Home Assignment 3

A switch/case statement allows multiway branching based on the value of an integer variable. In the following example, the switch variable test can assume one of the three values in [0, 2] and a different action is specified for each case.

```
switch(test) {
 case 0:
 a=a+1; break;
 case 1:
 a=a-1; break;
 case 2:
 b=2*b; break;
}
```

Assuming that **a** and **b** are stored in registers \$s2 and \$s3. You should read this code section carefully, understand how to implement switch/case statement.

```
#Laboratory Exercise 3, Home Assignment 3
.data
test: .word 1
.text
 la $s0, test #load the address of test variable
 lw $s1, 0($s0) #load the value of test to register $t1
 li $t0, 0 #load value for test case
 li $t1, 1
 li $t2, 2
 beq $s1, $t0, case_0
 beq $s1, $t1, case_1
 beq $s1, $t2, case_2
 j default
case_0: addi $s2, $s2, 1 #a=a+1
 j continue
case_1: sub $s2, $s2, $t1 #a=a-1
 j continue
case_2: add $s3, $s3, $s3 #b=2*b
 j continue
default:
continue:
```

### **Assignment 1**

Create a new project to implement the code in Home Assignment 1. Initialize for  $i$  and  $j$  variable. Compile and upload to the simulator. Run this program step by step, observe the changing of memory and the content of registers at each step.

### **Assignment 2**

Create a new project implementing the code in Home Assignment 2. Initialize for  $i$ ,  $n$ , step, sum variables and array  $A$ . Compile and upload to the simulator. Run this program step by step, observe the changing of memory and the content of registers by each step. Try to test with some more cases (change the value of variables).

### **Assignment 3**

Create a new project implementing the code in Home Assignment 3. Compile and upload to the simulator. Run this program step by step; observe the changing of memory and the content of registers by each step. Change the value of test variable and run this program some times to check all cases.

### **Assignment 4**

Modify the Assignment 1, so that the condition tested is

- a.  $i < j$
- b.  $i \geq j$
- c.  $i + j \leq 0$
- d.  $i + j > m + n$

### **Assignment 5**

Modify the Assignment 2, so that the condition tested at the end of the loop is

- a.  $i < n$
- b.  $i \leq n$
- c.  $\text{sum} \geq 0$
- d.  $A[i] == 0$

### **Assignment 6**

Using all of above instructions and statements, create a new project to implement this function: find the element with the largest absolute value in a list of integers. Assuming that this list is store in an integer array and we know the number of elements in it.

---

### **Conclusions**

Before you pass the laboratory exercise, think about the questions below:

- Which registers are affected by a branch instruction?



# Laboratory Exercise 4

## Arithmetic and Logical operation

---

### Goals

After this laboratory exercise, you should know how to use arithmetic, logical and shift instructions. In addition, you should also understand overflow in arithmetic operation and how to detect it.

---

### Literature

Behrooz Parhami (CAMS): Section 5.3

---

### Preparation

Before you start the exercise, you should review the textbook, section 5.3 and read this laboratory carefully. You should also review the Laboratory Exercise 2.

---

### Assignments at Home and at Lab

#### Home Assignment 1

The sum of two 32-bit integers may not be representable in 32 bits. In this case, we say that an overflow has occurred. Overflow is possible only with operands of the same sign. For two nonnegative (negative) operands, if the sum obtained is less (greater) than either operand, overflow has occurred. The following program detects overflow based on this rule. Two operands are stored in register \$s1 and \$s2, the sum is stored in register \$s3. If overflow occurs, \$t0 register is set to 1 and clear to 0 otherwise.

```
#Laboratory Exercise 4, Home Assignment 1
.text
start:
 li $t0,0 #No Overflow is default status
 addu $s3,$s1,$s2 # s3 = s1 + s2
 xor $t1,$s1,$s2 #Test if $s1 and $s2 have the same sign

 bltz $t1,EXIT #If not, exit
 slt $t2,$s3,$s1 #Test if $s1 and $s2 is negative?
 bltz $s1,NEGATIVE #s1 and $s2 are positive
 beq $t2,$zero,EXIT #s1 and $s2 are negative
 # if $s3 > $s1 then the result is not overflow
 j OVERFLOW
NEGATIVE:
 bne $t2,$zero,EXIT #s1 and $s2 are negative
 # if $s3 < $s1 then the result is not overflow
OVERFLOW:
 li $t0,1 #the result is overflow
EXIT:
```

## Home Assignment 2

The following program demonstrates how to use logical instructions to extract information from one register. We can extract one bit or more according to the mask we use. Read this example carefully and explain each lines of code.

```
#Laboratory Exercise 4, Home Assignment 2
.text
 li $s0, 0x0563 #load test value for these function
 andi $t0, $s0, 0xff #Extract the LSB of $s0
 andi $t1, $s0, 0x0400 #Extract bit 10 of $s0
```

## Home Assignment 3

This example show how the shift operations used to implement other instructions, such as multiply by a small power of 2.

```
#Laboratory Exercise 4, Home Assignment 3
.text
 li $s0, 1 #s0=1
 sll $s1, $s0, 2 #s1=s0*4
```

## Assignment 1

Create a new project to implement the Home Assignment 1. Compile and upload to simulator. Initialize two operands (register \$s1 and \$s2), run this program step by step, observe memory and registers value.

## Assignment 2

Write a program to do the following tasks:

- Extract MSB of \$s0
- Clear LSB of \$s0
- Set LSB of \$s0 (bits 7 to 0 are set to 1)
- Clear \$s0 (s0=0, must use logical instructions)

MSB: Most Significant Byte

LSB: Least Significant Byte

s0 = 0x 1 2 3 4 5 6 7 8  
          ↓          ↓  
        MSB      LSB

## Assignment 3

Pseudo instructions in MIPS are not-directly-run-on-MIPS-processor instructions which need to be converted to real-instructions of MIPS. Re-write the following pseudo instructions using real-instructions understood by MIPS processors:

- abs     \$s0, s1  
      s0 <= | \$s1 |
- move    \$s0, s1  
      s0 <= \$s1
- not     \$s0  
      s0 <= bit invert (s0)
- ble     \$s1, s2, L

```
if ($s1 <= $s2)
 j L
```

### **Assignment 4**

To detect overflow in additional operation, we also use other rule than the one in Assignment 1. This rule is: when add two operands that have the same sign, overflow will occur if the sum doesn't have the same sign with either operands. You need to use this rule to write another overflow detection program.

### **Assignment 5**

Write a program that implement multiply by a small power of 2. (2, 4, 8, 16, etc for example).

---

### **Conclusions**

Before you pass the laboratory exercise, think about the questions below:

- What is the difference between SLLV and SLL instructions?
- What is the difference between SRLV and SRL instructions?

# Laboratory Exercise 5

## Character string with SYSCALL function, and sorting

---

### Goals

After this laboratory exercise, you should understand the mechanism of storing ASCII and Unicode string. You will be able to program to process string and put string to console. In addition, you should know how to sort a list of elements.

---

### Literature

Patterson, Hennessy (COD): section 2.8, 2.13

---

### Preparation

Before you start the exercise, you should review the textbook, section 6.1 and read this laboratory carefully. You should also read the Mips Lab Environment Reference to find the usage of printf, putchar procedures ... and so on.

---

### About SYSCALL

A number of system services, mainly for input and output, are available for use by your MIPS program. They are described in the table below.

MIPS register contents are not affected by a system call, except for result registers as specified in the table below.

### How to use SYSCALL system services

1. Load the service number in register \$v0.
2. Load argument values, if any, in \$a0, \$a1, \$a2, or \$f12 as specified.
3. Issue the SYSCALL instruction.
4. Retrieve return values, if any, from result registers as specified.

#### Example: display the value stored in \$t0 on the console

```
li $v0, 1 # service 1 is print integer
li $a0, 0x307 # the interger to be printed is 0x307
syscall # execute
```

### Table of Frequently Available Services

| Service               | Code in \$v0 | Arguments               | Result |
|-----------------------|--------------|-------------------------|--------|
| print decimal integer | 1            | \$a0 = integer to print |        |

|                                     |       |                                                                                                               |                                                                                                                                                                               |
|-------------------------------------|-------|---------------------------------------------------------------------------------------------------------------|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| <b>print string</b>                 | 4     | \$a0 = address of null-terminated string to print                                                             |                                                                                                                                                                               |
| <b>read integer</b>                 | 5     |                                                                                                               | \$v0 contains integer read                                                                                                                                                    |
| <b>read string</b>                  | 8     | \$a0 = address of input buffer<br>\$a1 = maximum number of characters to read                                 | <i>See note below table</i>                                                                                                                                                   |
| <b>exit</b>                         | 10    | (terminate execution)                                                                                         |                                                                                                                                                                               |
| <b>print character</b>              | 11    | \$a0 = character to print                                                                                     | <i>See note below table</i>                                                                                                                                                   |
| <b>read character</b>               | 12    |                                                                                                               | \$v0 contains character read                                                                                                                                                  |
| <b>open file</b>                    | 13    | \$a0 = address of null-terminated string containing filename<br>\$a1 = flags<br>\$a2 = mode                   | \$v0 contains file descriptor (negative if error). <i>See note below table</i>                                                                                                |
| <b>read from file</b>               | 14    | \$a0 = file descriptor<br>\$a1 = address of input buffer<br>\$a2 = maximum number of characters to read       | \$v0 contains number of characters read (0 if end-of-file, negative if error). <i>See note below table</i>                                                                    |
| <b>write to file</b>                | 15    | \$a0 = file descriptor<br>\$a1 = address of output buffer<br>\$a2 = number of characters to write             | \$v0 contains number of characters written (negative if error). <i>See note below table</i>                                                                                   |
| <b>close file</b>                   | 16    | \$a0 = file descriptor                                                                                        |                                                                                                                                                                               |
| <b>exit2 (terminate with value)</b> | 17    | \$a0 = termination result                                                                                     | <i>See note below table</i>                                                                                                                                                   |
| <b>time (system time)</b>           | 30    |                                                                                                               | \$a0 = low order 32 bits of system time<br>\$a1 = high order 32 bits of system time.<br><i>See note below table</i>                                                           |
| <b>MIDI out</b>                     | 31    | \$a0 = pitch (0-127)<br>\$a1 = duration in milliseconds<br>\$a2 = instrument (0-127)<br>\$a3 = volume (0-127) | Generate tone and return immediately. <i>See note below table</i>                                                                                                             |
| <b>sleep</b>                        | 32    | \$a0 = the length of time to sleep in milliseconds.                                                           | Causes the MARS Java thread to sleep for (at least) the specified number of milliseconds. This timing will not be precise, as the Java implementation will add some overhead. |
| <b>MIDI out synchronous</b>         | 33    | \$a0 = pitch (0-127)<br>\$a1 = duration in milliseconds<br>\$a2 = instrument (0-127)<br>\$a3 = volume (0-127) | Generate tone and return upon tone completion. <i>See note below table</i>                                                                                                    |
| <b>print integer in hexadecimal</b> | 34    | \$a0 = integer to print                                                                                       | Displayed value is 8 hexadecimal digits, left-padding with zeroes if necessary.                                                                                               |
| <b>print integer in binary</b>      | 35    | \$a0 = integer to print                                                                                       | Displayed value is 32 bits, left-padding with zeroes if necessary.                                                                                                            |
| <b>print integer as unsigned</b>    | 36    | \$a0 = integer to print                                                                                       | Displayed as unsigned decimal value.                                                                                                                                          |
| <b>(not used)</b>                   | 37-39 |                                                                                                               |                                                                                                                                                                               |
| <b>set seed</b>                     | 40    | \$a0 = i.d. of pseudorandom                                                                                   | No values are returned. Sets the seed of the corresponding underlying Java                                                                                                    |

|                          |    |                                                                                                                                                                                                                                                                                                                                                                           |                                                                                                                                                                                                                                                                                                                                                                                                      |
|--------------------------|----|---------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
|                          |    | number generator (any int).<br>\$a1 = seed for corresponding pseudorandom number generator.                                                                                                                                                                                                                                                                               | pseudorandom number generator ( <code>java.util.Random</code> ). <i>See note below table</i>                                                                                                                                                                                                                                                                                                         |
| <b>random int</b>        | 41 | \$a0 = i.d. of pseudorandom number generator (any int).                                                                                                                                                                                                                                                                                                                   | \$a0 contains the next pseudorandom, uniformly distributed int value from this random number generator's sequence. <i>See note below table</i>                                                                                                                                                                                                                                                       |
| <b>random int range</b>  | 42 | \$a0 = i.d. of pseudorandom number generator (any int).<br>\$a1 = upper bound of range of returned values.                                                                                                                                                                                                                                                                | \$a0 contains pseudorandom, uniformly distributed int value in the range 0 = [int] [upper bound], drawn from this random number generator's sequence. <i>See note below table</i>                                                                                                                                                                                                                    |
| <b>ConfirmDialog</b>     | 50 | \$a0 = address of null-terminated string that is the message to user                                                                                                                                                                                                                                                                                                      | \$a0 contains value of user-chosen option<br>0: Yes<br>1: No<br>2: Cancel                                                                                                                                                                                                                                                                                                                            |
| <b>InputDialogInt</b>    | 51 | \$a0 = address of null-terminated string that is the message to user                                                                                                                                                                                                                                                                                                      | \$a0 contains int read<br>\$a1 contains status value<br>0: OK status<br>-1: input data cannot be correctly parsed<br>-2: Cancel was chosen<br>-3: OK was chosen but no data had been input into field                                                                                                                                                                                                |
| <b>InputDialogString</b> | 54 | \$a0 = address of null-terminated string that is the message to user<br>\$a1 = address of input buffer<br>\$a2 = maximum number of characters to read                                                                                                                                                                                                                     | <i>See Service 8 note below table</i><br>\$a1 contains status value<br>0: OK status. Buffer contains the input string.<br>-2: Cancel was chosen. No change to buffer.<br>-3: OK was chosen but no data had been input into field. No change to buffer.<br>-4: length of the input string exceeded the specified maximum. Buffer contains the maximum allowable input string plus a terminating null. |
| <b>MessageDialog</b>     | 55 | \$a0 = address of null-terminated string that is the message to user<br>\$a1 = the type of message to be displayed:<br>0: error message, indicated by Error icon<br>1: information message, indicated by Information icon<br>2: warning message, indicated by Warning icon<br>3: question message, indicated by Question icon<br>other: plain message (no icon displayed) | N/A                                                                                                                                                                                                                                                                                                                                                                                                  |



|                            |    |                                                                                                                                                                    |     |
|----------------------------|----|--------------------------------------------------------------------------------------------------------------------------------------------------------------------|-----|
| <b>MessageDialogInt</b>    | 56 | \$a0 = address of null-terminated string that is an information-type message to user<br>\$a1 = int value to display in string form after the first string          | N/A |
| <b>MessageDialogString</b> | 59 | \$a0 = address of null-terminated string that is an information-type message to user<br>\$a1 = address of null-terminated string to display after the first string | N/A |

## 1. print decimal integer

print an integer to standard output (the console).

