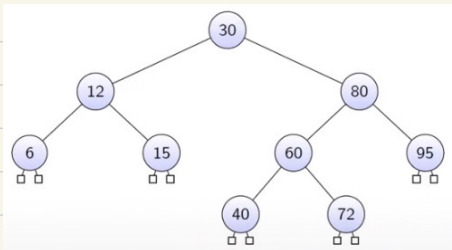
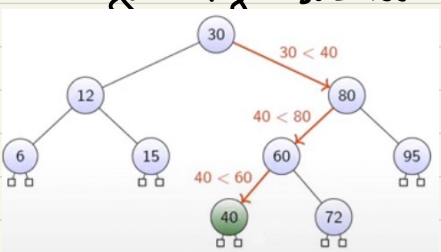
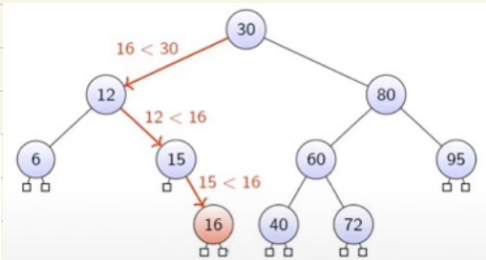
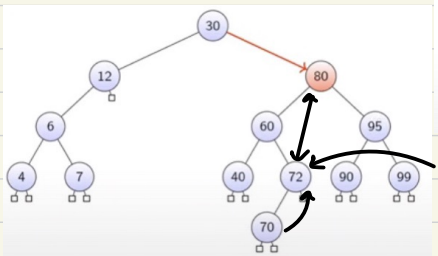
Video 1. Binary Search Tree

∀ node with key k:
→ keys in left subtree < k
→ keys in right subtree > k







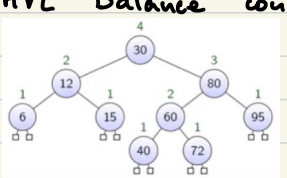


maximal node in left child
swap with node to be removed
recursively remove the node

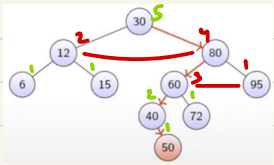
recurrence equation: $T(0) = c_0$ $T(n) = T(n-1) + c_1$
closed form: $T(n) = nc_1 + c_0$ ⇒ time complexity $O(n)$
height { best case $O(\log n)$
worst case $O(n)$

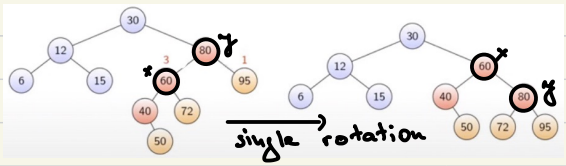
Video 2. AVL Tree

Adelson-Velskii and Landis
AVL Balance condition: the heights of the children $\Delta h \leq 1$
null references as leaves

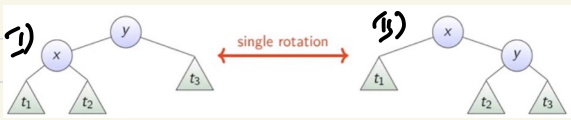


insert 50 ?



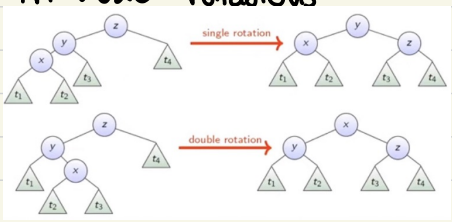


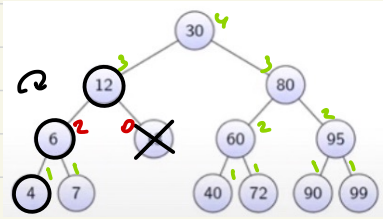
Two-node rotation

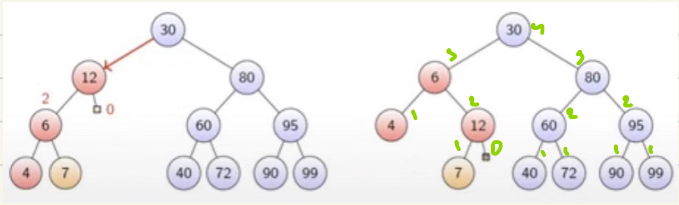


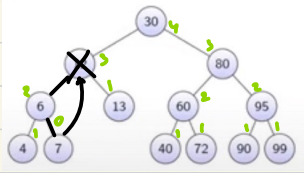
I -> II) $h(x) - h(t_3) > 1$
$h(t_1) - h(t_2) > 1$
II -> I) $h(y) - h(t_1) > 1$
$h(t_3) - h(t_2) > 1$

