Design and Implementation of a Distributed Database Management System Using Modern Technologies and Tools

Kembabazi Barbara Gamukama Yihong 2020280607

Yoke Kai Wen 2020280598 (Leader)

December 29, 2020

Abstract

This report summarizes the relevant knowledge covered in our DDBS module and presents the proposed design and implementation of a distributed database system built using MongoDB deployed on Docker containers. Through this project, we understand the reasons that pushed the evolution of NoSQL database platforms from SQL database platforms that had hitherto dominated the market. We specifically focus on some of the existing distributed database solutions, particularly highlighting their main features on enabling the integration of data among database management systems distributed at different sites. After reviewing existing solutions, we elucidate the problem definition which involves implementing a simplified distributed database system for a Douban-like social networking site. By applying key learning points from existing solutions to the given problem, we then decided to use MongoDB deployed on Docker containers to simulate a distributed database environment. We describe and explain in detail our command-line based implementation of our proposed solution, and finally conclude with future work that could be done. More details regarding the code implementation can be found in our github repostiory https://github. com/kwyoke/DDBS-2020, and also in the Appendix.

Keywords: distributed database systems, database management systems, integration, centralized

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1 PROBLEM BACKGROUND & MOTIVATION

The increasing usage of web and mobile applications today has led to the overwhelming need for secure and stable database systems. Database Management Systems (DBMSs) are an ubiquitous and critical component of modern computing, and the result of decades of research and development in both academia and industry.[2] Database systems are a common setup for all types of applications. Today, database system tasks have evolved from simple data processing and data maintenance in an application into centralized administration systems of data in multiple applications. Centralization of database management systems has enabled data independence with their main motivation as the desire to integrate the operational data and provide controlled access to that data within the enterprise. However, these systems, being either the mainframe or server of an enterprise's application, must satisfy all requests received by the system, otherwise they would become a bottleneck in case of any malfunction or failures. It is important to note that the success of centralized database systems lies in the ability to integrate the data and not the centralization.

1.1 Distributed database

Integration can be achieved through the implementation of both computer network technology and database system management systems during the development of the enterprise's system i.e. a distributed database system. A distributed database is a collection of multiple, logically interrelated databases distributed over a computer network; A distributed database management system is a software system that permits the management of the distributed database and makes the distribution transparent to the users[4]. Distributed database systems are more reliable and responsive as compared to the traditional database system owing to the fundamental reason of the distributed processing which copes efficiently with large-scale data management problems using the divide-and-conquer rule.

1.2 Centralized vs distributed database

The main difference between the centralized database and distributed database is the fact that a centralized database stores data at a single location only whereas a distributed database system consists of multiple databases which are connected with each other and are spread across different physical locations. Therefore, when system failure occurs at centralized system, the entire data will be destroyed, while it will not be the case for distributed database systems. Thus, our drive to build a system that achieves integration of data without the need for centralization, thereby eliminating the bottleneck faced by the traditional database system.

2 EXISTING SOLUTIONS

The serious problem with centralized databases / data centers i.e. system failures happening due to power outages, cooling failures, network failures, or natural disasters, results in the need for distributing and integrating data while at the same time maintaining its high performance, durability and /transparency among users. The study of distributed databases has been going on for some years now and has led to the rise of new database related technologies such as use of non-relational databases and NoSQL, unlike traditional databases that used only SQL and existed as relational databases.

2.1 Apache Cassandra

Development of Apache Cassandra was born from these studies. It was one of the existing solutions to distributed databases that we studied and used as basis for comparison during the development of our model. Apache Cassandra is a database platform that resembles a normal database and shares many design and implementation strategies. However, it does not support a full relational data model, and instead provides clients with a simple data model that supports dynamic control over data layout and format. In other words, it brings out the integration of data without the need to centralize the databases. Cassandra is a distributed storage system for managing very large amounts of structured data spread out across many commodity servers, while providing highly available service with no single point of failure.[3] Cassandra was developed at Facebook for the initial purpose of solving the problem of inbox search failure, thus it was constructed in a manner to treat failures as the norm rather than the exception. The core distributed systems techniques used in Cassandra are partitioning, replication,

membership, failure handling and scaling of the data and it runs on three main commands "insert", "get", "delete". [3] It partitions its data over a set of nodes and makes use of replicas, placing each into N hosts of each data item whereby N is a replication factor configured as the per-instance. There is a coordinator node in charge of the replication of the data items that fall within its range and replicates them into N-1 nodes in its ring. Since all nodes are aware of every other node in the ring and know their ranges, and amongst them one is selected as a leader using the Zookeeper system, therefore when one crashes, it knows what range it was responsible for immediately when it comes back up. This scheme of replicating across multiple data centers allows Cassandra to handle entire data center failures without any outage.