*Argument(s):*

\$v0 = 1  
\$a0 = number to be printed

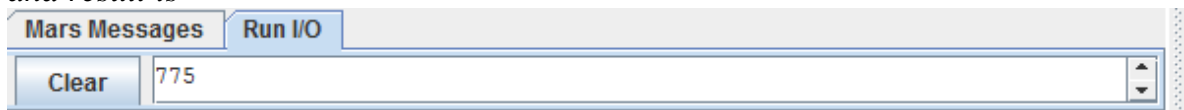
*Return value:*

none

*Example:*

```
li $v0, 1 # service 1 is print integer
li $a0, 0x307 # the interger to be printed is 0x307
syscall # execute
```

*and result is*



## 2. MessageDialogInt

show an integer to a information-type message dialog.

*Argument(s):*

\$v0 = 56  
\$a0 = address of the null-terminated message string  
\$a1 = int value to display in string form after the first

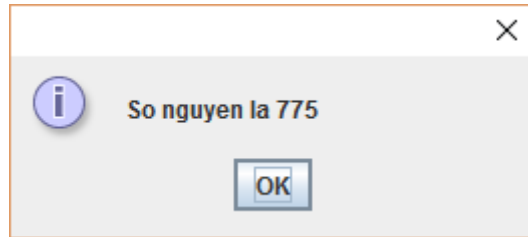
*string Return value:*

none

*Example:*

```
.data
Message: .asciiz "So nguyen la "
.text
li $v0, 56
la $a0, Message
li $a1, 0x307 # the interger to be printed is 0x307
syscall # execute
```

*and result is*



### 3. print string

Formatted print to standard output (the console).

*Argument(s):*

\$v0 = 1  
\$a0 = value to be printed

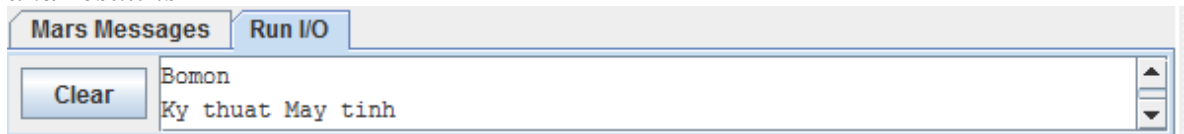
*Return value:*

none

*Example:*

```
.data
Message: .asciiz "Bomon \nKy thuat May tinh"
.text
li $v0, 4
la $a0, Message
syscall
```

*and result is*



### 4. MessageDialogString

Show a string to a information-type message dialog

*Argument(s):*

\$v0 = 59  
\$a0 = address of the null-terminated message string  
\$a1 = address of null-terminated string to display

*Return value:*

none

*Example:*

```
.data
Message: .asciiz "Bomon \nKy thuat May tinh:"
Address: .asciiz " phong 502, B1"
.text
li $v0, 59
la $a0, Message
la $a1, Address
syscall
```

*and result is*



## 5. read integer

Get a integer from standard input (the keyboard).

*Argument(s):*

\$v0 = 5

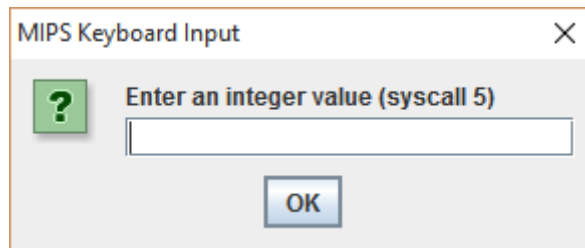
*Return value:*

\$v0 = contains integer read

*Example:*

```
li $v0, 5
syscall
```

and result is



## 6. InputDialogInt

Show a message dialog to read a integer with content parser

*Argument(s):*

\$v0 = 51

\$a0 = address of the null-terminated message string

*Return value:*

\$a0 = contains int read

\$a1 contains status value

0: OK status

-1: input data cannot be correctly parsed

-2: Cancel was chosen

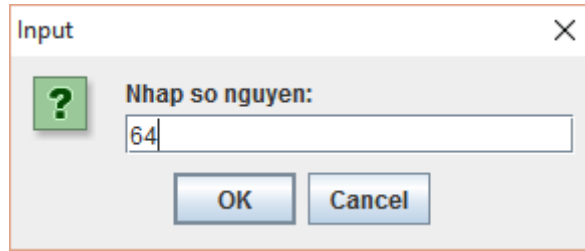
-3: OK was chosen but no data had been input into

field

*Example:*

```
.data
Message: .asciiz "Nhap so nguyen:"
.text
li $v0, 51
la $a0, Message
syscall
```

and result is



## 7. read string

Get a string from standard input (the keyboard).

*Argument(s):*

\$v0 = 8  
\$a0 = address of input buffer  
\$a1 = maximum number of characters to read

*Return value:*

none

*Remarks:*

For specified length n, string can be no longer than n-1.

- If less than that, adds newline to end.
- In either case, then pads with null byte

Just in special cases:

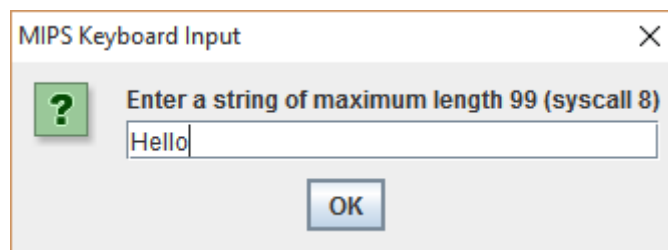
If n = 1, input is ignored and null byte placed at buffer address.

If n < 1, input is ignored and nothing is written to the buffer.

*Example:*

```
.data
Message: .space 100 # Buffer 100 byte chua chuoi ki tu can
.text
li $v0, 8
la $a0, Message
li $a1, 100
syscall
```

and result is



| Address    | Value (+0)  | Value (+4)  | Value (+8)  | Value (+c)  | Value (+10) |
|------------|-------------|-------------|-------------|-------------|-------------|
| 0x10010000 | 1 1 e H     | \0 \0 \n o  | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 |
| 0x10010020 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 |
| 0x10010040 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 |
| 0x10010060 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 | \0 \0 \0 \0 |

## 8. InputDialogString

Show a message dialog to read a string with content parser

*Argument(s):*

\$v0 = 54  
\$a0 = address of the null-terminated message string  
\$a1 = address of input buffer  
\$a2 = maximum number of characters to read

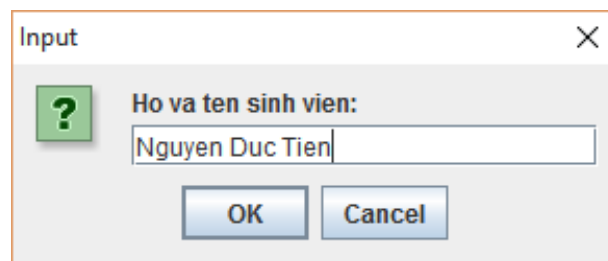
*Return value:*

\$a1 contains status value  
0: OK status  
-2: OK was chosen but no data had been input into field.  
No change to buffer.  
-3: OK was chosen but no data had been input into field  
-4: length of the input string exceeded the specified maximum. Buffer contains the maximum allowable input string plus a terminating null.

*Example:*

```
.data
Message: .asciiz "Ho va ten sinh vien:"
string: .space 100
.text
li $v0, 54
la $a0, Message
la $a1, string
la $a2, 100
syscall
```

*and result is*



## 9. print character

Print a character to standard output (the console).

*Argument(s):*

\$v0 = 11  
\$a0 = character to print (at the lowest significant byte)

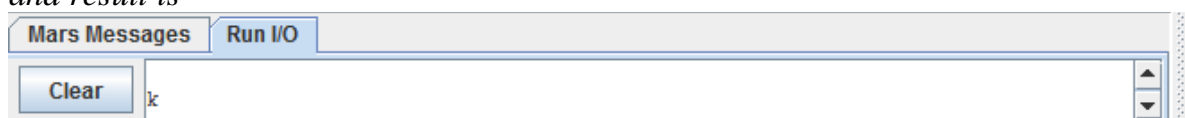
*Return value:*

none

*Example:*

```
li $v0, 11
li $a0, 'k'
syscall
```

*and result is*



## 10. read character

Get a character from standard output (the keyboard).

*Argument(s):*

\$v0 = 12

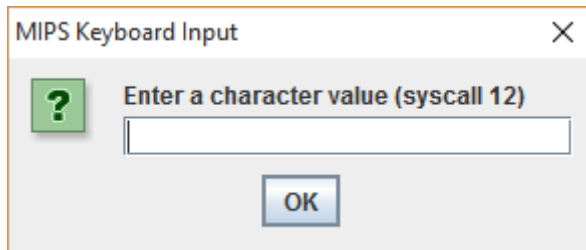
*Return value:*

\$v0 contains character read

*Example:*

```
li $v0, 12
syscall
```

*and result is*



## 11. ConfirmDialog

Show a message bog with 3 button: Yes | No | Cancel

*Argument(s):*

\$v0 = 50

\$a0 = address of the null-terminated message string

*Return value:*

\$a0 = contains value of user-chosen option

0: Yes

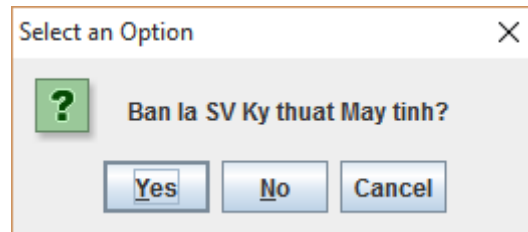
1: No

2: Cancel

*Example:*

```
.data
Message: .asciiz "Ban la SV Ky thuat May tinh?"
.text
li $v0, 50
la $a0, Message
syscall
```

*and result is*



## 12. MessageDialog

Show a message bog with icon and button OK only



*Argument(s):*

\$v0 = 55  
\$a0 = address of the null-terminated message string  
\$a1 = the type of message to be displayed:  
0: error message, indicated by Error icon  
1: information message, indicated by Information icon  
2: warning message, indicated by Warning icon  
3: question message, indicated by Question icon  
other: plain message (no icon displayed)

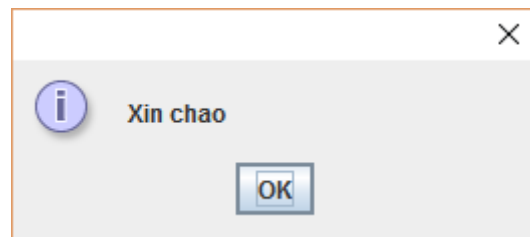
*Return value:*

none

*Example:*

```
.data
Message: .asciiz "Xin chao"
.text
li $v0, 55
la $a0, Message
syscall
```

*and result is*



### 13. MIDI out

Make a sound

*Argument(s):*

\$v0 = 31  
\$a0 = pitch (0-127)  
\$a1 = duration in milliseconds  
\$a2 = instrument (0-127)  
\$a3 = volume (0-127)

*Return value:*

Generate tone and return immediately

*Example:*

```
li $v0, 33
li $a0, 42 #pitch
li $a1, 2000 #time
li $a2, 0 #musical instrument
li $a3, 212 #volume
```

### 14. MIDI out synchronous

Make a sound

*Argument(s):*

\$v0 = 33

\$a0 = pitch (0-127)  
\$a1 = duration in milliseconds  
\$a2 = instrument (0-127)  
\$a3 = volume (0-127)

*Return value:*

Generate tone and return upon tone completion

*Example:*

```
li $v0, 33
li $a0, 42 #pitch
li $a1, 2000 #time
li $a2, 0 #musical instrusment
li $a3, 212 #volume
syscall
```

## 15. Exit

Terminated the software. Make sense that there is no EXIT instruction in the Instruction Set of any processors. Exit is a service belongs to Operating System.

*Argument(s):*

\$v0 = 10

*Return value:*

none

*Example:*

```
li $v0, 10 #exit
syscall
```

## 16. Exit with code

Terminated the software. Make sense that there is no EXIT instruction in the Instruction Set of any processors. Exit is a service belongs to Operating System.