Tri-node rotations







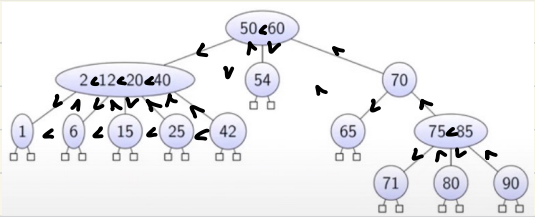


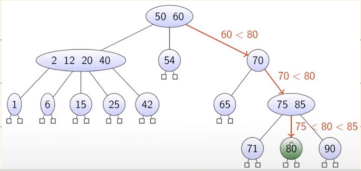
time-complexity (get, put, remove) = $O(\log n)$
tri-node restructuring = $O(1)$
get, put, remove -> $O(h)$ nodes + $O(1)$ restruc.
$h = O(\log n)$

Video 3. (2,4) Tree

internal node → at least 2 children
d children → ordered list of $k_1, \dots, k_{d-1}$ keys
$\forall k$ in i-th child → $k_{i-1} < k < k_i$ $k_0 = -\infty$ $k_d = \infty$

} multiway search tree



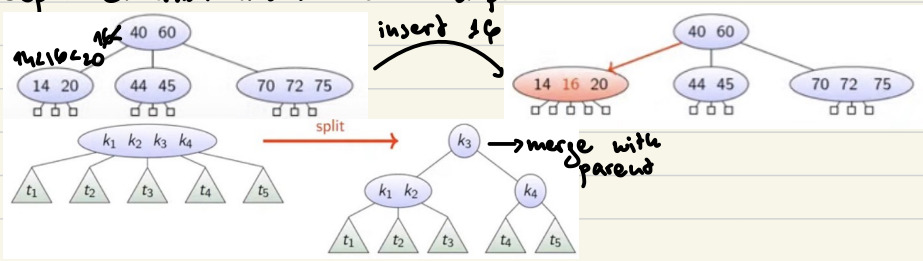


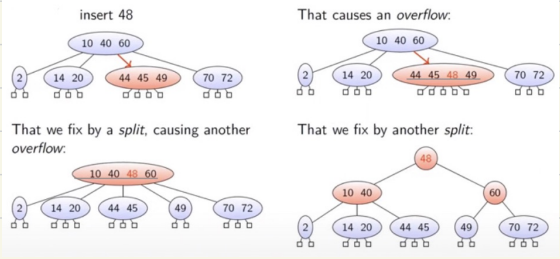
best case: $O(\log n)$

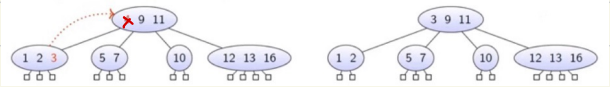
worst case: $O(n)$

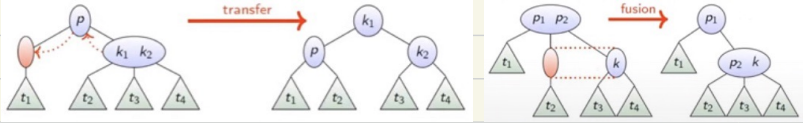
size: at most 4 children

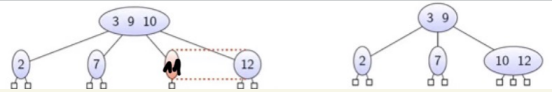
depth: external nodes -> same depth











get, put, remove = $O(\log n)$

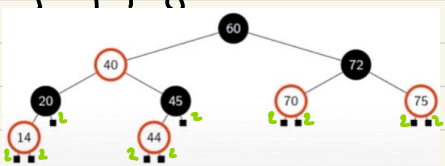
$h = O(\log n)$

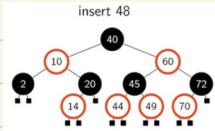
split, transfer, fusion = $O(1)$

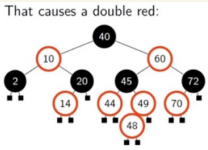
get, put, remove -> $O(h)$ nodes

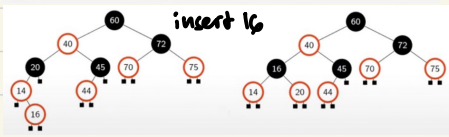
Video 4. Red-Black Tree

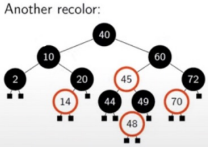
root property: the root is black
external property: every leaf is black
red property: the children of a red node are black
depth property: all leaves have the same black depth

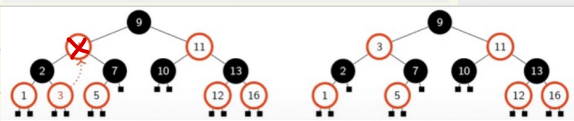


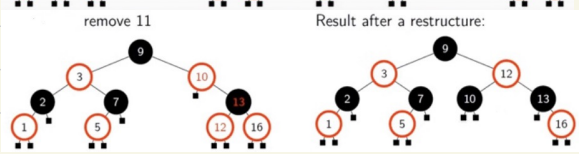


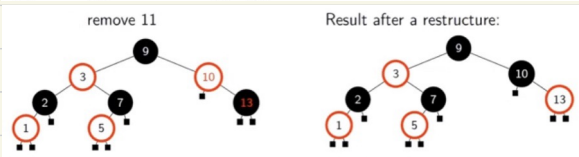












get, put, remove = O(log n)
height = O(log n)
recoloring = O(1)
get, put, remove -> O(h) nodes