2.2 Google Spanner

A study by Google on distributed database systems was the first system to distribute data at global scale and support externally-consistent distributed transactions[1]. This system is called Spanner, a globally-distributed database deployed at Google. Spanner also makes use of the replication factor as seen in the other studies. The replication configurations for data are said to be dynamically controlled at a fine grain by applications, and applications are able to set constraints on which data centers contain which data and the distance between users and data. The two unique features that Spanner executes are the ability to provide externally-consistent reads and writes, while at the same time provide globally consistent reads across all databases at a given timestamp. The ability to control concurrency is an important factor when dealing with distributed databases. The use of TrueTime API enables the above features to function since Spanner is able to keep globally-meaningful commit timestamps to transactions, when the transactions may be distributed; it is important to note that the TrueTime API is implemented by a set of time master machines per datacenter and a timeslave daemon per machine. The entire Spanner architecture mimics and implements the master-slave technique through its zones. A single deployment of a Spanner is called a universe which is made up of a set of zones. Each zone represents a unit of physical isolation; each having its own zone master and one hundred to thousand spanservers which are assigned data from the master and serve to the client. It deploys a universe master as the primary console that displays status information about all the zones for interactive debugging and a placement driver that handles automated movement of data across zones on the timescale of minutes. Replication and distributed transactions of data occurs within the spanservers through the implementation of a single Paxos state machine on top of each tablet(data structure), each set of replicas has a leader that also follows the master-slave architecture to deal with concurrency control during distributed transactions. [1]

2.3 Google F1

F1 is a distributed relational database system built at Google to support the Ad-Words business that was built on top of Spanner.[6] Both studies employ the master-slave architecture. However, F1 has a slave pool in each of its clusters and each cluster contains a spanner which can communicate with F1 servers outside its cluster. The shared slave pool consists of F1 processes that exist only to execute parts of distributed query plans on behalf of regular F1 servers. The entire system is based off the F1 servers / slaves and Spanner servers. The previous study on Spanner showed its utilization of replication and this is maintained as well in this study.

2.4 Key learning points form existing solutions

All existing solutions clarify the need for replication and partitioning of the databases / data items as a fault tolerance tactic. This is a task we aim to accomplish during the design and development of our system: to avoid system failure through the use of replicas.

3 PROBLEM DEFINITION

We are given 10GB of structured and unstructured data to be fragmented and allocated to various sites, which can then be efficiently modified and queried by users, with the entire distributed system being monitored by administrators. Essentially, the problem simulates the database component of a social networking application focused on article reading like Goodreads and Douban.

3.1 Data description

The structured data includes user.dat, article.dat, read.dat that contain tables recording structured information (see Figure 1), and the unstructured data comprises text, image and video data associated with each article.

```
    USER
        {"timestamp": "1506328859000", "id": "u0", "uid": "0", "name": "user0", "gender": "female", "email": "email0", "phone": "phone0", "dept": "dept2", "grade": "grade4", "language": "zh", "region": "Beijing", "role": "role0", "preferTags": "tags5", "obtainedCredits": "69"}

    ARTICLE
        {"id": "a0", "timestamp": "150600000000000", "aid": "0", "title": "title0", "category": "science", "abstract": "abstract of article 0", "articleTags": "tags8", "authors": "author1552", "language": "en", "text": "text_a0.txt", "image": "image_a0_0.jpg,image_a0_1.jpg,", "video": ""]

    READ
        {"timestamp": "1506332297000", "id": "r0", "uid": "15", "aid": "37", "readOrNot": "1", "readTimeLength": "4", "readSequence": "1", "agreeOrNot": "0", "commentOrNot": "0", "shareOrNot": "0", "commentDetail": "comments to this article: (15,37)"}
```

FIGURE 1: Example records in user, article, and read tables in JSON format

3.2 Populating Be-Read and Popular-Rank tables

Given the three tables (User, Article, Read) to be bulk-loaded into the data centre, we are to perform query, insert and join operations on them to populate two more tables: Be-Read and Popular-Rank.

- 1. **Be-Read** Displays user interactions for each article. Attributes are *id*, *times-tamp*, *aid*, *readNum*, *readUidList*, *commentNum*, *commentUidList*, *agreeNum*, *agreeUidList*, *shareNum*, *shareUidList*.
- 2. **Popular-Rank** Ranks the top five most popular articles for daily, weekly and monthly time intervals. Attributes are *id*, *timestamp*, *temporalGranularity*, *articleAidList*.