*Argument(s):*

\$v0 = 17  
\$a0 = termination result

*Return value:*

none

*Example:*

```
li $v0, 17 # exit
li $a0, 3 # with error code = 3
syscall
```

---

## Assignments at Home and at Lab

### Home Assignment 1

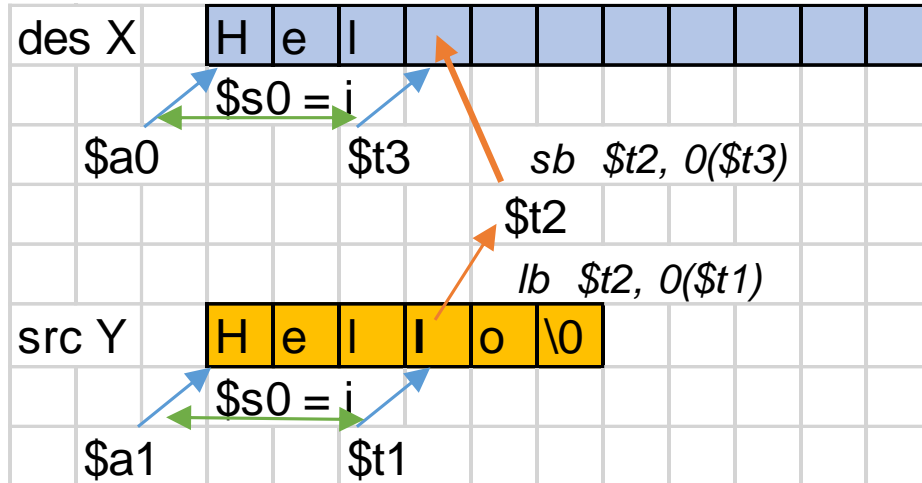
The following simple assembly program will display a welcome string. We use printf function for this purpose. Read this example carefully, pay attention to the way to pass parameters for printf function. Read Mips Lab Enviroment Reference for details.

```
#Laboratory Exercise 5, Home Assignment 1
.data
test: .asciiz "Hello World"
.text
```

```
li $v0, 4
la $a0, test
syscall
```

## Home Assignment 2

Procedure strcpy copies string y to string x using the null byte termination convention of C. Read this example carefully, try to understand all of this code section.



```
#Laboratory Exercise 5, Home Assignment 2
.data
x: .space 1000 # destination string x, empty
y: .asciiz "Hello" # source string y

.text
strcpy:
 add $s0,$zero,$zero #s0 = i=0
L1:
 add $t1,$s0,$a1 #t1 = s0 + a1 = i + y[0]
 # = address of y[i]
 lb $t2,0($t1) #t2 = value at t1 = y[i]
 add $t3,$s0,$a0 #t3 = s0 + a0 = i + x[0]
 # = address of x[i]
 sb $t2,0($t3) #x[i]= t2 = y[i]
 beq $t2,$zero,end_of_strcpy #if y[i]==0, exit
 nop
 addi bb$s0,$s0,1 #s0=s0 + 1 <-> i=i+1
 j L1 #next character
 nop
end_of_strcpy:
```

## Home Assignment 3

The following program count the length of a null-terminated string. Read this example carefully, analyse each line of code.

```
#Laboratory Exercise 5, Home Assignment 3
.data
string: .space 50
Message1: .asciiz "Nhap xau:"
Message2: .asciiz "Do dai la "
.text
main:
get_string: # TODO
```

```
get_length: la $a0, string # a0 = Address(string[0])
 xor $v0, $zero, $zero # v0 = length = 0
 xor $t0, $zero, $zero # t0 = i = 0
check_char: add $t1, $a0, $t0 # t1 = a0 + t0
 #= Address(string[0]+i)
 lb $t2, 0($t1) # t2 = string[i]
 beq $t2, $zero, end_of_str # Is null char?
 addi $v0, $v0, 1 # v0=v0+1->length=length+1
 addi $t0, $t0, 1 # t0=t0+1->i = i + 1
 j check_char
end_of_str:
end_of_get_length:
print_length: # TODO
```

## Assignment 1

Create a new project to implement the program in Home Assignment 1. Compile and upload to simulator. Run and observe the result. Go to data memory section, check how test string are stored and packed in memory.

## Assignment 2

Create a new project to print the sum of two register \$s0 and \$s1 according to this format:

“The sum of (s0) and (s1) is (result)”

## Assignment 3

Create a new project to implement the program in Home Assignment 2. Add more instructions to assign a test string for y variable, and implement *strcpy* function. Compile and upload to simulator. Run and observe the result.

## Assignment 4

Accomplish the Home Assignment 3 with *syscall* function to get a string from dialog, and show the length to message dialog.

## Assignment 5

Write a program that let user input a string. Input process will be terminated when user press Enter or then length of the string exceed 20 characters. Print the reverse string.

---

## Conclusions

Before you pass the laboratory exercise, think about the questions below:

- What the difference between the string in C and Java?
- In C, with 8 bytes, how many characters that we can store?
- In Java, with 8 bytes, how many characters that we can store?

# Laboratory Exercise 6

## Array and Pointer

### Goals

After this laboratory Exercise, you should understand how array and pointer are represented and you will be able to differentiate index and pointer using in stepping through an array of list.

### Literature

Patterson, Henessy (COD): section 2.8, 2.13

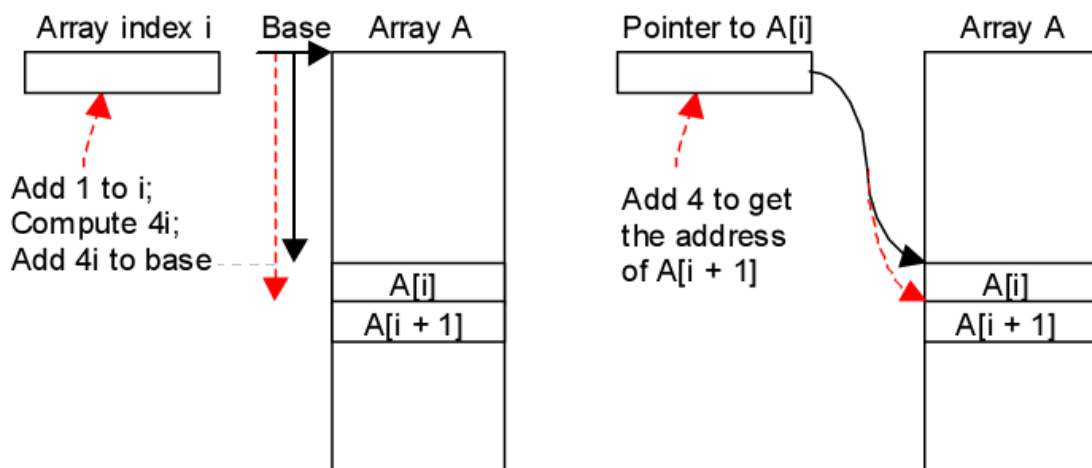
### Preparation

Before you start the exercise, you should review the textbook, section 6.1 and read this laboratory carefully.

### Array and Pointer

In a wide variety of programming tasks, it becomes necessary to step through an array or list, examining each of its elements in turn. For example, to determine the largest value in a list of integers, every element of the list must be examined. There are two basic ways of accomplishing this:

1. Index: use a register that holds the index  $i$  and increment the register in each step to effect moving from element  $i$  of the list to element  $i+1$
2. Pointer: use a register that points to (holds the address of) the list elements being examined and update it in each step to point to the next element.

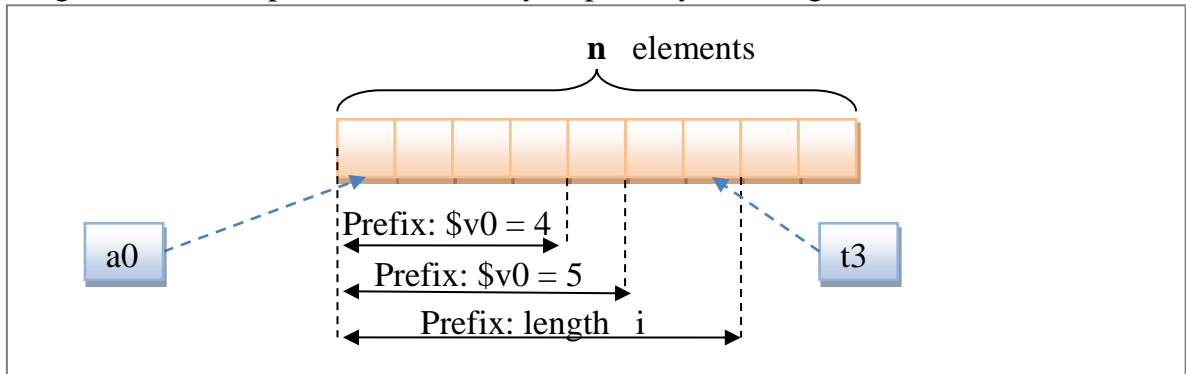


*Figure 4. Using the indexing method and the pointer updating method to step through the elements of an array*

## Assignments at Home and at Lab

### Home Assignment 1

Consider a list of integers of length  $n$ . A prefix of length  $i$  for the given list consist of the first  $i$  integers in the list, where  $0 \leq i \leq n$ . A maximum-sum prefix, as the name implies, is a prefix for which the sum of elements is the largest among all prefixes. For example, if the list is  $(2, -3, 2, 5, -4)$ , its maximum-sum prefix consists of the first four elements and the associated sum is  $2 - 3 + 2 + 5 = 6$ ; no other prefix of the given list has a larger sum. The following procedure uses indexing method to find the maximum-sum prefix in a list of integers. Read this procedure carefully, especially indexing method.



```
.data
A: .word -2, 6, -1, 3, -2

.text
main: la $a0,A
 li $a1,5
 j mspfx
 nop

continue:
lock: j lock
 nop

end_of_main:

#-----
#Procedure mspfx
@brief find the maximum-sum prefix in a list of integers
@param[in] a0 the base address of this list(A) need to be
processed
@param[in] a1 the number of elements in list(A)
@param[out] v0 the length of sub-array of A in which max sum
reaches.
@param[out] v1 the max sum of a certain sub-array
#-----
#Procedure mspfx
#function: find the maximum-sum prefix in a list of integers
#the base address of this list(A) in $a0 and the number of
#elements is stored in a1
mspfx: addi $v0,$zero,0 #initialize length in $v0 to 0
 addi $v1,$zero,0 #initialize max sum in $v1 to 0
 addi $t0,$zero,0 #initialize index i in $t0 to 0
 addi $t1,$zero,0 #initialize running sum in $t1 to 0
loop: add $t2,$t0,$t0 #put 2i in $t2
 add $t2,$t2,$t2 #put 4i in $t2
 add $t3,$t2,$a0 #put 4i+A (address of A[i]) in $t3
```

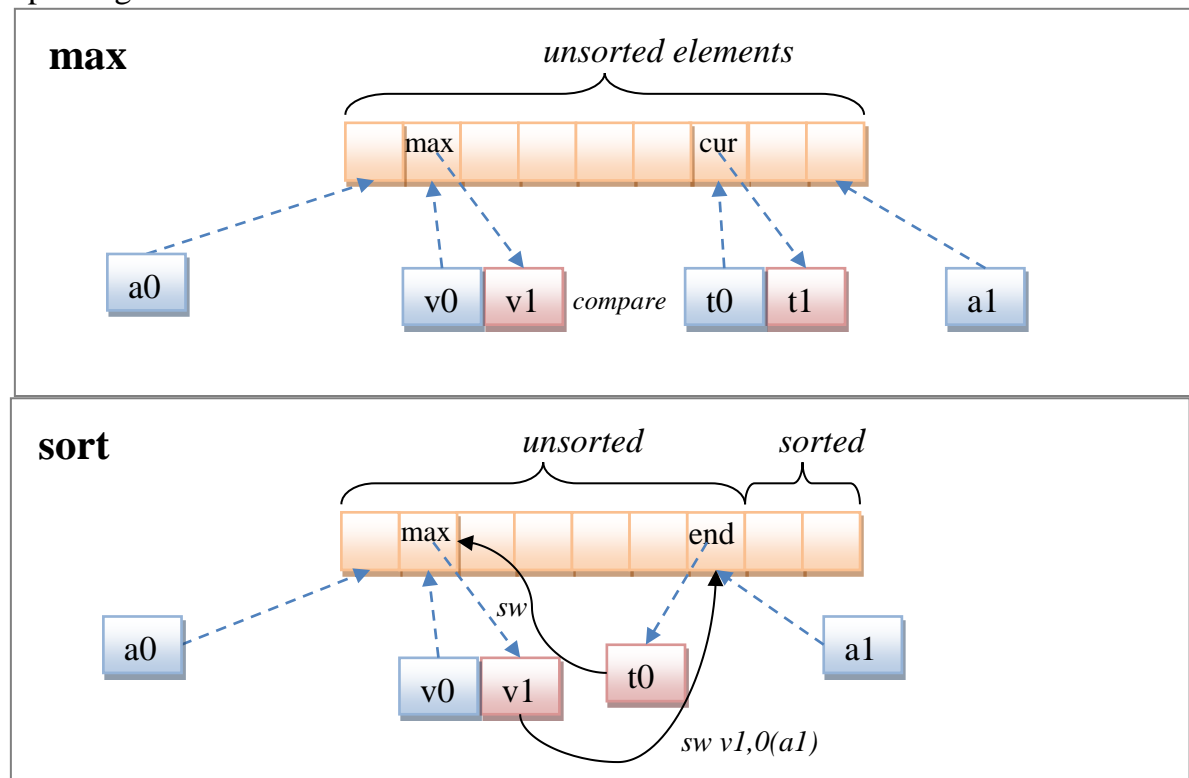
```

 lw $t4,0($t3) #load A[i] from mem(t3) into $t4
 add $t1,$t1,$t4 #add A[i] to running sum in $t1
 slt $t5,$v1,$t1 #set $t5 to 1 if max sum < new sum
 bne $t5,$zero,mdfy #if max sum is less, modify results
 j test #done?
mdfy: addi $v0,$t0,1 #new max-sum prefix has length i+1
 addi $v1,$t1,0 #new max sum is the running sum
test: addi $t0,$t0,1 #advance the index i
 slt $t5,$t0,$a1 #set $t5 to 1 if i<n
 bne $t5,$zero,loop #repeat if i<n
done: j continue
mspfx_end:

```

## Home Assignment 2

A given list of  $n$  numbers can be sorted in ascending order as follows. Find a largest number in the list (there may be more than one) and swap it with the last element in the list. The new last element is now in its proper position in sorted order. Now, sort the remaining  $n-1$  elements using the same step repeatedly. When only one element is left, sorting is complete. This method is known as selection sort. This example demonstrates this kind of sorting using pointer updating method. Read this procedure carefully and pay attention to pointer updating method.



```

.data
A: .word 7, -2, 5, 1, 5,6,7,3,6,8,8,59,5
Aend: .word

.text
main: la $a0,A #$a0 = Address(A[0])

```



```
 la $a1,Aend
 addi $a1,$a1,-4 #$a1 = Address(A[n-1])
 j sort #sort
after_sort: li $v0, 10 #exit
 syscall

end_main:
#-----
#procedure sort (ascending selection sort using pointer)
#register usage in sort program
#$a0 pointer to the first element in unsorted part
#$a1 pointer to the last element in unsorted part
#$t0 temporary place for value of last element
#$v0 pointer to max element in unsorted part
#$v1 value of max element in unsorted part
#-----
sort: beq $a0,$a1,done #single element list is sorted
 j max #call the max procedure
after_max: lw $t0,0($a1) #load last element into $t0
 sw $t0,0($v0) #copy last element to max location
 sw $v1,0($a1) #copy max value to last element
 addi $a1,$a1,-4 #decrement pointer to last element
 j sort #repeat sort for smaller list
done: j after_sort

#-----

#Procedure max
#function: fax the value and address of max element in the list
#$a0 pointer to first element
#$a1 pointer to last element
#-----

max:
 addi $v0,$a0,0 #init max pointer to first element
 lw $v1,0($v0) #init max value to first value
 addi $t0,$a0,0 #init next pointer to first
loop:
 beq $t0,$a1,ret #if next=last, return
 addi $t0,$t0,4 #advance to next element
 lw $t1,0($t0) #load next element into $t1
 slt $t2,$t1,$v1 #(next)<(max) ?
 bne $t2,$zero,loop #if (next)<(max), repeat
 addi $v0,$t0,0 #next element is new max element
 addi $v1,$t1,0 #next value is new max value
 j loop #change completed; now repeat
ret:
 j after_max
```

## Assignment 1

Create a new project to implement procedure in Home Assignment 1. Add code the main program and initialize data for the integer list. Compile and upload to simulator. Run this program step by step, observe the process of explore each element of the integer list using indexing method.

## Assignment 2

Create a new project to implement procedure in Home Assignment 2. Add code the main program and initialize data for the integer list. Compile and upload to simulator. Run this program step by step, observe the process of explore each element of the integer list using pointer updating method.

### **Assignment 3**

Write a procedure to implement bubble sort algorithm.

### **~~Assignment 4~~**

~~Write a procedure to implement insertion sort algorithm.~~

---

### **Conclusions**

Before you pass the laboratory exercise, think about the questions below:

- What the advantage and disadvantage of two methods: indexing and updating pointer?

# Laboratory Exercise 7

## Procedure calls, stack and parameters

### Goals

After this laboratory exercise, you should understand how to invoke a procedure and the mechanism of stack. In addition, you will be able to write some procedures that use stack to pass parameters and return results also.

### Literature

Behrooz Parhami (CAMS): Section 6.1

### Preparation

Before you start the exercise, you should review the textbook, section 6.1 and read this laboratory carefully.

### Procedure calls

A procedure is a subprogram that when called (initiated, invoked) performs a specific task, perhaps leading to one or more results, based on the input parameters (arguments) with which it is provided and returns to the point of call, having perturbed nothing else. In assembly language, a procedure is associated with a symbolic name that denotes its starting address. The jal instruction in MIPS is intended specifically for procedure calls: it performs the control transfer (unconditional jump) to the starting address of the procedure, while also saving the return address in register ra.

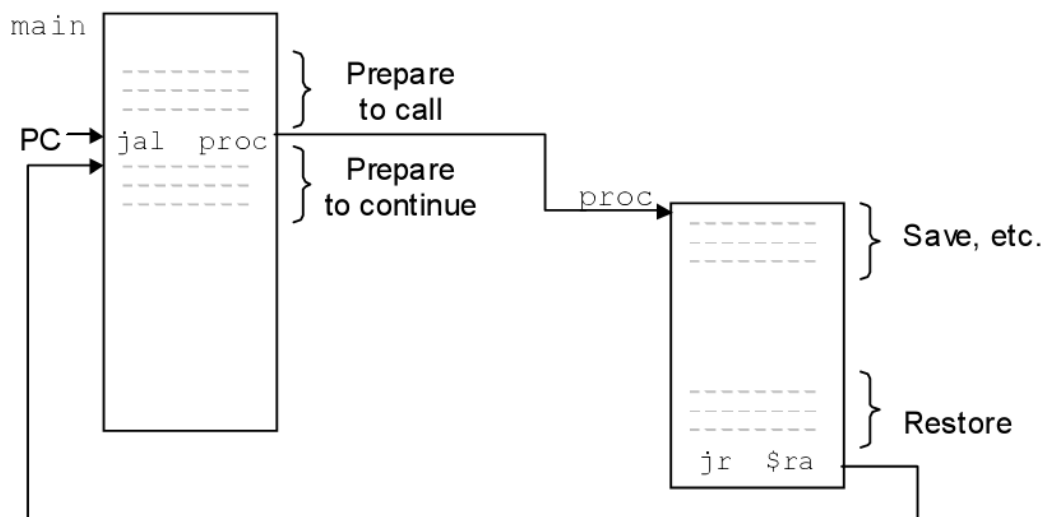


Figure 5. Relationship between the main program and a procedure

## Assignments at Home and at Lab

### Home Assignment 1

The following program uses abs procedure. This procedure is used to find the absolute value in an integer. This program uses two registers, \$a0 for input argument and \$v0 for result. You should read this example carefully, pay attention to how to write and invoke a procedure.

```
#Laboratory Exercise 7 Home Assignment 1
.text
main: li $a0,-45 #load input parameter
 jal abs #jump and link to abs procedure
 nop
 add $s0, $zero, $v0
 li $v0,10 #terminate
 syscall
endmain:
#-----
function abs
param[in] $a1 the interger need to be gained the absolute
value
return $v0 absolute value
#-----
abs:
 sub $v0,$zero,$a1 #put -(a0) in v0; in case (a0)<0

 bltz $a1,done #if (a0)<0 then done
 nop
 add $v0,$a1,$zero #else put (a0) in v0
done:
 jr $ra
```

### Home Assignment 2

In this example, procedure max has the function to find the largest elements in three of integers. These intergers are passed to max procedure through \$a0, \$a1 and \$a2 registers. Read this example carefully and try to explain each line of code.

```
#Laboratory Exercise 7, Home Assignment 2
.text
main: li $a0,2 #load test input
 li $a1,6
 li $a2,9
 jal max #call max procedure
 nop
endmain:
#-----
-
#Procedure max: find the largest of three integers
#param[in] $a0 integers
#param[in] $a1 integers
#param[in] $a2 integers
#return $v0 the largest value
#-----
-
max: add $v0,$a0,$zero #copy (a0) in v0; largest so far
 sub $t0,$a1,$v0 #compute (a1)-(v0)
```

```
 bltz $t0,okay #if (a1)-(v0)<0 then no change
 nop
 add $v0,$a1,$zero #else (a1) is largest thus far
okay: sub $t0,$a2,$v0 #compute (a2)-(v0)
 bltz $t0,done #if (a2)-(v0)<0 then no change
 nop
 add $v0,$a2,$zero #else (a2) is largest overall
done: jr $ra #return to calling program
```

### Home Assignment 3

The following program demonstrates the push and pop operations with stack. The value of two register \$s0 and \$s1 will be swap using stack. Read this example carefully, pay attention to adjust operation and the order of push and pop (sw and lw instructions).

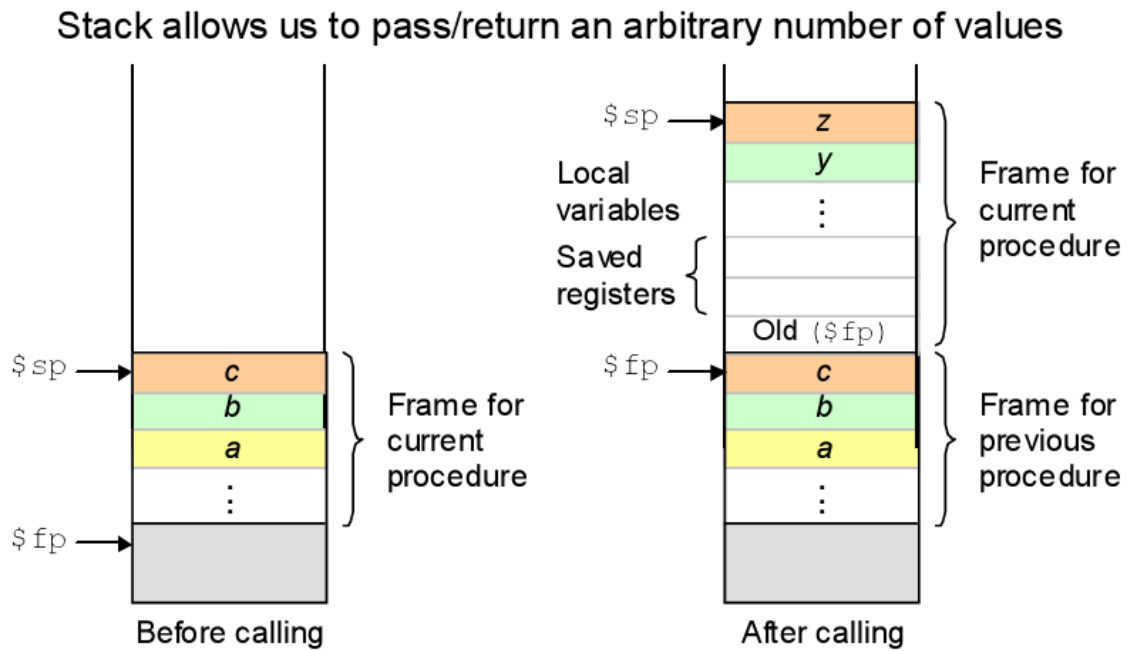
```
#Laboratory Exercise 7, Home Assignment 3
.text
push: addi $sp,$sp,-8 #adjust the stack pointer
 sw $s0,4($sp) #push $s0 to stack
 sw $s1,0($sp) #push $s1 to stack
work: nop
 nop
 nop
pop: lw $s0,0($sp) #pop from stack to $s0
 lw $s1,4($sp) #pop from stack to $s1
 addi $sp,$sp,8 #adjust the stack pointer
```

### Parameters and results

In this section, we answer the following unresolved questions relating to procedures

1. How to pass more than four input parameters to a procedure or receive more than two results from it?
2. Where does a procedure save its own parameters and immediate results when calling another procedure (nested calls)?

The stack is used in both cases. Before a procedure call, the calling program pushed the contents of any register that need to be save onto the top of the stack and follow these with any additional arguments for the procedure. The procedure can access these arguments in the stack. After the procedure terminates, the calling procedure expects to find the stack pointer undisturbed, thus allowing it to restore the save registers to their original states and proceed with its own computations. Thus, a procedure that uses the stack by modifying the stack pointer must save the content of the stack pointer at the outset and , at the end, restore sp to it original state. This is done by copying the stack pointer into the frame pointer register fp. Before doing this, however, the old contents of the frame pointer must be saved. Hence, while a procedure is executing, fp may hold the content of fp right before it was called; fp and sp together “frame” the area of the stack that is in use by the current procedure.



