Lalo Manual

1. Assumed Prior Lnowledge

11 Background knowledge

opcodes, instructions, subrantines, stacks, registers, program bits, hibbles, bytes endianness

x86-64 - little-endian machine

14 64-6it general purpose registers
(RAX, RBK, RCX, RDX, RDX, RSI) and R8-RIS)

64-6it stack pointer register (RSP) 4 contains the nemory address of the top of

the current program stack base pointer register (RBP) 4 used during subroutine execution

Von Neumann architecture

subsystems > random access main memory

CPU -> capable of executing instructions, residing in the main memory, which are binary codes

1.2 Essential Programs are stored in the computer memor sequences of instructions and data in binary machine language representation of a 01001000 0 48 c7 c0 01 00 00 00 # move the number 1 into the RAX register 48 c7 c1 01 00 00 00 # move the number 1 into the RCX register 11000111 1 # add the contents of RAX to RCX 11000000 00000001 punch cards - zeros and ones assemblers (1950s) - computer programs 00000000 00000000 00000000 L>translates text from symbolic 01001000 assembly language to machine code 11000111 11000001 instruction code - mnemonic, 10 00000001 represented in decimal or kex 11 00000000 12 00000000 each ardifecture -> own machine 13 00000000 14 01001000 assembly language - code DWL 00000001 15 \$1, %rax # Move the number 1 into the RAX register \$1, %rcx # Move the number 1 into the RCX register 11000001 16 %rax, %rcx # Add the contents of RAX to RCX Givary

2. Designing a Program

description -> psendocode -> assembly code

2.1. Description	input 0	output 1	
iuput≥0	1 2 3	1 3 3	
even -> +1	$\frac{10}{21}$	11 21	
	$\frac{42}{1041}$	43 1041	

2.2. Specification

2.3. I uplementation

movq %rbp , %rsp popq %rbp



3. Assembler Directives (. bss, ,-text, ; global) - stip, asciz, etc.) L'special functions tells the assembler to put a all subsequent code in a specific section ³makes certain labels visible memory space of a program > 3 different sections

text

hold all instructions, read-only < text .data initialized variables -. data .bss uninitialized variables . 6ss 3.2. Defining constants: equ .equ NAME, EXPRESSION defines symbolic names for expressions .equ FOO, 1024 pushq \$FOO # push 1024 3.3. Declaring variables: . byte, . word, . long, . quad .byte VALUE reserves and initialises memory for variables .word VALUE .long VALUE .quad VALUE FOO: .byte 0xAA, 0xBB, 0xCC # three bytes starting at address FOO BAR: .word 2718, 2818 # a couple of words BAZ: .long 0xDEADBEEF # a single long BAK: .quad 0xDEADBEEFBAADF00D # a single quadword FOO: .byte 0x0D, 0xF0, 0xAD, 0xBA, 0xEF, 0xBE, 0xAD, 0xDE FOO: .word 0xF00D, 0xBAAD, 0xBEEF, 0xDEAD FOO: .long 0xBADF00D, 0xDEADBEEF FOO: .quad 0xDEADBEEFBAADF00D 3.4. Reserving memory: . Skip AMOUNT BUFFER: .skip 1024 # reserve 1024 bytes of memory

STRING: ascii string reserves and initialises blocks of STRINGZ: asciz string ASCIS encoded characters

WELCOME: ascii "Hello!!" # A string.
byte 0x00 # ..followed by a 0-byte.

WELCOME: asciz "Hello!!" # A string followed by a 0-byte.

3.6. Global symbols: .flobal .global label

enters a label into the symbol table,

table of contents, contained in binary assembled filer

labels > useful for access by other programs

visible in debuyer <

other programs using subractives <

export main label -> very important

4 shows the operating system where to start

.global main

running the program

4. x86-64 Assembly Language

4.1. Instructions and Operands

what should the data cappen to act with

4.1.1. Operand Prefixes: Registers and Literal Values register names → % literal values → \$

4. L. 2 Instruction Postfixes (Specifying Operand Size)
b, w, l, q, modifiers
4 byte, word (26), long (46), quadword (86)
specifies the size of the operand

```
pushb $3 # Push one byte onto the stack (0x03) (NOTE: See below)
pushw $3 # Push two bytes onto the stack (0x0003)
pushl $3 # Push four bytes onto the stack (0x00000003)
pushq $3 # Push eight bytes onto the stack (0x000000000000000000)
```