The relationships between all five tables are demonstrated in Figure 2.

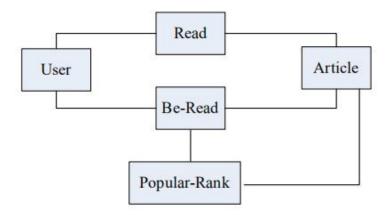


FIGURE 2: Relationships between five tables

3.3 Fragmentation and allocation requirements

We are to demonstrate fragmentation and allocation of structured data in a distributed environment using two DBMS sites, DBMS1 and DBMS2, according to the horizontal fragmentation and allocation scheme shown in Table 1. Unstructured data (text files, images, videos) are to be stored in Hadoop HDFS.

| | Attribute | DBMS1 | DBMS2 | |
|--------------|-----------|---------|------------|--|
| User | region | Beijing | Hong Kong | |
| Article | category | science | science | |
| Afficie | category | Science | technology | |
| Read | region | Beijing | Hong Kong | |
| Be-Read | antagory | science | science | |
| De-Reau | category | Science | technology | |
| Danulan Dank | antagami | science | science | |
| Popular-Rank | category | science | technology | |

Table 1: Table showing fragmentation and allocation scheme of five tables

3.4 All requirements summarised

- 1. Bulk data loading with data partitioning and replica consideration.
- 2. Efficient execution of data insert, update and queries, including retrieval of unstructured multimedia data.
- 3. Monitoring running status of DBMS servers (including amount and location of data, workload etc).

4 PROPOSED SOLUTION

To fulfil the requirements set out in the previous section, we used Docker containers to simulate the different DBMS sites, and MongoDB, a NoSQL document database as our database system. To query and modify the database, the Mongo shell (command line) is used. All scripts and instructions required for implementing our solution can be found in our Github repository https://github.com/kwyoke/DDBS-2020.

4.1 Tools used in this project

In this section, we describe and justify our choice of software tools for this project.

4.1.1 Docker



Containerisation has become a major trend in software development, with Docker being the most representative provider of container technology. The Docker platform allows different applications to run in separated isolated environments called 'containers' on the same host machine. Furthermore, the containers are lightweight (compared to virtual machines) since they run directly on the host machine's kernel. This makes Docker containers a perfect choice for deploying our distributed database environment to simulate different servers, given our limited computing resources.

4.1.2 MongoDB



MongoDB is a non-relational NoSQL document database that provides support for JSON-like storage, and is amenable to both structured and unstructured data which are stored in unstructured collections instead of relational tables, and hence suitable for our use case. Furthermore, MongoDB comes with a whole suite of functionalities such as sharding (automatic distribution of data across different servers), replica sets (allowing hot/cold standby DBMSs), PyMongo (python library interface), GridFS (specification for storing and retrieving large files) and monitoring capabilities that greatly enhances the efficiency of our project implementation. Below, we describe in greater detail the relevant functionalities of MongoDB.

• Sharding - MongoDB allows automatic distribution of data to different servers (shard clusters) based on how the shard clusters are configured, in a process called 'sharding'. This automatic routing is carried out by the mongos router server that tracks what data is on which shard by caching the metadata from the config servers, and then using the metadata to route

operations from applications and clients to the mongod instances (servers configured as shards).[5]

- **Replica sets** MongoDB allows convenient set up of replica sets that are groups of mongod processes (servers) that maintain the same dataset. Replica sets provide redundancy and high availability. In the event of system failure in the primary server, the secondary server in the same replica set will become the primary server, hence making the database system tolerant to fault. [5]
- **PyMongo** The PyMongo library allows us to work with MongoDB from Python, as we could not do everything within the Mongo shell and needed to write Python scripts to execute some of the required functions.
- **GridFS** We were expected to use Hadoop HDFS to store the unstructured multimedia data, but we found MongoDB's GridFS to be a better substitute for Hadoop HDFS as it was much easier to use and interface with our MongoDB database. GridFS is a MongoDB specification for storing and retrieving large files such as images, audio files, video files within MongoDB collections, and hence suitable for storing the large image and video files generated from the 10GB data.