*Figure 6. Use of the stack by a procedure*

## Home Assignment 4

The following program is a recursive procedure for computing  $n!$ . Read this program carefully and pay attention to register `sp` and `fp` at the start and the end of each recursive step.

```
#Laboratory Exercise 7, Home Assignment 4
.data
Message: .asciiz "Ket qua tinh giai thua la: "

.text
main: jal WARP

print: add $a1, $v0, $zero # $a0 = result from N!
 li $v0, 56
 la $a0, Message
 syscall

quit: li $v0, 10 #terminate
 syscall

endmain:

#-----
#
#Procedure WARP: assign value and call FACT
#-----
#
WARP: sw $fp, -4($sp) #save frame pointer (1)
 addi $fp, $sp, 0 #new frame pointer point to the top (2)
 addi $sp, $sp, -8 #adjust stack pointer (3)
 sw $ra, 0($sp) #save return address (4)

 li $a0, 6 #load test input N
 jal FACT #call fact procedure
 nop

 lw $ra, 0($sp) #restore return address (5)
 addi $sp, $fp, 0 #return stack pointer (6)
```

```

 lw $fp,-4($sp) #return frame pointer (7)
 jr $ra
wrap_end:

#-----
#
#Procedure FACT: compute N!
#param[in] $a0 integer N
#return $v0 the largest value
#-----
#
#
FACT: sw $fp,-4($sp) #save frame pointer
 addi $fp,$sp,0 #new frame pointer point to stack's
top
 addi $sp,$sp,-12 #allocate space for $fp,$ra,$a0 in
stack
 sw $ra,4($sp) #save return address
 sw $a0,0($sp) #save $a0 register

 slti $t0,$a0,2 #if input argument N < 2
 beq $t0,$zero,recursive #if it is false ((a0 = N) >=2)
 nop
 li $v0,1 #return the result N!=1
 j done
 nop
recursive:
 addi $a0,$a0,-1 #adjust input argument
 jal FACT #recursive call
 nop
 lw $v1,0($sp) #load a0
 mult $v1,$v0 #compute the result
 mflo $v0
done: lw $ra,4($sp) #restore return address
 lw $a0,0($sp) #restore a0
 addi $sp,$fp,0 #restore stack pointer
 lw $fp,-4($sp) #restore frame pointer
 jr $ra #jump to calling
fact_end:

```

At the begin of WRAP, \$sp = 7fffeffc

|          |
|----------|
|          |
|          |
|          |
|          |
|          |
|          |
|          |
|          |
| \$ra (4) |
| \$fp     |
|          |

7fffeff4 ← new \$sp (3) addi \$sp,\$sp,-8  
7fffeff8 (1) sw \$fp,-4(\$sp)  
7fffeffc ← new \$fp (2) addi \$fp,\$sp,0

At the end of WRAP, \$sp = 7fffeff4

|      |
|------|
|      |
| \$a0 |
| \$ra |
| \$fp |
| \$a0 |
| \$ra |
| \$fp |
| \$ra |
| \$fp |
|      |

← new \$fp  
← new \$fp  
7fffeff4 → restore \$ra (5)  
7fffeff8 → restore \$fp (7)  
7fffeffc → restore \$sq (6)



## **Assignment 1**

Create a new project to implement the program in Home Assignment 1. Compile and upload to simulator. Change input parameters and observe the memory when run the program step by step. Pay attention to register pc, \$rato clarify invoking procedure process (Refer to figure 7).

## **Assignment 2**

Create a new project to implement the program in Home Assignment 2. Compile and upload to simulator. Change input parameters (register a0, a1, a2) and observe the memory when run the program step by step. Pay attention to register pc, \$rato clarify invoking procedure process (Refer to figure 7).

## **Assignment 3**

Create a new project to implement the program in Home Assignment 3. Compile and upload to simulator. Pass test value to registers \$s0 and \$s1, observe run process, pay attention to stack pointer. Goto memory space that pointed by sp register to view push and pop operations in detail.

## **Assignment 4**

Create a new project to implement the program in Home Assignment 4. Compile and upload to simulator. Pass test input through register a0, run this program and test result in register v0. Run this program in step by step mode, observe the changing of register pc, ra, sp and fp. Draw the stack through this recursive program in case of n=3 (compute 3!).

## **Assignment 5**

Write a procedure to find the largest, the smallest and these positions in a list of 8 elements that are stored in registers \$s0 through \$s7. For example:

Largest: 9,3      -> The largest element is stored in \$s3, largest value is 9

Smallest: -3,6      -> The smallest element is stored in \$s6, smallest value is -3

Tips: using stack to pass arguments and return results.

---

## **Conclusions**

Before you pass the laboratory exercise, think about the questions below:

- What registers that the Caller need to save by convention?
- What registers that the Callee need to save by convention?
- In push label of Home Assignment 3, could we change the order of adjust stack and store word operations? If yes, what should we have to modify?
- What is stack pointer?
- What is frame pointer?

## **Laboratory Exercise 8, 9**

### **Mini-projects**

---

#### ***Goals***

Review elementary knowledge about assembly programming to solve whole problems. Learn how to work in a group to promote group spirit.

---

#### ***Time***

- Week 8 : stay at home and do mid-term practice exercises
- Mid-term Week : stay at home and do mid-term practice exercises
- Week 9 : teachers check the practice exercises

# Laboratory Exercise 10

## Goals

After this laboratory exercise, you should understand the method to control peripheral devices via simulators.

## Literature

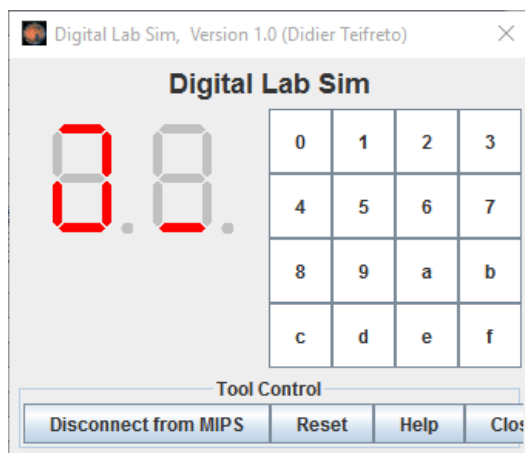
How does the CPU communicate with input and output devices such as the monitor or keyboard?

There are several ways. Intel machines have special instructions named in and out that communicate with I/O ports. These instructions are usually disabled for ordinary users, but they are used internally for communicating with I/O devices. This is called port-mapped I/O. However, we are going to look at a different method in which I/O devices have access to memory. The CPU can place data in memory that can be read by the I/O devices; likewise, the I/O devices can place data in memory for the CPU. This is called memory-mapped I/O or MMIO. (For more information, see P&H page 588 or Appendix B.8, or look it up online!)

## Assignments at Home and at Lab

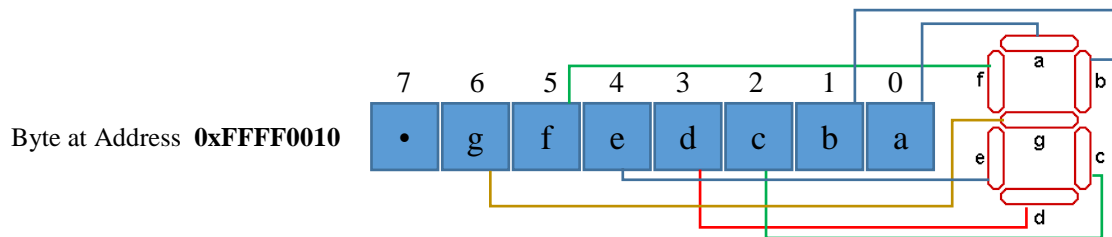
### Home Assignment 1 – LED PORT

Write a program using assembly language to show numbers from 0 to F to the 7-seg led.



To view the 7-segs, at the menu bar, click /Tools/Digi Lab Sim

Click Help to understand how to turn on the 7-seg led.



```
.eqv SEVENSEG_LEFT 0xFFFF0011 # Dia chi cua den led 7 doan trai.
 # Bit 0 = doan a;
 # Bit 1 = doan b; ...
 # Bit 7 = dau .
.eqv SEVENSEG_RIGHT 0xFFFF0010 # Dia chi cua den led 7 doan phai

.text
main:
 li $a0, 0x8 # set value for segments
 jal SHOW_7SEG_LEFT # show
 nop
 li $a0, 0x1F # set value for segments
 jal SHOW_7SEG_RIGHT # show
 nop
exit: li $v0, 10
 syscall
endmain:

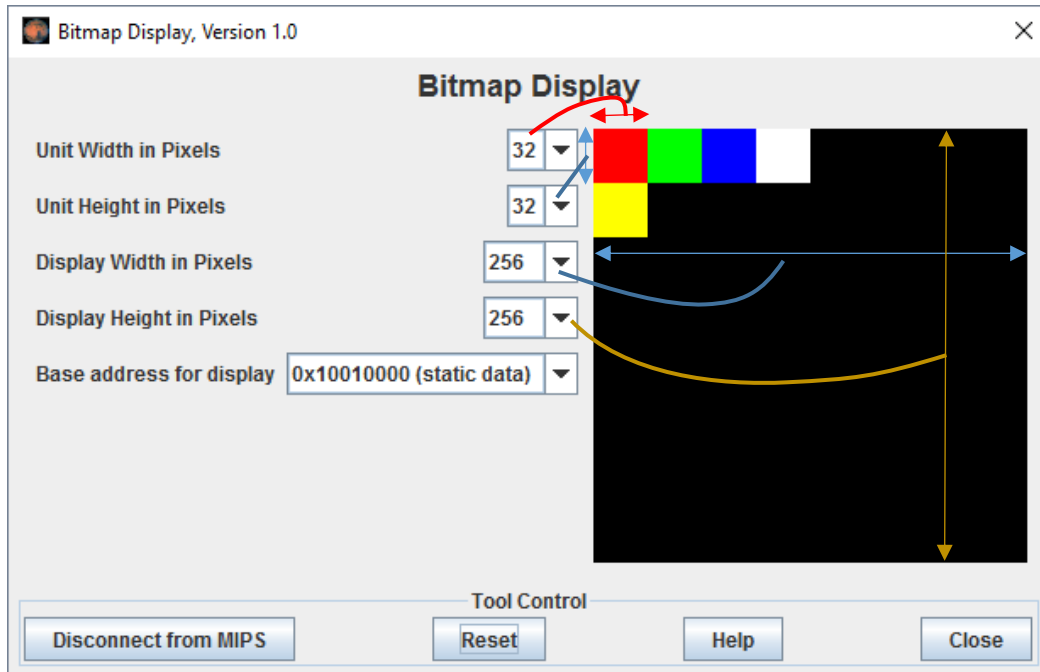
#-----
Function SHOW_7SEG_LEFT : turn on/off the 7seg
param[in] $a0 value to shown
remark $t0 changed
#-----
SHOW_7SEG_LEFT: li $t0, SEVENSEG_LEFT # assign port's address
 sb $a0, 0($t0) # assign new value
 nop
 jr $ra
 nop

#-----
Function SHOW_7SEG_RIGHT : turn on/off the 7seg
param[in] $a0 value to shown
remark $t0 changed
#-----
SHOW_7SEG_RIGHT: li $t0, SEVENSEG_RIGHT # assign port's address
 sb $a0, 0($t0) # assign new value
 nop
 jr $ra
 nop
```

## Home Assignment 2 - BITMAP DISPLAY

Bitmap Display like the graphic monitor, in which Windows OS draws windows, start button... In order to do that, developer should calculate color of all bitmap pixels on the screen and store these color value to the screen memory. Whenever we change a value in screen memory, the color of the respective pixel on the screen will be changed.

In MARS, in menu bar, click Tools / Bitmap Display to open the screen simulator



b

| 0  | R  | G  | B  |                      |
|----|----|----|----|----------------------|
| 00 | FF | 00 | 00 | 0x10010000 - pixel 0 |
| 00 | 00 | FF | 00 | 0x10010004 - pixel 1 |
| 00 | 00 | 00 | 00 | 0x10010008 - pixel 2 |
| 00 | FF | FF | FF | 0x1001000C - pixel 3 |

Each rectangular unit on the display represents one memory word in a contiguous address space starting with the specified base address (in above figure, base address is

0x10010000)

Value stored in that word will be interpreted as a 24-bit RGB

```
.eqv MONITOR_SCREEN 0x10010000 #Dia chi bat dau cua bo nho man hinh
.eqv RED 0x00FF0000 #Cac gia tri mau thuong su dung
.eqv GREEN 0x0000FF00
.eqv BLUE 0x000000FF
.eqv WHITE 0x00FFFFFF
.eqv YELLOW 0x00FFFF00
.text
 li $k0, MONITOR_SCREEN #Nap dia chi bat dau cua man hinh

 li $t0, RED
 sw $t0, 0($k0)
 nop

 li $t0, GREEN
 sw $t0, 4($k0)
 nop

 li $t0, BLUE
 sw $t0, 8($k0)
 nop

 li $t0, WHITE
 sw $t0, 12($k0)
```

```
nop

li $t0, YELLOW
sw $t0, 32($k0)
nop

li $t0, WHITE
lb $t0, 42($k0)
nop
```

### Home Assignment 3 – MARSBOT RIDER

The MarsBot is a virtual robot that has a very simple mode of operation. It travels around in two-dimensional space, optionally leaving a trail, or track, as it goes. It uses five words in memory:<sup>3</sup>

| Name       | Address    | Meaning                                               |
|------------|------------|-------------------------------------------------------|
| HEADING    | 0xffff8010 | Integer: An angle between 0 and 359                   |
| LEAVETRACK | 0xffff8020 | Boolean (0 or non-0): whether or not to leave a track |
| WHEREX     | 0xffff8030 | Integer: Current x-location of the MarsBot            |
| WHEREY     | 0xffff8040 | Integer: Current y-location of the MarsBot            |
| MOVING     | 0xffff8050 | Boolean: whether or not to move                       |

The CPU can place commands in the HEADING, LEAVETRACK, and MOVE locations; the robot can then change its direction of travel (using the HEADING value), turn on or turn off the pen" drawing the line (using the LEAVE-TRACK value), and can halt or resume moving (using the MOVING value).

```
.eqv HEADING 0xffff8010 # Integer: An angle between 0 and 359
 # 0 : North (up)
 # 90: East (right)
 # 180: South (down)
 # 270: West (left)
.eqv MOVING 0xffff8050 # Boolean: whether or not to move
.eqv LEAVETRACK 0xffff8020 # Boolean (0 or non-0):
 # whether or not to leave a track
.eqv WHEREX 0xffff8030 # Integer: Current x-location of
MarsBot
.eqv WHEREY 0xffff8040 # Integer: Current y-location of
MarsBot

.text
main: jal TRACK # draw track line
 nop
 addi $a0, $zero, 90 # Marsbot rotates 90* and start
running
 jal ROTATE
 nop
 jal GO
 nop
```

