

Quiz Submission

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Question 1

1. SuperSort(A)
sorts array A recursively
2. // base cases
3. if (length(A) == 1) return;
4. if (length(A) == 2)
5. swap the two elements if they're out of order;
6. return;
7. else // recursive calls
8. SuperSort the first two thirds of A;
9. SuperSort the second two thirds of A;
10. SuperSort the first two thirds of A again;
- 11.
- 12.

Answer:

so if length a = 1, we will return, if it's 2, we will swap elements(base), and then we will do 3 recursive calls on the first 2/3 of a, then 2nd, then 3rd.
therefore, Our recursive solution would be $T(n)=3T(2n/3)+O(n)$.

Question 2

Part B: Give a recurrence relation upper bound, tight up to constant factors, for the performance of your algorithm given in Part A. Don't forget the base case(s).

Answer:

$$T(n) = T(k/2) + O(1)$$

Question 3

Part C: Solve your recurrence relation given in Part B by providing an asymptotic solution tight up to constant factors, but do NOT use the master theorem or the master-master theorem/nuclear bomb. Show your work.

Answer:

the solution would just be adding the constant factors,

Question 4

Suppose you are given positive integers where for nonnegative integer . You would like to determine using only the elementary operations of addition, subtraction, multiplication, and division.

A brute force algorithm is described below.

The brute force algorithm requires time to find the correct value of . To help with this task, your friend designed a function called which takes input of the form and computes in time. In Part A, you will use to design a divide and conquer algorithm to find the correct value of with smaller asymptotic runtime than the above brute force algorithm. Hint: , so such that if is even or if is odd.

Part A: Design a divide and conquer algorithm that has an asymptotically faster performance than the given brute force algorithm.

Answer:

base : if $a = 1$, then $k = 0$. our design is, If let's say that we have a variable n and $n = \sqrt[k]{a}$ so the means that $n = b^{k/2}$. we want to decide if k is even or odd, if it's even then n would be $(b^{k/2})^2$ so that makes n^2 then $k = 2(k/2)$ then if it's odd then $a = b$ multiplied by $b \cdot (b^{k/2})^2$ and that would make $k = 2(k/2) + 1$. then the runtime would be $O(\log k)$