4.2 Distributed system architecture

We set up 8 docker containers in total (see Figure 3): 3 config servers in a replica set; 1 mongos router server; 2 DBMS servers shard, each being a single server replica set, for distributing structured data; 2 GridFS server shards, each being a single server replica set, for distributing multimedia data. This can be achieved by running our docker-compose files in our Github repository https://github.com/kwyoke/DDBS-2020, which sets up each Docker container as a mongod or mongos instance based on their respective purpose and configure them as shards. The output on the command line is shown in Figure 4. For greater availability and fault tolerance, we should have replica sets of at least three servers for our DBMS and GridFS shards, but due to lack of computer space, we decided to stick with only single server shards, and demonstrate the replica set concept via the config servers.

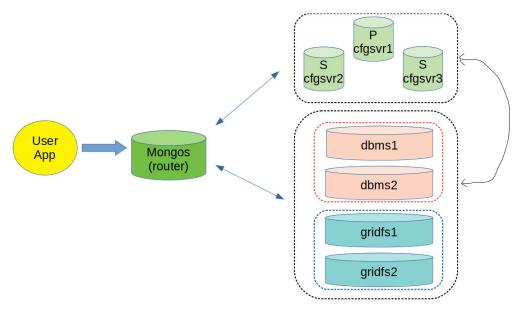


FIGURE 3: Architecture of proposed solution: distributed database system with MongoDB deployed on Docker containers

| ONTAINER ID | IMAGE | COMMAND | CREATED | STATUS | PORTS | NAMES |
|-------------|-------|-----------------------|--------------|---------------|--------------------------|----------|
| ab78d0809b5 | mongo | "docker-entrypoint.s" | 21 hours ago | Up 33 minutes | 0.0.0.0:50004->27017/tcp | grid2svr |
| 8a184abfeaa | mongo | "docker-entrypoint.s" | 21 hours ago | Up 33 minutes | 0.0.0.0:50003->27017/tcp | grid1svr |
| 245b1f89f36 | mongo | "docker-entrypoint.s" | 21 hours ago | Up 33 minutes | 0.0.0.0:50002->27017/tcp | dbms2svr |
| 67aa55c3b39 | mongo | "docker-entrypoint.s" | 21 hours ago | Up 33 minutes | 0.0.0.0:60000->27017/tcp | mongos |
| a9eae4c3e73 | mongo | "docker-entrypoint.s" | 21 hours ago | Up 32 minutes | 0.0.0.0:50001->27017/tcp | dbms1svr |
| a28b0909487 | mongo | "docker-entrypoint.s" | 21 hours ago | Up 32 minutes | 0.0.0.0:40003->27017/tcp | cfgsvr3 |
| 345df0fc8f2 | mongo | "docker-entrypoint.s" | 21 hours ago | Up 32 minutes | 0.0.0.0:40001->27017/tcp | cfgsvr1 |
| bca6251ae72 | mongo | "docker-entrypoint.s" | 21 hours ago | Up 32 minutes | 0.0.0.0:40002->27017/tcp | cfgsvr2 |

FIGURE 4: Output on command line after setting up docker containers

Figure 3 also shows the relationships between each server. The config servers store metadata on the volume and location of data in the DBMS and GridFS shards, which is accessed by the mongos router for routing operations from applications to the DBMS and GridFS shards.

4.3 Bulk loading

The user.dat, article.dat, read.dat data files are copied into the mongos docker container, and then imported into the mongos server's ddbs database using the mongoimport command.

4.4 Sharding

Sharding of collections is quite simple in MongoDB. There are two types of sharding: ranged and hashed. Ranged sharding allows us to assign documents of specific shard key values to specific target shards, while hashed sharding assigns documents more evenly (at random) based on hashed index. For this project, we mainly used ranged sharding because we have to direct documents based on a specific fragmentation rule (see Table 1). For the unstructured multimedia data, we employed hashed sharding as we are not concerned over which GridFS shard (grid1, grid2) the multimedia data chunks go to. Once we set the shard keys of each collection and configure the shard zones of each shard, the balancer running in the background will automatically distribute the database documents to the respective shard. These processes are completely transparent to the user who can simply interact with the mongos server as if it was a centralised database.

4.4.1 Summary of main steps for enabling and configuring sharding

- 1. Enable sharding for the target database.
- 2. Create index for collection to be sharded, and this index will be used as shard key. Shard key can be made up of multiple fields, and is critical for even distribution of data across shards.
- 3. Enable sharding for the collection, and set the collection index as shard key.
- 4. Configure the sharding zones (zone tags and zone range) for each shard, specific to the collection. This determines where each document in the collection will be assigned to.
- 5. Ensure that the balancer is running for that collection.

4.4.2 Limitations of sharding in MongoDB

While sharding in MongoDB is very easy to implement, it does not fulfil some of our requirements.