---

<sup>3</sup> <http://cs.allegheeny.edu/~rroos/cs210f2013>

```
sleep1: addi $v0,$zero,32 # Keep running by sleeping in 1000 ms
 li $a0,1000
 syscall

 jal UNTRACK # keep old track
 nop
 jal TRACK # and draw new track line
 nop

goDOWN: addi $a0, $zero, 180 # Marsbot rotates 180*
 jal ROTATE
 nop

sleep2: addi $v0,$zero,32 # Keep running by sleeping in 2000 ms
 li $a0,2000
 syscall

 jal UNTRACK # keep old track
 nop
 jal TRACK # and draw new track line
 nop

goLEFT: addi $a0, $zero, 270 # Marsbot rotates 270*
 jal ROTATE
 nop

sleep3: addi $v0,$zero,32 # Keep running by sleeping in 1000 ms
 li $a0,1000
 syscall

 jal UNTRACK # keep old track
 nop
 jal TRACK # and draw new track line
 nop

goASKEW: addi $a0, $zero, 120 # Marsbot rotates 120*
 jal ROTATE
 nop

sleep4: addi $v0,$zero,32 # Keep running by sleeping in 2000 ms
 li $a0,2000
 syscall

 jal UNTRACK # keep old track
 nop
 jal TRACK # and draw new track line
 nop

end_main:

#-----
GO procedure, to start running
param[in] none
#-----
GO: li $at, MOVING # change MOVING port
 addi $k0, $zero,1 # to logic 1,
 sb $k0, 0($at) # to start running
 nop
 jr $ra
 nop

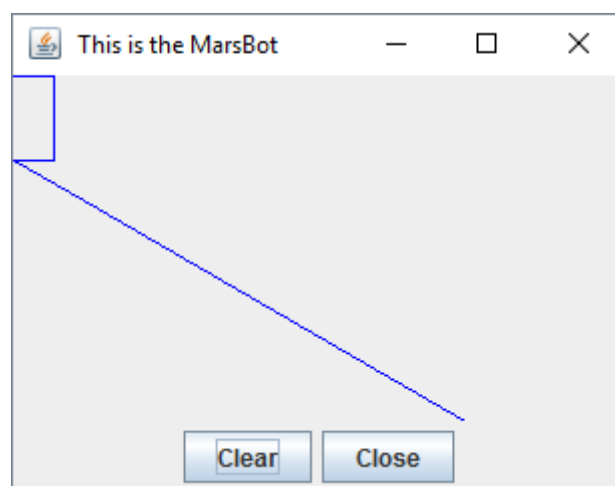
#-----
STOP procedure, to stop running
param[in] none
#-----
```

```
STOP: li $at, MOVING # change MOVING port to 0
 sb $zero, 0($at) # to stop
 nop
 jr $ra
 nop

#-----
TRACK procedure, to start drawing line
param[in] none
#-----
TRACK: li $at, LEAVETRACK # change LEAVETRACK port
 addi $k0, $zero, 1 # to logic 1,
 sb $k0, 0($at) # to start tracking
 nop
 jr $ra
 nop

#-----
UNTRACK procedure, to stop drawing line
param[in] none
#-----
UNTRACK:li $at, LEAVETRACK # change LEAVETRACK port to 0
 sb $zero, 0($at) # to stop drawing tail
 nop
 jr $ra
 nop

#-----
ROTATE procedure, to rotate the robot
param[in] $a0, An angle between 0 and 359
0 : North (up)
90: East (right)
180: South (down)
270: West (left)
#-----
ROTATE: li $at, HEADING # change HEADING port
 sw $a0, 0($at) # to rotate robot
 nop
 jr $ra
 nop
```

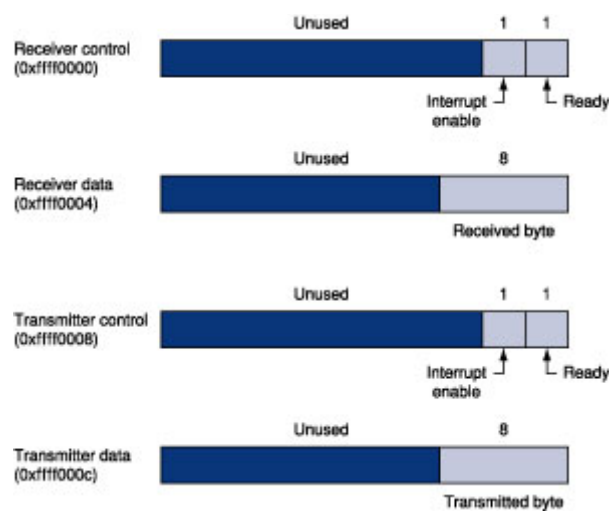




## Home Assignment 4 – KEYBOARD and DISPLAY MMIO

Use this program to simulate Memory-Mapped I/O (MMIO) for a keyboard input device and character display output device. It may be run either from MARS' Tools menu or as a stand-alone application.

While the tool is connected to MIPS, each keystroke in the text area causes the corresponding ASCII code to be placed in the Receiver Data register (low-order byte of memory word 0xffff0004), and the Ready bit to be set to 1 in the Receiver Control register (low-order bit of 0xffff0000). The Ready bit is automatically reset to 0 when the MIPS program reads the Receiver Data using an 'lw' instruction.



```
.eqv KEY_CODE 0xFFFF0004 # ASCII code from keyboard, 1 byte
.eqv KEY_READY 0xFFFF0000 # =1 if has a new keycode ?
 # Auto clear after lw

.eqv DISPLAY_CODE 0xFFFF000C # ASCII code to show, 1 byte
.eqv DISPLAY_READY 0xFFFF0008 # =1 if the display has already to do
 # Auto clear after sw

.text
 li $k0, KEY_CODE
 li $k1, KEY_READY

 li $s0, DISPLAY_CODE
 li $s1, DISPLAY_READY

loop: nop

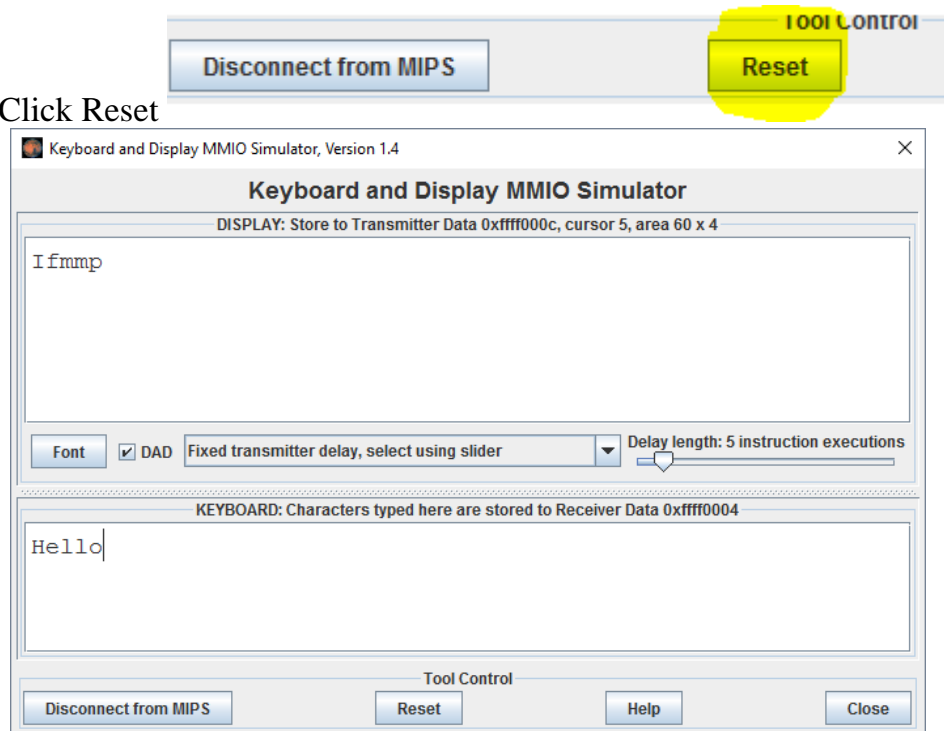
WaitForKey: lw $t1, 0($k1) # $t1 = [$k1] = KEY_READY
 nop
 beq $t1, $zero, WaitForKey # if $t1 == 0 then Polling
 nop
 #-----
ReadKey: lw $t0, 0($k0) # $t0 = [$k0] = KEY_CODE
 nop
 #-----
WaitForDis: lw $t2, 0($s1) # $t2 = [$s1] = DISPLAY_READY
 nop
```

```
 beq $t2, $zero, WaitForDis # if $t2 == 0 then Polling
 nop
 #-----
Encrypt: addi $t0, $t0, 1 # change input key
 #-----
ShowKey: sw $t0, 0($s0) # show key
 nop
 #-----
 j loop
 nop
```

### Warning: Must execute as below

1. Click Run 

2. Click Reset



### Assignment 1

Create a new project, type in, and build the program of Home Assignment 1.  
Show different values on LED . Write a program to make the led count from 0->9->0

### Assignment 2

Create a new project, type in, and build the program of Home Assignment 2.  
Draw something. Write a program to draw a red square, and moving the square from right to left

### Assignment 3

Create a new project, type in, and build the program of Home Assignment 3.  
Make the Bot run and draw a triangle by tracking and redraw the 1st triangle

### Assignment 4

Create a new project, type in, and build the program of Home Assignment 4.  
Read key char and terminate the application when receiving “exit” command.

# Laboratory Exercise 11

## Interrupts & IO programming

### Goals

After this laboratory exercise, you should understand the basic principles of interrupts and how interrupts can be used for programming. You should also know the difference between polling and using interrupts and the relative merits of these methods.

### Literature

■ Patterson and Hennessy: Chapter 2.7, 2.9, 2.10, 2.13, 5.7, Appendix A.6, A.7, A.10

### Polling or Interrupts

A computer can react to external events either by polling or by using interrupts. One method is simpler, while the other one is more systematic and also more efficient. You will study the similarities and differences of these methods using a simple “toy” example program.

Each peripheral device connects to the CPU via a few ports. CPU uses address to find out the respective port, and after that, CPU could read/write the new value to these port to get/control the device.

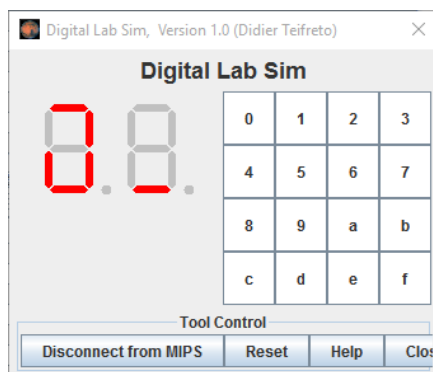
### Preparation

Study literature and these home assignments before coming into the class.

### Assignments at Home and at Lab

#### Home Assignment 1 - POOLING

Write a program using assembly language to detect key pressed in Digi Lab Sim and print the key number to console.



The program has a unlimited loop, to read the scan code of key button. This is POLLING.

In order to use the key matrix<sup>4</sup>, you should

1. assign experted row index into the byte at the address 0xFFFF0012

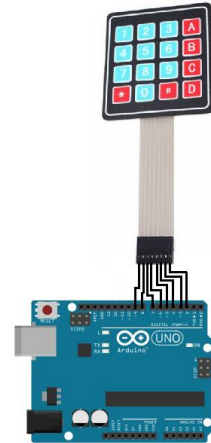
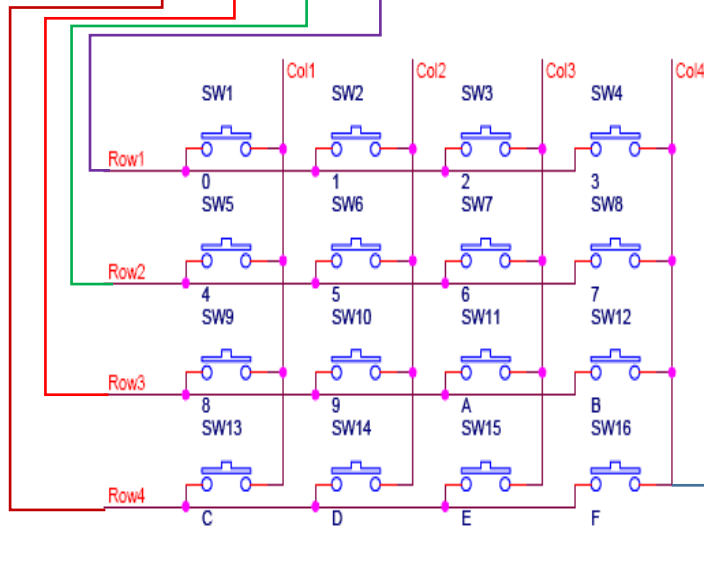
2. read byte at the address 0xFFFF0014, to detect which key button was pressed

<sup>4</sup> Key matrix animation: [http://hackyourmind.org/public/images/keypad12keys\\_anim.gif](http://hackyourmind.org/public/images/keypad12keys_anim.gif)

### IN\_ADRESS\_HEXA\_KEYBOARD

Address **0xFFFF0012**

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| 0 | 0 | 0 | 0 | R | R | R | R |



|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
| C | C | C | C | R | R | R | R |
| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |

### OUT\_ADRESS\_HEXA\_KEYBOARD

Address **0xFFFF0014**

```
#-----
col 0x1 col 0x2 col 0x4 col 0x8
#
row 0x1 0 1 2 3
0x11 0x21 0x41 0x81
#
row 0x2 4 5 6 7
0x12 0x22 0x42 0x82
#
row 0x4 8 9 a b
0x14 0x24 0x44 0x84
#
row 0x8 c d e f
0x18 0x28 0x48 0x88
#
#-----
command row number of hexadecimal keyboard (bit 0 to 3)
Eg. assign 0x1, to get key button 0,1,2,3
assign 0x2, to get key button 4,5,6,7
NOTE must reassign value for this address before reading,
eventhough you only want to scan 1 row
.eqv IN_ADRESS_HEXA_KEYBOARD 0xFFFF0012

receive row and column of the key pressed, 0 if not key pressed
Eg. equal 0x11, means that key button 0 pressed.
Eg. equal 0x28, means that key button D pressed.
.eqv OUT_ADRESS_HEXA_KEYBOARD 0xFFFF0014
```

```
.text
main: li $t1, IN_ADDRESS_HEX_A_KEYBOARD
 li $t2, OUT_ADDRESS_HEX_A_KEYBOARD
 li $t3, 0x08 # check row 4 with key C, D,
E, F

polling: sb $t3, 0($t1) # must reassign expected row
 lb $a0, 0($t2) # read scan code of key button
 print: li $v0, 34 # print integer (hexa)
 syscall
 sleep: li $a0, 100 # sleep 100ms
 li $v0, 32
 syscall
back_to_polling: j polling # continue polling
```

## Home Assignment 2 - INTERRUPT

Study the following assembly program, which **waits for an interrupt from keyboard matrix, and prints out a simple message**. Go over the code in detail and make sure that you understand everything, especially how to write and install an interrupt routine, how to enable an interrupt, and what happens when an interrupt is activated.