4.1.3. Partial Registers

0 judicect

addresses smaller parts of	the	64-k	oits (egister
movl %eax, %ebx # Copy 32 bits values between registers	8-byte	4-byte	2-byte	1-byte
movq %rax, %rbx # Copy 64 bits values between registers movw %ax, %bx # Copy only the lowest order 16 bits	%rax	%eax	%ax	%al
movb %al, %bl # Copy only the lowest order 8 bits	%rcx	%ecx	%cx	%cl
movb %ah, %al # Copy 8 bits within a single register	%rdx	%edx	%dx	%dl
	%rbx	%ebx	%bx	%bl
	%rsi	%esi	%si	%sil
u Addonasia May - ou	%rdi	%edi	%di	%dil
4. Addressing Memory • immediate	%rsp	%esp	%sp	%spl
ماده داد ماده	%rbp	%ebp	%bp	%bpl
o immediate	%r8	%r8d	%r 8 w	%r8b
1 1-1 - 10 1 1 10	%r9	%r9d	%r9w	%r9b
label → the value at the	%r10	%r10d	%r 10 w	%r10b
P. 10 . 1 . 1	%r11	%r11d	%r11w	%r11b
address of the label	%r12	%r12d	%r12w	%r12b
d 1	%r13	%r13d	%r13w	%r13b
\$label > the address	%r14	%r14d	%r14w	%r14b
0) 1:	%r15	%r15d	%r15w	%r15b

(%RAX) -> value at the memory address

 $-8(\% \text{ PBP}) \rightarrow \text{value}$ at the memory address stored in register + displacement table (%PDI, % RCX, &) \equiv

displacement (base judex, scale)
constant
expression registers 124 or 8

= displacement + base + index × scale

4.2. Eustruction Set

4.11

Mnemonic	Operands	Action	Description				Branching
	-,		Data Transfer	jmp	ADDRESS		Jump to address (or label).
mov.	SRC, DST	DST = SRC	Copy.	je	ADDRESS ADDRESS		Jump if equal.
pushq	SRC	%RSP -= 8, (%RSP) = SRC	Push a value onto the stack.	jne	ADDRESS		Jump if not equal. Jump if greater than.
popq	DST	DST = (%RSP), $%RSP += 8$	Pop a value from the stack.	jg jge	ADDRESS		Jump if greater than. Jump if greater or equal.
xchg.	A. B	TMP = A, A = B, B = TMP	Exchange two values.	11	ADDRESS		Jump if less than.
novzb.	SRC. DST	DST = SRC (one byte only)	Move byte, zero extended.	ile	ADDRESS		Jump if less or equal.
movzw.	SRC, DST	DST = SRC (one word only)	Move word, zero extended.	call	ADDRESS		Jump and push return address.
			Arithmetic	ret			Pop address and jump to it.
add.	SRC, DST	DST = DST + SRC	Addition.	loop	ADDRESS		decq %RCX, jump if not zero.
sub.	SRC, DST	DST = DST - SRC	Subtraction.				Logic and Shifting
inc.	DST	DST = DST + 1	Increment by one.	cmp.	A, B	sub A B (Only set flags)	Compare and set condition flags
dec.	DST	DST = DST - 1	Decrement by one.	xor.	SRC, DST	DST = SRC ^ DST	Bitwise exclusive or.
nul.	SRC	%RDX: %RAX = %RAX * SRC	Unsigned multiplication.	or.	SRC, DST SRC, DST	$DST = SRC \mid DST$ DST = SRC & DST	Bitwise inclusive or. Bitwise and.
imul.	SRC	%RDX: %RAX = %RAX * SRC	Signed multiplication.	and.	A. DST	DST = SRC & DST DST = DST << A	Shift left by one bit.
div.	SRC	%RAX = %RDX:%RAX / SRC	Unsigned division.	shr.		DST = DST << A DST = DST >> A	Shift right by one bit
		%RDX = %RDX: %RAX % SRC		out.	А, Бот	D31 = D31 >> A	Other
idiv.	SRC	%RAX = %RDX: %RAX / SRC	Signed division.	lea.	A. DST	DST = &A	Load effective address.
		%RDX = %RDX: %RAX % SRC	-	int	INT_NR		Software interrupt.

4.3. Registers & Variables variables

registers → fast and easy access

variables

stacks caller saved

main memory subroutines overwrite registers 4.4. The Stack 4.4.1. The Stack Pointer RSP — stack pointer register initialized by operating system at frontains the address of the first byte after programs memory space > "grows" downward into program's memory space L> push instruction → RSP value is decremented + pushed value stored at the new location at which the stack pointer then points

4.4.2. Cleaning up the Stack stack overflow - overwite program's code and data

4.4.3. The Base Pointer

RBP-base pointer rejister entry of subroutine > push value of RBP into the stack

end of subroutine > pop old RBP value odd the stack

during subroutine -> P-BP points at the base of

subroutines stack are a - find local variables and subroutine arguments

4.4.4. Subsortine Prologue and Epilogue

prologue - storing old base pointer and copying the stack

pointer to be a new base pointer

epilogue - restoring the old base pointer

stack frame - space between RBP and RSP

RBP points opposite of RSP

4.5. Subrontines
block of justructions starting at memory address (label)