1. **Assigning same data to both shards** - According to the allocation scheme in Table 1, we are supposed to assign science articles to both DBMS1 and DBMS2. However, in MongoDB, sharding does not allow two shards to

contain the same data, as all replica considerations are meant to be accounted for by the replica sets. To work around this, we copied article to another collection articlesci for just science articles, and assigned it to DBMS2. The same problem arose for beread and popRankSci, which were tackled in the same manner. See Table 2 for more details.

2. Collection must contain fields from shard key - Another issue is that sharding of a collection requires the collection to contain fields from the shard key. For instance, since we are required to shard read according to the "region" field in the user collection, we must also include the "region" field in the read table so that it can be sharded according to "region". Unfortunately, this adds extra overhead as we now have to store one more extra field. The same problem also arose for beread, bereadsci which requires us to include an extra "category" field for these two collections.

Taking into consideration sharding and its limitations in MongoDB, while trying to stick to the problem specification as closely as possible, we finally fragmented and allocated our collections in the scheme shown in Table 2.

| Collection | Shard key | DBMS1 zone tag | DBMS2 zone tag |
|-------------|-----------------|----------------|----------------|
| user | {region, uid} | BJ | НК |
| article | {category, aid} | SCI | TECH |
| articlesci | {category, aid} | - | SCI2 |
| read | {region, id} | BJ | НК |
| beread | {category, aid} | SCI | TECH |
| bereadsci | {category, aid} | - | SCI2 |
| popRank | {_id} | - | POPALL |
| popRankSci | {_id} | POPSCI | - |
| popRankSci2 | {_id} | | POPSCI2 |
| popRankTech | {_id} | - | TECH |

Table 2: Table showing sharding configuration for each collection

4.5 MongoDB aggregation pipeline

Aggregation processes multiple data documents and return computed records. MongoDB provides an aggregation pipeline that we use to populate the beread and popRank collections.

4.5.1 Populating beread

beread is formed from aggregating read - the read documents are grouped by the "aid" field, and then the number of reads, agrees, shares and comments are counted for each "aid" group, together with the relevant list of "uid"s involved in these interactions, and then finally output as the beread collection. The exact mongo shell command for the aggregation pipeline is shown in Figure 5. Note that before the aggregation, we already added the article category (and article timestamp) field to read to facilitate subsequent sharding for beread.

FIGURE 5: Aggregation pipeline to populate beread collection

Figure 6 shows a sample document from the populated beread collection.

```
mongos> db.beread.find().limit(1)

{    "_id" : "8145", "category" : "technology", "timestamp" : "1506000008145", "rea dNum" : 97, "readUidList" : [ "4621", "2699", "874", "2369", "1432", "2852", "74 77", "5699", "2483", "7015", "4448", "1146", "8545", "2621", "6881", "8785", "82 27", "5632", "982", "8522", "2169", "8940", "5604", "2135", "1508", "6786", "554 5", "1018", "5507", "4894", "3181", "9031", "1705", "1347", "6153", "33031", "222 1", "60003", "8594", "5413", "6000", "8760", "5841", "6469", "2585", "3567", "2924 1", "60003", "8594", "5413", "600", "8760", "5841", "6729", "8471", "7934", "348", "4142", "8070", "2055", "8997", "3154", "1699", "8762", "3187", "3353", "2 310", "1256", "8961", "6526", "9262", "9052", "1045", "3879", "6315", "2964", "3 710", "9023", "2326", "90", "8390", "7930", "1546", "4704", "3729", "8097", "243 19", "6536", "6763", "2828", "3908", "5478", "8722", "4437", "3900"], "commentNu 17, "commentUidList" : [ "2517", "13", "9023", "8762", "3031", "3154", "701 5", "4894", "1705", "5478", "5413", "982", "2964", "2924", "3710", "2135", "4437", "3908", "6469", "3900", "4024", "50", "9023", "4009", "4621", "2585", "6786", "1018", "1346", "3187", "6526", "1256", "2326", "8522", "3879", "1508", "5604", "5609", "113", "8471", "7015", "4142", "8070", "1146"], "shareNum" : 18, "shareUidList" : [ "4437", "5545", "1045", "2483", "8070", "1146"], "shareNum" : 18, "shareUidList" : [ "4437", "5545", "1045", "2483", "8070", "1146"], "shareNum" : 18, "shareUidList" : [ "4437", "5545", "5545", "1045", "2483", "8070", "1354", "1546", "348", "600", "5632", "3226", "6536", "6536", "2221", "6000", "3154", "1546", "348", "600", "5632", "2326", "6536", "5545", "1045", "5483", "8070", "3154", "1546", "348", "600", "5632", "3226", "6536", "5545", "1045", "5483", "8070", "3154", "1546", "348", "600", "5632", "3226", "6536", "5545", "1045", "5483", "8070", "3154", "1546", "348", "600", "5632", "3226", "6536", "6536", "2221", "3187", "5545", "3879", "2964", "8722"], "aid" : "a8145", "2326", "6536", "2221", "31
```