*Vietnamese support:*

*Cũng như các bộ xử lý khác, MIPS có 3 service với cùng một nguyên lý, nhưng khác nhau về mục đích sử dụng*

- *Exception: xảy ra khi có lỗi trong quá trình chạy, chẳng hạn tham chiếu bộ nhớ không hợp lệ.*
- *Trap: xảy ra bởi cách lệnh kiểm tra*
- *Interrupt: do các thiết bị bên ngoài kích hoạt*

*Cả 3 cơ chế trên đều được gọi chung là Exception.*

**Cách thức hoạt động:** *khi một exception xảy ra*

- *Khi có một Exception xảy ra, MIPS sẽ luôn nhảy tới địa chỉ cố định 0x80000180 để thực hiện chương trình con phục vụ ngắt. Để viết chương trình con phục vụ ngắt, sử dụng chỉ thị **.ktext** để viết code ở địa chỉ 0x80000180 nói trên.*
- *Bộ đồng xử lý C0, thanh ghi \$12 (status) sẽ bật bit 1*
- *Bộ đồng xử lý C0, thanh ghi \$13 (cause) sẽ thay đổi các bit 2~6 cho biết nguyên nhân gây ra ngắt*
- *Bộ đồng xử lý C0, thanh ghi \$14 (epc) sẽ chứa địa chỉ kế tiếp của chương trình chính, để quay trở về sau khi xử lý các đoạn mã Exception xong. (giống như thanh ghi \$ra)*
- *Trường hợp thanh ghi \$13 (cause) cho biết nguyên nhân làm tham chiếu địa chỉ bộ nhớ không hợp lệ, thanh ghi \$8 (vaddr) sẽ chứa địa chỉ lỗi đó.*
- *Nếu không có mã lệnh nào ở địa chỉ 0x80000180 (**.ktext**), chương trình sẽ hiện thông báo lỗi và tự kết thúc.*
- *Sau khi kết thúc chương trình con, sử dụng lệnh **eret** để quay trở lại chương trình chính. Lệnh **eret** sẽ gán nội dung thanh ghi PC bằng giá trị trong thanh ghi \$14 (epc).*

*Tuy nhiên, lưu ý rằng, trong MARS, thanh ghi PC vẫn chứa địa chỉ của lệnh mà ngắt xảy ra, tức là lệnh đã thực hiện xong, chứ không chứa địa chỉ của lệnh kế tiếp. Bởi vậy phải tự lập trình để tăng địa chỉ chứa trong thanh ghi epc bằng cách sử dụng 2 lệnh mfc0 (để đọc thanh ghi trong bộ đồng xử lý C0) và mtc0 (để ghi giá trị vào thanh ghi trong bộ đồng xử lý C0)*

```
mfc0 $at, $14 # $at <= Coproc0.$14 = Coproc0.epc
addi $at, $at, 4 # $at = $at + 4 (next instruction)
mtc0 $at, $14 # Coproc0.$14 = Coproc0.epc <= $at
```

- Các bit 8-15 của thanh ghi Cause, \$13 được sử dụng để xác định nguyên nhân gây ra ngắt. Hãy đọc thanh ghi này, kết hợp với thông tin chi tiết trong hướng dẫn sử dụng của từng thiết bị giả lập để biết được nguồn gốc gây ra ngắt.

**Cách thức viết chương trình phục vụ ngắt:** để viết chương trình con phục vụ ngắt khi có sự kiện ngắt xảy ra, có thể dùng một trong các phương pháp sau:

1. Viết chương trình con phục vụ ngắt trong cùng một file nguồn
2. Viết chương trình con phục vụ ngắt trong file nguồn độc lập, và lưu trữ trong cùng một thư mục với chương trình chính. Sau đó, sử dụng tính năng trong mục Setting là "Assemble all files in directory"
3. Viết chương trình con phục vụ ngắt trong file nguồn độc lập, và lưu trữ trong cùng một thư mục bất kì. Sau đó, sử dụng tính năng trong mục Setting là "Exception Handler.."

### **BUG: Ghi nhận các lỗi của công cụ MARS**

1. Giữa 2 lệnh syscall và lệnh jump, branch cần bổ sung thêm lệnh nop. Nếu không việc ghi nhận giá trị của thanh ghi PC vào EPC sẽ bị sai
2. Với các công cụ giả lập, nên bấm nút "Connect to MIPS" trước khi chạy giả lập. Nếu không, việc phát sinh sự kiện ngắt sẽ không xảy ra.

```
.eqv IN_ADDRESS_HEX keyboard 0xFFFF0012

.data
Message: .asciiz "Oh my god. Someone's pressed a button.\n"
#~~~~~
MAIN Procedure
#~~~~~
.text
main:
#-----
Enable interrupts you expect
#-----
Enable the interrupt of Keyboard matrix 4x4 of Digital Lab
Sim
li $t1, IN_ADDRESS_HEX keyboard
li $t3, 0x80 # bit 7 of = 1 to enable interrupt
```

```
 sb $t3, 0($t1)

 #-----
 # No-end loop, main program, to demo the effective of
interrupt
 #-----
Loop: nop
 nop
 nop
 nop
 b Loop # Wait for interrupt
end_main:

#~~~~~
GENERAL INTERRUPT SERVED ROUTINE for all interrupts
#~~~~~
.ktext 0x80000180
 #-----
 # Processing
 #-----
IntSR: addi $v0, $zero, 4 # show message
 la $a0, Message
 syscall
 #-----
 # Evaluate the return address of main routine
 # epc <= epc + 4
 #-----
next_pc: mfc0 $at, $14 # $at <= Coproc0.$14 = Coproc0.epc
 addi $at, $at, 4 # $at = $at + 4 (next instruction)
 mtc0 $at, $14 # Coproc0.$14 = Coproc0.epc <= $at
return: eret # Return from exception
```

### Home Assignment 3 – INTERRUPT & STACK

Study the following assembly program, in which

1. Main program enables 1 interrupt: from key matrix in Data Lab Sim
2. Main program only print a sequence number to console
3. Connect Data Lab Sim. Whenever user press a key button C, D, E, or F, an interrupt raises and print key scan-code to console

```
.eqv IN_ADDRESS_HEX4_KEYBOARD 0xFFFF0012
.eqv OUT_ADDRESS_HEX4_KEYBOARD 0xFFFF0014

.data
Message: .asciiz "Key scan code "

#~~~~~
MAIN Procedure
#~~~~~
.text
main:
 #-----
 # Enable interrupts you expect
 #-----
 # Enable the interrupt of Keyboard matrix 4x4 of Digital Lab
Sim
 li $t1, IN_ADDRESS_HEX4_KEYBOARD
 li $t3, 0x80 # bit 7 = 1 to enable
 sb $t3, 0($t1)
```

```
#-----
Loop an print sequence numbers
#-----
xor $s0, $s0, $s0 # count = $s0 = 0
Loop: addi $s0, $s0, 1 # count = count + 1
prn_seq: addi $v0, $zero, 1
 add $a0, $s0, $zero # print auto sequence number
 syscall
prn_eol: addi $v0, $zero, 11
 li $a0, '\n' # print endofline
 syscall
sleep: addi $v0, $zero, 32
 li $a0, 300 # sleep 300 ms
 syscall
 nop # WARNING: nop is mandatory here.
 b Loop # Loop
end_main:

#~~~~~
GENERAL INTERRUPT SERVED ROUTINE for all interrupts
#~~~~~
.ktext 0x80000180
#-----
SAVE the current REG FILE to stack
#-----
IntSR: addi $sp, $sp, 4 # Save $ra because we may change it later
 sw $ra, 0($sp)
 addi $sp, $sp, 4 # Save $ra because we may change it later
 sw $at, 0($sp)
 addi $sp, $sp, 4 # Save $ra because we may change it later
 sw $v0, 0($sp)
 addi $sp, $sp, 4 # Save $a0, because we may change it later
 sw $a0, 0($sp)
 addi $sp, $sp, 4 # Save $t1, because we may change it later
 sw $t1, 0($sp)
 addi $sp, $sp, 4 # Save $t3, because we may change it later
 sw $t3, 0($sp)
#-----
Processing
#-----
prn_msg: addi $v0, $zero, 4
 la $a0, Message
 syscall
get_cod: li $t1, IN_ADDRESS_HEX_A_KEYBOARD
 li $t3, 0x88 # check row 4 and re-enable bit 7
 sb $t3, 0($t1) # must reassign expected row
 li $t1, OUT_ADDRESS_HEX_A_KEYBOARD
 lb $a0, 0($t1)
prn_cod: li $v0, 34
 syscall
 li $v0, 11
 li $a0, '\n' # print endofline
 syscall
#-----
Evaluate the return address of main routine
epc <= epc + 4
#-----
next_pc: mfc0 $at, $14 # $at <= Coproc0.$14 = Coproc0.epc
 addi $at, $at, 4 # $at = $at + 4 (next instruction)
 mtc0 $at, $14 # Coproc0.$14 = Coproc0.epc <= $at
#-----
RESTORE the REG FILE from STACK
#-----
```



Thanh ghi số 13, status trong bộ đồng xử lý C0, chứa các thiết lập về tình trạng ngắt.

|                                                                  |    |    |    |    |    |    |  |                |   |   |   |   |     |    |
|------------------------------------------------------------------|----|----|----|----|----|----|--|----------------|---|---|---|---|-----|----|
|                                                                  |    |    |    |    |    |    |  | <i>ngắt.</i>   |   |   |   |   |     |    |
|                                                                  |    |    |    |    |    |    |  |                |   |   |   |   |     |    |
|                                                                  | 15 | 14 | 13 | 12 | 11 | 10 |  |                | 6 | 5 | 4 | 3 | 2   | 1  |
|                                                                  |    |    |    |    | KM | TC |  | Exception Code |   |   |   |   | K/U | IE |
| IE = 1 cho phép ngắt. IE = 0 vô hiệu hóa mọi hoạt động ngắt      |    |    |    |    |    |    |  |                |   |   |   |   |     |    |
| K/U=1 hoạt động ở chế độ Kernel. K/U=0 hoạt động ở chế độ User   |    |    |    |    |    |    |  |                |   |   |   |   |     |    |
| Ngoại lệ do syscall, overflow, lệnh tạo ngắt mềm như teq tegi... |    |    |    |    |    |    |  |                |   |   |   |   |     |    |
| Time Counter bộ đếm thời gian                                    |    |    |    |    |    |    |  |                |   |   |   |   |     |    |
| Key Matrix                                                       |    |    |    |    |    |    |  |                |   |   |   |   |     |    |
|                                                                  |    |    |    |    |    |    |  |                |   |   |   |   |     |    |