FIGURE 6: Sample document from beread collection

4.5.2 Populating popRank

popRank, popRankSci, popRankTech collections are formed from multiple aggregations and concatenations from the read collection. We have three collections for Popular-Rank to represent the most popular articles for both categories, for science category only and for technology articles only. For each popRank collection, we perform three aggregations, one for monthly, one for weekly and one for daily intervals. Each aggregation outputs a temporary collection, e.g. popRankMth, popRankWk, popRankDay which are then concatenated together to form popRank. In each aggregation pipeline, we extract the year, month, week, day information from the timestamp, group the documents by the respective interval (month, week or day), and then sum up all the interactions (read, agree, share, comment) for each article in each group to form a popularity score. Within each group, we then sort and rank the articles by their popularity score, and select the top five articles and push them into an array of field "articleAidList", finally outputting a temporary collection for that particular time interval (e.g. monthly). See Figure 7 for the full mongo shell command to aggregate and form popRankSciMth.

FIGURE 7: Aggregation pipeline to populate popRankSciMth (temporary collection)

Figure 8 shows a sample document from the populated collection popRank.

FIGURE 8: Sample document from popRank

4.6 MongoDB CRUD operations

CRUD operations like data insert, update and queries can be performed efficiently on our distributed database system, just like a centralised database due to MongoDB's sharding functionality. However, the issue arises when we want to reflect any new changes on for instance read to beread and popRank which are derived from the read collection. This is achieved using MongoDB's db.collection.watch() function which allows us to watch for changes in the read collection and subsequently update the beread and popRank collections accordingly. We implement this using a python script running in the background.

4.6.1 Python script to automatically refresh collections upon CRUD changes

In our Github Repository https://github.com/kwyoke/DDBS-2020, we include python scripts to watch and automatically refresh collections auto_refresh_onread.py, auto_refresh_onarticle.py. The python scripts leverage the PyMongo library to interface with MongoDB, watch out for CRUD changes in the target collections, and update derived collections accordingly. A code snippet demonstrating PyMongo and db.collection.watch() is shown in Figure 9 and the terminal output when the python script is running is shown in Figure 10. The required collections to watch and refresh are summarised in Figure 11.

FIGURE 9: Code fragment: leveraging PyMongo and db.collection.watch() to update collections

```
(ddbs) buyokegyku-device:-/Documents/THDMC/DDBs/proj_mongoblysharding$ python auto_refresh_onarticle.py
db_articlesci_updated {"_id': objectId('SfdbB327474dfBe7f9fffff'), 'id': 'a654321', 'timestamp': 'i500000000000', 'aid': '654321', 'title': '', 'category': 'science', 'abstract': '', 'articleTags': '', '
authors': '', 'language': 'zh', 'text': '', 'image': '', 'video': '')

U
```

FIGURE 10: PyMongo running in terminal and updating articlesci collection when changes to article are detected

| Collection to watch | Collection to update | Update to perform |
|---------------------|--|--|
| article | articlesci | insert document if category = "science" |
| read | beread, bereadsci | update fields: {readNum, readUidList, commentNum, commentUidList, agreeNum, agreeUidList, shareNum, shareUidList} for relevant documents with aid matching documents inserted into db.read |
| read | popRank, popRankSci, popRankTech | filter records in db.read with timestamps in the same month as new read, calculate popScoreAgg for each aid within the filtered records for each month/week/day, put top 5 for each month/week/day in arrays, form new documents and insert, replacing if timestamp overlaps |

FIGURE 11: List of refreshes to perform upon change in article, read collections

4.7 Storing of unstructured data in GridFS

The articles come with multimedia data that are not stored directly on our MongoDB DBMS servers due to their large size. Instead, our MongDB DBMS servers

only store the filenames of the multimedia data, which are then used to retrieve the image data from other storage servers (grid1, grid2) dedicated for multimedia data, which use the GridFS specification. We found GridFS to be a better substitute than Hadoop HDFS as they perfor the same function but it is much easier to interface GridFS with our MongoDB database system.

4.7.1 GridFS and sharding

We have two shard servers, grid1 and grid2 dedicated to storing multimedia data of articles. Before loading the multimedia to the data, we first set up sharding so that the (10GB) worth of data will be distributed to the two shards as it is uploaded to GridFS. GridFS is a specification for storing and retrieving large files of more than 16MB which is perfect for our use case in storing image, video and text files data. We apply hashed sharding here to allow more even distribution of multimedia data across the two shard servers.

GridFS automatically stores the data files in db.fs.chunks collection, and the filenames in db.fs.files collection. We only shard the db.fs.chunks collection since it is very large, while db.fs.files is not be sharded and remain on the mongos server.

4.7.2 Loading and retrieving files from GridFS

To load and retrieve files to/from GridFS, we use the mongofiles utility which allows us to manipulate GridFS objects from the command line.

4.8 Monitoring

MongoDB provides monitoring tools that allows administrators to easily monitor the running status of the DBMS servers, including its managed data (amount and location) as well as workload.

4.8.1 Monitoring managed data

Simply type sh.status() and db.collection.getShardDistribution() in the mongo shell of the mongos router container to observe the status of the various connected shard clusters (see Figure 12) and the distribution of documents in that collection (see Figure 13) respectively.

FIGURE 12: Output of sh.status()

```
Mongos> db.user.getShardDistribution()

Shard dbms1rs at dbms1rs/192.168.1.152:50001
data: 1.67MlB docs: 6043 chunks: 2
estimated data per chunk: 859KlB
estimated docs per chunk: 3021

Shard grid2rs at grid2rs/192.168.1.152:50004
data: 08 docs: 0 chunks: 1
estimated data per chunk: 08
estimated docs per chunk: 0

Shard dbms2rs at dbms2rs/192.168.1.152:50002
data: 1.1MlB docs: 3957 chunks: 1
estimated data per chunk: 1.1MlB
estimated data per chunk: 3957

Shard grid1rs at grid1rs/192.168.1.152:50003
data: 08 docs: 0 chunks: 1
estimated data per chunk: 08
estimated data per chunk: 08
estimated docs per chunk: 08
Shard dbms1rs contains 08 data, 08 docs in cluster, avg obj size on shard: 291B
Shard grid2rs contains 08 data, 08 docs in cluster, avg obj size on shard: 85
Shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 85
Shard grid2rs contains 08 data, 08 docs in cluster, avg obj size on shard: 85
Shard grid2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
Shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
Shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
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shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
shard dbms2rs contains 39.73% data; 39.57% docs in cluster, avg obj size on shard: 88
```

FIGURE 13: Output of db.user.getShardDistribution()

4.8.2 Monitoring workload

MongoDB also facilitates the monitoring of operation execution times, memory usage, CPU usage and operation counts. By using the mongostat utility in the command line, we can observe in real-time the number of operations of each server (see Figure 14).

FIGURE 14: Output of mongostat

Unfortunately, the Docker version of MongoDB does not allow free monitoring at the moment, but could be provided in the near future. In any case, we demonstrate that MongoDB does allow for monitoring of memory usage, CPU utilisation etc using a browser interface as shown in Figure 15.

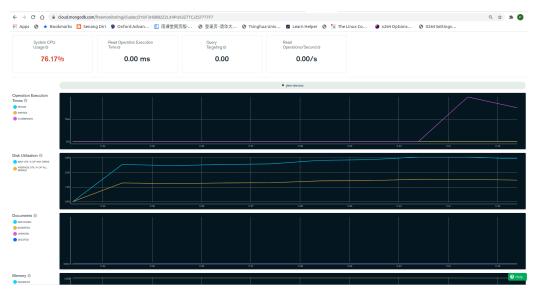


FIGURE 15: Browser demonstrating free monitoring

5 SOLUTION EVALUATION

In this section, we show that our system is able to fulfil the three main requirements set in the problem specification (section 3.4) of this project.

- 1. Bulk loading with data partitioning and replica consideration. We discussed this in-depth in Sections 4.3 and 4.4, and this can be verified by running Makefile commands make bulkload and make popNshard as explained in the Manual (appendix).
- 2. **Efficient execution of CRUD operations.** MongoDB easily facilitates CRUD operations on sharded collections as if they were all in one place. Furthermore, we also showed in section 4.6 that when one collection changes, another related collection will be updated via a python script. Moreover, in section 4.7, we show how multimedia data associated with an article can be retrieved.
- 3. **Monitoring running status of DBMS servers.** We demonstrated in section 4.8 how the monitoring utilities in MongoDB can be used to monitor the sharded data and the workload on each server.