1. Main program enables 2 interrupts simultaneously: from key matrix and time counter in Data Lab Sim
2. Main program do nothing with a deadlock
3. Connect Data Lab Sim. Whenever **user press any key** or **time interval reaches**, an interrupt raises and print key scan-code to console.

```
.eqv IN_ADDRESS_HEX_KEYBOARD 0xFFFF0012
.eqv COUNTER 0xFFFF0013 # Time Counter

.eqv MASK_CAUSE_COUNTER 0x00000400 # Bit 10: Counter interrupt
.eqv MASK_CAUSE_KEYMATRIX 0x00000800 # Bit 11: Key matrix
interrupt

.data
msg_keypress: .asciiz "Someone has pressed a key!\n"
msg_counter: .asciiz "Time interval!\n"

#~~~~~
MAIN Procedure
#~~~~~
.text
main:
#-----
Enable interrupts you expect
#-----
Enable the interrupt of Keyboard matrix 4x4 of Digital Lab

Sim
li $t1, IN_ADDRESS_HEX_KEYBOARD
li $t3, 0x80 # bit 7 = 1 to enable
sb $t3, 0($t1)
```

```
Enable the interrupt of TimeCounter of Digital Lab Sim
li $t1, COUNTER
sb $t1, 0($t1)

#-----
Loop an print sequence numbers
#-----
Loop: nop
 nop
 nop
sleep: addi $v0,$zero,32 # BUG: must sleep to wait for Time
Counter
 li $a0,200 # sleep 300 ms
 syscall
 nop # WARNING: nop is mandatory here.
 b Loop
end_main:

#~~~~~
GENERAL INTERRUPT SERVED ROUTINE for all interrupts
#~~~~~
.ktext 0x80000180
IntSR: #-----
 # Temporary disable interrupt
 #-----
dis_int: li $t1, COUNTER # BUG: must disable with Time Counter
 sb $zero, 0($t1)
 # no need to disable keyboard matrix interrupt
 #-----
 # Processing
 #-----
get_caus: mfc0 $t1, $13 # $t1 = Coproc0.cause
IsCount: li $t2, MASK_CAUSE_COUNTER # if Cause value confirm
Counter..
 and $at, $t1,$t2
 beq $at,$t2, Counter_Intr
IsKeyMa: li $t2, MASK_CAUSE_KEYMATRIX # if Cause value confirm Key..
 and $at, $t1,$t2
 beq $at,$t2, Keymatrix_Intr
others: j end_process # other cases

Keymatrix_Intr: li $v0, 4 # Processing Key Matrix Interrupt
 la $a0, msg_keypress
 syscall
 j end_process
Counter_Intr: li $v0, 4 # Processing Counter Interrupt
 la $a0, msg_counter
 syscall
 j end_process
end_process:
 mtc0 $zero, $13 # Must clear cause reg
en_int: #-----
 # Re-enable interrupt
 #-----
 li $t1, COUNTER
 sb $t1, 0($t1)
 #-----
 # Evaluate the return address of main routine
 # epc <= epc + 4
 #-----
next_pc: mfc0 $at, $14 # $at <= Coproc0.$14 = Coproc0.epc
 addi $at, $at, 4 # $at = $at + 4 (next instruction)
 mtc0 $at, $14 # Coproc0.$14 = Coproc0.epc <= $at
```

```
return: eret # Return from exception
```

## Home Assignment 5 – KEYBOARD

Vietnamese support:

- Bộ xử lý MIPS cho phép tạo ra ngắt mềm, bằng lệnh *teq*, hoặc *teqi*
- Thiết bị Keyboard không tự động tạo ra ngắt khi có một phím được bấm, mà người lập trình phải tự tạo ngắt mềm.

```
.eqv KEY_CODE 0xFFFF0004 # ASCII code from keyboard, 1 byte
.eqv KEY_READY 0xFFFF0000 # =1 if has a new keycode ?
 # Auto clear after lw

.eqv DISPLAY_CODE 0xFFFF000C # ASCII code to show, 1 byte
.eqv DISPLAY_READY 0xFFFF0008 # =1 if the display has already to do
 # Auto clear after sw

.eqv MASK_CAUSE_KEYBOARD 0x00000034 # Keyboard Cause

.text

 li $k0, KEY_CODE
 li $k1, KEY_READY

 li $s0, DISPLAY_CODE
 li $s1, DISPLAY_READY

loop: nop
WaitForKey: lw $t1, 0($k1) # $t1 = [$k1] = KEY_READY
 beq $t1, $zero, WaitForKey # if $t1 == 0 then Polling
MakeIntR: teqi $t1, 1 # if $t0 = 1 then raise an Interrupt
 j loop

#-----
Interrupt subroutine
#-----
.ktext 0x80000180
get_caus: mfc0 $t1, $13 # $t1 = Coproc0.cause
IsCount: li $t2, MASK_CAUSE_KEYBOARD # if Cause value confirm
Keyboard.. and $at, $t1, $t2
 beq $at, $t2, Counter_Keyboard
 j end_process

Counter_Keyboard:
ReadKey: lw $t0, 0($k0) # $t0 = [$k0] = KEY_CODE

WaitForDis: lw $t2, 0($s1) # $t2 = [$s1] = DISPLAY_READY
 beq $t2, $zero, WaitForDis # if $t2 == 0 then Polling

Encrypt: addi $t0, $t0, 1 # change input key

ShowKey: sw $t0, 0($s0) # show key
 nop

end_process:
next_pc: mfc0 $at, $14 # $at <= Coproc0.$14 = Coproc0.epc
 addi $at, $at, 4 # $at = $at + 4 (next instruction)
 mtc0 $at, $14 # Coproc0.$14 = Coproc0.epc <= $at
return: eret # Return from exception
```

### **Assignment 1**

Create a new project, type in, and build the program of Home Assignment 1.  
Upgrade the source code so that it could detect all 16 key buttons, from 0 to F.

### **Assignment 2**

Create a new project, type in, and build the program of Home Assignment 2.  
Upgrade the source code so that it could detect all 16 key buttons, from 0 to F.

### **Assignment 3**

Create a new project, type in, and build the program of Home Assignment 3.  
Upgrade the source code so that it could detect all 16 key buttons, from 0 to F.

### **Assignment 4**

Create a new project, type in, and build the program of Home Assignment 4.  
Upgrade the source code so that it could detect all 16 key buttons, from 0 to F.

### **Assignment 5**

Create a new project, type in, and build the program of Home Assignment 5.

---

### **Conclusions**

Before you finish the laboratory exercise, think about the questions below:

- What is polling?
- What are interrupts?
- What are interrupt routines?
- What are the advantages of polling?
- What are the advantages of using interrupts?
- What are the differences between interrupts, exceptions and traps?

# Laboratory Exercise 12

## Cache Memory

### Goals

After this laboratory exercise, you should understand the basic principles of cache memories and how different parameters of a cache memory affect the efficiency of a computer system.

### Literature

■ Patterson and Hennessy: Chapter 7.1–7.3

### Preparation

Read the literature and this laboratory exercise in detail and solve the home assignments.

### Assignments at Home and at Lab

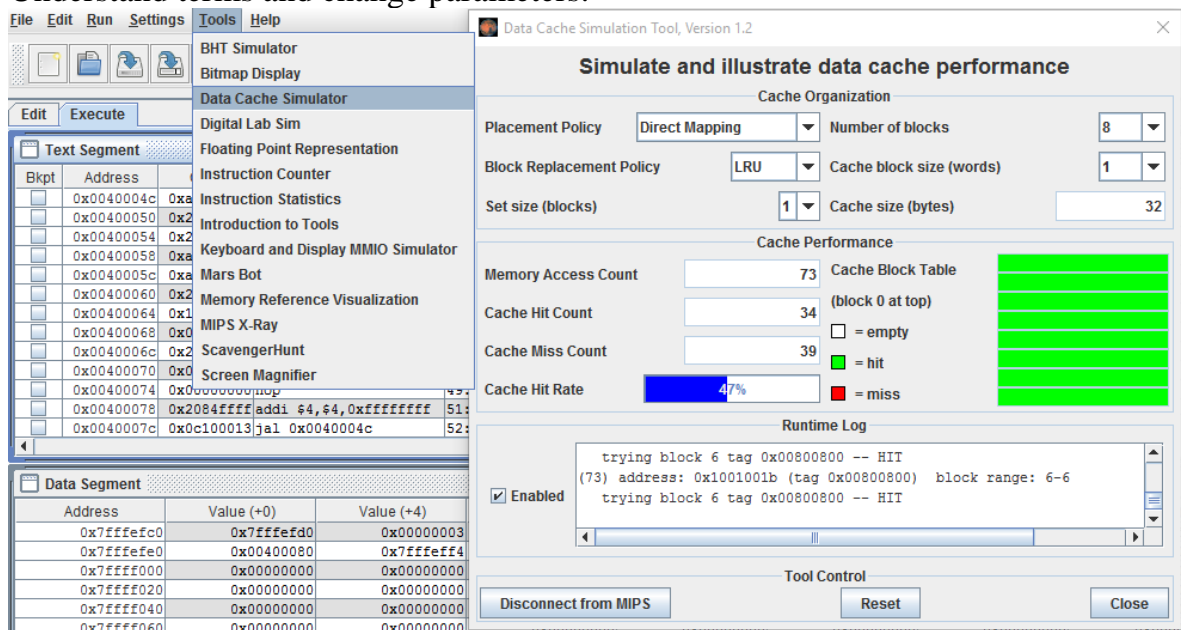
#### Home Assignment 1

Looking for and research information about Cache in CPU.

#### Home Assignment 2

Try to use Data Cache Simulation Tool in MARS.

Understand terms and change parameters.



### Function of a Cache Memory

A cache memory is a memory that is smaller but faster than the main memory. Due to the locality of memory references, the use of a cache memory can have the effect on the computer system that the apparent speed of the memory is that of the cache memory, while the size is that of the main memory.

## Assignment 1

Create a project, type in the program of Laboratory 7, Home Assignment 4 , build it and upload it to the cache simulator.

## Assignment 2

Run the program in the cache simulator and study how the instruction cache works. Then give *full* answers to the following questions.

- How is the full 32-bit address used in the cache memory?
- What happens when there is a cache miss?
- What happens when there is a cache hit?
- What is the block size?
- What is the function of the tag?

## Assignment 3

The parameters of the cache memory can be changed to test the effects of different cases. Investigate the effects of different parameter settings.

- Explain the following: cache size, block size, number of sets, write policy and replacement policy.
- If a cache is large enough that all the code within a loop fits in the cache, how many cache misses will there be during the execution of the loop? Is this good or bad?
- What should the code look like that would benefit the most from a large block size?

## Cache Efficiency

The actual efficiency gained by using a cache memory varies depending on cache size, block size and other cache parameters, but it also depends on the program and data.

---

## Conclusions

Before you pass the laboratory exercise, think about the questions below:

- What is the general idea with cache memory?
- What is a block?
- How does block size affect the efficiency of a cache?
- How fast is a cache memory? How fast is a DRAM?
- Do the optimal cache parameters depend on the program code?
- How can one select good cache parameters?
- Is it possible to change cache size on a PC? On a Mac?

## Track Change

| Ver | Date       | Details                                                 |
|-----|------------|---------------------------------------------------------|
| 5.3 | 2013/04/23 | More explanation for Lab Exercise 12, Home Assignment 2 |
| 6   | 2016/02    | Using MARS                                              |
| 6.2 | 2018/02    | Fix bug                                                 |
|     |            |                                                         |
|     |            |                                                         |

## Home Assignment 3

The following program sort a list of interger by ascending order. The swap procedure swap two successive elements. Read this example carefully, analyse each line of code, especially in sort procedure.

```
#Laboratory Exercise 5, Home Assignment 3
.data
test: .word 3,6,2,1
.text
main:
 la $a0,test
 li $a1,4
 jal sort
 nop
 la $a0,test
 lw $s0,0($a0) #load test array after sorting
 lw $s1,4($a0)
 lw $s2,8($a0)
 lw $s3,12($a0)
end_main:

swap:
 sll $t1,$a1,2 #reg $t1=k*4
 add $t1,$a0,$t1 #reg $t1=v+(k*4)
 #reg $t1 has the address of v[k]
 lw $t0,0($t1) #reg $t0 (temp)=v[k]
 lw $t2,4($t1) #reg $t2=v[k+1]
 #refers to next element of v
 sw $t2,0($t1) #v[k]=reg $t2
 sw $t0,4($t1) #v[k+1]=reg $t0 (temp)
 jr $ra
end_swap:

#Procedure sort
#function: sort a list of element ascending
sort:
 addi $sp,$sp,-20 #make room on stack for 5 registers
 sw $ra,16($sp) #save ra
 sw $s3,12($sp) #save $s3
 sw $s2,8($sp) #save $s2
 sw $s1,4($sp) #save $s1
 sw $s0,0($sp) #save $s0

 move $s2,$a0 #copy parameter $$a0 into $s2 (save $a0)
 move $s3,$a1 #copy parameter $a1 into $s3 (save a1)
 move $s0,$zero #i=0
for1tst:
 slt $t0,$s0,$s3 #reg $t0=0 if $s0>=s3 (i>=n)
 beq $t0,$zero,exit1 #go to exit1 if $s0>=s3 (i>=n)
 addi $s1,$s0,-1 #j=i-1
for2tst:
 slti $t0,$s1,0 #reg $t0=1 if $s1<0 (j<0)
 bne $t0,$zero,exit2 #go to exit2 if $s1<0 (j<0)
 sll $t1,$s1,2 #reg $t1=j*4
 add $t2,$s2,$t1 #reg $t2=v+(j*4)
 lw $t3,0($t2) #reg $t3=v[j]
 lw $t4,4($t2) #reg $t4=v[j+1]
 slt $t0,$t4,$t3 #reg $t0=0 if $t4>=t3
 beq $t0,$zero,exit2 #go to exit2 if $t4>=t3
 move $a0,$s2 #1st parameter of swap is v(old $a0)
 move $a1,$s1 #2nd parameter of swap is j
```



```
 jal swap
 addi $s1,$s1,-1 #j=j-1
 j for2tst #jump to test of inner loop
exit2:
 addi $s0,$s0,1 #i=i+1
 j for1tst #jump to test of outer loop
exit1:
 lw $s0,0($sp) #restore $s0
 lw $s1,4($sp) #restore $s1
 lw $s2,8($sp) #restore $s2
 lw $s3,12($sp) #restore $s3
 lw $ra,16($sp) #restore ra
 addi $sp,$sp,20 #restore stack pointer
 jr $ra
```

end\_sort:

## Assignment 4

Write a program that let user input a string (press Enter to terminate), sort this string by ascending order (based on ASCII code) and print the string after sorting.