6 CONCLUSION

6.1 Future Work

Although our system successfully implements the three main objectives, there are some functions advantageous to a distributed system that can be added. These include:

1. Hot Standby DBMS for fault tolerance.

MongoDB facilitates the setting up of replica sets which increase robustness against faults. At the moment, we only have a replica set for the config servers, but not for the data shards (DBMS1, DBMS2, grid1, grid2) due to lack of computer space, but this could be easily enabled by setting up extra Docker containers to simulate secondary servers to DBMS1, DBMS2. In case the primary server in one of the DBMS shard clusters fails, the system will automatically switch to the secondary server (which contains the same exact data as the primary server), thus allowing all processes to continue without any interruption.

2. Expansion at the DBMS-level allowing a new DBMS server to join.

Another docker container could be configured as another mongod instance and connected to the mongos router, a typical requirement when the volume of data gets too large. In fact, this was how we added the two GridFS servers (grid1, grid2) to manage the multimedia data. However, at the moment, the steps to configure another docker container as another DBMS and adding to the distributed system environment are quite convoluted, and in the future, a GUI could be designed for easy addition of a new DBMS server.

3. Dropping a DBMS server at will.

For the cases where a DBMS server is detected to have corrupt data/ missing data, it is unremunerative to maintain it in the system since the data would not be consistent among the replicas and lead to missing data in the new replicas. To maximize the space in our system, we need to enable the system to be able to detect and drop DBMS servers that fit the category.

4. Data migration from one data center to other.

Data center migration in the context of enterprises would refer to the physical movement of a data center from one location to another, and this tends to require proper planning, time and high costs. The goal is to minimize downtime and disruption while users experience no interruptions on their side during data migration which can be a result due to any change in our hardware systems. Therefore, there is need to ensure that whenever migration occurs, the data remains consistent between databases after a migration so that there is no case of corrupted or partially migrated data or in the worst case missing data.

6.1.1 **GUI**

Aside from the future work listed above, creating a user-friendly GUI to operating our database system would be ideal. At the moment our system can only be utilised through the terminal, and requires a certain level of technical knowledge to be able to use it. Furthermore, it takes quite a lot of time to key in all the required commands. Therefore there is a need to build a user friendly graphical interface where users can directly access the multimedia data from the articles at the front-end, and for administrators to easily manage the database system at the back-end.

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7 Appendix A: Manual

The detailed instructions and scripts for our project can be found in our Github repository https://github.com/kwyoke/DDBS-2020, but we still present the main steps required for the set-up in this section.

7.1 Prerequisites to Install

- 1. Ubuntu 18.04 machine
- 2. Docker, docker-compose
- 3. MongoDB Community Edition, mongofiles
- 4. PyMongo

We implemented the project on a Ubuntu 18.04 machine, but other machines should work fine as well. Docker and MongoDB should be installed according to their official documentations.

7.2 Generating dummy data

Run the executive_program/proj_data/genTable_mongoDB_toy.py to generate about 80MB of dummy data including user.dat, article.dat, read.dat and multimedia data. The 10GB version of the script can be found in mongoshell_tutorial/.

7.3 Running Makefile to set up everything

Navigate to executive_program/ which contains the Makefile. Simply run in the terminal:

```
make all # runs everything
```

Alternatively, to run only specific portions (see Makefile for more detail):

```
make setup # containers
make bulkload # copy and load user, article, tables
make popNshard # populate and shard collections
make store_multimedia # store multimedia data in GridFS servers
```

Now, all the required collections (user, article, read, beread, popRank) are populated and sharded, and queries can be performed in the mongos server's mongo shell.

Other relevant information:

The Makefile sets everything up, with the help of popNshard.js which contains the mongo shell commands for populating and sharding the collections, and bulk_store_multimedia.sh for storing multimedia files into GridFS.

Also, run auto_refresh_onarticle.py and auto_refresh_onread.py in terminals so that the collections will be updated when there are any changes.

8 Appendix B: Workload allocation

The project workload is allocated as shown in Figure 16, with Kai Wen as the leader. The colour intensity indicates the amount of contribution in each area.

| | | | Code implementation | |
|---------|------------|---------------------|---------------------|------------|
| | Discussion | Literature research | (MongoDB, Docker) | and report |
| Kai Wen | | | | |
| Barbara | | | | |
| | | | | |

FIGURE 16: Work allocation